



Solving competitive programming problems using Segment trees

Special class



Surya Kiran Adury



q1m13

- ICPC WF - 2014, 15 ~~#~~
- Google MTV 2017-20 ~~#~~
- Google LON 2015-17 ~~#~~
- B.Tech in ECE from IIT Roorkee





Objective

1. Class 1
 - a. Fenwick Trees
2. Class 2
 - a. Segment Trees
3. Class 3
 - a. Persistent Segment Trees
4. Class 4
 - a. Competitive programming problems ~~xx~~



What are we going to solve today?

1. Template problems

- a. Finding min of a range
- b. Finding number of elements in a range.
- c. Finding k'th of array.
- d. Sum of unique numbers of a range.

2. Actual problems

- a. [#330 Div 1 Problem - D - Codeforces](#)
- b. [SPOJ.com - Problem MKTHNUM](#)

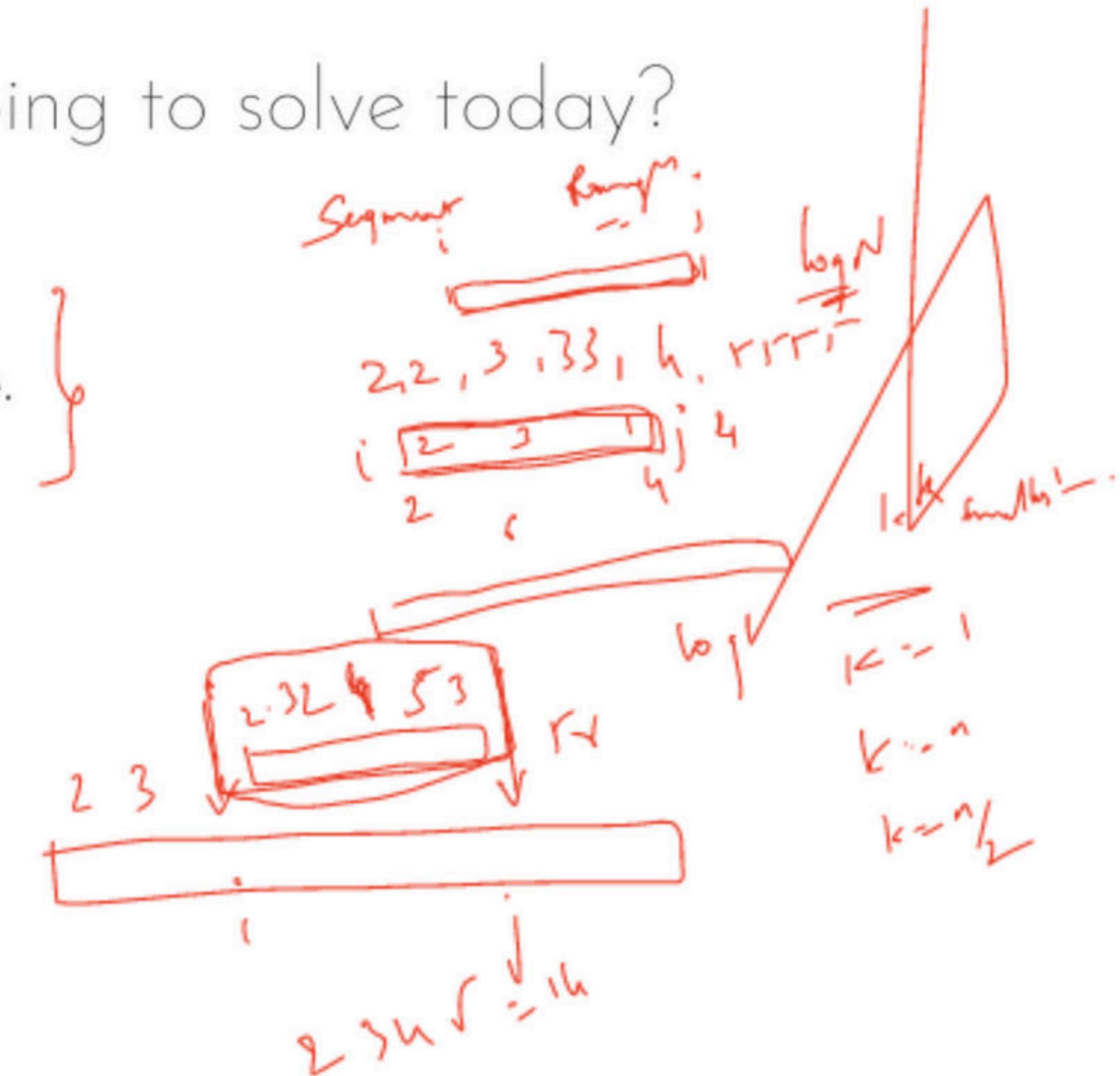
Medium

$$A = 9$$

$$B = 13$$

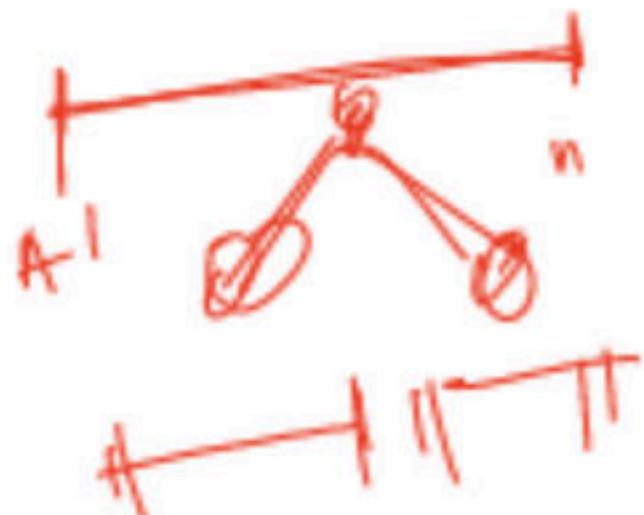
$$C = 11$$

$$D = 14$$





Finding min of a range.



```
struct SegMin {
    int left_child_id, right_child_id;
    int min_value;
} seg[N];
int used_id_cnt = 0;

int a[N];

inline void init(int seg_id, int seg_l, int seg_r) {
    if (seg_l == seg_r) {
        seg[seg_id].min_value = a[seg_l];
        return;
    }
    seg[seg_id].left_child_id = (++used_id_cnt);
    init(seg[seg_id].left_child_id, seg_l, (seg_l + seg_r) / 2);

    seg[seg_id].right_child_id = (++used_id_cnt);
    init(seg[seg_id].right_child_id, 1 + (seg_l + seg_r) / 2, r);

    seg[seg_id].min_value = min(seg[seg[seg_id].left_child_id].min_value,
                                seg[seg[seg_id].right_child_id].min_value);
}
```

init(0, l, n)

- A) O(N)
B) O(N log N)
C) O(N²)
D) N \approx 1 $n =$
 $N \cdot N \cdot N/h \dots$
 $= 2^r$



Finding min of a range.

$$a[id] = \sqrt{ch_w^2}$$

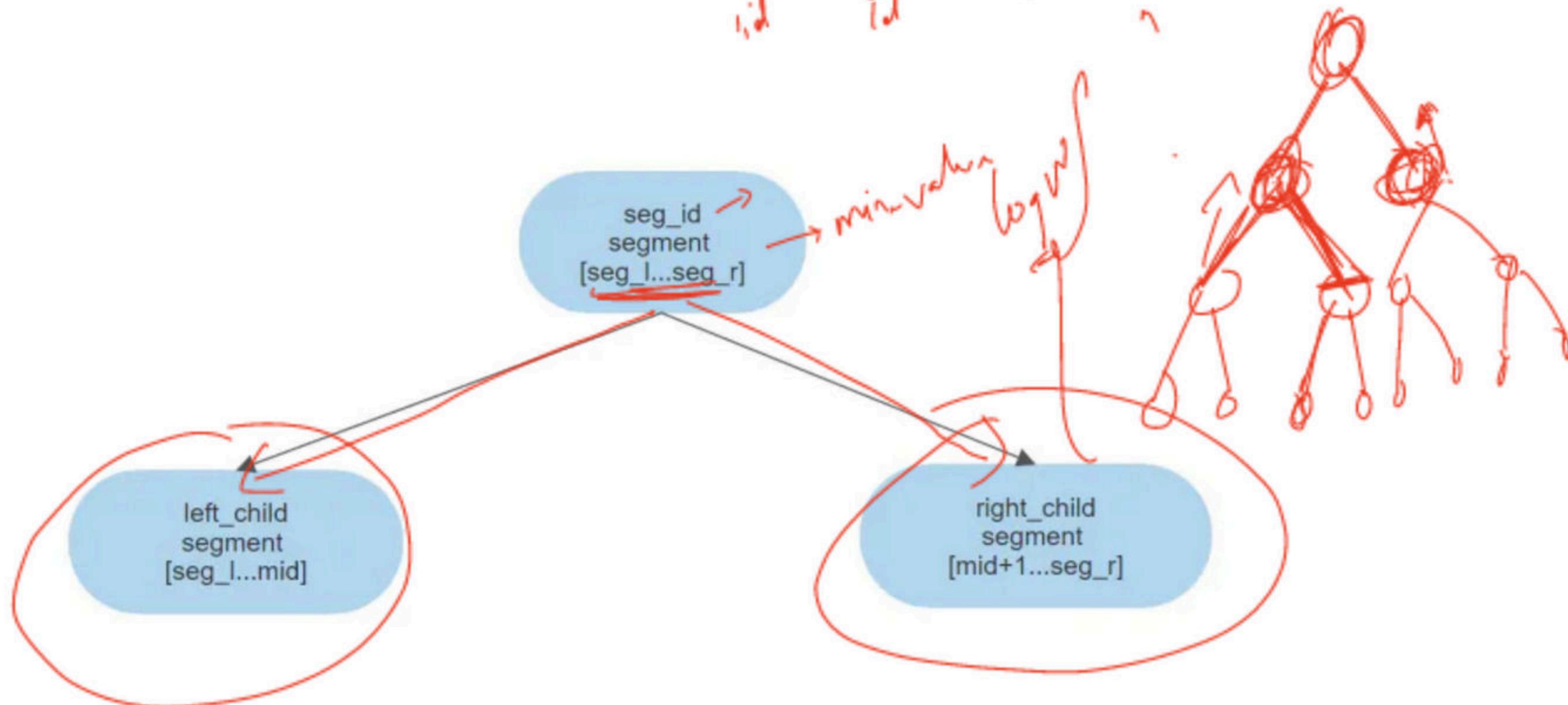
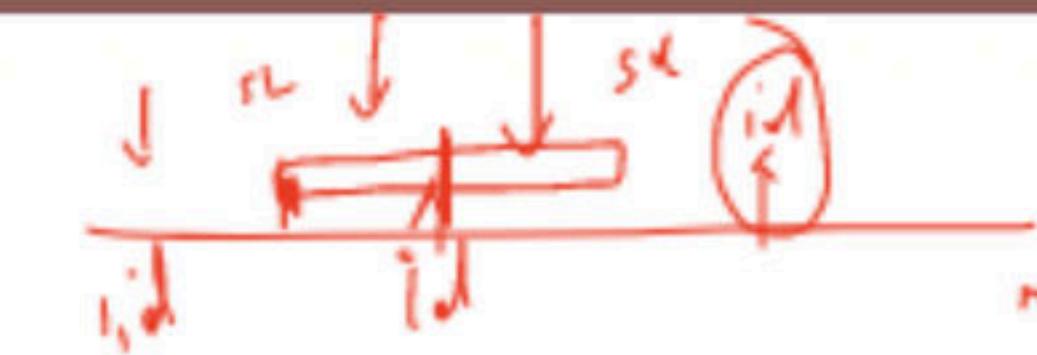
```
inline void update(int seg_id, int seg_l, int seg_r, int id, int value) {
    if (id < seg_l || id > seg_r) return;
    if (seg_l == seg_r) {
        seg[seg_id].min_value = value;
        a[id] = value;
        return;
    }
    update(seg[seg_id].left_child_id, seg_l, (seg_l + seg_r) / 2, id, value);
    update(seg[seg_id].right_child_id, 1 + (seg_l + seg_r) / 2, seg_r, id, value);

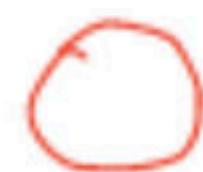
    seg[seg_id].min_value = min(seg[seg[seg_id].left_child_id].min_value,
                                seg[seg[seg_id].right_child_id].min_value);
}
```

