

Five Message Handshake Project in Spin

Alexander Steen and Max Wisniewski

Institut für Informatik, FU Berlin

Abstract. The distributed algorithm for the mutual exclusion problem proposed by Suzuki and Kasami [1] is checked with the model checker *Spin*. We present a modeling for the algorithm in Promela, the properties we want to check for this algorithm and a short error analysis, why the second algorithm of Suzuki and Kasami does not work.

1 Introduction

In the 1980s, Ricard and Agrawala published a distributed algorithm the solves the mutual exclusion problem for networks of independent agents. This algorithm requires $2N - 2$ message exchanges per invocation, where N is the number of agents in the network [2]. In 1985, Suzuki and Kasami proposed a different algorithm which is supposed to solve the problem of mutual exclusion yet only requiring N messages per invocation. A modification of this algorithm only uses bounded counters and requires asymptotically N messages per invocation as the number of invocations grow. Section 2 gives an overview of the problem and the solution proposed by Suzuki and Kasami. The remainder of this document describes how those algorithms can be modeled in promela and processed by the spin model checker.

2 Problem / Algorithm

The proposed algorithm is a solution for the distributed mutual exclusion problem.

Given N processes that can only communicate over messages and that do not have shared memory and a critical section we want to find a protocol such that at any time at most one process works in the critical section.

As in the concurrent case we want this protocol to satisfy *mutual exclusion*, absence of starvation, fairness and no unnecessary delay. The first one says that not more than one process may enter the critical section at any time. The second one says that if some processes want to enter the critical section than one may succeed. If a process wants to enter the critical section he will eventually do so which is the statement of the third property. The last one says that a process may enter the critical section without the help of another process.

The algorithm of Suzuki and Kasami works abstract as follows

```

‘Remainder of Code‘
‘Enter ‘:
if !hasPrivilege
    inc(#Attempts)
    send_all(REQUEST)
    recv(PRIVILEGE, waiting, suc_attempts)
fi

‘Critical Section ‘

inc(#suc_attempt)
add_atempting(waiting);
if !empty(waiting)
    send(head(waiting), PRIVILEGE, tail(waiting), suc_attempts)
fi

```

if every process keeps track of the number of times a process wanted to enter the critical section and the one with the current privilege knows the successful attempts and a queue of processes that should enter next.

If a process wants to enter he first checks if he already has the privilege. If not he increments his attempt and sends to all processes that he would like to enter. The process with the privilege will take note of that and send the privilege to him or at least at him to the queue of processes that wants to enter.

When receiving a process increments the request counter for that process and the one with the privilege may send this if itself is not in the critical section right now.

The second algorithm introduces a upper bound to the attempt counter. If a maximum is reached it will be set to 0 again in every process list.

3 Modeling

To determine whether the algorithms mentioned in section 2 indeed appropriately solves the problem of mutual exclusion, we model these algorithms in promela and apply ltl model checking using spin. Section 3.1 justifies our choice of the spin verification suite; the following sections describe the model itself. Finally, section 3.8 displays an example simulation run of the algorithms.

3.1 Verification system

Since the algorithms described by Suzuki and Kasami use an imperative style it appears intuitive to us that a model in promela, which itself is imperative, is best suited. Additionally, promela comes with native support for process communication via channels. Since we are checking a distributed algorithm, we use the channels of promela to simulate the message exchange.

3.2 Global and local variables

We chose to model almost all of the variables of a process as a global N -size array of variables, one entry per process. Of course the i -th entry of each of these arrays is only used by process i . This is due to (1) debug reasons and (2) to verification reasons. Firstly, at some points we needed to check the system state during model simulation. Here, we can use the `-w` option of spin to output the system state. Secondly, this way we could use the variables to construct ltl formulas. The only process local variables are j and n which are used to store message variables.

3.3 Type definitions

Supplementary types used to model the algorithm are included in the header files of the model.

Queue Since the algorithms use a queue, we implemented a simple queue without any checks for overflows or underflows. The type consists of an array and two counters, length and head. We use a standard implementation.

```
typedef Queue {
    int q[M];
    short head = 0;
    short length = 0;
}
/* result = queue.empty() */
inline qEmpty(result, queue) {...}
/* queue.append(elem) */
inline qAppend(queue, elem) {...}
/* result = queue.poll() */
inline qPoll(result, queue) {...}
/* result = queue.contains(elem) */
inline qElem(result, queue, elem) {...}
```

The size of the used array is $M > N$, which should be OK because the program a process id added at most once per process to the queue. To add some padding, we choose $M = N^2$.

Array Since arrays are also used, we need to model arrays of arrays. This is simply done by the following definition:

```
typedef Array{
    int a[N] = -1;
}
```

Since the array is initialized by -1 in the algorithm, we set the initial value in the type definition itself.

