

## **Centroid Decomposition**

Special class

# Centroid Decomposition

Course: https://unacademy.com/a/i-p-c-advanced-track

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## Objective

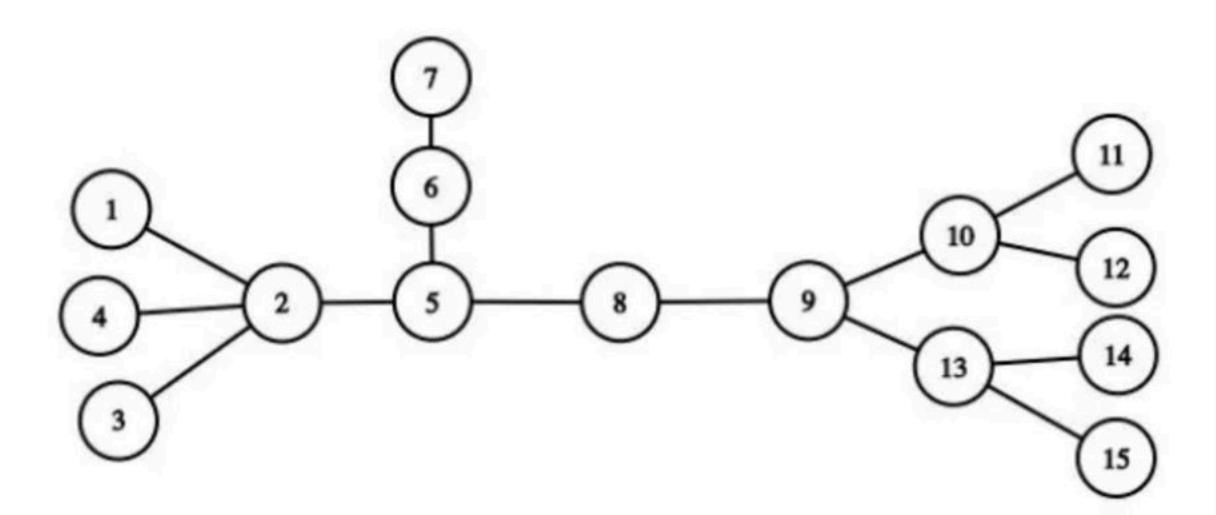
- Centroid of a Tree
  - Definition
  - How to find the centroid of a tree
- Decomposing the "Original Tree" to get "Centroid Tree"
  - Implementation
  - Visualization
- Properties of the Centroid Tree
  - Show any path A -- B in the tree can be written as A -- Centroid -- B.
  - Maintain some information for the O(NlogN) paths to answer generic path Queries
- Problem Discussion

#### Centroid of a Tree

 Centroid: Given a tree with N nodes, a centroid is a node whose removal splits the given tree into a forest of trees, where each of the resulting tree contains no more than N/2 nodes.

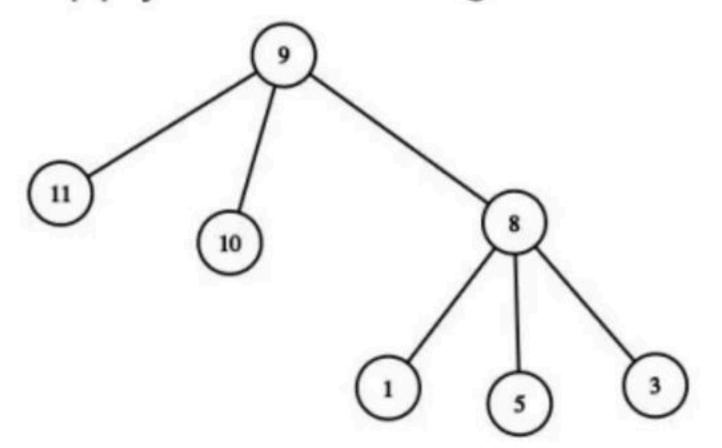
Question: Which node is the centroid of this tree (on the right)?