$O(n)$ $O(n)$ $O(n)$

↓ ↓ ↓

id id value

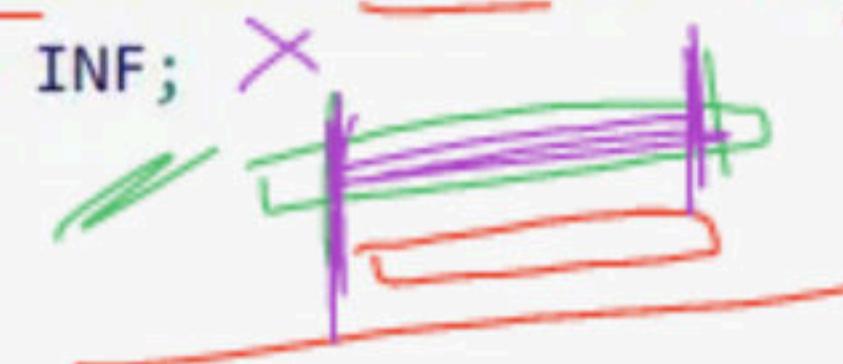


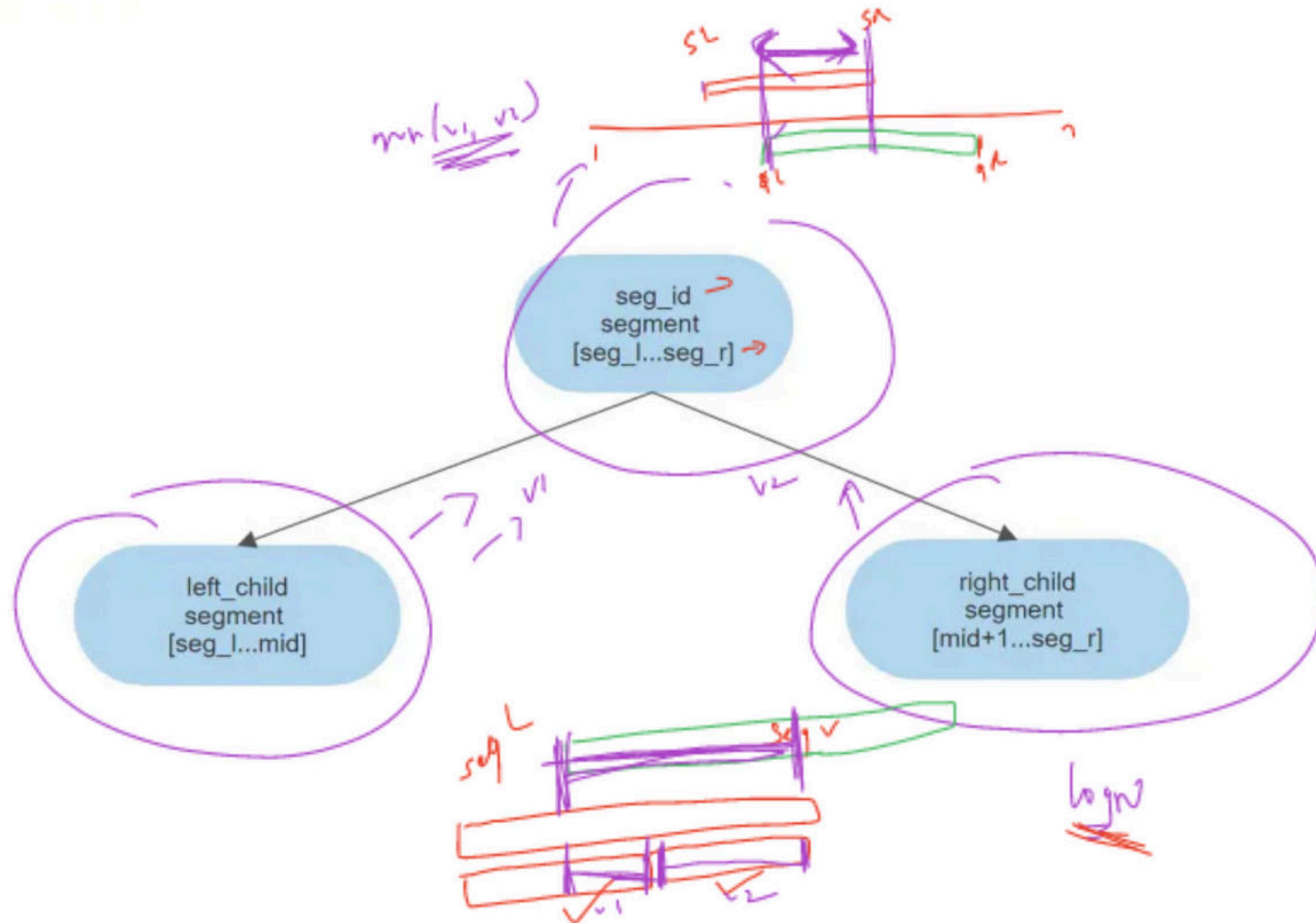


Finding min of a range.



```
inline int query_min(int seg_id, int seg_l, int seg_r, int query_l, int query_r) {
    if (query_l > seg_r || query_r < seg_l) return INF;
    if (query_l <= seg_l && seg_r <= query_r) {
        return seg[seg_id].min_value;
    }
    return min(query_min(seg[seg_id].left_child_id, seg_l, (seg_l + seg_r) / 2, query_l, query_r),
               query_min(seg[seg_id].right_child_id, 1 + (seg_l + seg_r) / 2, seg_r, query_l, query_r));
}
```





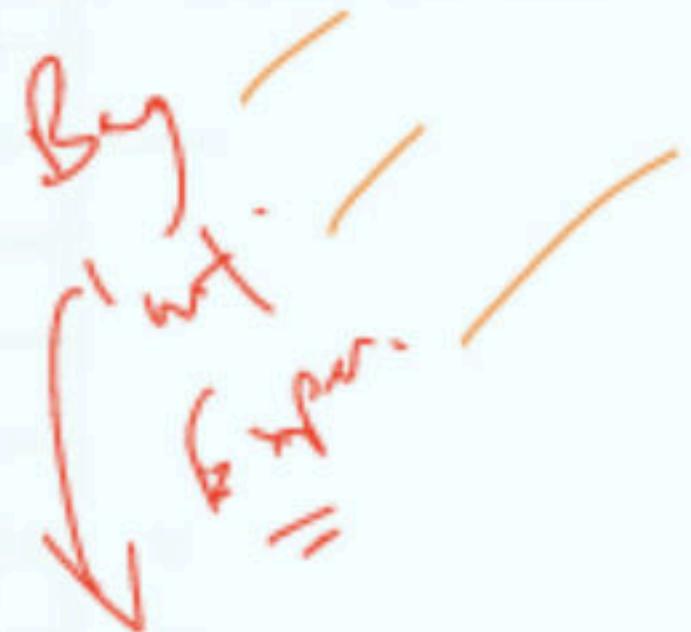




Questions?

