



# LINKEDLIST

School of Artificial Intelligence

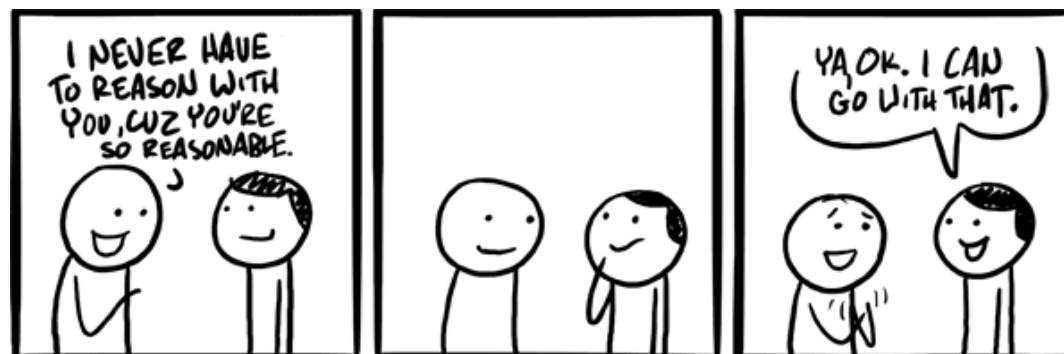
# 3RS

- Relax
- Reasonable
- Respect



RELAX

because sometimes a galactic war has to take a break.



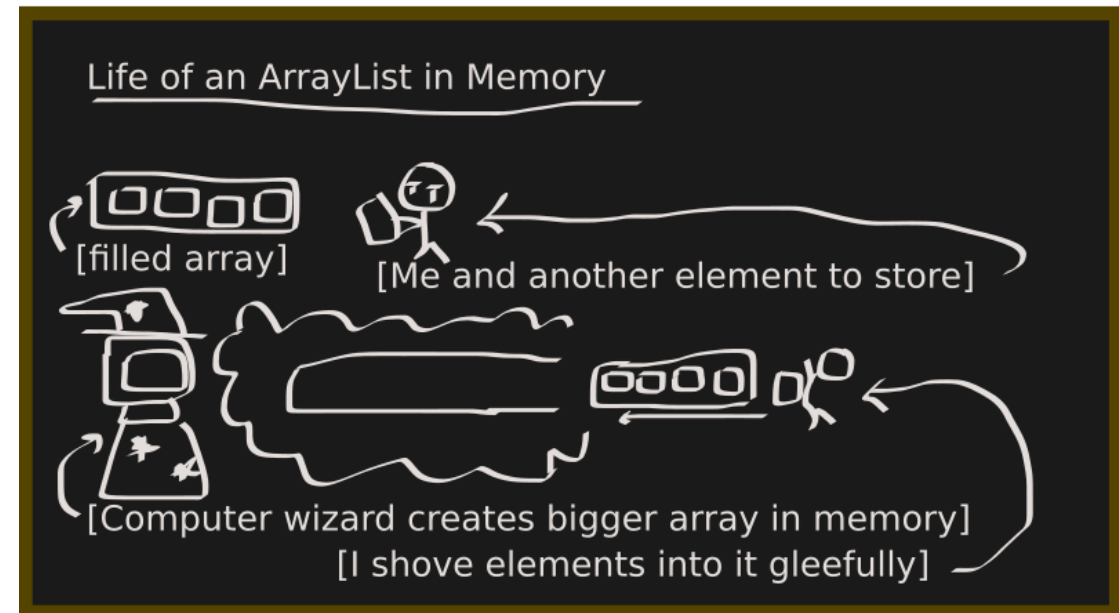
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Respect(尊重)

