9.1 Instruction Descriptions

ADDLW — Add Literal and W

Syntax: [labe] ADDLW k

Operands: $0 \le k \le 255$

Operation: $(W) + k \rightarrow (W)$

Status Affected: C, DC, Z

Encoding:

`11`	`111x`	`kkkk`	`kkkk`

Description: The contents of the W register are added to the eight bit literal 'k' and the result is placed in the W register.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process data	Write to W

Example: `ADDLW 0x15`

Before Instruction
W = 0x10
After Instruction
W = 0x25

ANDLW — AND Literal with W

Syntax: [labe] ANDLW k

Operands: $0 \le k \le 255$

Operation: (W) .AND. $(k) \rightarrow (W)$

Status Affected: Z

Encoding:

`11`	`1001`	`kkkk`	`kkkk`

Description: The contents of W register are AND'ed with the eight bit literal 'k'. The result is placed in the W register.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal "k"	Process data	Write to W

Example `ANDLW 0x5F`

Before Instruction
W = 0xA3
After Instruction
W = 0x03

ADDWF — Add W and f

Syntax: [labe] ADDWF f,d

Operands: $0 \le f \le 127$
$d \in [0,1]$

Operation: $(W) + (f) \rightarrow$ (destination)

Status Affected: C, DC, Z

Encoding:

`00`	`0111`	`dfff`	`ffff`

Description: Add the contents of the W register with the contents of register 'f'. If 'd' is 0 the result is stored in the W register. If 'd' is 1 the result is stored back in register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

Example `ADDWF FSR, 0`

Before Instruction
W = 0x17
FSR = 0xC2
After Instruction
W = 0xD9
FSR = 0xC2

ANDWF — AND W with f

Syntax: [labe] ANDWF f,d

Operands: $0 \le f \le 127$
$d \in [0,1]$

Operation: (W) .AND. $(f) \rightarrow$ (destination)

Status Affected: Z

Encoding:

`00`	`0101`	`dfff`	`ffff`

Description: AND the W register with contents of registe If 'd' is 0 the result is stored in the W register. If 'd' is 1 the result is stored back in register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

Example `ANDWF FSR, 1`

Before Instruction
W = 0x17
FSR = 0xC2
After Instruction
W = 0x17
FSR = 0x02

BCF — Bit Clear f

Syntax: [labe] BCF f,b

Operands: $0 \le f \le 127$
$0 \le b \le 7$

Operation: $0 \rightarrow (f<b>)$

Status Affected: None

Encoding:

`01`	`00bb`	`bfff`	`ffff`

Description: Bit 'b' in register 'f' is cleared.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write register 'f'

Example `BCF FLAG_REG, 7`

Before Instruction
FLAG_REG = 0xC7
After Instruction
FLAG_REG = 0x47

BSF — Bit Set f

Syntax: [labe] BSF f,b

Operands: $0 \le f \le 127$
$0 \le b \le 7$

Operation: $1 \rightarrow (f<b>)$

Status Affected: None

Encoding:

`01`	`01bb`	`bfff`	`ffff`

Description: Bit 'b' in register 'f' is set.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write register 'f'

Example `BSF FLAG_REG, 7`

Before Instruction
FLAG_REG = 0x0A
After Instruction
FLAG_REG = 0x8A

BTFSC — Bit Test, Skip if Clear

Syntax: [labe] BTFSC f,b

Operands: $0 \le f \le 127$
$0 \le b \le 7$

Operation: skip if $(f<b>) = 0$

Status Affected: None

Encoding:

`01`	`10bb`	`bfff`	`ffff`

Description: If bit 'b' in register 'f' is '1' then the next instruction is executed.
If bit 'b', in register 'f', is '0' then the next instruction is discarded, and a NOP is executed instead, making this a 2TCY instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	No-Operation

If Skip: (2nd Cycle)

Q1	Q2	Q3	Q4
No-Operation	No-Operation	No-Operation	No-Operation

Example

```
HERE    BTFSC  FLAG,1
FALSE   GOTO   PROCESS_CODE
TRUE    •
        •
        •
```

Before Instruction
PC = address `HERE`
After Instruction
if FLAG<1> = 0,
PC = address `TRUE`
if FLAG<1>=1,
PC = address `FALSE`

