

中间代码生成

(1. 表达式的翻译与控制流的翻译)

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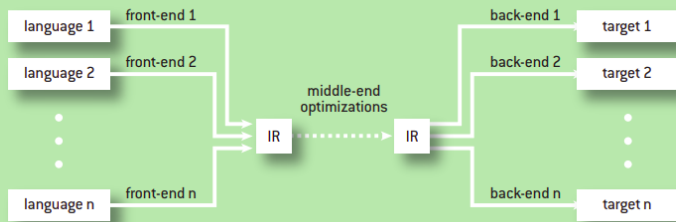
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FIGURE 2

A Compiler System Supporting Multiple Languages and Multiple Targets



The Increasing Significance of Intermediate Representations in Compilers (Fred Chow; 2013)

表达式的中间代码翻译

$$a = b + -c$$

| 产生式 | 语义规则 |
|---------------------------------|---|
| $S \rightarrow \text{id} = E ;$ | $S.code = E.code \parallel$ $\text{gen}(top.get(id.lexeme) \neq E.addr)$ |
| $E \rightarrow E_1 + E_2$ | $E.addr = \text{new Temp}()$ 临时变量: 虚拟寄存器 $E.code = E_1.code \parallel E_2.code \parallel$ $\text{gen}(E.addr \neq E_1.addr + E_2.addr)$ |
| $ - E_1$ | $E.addr = \text{new Temp}()$ $E.code = E_1.code \parallel$ $\text{gen}(E.addr \neq \text{'minus'} E_1.addr)$ |
| $ (E_1)$ | $E.addr = E_1.addr$ 生成代码 $E.code = E_1.code$ |
| $ \text{id}$ | $E.addr = top.get(id.lexeme)$ 符号表条目 $E.code = ''$ |

$t_1 = \text{minus } c$
 $t_2 = b + t_1$
 $a = t_2$

E.code: 中间代码

E.addr: 变量名 (包括中间变量)、常量

表达式的中间代码翻译 (增量式)

```

$$\begin{aligned} S &\rightarrow id = E ; \quad \{ gen(top.get(id.lexeme) \neq E.addr); \} \\ E &\rightarrow E_1 + E_2 \quad \{ E.addr = new Temp(); \\ &\quad \quad \quad gen(E.addr \neq E_1.addr '+' E_2.addr); \} \\ &\quad | - E_1 \quad \{ E.addr = new Temp(); \\ &\quad \quad \quad gen(E.addr \neq 'minus' E_1.addr); \} \\ &\quad | ( E_1 ) \quad \{ E.addr = E_1.addr; \} \\ &\quad | id \quad \{ E.addr = top.get(id.lexeme); \} \end{aligned}$$

```

```
int main() {  
    int a = 0, b = 1, c = 2;  
  
    a = b + -c;  
  
    return 0;  
}
```

```
%7 = sub nsw i32 0, %6  
%8 = add nsw i32 %5, %7  
store i32 %8, i32* %2, align 4
```

数组引用的中间代码翻译

声明 : $\text{int } a[2][3]$

数组引用 : $x = a[1][2]; a[1][2] = x$

需要计算 $a[1][2]$ 的相对于**数组基地址** a 的**偏移地址**

`int a[2][3]`

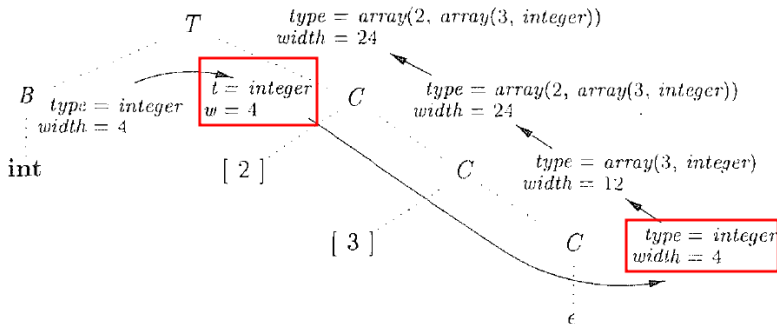


图 6-16 数组类型的语法制导翻译

数组类型声明

`int a[2][3]`

$array(2, array(3, integer))$

| | 类型 | 宽度 |
|-----------|-------------------------------|----|
| a | $array(2, array(3, integer))$ | 24 |
| $a[i]$ | $array(3, integer)$ | 12 |
| $a[i][j]$ | integer | 4 |

