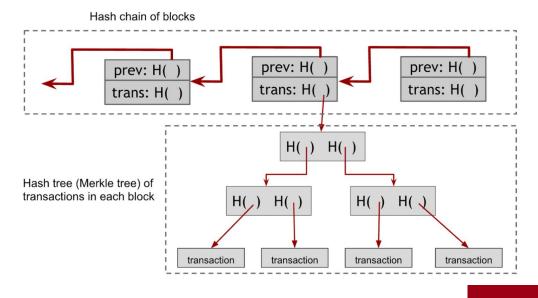
Table of contents

- Bitcoin miners
- Mining hardware
- Mining pool
- Changing protocol

1

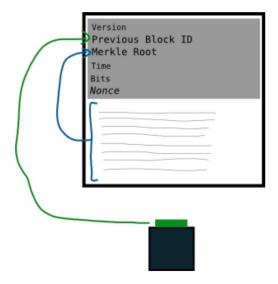
- The task of Bitcoin miners
 - Maintain blockchain
 - Validate Tx
 - Listen for new Tx
 - Assemble a candidate block
 - Find a valid nonce for the candidate block
 - Broadcast
 - Hope block is accepted
 - o Profit
 - Validate blocks from the network
 - Listen for new blocks
 - Validate these blocks

- Bitcoin block
 - O Bitcoin blockchain contains two different hash structures
 - Hash chain of blocks
 - Merkle tree of Tx within blocks



Bitcoin miners

- Bitcoin block
 - O Block header is a summary (metadata) of the data in the block



4

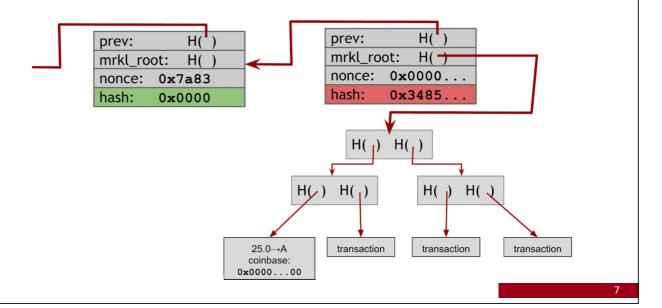
Bitcoin block

Field	Description
Version	Bitcoin version/rules-set
Previous Block ID	Hash of previous block, where chain is extending
Merkle Root	All Tx in this block hashed (single-line summary of all Tx) together
Time	Unix time, when a miner is trying to mine (hashing) this block the
Bits	A shortened version of target
Nonce	Miners change this 4 bytes (32 bits; 2 ³² possibilities starting from 0) value to get a hash of block header (block hash) below target

.

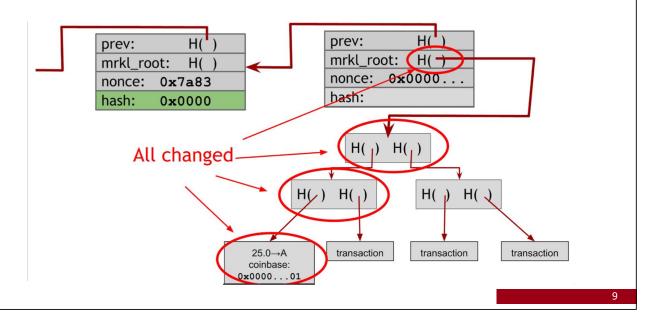
- Finding a valid block
 - Miners
 - Identify latest (previous) block
 - Compute the merkle root
 - Pick a set of pending Tx
 - O Take care of max size of block
 - Validate these Tx
 - Compile a coinbase Tx
 - Compute target
 - Pick a nonce value
 - Hash the block
 - Check if hash satisfies target
 - o If yes, broadcast
 - Otherwise, try again with a new nonce value

Finding a valid block



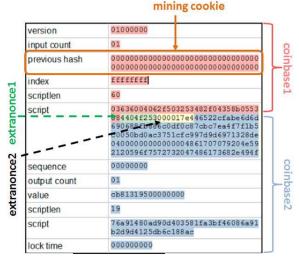
- Finding a valid block
 - O What if we exhaust nonce values?
 - o Changeable elements in block header
 - Nonce
 - We just exhausted nonce values
 - Time (rollnTime in getwork)
 - Limited scalability
 - Merkle root

- Change merkle root
 - Change Tx, costly
 - Change coinbase Tx



- Change coinbase Tx
 - O In regular Tx
 - Take scriptPubKey from previous Tx and scriptSig from current Tx
 - O Coinbase Tx has single blank input
 - No reference to previous output with scriptPubKey
 - scriptSig in current coinbase Tx can be arbitrary bytes
 - Of course, avoiding caveats (OP_CHECKSIG)
 - extraNonce solution
 - o extraNonce isn't part of the protocol
 - There is no extraNonce field in blocks or Tx

