

Parallel Quick Sort (Possible project task)

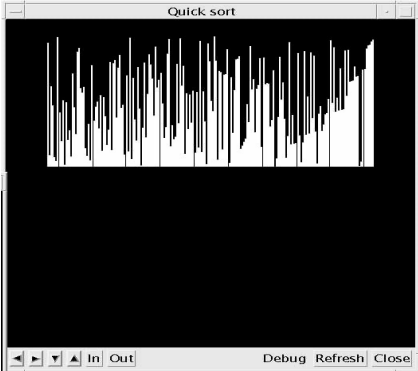
Algorithm:

1. Divide the data into p equal parts
2. Sort the data locally in each processor
3. Perform global sort
 - 3.1 Select pivot in each processor set
 - 3.2 In each processor, divide the data into two sets (smaller or larger)
 - 3.3 Split the processors into two groups and exchange data pair-wise
 - 3.4 Merge data into a sorted list in each processor
4. Repeat 3.1-3.4 recursively for each processor group

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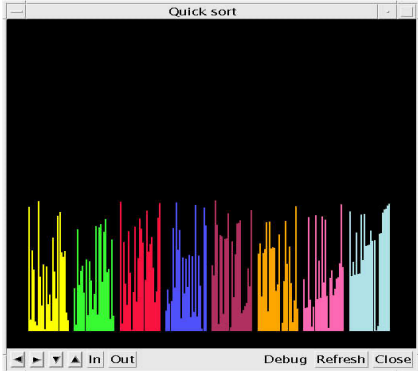
Unsorted data



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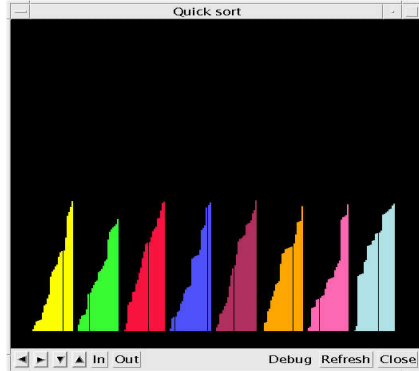
Step 1, Divide data into p equal parts



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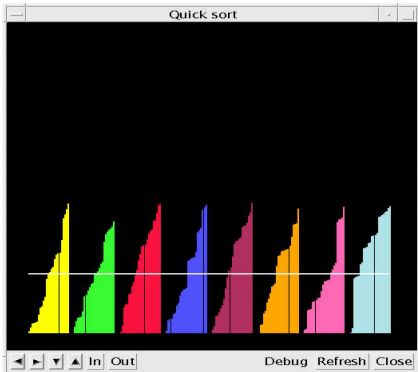
Step 2, Sort locally in each processor



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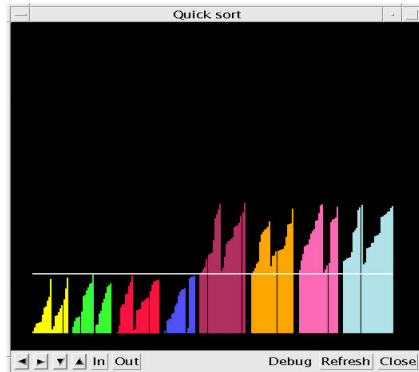
Step 3.1 Select pivot



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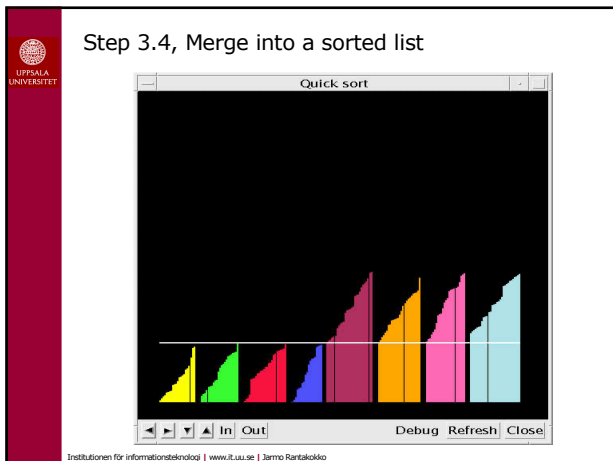
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Step 3.2, 3.3 Divide and exchange

