The University of Michigan Electrical Engineering & Computer Science EECS 281: Data Structures and Algorithms Spring 2024



PRACTICE MIDTERM EXAM **KEY 1**

Thursday, May 30, 2024 10:30AM – 12:30PM

Uniqname: Student ID:		
Name:		
Uniquame of person to your left:		
Uniquame of person to your right:		
Honor Pledge: "I have neither given nor received unauthorized aid on this examination, nor have I concealed any violations of the Honor Code." Signature:		

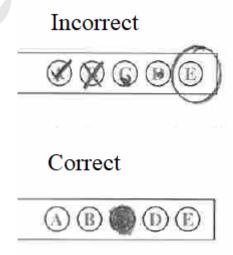
INSTRUCTIONS:

- The exam is closed book and notes except for one 8.5"x11" sheet of handwritten notes (both sides). All electronic device use is prohibited during the exam (calculators, phones, laptops, etc.).
- Print your name, student ID number, and uniquame **LEGIBLY**. Sign the Honor Pledge; this pledge applies to both the written and multiple choice sections. Your exam will not be graded without your signature.
- Record the **UNIQNAME** of the students to both your left and right. Write down nullptr if you are at the end of a row and there is no one on a given side.
- Record your **UNIQNAME** at the top of each odd-numbered page in the space provided. This will allow us to identify your exam if pages become separated during grading.
- This exam booklet contains **BOTH** the multiple choice and written portions of the exam. Instructions for each portion are provided at the beginning of their respective sections.
- **DO NOT** detach any pages from this booklet.
- There should be 20 pages in this book; if you are missing any pages, please let an instructor know.

Multiple Choice Portion [60 points]

MULTIPLE CHOICE INSTRUCTIONS:

- Select only **ONE** answer per multiple choice question.
- There are 24 multiple choice questions worth 2.5 points each.
- Record your choices for all multiple choice questions using the circles in the boxed area next to each question. Use a number 2 pencil, not a pen (so that you may change your answer). Fill the circle completely. There is no partial credit for multiple-choice questions. Make sure you bubble in the correct letter. See the example below for how to select your answer.



- There is no penalty for wrong answers: it is to your benefit to record an answer to each multiplechoice question, even if you're not sure what the correct answer is.
- The questions vary in difficulty try not to spend too much time on any one question. Use your time wisely.
- When asked about memory complexity, consider all possible sources (stack and heap).
- The term "heap" as a data structure refers to a binary heap, unless explicitly stated otherwise.
- The term "binary search tree" refers to a standard implementation that does not self-balance, unless explicitly stated otherwise.

1. Master Theorem

(A) (B) (C) (D) (E)

Which of the following recurrence relations can one solve by applying the Master Theorem?

A)
$$T(n) = \frac{1}{2}T(\frac{n}{4}) + \Theta(n)$$

B)
$$T(n) = T(\frac{n}{2}) + \Theta(n^{\log_2 4})$$

C)
$$T(n) = 4T(n-1) + \Theta(1)$$

D)
$$T(n) = 4T(\frac{n}{2}) + 2T(\frac{n}{4}) + \Theta(n^2)$$

E) None of the above.

2. Measuring Runtime



Consider the following function:

```
int func(double m, double n) {
2
        int c = 0;
3
        while (m > n) {
4
            m = m/2;
5
            for (int i = 1; i < n; i*=3;)</pre>
                 ++c;
6
7
        } // while
8
9
        return c;
10
     } // func
```

What best characterizes the running time of a call to this function?

- **A)** $O(n \log m)$
- **B)** $O(\frac{m}{n}\log n)$
- C) $O((\log \frac{m}{n}) \cdot (\log n))$
- **D)** $O(mn\log(mn))$
- **E)** $O(mn \cdot (\log m) \cdot (\log n))$

3. Math Foundations

(A) (B) (C) (D) (E)

Which of the following is **NOT** true regarding logarithms?

- **A)** $\log_2(n) = O(\log_3(n))$
- **B)** $\log_3(n) = O(\log_2(n))$
- **C)** $\log_2(n!) = O(n \log_2(n))$
- **D)** $\log_2(n) = \Theta(\log_3(n))$
- **E**) All of the above are true

4. Complexity Analysis



If f(n) = O(g(n)), and $f(n) = \Omega(h(n))$ and f(n), g(n), and h(n) are strictly increasing functions, which of the following MUST be true?

