

中间代码生成

(2. 回填技术)

魏恒峰

hfwei@nju.edu.cn

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$S \rightarrow \text{if} (B) S_1$

$\left\{ \begin{array}{l} B.true = \text{newlabel}() \\ B.false = S_1.next = S.next \\ S.code = B.code \parallel \text{label}(B.true) \parallel S_1.code \end{array} \right.$

B 还不知道 *S.next* 的指令地址, 如何跳转?

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再扫描一遍中间代码, 将标号替换成指令 (相对) 地址

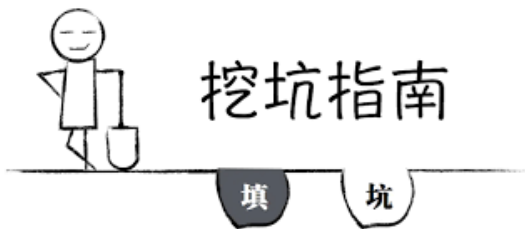
$$S \rightarrow \text{if} (B) S_1 \quad \left| \begin{array}{l} B.\text{true} = \text{newlabel}() \\ B.\text{false} = S_1.\text{next} = S.\text{next} \\ S.\text{code} = B.\text{code} || \text{label}(B.\text{true}) || S_1.\text{code} \end{array} \right.$$

B 还不知道 *S.next* 的指令地址, 如何跳转?

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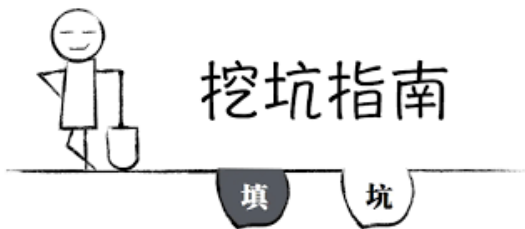
可否在生成中间代码的时候就填入指令地址?

回填 (Backpatching) 技术



子节点挖坑、祖先节点填坑

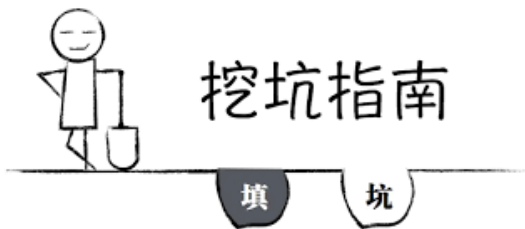
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子节点挖坑、祖先节点填坑

子节点暂时不指定跳转指令的目标
待祖先节点能够确定正确的目标地址时回头填充

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子节点挖坑、祖先节点填坑

子节点暂时不指定跳转指令的目标

待祖先节点能够确定正确的目标地址时回头填充

父节点通过**综合属性**收集子节点中具有相同目标的跳转指令

在自底向上的分析过程中

为左部非终结符 B 计算 $B.truelist$ 与 $B.falselist$

为左部非终结符 S 计算 $S.nextlist$

并为已能确定目标地址的跳转指令进行回填

针对布尔表达式的回填技术

- | | | |
|----|--------------------------------------|---|
| 1) | $B \rightarrow B_1 \parallel M B_2$ | { <i>backpatch</i> (<i>B</i> ₁ . <i>false</i> list, <i>M.instr</i>);
<i>B.true</i> list = <i>merge</i> (<i>B</i> ₁ . <i>true</i> list, <i>B</i> ₂ . <i>true</i> list);
<i>B.false</i> list = <i>B</i> ₂ . <i>false</i> list; } |
| 2) | $B \rightarrow B_1 \&\& M B_2$ | { <i>backpatch</i> (<i>B</i> ₁ . <i>true</i> list, <i>M.instr</i>);
<i>B.true</i> list = <i>B</i> ₂ . <i>true</i> list;
<i>B.false</i> list = <i>merge</i> (<i>B</i> ₁ . <i>false</i> list, <i>B</i> ₂ . <i>false</i> list); } |
| 3) | $B \rightarrow ! B_1$ | { <i>B.true</i> list = <i>B</i> ₁ . <i>false</i> list;
<i>B.false</i> list = <i>B</i> ₁ . <i>true</i> list; } |
| 4) | $B \rightarrow (B_1)$ | { <i>B.true</i> list = <i>B</i> ₁ . <i>true</i> list;
<i>B.false</i> list = <i>B</i> ₁ . <i>false</i> list; } |
| 5) | $B \rightarrow E_1 \text{ rel } E_2$ | { <i>B.true</i> list = <i>makelist</i> (<i>nextinstr</i>);
<i>B.false</i> list = <i>makelist</i> (<i>nextinstr</i> + 1);
<i>gen</i> ('if' <i>E</i> ₁ . <i>addr</i> <i>rel.op</i> <i>E</i> ₂ . <i>addr</i> 'goto -');
<i>gen</i> ('goto -'); } |
| 6) | $B \rightarrow \text{true}$ | { <i>B.true</i> list = <i>makelist</i> (<i>nextinstr</i>);
<i>gen</i> ('goto -'); } |
| 7) | $B \rightarrow \text{false}$ | { <i>B.false</i> list = <i>makelist</i> (<i>nextinstr</i>);
<i>gen</i> ('goto -'); } |
| 8) | $M \rightarrow \epsilon$ | { <i>M.instr</i> = <i>nextinstr</i> ; } |