$$addr(a[1][2]) = base + 1 \times 12 + 2 \times 4$$


```

S → id = E ; { gen( top.get(id.lexeme) '=' E.addr); }

| L = E ; { gen(L.array.base '[' L.addr ']' '=' E.addr); }

E → E1 + E2 { E.addr = new Temp();
                gen(E.addr '=' E1.addr '+' E2.addr); }

| id { E.addr = top.get(id.lexeme); }

| L { E.addr = new Temp();
      gen(E.addr '=' L.array.base '[' L.addr ']'); }

L → id [ E ] { L.array = top.get(id.lexeme);
               L.type = L.array.type.elem;
               L.addr = new Temp();
               gen(L.addr '=' E.addr '*' L.type.width); }

| L1 [ E ] { L.array = L1.array;
               L.type = L1.type.elem;
               t = new Temp();
               L.addr = new Temp();
               gen(t '=' E.addr '*' L.type.width);
               gen(L.addr '=' L1.addr '+' t); }

```

int a[2][3]

综合属性 $L.array.base$: 数组基地址 (即, 数组名)

$S \rightarrow id = E ; \quad \{ gen(top.get(id.lexeme) \neq E.addr); \}$

$\quad | \quad L = E ; \quad \{ gen(L.array.base '[' L.addr ']' \neq E.addr); \}$

$E \rightarrow E_1 + E_2 \quad \{ E.addr = new Temp();$
 $\quad \quad \quad gen(E.addr \neq E_1.addr '+' E_2.addr); \}$

$\quad | \quad id \quad \quad \quad \{ E.addr = top.get(id.lexeme); \}$

$\quad | \quad L \quad \quad \quad \{ E.addr = new Temp();$
 $\quad \quad \quad gen(E.addr \neq L.array.base '[' L.addr ']); \}$

综合属性 $L.addr$: 偏移地址

综合属性 $L.array$: 数组名 id 对应的符号表条目

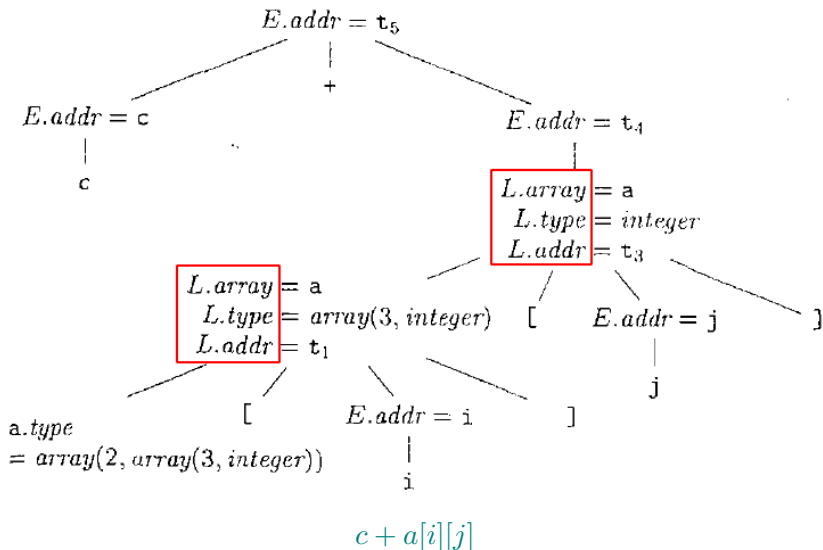
```
 $L \rightarrow id [ E ] \quad \{ L.array = top.get(id.lexeme);$   
     $L.type = L.array.type.elem;$   
     $L.addr = new Temp();$   
     $gen(L.addr '=' E.addr '*' L.type.width); \}$   
  
 $| \quad L_1 [ E ] \quad \{ L.array = L_1.array;$   
     $L.type = L_1.type.elem;$   
     $t = new Temp();$   
     $L.addr = new Temp();$   
     $gen(t '=' E.addr '*' L.type.width);$   
     $gen(L.addr '=' L_1.addr '+' t); \}$ 
```

综合属性 $L.type$: (当前) 元素类型

```
 $L \rightarrow id [ E ] \quad \{ L.array = top.get(id.lexeme);$   
 $\quad L.type = L.array.type.elem;$   
 $\quad L.addr = new Temp();$   
 $\quad gen(L.addr '=' E.addr '*' L.type.width); \}$   
  
 $| \quad L_1 [ E ] \quad \{ L.array = L_1.array;$   
 $\quad L.type = L_1.type.elem;$   
 $\quad t = new Temp();$   
 $\quad L.addr = new Temp();$   
 $\quad gen(t '=' E.addr '*' L.type.width);$   
 $\quad gen(L.addr '=' L_1.addr '+' t); \}$ 
```

