CS331 Tutorial 4

Java: CDS

A. Sahu

Dept of Comp. Sc. & Engg.

Indian Institute of Technology Guwahati

<u>Outline</u>

- Four Way to implement Lock in Java
- Concurrent List
- Concurrent Hash
- Assignment-II
 - Use may use C-List and C-Hash

Four ways to implement a synchronized counter in Java

- Suppose there is a Shared counter and every threads are attempting to manipulate
 - 1. Synchronized Block
 - 2. Atomic Variable
 - 3. Concurrent Lock
 - 4. Semaphore

First: Fine-Grained Synchronization

- Instead of using a single lock ...
- Split object into
 - Independently-synchronized components
 - Example: Hash: Modification
- Methods conflict when they access
 - The same component ...
 - At the same time

Second: Optimistic Synchronization

- Search without locking ...
- If you find it, lock and check ...
 - OK: we are done
 - Oops: start over
- Evaluation
 - Usually cheaper than locking, but
 - Mistakes are expensive

Third: Lazy Synchronization

- Postpone hard work
- Removing components is tricky
 - Logical removal
 - Mark component to be deleted
 - Physical removal
 - Do what needs to be done

Fourth: Lock-Free Synchronization

- Don't use locks at all
 - Use compareAndSet() & relatives ...

Fourth: Lock-Free Synchronization

- Don't use locks at all
 - Use compareAndSet() & relatives ...
- Advantages
 - No Scheduler Assumptions/Support
- Disadvantages
 - Complex
 - Sometimes high overhead

Concurrent List

Sequential List Based Set

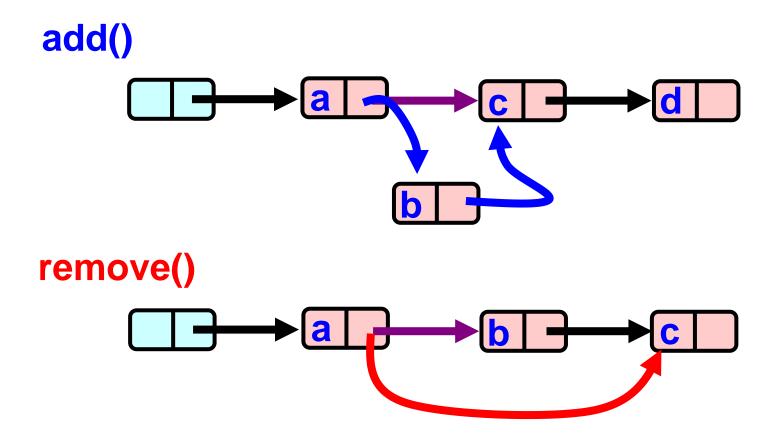
Add()



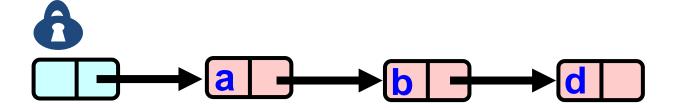
Remove()



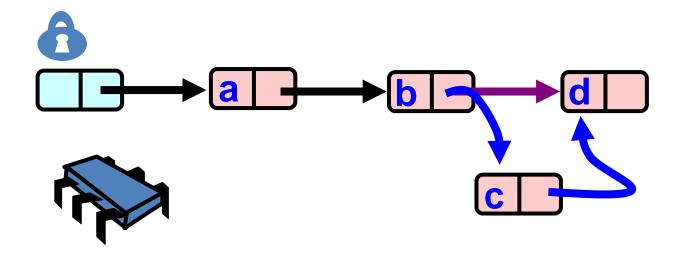
Sequential List Based Set



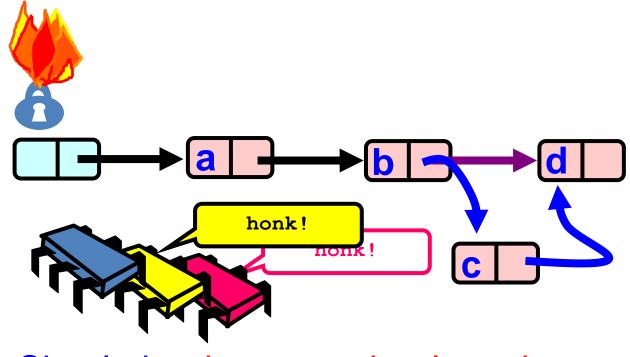
Coarse-Grained Locking



Coarse-Grained Locking



Coarse-Grained Locking



Simple but hotspot + bottleneck

List

- Suppose a list have Nodes, each node have
 - Key value : Customer Name
 - Item Value : Amount
- Node : key, value, next
- Operations frequencies
 - Balance check : 50%
 - Update amount: 49.8
 - Remove account : 0.1%
 - Add account : 0.1%