- Change coinbase Tx
 - o coinBase1/2 represent input/output (overlapping)
 - Change extraNonce (extraNonce1 || extraNonce2)
 - extraNonce1 is unique (pseudo-random) in Stratum (we'll see it later)
 - Only option is to change 4 bytes extraNonce2



11

- Change coinbase Tx
 - Overall increment/change extraNonce
 - Coinbase Tx = coinBase1 || extraNonce1 || extraNonce2 || coinBase2
 - Compute H(coinbase Tx)
 - Rebuild merkle root
 - Reset nonce, recompute hash

- Finding a valid block
 - o Is everyone solving the same puzzle?
 - No!
 - Miners choose Tx and build candidate block independently
 - What if different miners choose identical Tx?
 - Blocks would still differ
 - Miners specify their own address in coinbase Tx
 - What if they share the address?
 - Pooled mining
 - O Try different ranges of nonce, avoid duplicate work

13

- When you create and broadcast a Tx
 - O It goes to a node
 - Node stores/queues it in its mempool
 - Broadcast further
 - Tx reside in mempool until miner picks it
 - Prioritize tx based on fees/age
 - More fee, more chances to be picked
- What if miner never picks it?
 - o May be you paid a very low/no fee
 - Miner is least interested to pick it
 - Can Lincrease fee?
 - Replace By Fee (RBF)

- RBF
 - A node policy
 - Allows an unconfirmed Tx in a mempool to be replaced with a different Tx
 - Spends at least one of same inputs
 - Pays a higher Tx fee

15

- RBF
 - o Full RBF
 - Unconditionally allows a Tx to replace older ones as long as it pays a sufficient fee
 - Opt-in RBF
 - Only allows replacement when Tx being replaced have explicitly signaled to allow replacement
 - Via nSequence/sequence field
 - Both full and opt-in RBF may lead to double spend

RBF

- o First-seen-safe RBF
 - Only allows replacement if replacement/new Tx keeps same outputs as Tx being replaced
- Delayed RBF
 - Allows Tx to be replaced unconditionally, but only after a given number of blocks have been mined since Tx being replaced were first seen by node

17

- RBF
 - o nSequence
 - 4 bytes field in an Tx input
 - Usually only shown in rawTx
 - Send a Tx with a lock_time & sequence number 0
 - Tx is then not considered "final" by network
 - Can't include in a block until lock time
 - Before lock time expires
 - Replace Tx with as many new versions as you want
 - Newer versions have higher sequence numbers
 - To lock Tx permanently
 - Set sequence number to 0xFFFFFFF
 - Tx is considered final, even if lock_time has not reached

- Orphan Tx
 - When Tx are transmitted across network
 - May arrive to nodes in different order
 - A child might arrive before parent
 - Upon seeing a Tx (child)
 - That references a Tx (parent) that is not yet known
 - Nodes don't reject child
 - Put it in temporary orphan Tx pool
 - Where orphan Tx wait for parent's arrival
 - When parent arrives
 - Child Tx is retrieved, revalidated recursively, inserted in mempool

- Orphan block
 - When a fork occurs
 - A block may not be included in main chain
 - Block becomes orphaned
 - All Tx in it simply go back in node's queue/pool
 - o Revalidated, if their input has been spent
 - If so, evicted
 - Wait to be added to next block

Mining hardware

- Computing SHA-256 hash function is the core of miners' tasks
- CPU mining
 - o First generation mining was done on general purpose CPUs
 - O Simply compute SHA-256 in software, check the result
 - o Speed?
 - ~10-20 MH/s
 - Mining on a regular CPU at current difficulty level is impossible

21

Mining hardware

- GPU mining
 - Designed for processing videos/graphics
 - Large number of floating point units
 - Not used at all in SHA-256
 - Poor cooling characteristics (especially, in farms)
 - High electricity computation
 - o Speed?
 - ~200-300 MH/s
 - Ok, but not sufficient

Mining hardware

- FPGA mining
 - o Better performance over GPUs
 - "Bit fiddling" operations is trivial to specify
 - Use all transistors on the card
 - o Cooling is easier
 - o Speed?
 - ~1 GH/s
 - Better, but still not sufficient
 - Cost-per-performance was marginally improved over GPUs
 - Short-lived phenomenon

23

Mining hardware

- ASIC mining
 - Application-Specific Integrated Circuits
 - Optimized for the sole purpose of mining
 - Lifetime of ASICs was quite short initially
 - Due to rapidly increasing network hash rate
 - Now ASIC equipments have longer life time
 - After growth rate of hash power has stabilized

Mining hardware

- Economics of mining haven't been favorable to small miner
 - Mining cost = H/w cost + operation cost (electricity, cooling, etc.)
- Today, professional mining
 - o Mining farms
 - Mining pools
- Future, find alternatives

25

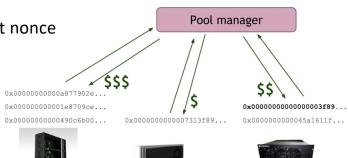
- A group of miners form a pool
- Attempt to mine a block
- Use a designated coinbase recipient
 - O Recipient is called pool manager
- No matter who actually finds nonce
 - o Pool manager will receive reward

- Pool manager takes that revenue
 - Distribute it among participants
 - Based on amount of work done by participants
 - Takes a cut
- How does pool manager know if miners are actually working?
 - o If yes, how to estimate amount of work?
 - Miners can cheat
 - Show proof?

