

Lets look at the code we are evaluating :

```
int i, j, k = 0;
for (i = n/2; i <= n; i++) {
    for (j = 2; j <= n; j = j * 2) {
        k = k + n/2;
    }
}
```

Now, lets just assume  $n = 8$  for now.

We will try to see, the values of j corresponding to each i.

```
i = 4, j = 2, 4, 8
i = 5, j = 2, 4, 8
i = 6, j = 2, 4, 8
i = 7, j = 2, 4, 8
i = 8, j = 2, 4, 8
```

If you notice, j keeps doubling till it is less than or equal to n. Number of times, you can double a number till it is less than n would be  $\log(n)$ .

Lets take more examples here to convince ourselves.

```
n = 16, j = 2, 4, 8, 16
n = 32, j = 2, 4, 8, 16, 32
```

So, j would run for  $O(\log n)$  steps.

i runs for  $n/2$  steps.

So, total steps ` =  $O(n/2 * \log(n)) = O(n \log n)$  `