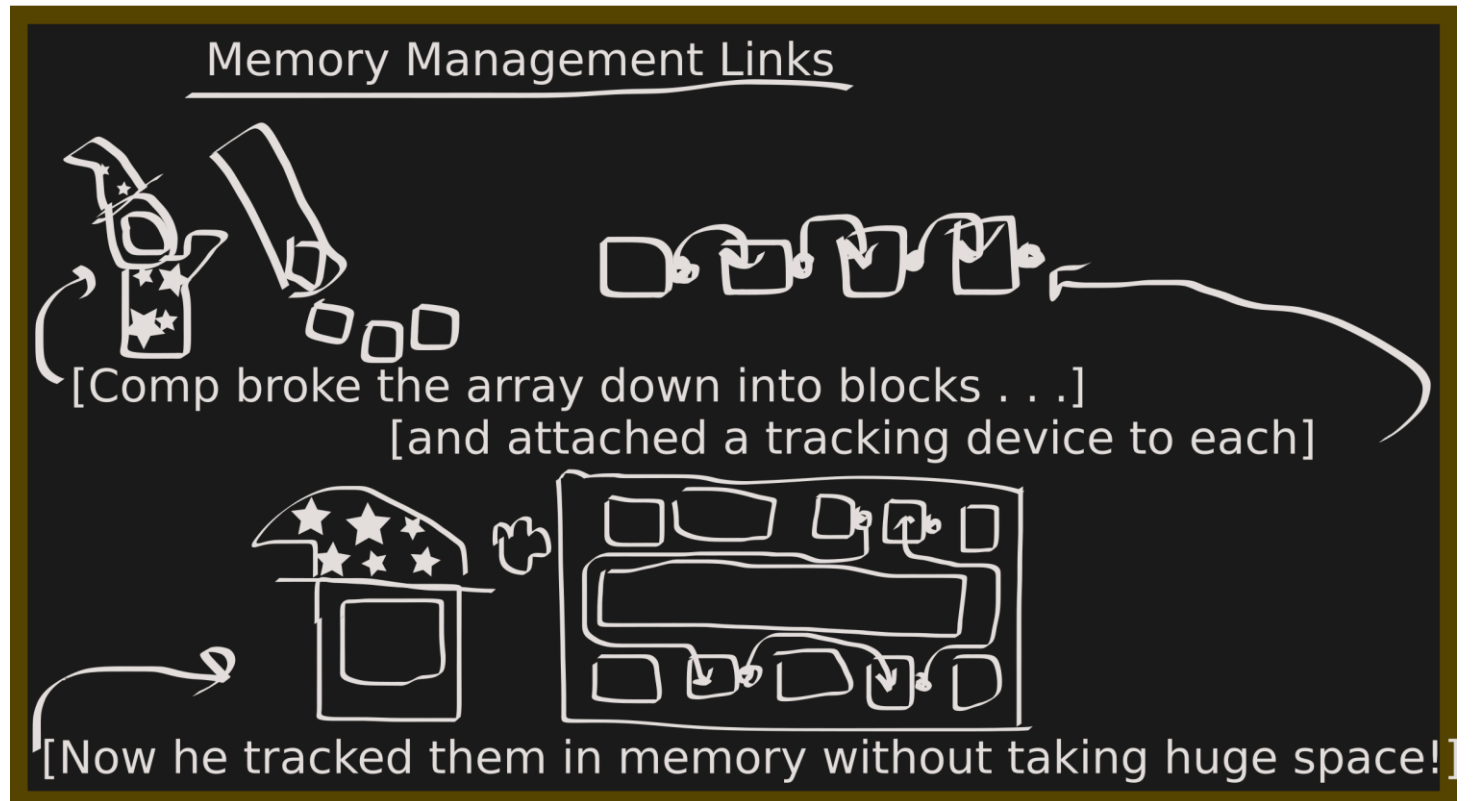
# PREVIOUSLY ON DS&A

- Array
  - Referential Array
  - Compact Array
- Sorting
  - Insertion Sort
  - Merge Sort
- Stack
- Queue
  - Implemented with circular array
- Deque
  - Double ended queue



# UP NEXT

- LinkedList (链表)
- Singly Linked Lists
- Circularly Linked Lists
- Doubly Linked Lists



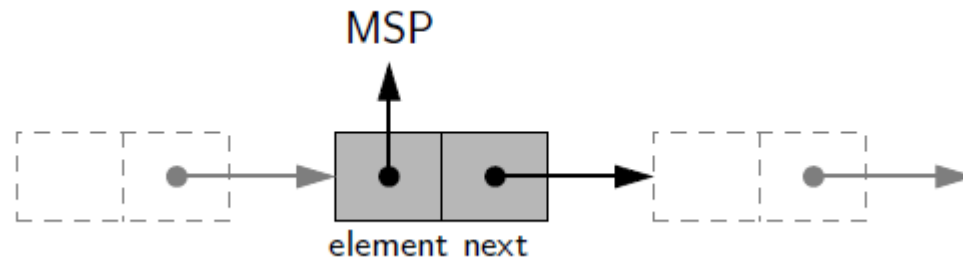
# LINKED LISTS (链表)

- Disadvantages of a dynamic array
  - The length of a dynamic array might be longer than the actual number of elements that it stores
  - Amortised bounds for operations may be unacceptable in real-time systems
  - Insertions and deletions at interior positions of an array are expensive
- Linkedlist
  - Stores values as an array does
  - Distributed representation
    - A lightweight object, known as a **node**, is allocated for each element
  - **Node** maintains a reference to its element + reference to its neighbouring **node**



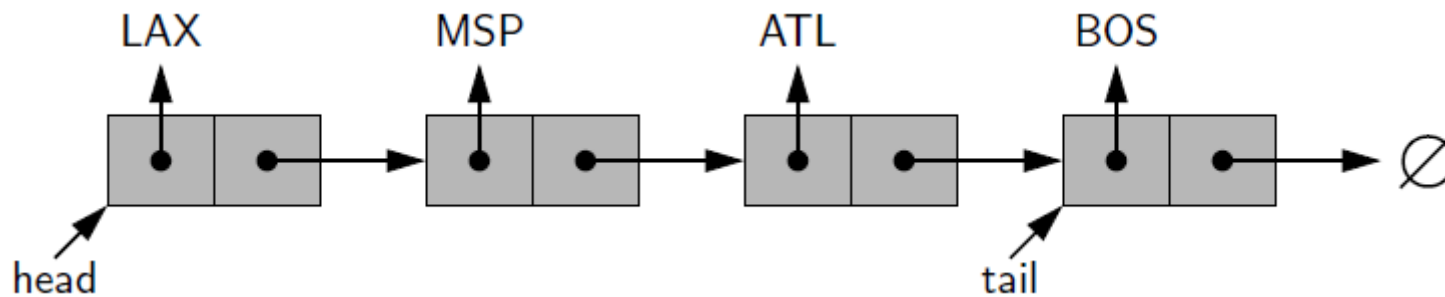
# SINGLY LINKED LISTS (单链表)

- Singly linked list
  - Collection of nodes that form a linear sequence
  - Each node:
    1. Reference to an object
    2. Reference to the *next* node



# SINGLY LINKED LISTS (单链表)

- Head (首) : first node in the linked list
- Tail (尾) : last node in the linked list
  - None as its *next* reference
- Linkedlist traversing (遍历):
  - Starting at the head and following the *next* reference until the tail
  - Often referred to as *link hopping* or *pointer hopping*



