# EECS 203: Discrete Mathematics Fall 2023 Homework 5

# Due Thursday, October. 12, 10:00 pm

No late homework accepted past midnight.

Number of Problems: 7 + 2 Total Points: 100 + 20

- Match your pages! Your submission time is when you upload the file, so the time you take to match pages doesn't count against you.
- Submit this assignment (and any regrade requests later) on Gradescope.
- Justify your answers and show your work (unless a question says otherwise).
- By submitting this homework, you agree that you are in compliance with the Engineering Honor Code and the Course Policies for 203, and that you are submitting your own work.
- Check the syllabus for full details.

# **Individual Portion**

# 1. Induction Construction [16 points]

Let P(n) be the statement that  $1 \cdot 1! + 2 \cdot 2! + \cdots + n \cdot n! = (n+1)! - 1$  whenever n is a positive integer. In this problem, we will prove this statement via weak induction.

- (a) What is the statement P(1)?
- (b) Show that P(1) is true, which is the base case for our inductive step.
- (c) In the base case we prove P(1); what do you need to prove in the inductive step?
- (d) What is the inductive hypothesis for your proof?
- (e) Complete the inductive step, indicating where you used the inductive hypothesis.
- (f) Explain why this proof shows P(n) is true for all positive integers n.

#### **Solution:**

- (a) P(1) is the statement that  $1 \cdot 1! = (1+1)! 1$
- (b)  $(1+1)! 1 = 2! 1 = 2 1 = 1 = 1 \cdot 1 = 1 \cdot 1!$ . Therefore P(1) is true.
- (c) We want to show that for all  $k \ge 1$ ,  $P(k) \to P(k+1)$ .
- (d) The inductive hypothesis is  $1 \cdot 1! + 2 \cdot 2! + \cdots + k \cdot k! = (k+1)! 1$  for some positive integer k.
- (e) Assume the inductive hypothesis. That is, let  $k \ge 1$  be an integer and assume  $1 \cdot 1! + 2 \cdot 2! + \cdots + k \cdot k! = (k+1)! 1$ . Then,

$$1 \cdot 1! + 2 \cdot 2! + \dots + k \cdot k! + (k+1) \cdot (k+1)!$$

$$= (k+1)! - 1 + (k+1) \cdot (k+1)!$$

$$= (k+1)! + (k+1) \cdot (k+1)! - 1$$

$$= (k+1)! [1 + (k+1)] - 1$$

$$= (k+1)!(k+2) - 1$$

$$= (k+2)! - 1$$
(by IH)

(f) Since we have shown P(1) is true, and  $P(k) \to P(k+1)$  for all positive integers k, we have shown  $1 \cdot 1! + 2 \cdot 2! + \cdots + n \cdot n! = (n+1)! - 1$  is true for all positive integers n by weak induction.

# Grading Guidelines [16 points]

### Parts a through d:

- +2 correct P(1)
- +2 correctly shows P(1) is true
- +3 correctly states we want to show that for all  $k \geq 1$ ,  $P(k) \rightarrow P(k+1)$
- +3 correct inductive hypothesis

#### Part e:

- +2 correct application of IH
- +2 simplification after application of IH

#### Part f:

+2 conclusion

# 2. Base Two Blues [14 points]

Prove using mathematical induction that  $\log_2(n) < n$  for every positive integer n. You may assume that the base-2 logarithm function is strictly increasing on its domain.

Fun Fact:  $\log_b(n) < n$  is actually true for every positive real number n and arbitrary base b > 1, but we're asking you to prove this by induction for the special case where b = 2 and n is a positive integer.

#### **Solution:**

# **Inductive Step:**

Assume  $\log_2(k) < k$ .

$$\begin{split} \log_2(k+1) & \leq \log_2(2k) & \text{(since } k \geq 1, \ k+1 \leq 2k) \\ & = \log_2(2) + \log_2(k) \\ & = 1 + \log_2(k) \\ & < k+1. & \text{(by IH)} \end{split}$$

#### Base case:

Let n = 1.  $\log_2(1) = 0 < 1$ . Therefore the base case holds.

By induction, we have proven that for every positive integer n,  $\log_2(n) < n$ .

# Grading Guidelines [14 points]

+3 assumes proper inequality  $\log_2(k) < k$ 

- +3 substituted 2k for k+1
- +3 applies log properties to break up  $\log_2(2k)$  into  $\log_2(2) + \log_2(k)$
- +3 applies IH
- +2 correct base case

# 3. Inductive Hypothe-six [15 points]

Prove by weak induction that 6 divides  $n^3 - n$  where n is a nonnegative integer. Don't include unneeded base cases.