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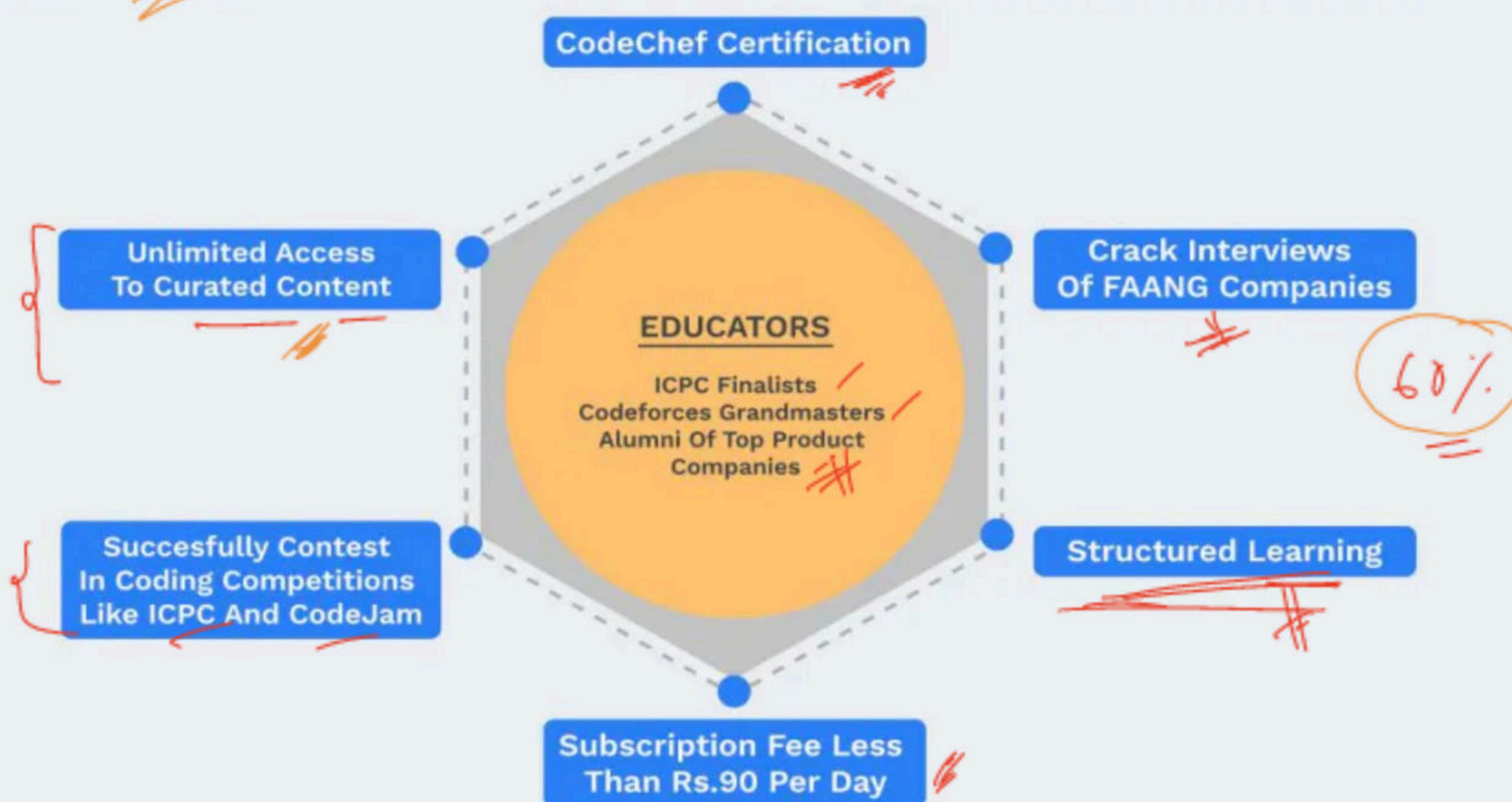
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- Conquest 2021: Year Long Journey for Intermediate Coders to Become Experts (C++) - Live on 8th Jan 2021
- C++: Conquest 2021: From Programming Fundamentals to Career Readiness - Live on 8th Jan 2021
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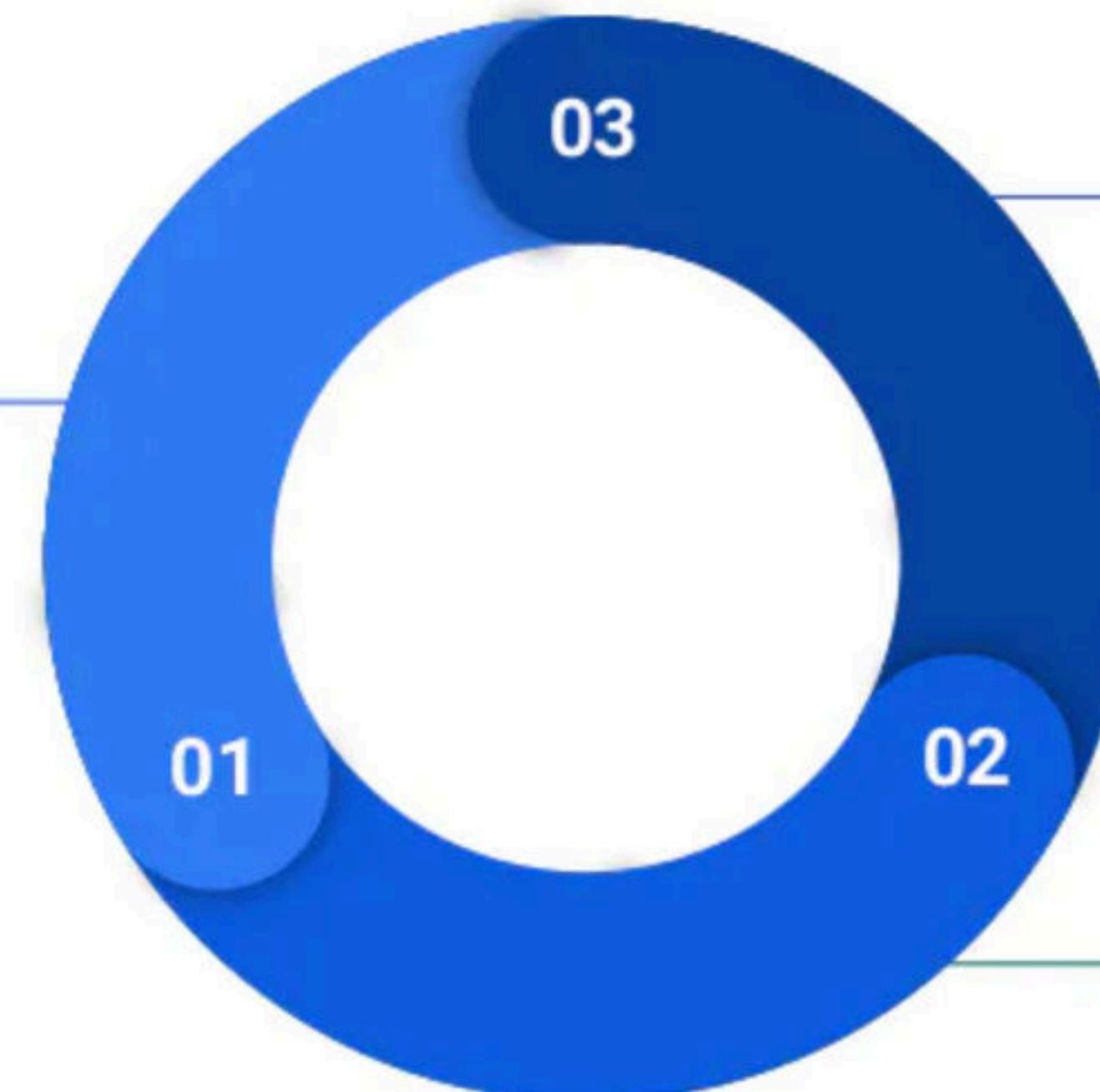




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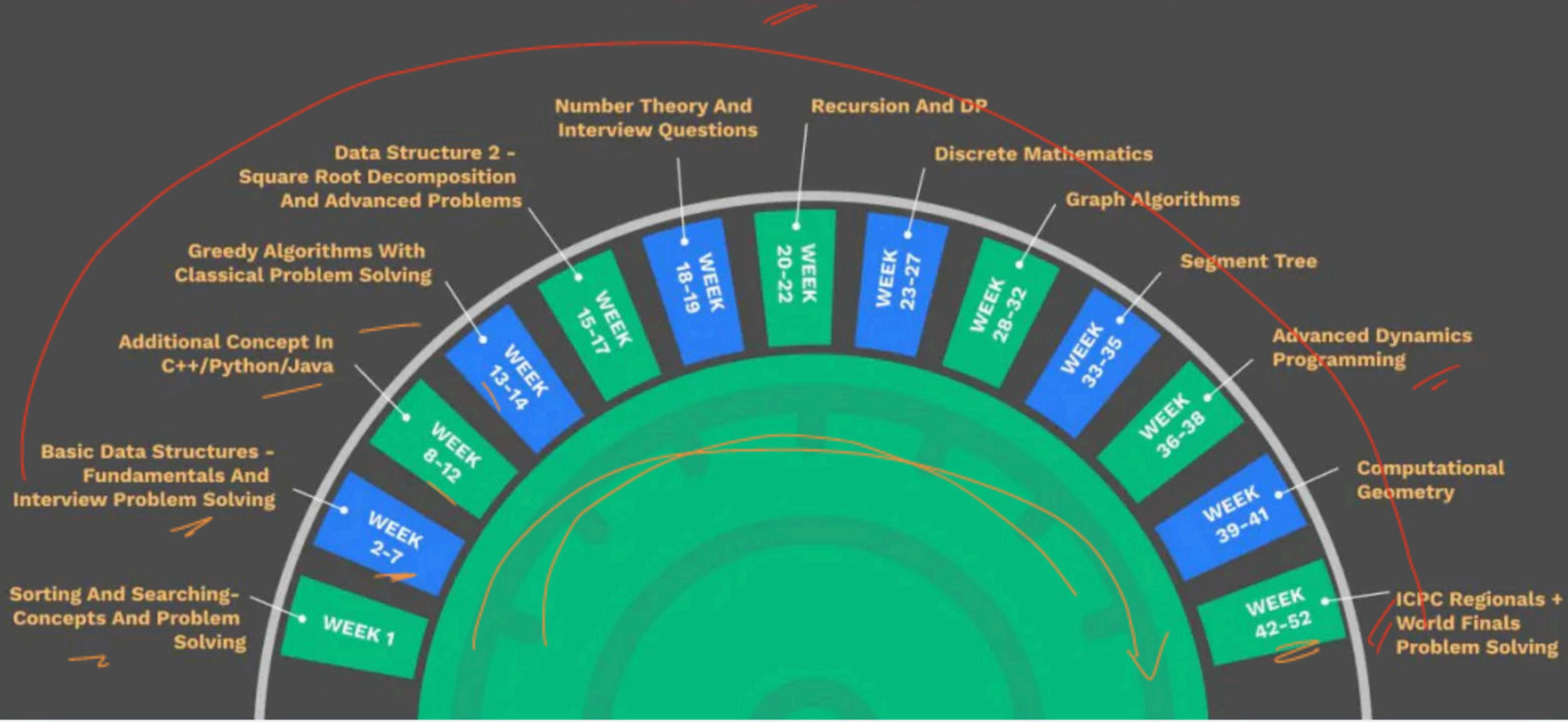
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Deepak Gour and 1 more		Sanket Singh and 1 more		Pulkit Chhabra	



HINDI	ENGLISH	HINDI	ENGLISH	HINDI	ENGLISH
Everest-Python : Complete Course on Competitive Programming		Everest-C++ : Complete Course on Competitive Programming		Everest-Java : Complete Course on Competitive Programming	
Started on Dec 14		Started on Dec 14		Started on Dec 14	
Sanket Singh		Deepak Gour and 1 more		Sanket Singh and 1 more	



Educators

**Tanuj Khattar**

ACM ICPC World Finalist - 2017, 2018. Indian IOI Team Trainer 2016-2018. Worked @ Google, Facebook, HFT. Quantum Computing Enthusiast.

