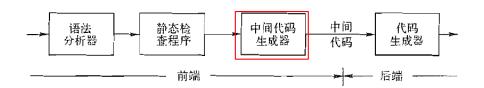
# 中间代码生成

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## Intermediate Representation (IR)



精确:不能丢失源程序的信息

独立: 不依赖特定的源语言与目标语言

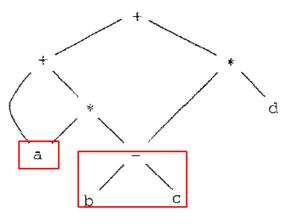
(如,没有复杂的寻址方式)

## Intermediate Representation (IR)



图 (抽象语法树)、三地址代码、C 语言

#### 表达式的有向无环图

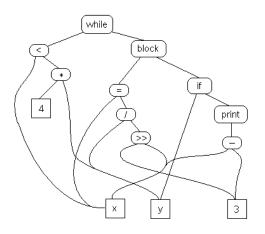


$$a + a * (b - c) + (b - c) * d$$

产生式		语义规则		
1)	$E  ightarrow \overline{E_1 + T}$	$E.node = \frac{\text{new Node}('+', E_1.node, T.node)}{}$		
2)	$E \rightarrow E_1 - T$	$E.node = $ $new Node('-', E_1.node, T.node)$		
3)	$E \to T$	E.node = T.node		
ļ	$T \rightarrow T_1 *F$	$T.node = $ <b>new</b> $Node('*', T_1.node, F.node)$		
4)	T  ightarrow ( $E$ )	T.node = E.node		
5)	$T  o \mathrm{id}$	$T.node = \frac{\mathbf{new}}{\mathbf{new}} Leaf(\mathbf{id}, \mathbf{id}.entry)$		
6)	$T \rightarrow \text{num}$	T.node =  new $Leaf(num, num.val)$		

在创建节点之前, 先判断是否已存在 (哈希表)

```
while (x < 4 * y) {
    x = y / 3 >> x;
    if (y) print x - 3;
}
```



# Definition (三地址代码 (Three-Address Code (TAC; 3AC))) 每个 **TAC** 指令**最多**包含三个操作数。

$$x = y \mathbf{op} z \tag{1}$$

 $x = \mathbf{op} \ y \tag{2}$ 

$$x = y \tag{3}$$

 $\mathbf{goto}\ L$  (4)

if x goto L (5)

if False x goto L (6)

if x relop y goto L (7)

# Definition (三地址代码 (Three-Address Code (TAC; 3AC)))

每个 TAC 指令最多包含三个操作数。

		$\mathtt{param}\ x_\mathtt{l}$
		$\mathtt{param}\ x_2$
$\mathbf{param}\;x$	(8)	
$\mathbf{call}\; p, n$	(9)	param $x_n$
$y=\mathbf{call}\; p,n$	(10)	call $p, n$
$\mathbf{return}\ y$	(11)	5-1-1 p, w
		$p(x_1, x_2, \ldots, x_n)$

# Definition (三地址代码 (Three-Address Code (TAC; 3AC)))

# 每个 TAC 指令最多包含三个操作数。

$$x = y[i] (12) x = &y (14)$$

$$x[i] = y (13) x = *y (15)$$

# 距离位置 y 处 i 个内存单元 \*x = y (16)

L: 
$$t_1 = i + 1$$
  
 $i = t_1$   
 $t_2 = i * 8$   
 $t_3 = a [t_2]$   
if  $t_3 < v$  goto L

```
100: t_1 = i + 1

101: i = t_1

102: t_2 = i * 8

103: t_3 = a [t_2]

104: if t_3 < v goto 100
```

#### 三地址代码的四元式表示

# Definition (四元式 (Quadruple))

一个四元式包含四个字段, 分别为 op、 $arg_1$ 、 $arg_2$  与 result。

$$a + a * (b - c) + (b - c) * d$$

$$t_1 = minus c$$
 $t_2 = b * t_1$ 
 $t_3 = minus c$ 
 $t_4 = b * t_3$ 
 $t_5 = t_2 + t_4$ 
 $a = t_5$ 