Fine Gran locking

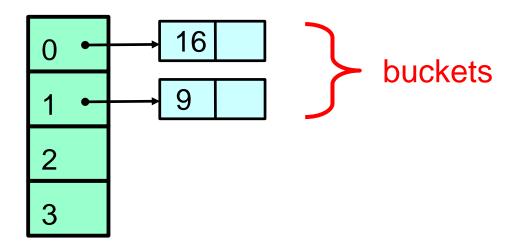
- Suppose a list have Nodes, each node have
 - Key value : Customer Name
 - Item Value : Amount
- Node : key, value, next
- Operations frequencies
 - Balance check: 50%: Lock not required
 - Update amount: 49.8: lock of the required node
 - Remove account : 0.1% : lock entire list
 - Add account : 0.1% : lock entire list

Fine Gran locking

- Suppose a list have Nodes, each node have
 - Key value : Customer Name
 - Item Value : Amount
- Node : key, value, next
- Operations frequencies
 - Balance check: 50%: Lock not required
 - Update amount: 49.8: lock of the required node
 - Remove account : 0.1% : lock entire list
 - Add account : 0.1% : lock entire list
- If the Add and remove account contribute: 90% of operation, we need to fine grain add and remove: Refer Book Art of Multiprocessor programming

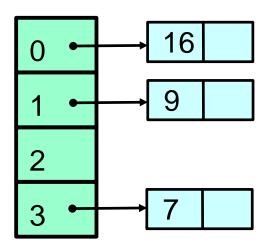
Concurrent Hash

Sequential Closed Hash Map



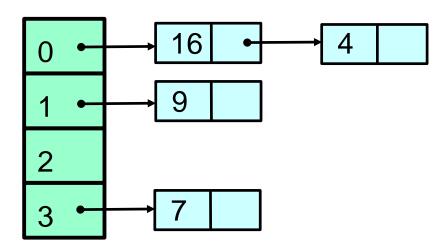
2 Items

