

## Exercise 04

Sunday, November 12, 2023 6:02 PM

### Problem 1 to hand in: Algorithm Example

Write an algorithm for the following problem:

Given a positive integer  $n$ , return a list that contains all prime numbers  $p$  with distance 1 to a power of two,  $2 \leq p \leq n$ .

a) Provide a representative test example.

① Input :  $n = 5$

Return should be [3]

② Input :  $n = 100$

Return should be [3, 5, 7, 17, 31]

b) Describe an algorithm that solves this problem intuitively.

Intuitively, the algorithm should suffice 3 requirements:

① main: able to iterate number from 2 to  $n$

(A simple for-loop might do the trick)

② prime number: able to distinguish a prime number

(An input: number, output: boolean function might satisfy the need)

③  $2^k - 1 / 2^k + 1$ : able to judge if a number

has a distance of 1 to any power of 2

(Similar to ②)

The two functions ② and ③ shall be integrated in the loop ①. When both conditions are fulfilled, we store the number in a array and in the end return the array after iteration.

c) Formulate the algorithm in Pseudocode. For a better learning effect, use as few Python-specific functions as possible.

Here we suppose the input is correct (positive integer  $\geq 2$ )

main( $n$ ):

$A \leftarrow \emptyset$

$i \leftarrow 2$

while True:

if  $2^i - 1 > n$ :

return  $A$

else:

if  $2^i + 1 < n$ :

$A = \text{add\_prime}(A, 2^i + 1)$

$A = \text{add\_prime}(A, 2^i - 1)$

~~add\_prime( $A, p$ ):~~

~~for  $i \leftarrow 2$  to  $p-1$ :~~

~~if  $p \bmod i \equiv 0$ :~~

~~return  $A$~~

add\_prime( $A, p$ ):

for  $i \leftarrow 2$  to  $(p-1) \text{ div } 2 + 1$

if  $p \bmod i \equiv 0$ :

return  $A$

note:

no need to go through

all the numbers  $< p$

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if p mod i ≡ 0 :
    return A
return A ∪ {p}

```

```

if p mod i ≡ 0 :
    return A
return A ∪ {p}

```

no need to  
go through  
all the numbers < p

d) Analyse the asymptotic worst-case running time of your algorithm.

In the worst case scenario:

The while-loop in the main algorithm has a running time in

$$i = \log_2(n+1) \in O(\log_2 n),$$

the add-prime algorithm has a running time

(suppose the  $p$  is a prime number) in

$$O(p) = O(2^i) = O(2^{\log_2(n)}) = O(n) \quad O\left(\frac{p}{2}\right) \approx O(n)$$

So in combination, the algorithm has a  
running time in  $O(n \log_2 n)$

e) Provide a proof sketch that the algorithm is correct.

1. Prove that loop invariant in add-prime is correct:

SE $[i]$ : at the beginning of the  $i$ -th iteration,

$p$  is proven not to be divisible by  $2, \dots, i-1$ .

2. Prove that add-prime returns  $A$  if  $p$  is composite num.

and  $A \cup \{p\}$  if  $p$  is prime num.

3. Prove the loop invariant in the main algorithm is correct:

SE $[i]$ : at the beginning of each loop,

$2^{i-1} - 1 < n$  and  $A$  contains prime numbers in

a form of  $2^k + 1$  or  $2^k - 1$  with  $1 \leq k \leq i-1$

4. Prove the main algorithm indeed return the required list.