

Problem: - 1

Solution of Matrix eqn with Gaussian Elimination

▣ The given Linear System is:

$$\begin{pmatrix} 4 & 1 & 2 \\ 2 & 4 & 7 \\ 1 & 1 & -3 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 9 \\ -5 \\ -9 \end{pmatrix}$$

In a two bit floating point decimal computer when the machine stores data in mantissa and exponent form; the same eq reads:

$$\begin{pmatrix} 0.4 \times 10^1 & 0.1 \times 10^1 & 0.2 \times 10^1 \\ 0.2 \times 10^1 & 0.4 \times 10^1 & -0.1 \times 10^1 \\ 0.1 \times 10^1 & 0.1 \times 10^1 & -0.3 \times 10^1 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 0.9 \times 10^1 \\ -0.5 \times 10^1 \\ -0.9 \times 10^1 \end{pmatrix}$$

(There is '0' at 2nd decimal for each mantissa which I've not written down.)

Now start Gaussian elimination.

$$\text{Step: 1} \rightarrow R_1 \rightarrow \frac{R_1}{0.4 \times 10^1}$$

$$\begin{pmatrix} 0.10 \times 10^1 & 0.25 \times 10^0 & 0.50 \times 10^0 \\ 0.20 \times 10^1 & 0.40 \times 10^1 & 0.10 \times 10^1 \\ 0.10 \times 10^1 & 0.10 \times 10^1 & -0.30 \times 10^1 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 0.23 \times 10^1 \\ -0.50 \times 10^1 \\ -0.90 \times 10^1 \end{pmatrix}$$

$$\text{Step 2: } R_2 \rightarrow R_2 - 0.2 \times 10^1 \times R_1; R_3 \rightarrow R_3 - R_1$$

(In machine the steps successively are

$$i) R_2 \rightarrow R_2 - 0.2 \times 10^1 \times R_1$$

$$ii) R_3 \rightarrow R_3 - 0.1 \times 10^1 \times R_1$$

I've written them in one step)

$$\begin{pmatrix} 0.10 \times 10^1 & 0.25 \times 10^0 & 0.50 \times 10^0 \\ 0 & 0.35 \times 10^1 & -0.2 \times 10^1 \\ 0 & 0.75 \times 10^0 & -0.35 \times 10^1 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 0.23 \times 10^1 \\ -0.96 \times 10^1 \\ -0.11 \times 10^2 \end{pmatrix}$$

Step 3:  $R_2 \rightarrow \frac{R_2}{A_{22}}$  i.e.  $R_2 \rightarrow \frac{R_2}{0.35 \times 10^1}$

$$\begin{pmatrix} 0.10 \times 10^1 & 0.25 \times 10^0 & 0.50 \times 10^0 \\ 0 & 0.10 \times 10^1 & -0.57 \times 10^0 \\ 0 & 0.75 \times 10^0 & -0.35 \times 10^1 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 0.23 \times 10^1 \\ -0.27 \times 10^1 \\ -0.11 \times 10^2 \end{pmatrix}$$

Step: 4 ~~to~~:  $R_3 \rightarrow R_3 - 0.75 \times 10^0 \times R_2$

$$\begin{pmatrix} 0.10 \times 10^1 & 0.25 \times 10^0 & 0.50 \times 10^0 \\ 0 & 0.10 \times 10^1 & -0.57 \times 10^0 \\ 0 & 0 & -0.31 \times 10^1 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 0.23 \times 10^1 \\ -0.27 \times 10^1 \\ -0.9 \times 10^1 \end{pmatrix}$$

Back Substitution:

$$x_3 = - \frac{0.9 \times 10^1}{-0.31 \times 10^1} = 0.29 \times 10^1 = 2.9$$

$$x_2 - 0.57 \times 0.29 \times 10^1 = -0.27 \times 10^1$$

$$\Rightarrow x_2 = -0.10 \times 10^1 = -1.0$$

$$x_1 + 0.25 \times (-0.10 \times 10^1) + (0.5 \times 10^0 \times 0.29 \times 10^1) = 0.23 \times 10^1$$

$$\Rightarrow x_1 = 0.11 \times 10^1$$

So the solution is:  $(0.11 \times 10^1, -0.10 \times 10^1, 0.29 \times 10^1)$

i.e.  $(1.1, -1.0, 2.9)$  Ans

Problem: 3

## SAGAR DAM

For the operator eq  $Ax = b$

$$x = \begin{pmatrix} x_1 \\ \vdots \\ x_m \end{pmatrix}; b = \begin{pmatrix} b_1 \\ \vdots \\ b_m \end{pmatrix}$$