综合属性 $B.truelist$ 保存 需要跳转到 $B.true$ 的指令地址

- 6) $B \rightarrow \text{true}$ { $B.truelist = makelist(nextinstr);$
 $gen('goto _');$ }
- 7) $B \rightarrow \text{false}$ { $B.falselist = makelist(nextinstr);$
 $gen('goto _');$ }

综合属性 $B.falselist$ 保存 需要跳转到 $B.false$ 的指令地址

综合属性 $B.truelist$ 保存 需要跳转到 $B.true$ 的指令地址

- 6) $B \rightarrow true$ { $B.truelist = makelist(nextinstr);$
 $gen('goto _');$ }
- 7) $B \rightarrow false$ { $B.falselist = makelist(nextinstr);$
 $gen('goto _');$ }

综合属性 $B.falselist$ 保存 需要跳转到 $B.false$ 的指令地址

$B \rightarrow true$	$B.code = gen('goto' B.true$
$B \rightarrow false$	$B.code = gen('goto' B.false)$

5) $B \rightarrow E_1 \text{ rel } E_2$ { $B.truelist = makelist(nextinstr);$
 $B.falselist = makelist(nextinstr + 1);$
 $gen('if' E_1.addr \text{ rel.op } E_2.addr 'goto -');$
 $gen('goto -');$ }

$B \rightarrow E_1 \text{ rel } E_2$ { $B.code = E_1.code || E_2.code$
 || $gen('if' E_1.addr \text{ rel.op } E_2.addr 'goto' B.true$
 || $gen('goto' B.false)$ }

$$3) \quad B \rightarrow ! B_1$$

$$\{ \boxed{B.truelist} = B_1.falselist; \\ \boxed{B.falselist} = B_1.truelist; \}$$

$$4) \quad B \rightarrow (B_1)$$

$$\{ \boxed{B.truelist} = B_1.truelist; \\ \boxed{B.falselist} = B_1.falselist; \}$$

$$B \rightarrow ! B_1$$

$$\left| \begin{array}{l} B_1.true = B.false \\ B_1.false = B.true \\ B.code = B_1.code \end{array} \right.$$

2) $B \rightarrow B_1 \ \&\& \ M \ B_2 \quad \{ \text{backpatch}(B_1.\text{truelist}, M.\text{instr});$
 $B.\text{truelist} = B_2.\text{truelist};$
 $B.\text{falselist} = \text{merge}(B_1.\text{falselist}, B_2.\text{falselist}); \}$