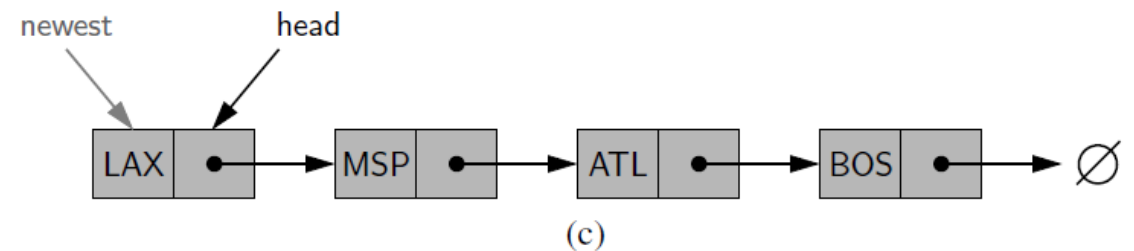
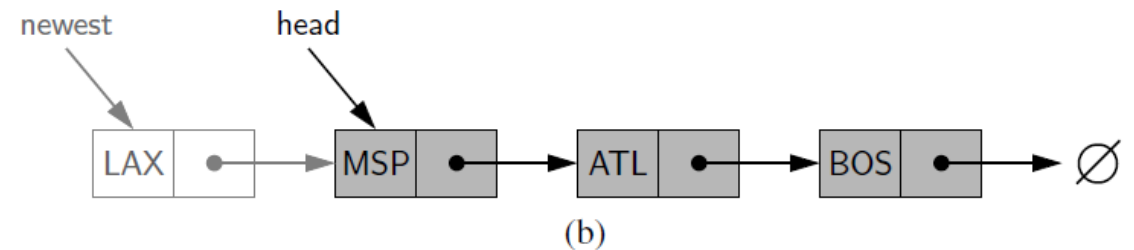
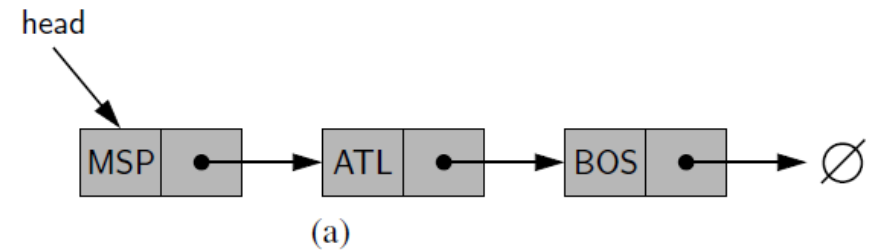
# SINGLY LINKED LISTS (单链表)

- In memory representation
  - Each node as a unique object
    - Reference to its element (value)
    - Reference to the *next* node (or None)
  - An object maintains the linkedlist
    - Reference to the head of the list – Why?
    - Reference to the tail of the list – why?
    - Number of elements in the linkedlists



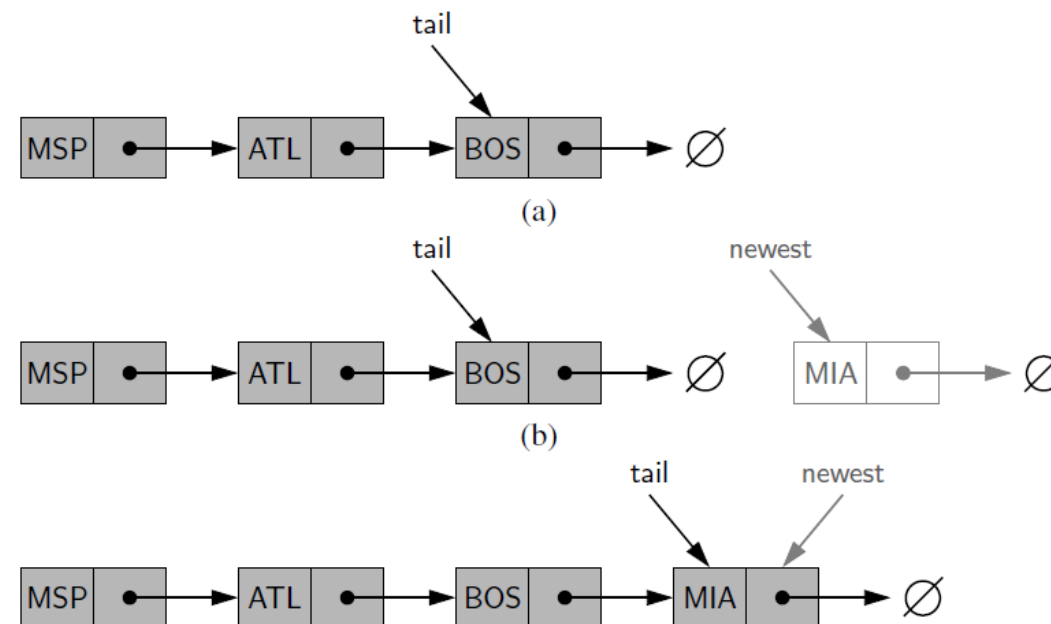
# SINGLY LINKED LISTS (单链表)

- Inserting an element at the head of a singly linked list
  1. Create a new node, with its value
  2. Set the *next* reference of the node to the head of the linkedlist
  3. Update the linkedlists' head
  4. Update the linkedlists' size



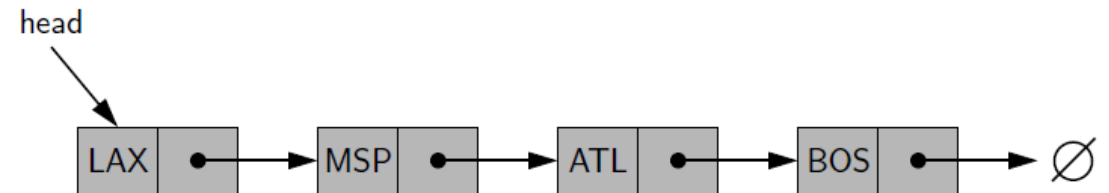
# SINGLY LINKED LISTS (单链表)

