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DEEMED TO BE UNIVERSITY
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CSE211 – Formal Languages and Automata Theory

U1L9 – Non-Deterministic Finite Automata

Dr. P. Saravanan

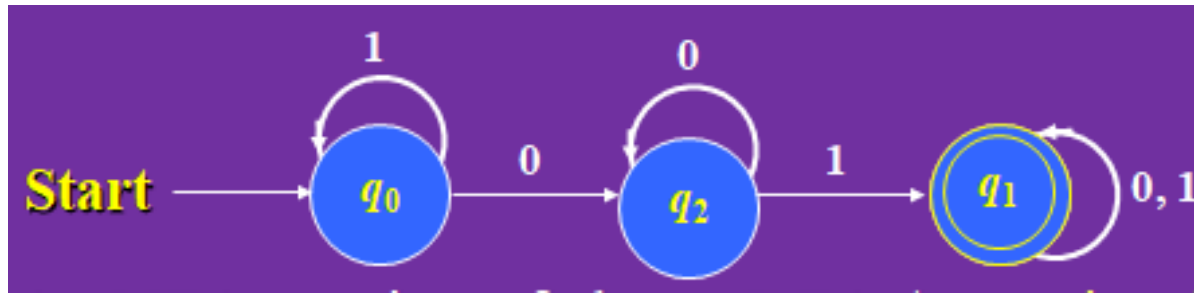
School of Computing
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Agenda

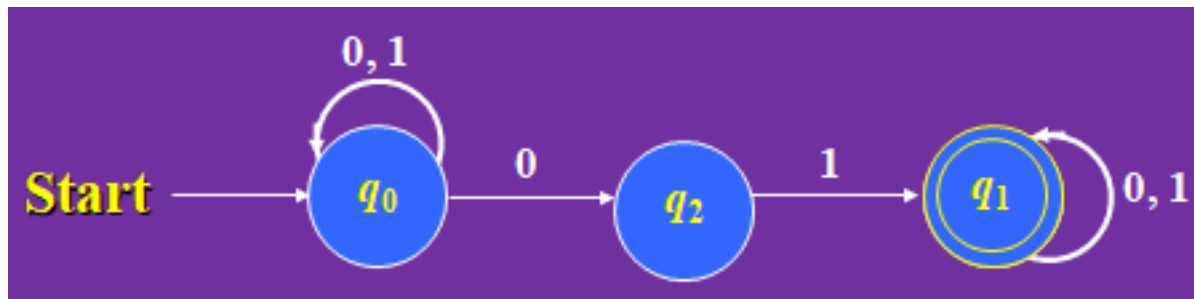
- Recap of previous class
- What is Non-deterministic Finite Automata(NFA)?
- Examples for NFA
- Definition of NFA
- Epsilon-NFA
- E-closure of a state in e-NFA
- Example for identifying e-closure

Nondeterministic Finite Automata (NFA)

- DFA for $L = \{x01y \mid x \text{ and } y \text{ are any strings of 0's and 1's}\}$

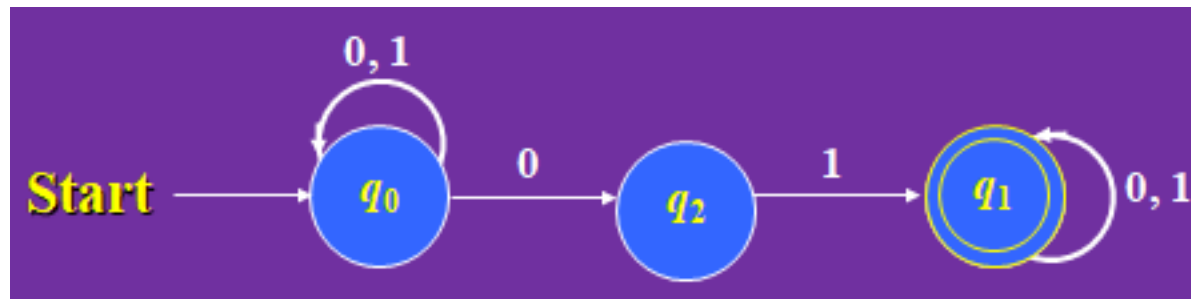


- NFA for the above DFA



Nondeterministic Finite Automata (NFA)

