

## BRUTE FORCE

We will consider each pair of painters using two nested loops and count all the sections that can be painted by removing each pair of painters.

Out of all the pairs, we will pick one which will result in maximum painted sections.

## Time Complexity

$O(N * Q * Q)$ , where  $N$  is the length of the fence, and  $Q$  is the number of painters.

Since we are iterating for every pair of painters and then linearly checking the total sections that can be painted without them, the time complexity is  $O(N * Q * Q)$ .

## Space Complexity

$O(N)$ , where  $N$  is the length of the fence.

This is the space used by the '*section*' array, which will be used to calculate the number of painters that can paint the  $i$ -th section.