



The RSA algorithm process.

The security of RSA is based on the belief that the encryption formula ($c = me \bmod n$) is a one-way function. The trapdoor that allows Bob to decrypt a Cipher text is the knowledge of factorization $n = pq$.

Example: Let $p=3$ and $q=11$ (both are prime numbers).

- Now, $n = p \cdot q = 3 \cdot 11 = 33$
- $\phi(n) = (p-1) \cdot (q-1) = (3-1) \cdot (11-1) = 2 \cdot 10 = 20$
- Value of e can be 7 since $1 < 7 < 20$ and $\gcd(20, 7) = 1$.
- Calculating $d = 7^{-1} \bmod 20 = 3$.
- Therefore, public key = $\{7, 33\}$ and private key = $\{3, 33\}$.

Suppose our message is $M=31$. You can encrypt and decrypt it using the RSA algorithm as follows:

Encryption: $C = (M^e) \bmod n = 31^7 \bmod 33 = 4$

Decryption: $M = (C^d) \bmod n = 4^3 \bmod 33 = 31$

Program:

```
import math

# step 1
p = 3
q = 7

# step 2
n = p*q
print("n =", n)

# step 3
phi = (p-1)*(q-1)

# step 4
e = 2
while(e<phi):
    if (math.gcd(e, phi) == 1):
        break
    else:
        e += 1

print("e =", e)

# step 5
k = 2
d = ((k*phi)+1)/e
print("d =", d)
print(f'Public key: {e, n}')
print(f'Private key: {d, n}')
```



```
# plain text
msg = 11
print(f'Original message:{msg}')

# encryption
C = pow(msg, e)
C = math.fmod(C, n)
print(f'Encrypted message: {C}')

# decryption
M = pow(C, d)
M = math.fmod(M, n)

print(f'Decrypted message: {M}')
```

Output:

```
n = 21
e = 5
d = 5.0
Public key: (5, 21)
Private key: (5.0, 21)
Original message:11
Encrypted message: 2.0
Decrypted message: 11.0
```

Result:

We gained knowledge of symmetric encryption and the RSA algorithm in this experiment. We also saw how the RSA algorithm was implemented in Python.