- Inserting an element at the tail of a singly linked list
  1. Create a new node, with its value
  2. Set the *next* reference of the node to None
  3. Set the *next* reference of the tail of the linkedlist to the new node
  4. Update the linkedlists' tail
  5. Update the linkedlists' size

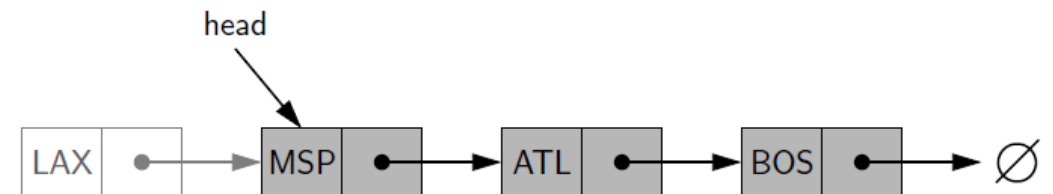


# SINGLY LINKED LISTS (单链表)

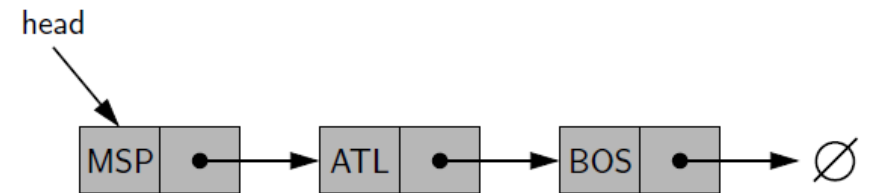
- Removing an element from the head of a singly linked list
  1. If L is empty, raise an error
  2.  $L.head = L.head.next$
  3.  $L.size = L.size - 1$
- What about removing the tail of a singly linked list?
  - Problems?
  - A traversal through the entire linkedlist is necessary with  $O(n)$
  - How do we solve this problem?



(a)



(b)



(c)

# STACK IMPLEMENTED WITH A SINGLY LINKED LIST

- Abstract Data Type (抽象数据类型 - ADT) of a Stack
  - `S.push(e)`: adds `e` to the top of stack `S`
  - `S.pop()`: removes and return the top element from stack `S`
    - Error when stack is empty
  - `S.top()`: returns a reference to the top element of stack `S` (but not to remove it)
  - `S.is_empty()`: returns `True` if `S` does not have any elements
  - `len(s)`: returns the number of elements in `S`
- Question: with a linkedlist, which side should be the top of the stack?

# STACK IMPLEMENTED WITH A SINGLY LINKED LIST

- Node class as the fundamental of the linkedlist
  - Element: to store the value of the node
  - Next: to store the reference to the next node
- The implementation of the Stack with linkedlist
  - push()
  - pop()
  - top()

# STACK IMPLEMENTED WITH A SINGLY LINKED LIST

- Node class as the fundamental of the linkedlist
  - Element: to store the value of the node
  - Next: to store the reference to the next node
- The implementation of the Stack with linkedlist
  - push()
  - pop() – raise error when stack is empty
  - top() – raise error when stack is empty



# STACK IMPLEMENTED WITH A SINGLY LINKED LIST

- Efficiency of Stack operations with LinkedList implementation?
  - S.push()
  - S.pop()
  - S.top()
  - Len(S)
  - S.is\_empty()
  - All  $O(1)$

# QUEUE IMPLEMENTED WITH A SINGLY LINKED LIST

- Abstract Data Type (ADT)
  - `Q.enqueue(e)`: add element `e` to the back of queue `Q` (入队)
  - `Q.dequeue()`: remove and return the first element from queue `Q` (出队)
  - `Q.first()`: returns a reference to the element at the front of queue `Q`
    - Error if `Q` is empty
  - `Q.is_empty()`: returns `True` if queue `Q` does not contain any elements
  - `len(Q)`: returns the number of elements in `Q`
- Question: which side of the queue should be the head of a linkedlist?
  - Front of the queue as the head of the linkedlist

# QUEUE IMPLEMENTED WITH A SINGLY LINKED LIST

```
1 class LinkedQueue:
2     """FIFO queue implementation using a singly linked list for storage."""
3
4     class _Node:
5         """Lightweight, nonpublic class for storing a singly linked node."""
6         (omitted here; identical to that of LinkedStack._Node)
7
8     def __init__(self):
9         """Create an empty queue."""
10        self._head = None
11        self._tail = None
12        self._size = 0                # number of queue elements
13
14    def __len__(self):
15        """Return the number of elements in the queue."""
16        return self._size
17
18    def is_empty(self):
19        """Return True if the queue is empty."""
20        return self._size == 0
21
22    def first(self):
23        """Return (but do not remove) the element at the front of the queue."""
24        if self.is_empty():
25            raise Empty('Queue is empty')
26        return self._head._element    # front aligned with head of list
```

```
27    def dequeue(self):
28        """Remove and return the first element of the queue (i.e., FIFO).
29
30        Raise Empty exception if the queue is empty.
31        """
32        if self.is_empty():
33            raise Empty('Queue is empty')
34        answer = self._head._element
35        self._head = self._head._next
36        self._size -= 1
37        if self.is_empty():           # special case as queue is empty
38            self._tail = None         # removed head had been the tail
39        return answer
40
41    def enqueue(self, e):
42        """Add an element to the back of queue."""
43        newest = self._Node(e, None)   # node will be new tail node
44        if self.is_empty():
45            self._head = newest        # special case: previously empty
46        else:
47            self._tail._next = newest
48            self._tail = newest        # update reference to tail node
49        self._size += 1
```