3.4 Messages and Channels

There are three kinds of messages used by the algorithms, namely

REQUEST containing two integers j, n , where j is the id of the sending process and n is the number of previous requests by that process. This message is used to indicate that one wants to enter the critical section.

PRIVILEGE containing two fields Q, LN , where Q is a queue of process ids and LN is an array of integers that indicate which process has entered the critical section so far. This message can be considered a token that authorizes its carrier to enter the critical section.

REPLY containing no further information. This message is sent as a confirmation to request messages.

These message types can be modeled as an enumeration type, in promela:

```
mtype = {REQUEST, PRIVILIGE, REPLY}
```

In our model, all of these messages types are subsumed by one message type (mtype, int, int, Queue, Array). We use a mailbox system, where each message to a specific process is put in one channel, yielding the definition

```
chan mailbox[N] = [N] of {mtype, int, int, Queue, Array}
```

Another possibility to model the channel system is to create a $N \times N$ -matrix of channels, one channel for each communication pair. We prefer our choice over this possibility because it avoids unnecessary code to iterate over each of the $N - 1$ input channels per process. Also, the information which process sent a particular message is encoded into the message itself where needed. Another benefit is the reduced number of channels needed.

As one can see, two message types only require two attributes, the other message does not need any additionally attributes. In order to be able to use one channel for all message types, we fill in the unneeded fields of a message with dummy variables, i.e. variables which carry no essential information.

3.5 Request Messages

The procedure p2, as called in the paper, is supposed to be executed "indivisibly" whenever a *REQUEST* message is received. We chose to put the code of this procedure in an atomic statement (so that its executed indivisibly), the procedure itself resides as inline function in the header file. This is due to our modeling: Since it is not possible to interrupt a process whenever a message arrives, we chose to process all pending *REQUEST* messages before entering the entry code. Since there is a line in procedure p1 where a process waits on the *PRIVILEGE* message (i.e. authorization to enter the criticle section), we here chose to process pending request messages while waiting.

3.6 Send and Receive in Spin

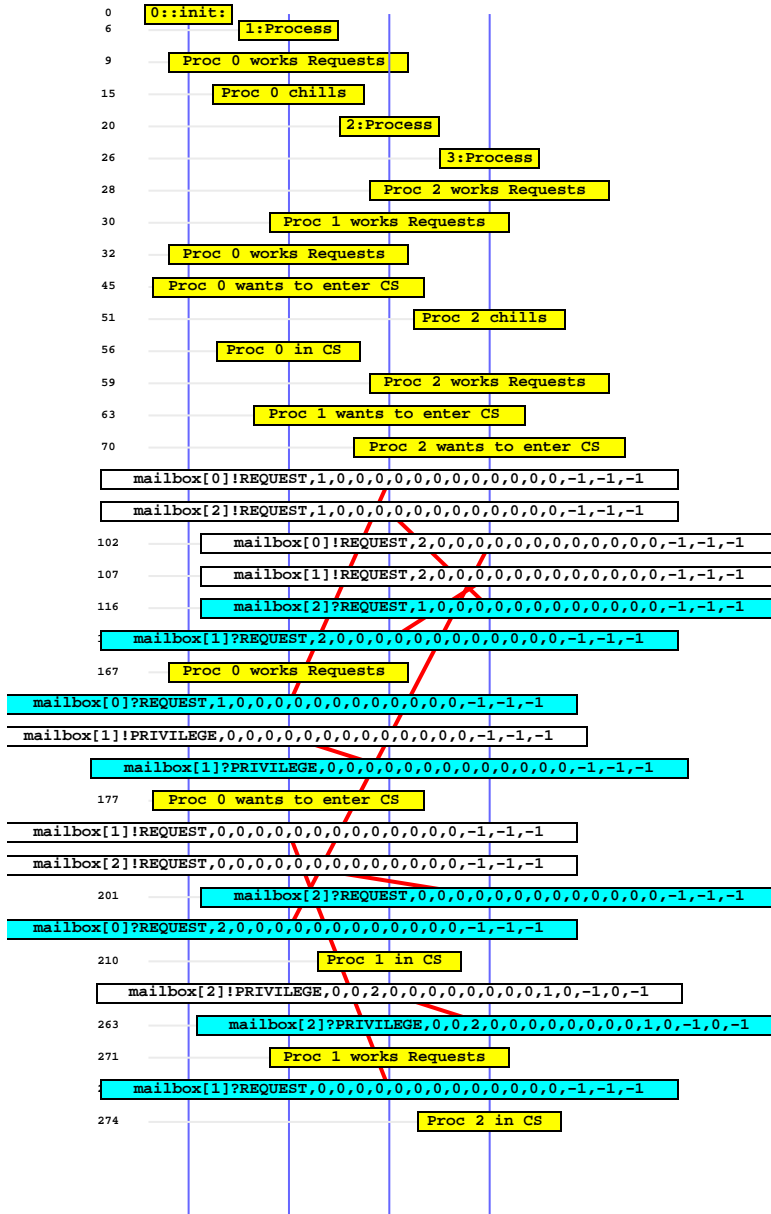
There occurred an ambiguous error in our implementation when we used a wrong number of matching variables in receiving a message. It happened some times that the received message differed from the send message. This way some of the requesting processes were dropped from the queue and were not considered for execution leading to a state where only one process was possible to enter the critical section.

