

## Exercise Sheet 6

for the lecture on

## Advanced Programming and Algorithms

Submission until **Monday, 4th December, 12:30 pm.**

Discussion in the exercise classes on 11th, 12th, and 15th December, 2023.

### Problem 1 to hand in: *Data Structure*

Describe a data structure that can store  $m$  values together with an update history and undo function.

The following methods should be managed by the data structure:

- `initialise( $m$ )`: set up an empty instance that can manage  $m$  values (that are initially set to 0) in  $\mathcal{O}(m)$ ,
- `update( $i, x$ )`: change the  $i$ th value to  $x$  in  $\mathcal{O}(1)$ ,
- `get( $i$ )`: return the current  $i$ th value  $\mathcal{O}(1)$ ,
- `undo()`: reset the latest update in  $\mathcal{O}(1)$  (a redo is not needed),
- `reset()`: reset all values to 0 and delete the update history in  $\mathcal{O}(1)$ ,
- `count_undo()`: count the number of available updates that can be undone in  $\mathcal{O}(1)$ .

Describe each of these methods (intuitively or in pseudocode) and explain why their running times are fulfilled.

### Problem 2 as a programming exercise: *Compound Data Types and Numpy*

Work on the current jupyter notebooks `lecture_06_compound_types` and `lecture_06_numpy` in order to learn how to use collective types such as `list` and `set` as well as numpy arrays.

### Problem 3 for discussion: *Set Problem*

Design an algorithm to solve the following problem: Given a set  $U = \{1, \dots, m\}$  and given sets  $S_1, S_2, \dots, S_n$  each of which is a subset of  $U$ , is there a pair of sets  $S_i, S_j$  that cover  $U$  completely, i.e.,  $S_i \cup S_j = U$ ?

**Problem 4 for discussion:** *Universal Hashing*

Let  $U$  be a set of possible keys,  $S \subseteq U$  be an arbitrary set of currently used keys with  $|S| = n$ . Moreover, let  $H$  be a universal family of hash functions mapping from  $U$  to  $T = \{0, \dots, m-1\}$ ,  $m \leq |U|$ .

Show that, for  $h \in H$  chosen uniformly, the total number of collisions is less than  $\frac{n(n-1)}{m}$  with probability at least  $\frac{1}{2}$ .

*Hint: Use an indicator random variable  $X_{xy}(h)$  that is 1 if  $h(x) = h(y)$ , and 0 otherwise. Show that the expected number of collisions is  $\frac{n(n-1)}{2m}$ . Then use Markov's inequality.*