**Sanket Singh**

Software Development Engineer @ LinkedIn | Former SDE @ Interviewbit | Google Summer of Code 2019 @ Harvard University | Former Intern @ISRO

**Pulkit Chhabra**

Codeforces: 2246 | Codechef: 2416 | Former SDE Intern @CodeNation | Former Intern @HackerRank

**Riya Bansal**

Software Engineer at Flipkart | Former SDE and Instructor @ InterviewBit | Google Women TechMakers Scholar 2018

**Triveni Mahatha**

Qualified ICPC 2016 World Final. Won multiple Codechef Long Challenges (India). ICPC Onsite Regionals' Problem setter and Judge. IIT Kanpur.

**Deepak Gour**

ICPC World Finalist 2020 | Former Instructor @InterviewBit | Software Engineer at AppDynamics



Educators

**Himanshu Singh**

World Finalist ICPC 2020, Winner Techgig Code Gladiators 2020, Winner TCC '19, 2020 CSE Graduate from IIT BHU, Works at Nutanix

**Murugappan S**

Software engineer at Google. Have won many programming contests. Max Rating of 2192 in codeforces and 2201 in codechef.

**Nishchay Manwani**

Hey I am Nishchay Manwani from CSE, IIT Guwahati and I'm a Seven star on Codechef and International Grandmaster on Codeforces.

**Vivek Chauhan**

Codechef: 7 stars (2612) India Rank 6, Codeforces: MASTER (2279), Won Codechef Long Challenges(India), TCO20 Southern Asia Runner up

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Course-wise Practice Problems



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Learn Competitive Programming with CodeChef

Trees and Graphs

Pulkit Chhabra

Starts on 21 Sep



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Name

Code

* Successful Submissions * Accuracy

Problems will be available in 6 days 7 hrs 23 mins 22 sec

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Contest Starts in:

6 7 23 22

Days Hrs Min Sec

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Awesome! You got 10% off!

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Proceed to pay

A) Beginner ✓

B) Novice ✓

C) Intermediate ✓

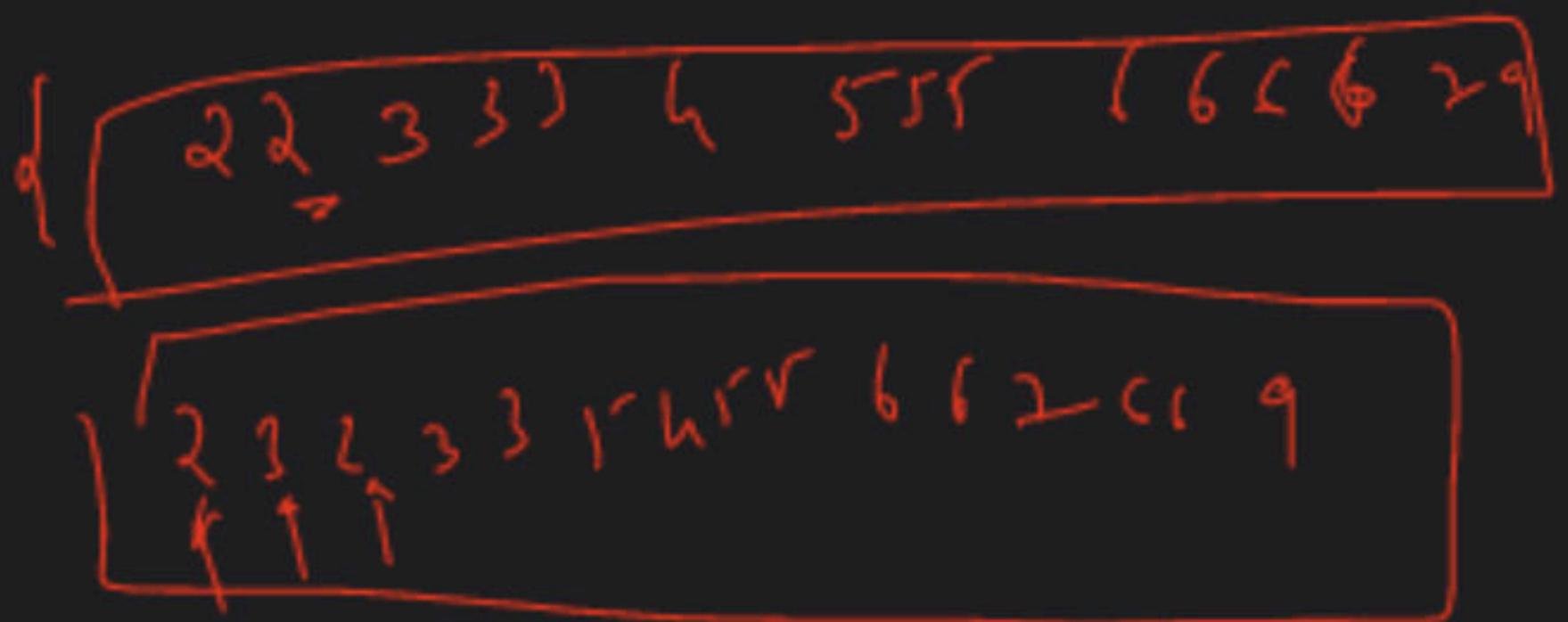
D) Advanced ✓



~~Finding number of elements in a range~~

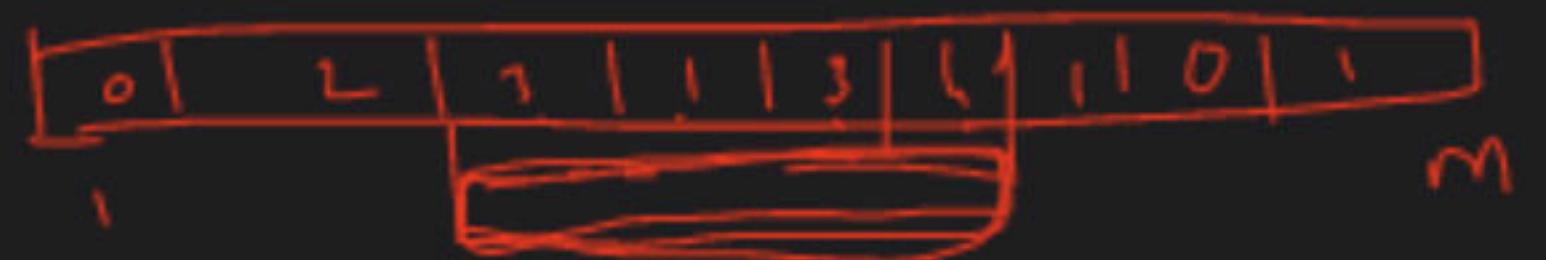
```
struct SegCnt {  
    int left_child_id, right_child_id;  
    int cnt;  
} seg_cnt[N];  
int used_id_cnt = 0;  
  
int a[N];  
  
inline void update_cnt(int seg_id, int seg_l, int seg_r, int id, int v) {  
    if (id < seg_l || id > seg_r) return;  
    seg_cnt[seg_id].cnt += v;  
    if (seg_l != seg_r) {  
        update_cnt(seg_cnt[seg_id].left_child_id, seg_l, (seg_l + seg_r) / 2, id, v);  
        update_cnt(seg_cnt[seg_id].right_child_id, 1 + (seg_l + seg_r) / 2, seg_r, id, v);  
    }  
}
```

~~i j = 31~~

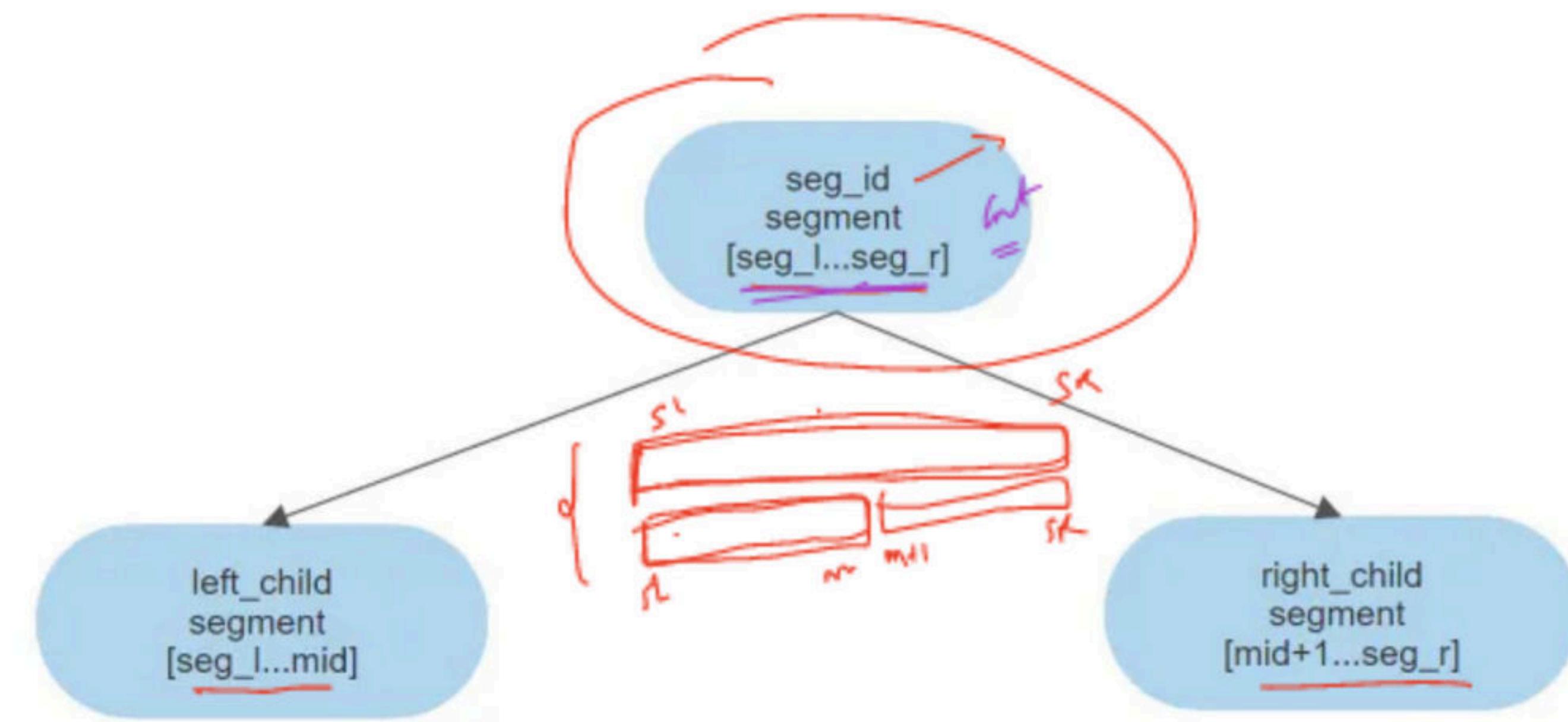


$i \leq x \leq j$

... 2 3 4 5 6 7 8 9



Char!