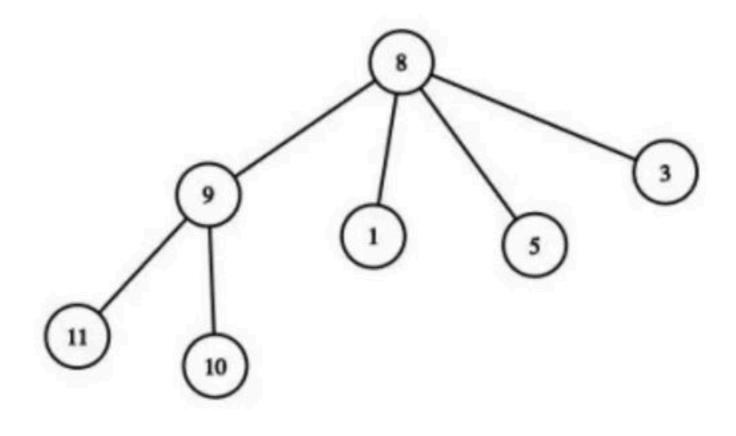
- A. 5
- B. 8
- C. 9
- D. 2



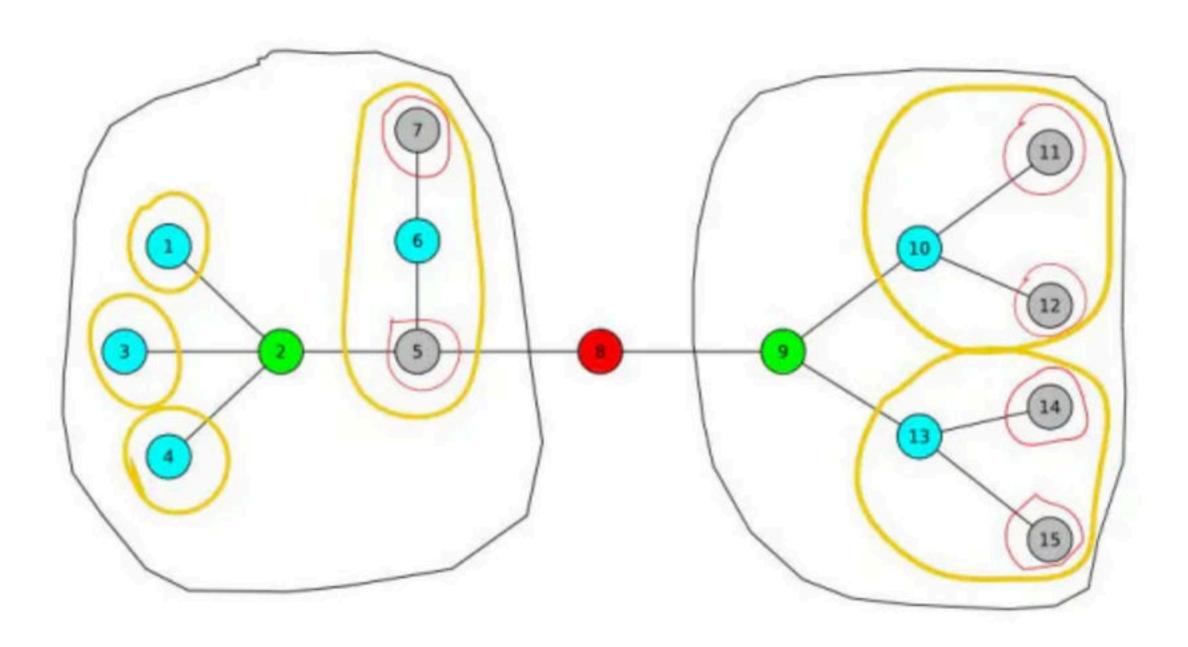
#### How to find a Centroid?

- Theorem (Jordan, 1869): For any given tree, the centroid always exists
- Proof: Start with any arbitrary vertex and check whether it satisfies the
  property. If yes, we're done, otherwise there exists only one adjacent subtree
  with more than N/2 nodes. Consider the adjacent vertex u in that subtree and
  apply the same argument. Repeat, till you find the centroid.





```
def decompose(root, centroid_parent = -1):
 centroid = find_centroid(root)
 if centroid_parent != -1:
   add_edge_in_centroid_tree(centroid_parent, centroid)
 for (adjecent_edge, adjacent_vertex) in G[centroid]:
   delete_edge(adjecent_edge)
   decompose(adjacent_vertex, centroid)
```