#### **Solution:**

Let P(n) be the predicate that 6 divides  $n^3 - n$ .

Base case: n = 0

P(0) means 6 divides  $0^3 - 0$ . Because  $0^3 - 0 = 0$ , and  $6 \mid 0$  since 6(0) = 0, P(0) is true.

# Inductive step:

For our inductive hypothesis, assume that P(k) is true for some  $k \ge 0$ . Then 6 divides  $k^3 - k$ , so there exists some integer a such that  $k^3 - k = 6a$ .

We want to show that P(k+1) is true, meaning we want to prove that 6 divides  $(k+1)^3 - (k+1)$  by demonstrating that there exists some integer b such that  $(k+1)^3 - (k+1) = 6b$ .

Our solution will take  $(k+1)^3 - (k+1)$ , work on it so we can apply the inductive hypothesis, then do more work so we end up writing it as 6 times some integer.

$$(k+1)^{3} - (k+1) = (k^{3} + 3k^{2} + 3k + 1) + (-k-1)$$

$$= (k^{3} - k) + (3k^{2} + 3k) + (1-1)$$

$$= (k^{3} - k) + 3k(k+1)$$

$$= 6a + 3k(k+1)$$

$$= 6a + 3(2c)$$

$$= 6a + 6c$$

$$= 6(a+c)$$
by IH
$$c \in \mathbb{Z} \text{ see } \bigstar$$

 $\bigstar$  Since either k or k+1 is even, and an even integer times any integer is even, k(k+1) is even.

So,  $(k+1)^3 - (k+1) = 6(a+c)$ . Because a+c is an integer, that means 6 divides  $(k+1)^3 - (k+1)$ , so P(k+1) is true.

Hence, by induction, 6 divides  $n^3 - n$  where n is a nonnegative integer.

# Grading Guidelines [15 points]

- +2 correct base case
- +3 correct inductive hypothesis (assuming P(n) or P(n-1))
- +4 correct application of inductive hypothesis
- +3 correct algebra up to substitution of inductive hypothesis
- +3 correct algebra to show P(n+1) if assuming P(n) or P(n) if assuming P(n-1)

# 4. Incorrect Strong Induction [14 points]

For each of the following **incorrect** strong induction proofs, note where the strong induction proof breaks down and is incorrect.

Hint: Consider where the inductive step breaks down.

(a) Proving for every nonnegative integer n, P(n): 3n = 0.

### **Inductive Step:**

Assume that P(j): 3j = 0 for all nonnegative integers j with  $0 \le j \le k$ . We wish to show P(k+1). We will rewrite k+1=a+b where a and b are nonnegative integers less than k+1. Thus,  $3 \cdot (k+1) = 3 \cdot (a+b) = 3a+3b = 0+0 = 0$ , therefore P(k+1) is proven.

**Base Case:**  $P(0): 3 \cdot 0 = 0$ 

Since we have shown the basis step and the inductive step, we have proved for every nonnegative integer n, P(n): 3n = 0.

(b) Proving that every cent value above 3 cents can be formed using just 3-cent and 4-cent stamps.

# **Inductive Step:**

Assume we can form cent values of j cents for all  $3 \le j \le k$  using just 3-cent and 4-cent stamps. We wish to show we can form k+1 cents using just 3-cent and 4-cent stamps. We can form a k+1 cent value by replacing 1 3-cent stamp with 1 4-cent stamp or by replacing 2 4-cent stamps with 3 3-cent stamps.

#### Base Case:

We can form cent values of 3-cents using one 3-cent stamp and we can form cent values of 4-cents using one 4-cent stamp. This covers our two base cases.

Since we have shown the basis step and the inductive step, we have proved every cent value above 3 cents can be formed using just 3-cent and 4-cent stamps.

#### **Solution:**

- (a) This inductive step is not valid for j = 1, since the only two nonnegative integers a and b that add up to 1 are 0 and 1, however you can not use P(1) in a proof to prove P(1).
- (b) This inductive step is not valid for 5-cent values, since for the 4-cent value composed of 1 4-cent stamp, there are no 3-cent stamps to turn to 4-cent stamps and no 2 4-cent stamps to turn into 3 3-cent stamps.