If  $A$  is of tridiagonal form i.e. only the diagonal, Superdiagonal & Subdiagonal elements are non zero.

$$\therefore A = \begin{pmatrix} a_{11} & a_{12} & 0 & 0 & 0 & \dots & 0 \\ a_{21} & a_{22} & a_{23} & 0 & 0 & \dots & 0 \\ 0 & a_{32} & a_{33} & a_{34} & \dots & \dots & 0 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & \dots & \dots & \dots & \dots & a_{m-1,m} & a_{mm} \end{pmatrix}$$

To solve eq  $Ax = b$  the Augmented matrix be

$$\tilde{A} = \left( \begin{array}{ccccccccc|c} a_{11} & a_{12} & 0 & 0 & 0 & \dots & 0 & b_1 \\ a_{21} & a_{22} & a_{23} & 0 & 0 & \dots & 0 & b_2 \\ 0 & a_{32} & a_{33} & a_{34} & 0 & \dots & 0 & \vdots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & \dots & \dots & \dots & \dots & a_{m-1,m} & a_{mm} & b_m \end{array} \right)$$

Now we apply Gaussian elimination.

$$\text{Step 1: } R_1 \rightarrow \frac{R_1}{a_{11}}; R_2 \rightarrow R_2 - a_{21} R_1$$



$$\tilde{A} = \left( \begin{array}{cccccc|c} 1 & \frac{a_{12}}{a_{11}} & 0 & 0 & \dots & 0 & b_1/a_{11} \\ 0 & a_{22} - \frac{a_{21} \cdot a_{12}}{a_{11}} & a_{23} - \frac{a_{21} \cdot a_{13}}{a_{11}} & \dots & 0 & 0 & b_2 - \frac{b_1}{a_{11}} \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots \\ 0 & 0 & \dots & \dots & \dots & 0 & b_m \end{array} \right)$$

Now the  $(n-1) \times (n-1)$  lower block is tridiagonal form

Applying the same rule  $n$  times for which in

$j$ th step  $R_j \rightarrow \frac{R_j}{a_{jj}}$ ;  $R_{j+1} \rightarrow R_{j+1} - a_{j,j+1} \times R_j$   
gives the ultimate Augmented matrix to be upper triangular form

$$\tilde{A}_{\text{ult}} = \left( \begin{array}{cccccc|c} 1 & a'_{12} & 0 & 0 & \dots & 0 & b'_1 \\ 0 & 1 & a'_{23} & 0 & \dots & 0 & b'_2 \\ 0 & 0 & 1 & a'_{34} & \dots & 0 & \vdots \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \dots & \dots & \dots & 1 & b'_m \end{array} \right)$$

in A part

Now to do this job we need 3 multiplication & 3 subtraction in any one of  $n$  steps.

(except  $j=1, n$  need 2). All other terms are zero.

So at  $j$ th step total arithmetic operations are  $3+3=6$ .

As there are  $n$  steps so total operations  
=  $6n$ .

but for  $j=1, n$  there are two multiplications & 2 subtractions.

$\therefore$  no. of steps till now  $N_1 = 6m - 2$ .

for the right side there are (in b part)  
( $m-1$ ) divisions & ( $m-1$ ) subtractions.

$$\therefore N_1 \rightarrow N_1 + 2(m-1) = 8m - 4.$$

For the back substitution (Now with the form we can do that) we need at

$i = 1$  th step  $\rightarrow$  no operation

all other step  $\rightarrow$  1 multiplication + 1 subtraction  
 $\underbrace{\hspace{10em}}_{m-1 \text{ steps}} \quad \underbrace{\hspace{10em}}_{2 \text{ operation.}}$

So total no. of ~~step~~ operations:

$$N_1 \rightarrow N = N_1 + 2(m-1) \\ = 10m - 6$$

i.e. at  $m \rightarrow \infty$  limit the complexity asymptotically goes to  $O(m)$  limit.

$\therefore$  For tridiagonal operator eq; the computational complexity varies as  
 $\sim O(m)$ .

Proved