综合属性 $L.addr$: (当前) 偏移地址

```
 $L \rightarrow id [ E ] \quad \{ L.array = top.get(id.lexeme);$   
     $L.type = L.array.type.elem;$   
     $L.addr = new Temp();$   
     $gen(L.addr '=' E.addr '*' L.type.width); \}$   
  
 $| L_1 [ E ] \quad \{ L.array = L_1.array;$   
     $L.type = L_1.type.elem;$   
     $t = new Temp();$   
     $L.addr = new Temp();$   
     $gen(t '=' E.addr '*' L.type.width);$   
     $gen(L.addr '=' L_1.addr '+' t); \}$ 
```



$$\begin{array}{l} t_1 = i * 12 \\ t_2 = j * 4 \\ t_3 = t_1 + t_2 \\ t_4 = a [t_3] \\ t_5 = c + t_4 \end{array}$$

$$c + a[i][j]$$

`%2 = alloca [2 x [3 x i32]], align 16`

```
int main() {  
    int a[2][3] = { 0 };  
  
    int i = 1, j = 2;  
    int c = 10, d = 20;  
  
    d = c + a[i][j];  
  
    return 0;  
}
```

```
%8 = load i32, i32* %5, align 4 %8: c  
%9 = load i32, i32* %3, align 4 %9: i  
%10 = sext i32 %9 to i64  
%11 = getelementptr inbounds [2 x [3 x i32]], [2 x [3 x i32]]* %2, i64 0, i64 %10  
%12 = load i32, i32* %4, align 4 %12: j  
%13 = sext i32 %12 to i64  
%14 = getelementptr inbounds [3 x i32], [3 x i32]* %11, i64 0, i64 %13  
%15 = load i32, i32* %14, align 4 %15: a[i][j]  
%16 = add nsw i32 %8, %15  
store i32 %16, i32* %6, align 4
```


控制流语句与布尔表达式的中间代码翻译

$$S \rightarrow \text{if } (B) \ S_1$$
$$S \rightarrow \text{if } (B) \ S_1 \ \text{else } S_2$$
$$S \rightarrow \text{while } (B) \ S_1$$

布尔表达式的作用: 布尔值 *vs.* 控制流跳转

$S \rightarrow \text{id} = E; \mid \text{if } (E) \ S \mid \text{while } (E) \ S \mid S \ S$
 $E \rightarrow E \parallel E \mid E \&\& E \mid E \text{ rel } E \mid E + E \mid (E) \mid \text{id} \mid \text{true} \mid \text{false}$

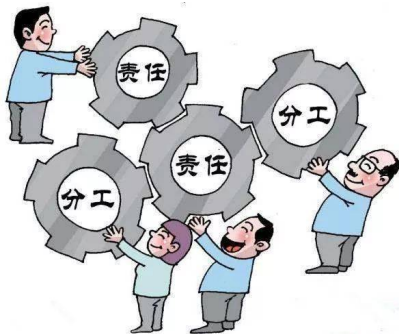
我们先关注“控制流跳转”

控制流语句与布尔表达式的中间代码翻译



分工明确

各司其职



父节点为子节点准备跳转指令的目标标签

子节点通过继承属性确定跳转目标

在自顶向下的分析过程中

为右部的每个 B 计算 $B.true$ 与 $B.false$

为右部的每个 S 计算 $S.next$

| 产生式 | 语义规则 |
|---------------------------------------|---|
| $P \rightarrow S$ | $S.next = newlabel()$ $P.code = S.code \parallel label(S.next)$ |
| $S \rightarrow assign$ | $S.code = assign.code$ |
| $S \rightarrow if (B) S_1$ | $B.true = newlabel()$ $B.false = S_1.next = S.next$ $S.code = B.code \parallel label(B.true) \parallel S_1.code$ |
| $S \rightarrow if (B) S_1 else S_2$ | $B.true = newlabel()$ $B.false = newlabel()$ $S_1.next = S_2.next = S.next$ $S.code = B.code$ $\parallel label(B.true) \parallel S_1.code$ $\parallel gen('goto' S.next)$ $\parallel label(B.false) \parallel S_2.code$ |
| $S \rightarrow while (B) S_1$ | $begin = newlabel()$ $B.true = newlabel()$ $B.false = S.next$ $S_1.next = begin$ $S.code = label(begin) \parallel B.code$ $\parallel label(B.true) \parallel S_1.code$ $\parallel gen('goto' begin)$ |
| $S \rightarrow S_1 S_2$ | $S_1.next = newlabel()$ $S_2.next = S.next$ $S.code = S_1.code \parallel label(S_1.next) \parallel S_2.code$ |