3.7 Privilege Receiving

We implemented the procedure that is called on an incoming *PRIVILEGE* message as a macro. We decided to implement it this way because we need the procedure at two points in the code. First in the remainder of code and secondly when a process waits for the *PRIVILEGE* message. Because Spin does not support a procedure we implemented it as a macro.

3.8 Simulation of the model

The figure below shows an example run of the first algorithm. The receive and send events, plus additional debug output is printed.



4 LTL Properties

A Mutual Exclusion Algorithm needs to satisfy the four properties

- Mutual Exclusion,
- Absence of Starvation,
- Fairness,
- No Unnecessary Delay

to be considered as correct.

In Spin we have to model for each of these properties one or more LTL - Properties.

4.1 Mutual Exclusion

We added a variable `incs` that is incremented before the critical section and decremented afterwards. If initialized to zero mutual exclusion is expressed by the property

$$\Box (\text{incs} \leq 1) \quad (1)$$

which is used in both implementations.

4.2 Absence of Starvation

For this property we need a label request. We already have an array to keep track of the request, but this flag is shortly after the critical section still set to one. Using the counter `incs` from the Mutual Exclusion property we can express deadlock freedom by

$$\Box \left(\left(\bigvee_{0 \leq i < N} \text{Process}[\text{p}[i]] @\text{request} \right) \Rightarrow \Diamond \text{incs} = 1 \right) \quad (2)$$

again for both algorithms.

4.3 Fairness

All processes do not differ except for their Identifier. Therefore we will check the fairness constraint for the first process and the second process. The first one because it has initially the privilege and the second one as a representative for every other process. This time we used a label at the critical section and an array for the process id's.

Fairness can be expressed by

$$\Box (\text{Process}[\text{p}[0]] @\text{request} \Rightarrow \Diamond \text{Process}[\text{p}[0]] @\text{critical} \quad (3)$$

$$\wedge \text{Process}[\text{p}[1]] @\text{request} \Rightarrow \Diamond \text{Process}[\text{p}[1]] @\text{critical}) \quad (4)$$

in both algorithms.

4.4 No Unnecessary Delay

If we checked the previous property we have already a solution for this property. This is because it holds that each process that wants to enter the critical section will do so eventually for each path in the transition system. If we want to check the property – if a process wants to enter the critical section and other process arrives at the request section, the first one will eventually enter the critical section – this only describes a subset of the paths considered in the fairness property. Therefore if we have proven that one, we have a proof for this one.

5 Checking the Properties

We could see that the first algorithm has unbounded character. We let the verifier run for some time and never saw an error (see APPENDIX).

5.1 Unbounded Variant

In this section we checked the first, unbounded variant of the algorithm. Because it is unbounded the checking will never end. We therefore checked until we reached a certain depth in the search tree and ended it afterwards.

Mutual Exclusion We used the above described LTL property

```
ltl claim1 { [] (incs <= 1)}
```

and run the verifier via

```
gcc -DVECTORSZ=4096 -DCOLLAPSE pan.c -o pan  
./pan -a
```

The verifier gave the output

```
State-vector 2072 byte, depth reached 9999, errors: 0  
40921117 states, stored  
37895505 states, matched  
78816622 transitions (= stored+matched)  
56986035 atomic steps  
hash conflicts: 29156326 (resolved)
```

which at least shows, that until this depth there was no error. The full output can be seen in A.1.

Absence of Starvation We used the LTL property

```
ltl claim2 { [] (( Process[p[0]] @request  
|| Process[p[1]] @request || Process[p[2]] @request) -><> (incs == 1)) }
```

and run the verifier via


```
gcc -DVECTORSZ=4096 -DCOLLAPSE pan.c -o pan
./pan -a -f.
```

This time we need the parameter '-f'. Otherwise it is possible for one process to never to something after wanting to enter the critical section. The output can be seen in A.1 and it assures us that there is no error in the searched depth.

Fairness Here we checked the two claims

```
ltl claim3 { [] ( Process[p[0]]@request
-> <> ( Process[p[0]]@cs)) }
ltl claim4 { [] ( Process[p[0]]@request
-> <> ( Process[p[1]]@cs)) }
```

which both could not be proven wrong after the execution of the verifier as in the last section. The output can be seen in A.2.