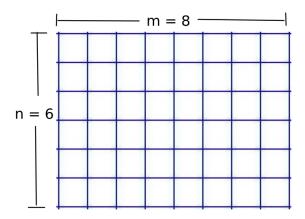
# Grading Guidelines [14 points]

#### Parts a and b:

- +3 correct identification of incorrect step
- +4 explanation of incorrect step

# 5. Chopping Ice [15 points]

Claire doesn't have an ice tray, so she makes ice by freezing water into a rectangle and then dividing the rectangle into grid-aligned cells. She would like to divide her block of ice into n rows and m columns quickly, before the ice melts! See the image below for an example.



- (a) State the number of cuts Claire needs to make to divide her ice block into  $n \times m$  cells. One cut means splitting a single rectangle into two rectangles. In other words, you may NOT make a single cut across multiple pieces of ice. You may use n and/or m in your answer.
- (b) Prove your answer from part (a).

#### **Solution:**

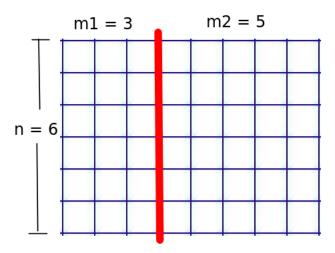
- (a) nm 1
- (b) Let P(A) mean "the minimum number of cuts to split a rectangle with area  $A = n \times m$  is nm 1." Use strong induction.

# **Inductive Hypothesis:**

Assume for all A' s.t.  $1 \le A' < A$ , P(A') is true (meaning for any rectangle of area A' = n'm' for some n' and m', n'm' - 1 cuts are the number of cuts necessary to split this block of ice).

**Inductive step:** We want to show P(A) given the inductive hypothesis.

Suppose  $A = n \times m$  for some integers n and m. Now, enumerate the first cut we made. Suppose we made a cut to divide our  $n \times m$  rectangle into two rectangles of area  $A_1$  and  $A_2$  rectangle where  $1 \leq A_1 < A$ , and  $1 \leq A_2 < A$ . Note that  $A_1 + A_2 = A$ . Suppose  $A_1 = n_1 \times m_1$  and  $A_2 = n_2 \times m_2$  for some integers  $n_1, n_2, m_1, m_2$ . By our inductive hypothesis, it would take  $(n_1m_1 - 1) + (n_2m_2 - 1)$  cuts to completely cut up both smaller rectangles. In total, it would take  $(n_1m_1 - 1) + (n_2m_2 - 1) + 1 = A_1 + A_2 - 2 + 1 = A - 1 = nm - 1$  cuts. Thus P(A) holds.



#### Base Case:

P(1): 0 cuts are necessary to split a 1 by 1 rectangle since one cube does not need to be split anymore.

#### **Alternate Base Cases:**

P(2): 1 cut is necessary to split an area 2 rectangle (1 by 2 or 2 by 1 rectangle) into two 1 by 1 pieces.

# Grading Guidelines [15 points]

- +3 states nm-1 (or equivalently area -1) in part (a)
- +1 states the correct base case: A = 1, or equivalently, m = n = 1. A = 2 (i.e. either m = 1, n = 2 or m = 2, n = 1) is acceptable if **both** combinations of m, n are explicitly specified values are correctly included.
- +2 correctly proves the base case(s) from the previous item
- +2 correct inductive hypothesis
- +2 correct "want to show" (can be implicit) for the inductive step
- +5 correctly shows P(A) using the inductive hypothesis and algebra

# 6. Pastry Recurrence [12 points]

A baker decorates a cookie in 2 minutes, a cupcake in 3 minutes, and a pie in 3 minutes. Let  $a_n$  denote the number of distinct ways the baker decorates pastries in exactly n minutes for  $n \ge 0$  (where order matters).

- (a) Find a recurrence relation for  $a_n$ .
- (b) What are the initial conditions? Use the fewest initial conditions necessary.

### **Solution:**

# (a) Solution: "Walking Backwords"

If the last pastry decorated was a cookie, then there are  $a_{n-2}$  possible ways to decorate. If the last pastry decorated was a cupcake, there are  $a_{n-3}$  ways. If the last pastry decorated was a pie, there are  $a_{n-3}$  ways. Therefore, our recurrence is  $a_n = a_{n-2} + a_{n-3} + a_{n-3} = a_{n-2} + 2a_{n-3}$ .

