# **CSC 3150 Assignment 2 Report**

120090472 Tianhao SHI

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### **Environment and Execution**

#### **Environment Info**

This assignment is written, compiled, tested in the following environment

```
vagrant@csc3150:~
                                                                           ₹921
                                             with vagrant@csc3150 at 03:23:37 ©
 Q ~ ~
  screenfetch
                               vagrant@csc3150
                               OS: Ubuntu 16.04 xenial
            ://+////-ууууууо
                               Kernel: x86_64 Linux 5.10.50
                               Uptime: 6d 21h 20m
       .:++0: /++++++/:-
                               Packages: 602
Shell: zsh 5.1.1
      0:+0+:++.
                  .-/00++++/
                               CPU: AMD Ryzen 7 4800H with Radeon Graphics @ 2.894GHz
                     `+sssoo+/
      .:+0:+0/.
                               RAM: 321MiB / 3933MiB
  .++/+:+00+0:
                      /sssooo.
 /+++//+:<u>`oo+o</u>
 \+/+0+++<sup>0</sup>++0
  .++.0+++00+:
      \+.++0+0``
        :o+++ `ohhhhhhhhyo++os:
         .o:`.syhhhhhhh/.oo++o
                     `00++.
  ++ --version | grep g++
g++ (Ubuntu 5.4.0-6ubuntu1~16.04.12) 5.4.0 20160609
  gcc --version | grep gcc
gcc (Ubuntu 5.4.0-6ubuntu1~16.04.12) 5.4.0 20160609
```

## Steps to execute

#### For the frog game:

```
$cd ~/Assignment_2_120090472/source #navigate to the code directory
$make # compile and run

# (optional)
#$ make build # compile only
#$./hw2 #run
```

#### For the thread pool

### Task 1

#### **Core Data Structures**

#### The frog struct

I modified the provided frog struct. Its members x and y are changed to row and col for consistency conserns.

#### The logs\_left\_arr

This global array of integers with size 9 stores the column information of each log's left edge. The right edge's location can be acquired using the log\_right macro.

#### The board

The game board is a 2d array with size ROW+1 x COL. This corresponds with 2 banks and 9 logs

### **Arrangements of threads**

This program will create two threads other than the main thread, the <code>logs\_thread</code> and the <code>frog\_thread</code>. The former thread controls the movements of logs. It will also let the frog follow the log's movement when the frog jumps on a particular log. The latter thread is used to listen to keyboard inputs that allows players to control the frog. The <code>frog\_thrad</code> thread is also responsible for updating the game status. Both threads will update the game board whenever the position of frog or log is updated.

Both threads will be quitted based on the value of global variable status. The value of status corresponds to the following:

- 1: running
- 0: game lost
- 2: game won
- 3: user quit

The main thread will first initialize the frog and log's position. Then it will first print the game board. It will then create the two threads and a mutex lock (to be discussed later in the report). The two threads are then joined, and main thread will be blocked until these two threads are exited. It will then analyze the game status and print the corresponding message.

A mutex is created to guaranteen thread safety. For the frog thread, it will acquire this mutex lock before examining the kbhit() and release it after completing drawing the map. For the logs\_thread, it will acquire a lock before updating the log position and after completing drawing of the map.

The following graph briefly explains the arrangement of threads in this program:

## Control of the frog

the frog is controlled by the keyboard. During execution, the frog\_thread will query if a key have been pressed by calling the provided kbhit() function. When this function returns 1, the pressed character will be read from the stdin using getchar().

## Movement of logs and frog

The logs' position will be update for a fixed time interval (defined by the Log\_SPEED). The elements in the logs\_left\_arr will all be increased/decreased by 1 every time they get updated. The actual increase or decrease is determined by the parit of the row index. The logs on the odd rows (that is, on the 1, 3, 5, 7, 9) will move right (increase column index), while the logs on the even rows will move left. Whenever the log's left bound hits index 0 or Row-1, it will be rounded to the other side of the board. The design of the log right macro also take the boundary into consideration.

When a frog is on one particular log, its position will be incremented or decremented the same as the log it rests on.

## Judging the game status

In the frog\_thread, after analyzing the input key and updating the frog's position, the update\_game\_status() function will be called. The game status is determined by: 1) Whether the user want to quit; 2) Whether the frog have reached the upper bank or falls into water.