5.2 Bounded Variant

This algorithm has a bounded state space and can, in principle, be fully checked.

Absence of Starvation The second algorithm cannot be correct as proposed in the paper. We have executed the algorithm several times and always reached a state in which could no longer be made any progress. Because with a liveness property we could not perform a BFS to this event, we constructed the error from the error traces we saw.

Consider three processes p1, p2, p3 where p1 initially holds the privilege. Now p1 and p2 enter the critical section swapping each round until both entered the critical section L times. It holds that $RN[0] = RN[1] = L - 1$ in both processes and let be p1 be the process that acquired the privilege last.

If p1 now leaves the critical section $RN[0] = L - 1$ holds and p1 waits for two REPLAY messages which he will get because he has sent two REQUEST messages while his requestCount was L . He receives those two messages and now attempts to enter the critical section again. Because no other process wants to enter and he still holds the privilege he can do so.

On leaving the critical section again $RN[0] = L - 1$ holds because RN is only changed if p1 does not have the privilege. He now waits for two REPLY messages which he will never receive because he never send REQUEST messages in the first place. Finally because he still holds the privilege no other process can enter the critical section and we reached a deadlock state.

Fairness and No Unnecessary Delay If we have an algorithm that contains a Deadlock (Lifelock in this case) then neither Fairness nor 'No Unnecessary Delay' can be satisfied because they both depend on the fact that it is possible to enter the critical section.

Mutual Exclusion We assume the algorithm satisfies at least mutual exclusion because in this matter it does not differ from the first algorithm. But Spin couldn't check the property because he could not perform the receive that led to the error mentioned before.

Hence we could not check this property successful because Spin can not do so if a receive statement runs into a deadlock state.

6 Conclusion

We modeled both algorithms of Suzuki and Kasami in promela. We were not able to verify either one of both models to be correct. The first, of course, contains unbounded counters and thus cannot be fully checked. Nevertheless, we were able to show that for a certain search depth, the first algorithm does not contain any reachable errors. Here, we were able to show that the first algorithm is unlikely to contain errors regarding safety and liveness properties. However, the second algorithm did trigger a number of errors during model checking. The only property we could not refute is the safety property, yet we were again not able to perform an exhaustive state search. We are convinced that the second algorithm contains a deadlock and thus does not solve the mutual exclusion problem.

References

1. Suzuki, I., Kasami, T.: A distributed mutual exclusion algorithm. *ACM Trans. Comput. Syst.* **3**(4) (November 1985) 344–349
2. Ricart, G., Agrawala, A.K.: An optimal algorithm for mutual exclusion in computer networks. *Commun. ACM* **24**(1) (January 1981) 9–17