- **A)** g(n) = O(h(n))
- **B)** $g(n) = \Theta(h(n))$
- C) $g(n) = \Omega(h(n))$
- **D)** $g(n) = \Theta(f(n))$
- **E**) None of the above

5. Recursion

(A) (B) (C) (D) (E)

Consider the following recurrence relation:

$$T(n) = \begin{cases} c_0, & n = 1\\ nT(n-1) + c, & n > 1. \end{cases}$$

What is the complexity of T(n)?

- **A)** $\Theta(n \log n)$
- **B)** $\Theta(n^2)$
- C) $\Theta(n^3)$
- **D)** $\Theta(c^n)$
- E) $\Theta(n!)$

6. Queues



What is the best possible space complexity for a function printing out the contents of a std::queue of size n in order, such that the contents and order of the queue are the same before and after the function is called and executed? This means the auxiliary, or additional, space required to make this occur, not the space taken up by the data that was in the queue to begin with. You can assume that the queue is passed by reference.

- **A)** O(1)
- **B)** O(n)
- **C**) $O(n^2)$
- **D)** $O(n^3)$

7. Code Analysis

NOTE: Same code as next page, different question

Consider an array data of size n which contains 0s and 1s. Each element of this array containing 0 appears before all those elements containing 1. For example, [0, 0, 0, 0, 1] is such an array.

```
int myFunction(int data[], int length) {
2
         int mid = 0, left = 0; int right = length - 1;
3
         if (data[0] == 1)
             return 0;
4
5
         else if (data[length - 1] == 0)
6
             return length;
7
         else {
8
             bool isFound = false;
9
             while (!isFound) {
10
                 mid = left + (right - left) / 2;
                 if (data[mid] == 0)
11
12
                     left = mid + 1;
13
                 else {
14
                     if (data[mid - 1] == 0)
15
                          isFound = true;
16
                     else
17
                          right = mid - 1;
18
                 } // else
19
             } // while
20
             return mid;
21
         } // else
22
    } // myFunction()
```

In general, what does the function return when it is given such an array?

- **A)** Number of 0's in data[]
- **B)** Number of 1's in data[]
- C) (Number of 0's in data[]) + 1
- **D)** (Number of 1's in data []) + 1
- **E)** (Number of 1's in data[]) 1

8. Complexity of Code



NOTE: Same code as previous page, different question

Consider an array data of size n which contains 0s and 1s. Each element of this array containing 0 appears before all those elements containing 1. For example, [0, 0, 0, 0, 1] is such an array.

```
int myFunction(int data[], int length) {
2
         int mid = 0, left = 0; int right = length - 1;
3
         if (data[0] == 1)
             return 0;
4
         else if (data[length - 1] == 0)
5
6
             return length;
7
         else {
8
             bool isFound = false;
9
             while (!isFound) {
10
                 mid = left + (right - left) / 2;
11
                 if (data[mid] == 0)
12
                     left = mid + 1;
13
                 else {
14
                     if (data[mid - 1] == 0)
15
                          isFound = true;
16
                     else
17
                          right = mid - 1;
18
                 } // else
19
             } // while
20
             return mid;
21
         } // else
22
    } // myFunction()
```

What is the complexity of the above algorithm?

- **A)** $\Theta(1)$
- **B**) $\Theta(\log(\log n))$
- C) $\Theta(\log n)$
- **D)** $\Theta(n)$
- **E)** $\Theta(n \log n)$

9. Heaps and Heapsort

(A) (B) (C) (D) (E)

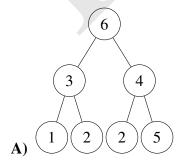
Which of the following is **TRUE** about binary heaps and heapsort?

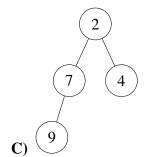
- A) Finding the maximum element in a min heap tree has a complexity of $\Theta(\log n)$
- **B)** Turning a random set of elements into a heap has a complexity of $\Theta(\log n)$
- C) If the priority of one element in a heap changes, fixing the heap requires a minimum of $\Theta(n)$ operations
- **D)** The worst-case time complexity of heapsort is $\Theta(n^2)$
- E) Heapsort can be implemented with O(1) additional space

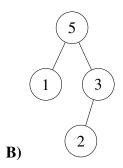
10. Binary Heaps

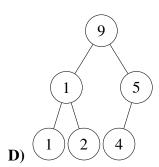
A B C D

Which of the following represents a valid binary heap?