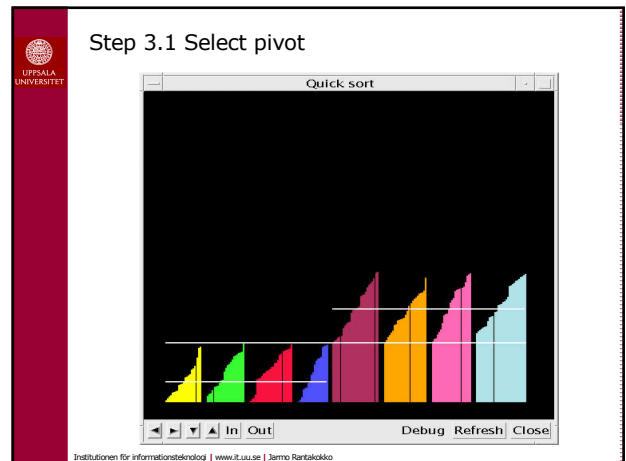


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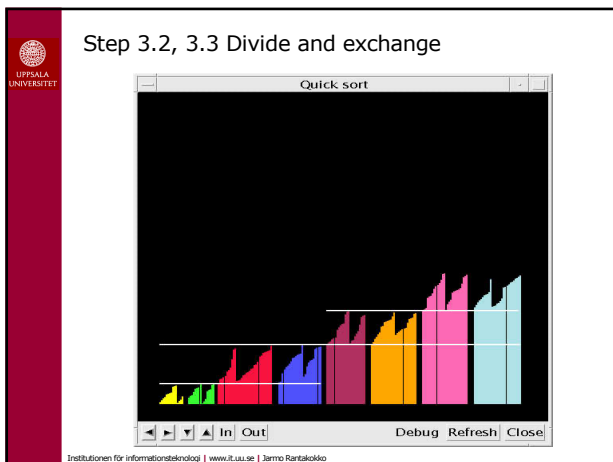
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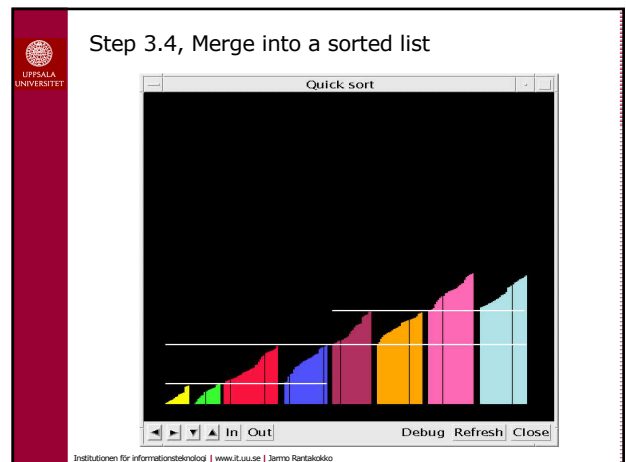
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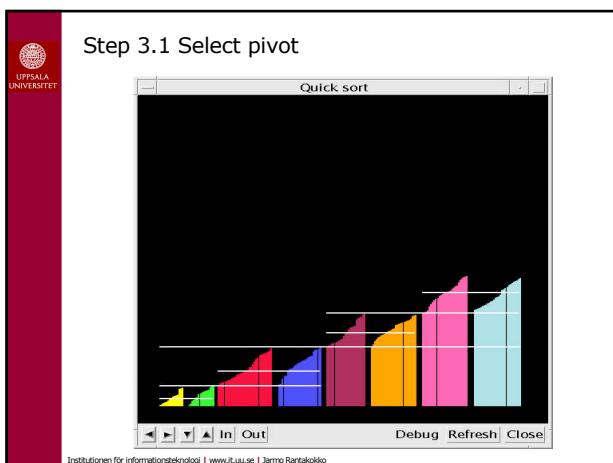
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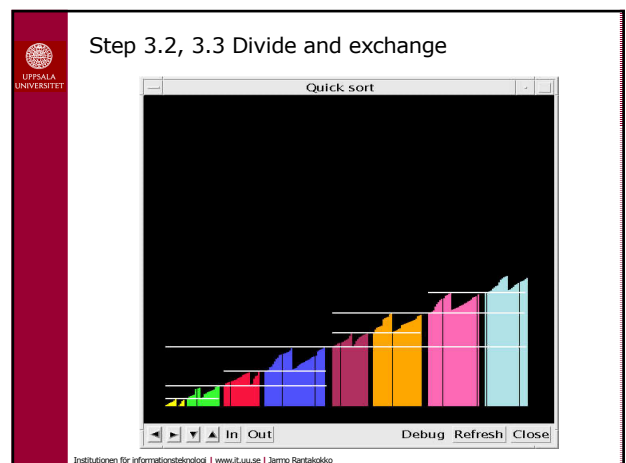
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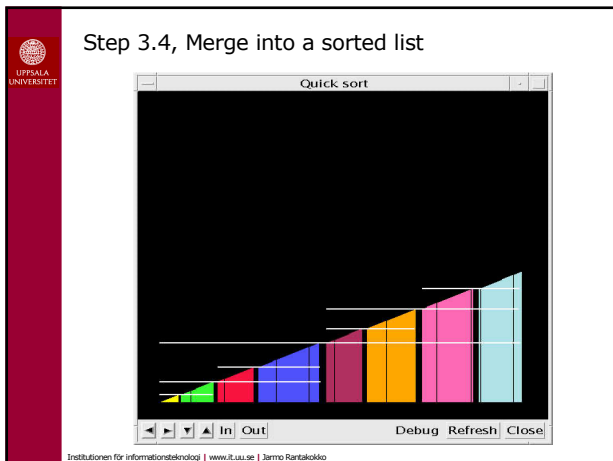
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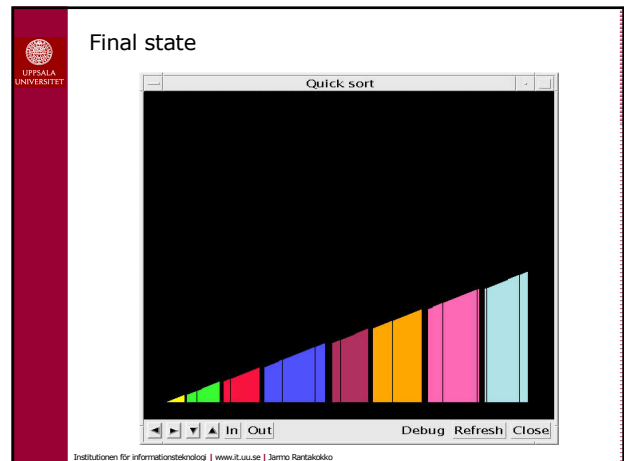
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Pseudo code (global sort):

```

void Global_sort(data,size,myid){
  if size==1 return
  locid=myid%size
  group=myid/size

  if locid==0 pivot[group]=select(data[myid])
  Synchronize group
  splitpoint=findsplit(data[myid],pivot[group])
  wait neighbor split
  if locid<size/2
    Merge(data[myid],data[myid+size/2],lowerparts)
  else
    Merge(data[myid],data[myid-size/2],upperparts)
  wait neighbor merge
  Global_sort(data,size/2,myid)
}

```

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Pivot selection strategies:

Strategy 1: Select median in processor 0 in each processor set (communicator) and step.
(OK if data equally rand, bad if almost sorted)

Strategy 2: Select the mean of all medians in respective processor set and step.
(Can give too much weight to extreme medians)

Strategy 3: Sort the medians and select the mean value of the two middlemost medians in each processor set and step.
(Independent of dist but more costly strategy)

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