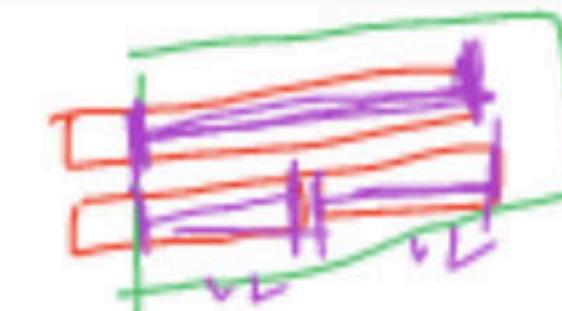




Finding number of elements in a range

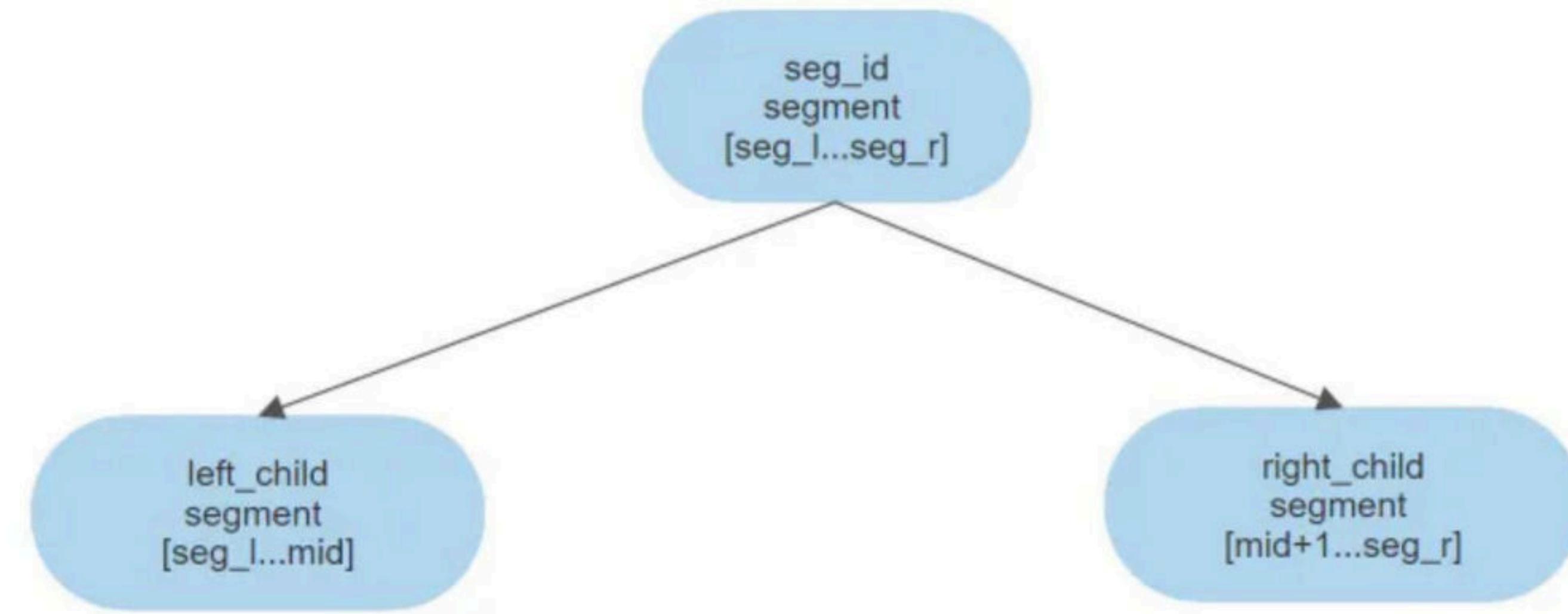
query cnt (i, l, m, i, j)

```
// Returns count of elements in the range intersection of (seg_l, seg_r) and (query_l, query_r)
inline int query_cnt(int seg_id, int seg_l, int seg_r, int query_l, int query_r) {
    if (query_l > seg_r || seg_l > query_r) return 0;
    if (query_l <= seg_l && seg_r <= query_r) return seg_cnt[seg_id].cnt;
    return query_cnt(seg_cnt[seg_id].left_child_id, seg_l, (seg_l + seg_r) / 2, query_l, query_r)
           + query_cnt(seg_cnt[seg_id].right_child_id, 1 + (seg_l + seg_r) / 2, seg_r, query_l, query_r);
}
```





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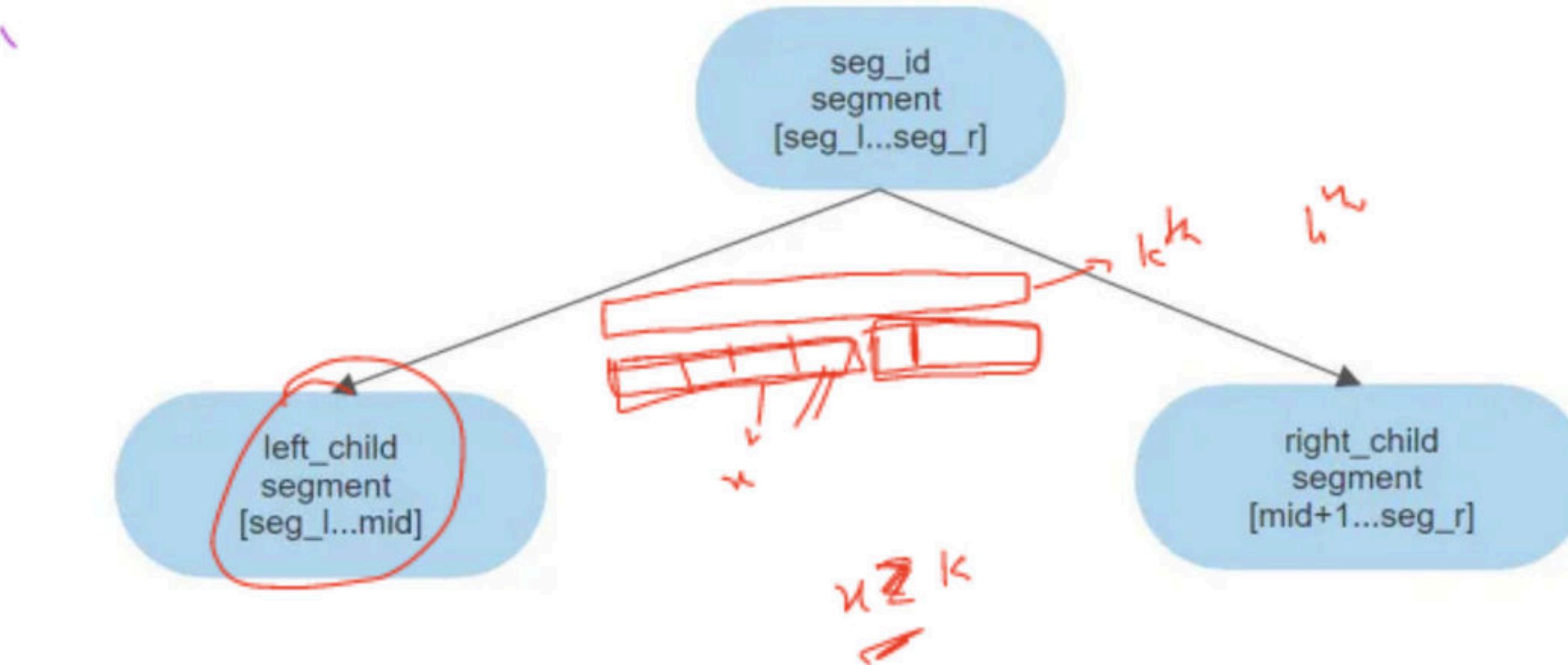




Finding k'th element in a range

```
// Returns k'th element in the range (seg_l, seg_r)
inline int query_kth(int seg_id, int seg_l, int seg_r, int k) {
    if (seg_r == seg_l) {
        return seg_l;
    }
    int left_cnt = seg_cnt[seg_cnt[seg_id].left_child_id].cnt;
    if (left_cnt >= k) {
        return query_kth(seg_cnt[seg_id].left_child_id, seg_l, (seg_l + seg_r) / 2, k);
    } else {
        return query_kth(seg_cnt[seg_id].right_child_id, 1 + (seg_l + seg_r) / 2, seg_r, k - left_cnt);
    }
}
```