	о́р	arg,	$arg_2$	result
0	minus	С	ı	tı
1	*	Ъ	t <sub>1</sub>	t <sub>2</sub>
2	minus	C	[	t <sub>3</sub>
3	*	ь	t <sub>3</sub>	t4
4	+	$t_2$	t4	t <sub>5</sub>
5	=	$t_5$	1	, a
			•	

$$x = y[i]$$
$$x[i] = y$$

$$= [ ] \qquad y \qquad i \qquad x$$

$$[ ] = \qquad i \qquad y \qquad x$$

$$x = &y$$
$$x = *y$$
$$*x = y$$

$$= & y & x \\ = * & y & x \\ * = & y & x$$

#### 表达式的中间代码翻译

产生式	语义规则
$S \rightarrow id = E$ ;	S.code = E.code   $gen(top.get(id.lexeme))' = 'E.addr)$
$E \rightarrow E_1 + E_2$	$E.addr = \mathbf{new} \ Temp()$ $E.code = E_1.code \mid\mid E_2.code \mid\mid$ $gen(E.addr'='E_1.addr'+'E_2.addr)$
- E <sub>i</sub>	$E.addr = \mathbf{new} \ Temp() \ E.code = E_1.code \mid \mid \ gen(E.addr'=' 'minus' \ E_1.addr)$
[ (E <sub>1</sub> )	$E.addr = E_1.addr$ $E.code = E_1.code$
id	E.addr = top.get(id.lexeme) 符号表条目 E.code = ''

# 综合属性 E.code 与 E.addr

产生式	语义规则
$S \rightarrow id = E$ ;	$S.code = E.code \mid gen(top.get(id.lexeme))' = 'E.addr)$
$E \rightarrow E_1 + E_2$	$ \begin{split} E.addr &= \mathbf{new} \ Temp  () \\ E.code &= E_1.code \mid\mid E_2.code \mid\mid \\ gen(E.addr'=' E_1.addr'+' E_2.addr) \end{split} $
- E <sub>i</sub>	$E.addr = \mathbf{new} \; Temp  ()$ $E.code = E_1.code \mid   gen(E.addr'=' 'minus' \; E_1.addr)$
[ (E <sub>1</sub> )	$E.add\tau = E_1.add\tau$ $E.code = E_1.code$
id	E.addr = top.get(id.lexeme) 符号表条目 E.code = ''

$$t_1 = minus c$$
  
 $t_2 = b + t_1$   
 $a = t_2$ 

$$a = b + -c$$

## 表达式的中间代码翻译 (增量式)

$$S \rightarrow id = E$$
; {  $gen(top.get(id.lexeme) '=' E.addr)$ ; }

 $E \rightarrow E_1 + E_2$  {  $E.addr = new Temp()$ ;  $gen(E.addr '=' E_1.addr '+' E_2.addr)$ ; }

|  $-E_1$  {  $E.addr = new Temp()$ ;  $gen(E.addr '=' minus' E_1.addr)$ ; }

|  $(E_1)$  {  $E.addr = E_1.addr$ ; }

|  $id$  {  $E.addr = top.get(id.lexeme)$ ; }

#### 综合属性 E.addr

#### 数组引用的中间代码翻译

声明: int a[2][3]

数组引用: x = a[1][2]; a[1][2] = x

需要计算 a[1][2] 的相对于数组基地址 a 的偏移地址

#### 数组引用的中间代码翻译

# int a[2][3]

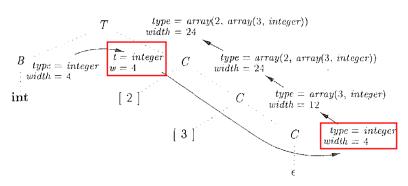


图 6-16 数组类型的语法制导翻译

## 数组类型声明

# int a[2][3]

array(2, array(3, integer))