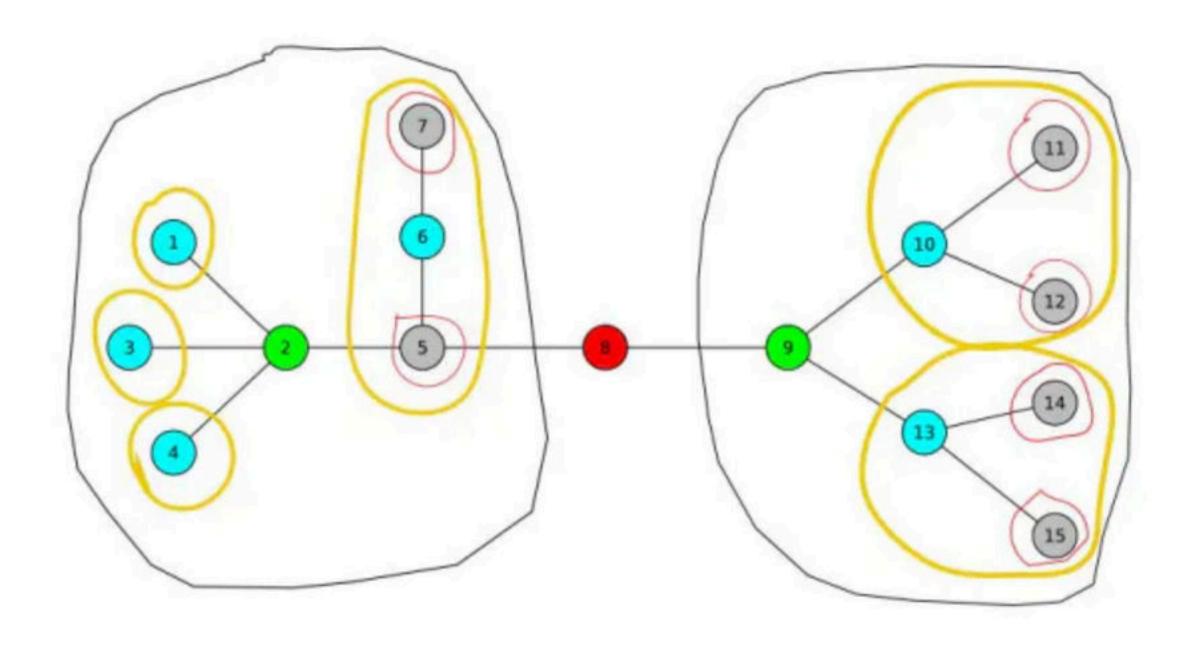
Question: Which node will be the root node of the Centroid Tree?