$|c \leq k$ $k > c$





<-->

<-->

Questions?
~~Not~~



Sum of unique numbers of a range

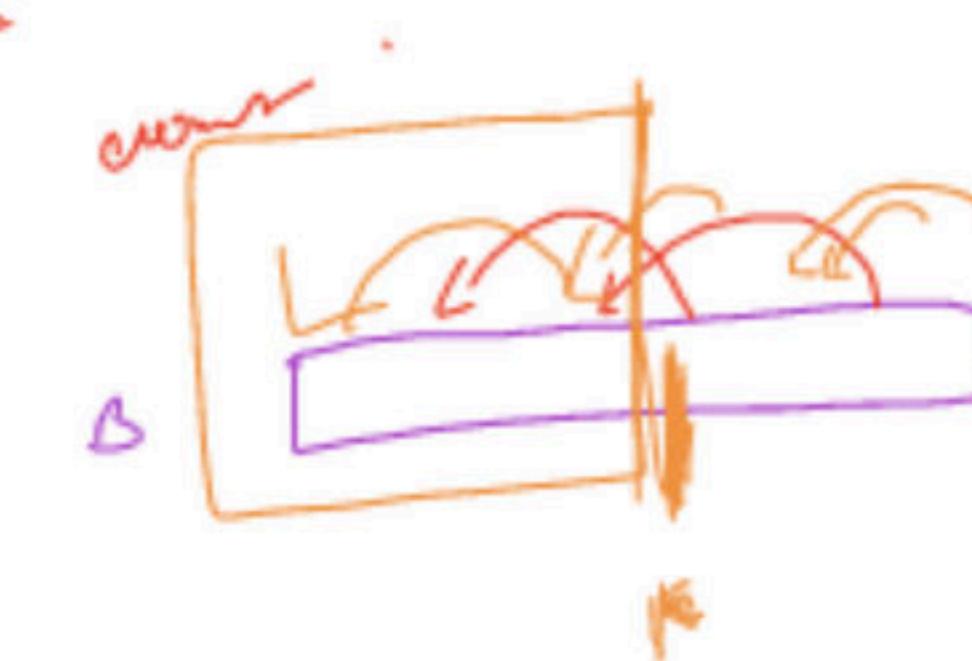
$A = [1 \ 2 \ 3 \ \dots \ N]$ $N_{\text{elements}} =$



$$25 + 31 + 37 = 93$$

$$v_2 - v_1 =$$

Persistent Segment Tree





Sum of unique numbers of a range

```
int last_occ[N];
int a[N];

inline void solve () {
    int n;
    cin >> n;
    rep (i, 1, n) {
        cin >> a[i];
    }
    root_id[0] = init(0, n);
    for (int i = 1; i <= n; ++i) {
        root_id[i] = update(root_id[i - 1], 0, n, last_occ[a[i]], a[i]);
        last_occ[a[i]] = i;
    }
    int q;
    cin >> q;
    while (q--) {
        int l, r;
        cin >> l >> r;
        int prefix_sum_1 = query(root_id[r], 0, n, l - 1);
        int prefix_sum_2 = query(root_id[l - 1], 0, n, l - 1);
        cout << prefix_sum_1 - prefix_sum_2 << endl;
    }
}
```

B 3
12
C =

Bat ~~her~~
= ~~lions~~
vomit

posting \downarrow vise



Sum of unique numbers of a range

```
Struct SegSum {
    int left_child_id;
    int right_child_id;
    int sum;
} seg_sum[S];

int root_id[N];
int left_child_id[N];
int right_child_id[N];
int used_seg_cnt = 0;

inline int init(int l, int r) {
    int seg_id = ++used_seg_cnt;
    if (l < r) {
        seg_sum[seg_id].left_child_id = init(l, (l + r) / 2);
        seg_sum[seg_id].right_child_id = init(1 + (l + r) / 2, r);
    }
    return seg_id;
}
```



Sum of unique numbers of a range

```
inline int make_parent(int left_child_id, int right_child_id) {
    int new_seg_id = ++used_seg_cnt;
    seg_sum[new_seg_id].sum = seg_sum[left_child_id].sum + seg_sum[right_child_id].sum;
    seg_sum[new_seg_id].left_child_id = left_child_id;
    seg_sum[new_seg_id].right_child_id = right_child_id;
    return new_seg_id;
}

inline int make_leaf(int v) {
    int new_seg_id = ++used_seg_cnt;
    seg_sum[new_seg_id].sum = v;
    return new_seg_id;
}

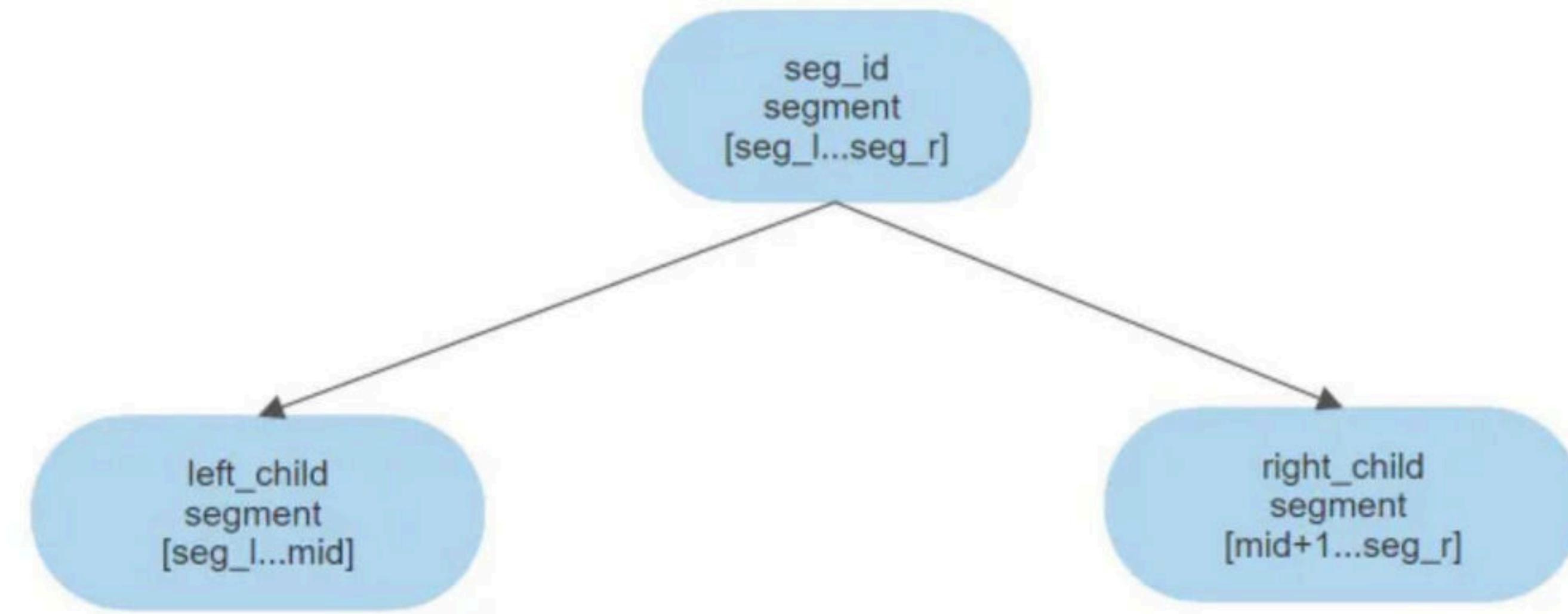
inline int update(int seg_id, int seg_l, int seg_r, int id, int v) {
    if (id < seg_l || id > seg_r) return seg_id;

    if (seg_l == seg_r) return make_leaf(id, v);

    return make_parent(update(seg_sum[seg_id].left_child_id, seg_l, (seg_l + seg_r) / 2, id, v),
                      update(seg_sum[seg_id].right_child_id, 1 + (seg_l + seg_r) / 2, seg_r, id, v));
}
```



<->



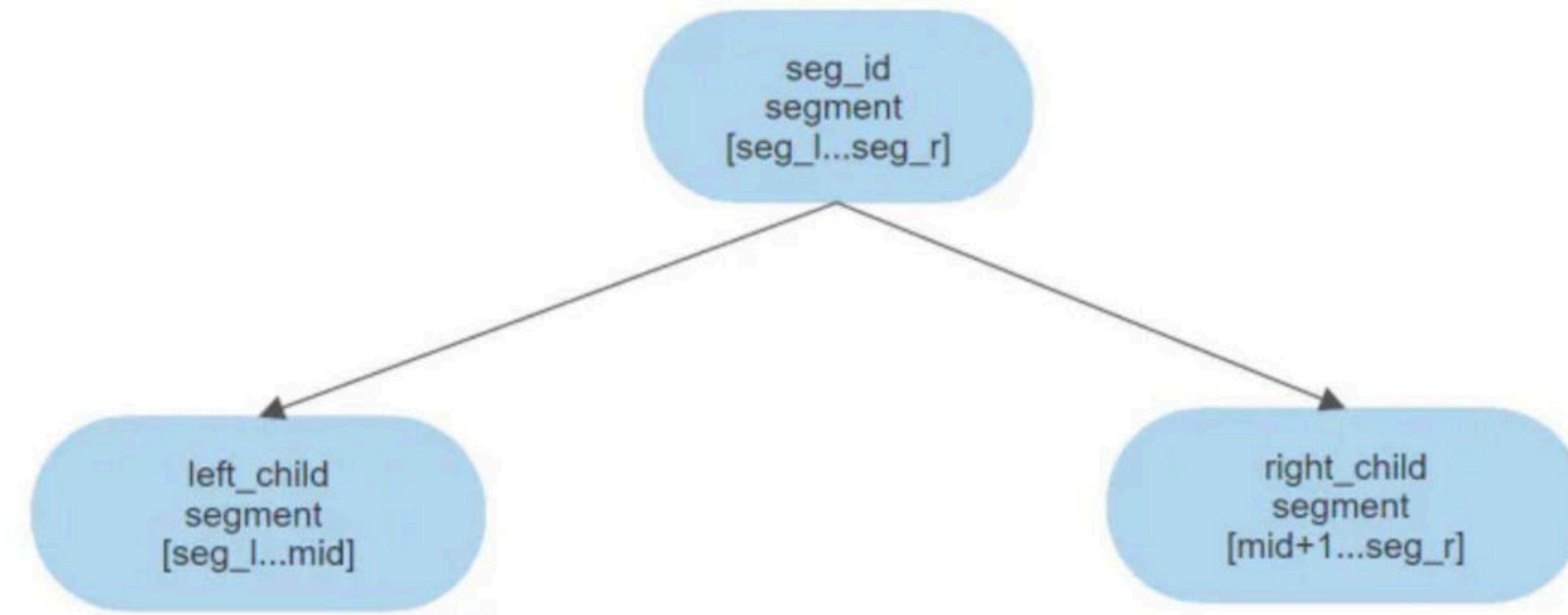


Sum of unique numbers of a range

```
int query(int seg_id, int seg_l, int seg_r, int query_prefix_id) {
    if (query_prefix_id < seg_l) return 0;
    if (query_prefix_id >= seg_r) return seg_sum[seg_id].sum;
    return query(seg[seg_id].left_child_id, seg_l, (seg_l + seg_r) / 2, query_prefix_id)
        + query(seg[seg_id].right_child_id, 1 + (seg_l + seg_r) / 2, seg_r, query_prefix_id);
}
```



<->





Questions?



