中间代码生成(2. 回填技术)

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$$S \rightarrow \mathbf{if} (B) S_1$$

```
\begin{array}{lll} B.true &=& newlabel() \\ B.false &=& S_1.next = S.next \\ S.code &=& B.code \mid\mid label(B.true) \mid\mid S_1.code \end{array}
```

B 还不知道 S.next 的指令地址, 如何跳转?

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B 还不知道 S.next 的指令地址, 如何跳转?

再扫描一遍中间代码,将标号替换成指令(相对)地址



$$S \rightarrow \mathbf{if} (B) S_1$$

$$B.true = newlabel()$$

 $B.false = S_1.next = S.next$
 $S.code = B.code || label(B.true) || S_1.code$

B 还不知道 S.next 的指令地址, 如何跳转?

再扫描一遍中间代码,将标号替换成指令(相对)地址

可否在生成中间代码的时候就填入指令地址?

回填 (Backpatching) 技术



子节点挖坑、祖先节点填坑

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子节点挖坑、祖先节点填坑

子节点暂时不指定跳转指令的目标 待祖先节点能够确定正确的目标地址时回头填充

回填 (Backpatching) 技术



子节点挖坑、祖先节点填坑

子节点暂时不指定跳转指令的目标 待祖先节点能够确定正确的目标地址时回头填充

父节点通过综合属性收集子节点中具有相同目标的跳转指令

在自底向上的分析过程中

为左部非终结符 B 计算 B.truelist 与 B.falselist 为左部非终结符 S 计算 S.nextlist

并为已能确定目标地址的跳转指令进行回填

针对布尔表达式的回填技术

```
1) B \rightarrow B_1 \parallel M B_2
                               { backpatch(B_1,falselist,M.instr);
                                  B.truelist = merge(B_1.truelist, B_2.truelist):
                                  B.falselist = B_2.falselist; 
     B \rightarrow B_1 \&\& M B_2
                                  backpatch(B_1.truelist.M.instr):
                                  B.truelist = B_{\uparrow}.truelist;
                                  B.falselist = merge(B_1.falselist, B_2.falselist); 
                               \{B.truelist = B_1.falselist;
                                  B.falselist = B_1.truelist; }
     B \rightarrow (B_{\perp})
                               { B.truelist = B_1.truelist;}
                                  B.falselist = B_1.falselist; 
    B \to E_1 \text{ rel } E_2 { B.truelist = makelist(nextinstr):
                                  B.falsclist = makelist(nextinstr + 1);
                                  gen('if' E1.addr rel.op E2.addr 'goto _'):
                                  gen('goto _'): }
     B \to \mathbf{true}
                               \{ B.truelist = makelist(nextinstr); \}
                                 gen('goto _'); }
     B \to \mathbf{false}
                               { B.falselist = makelist(nextinstr):
                                 gen('goto _'); }
                               \{ M.instr = nextinstr. \}
```

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综合属性 B.truelist 保存 需要跳转到 B.true 的指令地址

6) $B \rightarrow \mathbf{true}$ { $B.truelist = makelist(nextinstr); \\ <math>gen('goto \ ');$ }

7) $B \rightarrow \mathbf{false}$ { $B.falselist = makelist \ nextinstr); \\ <math>gen('goto \ ');$ }

综合属性 B.falselist 保存 需要跳转到 B.false 的指令地址

综合属性 B.truelist 保存 需要跳转到 B.true 的指令地址

6)
$$B \rightarrow \text{true}$$
 { $B.truelist = makelist(nextinstr); } gen('goto _'); }$

7)
$$B \rightarrow \text{false}$$
 { $B.\text{falselist} = makelist next instr); } gen('goto _'); }$

综合属性 B.falselist 保存 需要跳转到 B.false 的指令地址

$$B o ext{true}$$
 $B.code = gen('goto' B.true)$ $B o ext{false}$ $B.code = gen('goto' B.false)$

```
5) B \rightarrow E_1 \text{ rel } E_2
```

```
B.truelist = makelist(nextinstr):
B.falselist = makelist(nextinstr + 1);
gen('if' E<sub>1</sub>.addr rel.op E<sub>2</sub>.addr 'goto _'):
gen('goto _'): }
```

$$B \rightarrow E_1 \text{ rel } E_2$$
 | $B.code = E_1.code \parallel E_2.code$ | $||gen('if' E_1.addr \text{ rel.op } E_2.addr 'goto' B.true$ | $||gen('goto' B.false)$

$$3) \quad B \to 1 B_1$$

$$A$$
) $B \rightarrow (B_1)$

$$B \rightarrow \pm B_1$$

{
$$B.truelist = B_1.falselist;$$

 $B.falselist = B_1.truelist;$ }
{ $B.truelist = B_1.truelist;$
 $B.falselist = B_1.falselist;$ }

$$B_1.true = B.false$$

 $B_1.false = B.true$
 $B.code = B_1.code$

```
2) B \rightarrow B_1 \&\& M B_2 { backpatch(B_1.truelist, M.instr); B.truelist = B_2.truelist; B.falselist = merge(B_1.falselist, B_2.falselist); }
```