A. 2

B. 5

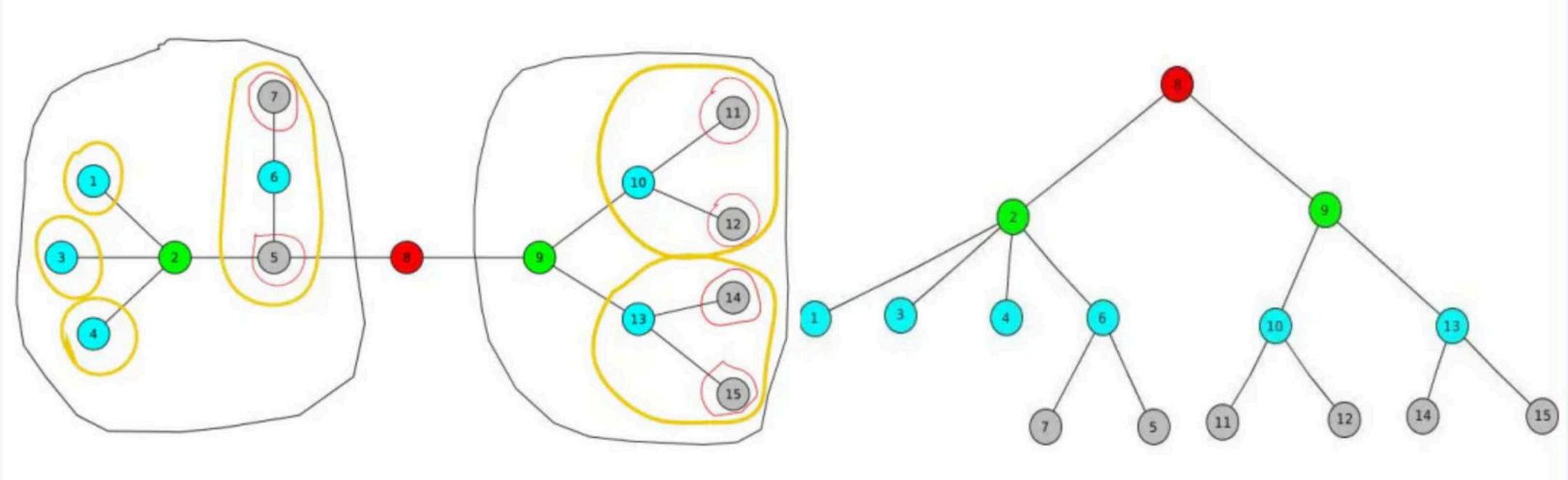
C. 8

D. 9



Question: How many children nodes would node 8 have in the Centroid Tree?