Appendix

A Unbonded Algorithm Output

A.1 Mutual Exclusion

```
Alex@hildegunst ~  
$ ./pan -a  
error: max search depth too small  
Depth=    9999 States=    1e+06 Transitions= 1.92e+06 Memory=  
146.812 t=    7.21 R=    1e+05  
Depth=    9999 States=    2e+06 Transitions= 3.83e+06 Memory=  
228.453 t=    14.6 R=    1e+05  
Depth=    9999 States=    3e+06 Transitions= 5.74e+06 Memory=  
311.656 t=    22.4 R=    1e+05  
Depth=    9999 States=    4e+06 Transitions= 7.68e+06 Memory=  
402.281 t=    30.4 R=    1e+05  
Depth=    9999 States=    5e+06 Transitions= 9.63e+06 Memory=  
503.453 t=    38.3 R=    1e+05  
Depth=    9999 States=    6e+06 Transitions= 1.16e+07 Memory=  
582.750 t=    45.8 R=    1e+05  
Depth=    9999 States=    7e+06 Transitions= 1.35e+07 Memory=  
670.250 t=    53.1 R=    1e+05  
Depth=    9999 States=    8e+06 Transitions= 1.54e+07 Memory=  
744.859 t=    60.6 R=    1e+05  
Depth=    9999 States=    9e+06 Transitions= 1.73e+07 Memory=  
827.672 t=    68.3 R=    1e+05  
Depth=    9999 States=   1e+07 Transitions= 1.93e+07 Memory=  
914.000 t=    76.3 R=    1e+05  
Depth=    9999 States=   1.1e+07 Transitions= 2.12e+07 Memory=  
997.203 t=    83.9 R=    1e+05  
Depth=    9999 States=   1.2e+07 Transitions= 2.31e+07 Memory=  
1082.750 t=    91.5 R=    1e+05  
Depth=    9999 States=   1.3e+07 Transitions= 2.51e+07 Memory=  
1170.250 t=    99.2 R=    1e+05  
Depth=    9999 States=   1.4e+07 Transitions= 2.7e+07 Memory=  
1255.016 t=   107 R=    1e+05  
Depth=    9999 States=   1.5e+07 Transitions= 2.9e+07 Memory=  
1340.172 t=   115 R=    1e+05  
Depth=    9999 States=   1.6e+07 Transitions= 3.09e+07 Memory=  
1423.766 t=   123 R=    1e+05  
Depth=    9999 States=   1.7e+07 Transitions= 3.28e+07 Memory=  
1507.359 t=   131 R=    1e+05
```

```

Depth=    9999 States=  1.8e+07 Transitions= 3.47e+07 Memory=
1593.297 t=    139 R=    1e+05
Depth=    9999 States=  1.9e+07 Transitions= 3.67e+07 Memory=
1677.672 t=    146 R=    1e+05
Depth=    9999 States=  2e+07 Transitions= 3.86e+07 Memory=
1756.969 t=    154 R=    1e+05
Depth=    9999 States=  2.1e+07 Transitions= 4.05e+07 Memory=
1840.172 t=    162 R=    1e+05
Depth=    9999 States=  2.2e+07 Transitions= 4.25e+07 Memory=
1921.422 t=    170 R=    1e+05
Depth=    9999 States=  2.3e+07 Transitions= 4.44e+07 Memory=
2001.891 t=    178 R=    1e+05
Depth=    9999 States=  2.4e+07 Transitions= 4.63e+07 Memory=
2084.703 t=    186 R=    1e+05
Depth=    9999 States=  2.5e+07 Transitions= 4.83e+07 Memory=
2168.297 t=    194 R=    1e+05
Depth=    9999 States=  2.6e+07 Transitions= 5.02e+07 Memory=
2253.453 t=    202 R=    1e+05
Depth=    9999 States=  2.7e+07 Transitions= 5.21e+07 Memory=
2321.812 t=    210 R=    1e+05
Depth=    9999 States=  2.8e+07 Transitions= 5.4e+07 Memory=
2395.641 t=    218 R=    1e+05
Depth=    9999 States=  2.9e+07 Transitions= 5.59e+07 Memory=
2474.156 t=    226 R=    1e+05
Depth=    9999 States=  3e+07 Transitions= 5.79e+07 Memory=
2560.094 t=    234 R=    1e+05
Depth=    9999 States=  3.1e+07 Transitions= 5.98e+07 Memory=
2642.906 t=    242 R=    1e+05
Depth=    9999 States=  3.2e+07 Transitions= 6.17e+07 Memory=
2726.109 t=    251 R=    1e+05
Depth=    9999 States=  3.3e+07 Transitions= 6.36e+07 Memory=
2804.234 t=    259 R=    1e+05
Depth=    9999 States=  3.4e+07 Transitions= 6.55e+07 Memory=
2880.797 t=    267 R=    1e+05
pan: resizing hashtable to -w26.. done
Depth=    9999 States=  3.5e+07 Transitions= 6.74e+07 Memory=
3211.406 t=    282 R=    1e+05
Depth=    9999 States=  3.6e+07 Transitions= 6.94e+07 Memory=
3302.031 t=    290 R=    1e+05
Depth=    9999 States=  3.7e+07 Transitions= 7.13e+07 Memory=
3382.500 t=    298 R=    1e+05
tate-vector 2072 byte, depth reached 9999, errors: 0
  40921117 states, stored
  37895505 states, matched
  78816622 transitions (= stored+matched)

```

```

56986035 atomic steps
hash conflicts: 29156326 (resolved) State-vector 2072 byte,
    depth reached 9999, errors: 0
40921117 states, stored
37895505 states, matched
78816622 transitions (= stored+matched)
56986035 atomic steps
hash conflicts: 29156326 (resolved)
Depth= 9999 States= 3.9e+07 Transitions= 7.51e+07 Memory=
3546.172 t= 314 R= 1e+05
Depth= 9999 States= 4e+07 Transitions= 7.71e+07 Memory=
3636.016 t= 322 R= 1e+05
pan: out of memory

hint: to reduce memory, recompile with
    -DMA=2072 # better/slower compression, or
    -DBITSTATE # supertrace, approximation

(Spin Version 6.2.3 — 24 October 2012)
Warning: Search not completed
    + Partial Order Reduction
    + Compression

Full statespace search for:
    never claim + (claim1)
    assertion violations + (if within scope of claim)
    acceptance cycles + (fairness disabled)
    invalid end states - (disabled by never claim)

State-vector 2072 byte, depth reached 9999, errors: 0
40921117 states, stored
37895505 states, matched
78816622 transitions (= stored+matched)
56986035 atomic steps
hash conflicts: 29156326 (resolved)

Stats on memory usage (in Megabytes):
81641.175 equivalent memory usage for states
    (stored*(State-vector + overhead))
3442.784 actual memory usage for states (compression: 4.22%)
    state-vector as stored = 68 byte + 20 byte overhead
256.000 memory used for hash table (-w26)
0.382 memory used for DFS stack (-m10000)
1.840 memory lost to fragmentation
3697.344 total actual memory usage