8)
$$M \to \epsilon$$
 { $M.instr = nextinstr$, }

$$B \rightarrow B_1 \&\& B_2$$
 $| B_1.true = newlabel() | B_1.false = B.false | B_2.true = B.true | B_2.false | B.false | B.false | B.code = B_1.code || label(B_1.true) || B_2.code || B_2.code || B_2.true || B_2.code || B_2.true || B_2.code || B_2.true || B_2$

1)
$$B \rightarrow B_1 \parallel M B_2 = \{ \begin{array}{ll} backpatch(B_1, falselist, M.instr); \\ B.truelist = merge(B_1, truelist, B_2, truelist); \\ B.falselist = B_2, falselist; \} \end{array}$$

8)
$$M \to \epsilon$$
 { $M.instr = nextinstr$, }

$$B \rightarrow B_1 \mid \mid B_2$$

$$B_1.true = B.true$$

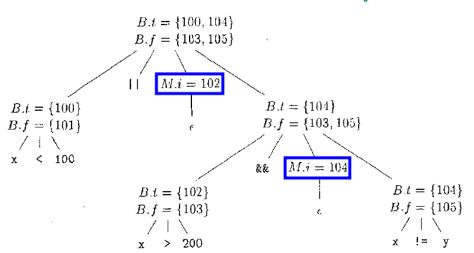
$$B_1.false = newlabel()$$

$$B_2.true = B.true$$

$$B_2.false = B.false$$

$$B.code = B_1.code \mid \mid label(B_1.false) \mid \mid B_2.code$$

$x < 100 \mid \mid x > 200 \&\& x != y$



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```
100: if x < 100 goto _

101: goto _

102: if x > 200 goto 104

103: goto _

104: if x != y goto _

105: goto _
```

a) 将 104 回填到指令 102 中之后

```
100: if x < 100 goto _

101: goto 102

102: if x > 200 goto 104

103: goto _

104: if x != y goto _

105: goto _
```

b) 将 102 回填到指令 101 中之后

$$S \rightarrow \text{if } (B) S \mid \text{if } (B) S \text{ else } S \mid \text{ while } (B) S \mid \{L\} \mid A;$$

 $L \rightarrow L S \mid S$

```
 S → if (B) M S<sub>1</sub> { backpatch(B.truelist, M.instr);

                           S.nextlist = merge(B.falselist, S_1.nextlist);
2) S \rightarrow \text{if } (B) M_1 S_1 N \text{ else } M_2 S_2
                         { backpatch(B.truelist, M<sub>1</sub>.instr);
                            backpatch(B.falselist, M_2.instr);
                            temp = merge(S_1.nextlist, N.nextlist);
                           S.nextlist = merge(temp, S_2.nextlist);
3) S \rightarrow \text{ while } M_1 (B) M_2 S_1
                         { backpatch(S<sub>1</sub>.nextlist, M<sub>1</sub>.instr);
                            backpatch(B.truelist, M_2.instr);
                           S.nextlist = B.falselist;
                           gen ('goto' Mi.instr); }
                     \{S.nextlist = L.nextlist;\}
                     { S.nextlist = null; }
                         \{M.instr = nextinstr, \}
                          \{ N.nextlist = makelist(nextinstr); \}
                           gen('goto _'); }
                      { backpatch(L<sub>1</sub>.nextlist, M.instr);
                        L.nextlist = S.nextlist; }

 L → S

                         \{L.nextlist = S.nextlist;\}
```

1)
$$S \to if$$
 (B) MS_1 { $backpatch$ [B.truelist, M.instr); $S.nextlist$ = $merge(B.falselist, S_1.nextlist)$; }

6)
$$M \to \epsilon$$
 { $M.instr = nextinstr$, }

1)
$$S \to if(B) M S_1 \{ backpatch B.truelist, M.instr);$$

 $S.nextlist = merge(B.falselist, S_1.nextlist); \}$

6)
$$M \to \epsilon$$
 { $M.instr = nextinstr$, }

$$S \rightarrow \mathbf{if} (B) S_1$$

$$B.true = \underbrace{newlabel()}_{B.false} = \underbrace{S_1.next}_{S.code} = S.next$$

$$S.code = B.code || label(B.true) || S_1.code$$