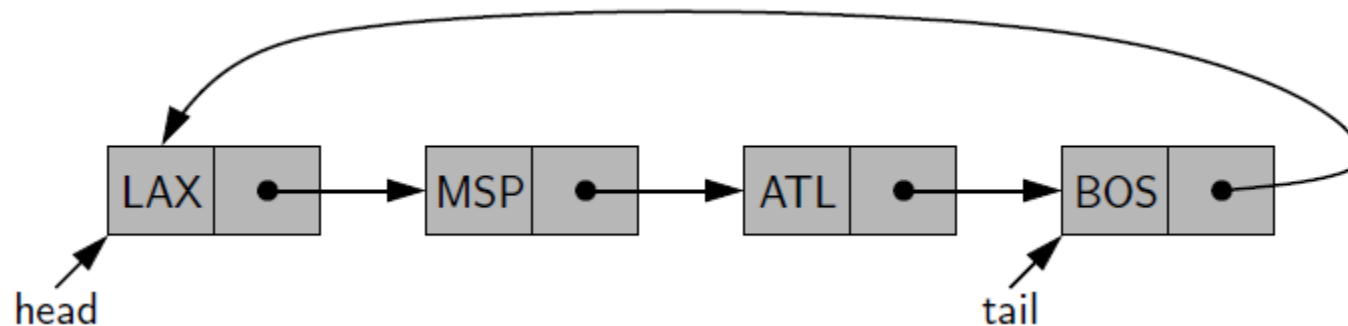
# QUEUE IMPLEMENTED WITH A SINGLY LINKED LIST

- Efficiency of Stack operations with LinkedList implementation?
  - Q.enqueue(e)
  - Q.dequeue()
  - Q.first()
  - Q.is\_empty()
  - len(Q)

# CIRCULARLY LINKEDLISTS

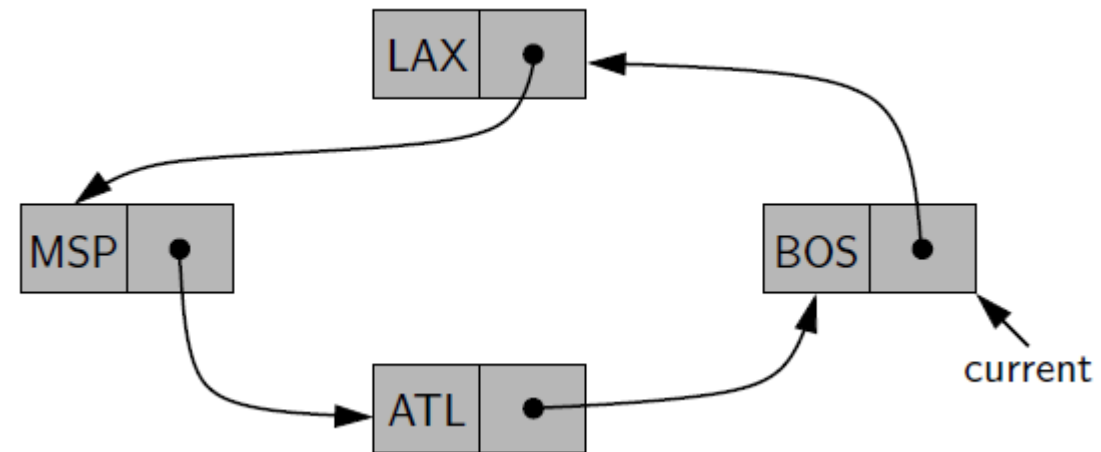
## (环形链表)

- Last week we talked about “circular” array
- We have seen that circular array is artificial
  - Need to maintain the reference to the front of the array
  - Need to compute the available slot of the array (tail)
  - Need to grow/shrink the array when necessary
- Circular linked list
  - Tail's *next* refers to head