Mining pool

- Mining shares
 - Miners probabilistically prove amount of work done
 - Consider a target begins with 30 zeros
 - Block's valid hash must have >=30 zeros
 - Miners find many hashes, some are close to target
 - Miners show these nearly valid hashes as proof
 - Threshold can be set, e.g., >=20 zeros
 - More work, more profit

No bonus for finding right nonce









- Mining shares
 - Pay-per-share model
 - Pool manager pays a flat fee for shares above threshold
 - Best for miners
 - Pool manager absorbs the risk
 - Must pay miners even if right nonce is not found
 - Proportional model
 - Payment depends on whether or not pool actually finds right nonce
 - If found, reward is distributed based on contributions
 - Otherwise, no/low reward
 - Miners do pool hopping to optimize revenue

29

- Are mining pools good?
 - o Pros
 - Easier for smaller miners to get involved
 - Easier to upgrade network
 - Central pool manager can force miners to upgrade
 - o Cons
 - Issue of centralization
 - 51% mining pools
 - All are miners, less full nodes

- Mining pool protocols
 - Three major protocols
 - getwork
 - getblocktemplate (gbt)
 - Stratum

31

- Mining pool protocols
 - o getwork
 - Original mining method
 - HTTP-based requests and responses
 - JSON Remote Procedure Calls (RPC) to request a block
 - Server gives only block headers to client
 - No Tx or way to modify block (except nonce)
 - Client is limited to try all nonce values
 - If exhausted, get more work from server
 - High bandwidth usage for modern mining h/w
 - Worst approach in today's scenario
 - rollnTime extension
 - Allows a client to generate own work by modifying timestamp on block header
 - Limited scalability

- Mining pool protocols
 - o getblocktemplate (gbt)
 - HTTP-based
 - JSON RPC for communication between client and poolserver
 - Server gives templates to client
 - Client uses it to generate own work
 - Full block data is given
 - Allows the client to modify the block
 - Miner can choose which Tx to include in block
 - Very high bandwidth usage
 - Well documented, detailed specifications

33

- Mining pool protocols
 - o Stratum
 - Not HTTP-based
 - Direct TCP connections
 - JSON RPC for communication between client and poolserver
 - Uses parts of gbt's message specification
 - Server gives templates to client
 - Client uses it to generate own work
 - Block headers and only Tx hashes are given
 - Change extraNonce
 - Least bandwidth usage
 - Very fast, efficient switching to new work
 - No real specification

Mining pool protocols

o Stratum

Changing protocol

- How to introduce new features in Bitcoin protocol?
 - O Just release an update and ask all nodes to upgrade?
 - Not really!
 - Not every node would upgrade
 - Some nodes would fail to get or to get it in time
 - Two types of changes
 - One causes a hard fork
 - Other causes a soft fork

Changing protocol

- Hard fork
 - o Introduces new features that were previously considered invalid
 - New nodes would accept blocks as valid that old nodes would reject
 - After majority of nodes have upgrade
 - Longest branch will contain blocks that are invalid for old nodes
 - So, old nodes will consider their (shorter) branch to be longest valid branch
 - Such changes cause a hard fork
 - Blockchain splits
 - Branches will never join together again

37

Changing protocol

- Hard fork
 - o E.g., fixing the bug in *OP_CHECKMULTISIG* would need a hard fork
 - What can we add with a hard fork
 - Adding new opcodes to protocol
 - Changing limits on block/Tx size

Changing protocol

- Soft fork
 - Adding features that make validation rules stricter
 - Old node would accept all blocks while new nodes would reject some
 - After majority of nodes have upgraded
 - New nodes will enforce new, tighter rules
 - Old miners might mine invalid blocks
 - Include some Tx invalid under new rules
 - But, they will figure out soon after their blocks are rejected
 - Without understanding reason
 - Ask their operator to check/upgrade
 - Such changes cause a soft fork
 - There will be many small, temporary forks

39

Changing protocol

- Soft fork
 - o E.g., introduction of P2SH scripts
 - Old nodes would still verify a valid P2SH correctly
 - Hash the data value, check if hash matches value in scriptPubKey
 - OP_EQUAL pushes 1 on top of stack
 - Old nodes don't know (now required) step of running that value itself
 - What can we add with a soft fork
 - New cryptographic schemes
 - Some extra metadata in coinbase parameter
 - Some specific formatting to a field
 - o 2016 "The DAO" hack
 - Soft/hard fork proposals

40