继承属性 $S.next$

$P \rightarrow S$

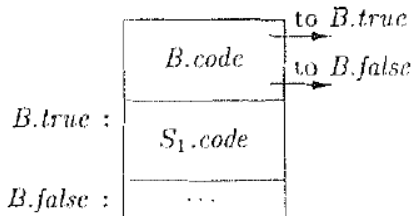
$$\left| \begin{array}{l} S.next = newlabel() \\ P.code = S.code || label(S.next) \end{array} \right.$$

$S.next$ 为语句 S 指明了“跳出” S 的目标

$$S \rightarrow \text{assign} \quad | \quad S.code = \text{assign}.code$$

代表了表达式的翻译, 包括数组引用

$$S \rightarrow \text{if}(B) S_1$$

$$\left\{ \begin{array}{l} B.true = \text{newlabel}() \\ B.false = S_1.next = S.next \\ S.code = B.code \parallel \text{label}(B.true) \parallel S_1.code \end{array} \right.$$


if (true)

if (false) assign

$$P \rightarrow S$$

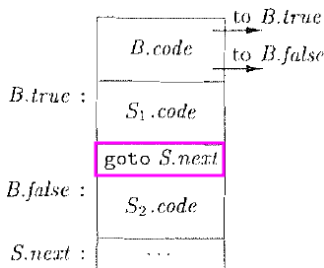
$$B \rightarrow \text{true}$$

$$B \rightarrow \text{false}$$

$$B.code = \text{gen}('goto' B.true)$$

$$B.code = \text{gen}('goto' B.false)$$

| | |
|---|---|
| $S \rightarrow \text{if} (B) S_1 \text{ else } S_2$ | $B.true = \text{newlabel}()$ $B.false = \text{newlabel}()$ $S_1.next = S_2.next = S.next$ $S.code = B.code$ $ \text{label}(B.true) S_1.code$ $ \text{gen}('goto' S.next)$ $ \text{label}(B.false) S_2.code$ |
|---|---|

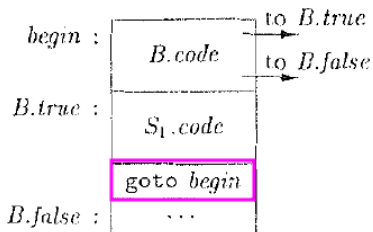


```

if (true)
    if (true) assign else assign
else
    assign
  
```

$S \rightarrow \text{while} (B) S_1$

```
begin = newlabel()
B.true = newlabel()
B.false = S.next
S1.next = begin
S.code = label(begin) || B.code
           || label(B.true) || S1.code
           || gen('goto' begin)
```



```
while (true)
    if (false) assign else assign
```

$S \rightarrow S_1 S_2$

$\begin{array}{|l} S_1.next = newlabel() \\ S_2.next = S.next \\ S.code = S_1.code || label(S_1.next) || S_2.code \end{array}$

if (true) assign else assign assign

布尔表达式的中间代码翻译

| 产生式 | 语义规则 |
|----------------------------------|--|
| $B \rightarrow B_1 \ \ B_2$ | $B_1.true = B.true$ $B_1.false = newlabel()$ $B_2.true = B.true$ $B_2.false = B.false$ $B.code = B_1.code \ \ label(B_1.false) \ \ B_2.code$ |
| $B \rightarrow B_1 \ \&\& \ B_2$ | $B_1.true = newlabel()$ $B_1.false = B.false$ $B_2.true = B.true$ $B_2.false = B.false$ $B.code = B_1.code \ \ label(B_1.true) \ \ B_2.code$ |
| $B \rightarrow ! B_1$ | $B_1.true = B.false$ $B_1.false = B.true$ $B.code = B_1.code$ |
| $B \rightarrow E_1 \ rel \ E_2$ | $B.code = E_1.code \ \ E_2.code$ $\quad \ \ gen('if' \ E_1.addr \ rel \ op \ E_2.addr \ 'goto' \ B.true)$ $\quad \ \ gen('goto' \ B.false)$ |
| $B \rightarrow true$ | $B.code = gen('goto' \ B.true)$ |
| $B \rightarrow false$ | $B.code = gen('goto' \ B.false)$ |

$B \rightarrow \text{true}$ $B.\text{code} = \text{gen}(\text{'goto' } B.\text{true})$

