

- Obj: 1. Technical Vocab: logical equivalence; <sup>① 逆否</sup> contrapositive; <sup>② 永真命题, 重言式</sup> tautology; <sup>③ 矛盾</sup> contradiction; <sup>④ 谓词</sup> predicate, quantifier
2. Key Logical equivalence rules: DeMorgan's, distributive law, implication breakout.
3. the contrapositive of an if-then statement
4. Negate an if-then statement
5. A compound proposition is a tautology, a contradiction, or neither
6. Logical symbols for "there exists" and "for all" ( $\exists$ ,  $\forall$ )

9/5/23

② Double Negation law:  $\neg(\neg p) \equiv p$

### Lecture 3 Handout: Logical Equivalence

p = "I pet my cat" h = "my cat is happy"

Translate each English statement to logic, then complete the truth table for each.

1. If I pet my cat then she is happy.  $p \rightarrow h$
2. If my cat is unhappy then I didn't pet her.  $\neg h \rightarrow \neg p$
3. I didn't pet my cat or she is happy.  $\neg p \vee h$

p	h	$\neg p$	$\neg h$	$p \rightarrow h$	$\neg h \rightarrow \neg p$	$\neg p \vee h$
T	T	F	F	T	T	T
T	F	F	T	F	F	F
F	T	T	F	T	T	T
F	F	T	T	T	T	T

### Logical Equivalence

Two compound propositions A, B are **logically equivalent** if they have the **same truth value** for any instantiation of their input truth values (i.e., if they have the **same truth table**)

Some notations for logical equivalence:

- "p, q are logically equivalent"
- $p \equiv q$
- $p \leftrightarrow q$
- "p if and only if q" (充要命题 即为逻辑等价)
- "p is necessary and sufficient for q"
- "p iff q"

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### DeMorgans Laws in Action (Words)

Negate each of the following:

Proposition 1: I will go to the store or I will go to the park.  $s \vee p$   
 Negation: It's not true that I will go to S or I will go to P  
 Simplify: I will not go to S and I will not go to P.

③ DeMorgan's Law #1:  $\neg(s \vee p) \equiv \neg s \wedge \neg p$   
 (negating an or statement)

Proposition 2: I am 18 years old and I live in West Quad.  $y \wedge w$   
 Negation: It's not true that I am 18 and I live in West Quad.  
 Simplify: I am not 18 and I don't live in West Quad.

④ DeMorgan's Law #2:  $\neg(y \wedge w) \equiv \neg y \vee \neg w$   
 (negating an and statement)

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### Useful Logical Equivalence Rules

$p \vee F \equiv p$  ⑤ Identity Law  
 $p \vee T \equiv T$   
 $p \wedge F \equiv F$   
 $p \wedge T \equiv p$

Distributive Laws: ⑥ Distributive Law #1

$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$

$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$

⑦ Distributive Law #2

DeMorgan's Laws:

$\neg(p \vee q) \equiv \neg p \wedge \neg q$

$\neg(p \wedge q) \equiv \neg p \vee \neg q$

"Implication breakout" rule: ⑩

$p \rightarrow q \equiv \neg p \vee q$

Contrapositive: ⑨

$p \rightarrow q \equiv \neg q \rightarrow \neg p$

Negating an "implies": ⑪

$\neg(p \rightarrow q) \equiv p \wedge \neg q$

Which of the following ALWAYS has the same truth value as  $p \rightarrow q$ ?

- A) Converse:  $q \rightarrow p$   
 B) Inverse:  $\neg p \rightarrow \neg q$   
 C) Contrapositive:  $\neg q \rightarrow \neg p$

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⑧  $(p \vee q) \vee r = p \vee (q \vee r)$   
 $(p \wedge q) \wedge r = p \wedge (q \wedge r)$

## Contrapositives

Statement: "If p, then q"  $p \rightarrow q$

Contrapositive of statement: "If not q, then not p"  $\neg q \rightarrow \neg p$

$$p \rightarrow q \equiv \neg q \rightarrow \neg p$$

Find the contrapositive of each statement:

- If it's Tuesday, then we have EECS 203 class. (easy)
- If you don't live in Michigan, then you don't live in Ann Arbor.
- If *not* p, then q. (Here p, q can stand for any propositions.)

Negate both sides and switch the order

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Exercise:  
Negating an "implies" statement

$$\neg(p \rightarrow q) \equiv$$

Original statement:  
"If  $x^2$  is an integer, then  $x$  is an integer."  
 $p \rightarrow q$

Find the negation of the statement, and express it in logic and in English.

$$\neg(p \rightarrow q) \equiv \neg(\neg p \vee q) \Leftarrow \text{Implication breakdown}$$

$$\equiv p \wedge \neg q \Leftarrow \text{DeMorgan's}$$

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## Tautologies and Contradictions

A compound proposition is a **tautology** iff

(1) it's always true regardless of the input values

Example:  $p \vee \neg p$

A compound proposition is a **contradiction** iff

(2) it's always false regardless of the input value

Example:  $p \wedge \neg p$

Additional Exercises: Determine whether each statement is a tautology.