	元素类型	元素宽度
a[]	array(3, integer)	12
a[][]	integer	4

$$addr(a[1][2]) = base + 1 \times 12 + 2 \times 4$$

```
S \rightarrow id = E; { gen(top.get(id.lexeme)' = 'E.addr); }
       L = E:
                   \{ gen(L.array.base' ['L.addr']' '='E.addr); \}
E \rightarrow E_1 + E_2 + E_3 { E.addr = new Temp();
                    gen(E.addr'='E_1.addr'+'E_2.addr);
      id
                  \{E.addr = top.get(id.lexeme);\}
    \mid L \mid
                   \{E.addr = new Temp();
                    gen(E.addr'=' L.array.base'[' L.addr']'); }
L \rightarrow \operatorname{id} [E]
                   \{L.array = top.get(id.lexeme):
                    L.type = L.array.type.elem;
                    L.addr = new Temp();
                    qen(L.addr'='E.addr'*'L.type.width);
                   \{L.array = L_1.array:
                    L.type = L_1.type.elem;
                    t = new Temp():
                    L.addr = new Temp();
                    qen(t'='E.addr'*'L.type.width);
                    qen(L.addr'='L_1.addr'+'t);
```

# 综合属性 L.array.base: 数组基地址 (即,数组名)

```
S \rightarrow id = E; { gen(top.get(id.lexeme)' = 'E.addr); }
    | L = E ; { gen(L.array.base' ['L.addr']' '='E.addr);
E \rightarrow E_1 + E_2 + E_2 { E.addr = new Temp();
                      gen(E,addr'='E_1,addr'+'E_2,addr);
       id
                    \{E.addr = top.get(id.lexeme);\}
                    { E.addr = new \ Temp();

gen(E.addr'=' \ L.array.base'[' \ L.addr']'); }
    L
```

#### 综合属性 L.addr: 偏移地址

# 综合属性 L.array: 数组名 id对应的符号表条目

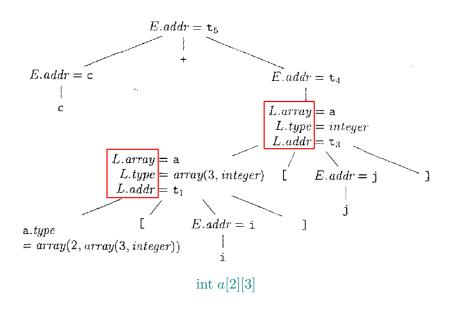
```
L \rightarrow id [E] \{L.array = top.get(id.lexeme);
                    L.type = L.array.type.elem;
                    L.addr = new Temp();
                    qen(L.addr'='E.addr'*'L.type.width); }
   L_1 \ [E] \ \{L.array = L_1.array;
                    L.type = L_1.type.elem:
                    t = \mathbf{new} \ Temp();
                    L.addr = new Temp();
                    qen(t'='E.addr'*'L.type.width);
                    qen(L.addr'='L_1.addr'+'t);
```

## 综合属性 L.type: (当前) 元素类型

```
L \rightarrow id [E] \{L.array = top.get(id.lexeme);
                    L.type = L.array.type.elem;
                    L.addr = \mathbf{new} \ Temp();
                    qen(L.addr'='E.addr'*'L.type,width);}
   L_1 [E] \{L.array = L_1.array;
                    L.type = L_1.type, elem;
                    t = \mathbf{new} \ Temp():
                    L.addr = new Temp():
                    gen(t'='E.addr'*'L.type.width);
                    gen(L.addr'='L_1.addr'+'t);
```

## 综合属性 L.addr: (当前) 偏移地址

```
L \rightarrow id [E] \{L.array = top.get(id.lexeme);
                    L.type = L.array.type.elem;
                    L.addr = \mathbf{new} \ Temp();
                    gen(L.addr'='E.addr'*'L.type.width); 
    L_1 [E] \{L.array = L_1.array;
                    L.type = L_1.type.elem;
                    t = \mathbf{new} \ Temp();
                    L.addr = new Temp();
                    gen(t'='E.addr'*'L.type.width);
                    gen(L.addr'='L_1.addr'+'t);
```



$$t_1 = i * 12$$
 $t_2 = j * 4$ 
 $t_3 = t_1 + t_2$ 
 $t_4 = a [t_3]$ 
 $t_5 = c + t_4$ 

int a[2][3]