11. Heaps



Suppose that we wanted to add a function popRandom() to our priority queues in Project 2. When implementing it with the binary heap, we could generate a random index, swap values between that index and the last element of the data vector, and call the vector's pop_back() function. What else would we have to do to preserve the binary heap property? Choose the answer with the **smallest worst-case** complexity.

- A) Call update_priorities().
- **B)** Call fixDown () followed by fixUp(), both on the random index.
- C) Call fixDown (), starting at the root of the heap.
- **D)** Call fixUp(), starting at the last element of the heap.
- E) Call exactly one of fixUp() or fixDown(), depending on the values swapped.

12. Reversing an Array



You are asked to reverse the order of an arbitrary array of length n in constant memory. Your code skeleton looks like:

```
1     for (...) {
2         std::swap(array[i], array[n - i - 1]);
3     }
```

What belongs in the for loop header?

- A) for (size_t i = 0; i < n; ++i)
- **B)** for $(size_t i = 0; i < (n / 2); ++i)$
- C) for (size_t i = n; $i \ge 0$; --i)
- **D**) None of the above

13. Linked Lists 1

Which of the following is a reason that you would use a doubly linked list over a vector?

- A) You need random access
- **B**) You need perform binary search on a data set
- C) You need to sort a large number of characters
- D) You're given a pointer to an element and need to insert an element right in front of it

14. Linked Lists 2



For which operation does a doubly linked list outperform a singly linked list? Assume the worst case, and that neither version of the linked list contains a tail pointer.

- A) Accessing element at random index
- **B**) Appending an element to the end
- C) Deleting an element given a pointer to the node that contains the element
- **D)** Swapping two elements that are not next to each other, given pointers to the two nodes

15. Array-based Containers

(A) (B) (C) (D)

Suppose we implement a vector with an underlying dynamic array and a special push_back () method. This push_back () method works normally when the array is not full. When the array fills up, the push_back () method creates a new dynamic array with space for 100 more elements than the old array, copies the elements from the old array to the new array, and then deletes the old array. What is the amortized complexity of this special push_back () operation?

- **A)** O(1)
- **B)** $O(\log n)$
- C) $O(\sqrt{n})$
- **D)** O(n)

16. Container Data Structures



Which data structure is best for implementing a sorted container that supports binary search?

- A) Stack
- B) Vector
- C) Linked List
- D) Binary Heap

17. Contiguous Containers



Which of the following statements about native (C-style) array and std::vector is TRUE?

- A) A vector can store references while a native array cannot.
- **B)** It is valid for a function to return a vector declared on its own stack, but invalid to return an array declared on its own stack
- C) Since a vector stores data on the heap, you must always explicitly destroy them using the delete operator
- D) A vector may be multi-dimensional, an array cannot

18. The STL



Which of the following statements is **TRUE**?

- A) If you define a class with a std::vector member variable, you must define the Big Three (copy constructor, assignment operator, destructor) for the class because vectors use dynamic memory
- B) STL container classes (such as std::vector) may only contain objects, not primitive types
- C) To use a std::priority_queue<T>, type T must support operator<()
- **D**) None of the above

19. Ordered Arrays and Related Algorithms

ABCD

Given a C-string, what is the best space complexity of an algorithm to determine whether the string contains duplicate characters?

- **A)** O(1)
- **B)** O(n)
- C) $O(n \log n)$
- **D)** $O(n^2)$

20. Ordered Arrays



Given two sorted arrays, each with n numbers, you are asked to find the median of all of the elements from the two sorted arrays. What is the time complexity of the best solution?

- **A)** O(1)
- **B)** $O(\log n)$
- C) O(n)
- **D)** $O(n \log n)$
- **E)** $O(n^2)$

http://www.geeksforgeeks.org/median-of-two-sorted-arrays/

21. Sorting



On a nearly-sorted singly linked list, which of the following sorting algorithms is asymptotically fastest on large input?