```

```

nr of templates: [ 0:globals 1:chans 2:procs ]
collapse counts: [ 0:4015087 2:13 3:3852 4:1 ]

pan: elapsed time 329 seconds
pan: rate 124449.14 states/second

```

Absence of Starvation (

```

Alex@hildegunst ~
$ gcc -DVECTORSZ=4096 -DCOLLAPSE pan.c -o pan

Alex@hildegunst ~
$ ./pan -a -f
error: max search depth too small
Depth=    9999 States=    1e+06 Transitions= 2.55e+06 Memory=
95.250 t=    10.2 R=    1e+05
Depth=    9999 States=    2e+06 Transitions= 4.91e+06 Memory=
120.641 t=    19.2 R=    1e+05
Depth=    9999 States=    3e+06 Transitions= 7.25e+06 Memory=
145.250 t=    28.3 R=    1e+05
Depth=    9999 States=    4e+06 Transitions= 9.61e+06 Memory=
166.344 t=    37.5 R=    1e+05
Depth=    9999 States=    5e+06 Transitions= 1.2e+07 Memory=
194.859 t=    47.5 R=    1e+05
Depth=    9999 States=    6e+06 Transitions= 1.43e+07 Memory=
222.203 t=    57.8 R=    1e+05
Depth=    9999 States=    7e+06 Transitions= 1.67e+07 Memory=
246.031 t=    66.6 R=    1e+05
Depth=    9999 States=    8e+06 Transitions= 1.9e+07 Memory=
268.297 t=    75.6 R=    1e+05
Depth=    9999 States=    9e+06 Transitions= 2.13e+07 Memory=
289.000 t=    84.1 R=    1e+05
Depth=    9999 States=   1e+07 Transitions= 2.36e+07 Memory=
313.219 t=    93.6 R=    1e+05
Depth=    9999 States=   1.1e+07 Transitions= 2.6e+07 Memory=
333.141 t=   103 R=    1e+05
Depth=    9999 States=   1.2e+07 Transitions= 2.83e+07 Memory=
358.531 t=   113 R=    1e+05
Depth=    9999 States=   1.3e+07 Transitions= 3.08e+07 Memory=
383.141 t=   123 R=    1e+05
Depth=    9999 States=   1.4e+07 Transitions= 3.31e+07 Memory=
407.359 t=   131 R=    1e+05