Add an Item



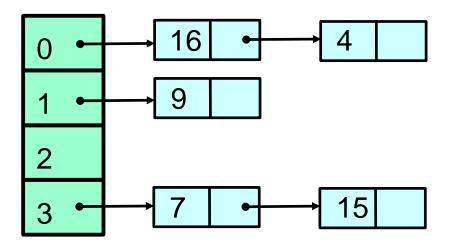
3 Items

Add Another: Collision



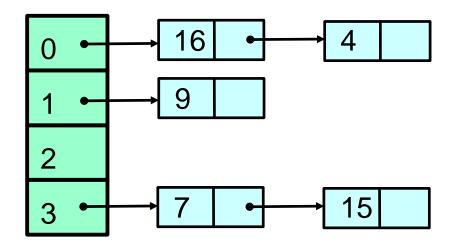
4 Items

More Collisions



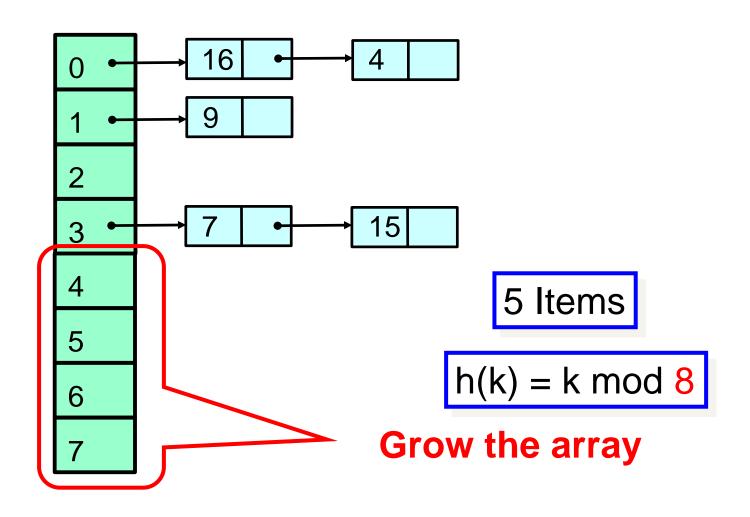
5 Items

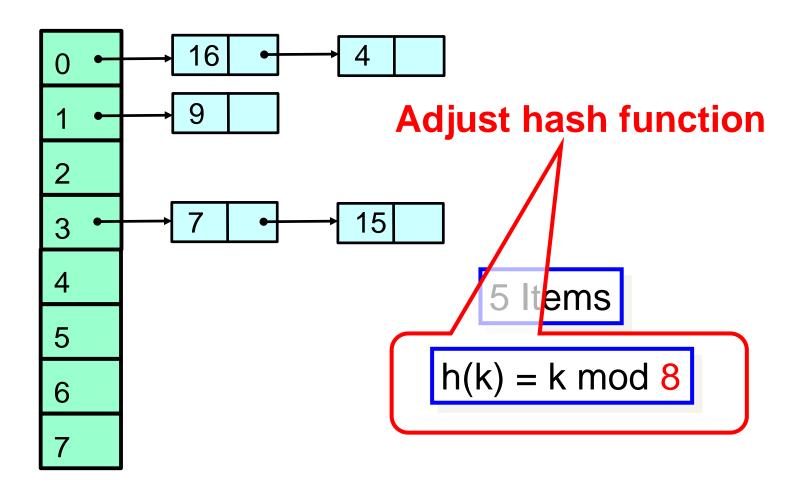
More Collisions

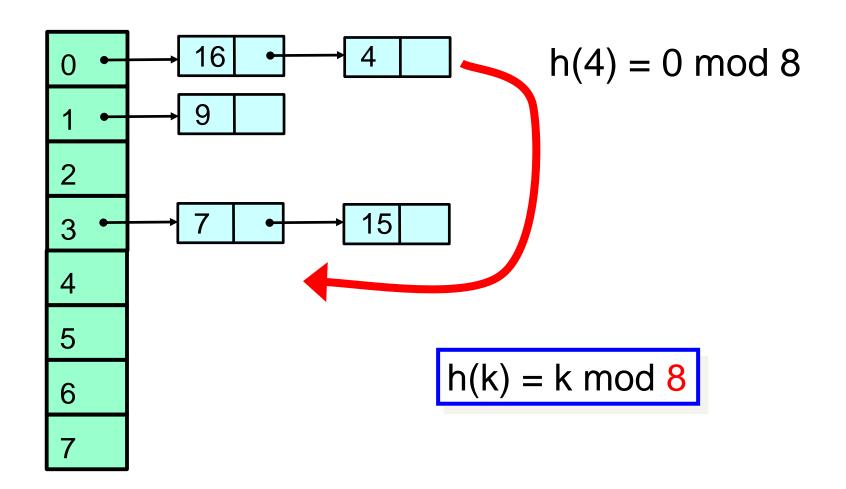


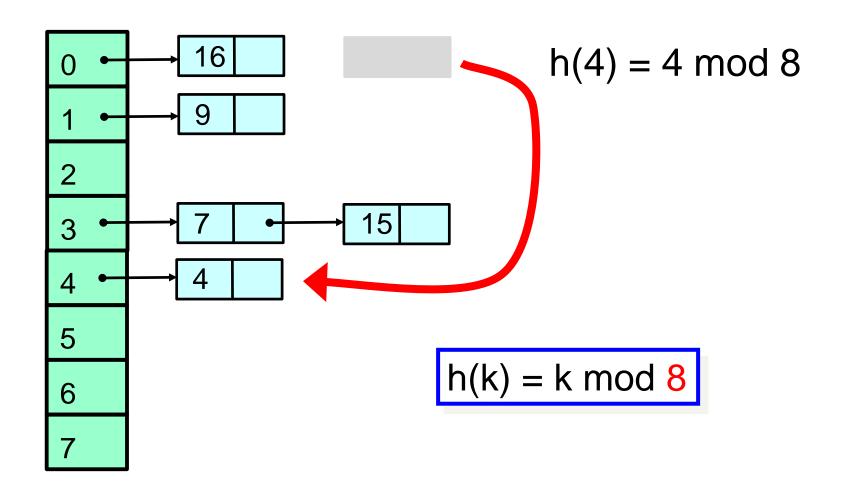
Problem: buckets getting too long

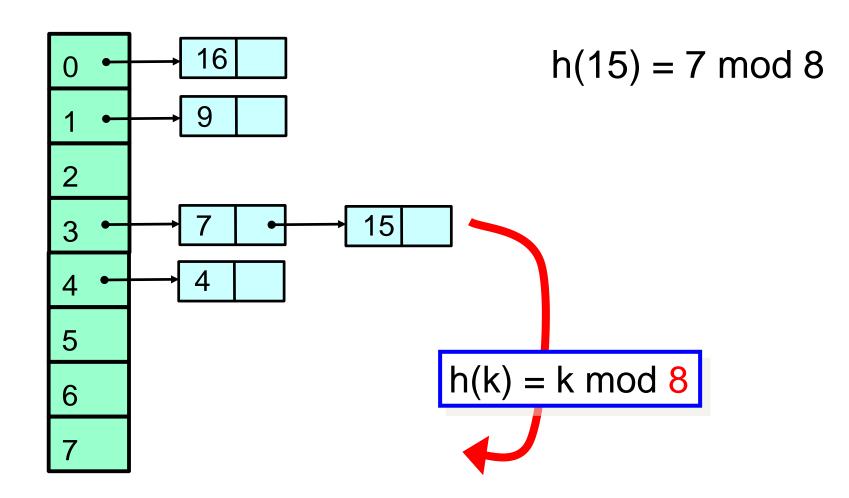
5 Items

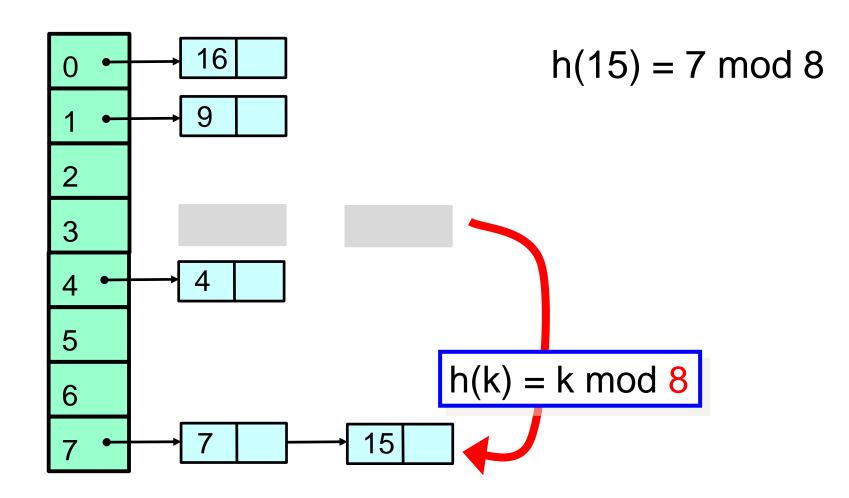












Fields

```
public class SimpleHashSet {
  protected LockFreeList[] table;

public SimpleHashSet(int capacity) {
  table = new LockFreeList[capacity];
  for (int i = 0; i < capacity; i++)
     table[i] = new LockFreeList();
}

Array of lock-free lists</pre>
```

Constructor

```
public class SimpleHashSet {
   protected LockFreeList[] table;

public SimpleHashSet(int capacity) {
   table = new LockFreeList[capacity];
   for (int i = 0; i < capacity; i++)
     table[i] = new LockFreeList();
}
...
Initial size</pre>
```

Constructor

```
public class SimpleHashSet {
   protected LockFreeList[] table;

public SimpleHashSet(int capacity) {
   table = new LockFreeList[capacity];
   for (int i = 0; i < capacity; i++)
     table[i] = new LockFreeList();
}

Allocate memory</pre>
```

Constructor

```
public class SimpleHashSet {
  protected LockFreeList[] table;
  public SimpleHashSet(int capacity) {
    table = new LockFreeList[capacitv];
   for (int i = 0; i < capacity; i++)
      table[i] = new LockFreeList();
```

Initialization

Add Method

```
public boolean add(Object key) {
  int hash =
   key.hashCode() % table.length;
  return table[hash].add(key);
}
```

Add Method

```
public boolean add(Object key) {
  int hash =
    key.hashCode() % table.length;
  return table[hash].add(key);
}
```

Use object hash code to pick a bucket

Add Method

```
public boolean add(Object key) {
  int hash =
    key.hashCode() % table.length;
  return table[hash].add(key);
}
```

Call bucket's add() method

Is Resizing Necessary?