- A) Heap sort
- **B**) Insertion sort
- C) Merge sort
- **D**) All of these would be linear

22. Quicksort



Suppose we run quicksort on the following list [4, 12, 9, 1, 2, 5, 3, 8, 7]. We choose the last element, 7, as the pivot. After partitioning as described in lecture (but **BEFORE** we run the recursive function calls) what is the resulting list?

- **A)** [4, 3, 5, 1, 2, 7, 12, 9, 8]
- **B**) [4, 3, 5, 1, 2, 9, 12, 8, 7]
- **C**) [4, 3, 1, 5, 2, 7, 12, 8, 9]
- **D**) [4, 3, 5, 1, 2, 7, 12, 8, 9]

23. Sorting Students



Given the following code:

```
enum class Standing_e : int { FRESHMAN, SOPHOMORE, JUNIOR, SENIOR };
2
3
    struct Student {
        int age;
4
5
        Standing_e standing;
        double gpa;
6
7
    } ;
8
    bool operator<(const Student& s1, const Student& s2) const {
9
10
        return static_cast<int>(s1.standing) < static_cast<int>(s2.standing);
11
```

Given an array of Student objects of size n (where n is very large), what would be the best possible worst-case time complexity to sort this array by class standing?

- A) $\Theta(n)$
- **B)** $\Theta(n \log n)$
- C) $\Theta(n^2)$
- **D)** $\Theta(n^3)$
- **E**) None of the above

24. Elementary Sorts



Executing a selection sort algorithm on a singly linked list has an asymptotically worse average-case time complexity than executing a selection sort on an array with random access.

- A) True
- **B**) False

Written Portion [40 points]

WRITTEN PORTION INSTRUCTIONS:

• There are **two** questions in this section, with the following point distribution:

Question 25	20 points
Question 26	20 points

- Please write your solution legibly in the space provided. Solutions that are difficult to read may receive fewer points.
- Use the back of the question page to work out solutions for each problem and then rewrite your final answer in the space provided.
- The directions for each problem will specify which STL algorithms/functions may be used in your solution.
- Solutions that exceed the allowed line count will lose one (1) point per line over the limit.
- Partial credit will be awarded for partial solutions.
- Errors in program logic or syntax may result in a reduced score.
- Be sure to consider all possible sources of memory (stack and heap) and time complexity.

25. Programming: Linear-time Algorithms and the STL [20 Points]

Implement STL's unique_copy () function according to its official interface and description.

Quoting from http://www.sgi.com/tech/stl/unique_copy.html:

unique_copy() copies elements from the range [first, last) to a range beginning with result, except that in a consecutive group of duplicate elements only the first one is copied. The return value is the end of the range to which the elements are copied.

Complexity: Linear. For elements in the ranges, exactly last - first applications of operator==() and at most last - first assignments.

For example, if you executed this code:

```
int data[6] = {1, 3, 3, 1, 1, 0}, output[6];
unique_copy(data, data + 6, output);
```

You would have this in array output:

```
1 3 1 0
```

Implementation: Use the back of this page as a working area, then rewrite **neatly** on the front. Limit: 15 lines of code (points deducted if longer). You may **NOT** use other STL algorithms/functions.

This page is intentionally left blank. You may use this page as working space.

26. Programming: Finding Max Elements [20 Points]

Suppose you are given an array of int, size n and a number k. Return the k largest elements Output does not need to be sorted. You can assume that k < n.

Requirements: Your solution must be faster than $O(n \log n)$ time. You may use up to O(n) auxiliary space.

Implementation: Use the back of this page as a working area, then rewrite neatly on the front. Limit: 15 lines of code (points deducted if longer). You may use any C, C++, or STL function/algorithm that you wish the process of the process of the page as a working area, then rewrite neatly on the front. Limit: 15 lines of code (points deducted if longer). You may use any C, C++, or STL function/algorithm that you wish the page as a working area, then rewrite neatly on the front. Limit: 15 lines of code (points deducted if longer). You may use any C, C++, or STL function/algorithm that you wish the page as a working area, then rewrite neatly on the front. Limit: 15 lines of code (points deducted if longer). You may use any C, C++, or STL function/algorithm that you wish the page as a working area, then rewrite neatly on the front.

This page is intentionally left blank. You may use this page as working space.