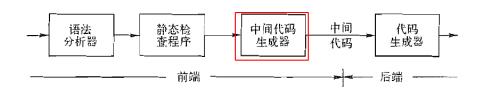
中间代码生成

魏恒峰

hfwei@nju.edu.cn

2020年12月22日





Intermediate Representation (IR)



Intermediate Representation (IR)



精确:不能丢失源程序的信息

独立: 不依赖特定的源语言与目标语言

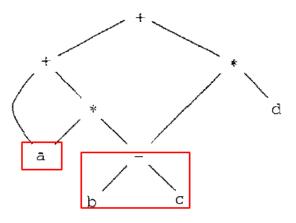
(如,没有复杂的寻址方式)

Intermediate Representation (IR)



图 (抽象语法树)、三地址代码、C 语言

表达式的有向无环图



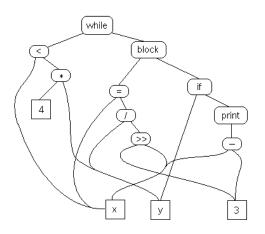
$$a + a * (b - c) + (b - c) * d$$

产生式		语义规则	
1)	$E \mapsto E_1 + T$	$E.node = \frac{\text{new Node}('+', E_1.node, T.node)}{}$	
2)	$E o E_1 - T$	$E.node = \frac{\text{new }Node('-', E_1.node, T.node)}{}$	
3)	$E \to T$	E.node = T.node	
	$T \rightarrow T_1 * F$	$T.node = $ new $Node('*', T_1.node, F.node)$	
4)	T ightarrow (E)	T.node = E.node	
5)	$T o \mathbf{id}$	$T.node = \frac{\mathbf{new}}{\mathbf{new}} Leaf(\mathbf{id}, \mathbf{id}.entry)$	
6)	$T \rightarrow \text{num}$	T.node = new $Leaf(num, num.val)$	

产生式		语义规则	
1)	$E ightarrow E_1 + T$	$E.node = \frac{\text{new } Node('+', E_1.node, T.node)}{}$	
2)	$E \rightarrow E_1 - T$	$E.node = \frac{\mathbf{new}}{\mathbf{Node}('-', E_1.node, T.node)}$	
3)	$E \to T$	E.node = T.node	
	$T \rightarrow T_1 *F$	$T.node = $ new $Node('*', T_1.node, F.node)$	
4)	T ightarrow (E)	T.node = E.node	
5)	$T o \mathrm{id}$	T.node = new $Leaf(id, id.entry)$	
6)	$T \rightarrow \text{num}$	T.node = $new $ $Leaf(num, num.val)$	

在创建节点之前, 先判断是否已存在 (哈希表)

```
while (x < 4 * y) {
    x = y / 3 >> x;
    if (y) print x - 3;
}
```



Definition (三地址代码 (Three-Address Code (TAC; 3AC))) 每个 **TAC** 指令最多包含三个操作数。

$$x = y \mathbf{op} z \tag{1}$$

$$x = \mathbf{op} \ y \tag{2}$$

$$x = y \tag{3}$$

Definition (三地址代码 (Three-Address Code (TAC; 3AC)))

每个 TAC 指令最多包含三个操作数。

$$x = y \mathbf{op} z \tag{1}$$

$$x = \mathbf{op} \ y \tag{2}$$

$$x = y \tag{3}$$

goto L (4)

if x goto L (5)

if False x goto L (6)

if x relop y goto L (7)

Definition (三地址代码 (Three-Address Code (TAC; 3AC)))

每个 TAC 指令最多包含三个操作数。

		$\mathtt{param}\ x_\mathtt{l}$
		$\mathtt{param}\ x_2$
$\mathbf{param}\;x$	(8)	
$\mathbf{call}\; p, n$	(9)	param x_n
$y=\mathbf{call}\; p,n$	(10)	call p, n
$\mathbf{return}\;y$	(11)	7-22 P, W
		$p(x_1, x_2, \ldots, x_n)$

Definition (三地址代码 (Three-Address Code (TAC; 3AC))) 每个 **TAC** 指令最多包含三个操作数。

$$x = y[i] \tag{12}$$

$$x[i] = y \tag{13}$$

距离位置 y 处 i 个内存单元

Definition (三地址代码 (Three-Address Code (TAC; 3AC)))

每个 TAC 指令最多包含三个操作数。

$$x = y[i] (12) x = &y (14)$$

$$x[i] = y (13) x = *y$$

距离位置 y 处 i 个内存单元 *x = y (16)

(15)

L:
$$t_1 = i + 1$$

 $i = t_1$
 $t_2 = i * 8$
 $t_3 = a [t_2]$
if $t_3 < v \text{ goto } L$

L:
$$t_1 = i + 1$$

 $i = t_1$
 $t_2 = i * 8$
 $t_3 = a [t_2]$
if $t_3 < v$ goto L

```
100: t_1 = i + 1

101: i = t_1

102: t_2 = i * 8

103: t_3 = a [t_2]

104: if t_3 < v goto 100
```