# CIRCULARLY LINKEDLISTS (环形链表)

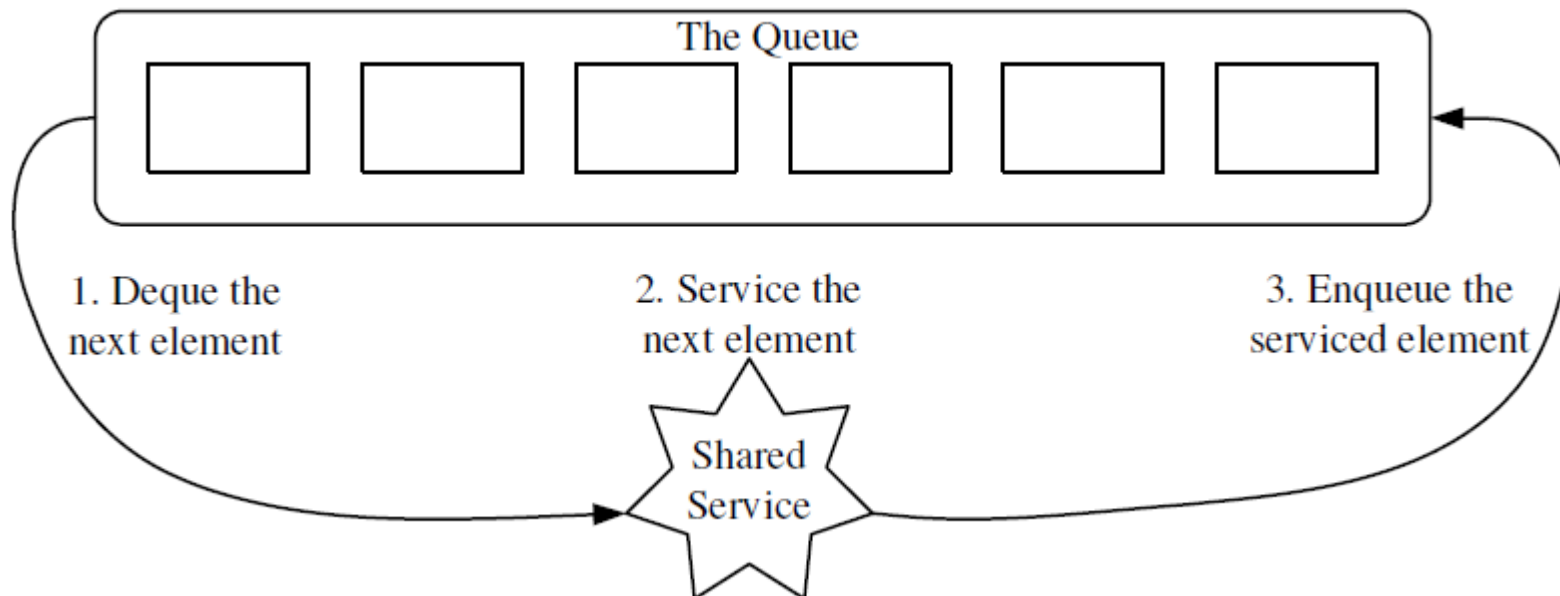
- Often used to represent cyclic data sets
  - No clear notion of a beginning and an end
  - Circular underground (地铁)
  - The order which players take turns during a game
- As data structure designers
  - Must maintain a reference to a particular node to make use of the list
  - We use **current**





# ROUND-ROBIN SCHEDULERS (轮询调度器)

- Example of a circularly linked list
- Iterates through a collection of elements in a circular fashion
- “serves” each element by performing a given action on it
- Used to allocate resource that must be shared by a collection of clients (processes)
  - E.g. CPU time
    1. `e = Q.dequeue()`
    2. Serve element `e`
    3. `Q.enqueue(e)`



# ROUND-ROBIN SCHEDULERS (轮询调度器)

- Implemented with a circularly linked list

```
1 class CircularQueue:
2     """ Queue implementation using circularly linked list for storage."""
3
4     class _Node:
5         """ Lightweight, nonpublic class for storing a singly linked node."""
6         (omitted here; identical to that of LinkedStack._Node)
7
8     def __init__(self):
9         """ Create an empty queue."""
10        self._tail = None           # will represent tail of queue
11        self._size = 0             # number of queue elements
12
13    def __len__(self):
14        """ Return the number of elements in the queue."""
15        return self._size
16
17    def is_empty(self):
18        """ Return True if the queue is empty."""
19        return self._size == 0
```

# ROUND-ROBIN SCHEDULERS (轮询调度器)

- Implemented with a circularly linked list

```
20 def first(self):
21     """Return (but do not remove) the element at the front of the queue.
22
23     Raise Empty exception if the queue is empty.
24     """
25     if self.is_empty():
26         raise Empty('Queue is empty')
27     head = self._tail._next
28     return head._element
29
30 def dequeue(self):
31     """Remove and return the first element of the queue (i.e., FIFO).
32
33     Raise Empty exception if the queue is empty.
34     """
35     if self.is_empty():
36         raise Empty('Queue is empty')
37     oldhead = self._tail._next
38     if self._size == 1:          # removing only element
39         self._tail = None       # queue becomes empty
40     else:
41         self._tail._next = oldhead._next  # bypass the old head
42     self._size -= 1
43     return oldhead._element
```

# ROUND-ROBIN SCHEDULERS (轮询调度器)

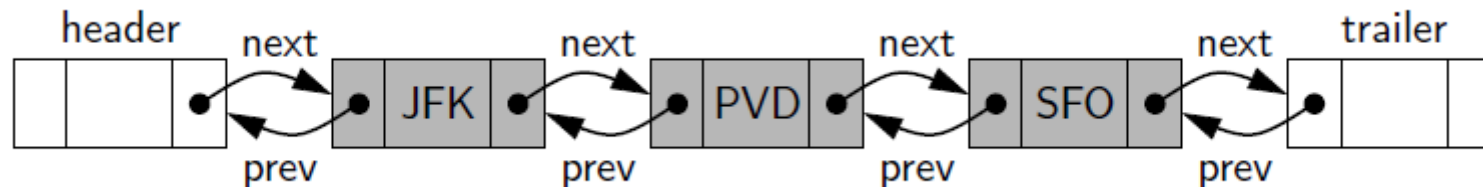
- Implemented with a circularly linked list

```
45  def enqueue(self, e):
46      """ Add an element to the back of queue."""
47      newest = self._Node(e, None)      # node will be new tail node
48      if self.is_empty():
49          newest._next = newest          # initialize circularly
50      else:
51          newest._next = self._tail._next # new node points to head
52          self._tail._next = newest      # old tail points to new node
53          self._tail = newest            # new node becomes the tail
54          self._size += 1
55
56  def rotate(self):
57      """ Rotate front element to the back of the queue."""
58      if self._size > 0:
59          self._tail = self._tail._next # old head becomes new tail
```

# DOUBLY LINKED LISTS

## (双链表)

- Limitations on Singly linked list
  - Delete the tail of the linked list is inefficient – always  $O(n)$
- Doubly linkedlist
  - Node – reference to both *previous* and *next* of it
  - Doubly linked list
- Structure: similar to singly linked list, with one exception
  - Head and trailer sentinels
  - A header node and a trailer node: empty 'dummy' nodes (哑元节点)
    - Also known as **sentinels** (哨兵节点)