BTFSS — Bit Test f, Skip if Set

Syntax: [labe] BTFSS f,b

Operands: $0 \le f \le 127$
$0 \le b < 7$

Operation: skip if $(f<b>) = 1$

Status Affected: None

Encoding:

`01`	`11bb`	`bfff`	`ffff`

Description: If bit 'b' in register 'f' is '0' then the next instruction is executed.
If bit 'b' is '1', then the next instruction is discarded and a NOP is executed instead, making this a 2TCY instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	No-Operation

If Skip: (2nd Cycle)

Q1	Q2	Q3	Q4
No-Operation	No-Operation	No-Operation	No-Operation

Example

```
HERE    BTFSC  FLAG,1
FALSE   GOTO   PROCESS_CODE
TRUE    •
        •
        •
```

Before Instruction
PC = address `HERE`
After Instruction
if FLAG<1> = 0,
PC = address `FALSE`
if FLAG<1> = 1,
PC = address `TRUE`

CALL — Call Subroutine

Syntax: [label] CALL k

Operands: $0 \le k \le 2047$

Operation: $(PC)+ 1 \rightarrow$ TOS,
$k \rightarrow PC<10:0>$,
$(PCLATH<4:3>) \rightarrow PC<12:11>$

Status Affected: None

Encoding:

`10`	`0kkk`	`kkkk`	`kkkk`

Description: Call Subroutine. First, return address (PC+1) is pushed onto the stack. The eleven bit immediate address is loaded into PC bits <10:0>. The upper bits of the PC are loaded from PCLATH.`CALL` is a two cycle instruction.

Words: 1

Cycles: 2

Q Cycle Activity:

	Q1	Q2	Q3	Q4
1st Cycle	Decode	Read literal 'k', Push PC to Stack	Process data	Write to PC
2nd Cycle	No-Operation	No-Operation	No-Operation	No-Operation

Example `HERE CALL THERE`

Before Instruction
PC = Address `HERE`
After Instruction
PC = Address `THERE`
TOS = Address `HERE+1`

CLRF — Clear f

Syntax: [labe] CLRF f

Operands: $0 \le f \le 127$

Operation: $00h \rightarrow (f)$
$1 \rightarrow Z$

Status Affected: Z

Encoding:

`00`	`0001`	`1fff`	`ffff`

Description: The contents of register 'f' are cleared and the Z bit is set.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write register 'f'

Example `CLRF FLAG_REG`

Before Instruction
FLAG_REG = 0x5A
After Instruction
FLAG_REG = 0x00
Z = 1

CLRW — Clear W

Syntax: [label] CLRW

Operands: None

Operation: $00h \rightarrow (W)$
$1 \rightarrow Z$

Status Affected: Z

Encoding:

`00`	`0001`	`0xxx`	`xxxx`

Description: W register is cleared. Zero bit (Z) is set.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	No-Operation	Process data	Write to W

Example `CLRW`

Before Instruction
W = 0x5A
After Instruction
W = 0x00
Z = 1

CLRWDT — Clear Watchdog Timer

Syntax: [label] CLRWDT

Operands: None

Operation: $00h \rightarrow$ WDT
$0 \rightarrow$ WDT prescaler,
$1 \rightarrow \overline{TO}$
$1 \rightarrow \overline{PD}$

Status Affected: $\overline{TO}$, $\overline{PD}$

Encoding:

`00`	`0000`	`0110`	`0100`

Description: `CLRWDT` instruction resets the Watchdog Timer. It also resets the prescaler of the WDT. Status bits $\overline{TO}$ and $\overline{PD}$ are set.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	No-Operation	Process data	Clear WDT Counter

Example `CLRWDT`

Before Instruction
WDT counter = ?
After Instruction
WDT counter = 0x00
WDT prescaler= 0
$\overline{TO}$ = 1
$\overline{PD}$ = 1

COMF **Complement f**

Syntax: [label] COMF f,d

Operands: 0 ≤ f ≤ 127
d ∈ [0,1]