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#### 控制流语句的中间代码翻译

$$S \rightarrow \text{ if } (B) S_1$$
  
 $S \rightarrow \text{ if } (B) S_1 \text{ else } S_2$   
 $S \rightarrow \text{ while } (B) S_1$ 

产生式	语义规则
$P \rightarrow S$	S.next = newlabel() $P.code = S.code \mid\mid label(S.next)$
$S \rightarrow assign$	S.code = assign.code
$S \rightarrow \text{if } (B) S_1$	$\begin{array}{lll} B.true &= newlabel() \\ B.false &= S_1.next &= S.next \\ S.code &= B.code \mid\mid label(B.true) \mid\mid S_1.code \end{array}$
$S \rightarrow \text{if } (B) S_1 \text{ else } S_2$	$B.true = newlabel() \\ B.false = newlabel() \\ S_1.next = S_2.next = S.next \\ S.code = B.code \\    label(B.true)    S_1.code \\    gen('goto' S.next) \\    label(B.false)    S_2.code$
$S \rightarrow $ while $(B) S_1$	begin = newlabel()
	B.truc = newlabel()   B.false = S.next
	$S_1.next = S.next$ $S_1.next = begin$
	S.code = label(begin)    B.code $   label(B.true)    S_1.code$    gen('goto' begin)
$S \rightarrow S_1 S_2$	$S_1.next = newlabel()$ $S_2.next = S.next$ $S.code = S_1.code \mid\mid label(S_1.next) \mid\mid S_2.code$

#### 继承属性 S.next: S 的下一条指令

$$P \rightarrow S$$
  $S.next = newlabel()$   $P.code = S.code || label(S.next)$ 

S.next 为语句 S 指明了"跳出"S 的目标

 $S \rightarrow assign$ 

S.code = assign.code

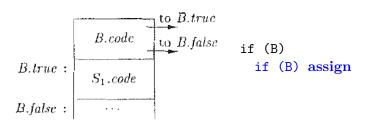
表达式语句的翻译,包括数组引用

$$S \rightarrow \mathbf{if} (B) S_1$$

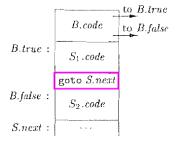
$$B.true = newlabel()$$

$$B.false = S_1.next = S.next$$

$$S.code = B.code || label(B.true) || S_1.code$$



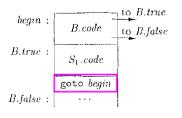
```
S 	o 	ext{if } (B) S_1 	ext{ else } S_2 \ egin{array}{ll} B. true &= newlabel() \ B. false &= newlabel() \ \hline S_1.next &= S_2.next &= S.next \ S.code &= B.code \ &\parallel label(B.true) \parallel S_1.code \ &\parallel gen('goto' \ S.next) \ &\parallel label(B.false) \parallel S_2.code \ \end{array}
```



if (B) if (B) assign else assign else assign

```
S \rightarrow \text{ while } (B) S_1
```

```
\begin{array}{ll} begin = newlabel() \\ B.true = newlabel() \\ B.false = S.next \\ \hline S_1.next = begin \\ S.code = label(begin) \mid\mid B.code \\ \mid\mid label(B.true) \mid\mid S_1.code \\ \mid\mid gen('goto'\ begin) \end{array}
```



while (B)
if (B) assign else assign

$$S \rightarrow S_1 S_2$$

$$\begin{array}{ll} S_1.next &=& newlabel() \\ S_2.next &=& S.next \\ S.code &=& S_1.code \mid\mid label(S_1.next) \mid\mid S_2.code \end{array}$$

产生式	语义规则
$P \rightarrow S$	
$S \rightarrow assign$	S.code = assign.code
$S \rightarrow \mathbf{if}(B) S_1$	$B.true = newlabel()$ $B.false =  S_1.next  = S.next$ $S.code =  B.code    label(B.true)    S_1.code $
$S \rightarrow \text{if } (B) S_1 \text{ else } S_2$	$B.true = newlabel() \\ B.false = newlabel() \\ [S_1.next = S_2.next] = S.next \\ S.code = B.code \\    label(B.true)    S_1.code \\    gen('goto' S.next) \\    label(B.false)    S_2.code$
$S \rightarrow $ while $(B) S_1$	begin = newlabel()
	$B.true = newlabel()$ $B.false = S.next$ $\boxed{S_1.next} = begin$ $S.code = label(begin)    B.code$ $   label(B.true)    S_1.code$ $   gen('goto' begin)$
$S \rightarrow S_1 S_2$	