#330 Div 1 Problem - D - Codeforces

D. REQ



time limit per test: 3 seconds

memory limit per test: 256 megabytes

input: standard input

output: standard output

Today on a math lesson the teacher told Vovochka that the Euler function of a positive integer $\varphi(n)$ is an arithmetic function that counts the positive integers less than or equal to n that are relatively prime to n . The number 1 is coprime to all the positive integers and $\varphi(1) = 1$.

Now the teacher gave Vovochka an array of n positive integers a_1, a_2, \dots, a_n and a task to process q queries $l_i r_i$ — to calculate and print

$$\varphi\left(\prod_{i=l}^r a_i\right)$$

modulo $10^9 + 7$. As it is too hard for a second grade school student, you've decided to help Vovochka.

$$n = p_1^{e_1} p_2^{e_2} \cdots p_k^{e_k}$$

$$\varphi(n) = n \cdot \frac{\left(1 - \frac{1}{p_1}\right) \left(1 - \frac{1}{p_2}\right) \cdots \left(1 - \frac{1}{p_k}\right)}{\left(1 - \frac{1}{p_1}\right)}$$

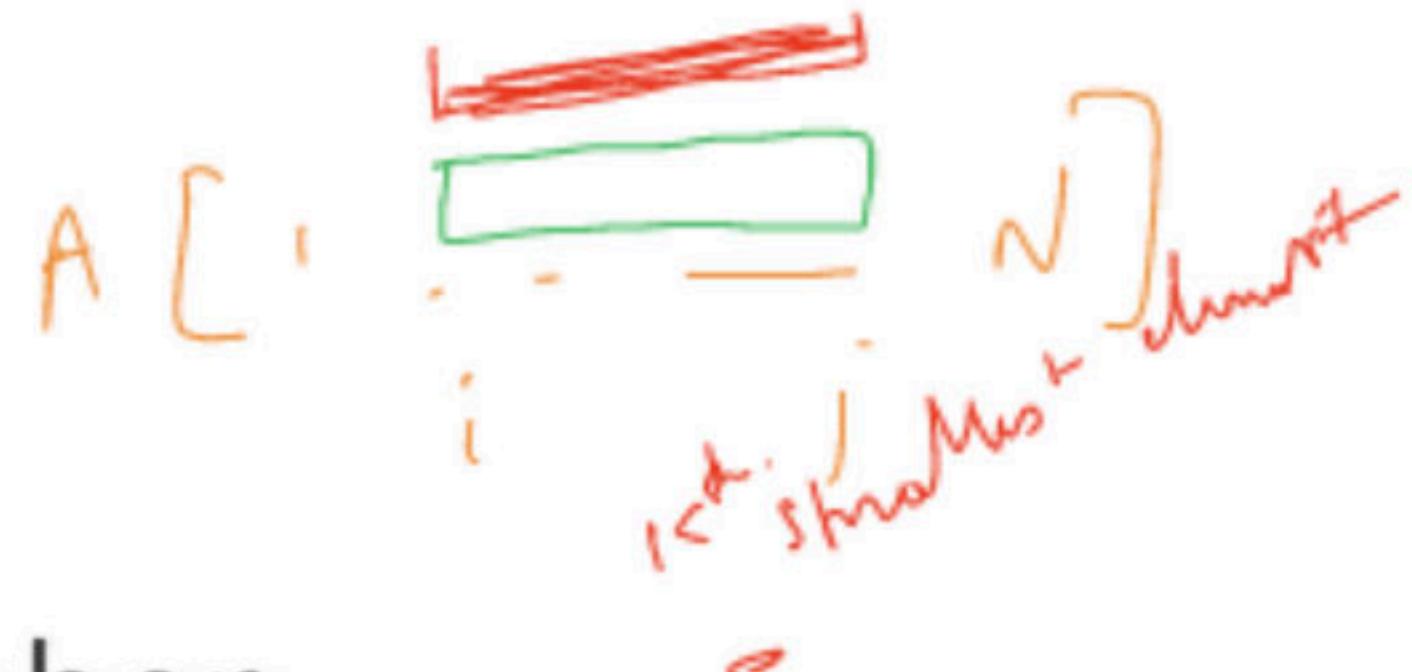




Questions?



MKTHNUM



MKTHNUM - K-th Number

#sorting #tree

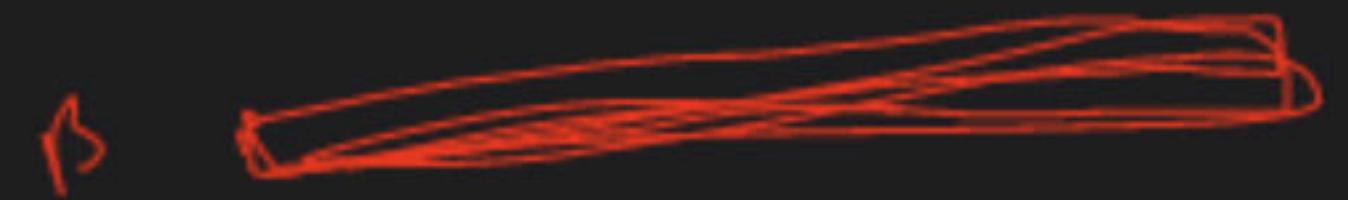
English

Vietnamese

You are working for Macrohard company in data structures department. After failing your previous task about key insertion you were asked to write a new data structure that would be able to return quickly k-th order statistics in the array segment.

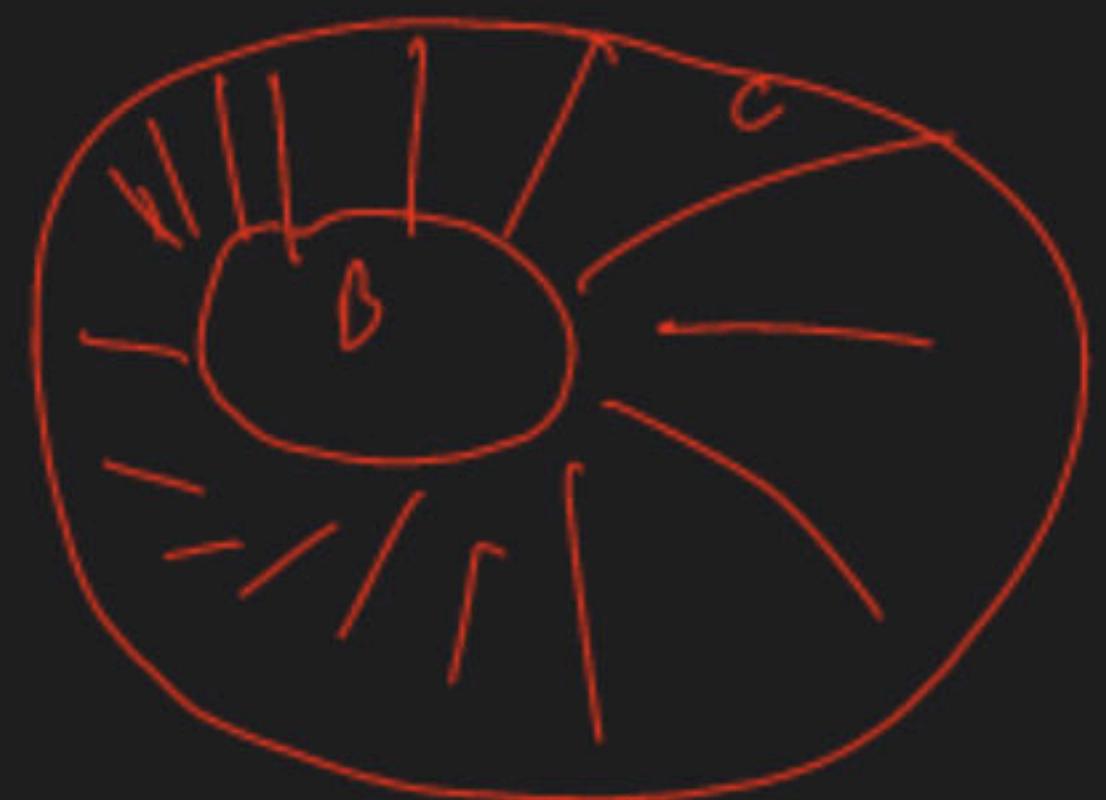
That is, given an array $a[1 \dots n]$ of different integer numbers, your program must answer a series of questions $Q(i, j, k)$ in the form: "What would be the k-th number in $a[i \dots j]$ segment, if this segment was sorted?"

For example, consider the array $a = (1, 5, 2, 6, 3, 7, 4)$. Let the question be $Q(2, 5, 3)$. The segment $a[2 \dots 5]$ is $(5, 2, 6, 3)$. If we sort this segment, we get $(2, 3, 5, 6)$, the third number is 5, and therefore the answer to the question is 5.