Operation: $(\bar{f}) \rightarrow$ (destination)

Status Affected: Z

Encoding:

00	1001	dfff	ffff

Description: The contents of register 'f' are complemented. If 'd' is 0 the result is stored in W. If 'd' is 1 the result is stored back in register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

Example

```
COMF      REG1,0
```

Before Instruction
REG1 = 0x13
After Instruction
REG1 = 0x13
W = 0xEC

DECF **Decrement f**

Syntax: [label] DECF f,d

Operands: 0 ≤ f ≤ 127
d ∈ [0,1]

Operation: (f) - 1 → (destination)

Status Affected: Z

Encoding:

00	0011	dfff	ffff

Description: Decrement contents of register 'f'. If 'd' is 0 the result is stored in the W register. If 'd' is 1 the result is stored back in register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

Example

```
DECF    CNT,  1
```

Before Instruction
CNT = 0x01
Z = 0
After Instruction
CNT = 0x00
Z = 1

DECFSZ **Decrement f, Skip if 0**

Syntax: [label] DECFSZ f,d

Operands: 0 ≤ f ≤ 127
d ∈ [0,1]

Operation: (f) - 1 → (destination);
skip if result = 0

Status Affected: None

Encoding:

00	1011	dfff	ffff

Description: The contents of register 'f' are decremented. If 'd' is 0 the result is placed in the W register. If 'd' is 1 the result is placed back in register 'f'.
If the result is not 0 ~~If the result is 1~~, the next instruction, is executed. If the result is 0, then a NOP is executed instead making it a 2TCY instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

If Skip: (2nd Cycle)

Q1	Q2	Q3	Q4
No-Operation	No-Operation	No-Operation	No-Operation

Example

```
HERE      DECFSZ   CNT, 1
          GOTO     LOOP
CONTINUE  •
          •
          •
```

Before Instruction
PC = address HERE
After Instruction
CNT = CNT - 1
if CNT = 0,
PC = address CONTINUE
if CNT ≠ 0,
PC = address HERE+1

GOTO **Unconditional Branch**

Syntax: [label] GOTO k

Operands: 0 ≤ k ≤ 2047

Operation: k → PC<10:0>
PCLATH<4:3> → PC<12:11>

Status Affected: None

Encoding:

10	1kkk	kkkk	kkkk

Description: GOTO is an unconditional branch. The eleven bit immediate value is loaded into PC bits <10:0>. The upper bits of PC are loaded from PCLATH<4:3>. GOTO is a two cycle instruction.

Words: 1

Cycles: 2

Q Cycle Activity:

	Q1	Q2	Q3	Q4
1st Cycle	Decode	Read literal 'k'	Process data	Write to PC
2nd Cycle	No-Operation	No-Operation	No-Operation	No-Operation

Example

```
GOTO THERE
```

After Instruction
PC = Address THERE

INCF **Increment f**

Syntax: [label] INCF f,d

Operands: 0 ≤ f ≤ 127
d ∈ [0,1]

Operation: (f) + 1 → (destination)

Status Affected: Z

Encoding:

00	1010	dfff	ffff

Description: The contents of register 'f' are incremented. If 'd' is 0 the result is placed in the W register. If 'd' is 1 the result is placed back in register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

Example

```
INCF    CNT,  1
```

Before Instruction
CNT = 0xFF
Z = 0
After Instruction
CNT = 0x00
Z = 1

INCFSZ **Increment f, Skip if 0**

Syntax: [label] INCFSZ f,d

Operands: 0 ≤ f ≤ 127
d ∈ [0,1]

Operation: (f) + 1 → (destination),
skip if result = 0

Status Affected: None

Encoding:

00	1111	dfff	ffff

Description: The contents of register 'f' are incremented. If 'd' is 0 the result is placed in the W register. If 'd' is 1 the result is placed back in register 'f'.
If the result is not 0 ~~If the result is 1~~, the next instruction is executed. If the result is 0, a NOP is executed instead making it a 2TCY instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

If Skip: (2nd Cycle)