# 布尔表达式的中间代码翻译

产生式	语义规则		
$B \rightarrow B_1 \mid \mid B_2$	$B_1.true = B.true$ $B_1.false = newlabel()$		
	$B_2.truc = B.true$ $B_2.false = B.false$ $B.code = B_1.code \mid\mid label(B_1.false) \mid\mid B_2.code$		
$B \rightarrow B_1 \&\& B_2$	$B_1.true = newlabel()$ $B_1.false = B.false$ $B_2.true = B.true$ $B_2.false = B.false$ $B.code = B_1.code \mid\mid label(B_1.true) \mid\mid B_2.code$		
$B \rightarrow ! B_1$	$B_1.true = B.false$ $B_1.false = B.true$ $B.code = B_1.code$		
$B \rightarrow E_1 \text{ rel } E_2$	$B.code = E_1.code \mid\mid E_2.code \mid\mid gen('if' E_1.addr rel.op E_2.addr 'goto' B.true) \mid\mid gen('goto' B.false)$		
$B \rightarrow  ext{true}$	B.code = gen('goto' B.true)		
$B \rightarrow false$	B.code = gen('goto' B.false)		

$$B \rightarrow \text{true}$$

$$B \rightarrow \mathbf{false}$$

$$B.code = gen('goto' B.true)$$

$$B.code = gen('goto' B.false)$$

#### if (true) assign

$$S \rightarrow \mathbf{if} (B) S_1$$

$$\begin{array}{ll} B.true &= newlabel() \\ B.false &= S_1.next \\ S.code &= B.code \mid\mid label(B.true) \mid\mid S_1.code \end{array}$$

## if (false) assign

$$B \rightarrow ! B_1$$

$$B_1.true = B.false$$
  
 $B_1.false = B.true$   
 $B.code = B_1.code$ 

#### if (!true) assign

$$S \rightarrow \mathbf{if} (B) S_1$$

```
\begin{array}{lll} B.true &= & newlabel() \\ B.false &= & S_1.next \\ S.code &= & B.code \mid \mid label(B.true) \mid \mid S_1.code \end{array}
```

#### if (!false) assign

#### 短路求值

# if (true || false) assign

$$S \rightarrow \mathbf{if} (B) S_1$$

$$B.true = \underbrace{newlabel()}_{B.false} = \underbrace{S_1.next}_{S.code} = S.next$$

$$S.code = B.code || label(B.true) || S_1.code$$

#### if (false || true) assign

#### 短路求值

$$B \rightarrow B_1 \&\& B_2 \mid B_1.true = newlabel() \ \mid B_1.false \mid B.false \mid B_2.true = B.true \mid B_2.false \mid B.false \mid B.code \mid B_1.code \mid abel(B_1.true) \mid B_2.code \mid abel(B_1.true) \mid B_2.code \mid abel(B_1.true) \mid abel(B_1$$

## if (true && false) assign

$$S \rightarrow \mathbf{if} (B) S_1$$

$$B.true = \underbrace{newlabel()}_{B.false} = \underbrace{S_1.next}_{S.code} = S.next$$

$$S.code = B.code || label(B.true) || S_1.code$$

## if (false && true) assign

$$B \rightarrow E_1 \text{ rel } E_2$$
 |  $B.code = E_1.code \mid\mid E_2.code$  |  $||gen('if' E_1.addr \text{ rel.}op E_2.addr 'goto' B.true)$  |  $||gen('goto' B.false)$ 

```
if (x < 100 \mid | x > 200 \&\& x != y) x = 0;
```

```
if x < 100 goto L_2
       goto \mathsf{L}_3
L_3: if x > 200 goto L_4 goto L_1
 L_4: if x != y goto L_2
goto L_1
L_2: x = 0
```

# Thank You!



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