### Solution: "Walking Forwards"

If the first decorated pastry was a cookie, then the first 2 minutes is used to decorate a cookie, and there are  $a_{n-2}$  ways to use the rest of the minutes. If the first decorated pastry was a cupcake or pie, then there are 2 options for the first 3 minutes, and there are  $a_{n-3}$  ways to use the rest of the minutes. This would mean there are  $2a_{n-3}$  ways in this case. Therefore, our recurrence is  $a_n = a_{n-2} + a_{n-3} + a_{n-3} = a_{n-2} + 2a_{n-3}$ .

(b) For 0 minutes, there is one possibility (no baked goods decorated) so  $a_0 = 1$ . For 1 minute, there are zero possibilities (you can't use exactly 1 minute to decorate) so  $a_1 = 0$ . For 2 minutes, there is one possibility (a cookie) so  $a_2 = 1$ .

# Grading Guidelines [12 points]

#### Part a:

- +3 correctly handles the case when the last pastry decorated was a cookie
- +3 correctly handles the case when the last pastry decorated was a cupcake
- +3 correctly handles the case when the last pastry decorated was a pie

#### Part b:

- +1 correct number of initial conditions
- +2 correct value of initial conditions

# 7. Raven's Wrestlers [14 points]

Raven has n weeks to build her wrestling figure collection. Every week, Raven buys one item to add to her collection. There are 4 different types of things she can buy: Figures, T-shirts for her wrestlers to wear, Weapons for them to fight with, or Display Stands to show them off on her shelves.

- Her shelves can fit 2 Stands nicely, so when she buys a Display Stand, she will always buy a second one the next week to finish the shelf. Additionally, the week after buying the second Stand, she will buy something other than a Display Stand (they aren't as exciting to buy)
- When she buys a Figure, she gets very excited about it and wants to buy a new T-shirt for it to wear the following week.

Let  $a_n$  represent the number of ways Raven can buy items across the n weeks (where  $n \geq 0$ )

- (a) Find a recurrence relation for  $a_n$ .
- (b) Which terms would need to be defined with initial conditions (no need to find the value, just which terms)

Note 1: Buying the same items in a different order counts as a different way of buying items. We treat all items in a category as identical.

Note 2: on week n, Raven will not buy a Figure (because she knows she will miss buying a T-shirt) or a Stand (what a sad way to end the collection). This information is not needed for the simplest solutions, but some alternate solutions may need to know this.

Solution:			

### (a) Forward solution:

On week 1, Raven can purchase any of the 4 items.

#### Case 1: T-shirt

There are no further requirements, so on week 2, she can again buy anything. This means the number of ways to select items to buy for the remaining n-1 weeks is  $a_{n-1}$ , as you can just ignore week 1.

### Case 2: Weapon

Similarly, there are no further requirements, so there are  $a_{n-1}$  ways to select the remaining weeks.

# Case 3: Figure

In this case, Raven will buy a T-shirt on week 2. After that, though, we are back to being able to buy anything, so there are  $a_{n-2}$  ways to select for the remaining n-2 weeks.

#### Case 4: Stand

In this case, Raven will have to buy a stand on week 2, and for week 3, we know she will not buy another stand, so we can't just use  $a_{n-2}$ . Instead, we need subcases for the other options:

Case 4a: Stand, Stand, T-shirt.  $a_{n-3}$ 

Case 4b: Stand, Stand, Weapon.  $a_{n-3}$ 

Case 4c: Stand, Stand, Figure

In this case, she will buy a T-shirt on week 4, leaving only n-4 weeks to decide with no requirements. This means there are  $a_{n-4}$  for this subcase.

Putting all this together, we get  $a_n = 2a_{n-1} + a_{n-2} + 2a_{n-3} + a_{n-4}$ 

#### Backward solution:

On week n, Raven can only purchase T-shirts or Weapons. This means that if we ever find ourselves with more options than just these 2, we will need to branch those off as subcases.

#### Case 1: Weapon

If Raven buys a Weapon on week n, she still cannot buy a Figure on week n-1 (otherwise, she'd have to buy a T-shirt on week n). She could, however, buy a Stand on week n-1.

# Case 1a: Weapon, T-shirt or Weapon

If we know she bought a T-shirt or Weapon on week n-1, that actually means that the first n-1 weeks have the same constraints as if there were no week n (because she can't buy Stands or Figures on the last week). This means there are  $a_{n-1}$  ways to buy items with a Weapon on week n and either a T-shirt or Weapon

on week n-1.