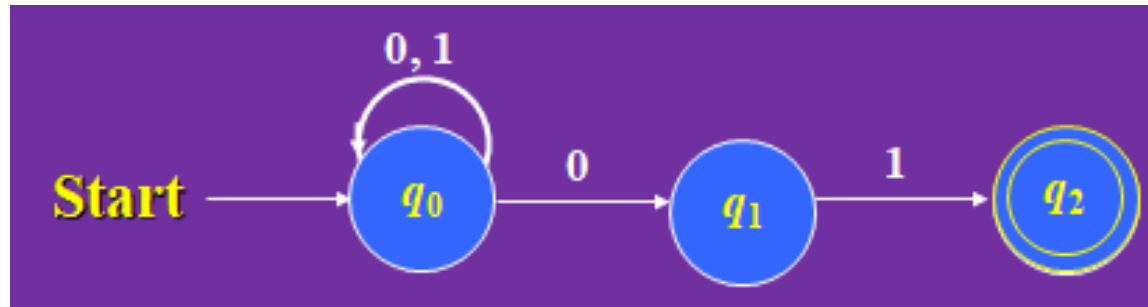
- Some properties of NFA's



- Some transitions may “die,” like $\delta(q_2, 0)$.
- Some transitions have multiple choices, like $\delta(q_0, 0) = q_0$ and q_2 .

NFA - Example

- Design an NFA accepting the following language
 $L = \{w \mid w \in \{0, 1\}^* \text{ and ends in } 01\}.$

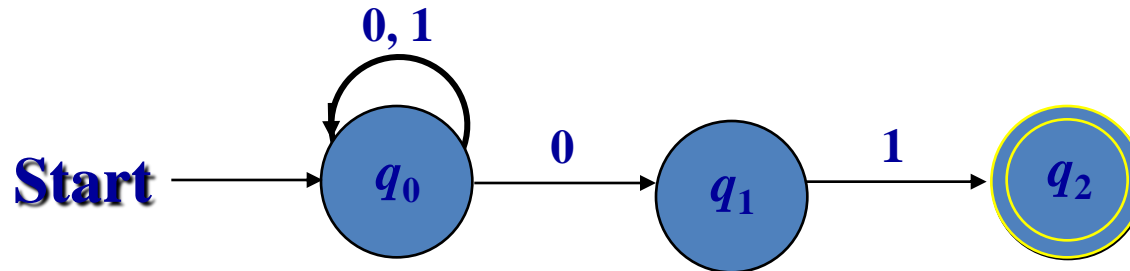


- Non-determinism creates many transition paths
- But if there is one path leading to a final state, then the input is accepted
- Check the input $x = 1000101$

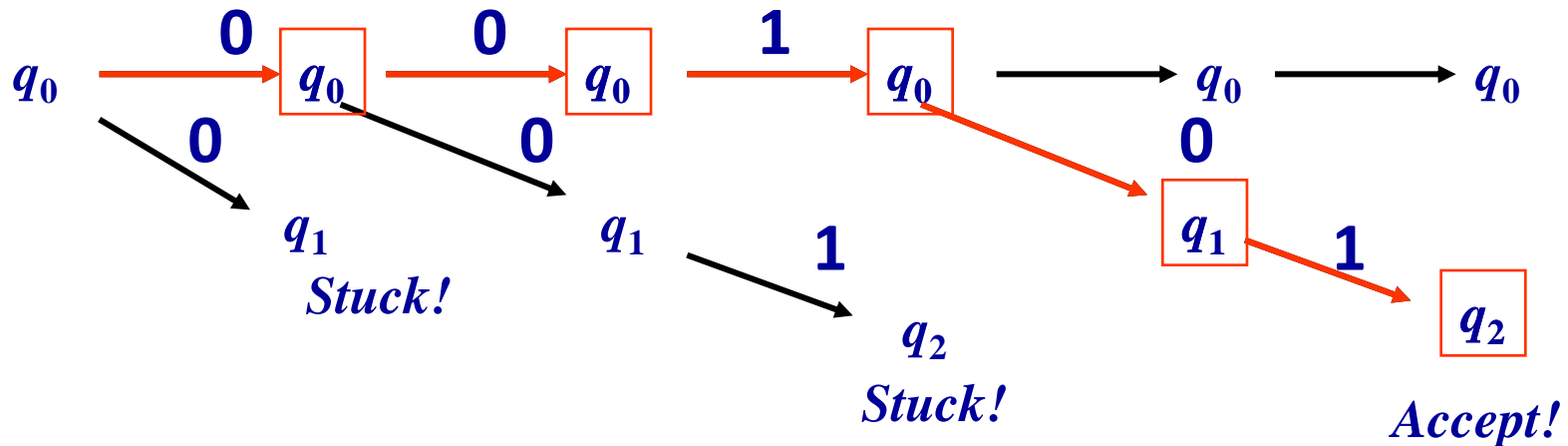
Dr.PS

NFA Example

■ Example



- When input $x = 00101$, the NFA processes x in the following way:



Nondeterministic Finite Automata

Definition of NFA

- An NFA **A** is a 5-tuple $A = (Q, \Sigma, \delta, q_0, F)$ where
 - Q = a finite (nonempty) set of states;
 - Σ = a finite (nonempty) set of input symbols;
 - q_0 = a start state;
 - F = a set of (nonempty) final or accepting states;
 - $\delta: Q \times \Sigma \rightarrow 2^Q$ is a **transition function**

Finite Automata with Epsilon-Transitions

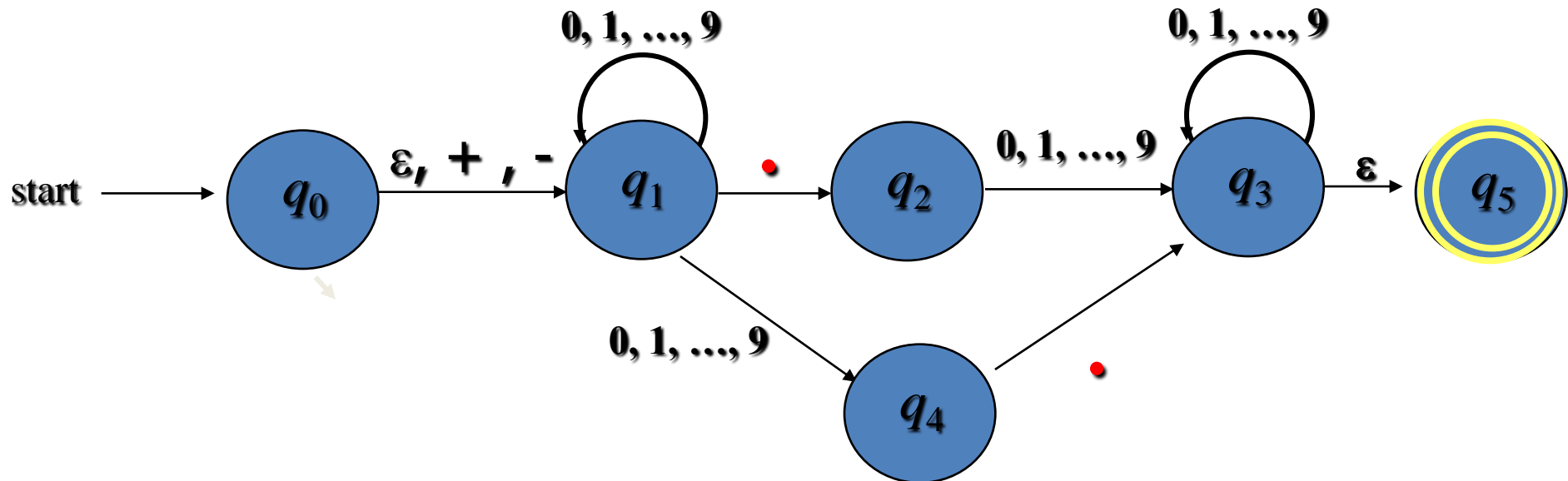


■ Use of ϵ -transitions

- We **allow** the automaton to **accept** the empty string ϵ .
- This means that a transition is allowed to occur **without** reading in a symbol.
- The resulting NFA is called ϵ -NFA.
- It adds “**programming (design) convenience**” (more intuitive for use in designing FA's)
- But,
 - ϵ cannot be used as an input symbol, but can be accepted to yield a transition!
 - $\delta(q_i, \epsilon)$ is defined for every state q_i

Finite Automata with Epsilon-Transitions

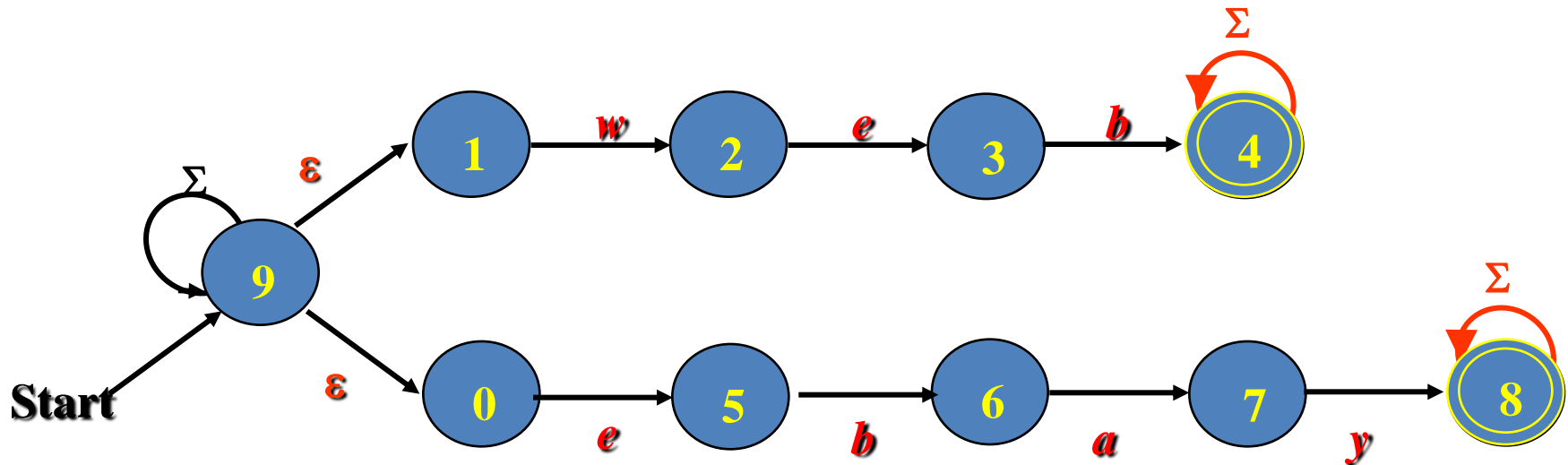
- Example : an ϵ -NFA accepting decimal numbers like 2.15, .125, +1.4, -0.501...