8) $M \rightarrow \epsilon \quad \{ M.\text{instr} = \text{nextinstr}; \}$

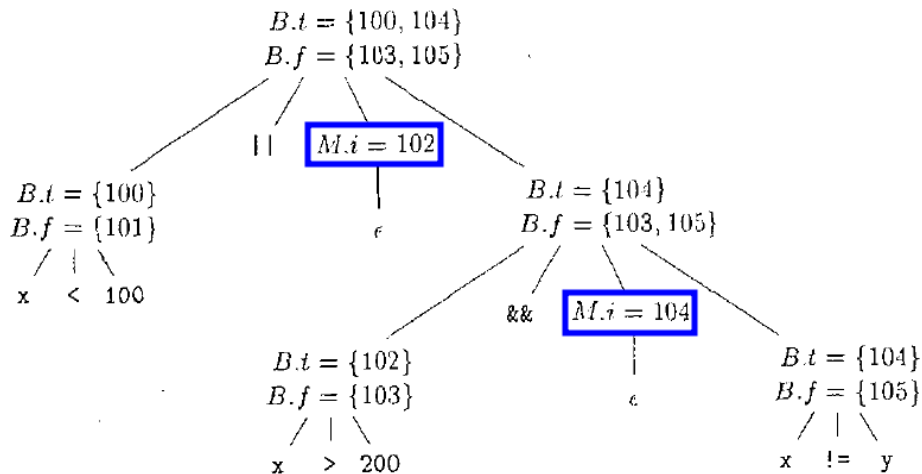
$B \rightarrow B_1 \ \&\& \ B_2 \quad \begin{cases} B_1.\text{true} = \text{newlabel}() \\ B_1.\text{false} = B.\text{false} \\ B_2.\text{true} = B.\text{true} \\ B_2.\text{false} = B.\text{false} \\ B.\text{code} = B_1.\text{code} \ || \ \text{label}(B_1.\text{true}) \ || \ B_2.\text{code} \end{cases}$

1) $B \rightarrow B_1 \parallel M B_2$ { $\text{backpatch}(B_1.\text{falselist}, M.\text{instr});$
 $B.\text{truelist} = \text{merge}(B_1.\text{truelist}, B_2.\text{truelist});$
 $B.\text{falselist} = B_2.\text{falselist};$ }

8) $M \rightarrow \epsilon$ { $M.\text{instr} = \text{nextinstr};$ }

$B \rightarrow B_1 \parallel B_2$ { $B_1.\text{true} = B.\text{true}$
 $B_1.\text{false} = \text{newlabel}()$
 $B_2.\text{true} = B.\text{true}$
 $B_2.\text{false} = B.\text{false}$
 $B.\text{code} = B_1.\text{code} \parallel \text{label}(B_1.\text{false}) \parallel B_2.\text{code}$ }

$x < 100 \ || \ x > 200 \ \&\& \ x \neq y$




```
100:  if x < 100 goto -  
101:  goto -  
102:  if x > 200 goto 104  
103:  goto -  
104:  if x != y goto -  
105:  goto -
```

a) 将 104 回填到指令 102 中之后

```
100:  if x < 100 goto -  
101:  goto 102  
102:  if x > 200 goto 104  
103:  goto -  
104:  if x != y goto -  
105:  goto -
```

b) 将 102 回填到指令 101 中之后

$$\begin{aligned} S &\rightarrow \text{if}(B) S \mid \text{if}(B) S \text{ else } S \mid \text{while}(B) S \mid \boxed{\{L\}} \mid A ; \\ L &\rightarrow L S \mid S \end{aligned}$$

- 1) $S \rightarrow \text{if}(B) M S_1 \{ \text{backpatch}(B.\text{truelist}, M.\text{instr});$
 $S.\text{nextlist} = \text{merge}(B.\text{falselist}, S_1.\text{nextlist}); \}$
- 2) $S \rightarrow \text{if}(B) M_1 S_1 N \text{ else } M_2 S_2$
 $\{ \text{backpatch}(B.\text{truelist}, M_1.\text{instr});$
 $\text{backpatch}(B.\text{falselist}, M_2.\text{instr});$
 $\text{temp} = \text{merge}(S_1.\text{nextlist}, N.\text{nextlist});$
 $S.\text{nextlist} = \text{merge}(\text{temp}, S_2.\text{nextlist}); \}$
- 3) $S \rightarrow \text{while } M_1 (B) M_2 S_1$
 $\{ \text{backpatch}(S_1.\text{nextlist}, M_1.\text{instr});$
 $\text{backpatch}(B.\text{truelist}, M_2.\text{instr});$
 $S.\text{nextlist} = B.\text{falselist};$
 $\text{gen}(\text{'goto' } M_1.\text{instr}); \}$
- 4) $S \rightarrow \{ L \} \quad \{ S.\text{nextlist} = L.\text{nextlist}; \}$
- 5) $S \rightarrow A ; \quad \{ S.\text{nextlist} = \text{null}; \}$
- 6) $M \rightarrow \epsilon \quad \{ M.\text{instr} = \text{nextinstr}; \}$
- 7) $N \rightarrow \epsilon \quad \{ N.\text{nextlist} = \text{makelist}(\text{nextinstr});$
 $\text{gen}(\text{'goto' } -); \}$
- 8) $L \rightarrow L_1 M S \quad \{ \text{backpatch}(L_1.\text{nextlist}, M.\text{instr});$
 $L.\text{nextlist} = S.\text{nextlist}; \}$
- 9) $L \rightarrow S \quad \{ L.\text{nextlist} = S.\text{nextlist}; \}$