Q1	Q2	Q3	Q4
No-Operation	No-Operation	No-Operation	No-Operation

Example

```
HERE      INCFSZ     CNT, 1
          GOTO       LOOP
CONTINUE  •
          •
          •
```

Before Instruction
PC = address HERE
After Instruction
CNT = CNT + 1
if CNT= 0,
PC = address CONTINUE
if CNT≠ 0,
PC = address HERE +1

IORLW **Inclusive OR Literal with W**

Syntax: [label] IORLW k

Operands: 0 ≤ k ≤ 255

Operation: (W) .OR. k → (W)

Status Affected: Z

Encoding:

11	1000	kkkk	kkkk

Description: The contents of the W register is OR'ed with the eight bit literal 'k'. The result is placed in the W register.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process data	Write to W

Example

```
IORLW   0x35
```

Before Instruction
W = 0x9A
After Instruction
W = 0xBF
Z = ~~1~~

Z = 0 !!!

IORWF **Inclusive OR W with f**

Syntax: [label] IORWF f,d

Operands: 0 ≤ f ≤ 127
d ∈ [0,1]

Operation: (W) .OR. (f) → (destination)

Status Affected: ~~$\overline{Z}$~~ Z !!!

Encoding:

00	0100	dfff	ffff

Description: Inclusive OR the W register with contents of register 'f'. If 'd' is 0 the result is placed in the W register. If 'd' is 1 the result is placed back in register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

Example

```
IORWF          RESULT, 0
```

Before Instruction
RESULT = 0x13
W = 0x91
After Instruction
RESULT = 0x13
W = 0x93
Z = ~~1~~

Z = 0 !!!

MOVF **Move f**

Syntax: [label] MOVF f,d

Operands: 0 ≤ f ≤ 127
d ∈ [0,1]

Operation: (f) → (destination)

Status Affected: Z

Encoding:

00	1000	dfff	ffff

Description: The contents of register f is moved to a destination dependant upon the status of d. If d = 0, destination is W register. If d = 1, the destination is file register f itself. d = 1 is useful to test a file register since status flag Z is affected.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

Example

```
MOVF    FSR,  0
```

After Instruction
W = value in FSR register
Z = 1

MOVLW **Move Literal to W**

Syntax: [label] MOVLW k

Operands: 0 ≤ k ≤ 255

Operation: k → (W)

Status Affected: None

Encoding:

11	00xx	kkkk	kkkk

Description: The eight bit literal 'k' is loaded into W register. The don't cares will assemble as 0's.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process data	Write to W

Example

```
MOVLW   0x5A
```

After Instruction
W = 0x5A

MOVWF **Move W to f**

Syntax: [label] MOVWF f

Operands: 0 ≤ f ≤ 127

Operation: (W) → (f)

Status Affected: None

Encoding:

00	0000	1fff	ffff

Description: Move data from W register to register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write register 'f'

Example

```
MOVWF    OPTION_REG
```

Before Instruction
OPTION = 0xFF
W = 0x4F
After Instruction
OPTION = 0x4F
W = 0x4F

NOP — No Operation

Syntax:	[label] NOP
Operands:	None
Operation:	No operation
Status Affected:	None
Encoding:	00 | 0000 | 0xx0 | 0000
Description:	No operation.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	No-Operation	No-Operation	No-Operation

Example `NOP`

RETFIE — Return from Interrupt

Syntax:	[label] RETFIE
Operands:	None
Operation:	TOS → PC, 1 → GIE
Status Affected:	None
Encoding:	00 | 0000 | 0000 | 1001
Description:	Return from Interrupt. Stack is POPed and Top of Stack (TOS) is loaded in the PC. Interrupts are enabled by setting Global Interrupt Enable bit, GIE (INTCON<7>). This is a two cycle instruction.
Words:	1
Cycles:	2

Q Cycle Activity:

	Q1	Q2	Q3	Q4
1st Cycle	Decode	No-Operation	Set the GIE bit	Pop from the Stack
2nd Cycle	No-Operation	No-Operation	No-Operation	No-Operation

Example `RETFIE`

After Interrupt
PC = TOS
GIE = 1

OPTION — Load Option Register

Syntax:	[label] OPTION
Operands:	None
Operation:	(W) → OPTION
Status Affected:	None
Encoding:	00 | 0000 | 0110 | 0010
Description:	The contents of the W register are loaded in the OPTION register. This instruction is supported for code compatibility with PIC16C5X products. Since OPTION is a readable/writable register, the user can directly address it.
Words:	1
Cycles:	1

Example

To maintain upward compatibility with future PIC16CXX products, do not use this instruction.