### Case 1b: Weapon, Stand

If Raven buys a Weapon on week n and a Stand on week n-1, then she must have also bought a stand on week n-2. Before that, however, we are back to the default of only being able to purchase a Weapon or T-shirt on the last week. That means there are  $a_{n-3}$  ways to end with a Weapon on week n and a Stand on week n-1.

#### Case 2: T-shirt

If Raven buys a T-shirt on week n, she could actually buy anything at all the previous week. As with case 1, a recursive term would only deal with the T-shirt or Weapon options, but we will need other subcases for Stand and Figure

# Case 2a: T-shirt, T-shirt or Weapon

As with case 1a, this is a perfect setup to use  $a_{n-1}$ , as the requirements on week n-1 match the original week n.

### Case 2b: T-shirt, Stand

As with case 1b, this means Raven must buy a Stand on week n-2 as well, and week n-3 could not be a third Stand in a row, nor a Figure, so the remaining n-3 weeks can be decided by  $a_{n-3}$ 

### Case 2c: T-shirt, Figure

If Raven buys a T-shirt on week n and a Figure on week n-1 (which is valid because she did buy a T-shirt the next week), then the options for week n-2 are T-shirt, Weapon, or Stand. Obnoxiously enough, once again, a recursive call can only deal with the T-shirt and Weapon options, so we will need sub sub cases:

Case 2ci: T-shirt, Figure, Weapon or T-shirt

Similar to cases 1a and 2a, we have  $a_{n-2}$  ways to select what to buy.

#### Case 2cii: T-shirt, Figure, Stand

Similar to cases 1b and 2b, we know that she must buy a second Stand on week n-3, leaving the remaining n-4 weeks free to be anything (with the requirement that week n-4 be something that can come before a stand, which is only T-shirts or Weapons, just like we wane). This means there are  $a_{n-4}$  options here.

Adding all this up, we get exactly the same recurrence as the other solution:  $a_n = 2a_{n-1} + a_{n-2} + 2a_{n-3} + a_{n-4}$ 

(b) Regardless of which solution we use, the furthest back recurrence term is  $a_{n-4}$ , so we need 4 initial conditions:  $a_0, a_1, a_2, a_3$ .

# Grading Guidelines [14 points]

#### Part a (forward solution):

- +3 correctly handles T-shirt case
- +3 correctly handles Weapon case

- +3 correctly handles Figure case
- +3 correctly handles Stand case

# Part a (backward solution):

- +3 identifies correct cases (Weapon and T-shirt)
- +3 identifies correct sub-cases (1a, 1b, 2a, 2b, and 2c)
- +3 obtains correct terms from Case 1
- +3 obtains correct terms from Case 2

## Part b:

+2 correct initial conditions (does not need to provide values)

# Groupwork

# 1. Grade Groupwork 4

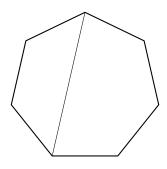
Using the solutions and Grading Guidelines, grade your Groupwork 4:

- Mark up your past groupwork and submit it with this one.
- Write whether your submission achieved each rubric item. If it didn't achieve one, say why not.
- Use the table below to calculate scores.
- For extra credit, write positive comment(s) about your work.
- You don't have to redo problems correctly, but it is recommended!
- What if my group changed?
  - If your current group submitted the same groupwork last time, grade it together.
  - If not, grade your version, which means submitting this groupwork assignment separately. You may discuss grading together.

	(i)	(ii)	(iii)	(iv)	(v)	(vi)	(vii)	(viii)	(ix)	(x)	(xi)	Total:
Problem 2												/20
Problem 3												/30
Total:												/50

# 2. Polly Gone [12 points]

A convex polygon is a 2D shape with all straight edges such that any line segment between any two non-adjacent vertices passes entirely through its interior (see example picture below). Show via induction that the sum of the interior angles of a convex polygon with n sides is  $(n-2) \cdot 180^{\circ}$ . Don't include unneeded base cases.



**Hint 1**: It is helpful to know that a triangle's interior angles always sum to 180°. You may assume this is true for the problem.

**Hint 2**: In order to apply your inductive hypothesis to a convex polygon, you'll need to think of it in terms of smaller convex polygons. How can you make smaller polygons out of a big one?