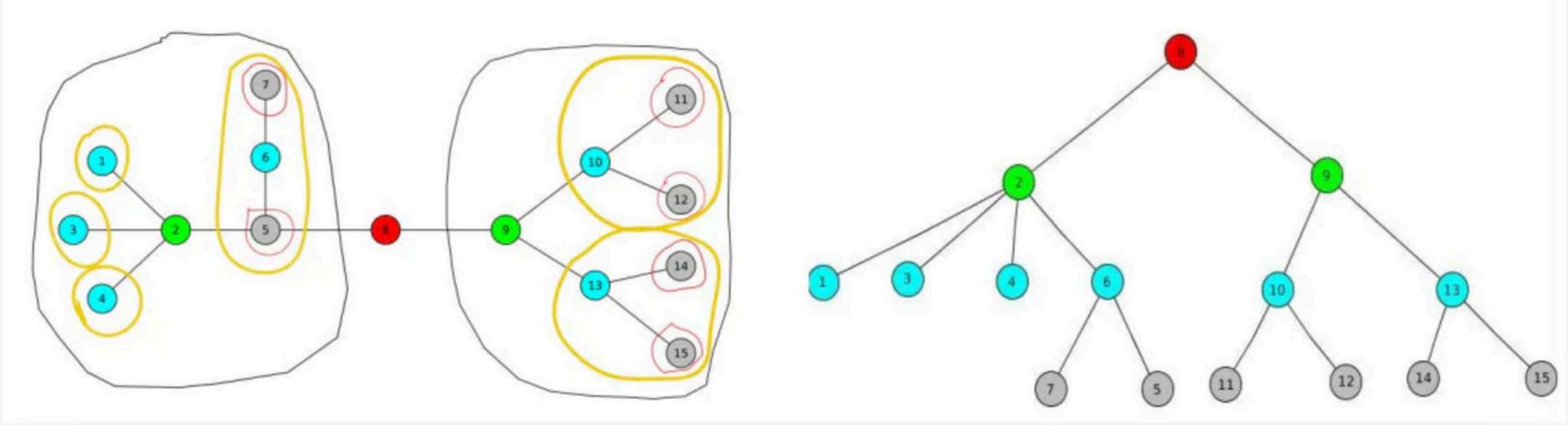
- A. 1
- B. 2
- C. 3
- D. 4



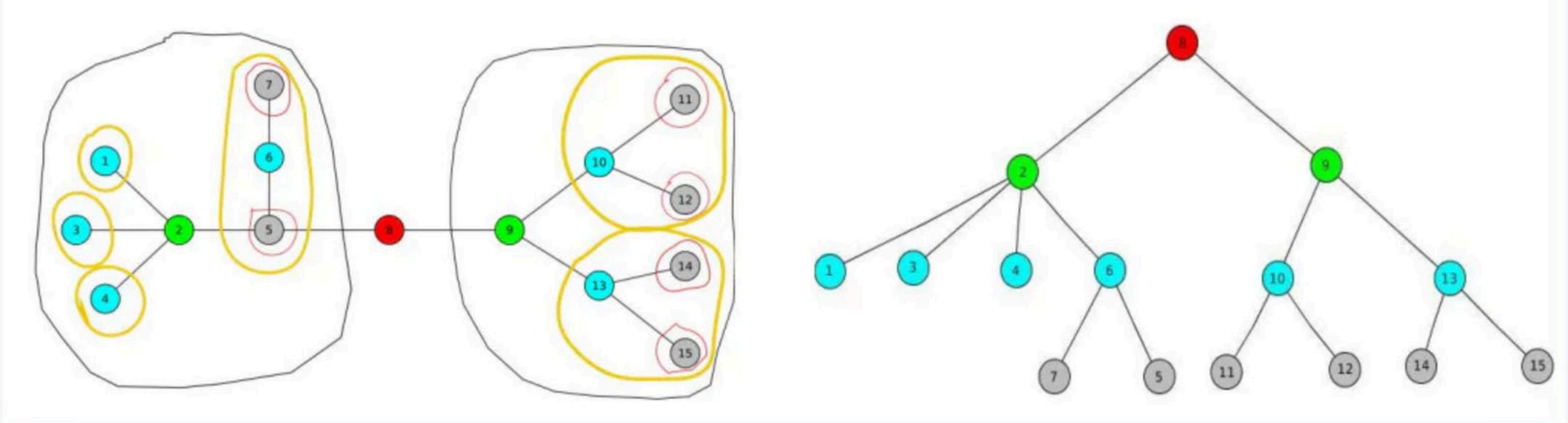
**Original Tree** 

**Centroid Tree** 

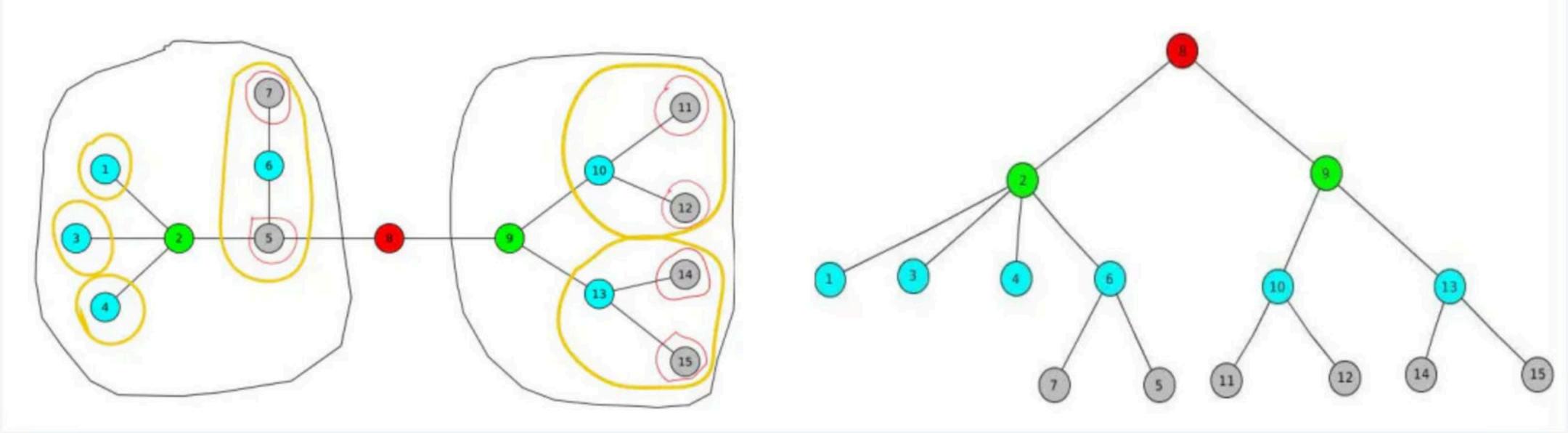
- Property-1: The Centroid Tree contains all N nodes of the original tree
- Since each node will become a centroid of some smaller tree (maybe a tree consisting only of that one single node), Hence the centroid tree formed would contain all the N nodes of the original tree.



- Property-2: The height of the centroid tree is at most O(logN)
- Since at each step, the new trees formed by removing the centroid have size at-most N/2, the maximum no of levels would be O(logN). Hence, the height of the centroid tree would be at most O(logN).