三地址代码的四元式表示

Definition (四元式 (Quadruple))

一个四元式包含四个字段, 分别为 op、 arg_1 、 arg_2 与 result。

三地址代码的四元式表示

Definition (四元式 (Quadruple))

一个四元式包含四个字段, 分别为 op、 arg_1 、 arg_2 与 result。

$$a + a * (b - c) + (b - c) * d$$

$$t_1 = minus c$$
 $t_2 = b * t_1$
 $t_3 = minus c$
 $t_4 = b * t_3$
 $t_5 = t_2 + t_4$
 $a = t_5$

	о́р	arg ₁	arg_2	result
0	minus	С	,	tı
1	*	Ъ	t ₁	t_2
2	minus	С	(t ₃
3	*	b	t ₃	t4
4	+	t ₂	t4	t ₅
5	=	t_5		, a

$$x = y[i]$$
$$x[i] = y$$

$$= [] \qquad y \qquad i \qquad x$$
$$[]= \qquad i \qquad y \qquad x$$

$$x = y[i]$$
$$x[i] = y$$

$$= [] \qquad y \qquad i \qquad x$$

$$[] = \qquad i \qquad y \qquad x$$

$$x = \&y$$
$$x = *y$$
$$*x = y$$

$$x = y[i]$$
$$x[i] = y$$

$$= [] \qquad y \qquad i \qquad x$$

$$[] = \qquad i \qquad y \qquad x$$

$$x = \&y$$
$$x = *y$$
$$*x = y$$

$$= & y & x \\ = * & y & x \\ * = & y & x$$

表达式的中间代码翻译

产生式	语义规则
$S \rightarrow id = E$;	S.code = E.code $gen(top.get(id.lexeme))' = 'E.addr)$
$E \rightarrow E_1 + E_2$	$ E.addr = \mathbf{new} \ Temp() \\ E.code = E_1.code \mid\mid E_2.code \mid\mid \\ gen(E.addr'='E_1.addr'+'E_2.addr) $
- E _i	$E.addr = \mathbf{new} \ Temp() \ E.code = E_1.code \mid \ gen(E.addr'=' 'minus' \ E_1.addr)$
[(E ₁)	$E.addr = E_1.addr \ E.code = E_1.code$
id	E.addr = top.get(id.lexeme) 符号表条目 E.code = ''

综合属性 E.code 与 E.addr

产生式	语义规则
$S \rightarrow id = E$;	$S.code = E.code \mid gen(top.get(id.lexeme))' = 'E.addr)$
$E \rightarrow E_1 + E_2$	$ \begin{split} E.addr &= \mathbf{new} \ Temp () \\ E.code &= E_1.code \mid \mid E_2.code \mid \mid \\ gen(E.addr'=' E_1.addr'+' E_2.addr) \end{split} $
- E _i	$E.addr = \mathbf{new} \ Temp()$ $E.code = E_1.code \mid gen(E.addr'=' 'minus' E_1.addr)$
[(E ₁)	$E.addr = E_1.addr$ $E.code = E_1.code$
id	E.addr = top.get(id.lexeme) 符号表条目 E.code = ''

$$t_1 = minus c$$

 $t_2 = b + t_1$
 $a = t_2$

$$a = b + -c$$

表达式的中间代码翻译 (增量式)

```
S \rightarrow id = E; { gen(top.get(id.lexeme)'='E.addr): }
E \rightarrow E_1 + E_2 \quad \{ E.addr = new Temp() : \}
                    gen(E.addr'='E_1,addr'+'E_2,addr);
      -E_1 { E.addr = new Temp();
                    gen(E.addr'=''minus' E_1.addr); 
      \{E_1 \mid E_1 \mid E_2 \mid E_1 \mid addr; \}
               \{E.addr = top.get(id.lexeme); \}
      id
```

综合属性 E.addr

数组引用的中间代码翻译

声明: int a[2][3]

数组引用: x = a[1][2]; a[1][2] = x

数组引用的中间代码翻译

声明: int a[2][3]

数组引用: x = a[1][2]; a[1][2] = x

需要计算 a[1][2] 的相对于**数组基地址** a 的**偏移地址**

数组引用的中间代码翻译

int a[2][3]

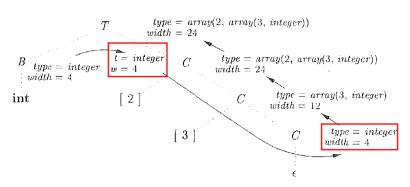


图 6-16 数组类型的语法制导翻译

数组类型声明

int a[2][3]

array(2, array(3, integer))