1) $S \rightarrow \text{if}(B) M S_1 \{ \text{backpatch}(B.\text{truelist}, M.\text{instr});$
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$S \rightarrow \text{if}(B) S_1 \quad \left\{ \begin{array}{l} B.\text{true} = \text{newlabel}() \\ B.\text{false} = S_1.\text{next} = S.\text{next} \\ S.\text{code} = B.\text{code} \parallel \text{label}(B.\text{true}) \parallel S_1.\text{code} \end{array} \right.$

$$S \rightarrow \text{if}(B) M_1 S_1 N \text{ else } M_2 S_2$$

```

{
    backpatch(B.truelist, M1.instr);
    backpatch(B.falselist, M2.instr);
    temp = merge(S1.nextlist, N.nextlist);
    S.nextlist = merge(temp, S2.nextlist);
}

```

6) $M \rightarrow \epsilon$ { $M.instr = nextinstr$; }

7) $N \rightarrow \epsilon$ { $N.nextlist = makelist(nextinstr$;
 $gen('goto -')$); }

$$S \rightarrow \text{if}(B) M_1 S_1 N \text{ else } M_2 S_2$$

```

{
    backpatch(B.truelist, M1.instr);
    backpatch(B.falselist, M2.instr);
    temp = merge(S1.nextlist, N.nextlist);
    S.nextlist = merge(temp, S2.nextlist);
}

```

6) $M \rightarrow \epsilon$ $\{ M.instr = nextinstr; \}$

7) $N \rightarrow \epsilon$ $\{ N.nextlist = makelist(nextinstr);$
 $gen('goto -'); \}$

$S \rightarrow \text{if}(B) S_1 \text{ else } S_2$	$B.true = newlabel()$ $B.false = newlabel()$ $S_1.next = S_2.next = S.next$ $S.code = B.code$ $ \text{label}(B.true) S_1.code$ $ gen('goto' S.next)$ $ \text{label}(B.false) S_2.code$
--	--

3) $S \rightarrow \text{while } M_1 (B) M_2 S_1$

```

{ backpatch( $S_1.nextlist$ ,  $M_1.instr$ );
  backpatch( $B.truelist$ ,  $M_2.instr$ );
   $S.nextlist = B.falselist$ ;
  gen('goto'  $M_1.instr$ ); }

```

6) $M \rightarrow \epsilon$ $\{ M.instr = nextinstr; \}$

3) $S \rightarrow \text{while } M_1 (B) M_2 S_1$

```

{
  backpatch( $S_1.nextlist$ ,  $M_1.instr$ );
  backpatch( $B.truelist$ ,  $M_2.instr$ );
   $S.nextlist = B.falselist$ ;
  gen('goto'  $M_1.instr$ );
}
```

6) $M \rightarrow \epsilon$ $\{ M.instr = nextinstr; \}$

$S \rightarrow \text{while } (B) S_1$	<pre> begin = newlabel() B.true = newlabel() B.false = S.next $S_1.next = begin$ S.code = label(begin) B.code label(B.true) $S_1.code$ gen('goto' begin)</pre>
---	--

4) $S \rightarrow \{ L \}$ $\{ S.nextlist = L.nextlist; \}$

5) $S \rightarrow A ;$ $\{ S.nextlist = \text{null}; \}$

6) $M \rightarrow \epsilon$ $\{ M.instr = nextinstr; \}$

8) $L \rightarrow L_1 M S$ $\{ \text{backpatch}(L_1.nextlist, M.instr);$
 $L.nextlist = S.nextlist; \}$