- Constant-time method calls require
 - Constant-length buckets
 - Table size proportional to set size
 - As set grows, must be able to resize

Set Method Mix

- Typical load
 - 90% contains() //Read operation
 - 9% add ()
 - 1% remove()
- Growing is important
- Shrinking not so much

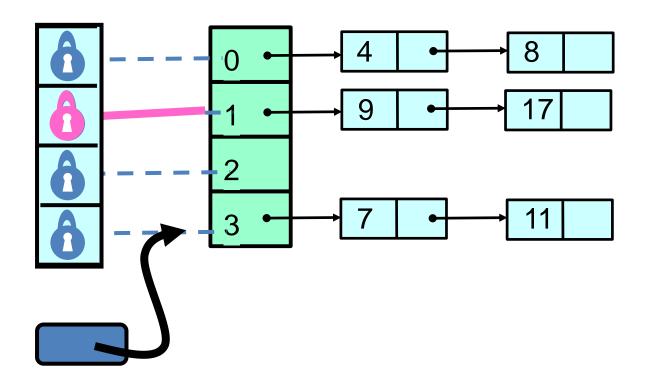
When to Resize?

- Many reasonable policies. Here's one.
- Pick a threshold on num of items in a bucket
- Global threshold
 - When ≥ ¼ buckets exceed this value
- Bucket threshold
 - When any bucket exceeds this value