RETLW — Return with Literal in W

Syntax:	[label] RETLW k
Operands:	0 ≤ k ≤ 255
Operation:	k → (W); TOS → PC
Status Affected:	None
Encoding:	11 | 01xx | kkkk | kkkk
Description:	The W register is loaded with the eight bit literal 'k'. The program counter is loaded from the top of the stack (the return address). This is a two cycle instruction.
Words:	1
Cycles:	2

Q Cycle Activity:

	Q1	Q2	Q3	Q4
1st Cycle	Decode	Read literal 'k'	No-Operation	Write to W, Pop from the Stack
2nd Cycle	No-Operation	No-Operation	No-Operation	No-Operation

Example

```
      CALL TABLE  ;W contains table
                  ;offset value
      •            ;W now has table value
      •
      •
TABLE ADDWF PC   ;W = offset
      RETLW k1   ;Begin table
      RETLW k2   ;
      •
      •
      •
      RETLW kn   ; End of table
```

Before Instruction
W = 0x07
After Instruction
W = value of k8

RETURN — Return from Subroutine

Syntax:	[label] RETURN
Operands:	None
Operation:	TOS → PC
Status Affected:	None
Encoding:	00 | 0000 | 0000 | 1000
Description:	Return from subroutine. The stack is POPed and the top of the stack (TOS) is loaded into the program counter. This is a two cycle instruction.
Words:	1
Cycles:	2

Q Cycle Activity:

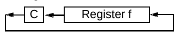
	Q1	Q2	Q3	Q4
1st Cycle	Decode	No-Operation	No-Operation	Pop from the Stack
2nd Cycle	No-Operation	No-Operation	No-Operation	No-Operation

Example `RETURN`

After Interrupt
PC = TOS

RLF — Rotate Left f through Carry

Syntax:	[label] RLF f,d
Operands:	0 ≤ f ≤ 127 d ∈ [0,1]
Operation:	See description below
Status Affected:	C
Encoding:	00 | 1101 | dfff | ffff
Description:	The contents of register 'f' are rotated one bit to the left through the Carry Flag. If 'd' is 0 the result is placed in the W register. If 'd' is 1 the result is stored back in register 'f'.



Words:	1
Cycles:	1

Q Cycle Activity:

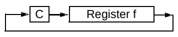
Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

Example `RLF REG1,0`

Before Instruction
REG1 = 1110 0110
C = 0
After Instruction
REG1 = 1110 0110
W = 1100 1100
C = 1

RRF — Rotate Right f through Carry

Syntax:	[label] RRF f,d
Operands:	0 ≤ f ≤ 127 d ∈ [0,1]
Operation:	See description below
Status Affected:	C
Encoding:	00 | 1100 | dfff | ffff
Description:	The contents of register 'f' are rotated one bit to the right through the Carry Flag. If 'd' is 0 the result is placed in the W register. If 'd' is 1 the result is placed back in register 'f'.



Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

Example `RRF REG1,0`

Before Instruction
REG1 = 1110 0110
C = 0
After Instruction
REG1 = 1110 0110
W = 0111 0011
C = 0

SLEEP

Syntax:	[label] SLEEP
Operands:	None
Operation:	00h → WDT, 0 → WDT prescaler, 1 → $\overline{TO}$, 0 → $\overline{PD}$
Status Affected:	$\overline{TO}$, $\overline{PD}$
Encoding:	00 | 0000 | 0110 | 0011
Description:	The power-down status bit, $\overline{PD}$ is cleared. Time-out status bit, $\overline{TO}$ is set. Watchdog Timer and its prescaler are cleared. The processor is put into SLEEP mode with the oscillator stopped. See Section 14.8 for more details.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	No-Operation	No-Operation	Go to Sleep

Example: `SLEEP`

SUBLW — Subtract W from Literal

Syntax:	[label] SUBLW k
Operands:	0 ≤ k ≤ 255
Operation:	k - (W) → (W)
Status Affected:	C, DC, Z
Encoding:	11 | 110x | kkkk | kkkk
Description:	The contents of W register is subtracted (2's compl ment method) from the eight bit literal 'k'. The result is placed in the W register.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process data	Write to W

Example 1: `SUBLW 0x02`

Before Instruction
W = 1
C = ?
Z = ?