9) $L \rightarrow S$ $\{ L.nextlist = S.nextlist; \}$

- 1) $S \rightarrow \text{if}(B) M S_1 \{ \text{backpatch}(B.\text{truelist}, M.\text{instr});$
 $S.\text{nextlist} = \text{merge}(B.\text{falselist}, S_1.\text{nextlist}); \}$
- 2) $S \rightarrow \text{if}(B) M_1 S_1 N \text{ else } M_2 S_2$
 $\{ \text{backpatch}(B.\text{truelist}, M_1.\text{instr});$
 $\text{backpatch}(B.\text{falselist}, M_2.\text{instr});$
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 $S.\text{nextlist} = \text{merge}(\text{temp}, S_2.\text{nextlist}); \}$
- 3) $S \rightarrow \text{while } M_1 (B) M_2 S_1$
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 $S.\text{nextlist} = B.\text{falselist};$
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- 4) $S \rightarrow \{ L \} \quad \{ S.\text{nextlist} = L.\text{nextlist}; \}$
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 $L.\text{nextlist} = S.\text{nextlist}; \}$
- 9) $L \rightarrow S \quad \{ L.\text{nextlist} = S.\text{nextlist}; \}$

只有 (3) 与 (7) 生成了新的代码, 控制流语句的主要目的是“控制”流。

```
switch (  $E$  ) {  
    case  $V_1$ :  $S_1$   
    case  $V_2$ :  $S_2$   
    ...  
    case  $V_{n-1}$ :  $S_{n-1}$   
    default:  $S_n$   
}
```

非 C 语言语义 (break)

```

switch (  $E$  ) {
    case  $V_1$ :  $S_1$ 
    case  $V_2$ :  $S_2$ 
        ...
    case  $V_{n-1}$ :  $S_{n-1}$ 
    default:  $S_n$ 
}

```

非 C 语言语义 (break)

```

code to evaluate  $E$  into  $t$ 
goto test
L1:    code for  $S_1$ 
        goto next
L2:    code for  $S_2$ 
        goto next
...
Ln-1:  code for  $S_{n-1}$ 
        goto next
Ln:    code for  $S_n$ 
        goto next
test:   if  $t = V_1$  goto L1
        if  $t = V_2$  goto L2
        ...
        if  $t = V_{n-1}$  goto Ln-1
        goto Ln
next:

```

```

switch (  $E$  ) {
    case  $V_1$ :  $S_1$ 
    case  $V_2$ :  $S_2$ 
        ...
    case  $V_{n-1}$ :  $S_{n-1}$ 
    default:  $S_n$ 
}

```

$V_i : L_i$ 队列

```

code to evaluate  $E$  into  $t$ 
goto test
L1:   code for  $S_1$ 
      goto next
L2:   code for  $S_2$ 
      goto next
      ...
L $n-1$ : code for  $S_{n-1}$ 
      goto next
L $n$ :  code for  $S_n$ 
      goto next
test:  if  $t = V_1$  goto L1
      if  $t = V_2$  goto L2
      ...
      if  $t = V_{n-1}$  goto L $n-1$ 
      goto L $n$ 
next:

```

```
code to evaluate  $E$  into  $t$   
goto test
```

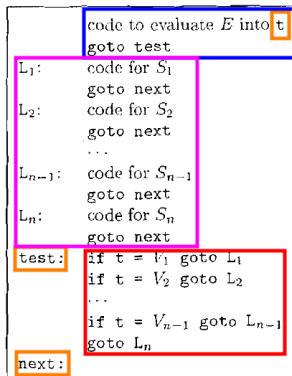
```
L1: code for  $S_1$   
goto next  
L2: code for  $S_2$   
goto next  
...  
L $n-1$ : code for  $S_{n-1}$   
goto next  
L $n$ : code for  $S_n$   
goto next
```

```
test: if  $t = V_1$  goto L1  
if  $t = V_2$  goto L2  
...  
if  $t = V_{n-1}$  goto L $n-1$   
goto L $n$ 
```

```
next:
```

```
case  $t$   $V_1$  L1  
case  $t$   $V_2$  L2  
...  
case  $t$   $V_{n-1}$  L $n-1$   
case  $t$   $t$  L $n$   
next:
```