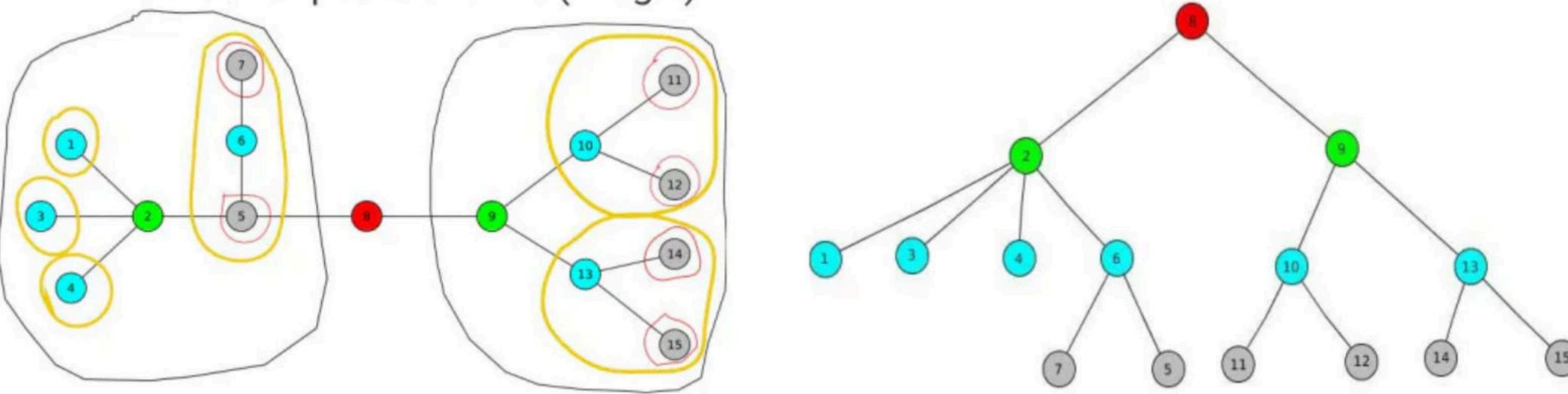
- Property-3: Consider any two arbitrary vertices A and B and the path between them (in the original tree) can be broken down into A-->C and C-->B where C is LCA of A and B in the centroid tree.
- In the original tree, the first time A and B got disconnected was when we removed vertex C. Hence, A-B = A-C + C-B.



 Property-4: Hence, we decompose the given tree into O(NlogN) different paths (from each centroid to all the vertices in the corresponding part) such that any path is a concatenation of two different paths from this set.

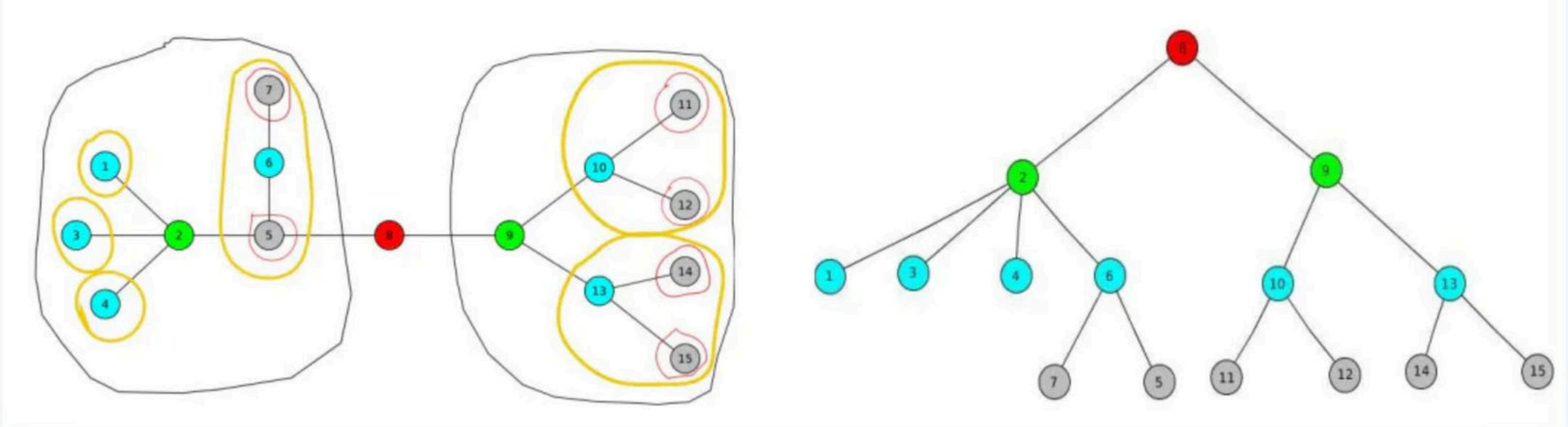
 There are at-most O(N) special paths at each level - every path starts at the centroid of that subtree and ends in a vertex. Since there are logN levels, total

number of paths are  $\sim O(NlogN)$ .



