

Problem 1: We can use array to do this:

① initialize m :

create an array values of size m
set all elements in values to 0

• The running time is $O(m)$ because it creates an array of size " m ", which takes linear times related to the size of array.

② update (i, x) :

values $[i] \leftarrow x$

updateValues $[i] \leftarrow x$

updateCount $[i] += 1$

latestUpdateIndex $\leftarrow i$

• The running time is $O(1)$ because the update operation involves direct access to the array index.

③ get (i) :

return values $[i]$

• The running time is $O(1)$ because accessing the value at a specific index in an array is a constant-time operation.

④ undo():

$i \leftarrow \text{latestUpdateIndex}$

values $[i] \leftarrow \text{updateValues}[i]$

updateCount $[i] -= 1$

The running time is $O(1)$ because it directly accesses the information about the latest update.

⑤ `reset()`:

`values ← [0] · len(values)`

`updateValues ← [0] · len(updateValues)`

`updateCount ← [0] · len(updateCount)`

· The running time is $O(1)$ because it involves creating new arrays and the size of the array is constant.

⑥ `count_undo()`:

`return sum(updateCount)`

· The running time is $O(1)$ because it involves summing the elements in the "updateCount" array, which is constant time.