#### **Game Quit**

When the user presses q or Q, the status is directly updated to 3, and the while loop in the frog\_move is breaked. The mutex lock is released, allowing logs\_thread to respond and quit.

#### **Game Won**

The game can only be won when the frog reaches anywhere on the upper bank. Therefore, by checking whether frog.row==0 we can know if the game is won or not.

#### **Game Lost**

The game can be lost in two ways:

- 1. The frog hits the left or right boundary of the game
- 2. The frog falls into water

If the frog is in the river (0< frog.row <ROW) and hits the boundary (frog.col ==0 || frog.col ==ROW-1), then it hits the boundary.

The second cindition can be checked by examining relative position between the frog and the log on the same row. If the log doesn't wraps around the board, then when the frog's column number is smaller than the left edge of the log, or when it's greater than the right edge of the log (frog.col < logs\_left[frog.col-1] | | frog.col > log\_right(logs\_left[frog.col-1])), then it will fall into river. Or, if the log wraps around the board, then the frog will only fall into river when its column number is both smaller than the left of log **AND** greater than the right of log. The following graph illustrates my point:

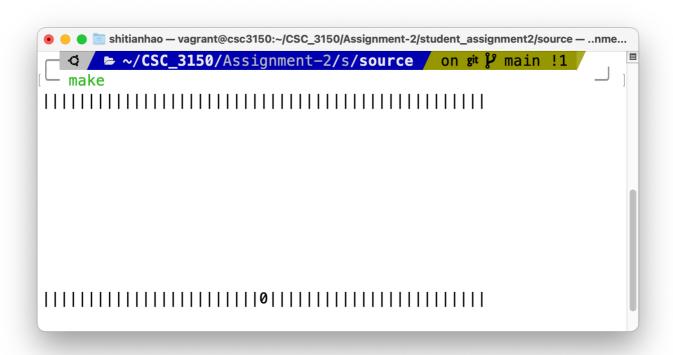
Whenever the game is won or lost, the <code>update\_game\_status()</code> will change the value of the global variable <code>status.lf</code> status!=1, the <code>frog thread</code> will release the mutex lock and break from the loop.

## Print out the message

This step is fairly straightforward. After the frog\_thread and the logs\_thread is quitted, the main thread will print out message based on the value of the status variable.

### **Program Output**

The static output of program (before any modification):



The output after starting the game

● ● □ CSC 3001 — ~/CSC_3150/	Assignment-2/student_assignme
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<u>                                      </u>	
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## Losing the game

```
CSC 3001 — vagrant@csc3150:~/CSC_3150/Assignment-2/...

You lose the game!

CSC_3150/Assignment-2/s/source on git ** de v **1 !1 ?4
```

## Winning the game

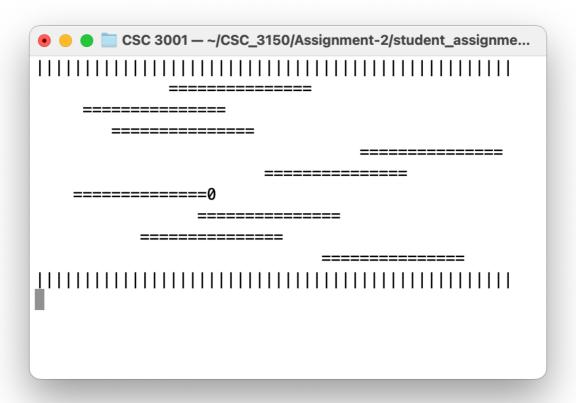
```
CSC 3001 — vagrant@csc3150:~/CSC_3150/Assignment-2/...

You win the game!

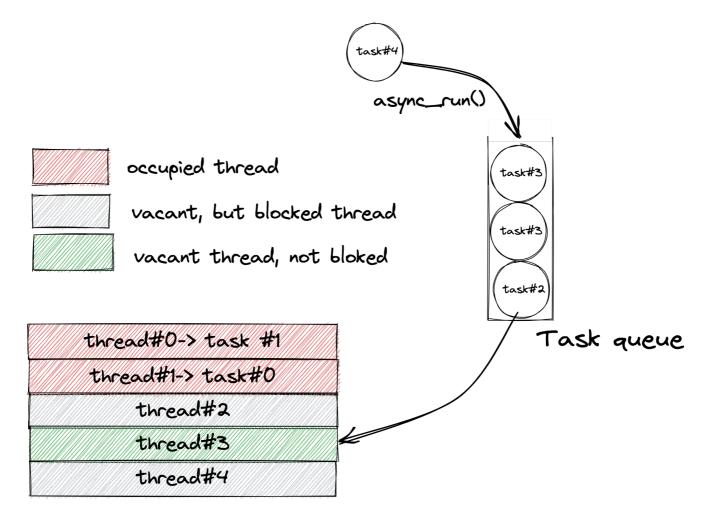
O > ~/CSC_3150/Assignment-2/s/source on git p de v +1 !1 ?4
```