# DOUBLY LINKED LISTS

## (双链表)

- Advantages of using sentinels
  - Header and trailer never change, only the nodes between them change
  - All insertions can be treated in a unified manner
    - A node is always placed between a pair of existing nodes
  - All deletions can also be treated in a unified manner
  - There is no special care for cases when the list is empty

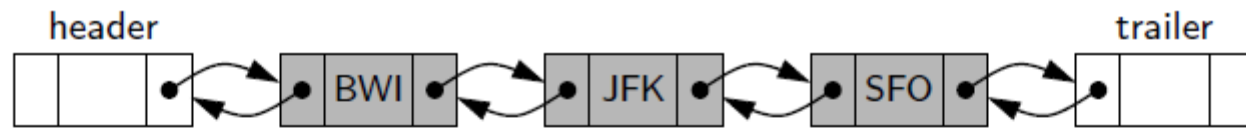
```
new_node = new Node(e, None)
if(self.is_empty()) {
    self_head = new_node
}
Else {
    self._tail._next = new_node
}
...
```



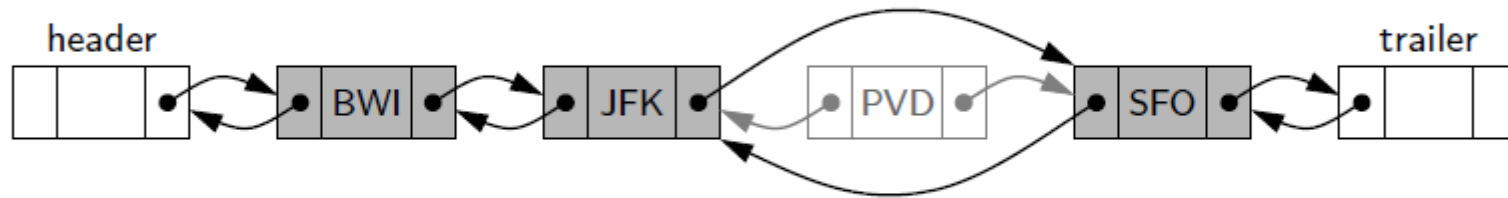
# DOUBLY LINKED LISTS

## (双链表)

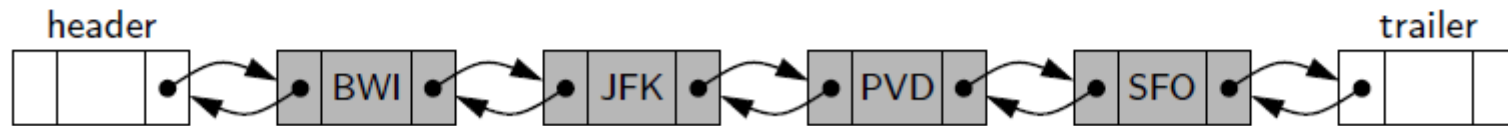
- Inserting into a doubly linked list



(a)



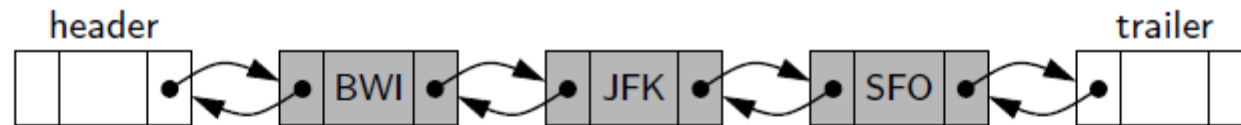
(b)



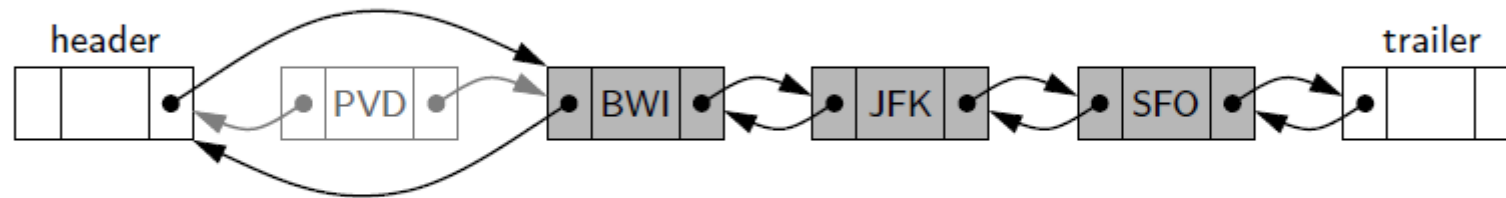
(c)

# DOUBLY LINKED LISTS (双链表)

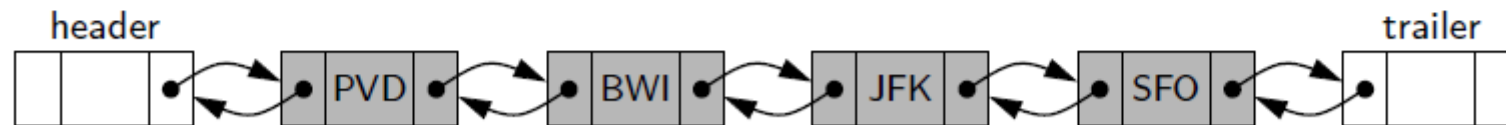
- Inserting into a doubly linked list



(a)



