

p	q	$p \wedge q$	$p \vee q$	$p \oplus q$	$p \rightarrow q$	$p \leftrightarrow q$
T	T	T	T	F	T	T
T	F	F	T	T	F	F
F	T	F	T	T	T	F
F	F	F	F	F	T	T

TABLE 7 Logical Equivalences Involving Conditional Statements.

$p \rightarrow q \equiv \neg p \vee q$
 $p \rightarrow q \equiv \neg q \rightarrow \neg p$
 $p \vee q \equiv \neg p \rightarrow q$
 $p \wedge q \equiv \neg(p \rightarrow \neg q)$
 $\neg(p \rightarrow q) \equiv p \wedge \neg q$
 $(p \rightarrow q) \wedge (p \rightarrow r) \equiv (p \rightarrow (q \wedge r))$
 $(p \rightarrow r) \wedge (q \rightarrow r) \equiv (p \vee q) \rightarrow r$
 $(p \rightarrow q) \vee (p \rightarrow r) \equiv p \rightarrow (q \vee r)$
 $(p \rightarrow r) \vee (q \rightarrow r) \equiv (p \wedge q) \rightarrow r$

$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$ $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$	Distributive laws
$\neg(p \wedge q) \equiv \neg p \vee \neg q$ $\neg(p \vee q) \equiv \neg p \wedge \neg q$	De Morgan's laws
$p \vee (p \wedge q) \equiv p$ $p \wedge (p \vee q) \equiv p$	Absorption laws

Diagonalization

General format for **diagonalization proofs**:

- Arrange items in a table
- Look at the elements along the main diagonal
- Manipulate each diagonal element and put them together to create a **new** row that is **guaranteed not** to be in the table
 - * Is the new row the same as row i of the table?
 - * No: its i^{th} bit is different from the string in row i .

New:

1	0	0	0	1	0
2	0	1	0	0	0
3	0	0	1	0	0
4	0	0	0	0	0
5	0	0	0	0	1

i	$f(i)$	digit i of $f(i)$	digit i of $f(i)$	digit i of $f(i)$	digit i of $f(i)$...
1	0	0	0	0	0	...
2	$\pi/10$	0	3	1	4	...
3	0.149	0	1	4	9	...
4	$e/10$	0	2	7	1	...
...

Mod

Recall: $a \equiv b \pmod{m}$ means $a = b + km$ for some integer k (and assuming m is a positive integer)

Suppose $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$.

Claim: $a+c \equiv b+d \pmod{m}$ (Addition works!)

Claim: $a-c \equiv b-d \pmod{m}$ (Subtraction works!)

Claim: $ac \equiv bd \pmod{m}$ (Multiplication works!)

Some proofs. Let $a = b + km$ and $c = d + jm$.

So $a+c = b+km+d+jm = (b+d) + (k+j)m \equiv b+d \pmod{m}$

So $ac = (b+km)(d+jm) = bd + (bj+dk+kjm)m \equiv bd \pmod{m}$

If p is prime, then for any positive int $a < p$, there exists a unique positive $a^{-1} < p$ s.t. $aa^{-1} \equiv 1 \pmod{p}$

Inverse of a can only be calculated if p and a are relatively prime!

Guide for Induction Proofs

- Restate the claim you are trying to prove
- **Base case:** Prove the claim holds for the "first" value of n
 - Prove $P(n_0)$ is true
- **Inductive Step:** Prove that $P(k) \rightarrow P(k+1)$ for an arbitrary integer k in the desired range.
 - Let k be an arbitrary integer with $k \geq n_0$
 - Assume $P(k)$
 - Show that $P(k+1)$ holds
- **Conclusion:** explain that you've proven the desired claim.

Remember that the inductive step is like "climbing up the ladder", so it needs to cover the first rung we get on (the base case) as well!