- To accept a number like "+5." (nothing after the decimal point), we have to add q_4 .

Finite Automata with Epsilon-Transitions

- Example : a more intuitive ε -NFA for Example 2.14



Design an NFA with ε -transitions.

Finite Automata with Epsilon-Transitions

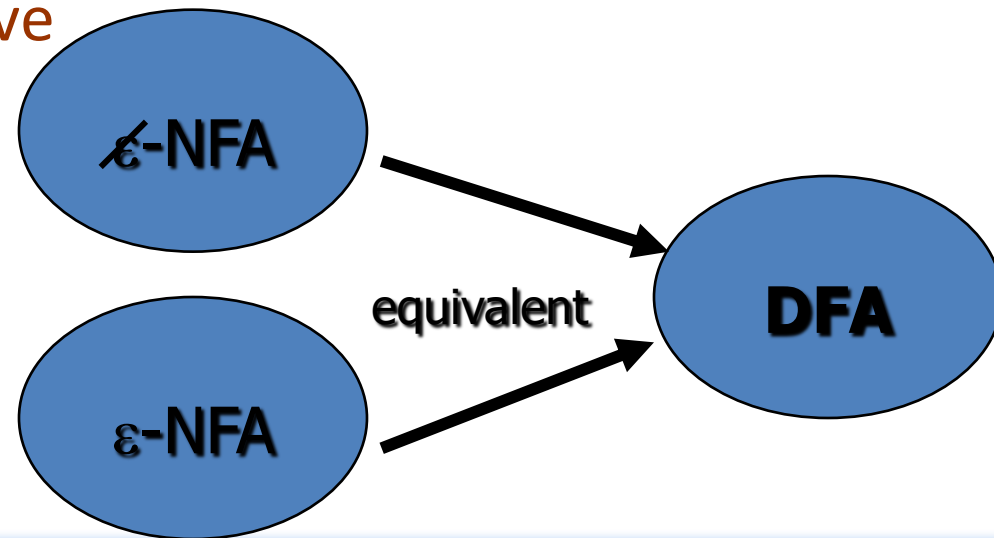
■ Formal Notation for an ε -NFA

- Definition: an ε -NFA A is denoted by $A = (Q, \Sigma, \delta, q_0, F)$
 - Q = a finite (nonempty) set of states;
 - Σ = a finite (nonempty) set of input symbols;
 - q_0 = a start state;
 - F = a set of (nonempty) final or accepting states;
 - $\delta: Q \times \Sigma \cup \{\varepsilon\} \rightarrow 2^Q$ is a **transition function**

Review

■ 3 Types of Automata

- DFA ---good for soft/hardware implementation
 - $\delta: Q \times \Sigma \rightarrow Q$ is the **transition function**
- NFA ---intermediately intuitive
 - $\delta: Q \times \Sigma \rightarrow 2^Q$ is the **transition function**
- ϵ -NFA ---most intuitive
 - $\delta: Q \times \Sigma \cup \{\epsilon\} \rightarrow 2^Q$ is the **transition function**



Summary

- What is Non-deterministic Finite Automata(NFA)?
- Examples for NFA
- Definition of NFA
- Epsilon-NFA
- E-closure of a state in e-NFA
- Example for identifying e-closure

References

- John E. Hopcroft, Rajeev Motwani and Jeffrey D. Ullman, *Introduction to Automata Theory, Languages, and Computation*, Pearson, 3rd Edition, 2011.
- Peter Linz, *An Introduction to Formal Languages and Automata*, Jones and Bartle Learning International, United Kingdom, 6th Edition, 2016.

Next Class:

Converting NFA or e-NFA to DFA

THANK YOU.