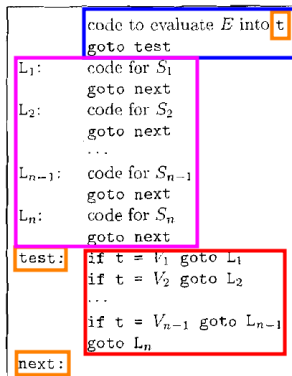
case 三地址代码



```

case t  $V_1$   $L_1$ 
case t  $V_2$   $L_2$ 
...
case t  $V_{n-1}$   $L_{n-1}$ 
case t t  $L_n$ 
next:

```

```

case t V1 L1
case t V2 L2
...

case t Vn-1 Ln-1
case t t Ln
next:

```

Jump Table Structure

C code:

```

switch(x) {
  case 1: <some code>
          break;
  case 2: <some code>
          break;
  case 3: <some code>
          break;
  case 5: <some code>
          break;
  case 6: <some code>
          break;
  default: <some code>
}

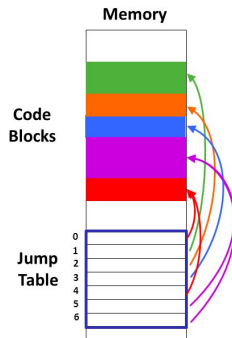
```

We can use the jump table when $x \leq 6$:

```

if (x <= 6)
  target = JTab[x];
  goto *target;
else
  goto default;

```



Winter 2013

x86 Programming III

5

Jump Table 优化

函数/过程的中间代码翻译

$n = f(a[i])$

1) $t_1 = i * 4$

2) $t_2 = a[t_1]$

3) **param t_2**

4) $t_3 =$ **call f, 1**

5) $n = t_3$

新增文法以支持函数定义与调用

$$\begin{aligned} D &\rightarrow \text{define } T \text{ id } (F) \{ S \} \\ F &\rightarrow \epsilon \mid T \text{ id } , F \\ S &\rightarrow \text{return } E ; \\ E &\rightarrow \text{id } (A) \\ A &\rightarrow \epsilon \mid E , A \end{aligned}$$

函数定义

$$\begin{array}{ll} D & \rightarrow \text{define } T \text{ id } (F) \{ S \} \\ F & \rightarrow \epsilon \mid T \text{ id } , F \\ S & \rightarrow \text{return } E ; \end{array}$$

函数名 `id` 放入当前符号表, 建立新的符号表, 处理形参 F 与函数体 S

函数调用

$$\begin{aligned} E &\rightarrow \text{id} (A) \\ A &\rightarrow \epsilon \mid E , A \end{aligned}$$

```
param  $x_1$   
param  $x_2$   
...  
param  $x_n$   
call  $p, n$ 
```

函数调用

```
S::=CALL id(Elist) { S.code := Elist.code  
A      || gencode("CALL", id.place, Elist.number) }  
Elist::=Elist1, E { Elist.code := E.code || Elist1.code 逆序  
      || gencode("PARAM", E.place);  
      Elist.number := Elist1.number + 1 }  
Elist::=E { Elist.code := E.code || gencode("PARAM", E.place);  
      Elist.number := 1 }
```

C 语言并未规定参数计算的顺序

函数调用

```
S ::= CALL id(Elist) A
{ Count := 0; S.code := Elist.code;
  while NOT EmptyQ(q) do
  begin
    t := HeadQ(q);
    S.code := S.code || gencode("PARAM", t);
    DelQ(q); Count := Count + 1
  end;
  S.code := S.code || gencode("CALL", id.place, Count)
}
```

逆序

```
Elist ::= Elist1, E { Elist.code := E.code || Elist1.code;
                      EnterQ(E.place, q) }
Elist ::= E { Elist.code := E.code; CreateQ(q);
             EnterQ(E.place, q) }
```

集中生成 **param** 指令, 代码更紧凑

Thank
You!



Office 926

hfwei@nju.edu.cn