```
S \rightarrow \mathbf{if}(B) M_1 S_1 N \mathbf{else} M_2 S_2
\{ \begin{array}{c} backpatch B.truelist, M_1.instr); \\ backpatch B.falselist, M_2.instr); \\ temp = merge(S_1.nextlist, N.nextlist); \\ S.nextlist = merge(temp, S_2.nextlist); \} \end{array}
```

```
6) M \to \epsilon { M.instr = nextinstr; }

7) N \to \epsilon { N.nextlist = makelist(nextinstr); gen('goto _'); }
```

```
S \rightarrow \mathbf{if}(B) M_1 S_1 N \text{ else } M_2 S_2
                                     { backpatch B.truelist, M_1.instr); backpatch B.falselist, M_2.instr);}
                                         \underline{temp} = \underline{merge}(S_1.nextlist, N.nextlist);
                                        S.nextlist = merge(temp, S_2.nextlist); 
            6) M \to \epsilon
                                                  \{ M.instr = nextinstr, \}
            7) N \to \epsilon
                                                  \{ N.nextlist = makelist(nextinstr); \}
                                                     gen('goto _'); }
                 S \rightarrow \text{if } (B) S_1 \text{ else } S_2
B.true = newlabel()
B.false = newlabel()
S_1.next = S_2.next = S.next
S.code = B.code
|| label(B.true) || S_1.code
|| gen('goto' S.next) || S_1.code
                                                                      || label(B.false) || S_2.code
```

6)
$$M \rightarrow \epsilon$$
 { $M.instr = nextinstr$, }

```
3) S \rightarrow \text{ while } M_1 (B) M_2 S_1
                                                            \{\begin{array}{ll} backpatch & S_1.nextlist, & M_1.instr); \\ backpatch & B.truelist, & M_2.instr); \\ S.nextlist & = & B.falselist; \\ gen('goto' & M_1.instr); & \} \end{array}
                                                                        \{ M.instr = nextinstr, \}
        6) M \rightarrow \epsilon
```

$$S \rightarrow \text{while } (B) S_1$$

$$\begin{array}{c} begin = newlabel() \\ B.true = newlabel() \\ B.false = S.next \\ \hline S_1.next = begin \\ S.code = label(begin) \mid\mid B.code \\ \mid\mid label(B.true) \mid\mid S_1.code \\ \mid\mid gen('goto' begin) \end{array}$$

4)
$$S \rightarrow \{L\}$$

 $\{ S.nextlist = L.nextlist; \}$

5)
$$S \rightarrow A$$
;

 $\{S.nextlist = null; \}$

6)
$$M \to \epsilon$$

 $\{ M.instr = nextinstr, \}$

8)
$$L \rightarrow L_1 M S$$

{ $backpatch(L_1.nextlist, M.instr);$ L.nextlist = S.nextlist; }

9)
$$L \rightarrow S$$

 $\{L.nextlist = S.nextlist;\}$

```
 S → if (B) M S<sub>1</sub> { backpatch(B.truelist, M.instr);

                             S.nextlist = merge(B.falselist, S_1.nextlist);
2) S → if (B) M<sub>1</sub> S<sub>1</sub> N else M<sub>2</sub> S<sub>2</sub>
                           { backpatch(B.truelist, M<sub>1</sub>.instr);
                             backpatch(B.falselist, M_2.instr);
                             temp = merge(S_1.nextlist, N.nextlist);
                             S.nextlist = merge(temp, S_2.nextlist);
3) S \rightarrow while M_1 (B) M_2 S_1
                           { backpatch(S<sub>1</sub>.nextlist, M<sub>1</sub>.instr);
                             backpatch(B.truelist, M_2.instr);
                            S.nextlist = B.falselist;
                            qen('goto' M<sub>1</sub>.instr);
4) S \rightarrow \{L\}
                          \{S.nextlist = L.nextlist;\}
5) S → A :
                          \{ S.nextlist = null; \}

 M → ϵ

                           \{M.instr = nextinstr, \}
7) N \rightarrow \epsilon
                            \{ N.nextlist = makelist(nextinstr); \}
                            gen('goto _'); }
8) L → L<sub>1</sub> M S
                           { backpatch(L<sub>1</sub>.nextlist, M.instr);
                             L.nextlist = S.nextlist;

 L → S

                           \{L.nextlist = S.nextlist;\}
```

只有(3)与(7)生成了新的代码,控制流语句的主要目的是"控制"流。

Thank You!



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