1.  $(\neg q \rightarrow \neg p) \rightarrow (p \rightarrow q)$

2.  $(\neg p \rightarrow F) \rightarrow p$

$$\equiv (p \rightarrow q) \rightarrow (q \rightarrow p)$$

$$\equiv T$$

$$\equiv (T \rightarrow p) \rightarrow p$$

$$\equiv T$$

$\therefore$  tautology

T	P	$T \rightarrow p$	$(T \rightarrow p) \rightarrow p$
T	T	T	T
T	F	F	T

$\therefore$  tautology

## Predicates & Quantifiers

A **predicate** is a statement with some variables unspecified (13)

Ex.  $P(x) = "x \text{ likes Bubble Tea}"$ ,  $Q(x) = "2x = 7"$

Quantifiers:

(14)  $\forall$  means "for all", "for each", "for every"  
(like a big chain of ANDs)

(15)  $\exists$  means "there exists"  
(like a big chain of ORs)

	True iff ...	False iff ...
$\forall x P(x)$	$P(x)$ is <u>true</u> for <u>every x</u> in the domain	$P(x)$ is <u>false</u> for <u>at least one x</u> in the domain
$\exists x P(x)$	$P(x)$ is <u>true</u> for <u>at least one x</u> in the domain	$P(x)$ is <u>false</u> for <u>every x</u> in the domain

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$P(x) = "x \text{ likes Bubble Tea}"$

$\forall x P(x)$ : Everyone likes ~

$\exists x P(x)$ : Someone likes ~

总结:

1. De Morgan's And/Or Laws:

$$\neg(p \wedge q) \equiv \neg p \vee \neg q$$

$$\neg(p \vee q) \equiv \neg p \wedge \neg q$$

2. Identity Laws

$$p \wedge T \equiv p$$

$$p \vee F \equiv p$$

3. Implication Breakout

$$p \rightarrow q \equiv \neg p \vee q$$

4. Commutative Laws

$$p \wedge q \equiv q \wedge p$$

$$p \vee q \equiv q \vee p$$

5. Associative Laws

$$p \wedge (q \wedge r) \equiv (p \wedge q) \wedge r$$

$$p \vee (q \vee r) \equiv (p \vee q) \vee r$$

6. Distributive Laws

$$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$$

$$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$$

7. The Contrapositive

$$p \rightarrow q \equiv \neg q \rightarrow \neg p$$

8. Tautology: always T

Contradiction: always F

satisfiable: can be T

Eg1: Is this a tautology?

$$\underline{[p \wedge (p \rightarrow q)] \rightarrow \neg q}$$

Sol 1. (Truth table)

Sol 2. There is only one situation where  $A \rightarrow B \equiv F$ ,  
that is,  $A \equiv T, B \equiv F$ .

Assume,  $\neg q \equiv F$ , then  $q \equiv T$

then no matter  $p \equiv T$  or  $p \equiv F$ ,  $(p \rightarrow q) \equiv T$

then if  $p \equiv T$ ,  $\neg p \equiv F$ , then  $[\neg p \wedge (p \rightarrow q)] \equiv F$ ,

$$\underline{[\neg p \wedge (p \rightarrow q)] \rightarrow \neg q \equiv (F \rightarrow F) \equiv T}$$

else if  $p \equiv F$ ,  $\neg p \equiv T$ , then  $[\neg p \wedge (p \rightarrow q)] \equiv T$ ,

Eg2: Is it a tautology?

$$\begin{aligned} \text{then } [\neg p \wedge (p \rightarrow q)] \rightarrow \neg q &\equiv (T \rightarrow F) \\ &\equiv F \\ \therefore \text{not tautology} \end{aligned}$$

$$b \quad [p \wedge (p \rightarrow q)] \rightarrow q$$

$$\equiv p \wedge (\neg p \vee q)$$

$$\equiv (p \wedge \neg p) \vee (p \wedge q)$$

$$F \vee (p \wedge q)$$

$$\equiv (p \wedge q)$$

$$\equiv T, \therefore \text{tautology}$$

Eg 3

$Q(x) = x+1 > 2x$ , domain: all integers

a)  $Q(-1) \equiv T$

b)  $Q(1) \equiv F$

c)  $\exists x Q(x) \equiv T$

d)  $\forall x Q(x) \equiv F$

e)  $\exists x \neg Q(x) \equiv T$

f)  $\forall x \neg Q(x) \equiv F$