Important note: Your base case should be included at the beginning of the domain of the inductive step

Equivalently: Show $P(k-1) \rightarrow P(k)$

To prove $A=B$:

Definition of Set Minus

$$A - B = A \cap \bar{B}$$

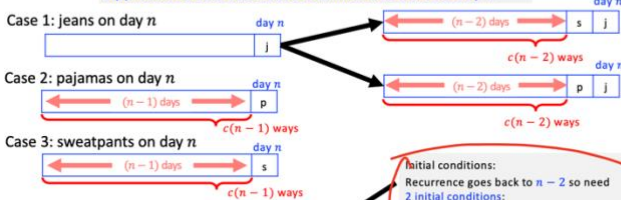
1. Prove that $A \subseteq B$
 - a. assume an arbitrary element x in A
 - b. show that this element must also be in B
2. Prove that $B \subseteq A$
 - a. assume an arbitrary element x in B
 - b. show that this element must also be in A

Example: Katie's Pandemic Outfits

Katie has three outfit choices while social distancing. Each day she wears either pajamas, sweatpants, or jeans. Her only rule is that she never wears jeans on two or more days in a row.

Let $c(n)$ represent the number of ways Katie can choose outfits across n days, $n \geq 0$. Find a recurrence relation for $c(n)$ including the initial conditions

Approach: Create cases based on what she wore on day n



Sum counts across all cases to get:

$$c(n) = 2c(n-1) + c(n-2)$$

With initial conditions:

$$c(0) = 1, \quad c(1) = 3$$

$$c(n) = \begin{cases} \text{Case 1: } c(n-2) + c(n-2) = 2 \\ \text{Case 2: } c(n-1) = 1 \\ \text{Case 3: } c(n-1) = 1 \end{cases}$$

Strong Induction as Dominoes

Let $P(n)$ be a predicate.

Goal: Prove that $P(n)$ is true for all $n \in \mathbb{Z}^+$

Step 1: [Strong] Inductive Step

If you can knock down **all the previous** dominos, then you can knock down the **next one**.

$$(P(1) \wedge P(2) \wedge \dots \wedge P(k)) \rightarrow P(k+1)$$

Step 2: Base Case(s)

You can knock down the first domino(s)

Possibly also:
 $P(2)$ and $P(3)$ and more

Therefore, you can knock down all dominos.

Weak induction vs Strong Induction

Basic Step: $P(1)$ Same $P(1)$
 Inductive Step: $\forall k \in \mathbb{Z}^+ [P(k) \rightarrow P(k+1)]$ $\forall k \in \mathbb{Z}^+ [P(1) \wedge \dots \wedge P(k) \rightarrow P(k+1)]$
 Conclusion: $\forall n \in \mathbb{Z}^+ P(n)$ Same $\forall n \in \mathbb{Z}^+ P(n)$

Strong Induction 有更强的条件
所以能 strong Induction
反之亦然

1. cardinality) $|S|$: the number of elements in S .
2. (Cartesian Product) $A \times B = \{(a,b) | a \in A, b \in B\}$
 $A \times B \times \dots \times N = \{(a,b,\dots,n) | a \in A, b \in B, \dots, n \in N\}$
3. $S \cap T$, $S \cup T$, \bar{S} , $S - T$
4. subset, proper set, disjoint
5. Power set: $PCS = \{A | A \subseteq S\}$
6. $A \subseteq B \wedge B \subseteq A \Rightarrow A = B$
7. $\overline{A \cup B} = \bar{A} \cap \bar{B}$
8. $|S| = n \Rightarrow |PCS| = 2^n$
9. $|A \cup B| = |A| + |B| - |A \cap B|$
 $|A \cap B| = |A| + |B| - |A \cup B|$
10. $|A + B + C| = |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cap B \cap C|$
 $|U_i A_i| = \{\text{individual sizes}\} - \{\text{pairwise } \cap \text{ sizes}\} + \{\text{3-wise } \cap\}$

$\{ \text{individual sizes} \} - \{ \text{pairwise } N \} + \{ \text{3-wise } N \} - \{ \text{4-wise } N \} + \{ \text{5-wise } N \} - \dots$

If $f: A \rightarrow B$ is onto, then $|A| \geq |B|$

To show that f is not surjective Find a particular $y \in B$ such that $f(x) \neq y$ for all $x \in A$.

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Definition: A set S is **countable** iff $|S| \leq |\mathbb{Z}^+|$
(otherwise it's *uncountable*)