```

Depth=	9999	States=	1.5e+07	Transitions=	3.54e+07	Memory=
430.406	t=	140	R=	1e+05		
Depth=	9999	States=	1.6e+07	Transitions=	3.77e+07	Memory=
453.453	t=	149	R=	1e+05		
Depth=	9999	States=	1.7e+07	Transitions=	4e+07	Memory=
476.500	t=	158	R=	1e+05		
Depth=	9999	States=	1.8e+07	Transitions=	4.24e+07	Memory=
497.984	t=	168	R=	1e+05		
Depth=	9999	States=	1.9e+07	Transitions=	4.48e+07	Memory=
520.641	t=	180	R=	1e+05		
Depth=	9999	States=	2e+07	Transitions=	4.71e+07	Memory=
544.859	t=	189	R=	1e+05		
Depth=	9999	States=	2.1e+07	Transitions=	4.94e+07	Memory=
568.688	t=	198	R=	1e+05		
Depth=	9999	States=	2.2e+07	Transitions=	5.18e+07	Memory=
592.906	t=	207	R=	1e+05		
Depth=	9999	States=	2.3e+07	Transitions=	5.4e+07	Memory=
615.172	t=	215	R=	1e+05		
Depth=	9999	States=	2.4e+07	Transitions=	5.64e+07	Memory=
640.172	t=	225	R=	1e+05		
Depth=	9999	States=	2.5e+07	Transitions=	5.88e+07	Memory=
658.922	t=	235	R=	1e+05		
Depth=	9999	States=	2.6e+07	Transitions=	6.11e+07	Memory=
681.969	t=	243	R=	1e+05		
Depth=	9999	States=	2.7e+07	Transitions=	6.34e+07	Memory=
704.234	t=	252	R=	1e+05		
Depth=	9999	States=	2.8e+07	Transitions=	6.58e+07	Memory=
730.406	t=	263	R=	1e+05		
Depth=	9999	States=	2.9e+07	Transitions=	6.81e+07	Memory=
752.281	t=	272	R=	1e+05		
Depth=	9999	States=	3e+07	Transitions=	7.05e+07	Memory=
775.328	t=	281	R=	1e+05		
Depth=	9999	States=	3.1e+07	Transitions=	7.28e+07	Memory=
795.641	t=	290	R=	1e+05		
Depth=	9999	States=	3.2e+07	Transitions=	7.51e+07	Memory=
817.125	t=	299	R=	1e+05		
Depth=	9999	States=	3.3e+07	Transitions=	7.74e+07	Memory=
837.438	t=	309	R=	1e+05		
Depth=	9999	States=	3.4e+07	Transitions=	7.97e+07	Memory=
861.656	t=	318	R=	1e+05		
pan: resizing hashtable to -w26.. done						
Depth=	9999	States=	3.5e+07	Transitions=	8.21e+07	Memory=
1125.859	t=	329	R=	1e+05		
Depth=	9999	States=	3.6e+07	Transitions=	8.44e+07	Memory=
1150.859	t=	339	R=	1e+05		

```

Depth=    9999 States=  3.7e+07 Transitions= 8.67e+07 Memory=
1168.047 t=      348 R=    1e+05
Depth=    9999 States=  3.8e+07 Transitions=  8.9e+07 Memory=
1189.531 t=      358 R=    1e+05
.....

Depth=    9999 States= 1.46e+08 Transitions= 3.43e+08 Memory=
3616.094 t= 1.41e+03 R=    1e+05
Depth=    9999 States= 1.47e+08 Transitions= 3.45e+08 Memory=
3641.094 t= 1.42e+03 R=    1e+05
Depth=    9999 States= 1.48e+08 Transitions= 3.47e+08 Memory=
3663.359 t= 1.43e+03 R=    1e+05
Depth=    9999 States= 1.49e+08 Transitions=  3.5e+08 Memory=
3683.281 t= 1.44e+03 R=    1e+05
pan: out of memory
hint: to reduce memory, recompile with
      -DMA=2084  # better/slower compression , or
      -DBITSTATE # supertrace , approximation

(Spin Version 6.2.3 — 24 October 2012)
Warning: Search not completed
        + Partial Order Reduction
        + Compression

Full statespace search for:
      never claim                + (claim2)
      assertion violations      + (if within scope of claim)
      acceptance  cycles       + (fairness enabled)
      invalid end states       - (disabled by never claim)

State-vector 2084 byte, depth reached 9999, errors: 0
  51599459 states , stored (1.49533e+08 visited)
  2.0147333e+08 states , matched
  3.5100596e+08 transitions (= visited+matched)
  3.592431e+08 atomic steps
hash conflicts: 41393391 (resolved)

Stats on memory usage (in Megabytes):
103535.902      equivalent memory usage for states
                 (stored*(State-vector + overhead))
  3441.959      actual memory usage for states (compression: 3.32%)
                 state-vector as stored = 50 byte + 20 byte overhead
   256.000      memory used for hash table (-w26)
    0.382       memory used for DFS stack (-m10000)
    1.403       memory lost to fragmentation

```



```

3696.953          total actual memory usage

nr of templates: [ 0:globals 1:chans 2:procs ]
collapse counts: [ 0:3040490 2:13 3:3918 4:2 ]

pan: elapsed time 1.44e+03 seconds
pan: rate 103649.56 states/second

```

A.2 Fairness

```

Alex@hildegunst ~
$ gcc -DVECTORSZ=4096 -DCOLLAPSE -DMEMLIM=2048 pan.c -o pan

Alex@hildegunst ~
$ ./pan -a -f
Depth=    9999 States=    1e+06 Transitions= 2.51e+06 Memory=
87.828 t=    9.83 R=    1e+05
Depth=    9999 States=    2e+06 Transitions= 4.96e+06 Memory=
108.922 t=    19.4 R=    1e+05
Depth=    9999 States=    3e+06 Transitions= 7.48e+06 Memory=
129.625 t=    29.5 R=    1e+05
Depth=    9999 States=    4e+06 Transitions= 9.92e+06 Memory=
150.328 t=     39 R=    1e+05
Depth=    9999 States=    5e+06 Transitions= 1.24e+07 Memory=
172.984 t=     49 R=    1e+05
Depth=    9999 States=    6e+06 Transitions= 1.49e+07 Memory=
197.594 t=    59.8 R=    1e+05
Depth=    9999 States=    7e+06 Transitions= 1.75e+07 Memory=
221.031 t=    71.3 R=    1e+05
Depth=    9999 States=    8e+06 Transitions= 2.01e+07 Memory=
243.688 t=    81.6 R=    1e+05
Depth=    9999 States=    9e+06 Transitions= 2.26e+07 Memory=
266.734 t=    91.8 R=    1e+05
Depth=    9999 States=   1e+07 Transitions= 2.51e+07 Memory=
288.609 t=   101 R=    1e+05
Depth=    9999 States=   1.1e+07 Transitions= 2.76e+07 Memory=
308.922 t=   112 R=    1e+05
Depth=    9999 States=   1.2e+07 Transitions= 3.01e+07 Memory=
332.359 t=   121 R=    1e+05
Depth=    9999 States=   1.3e+07 Transitions= 3.26e+07 Memory=
354.234 t=   131 R=    1e+05
Depth=    9999 States=   1.4e+07 Transitions= 3.51e+07 Memory=
377.672 t=   141 R=    1e+05