path
Kⁿ element

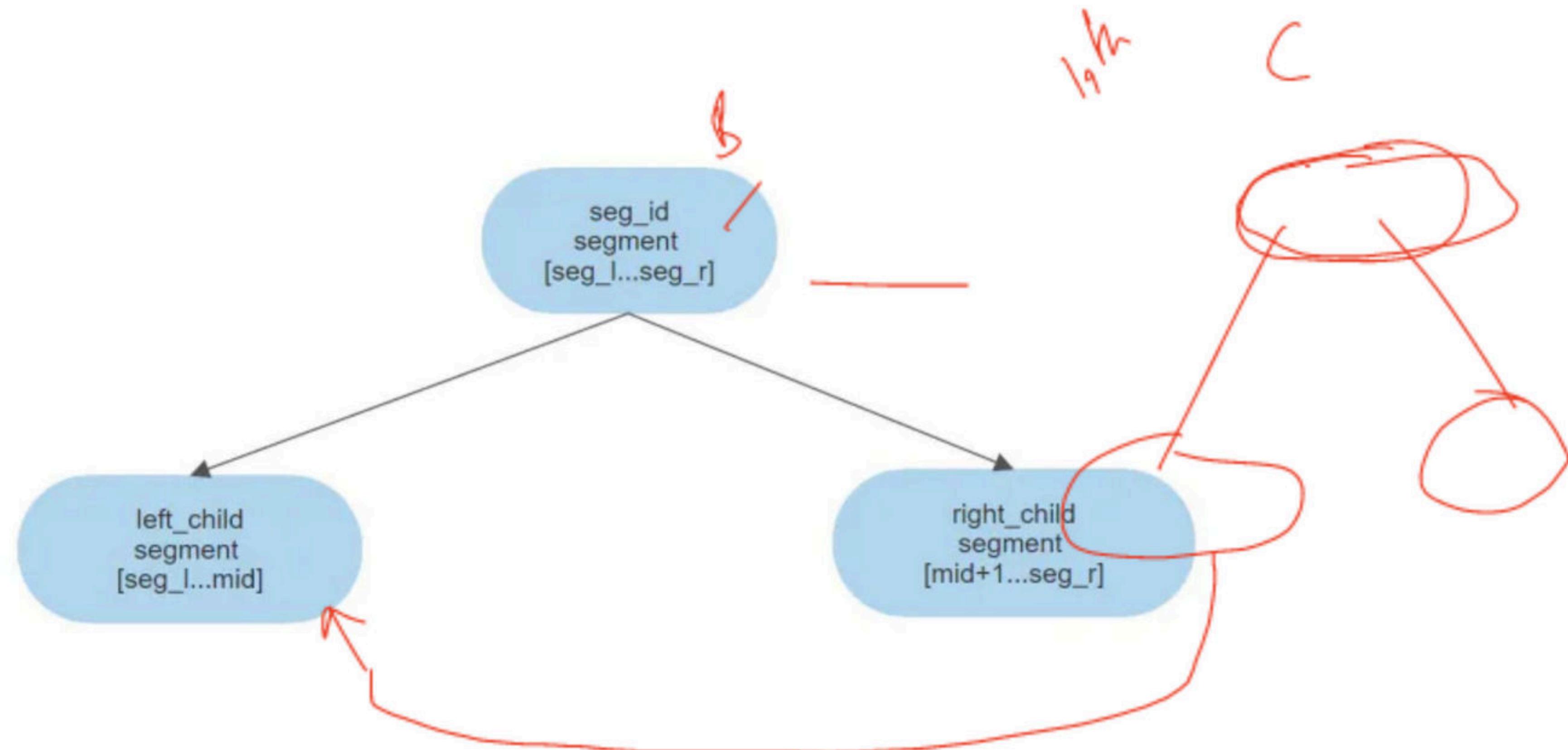
B < C





MKTHNUM

```
inline int query(int seg_id1, int seg_id2, int seg_l, int seg_r, int k) {
    if (seg_l == seg_r) return seg_l;
    int left_cnt = seg_cnt[seg_cnt[seg_id1].left_child_id].cnt
            - seg_cnt[seg_cnt[seg_id2].left_child_id].cnt;
    if (left_cnt >= k) {
        return query_kth(seg_cnt[seg_id1].left_child_id, seg_cnt[seg_id2].left_child_id, seg_l, (seg_l + seg_r) / 2, k);
    } else {
        return query_kth(seg_cnt[seg_id1].right_child_id, seg_cnt[seg_id2].right_child_id, 1 + (seg_l + seg_r) / 2, seg_r, k - left_cnt);
    }
}
```





MKTHNUM

```
inline int make_parent(int left_child_id, int right_child_id) {
    int new_seg_id = ++used_seg_cnt;
    seg_sum[new_seg_id].sum = seg_sum[left_child_id].sum + seg_sum[right_child_id].sum;
    seg_sum[new_seg_id].left_child_id = left_child_id;
    seg_sum[new_seg_id].right_child_id = right_child_id;
    return new_seg_id;
}

inline int make_leaf(int v) {
    int new_seg_id = ++used_seg_cnt;
    seg_sum[new_seg_id].sum = v;
    return new_seg_id;
}

inline void update(int seg_id, int seg_l, seg_r, int id) {
    if (id < seg_l || id > seg_r) return seg_id;

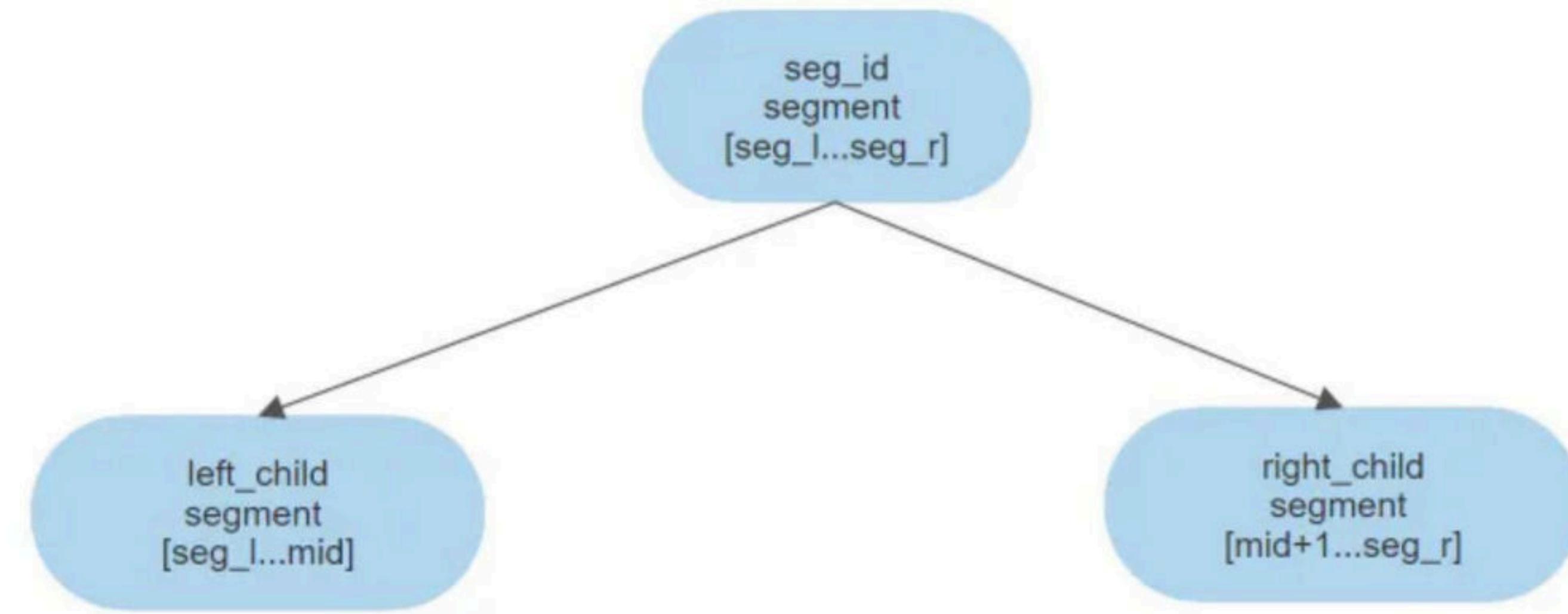
    if (seg_l == seg_r) return make_leaf(id, seg_cnt[seg_id].cnt + 1);

    return make_parent(update(seg_sum[seg_id].left_child_id, seg_l, (seg_l + seg_r) / 2, id),
                       update(seg_sum[seg_id].right_child_id, 1 + (seg_l + seg_r) / 2, seg_r, id));
}
```

AS^T



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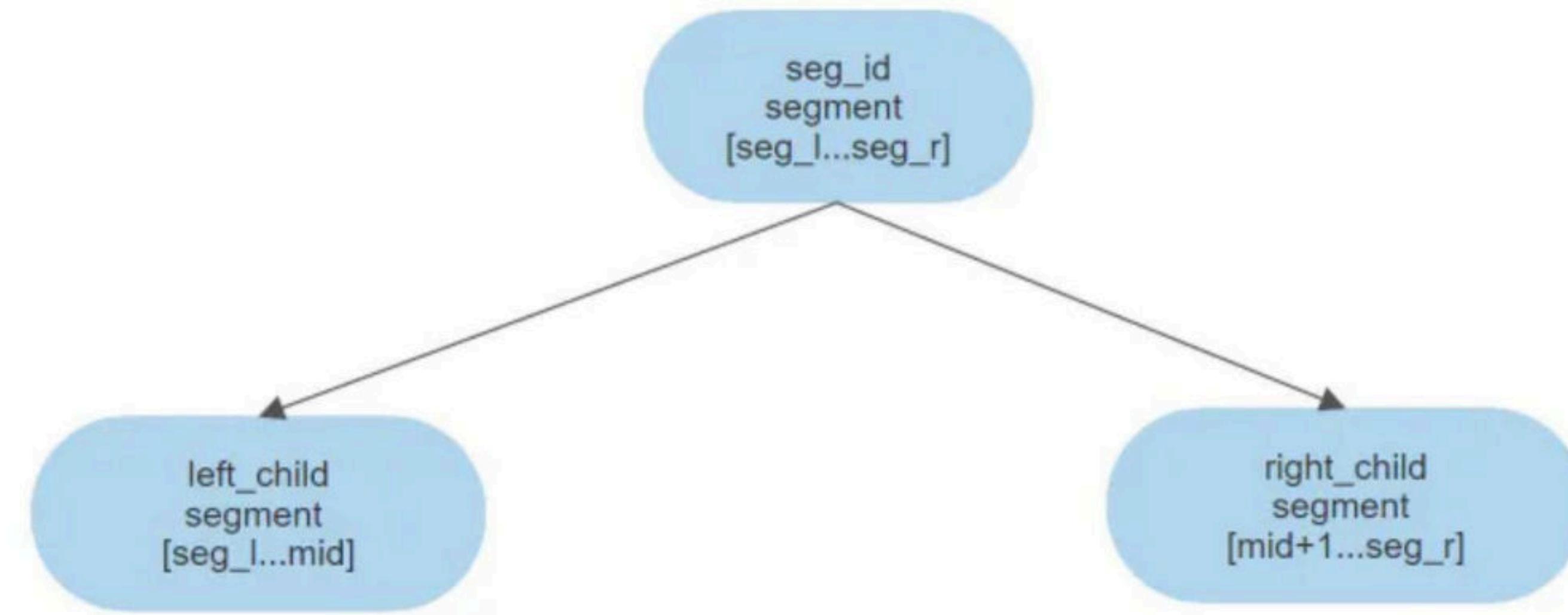


MKTHNUM

```
inline void solve() {
    int n;
    cin >> n;
    root_id[0] = init(1, M);
    for (int i = 1; i <= n; ++i) {
        cin >> a[i];
        cnt[a[i]]++;
        root_id[i] = update(root_id[i - 1], 1, M, a[i]);
    }
    int q;
    cin >> q;
    while (q--) {
        int i, j, k;
        cin >> i >> j >> k;
        cout << query(root_id[j], root_id[i - 1], 1, M, k) << endl;
    }
}
```



<->





Questions?

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