After Instruction
W = 1
C = 1; result is positive
Z = 0

Example 2: Before Instruction
W = 2
C = ?
Z = ?

After Instruction
W = 0
C = 1; result is zero
Z = 1

Example 3: Before Instruction
W = 3
C = ?
Z = ?

After Instruction
W = 0xFF
C = 0; result is negative
Z = 0

SUBWF **Subtract W from f**

Syntax: [label] SUBWF f,d

Operands: 0 ≤ f ≤ 127
d ∈ [0,1]

Operation: (f) - (W) → (destination)

Status Affected: C, DC, Z

Encoding:

00	0010	dfff	ffff

Description: Subtract (2's complement method) contents of W register from register 'f'. If 'd' is 0 the result is stored in the W register. If 'd' is 1 the result is stored back in register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

Example 1: `SUBWF REG1,1`

Before Instruction

REG1	=	3
W	=	2
C	=	?
Z	=	?

After Instruction

REG1	=	1
W	=	2
C	=	1; result is positive
Z	=	0

Example 2: Before Instruction

REG1	=	2
W	=	2
C	=	?
Z	=	?

After Instruction

REG1	=	0
W	=	2
C	=	1; result is zero
Z	=	1

Example 3: Before Instruction

REG1	=	1
W	=	2
C	=	?
Z	=	?

After Instruction

REG1	=	0xFF
W	=	2
C	=	0; result is negative
Z	=	0

SWAPF **Swap Nibbles in f**

Syntax: [label] SWAPF f,d

Operands: 0 ≤ f ≤ 127
d ∈ [0,1]

Operation: (f<3:0>) → (destination<7:4>),
(f<7:4>) → (destination<3:0>)

Status Affected: None

Encoding:

00	1110	dfff	ffff

Description: The upper and lower nibbles of contents of register 'f' are exchanged. If 'd' is 0 the res is placed in W register. If 'd' is 1 the result is placed in register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

Example `SWAPF REG, 0`

Before Instruction

REG1	=	0xA5

After Instruction

REG1	=	0xA5
W	=	0x5A

TRIS **Load TRIS Register**

Syntax: [label] TRIS f

Operands: 5 ≤ f ≤ 7

Operation: (W) → TRIS register f;

Status Affected: None

Encoding:

00	0000	0110	0fff

Description: The instruction is supported for code compatibility with the PIC16C5X products. Since TRIS registers are readable and writable, the user can directly address them.

Words: 1

Cycles: 1

Example

To maintain upward compatibility with future PIC16CXX products, do not use this instruction.

XORLW **Exclusive OR Literal with W**

Syntax: [label] XORLW k

Operands: 0 ≤ k ≤ 255

Operation: (W) .XOR. k → (W)

Status Affected: Z

Encoding:

11	1010	kkkk	kkkk

Description: The contents of the W register are XOR'ed with the eight bit literal 'k'. The result is placed in the W register.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process data	Write to W

Example: `XORLW 0xAF`

Before Instruction

W	=	0xB5

After Instruction

W	=	0x1A

XORWF **Exclusive OR W with f**

Syntax: [label] XORWF f,d

Operands: 0 ≤ f ≤ 127
d ∈ [0,1]

Operation: (W) .XOR. (f) → (destination)

Status Affected: Z

Encoding:

00	0110	dfff	ffff

Description: Exclusive OR the contents of the W register with contents of register 'f'. If 'd' is 0 the result is stored in the W register. If 'd' is 1 the result is stored back in register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

Example `XORWF REG 1`

Before Instruction

REG	=	0xAF
W	=	0xB5

After Instruction

REG	=	0x1A
W	=	0xB5