**Quitting the game** 

## **During the game**



## The idea of a thread pool



The thread poo consists of two structures, the threads that run tasks concurantly, and a task queue that holds the tasks to be executed. Whenever user calls <code>async\_run</code>, a new task is added to the tail of the task queue. Whenever there is a vacant thread, it blocks other potentially vacant thread and pops the head task from the queue (given the queue is not empty).

#### **Core Data Structure**

#### The task item

Each task item represents a task to be executed by the thread pool. It contains the function to be executed by a thread and its corresponding arguments. It also contains a next and prev pointer, which is used when it's added to the task queue

#### The task queue

Different tasks jointly forms the task queue. This queue will be located in the heap memory. Its size grow whenever the <code>async\_run</code> function calls. The function adds a new node to the task queue. The thread pool pops item from the queue whenever there are vacant threads in the pool (and there are tasks to do in the queue)

#### The work container() function

When threads are created, a start routine must be provided. This function will be executed during the whole life span of the thread. However, this start routine is NOT the tasks that a thread needs to execute.

Therefore, I introduced a function called work container. This function does the following things:

- 1. Acquire a mutex lock (prevent other vacant threads from popping queue at the same time)
- 2. If the queue is empty, block until it recceives the signal that the queue is no longer empty (this prevents busy waiting, and will be discussed later)
- 3. Pop a task item from the queue
- 4. Release the lock
- 5. Execute the given task

## **Code design**

#### The async\_init() function

This function does the following things:

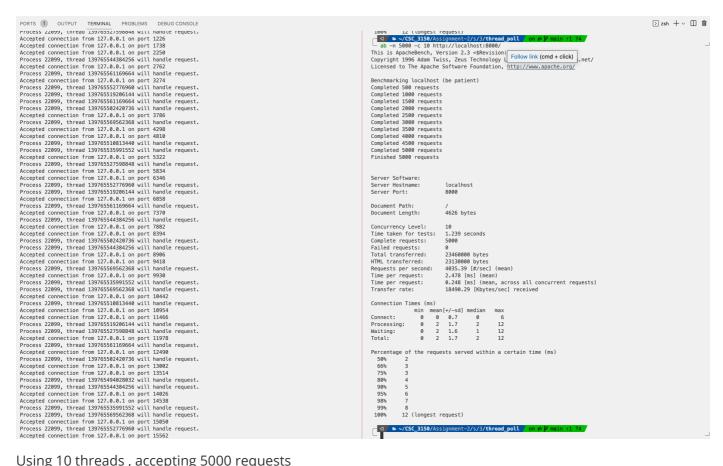
- 1. Initialize the task queue in the heap memory
- 2. Create the threads and detach them
  - These threads are detached because there is no point joining them, and thus they must free their resources when exited (although in this homework we didn't write methods to terminate a thread pool)

### The async\_run() function

This function is called whenever user/server receives requests and need to create a new task. Specifically speaking, this function does the following tasks:

- 1. Acquire a lock so that the queue is not accessed during the addition of task
- 2. Create a new task containing the function to be executed and its arguments
- 3. Append the new task to the end of gueue
- 4. Release the lock

## **Program Output**



Using 10 threads, accepting 5000 requests

## **Learning Outcome**

After this homework, I learned how to create POSIX threads. The creation of different threads opens the door of multithread programming. I also learned how to use mutex locks to wrire thread-safe programs, as well as use conditional variables for communication between different threads.