$B \rightarrow \text{false}$ $B.\text{code} = \text{gen}(\text{'goto' } B.\text{false})$

if (true) assign

$S \rightarrow \text{if}(B) S_1$ $\left\{ \begin{array}{l} B.\text{true} = \text{newlabel}() \\ B.\text{false} = S_1.\text{next} = S.\text{next} \\ S.\text{code} = B.\text{code} || \text{label}(B.\text{true}) || S_1.\text{code} \end{array} \right.$

if (false) assign

$B \rightarrow ! B_1$

$\left\{ \begin{array}{l} B_1.true = B.false \\ B_1.false = B.true \\ B.code = B_1.code \end{array} \right.$

if (!true) assign

$S \rightarrow \text{if}(B) S_1$

$\left\{ \begin{array}{l} B.true = \text{newlabel}() \\ B.false = S_1.next = S.next \\ S.code = B.code || \text{label}(B.true) || S_1.code \end{array} \right.$

if (!false) assign

短路求值

$$B \rightarrow B_1 \ || \ B_2 \quad \left| \begin{array}{l} B_1.true = B.true \\ B_1.false = newlabel() \\ B_2.true = B.true \\ B_2.false = B.false \\ B.code = B_1.code \ || \ label(B_1.false) \ || \ B_2.code \end{array} \right.$$

if (true || false) assign

$$S \rightarrow \text{if} (B) S_1 \quad \left| \begin{array}{l} B.true = newlabel() \\ B.false = S_1.next = S.next \\ S.code = B.code \ || \ label(B.true) \ || \ S_1.code \end{array} \right.$$

if (false || true) assign

短路求值

$$B \rightarrow B_1 \ \&\& \ B_2 \quad \left| \begin{array}{l} B_1.true = newlabel() \\ B_1.false = B.false \\ B_2.true = B.true \\ B_2.false = B.false \\ B.code = B_1.code \ || \ label(B_1.true) \ || \ B_2.code \end{array} \right.$$

if (true && false) assign

$$S \rightarrow \text{if} (B) S_1 \quad \left| \begin{array}{l} B.true = newlabel() \\ B.false = S_1.next = S.next \\ S.code = B.code \ || \ label(B.true) \ || \ S_1.code \end{array} \right.$$

if (false && true) assign

$$B \rightarrow E_1 \text{ rel } E_2 \quad \left| \quad \begin{array}{l} B.code = E_1.code \parallel E_2.code \\ \parallel \text{gen('if' } E_1.addr \text{ rel.op } E_2.addr \text{ 'goto' } B.true) \\ \parallel \text{gen('goto' } B.false) \end{array} \right.$$

```
if (x < 100 || x > 200 && x != y) x = 0;
```

```
        if x < 100 goto L2  
        goto L3  
L3:    if x > 200 goto L4  
        goto L1  
L4:    if x != y goto L2  
        goto L1  
L2:    x = 0  
L1:
```

布尔表达式的作用: 布尔值 *vs.* 控制流跳转

$S \rightarrow \text{id} = E; \mid \text{if} (E) S \mid \text{while} (E) S \mid S S$
 $E \rightarrow E \parallel E \mid E \&\& E \mid E \text{ rel } E \mid E + E \mid (E) \mid \text{id} \mid \text{true} \mid \text{false}$

根据 E 所处的上下文判断 E 所扮演的角色, 调用不同的代码生成函数

函数 $\text{jump}(t, f)$: 生成控制流代码

函数 $\text{rvalue}()$: 生成计算布尔值的代码, 并将结果存储在临时变量中

| 产生式 | 语义规则 |
|---------------------------|--|
| $S \rightarrow id = E ;$ | $S.code = E.code \parallel$ $gen(top.get(id.lexeme) '=' E.addr)$ |
| $E \rightarrow E_1 + E_2$ | $E.addr = new Temp()$ 临时变量: 虚拟寄存器 $E.code = E_1.code \parallel E_2.code \parallel$ $gen(E.addr '=' E_1.addr '+' E_2.addr)$ |
| $ - E_1$ | $E.addr = new Temp()$ $E.code = E_1.code \parallel$ $gen(E.addr '=' 'minus' E_1.addr)$ |
| $ (E_1)$ | $E.addr = E_1.addr$ 生成代码 $E.code = E_1.code$ |
| $ id$ | $E.addr = top.get(id.lexeme)$ 符号表条目 $E.code = ''$ |

$$E \rightarrow E_1 \&\& E_2$$

为 E 生成跳转代码, 在真假出口处将 true 或 false 存储到临时变量

`x = a < b && c < d`

```
    ifFalse a < b goto L1
    ifFalse c < d goto L1
    t = true
    goto L2
L1: t = false
L2: x = t
```

Thank
You!



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