```

Depth=	9999	States=	1.5e+07	Transitions=	3.76e+07	Memory=	
402.672	t=	152	R=	1e+05			
Depth=	9999	States=	1.6e+07	Transitions=	4e+07	Memory=	
423.766	t=	161	R=	1e+05			
Depth=	9999	States=	1.7e+07	Transitions=	4.26e+07	Memory=	
444.469	t=	171	R=	1e+05			
Depth=	9999	States=	1.8e+07	Transitions=	4.51e+07	Memory=	
465.172	t=	181	R=	1e+05			
Depth=	9999	States=	1.9e+07	Transitions=	4.77e+07	Memory=	
485.484	t=	192	R=	1e+05			
Depth=	9999	States=	2e+07	Transitions=	5.02e+07	Memory=	
511.266	t=	203	R=	1e+05			
Depth=	9999	States=	2.1e+07	Transitions=	5.27e+07	Memory=	
532.750	t=	213	R=	1e+05			
Depth=	9999	States=	2.2e+07	Transitions=	5.52e+07	Memory=	
554.234	t=	223	R=	1e+05			
Depth=	9999	States=	2.3e+07	Transitions=	5.78e+07	Memory=	
572.594	t=	233	R=	1e+05			
Depth=	9999	States=	2.4e+07	Transitions=	6.03e+07	Memory=	
588.219	t=	244	R=	1e+05			
Depth=	9999	States=	2.5e+07	Transitions=	6.27e+07	Memory=	
607.750	t=	253	R=	1e+05			
Depth=	9999	States=	2.6e+07	Transitions=	6.52e+07	Memory=	
630.406	t=	264	R=	1e+05			
Depth=	9999	States=	2.7e+07	Transitions=	6.77e+07	Memory=	
651.109	t=	273	R=	1e+05			
Depth=	9999	States=	2.8e+07	Transitions=	7.02e+07	Memory=	
672.594	t=	284	R=	1e+05			
Depth=	9999	States=	2.9e+07	Transitions=	7.27e+07	Memory=	
694.078	t=	295	R=	1e+05			
Depth=	9999	States=	3e+07	Transitions=	7.52e+07	Memory=	
717.125	t=	306	R=	1e+05			
Depth=	9999	States=	3.1e+07	Transitions=	7.76e+07	Memory=	
736.656	t=	315	R=	1e+05			
Depth=	9999	States=	3.2e+07	Transitions=	8.01e+07	Memory=	
756.578	t=	326	R=	1e+05			
Depth=	9999	States=	3.3e+07	Transitions=	8.26e+07	Memory=	
776.891	t=	336	R=	1e+05			
Depth=	9999	States=	3.4e+07	Transitions=	8.52e+07	Memory=	
799.547	t=	347	R=	1e+05			
pan: resizing hashtable to -w26.. done							
Depth=	9999	States=	3.5e+07	Transitions=	8.78e+07	Memory=	
1058.281	t=	360	R=	1e+05			
....							

```

Depth=      9999 States=  1.1e+08 Transitions= 2.75e+08 Memory=
2727.812 t= 1.17e+03 R=    9e+04
Interrupted

(Spin Version 6.2.3 — 24 October 2012)
Warning: Search not completed
        + Partial Order Reduction
        + Compression

Full statespace search for:
        never claim          + (claim3)
        assertion violations  + (if within scope of claim)
        acceptance cycles    + (fairness enabled)
        invalid end states    - (disabled by never claim)

State-vector 2084 byte, depth reached 9999, errors: 0
 35000887 states, stored (1.10849e+08 visited)
1.6674613e+08 states, matched
2.7759498e+08 transitions (= visited+matched)
2.8812403e+08 atomic steps
hash conflicts: 22501231 (resolved)

Stats on memory usage (in Megabytes):
70230.356      equivalent memory usage for states
                (stored*(State-vector + overhead))
 2492.384      actual memory usage for states (compression: 3.55%)
                state-vector as stored = 55 byte + 20 byte overhead
   256.000     memory used for hash table (-w26)
    0.382      memory used for DFS stack (-m10000)
    1.435      memory lost to fragmentation
 2747.344      total actual memory usage

nr of templates: [ 0:globals 1:chans 2:procs ]
collapse counts: [ 0:2257676 2:13 3:4776 4:2 ]

pan: elapsed time 1.18e+03 seconds
pan: rate 93661.735 states/second

```