Coarse-Grained Locking

- Good parts
 - Simple
 - Hard to mess up
- Bad parts
 - Sequential bottleneck

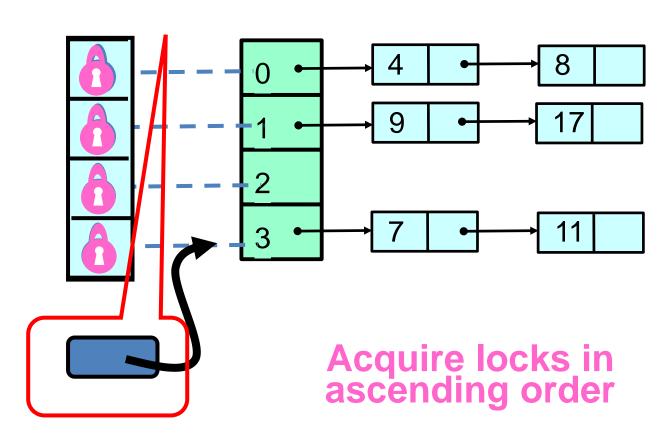
Fine-grained Locking



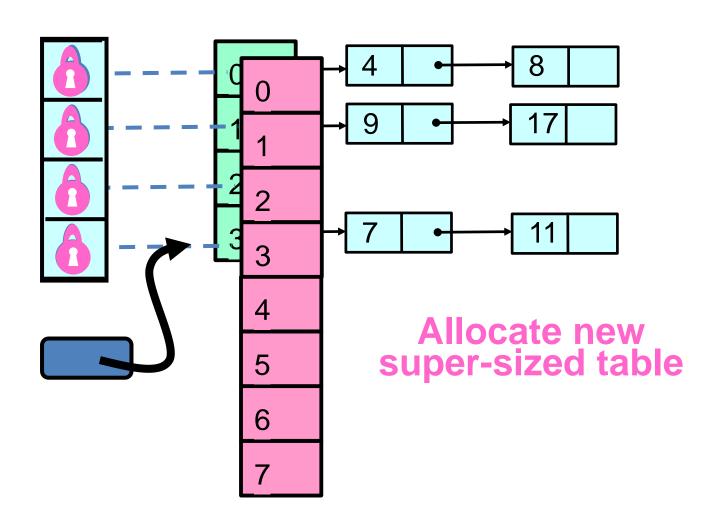
Each lock associated with one bucket

Resize This

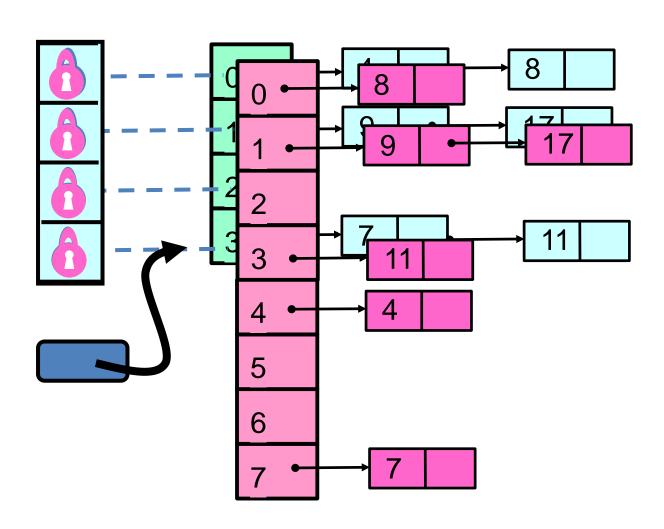
Make sure table reference didn't change between resize decision and lock acquisition



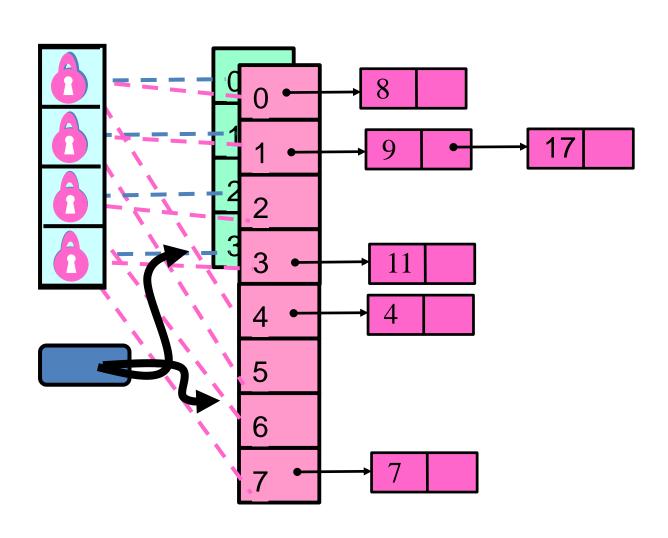
Resize This



Resize This



Striped Locks: each lock now associated with two buckets Resize This



Observations

- We grow the table, but not locks
 - Resizing lock array is tricky ...
- We use sequential lists
 - If we're locking anyway, why pay?

- Design and implement a Java multithreaded program
- To simulate a bank money transaction system for up to 10⁶ users.
- Name of the bank is Guwahati National Bank (GNB)
 - the bank have 10 branches at different location of India and each branch have 10 updaters.
- Initially, 10⁴ users will be there for each branch with a random amount of money in their accounts.
- Every updater of the GNB can be modeled as a separate thread.

- The updater gets a request
 - To Cash Deposit, Cash withdrawal (withdrawal amount should less than amount money in the account).
 - To Transfer Money from one customer account to another customer account. The source account and destination account may be in different branches of GNB.
 - to add a customer with some initial money in that account,
 added customer will be in the updater branch of GNB.
 - To delete a customer from the system (or close the account of a user of the GNBs from any branch).
 - To transfer customer account from one branch to another branch of GNB.

- Assume the probability of getting cash deposit, cash withdrawal, money transfer, add a customer, delete a customer, transfer a customer to updater are 0.33, 0.33, 0.33, 0.003, 0.003, and 0.004 respectively.
- Simulate up to 10⁶ transactions per updater to test
 - the correctness of your implementation.
 - Report the execution time of your simulation program.

- Suppose information about all the customer accounts of the GNB
 - is maintained by an array of linked lists.
 - Each linked list represent (or hold data of a) branch of GNB, so there are 10 linked list and the array of linked list is maintained by a hash data structure.
 - Every customer account number is represented by 10 digits and the first digit of the customer account number identify the branches of GNB.
- You to be used linked list and hash should be
 - thread-safe, and throughput should be high.
 - You are allowed to use any inbuilt data structure, locking protocol, synchronized functions for the same.

Thanks