#### Solution:

**Base case:** A triangle's interior angles sum to  $(3-2)180^{\circ} = 180^{\circ}$ .

### Inductive step:

We claim that for k < n, the sum of the interior angles is  $(k-2)180^{\circ}$ . Now, we can note that any convex n-sided polygon can be divided into a triangle and a n-1-sided polygon i.e. for an n-sided polygon, we can take any three consecutive vertices and draw a line between the first and the third of those. The sum of the interior angles in the triangle and n-1-gon equal the sum the sum of the angles in the n-gon. From our inductive hypothesis, we know the triangle's interior angles sum to  $180^{\circ}$  and the n-1-gon's interior angles sum to  $((n-1)-2)180^{\circ}$ . Their sum is

$$180^{\circ} + ((n-1) - 2)180^{\circ} = 180^{\circ} + (n-3)180^{\circ}$$
$$= (1 + (n-3))180^{\circ}$$
$$= (n-2)180^{\circ}.$$

Therefore, by induction, we've shown that the sum of the interior angles of a convex n-gon is  $(n-2)180^{\circ}$ .

Alternatively, we can solve this with strong induction, where we divide the polygon along a line between any two non-consecutive vertices. For the inductive step, we divide the polygon into a (n-i+1)-gon and an (i+1)-gon for some  $2 \le i \le n-2$ . Since n-i+1 and i+1 are both less than n, by our inductive hypothesis the sum of the interior angles is

$$(n-i+1-2)180^{\circ} + (i+1-2)180^{\circ} = ((n-i+1-2)+(i+1-2))180^{\circ}$$
  
=  $(n-i+i+1+1-2-2)180^{\circ}$   
=  $(n-2)180^{\circ}$ .

# Draft Grading Guidelines [12 points]

- (i) +4 specifies triangle as the only correct base case
- (ii) +4 correct inductive hypothesis
- (iii) +4 correctly subdivides polygons based on induction type

# 3. Running Recurrence [8 points]

An EECS 203 student goes to lecture everyday. On each day, she always chooses exactly one method of transportation, and always chooses to walk, bike, or take a bus. She also follows additional rules:

- She never walks two days in a row.
- If she takes the bus, she must have biked two days ago and walked a day ago.

Let  $a_n$  denote the number of ways she can go to EECS 203 lecture across n days for  $n \geq 0$ .

- (a) Find a recurrence relation for  $a_n$ .
- (b) What are the initial conditions? Use the fewest initial conditions necessary.

#### Solution:

## (a) Solution: "Walking Backwards"

We consider the different cases for what method of transportation was chosen on the n-th day.

Case 1 Bike: If the student biked on the n-th day, there are no restrictions on the previous day. Because the last day is set and there are no restrictions on the (n-1)-th day, there are  $a_{n-1}$  ways for this case.

Case 2 Bus: If the student took a bus on the n-th day, she must have biked on the (n-2)-th day and walked on the (n-1)-th day. Because the last three days are set, and there are no restrictions for the (n-3)-th day, there are  $a_{n-3}$  ways for this case.

Case 3 Walk: If the student walked on the n-th day, she must have not walked the day before, and there are two cases for the (n-1)-th day.

- If she biked on the (n-1)-th day, there are no restrictions for the (n-2)-th day, and there are  $a_{n-2}$  ways for this case.
- If she took a bus on the (n-1)-th day, she must have biked on the (n-3)-th day and walked on the (n-2)-th day. Because the last 4 days are set, there are  $a_{n-4}$  ways for this case.

Therefore, if the student walked on the *n*-th day, there are  $a_{n-2} + a_{n-4}$  ways.

The recurrence is  $a_n = a_{n-1} + a_{n-2} + a_{n-3} + a_{n-4}$ .

(b) There is 1 way to go to lecture 0 times so  $a_0 = 1$ . There are two ways to go to lecture on the first day (bike, walk) so  $a_1 = 2$ . There are 3 ways to go to lecture by the second day (bike bike, bike walk, walk bike) so  $a_2 = 3$ . There are 6 ways to

go to lecture by the third day (bike bike bike, bike bike walk, bike walk bike, walk bike walk, walk bike bike) so  $a_3 = 6$ .

# Draft Grading Guidelines [8 points]

### Part a:

- (i) +2 correctly handles Case 1
- (ii) +2 correctly handles Case 2
- (iii) +2 correctly handles Case 3

### Part b:

- (iv) +1 correct number of initial conditions
- (v) +1 correct value of initial conditions