Q1: Given a tree with N nodes and Q queries of the form u, v - return the sum of elements on path from u to v.

- Do Centroid Decomposition and precompute dist[LOGN][N]
- Find LCA of u,v in the centroid tree: ans = dist[level(LCA)][u] + dist[level(LCA)][v]

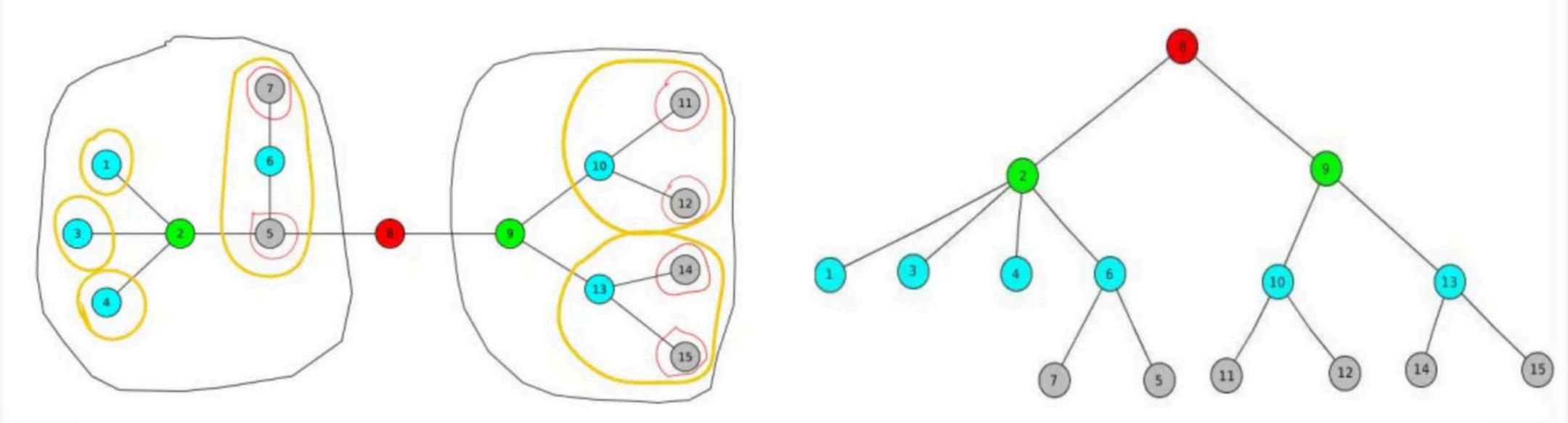


Q2: Given a tree with N blue nodes, there are Q queries of the form

Update v: Paint node v red.

Query v: Find the distance of closest red node to node v

- Let ans[i] denote the min distance to a red node for the centroid "i" in its corresponding part.
- For each update, to paint a node  $\mathbf{u}$  red, we move up to all the ancestors  $\mathbf{x}$  of  $\mathbf{u}$  in the centroid tree and update their ans as ans[x] = min(ans[x], dist(x,u)) because node u will be in the part of all the ancestors of u.
- For each query, to get the closest red node to node u, we again move up to all the ancestors of u in the centroid tree and take the minimum as mn = min(mn, dist(x,u) + ans[x]);



Q2: Given a weighted tree, find the no. of pairs of nodes the distance between which is a prime number.

#### **Naive Solution**

- Do a DFS and for every node, count the prime paths that pass through this node.
- O(SubtreeSize(x) \* NumberOfPrimes) for every node x.
- In worse case, this could be O(N<sup>2</sup> \* P).

Q2: Given a weighted tree, find the no. of pairs of nodes the distance between which is a prime number.

#### **Solution Using Centroid Decomposition**

- Do the same thing as earlier, but now by smartly rooting the tree only at centroids.
- For each centroid, we find the no of nodes at distance "i" from the centroid in it's part and store it in dist[i].
- Iterate over all primes and find number of "matching paths" from other subtrees by looking at dist[Prime[j] – distance(i,centroid)].
- At each level, we spend O(N \* P) therefore O(N \* P \* logN) total.

## Further Reading

 https://tanujkhattar.wordpress.com/2016/01/10/centroid-decomposition-ofa-tree/