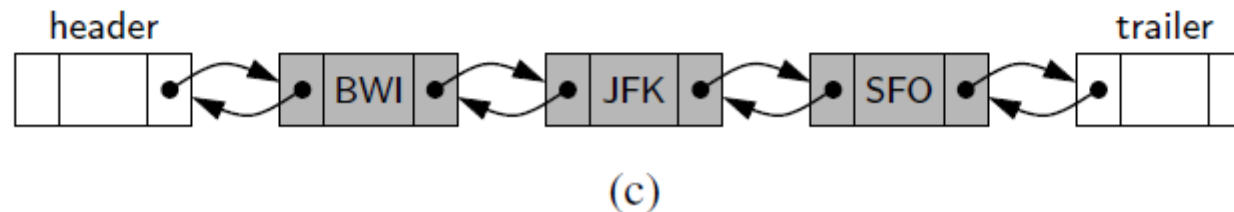
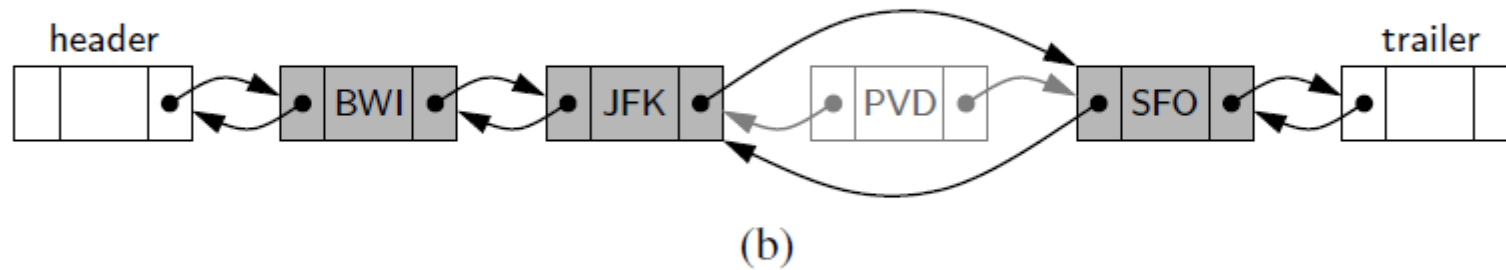
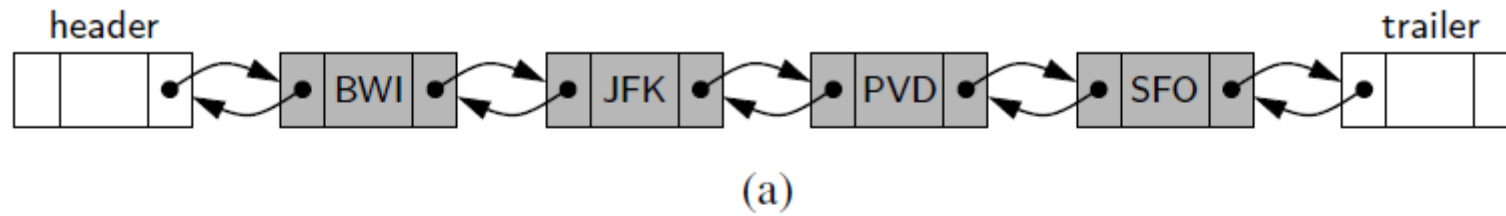
(b)



(c)

# DOUBLY LINKED LISTS (双链表)

- Inserting into a doubly linked list



# IMPLEMENTATION OF DOUBLY LINKED LISTS

- Inserting into a doubly linked list

```
class _Node:
    """Lightweight, nonpublic class for storing a doubly linked node."""
    __slots__ = '_element', '_prev', '_next' # streamline memory

    def __init__(self, element, prev, next): # initialize node's fields
        self._element = element           # user's element
        self._prev = prev                  # previous node reference
        self._next = next                  # next node reference
```

# IMPLEMENTATION OF DOUBLY LINKED LISTS

- Inserting into a doubly linked list

```
1 class _DoublyLinkedBase:
2     """ A base class providing a doubly linked list representation."""
3
4     class _Node:
5         """ Lightweight, nonpublic class for storing a doubly linked node."""
6         (omitted here; see previous code fragment)
7
8     def __init__(self):
9         """ Create an empty list."""
10        self._header = self._Node(None, None, None)
11        self._trailer = self._Node(None, None, None)
12        self._header._next = self._trailer # trailer is after header
13        self._trailer._prev = self._header # header is before trailer
14        self._size = 0 # number of elements
15
16    def __len__(self):
17        """ Return the number of elements in the list."""
18        return self._size
19
20    def is_empty(self):
21        """ Return True if list is empty."""
22        return self._size == 0
```

```
24    def _insert_between(self, e, predecessor, successor):
25        """ Add element e between two existing nodes and return new node."""
26        newest = self._Node(e, predecessor, successor) # linked to neighbors
27        predecessor._next = newest
28        successor._prev = newest
29        self._size += 1
30        return newest
31
32    def _delete_node(self, node):
33        """ Delete nonsentinel node from the list and return its element."""
34        predecessor = node._prev
35        successor = node._next
36        predecessor._next = successor
37        successor._prev = predecessor
38        self._size -= 1
39        element = node._element # record deleted element
40        node._prev = node._next = node._element = None # deprecate node
41        return element # return deleted element
```



# QUIZ FOR THIS WEEK

- Four people need to cross a rope bridge to get back to their camp
- They only have one flashlight, it only has enough power left to light for 17 minutes
- To cross the bridge, they must use a flashlight
- The bridge is only strong enough to support two people at any time
- Each of the campers walks at a different speed
  - 1 min
  - 2 mins
  - 5 mins
  - 10 mins
- How can the campers make it across in exactly 17 minutes?





# THANKS

See you in the next session!