	元素类型	元素宽度
a[]	array(3, integer)	12
a[][]	integer	4

int a[2][3]

array(2, array(3, integer))

	元素类型	元素宽度
a[]	array(3, integer)	12
a[][]	integer	4

$$addr(a[1][2]) = base + 1 \times 12 + 2 \times 4$$

```
S \rightarrow id = E; { gen(top.get(id.lexeme)' = 'E.addr); }
      L = E:
                  \{ gen(L.array.base' ['L.addr']' '='E.addr); \}
E \rightarrow E_1 + E_2 + E_3 { E.addr = new Temp();
                    gen(E.addr'='E_1.addr'+'E_2.addr);
      id
                 \{E.addr = top.get(id.lexeme);\}
    \mid L \mid
                  \{E.addr = new Temp();
                    gen(E.addr'=' L.array.base'[' L.addr']'); }
                  { L.array = top.qet(id.lexeme):
                    L.type = L.array.type.elem;
                    L.addr = new Temp();
                    qen(L.addr'='E.addr'*'L.type.width);
                   \{L.array = L_1.array;
                    L.type = L_1.type.elem;
                    t = new Temp():
                    L.addr = new Temp();
                    qen(t'='E.addr'*'L.type.width);
                    qen(L.addr'='L_1.addr'+'t);
```

语义分析

综合属性 L.array.base: 数组基地址 (即,数组名)

```
S \rightarrow id = E; { gen(top.get(id.lexeme)' = 'E.addr); }
    | L = E ; { gen(L.array.base' ['L.addr']' '='E.addr);
E \rightarrow E_1 + E_2 + E_2 { E.addr = new Temp();
                      gen(E,addr'='E_1,addr'+'E_2,addr);
       id
                    \{E.addr = top.get(id.lexeme);\}
                    { E.addr = new \ Temp();

gen(E.addr'=' \ L.array.base'[' \ L.addr']'); }
    L
```

综合属性 L.addr: 偏移地址

综合属性 L.array: 数组名 id对应的符号表条目

```
L \rightarrow id [E] \{L.array = top.get(id.lexeme);
                    L.type = L.array.type.elem;
                    L.addr = new Temp();
                    qen(L.addr'='E.addr'*'L.type.width); }
   L_1 \ [E] \ \{L.array = L_1.array;
                    L.type = L_1.type.elem:
                    t = \mathbf{new} \ Temp();
                    L.addr = new Temp();
                    qen(t'='E.addr'*'L.type.width);
                    qen(L.addr'='L_1.addr'+'t);
```

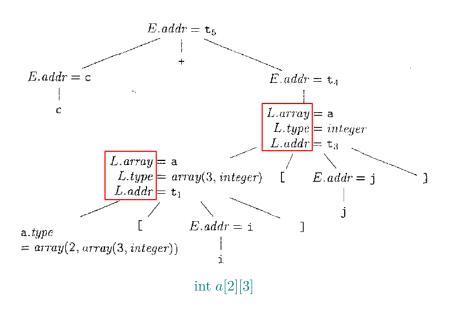
综合属性 L.type: (当前) 元素类型

```
L \rightarrow id [E] \{L.array = top.get(id.lexeme);
                    L.type = L.array.type.elem;
                    L.addr = \mathbf{new} \ Temp();
                    qen(L.addr'='E.addr'*'L.type.width); \}
    L_1 [E] \{L.array = L_1.array;
                    L.type = L_1.type, elem;
                    t = \mathbf{new} \ Temp():
                    L.addr = new Temp():
                    gen(t'='E.addr'*'L.type.width);
                    gen(L.addr'='L_1.addr'+'t);
```

综合属性 L.addr: (当前) 偏移地址

```
L \rightarrow id [E] \{L.array = top.get(id.lexeme)\}
                   L.type = L.array.type.elem;
                  L.addr = new Temp();
                   gen(L.addr'='E.addr'*'L.type.width);
   L_1 [E] \{L.array = L_1.array;
                   L.type = L_1.type.elem;
                   t = new Temp();
                  L.addr = new Temp();
                   gen(t'='E.addr'*'L.type.width);
                  gen(L.addr'='L_1.addr'+'t);
```

24 / 42



$$t_1 = i * 12$$
 $t_2 = j * 4$
 $t_3 = t_1 + t_2$
 $t_4 = a [t_3]$
 $t_5 = c + t_4$

int a[2][3]

26 / 42

控制流语句的中间代码翻译

$$S \rightarrow \text{ if } (B) S_1$$

 $S \rightarrow \text{ if } (B) S_1 \text{ else } S_2$
 $S \rightarrow \text{ while } (B) S_1$

产生式	语义规则
$P \rightarrow S$	S.next = newlabel() $P.code = S.code \mid label(S.next)$
$S \rightarrow assign$	S.code = assign.code
$S \rightarrow \mathbf{if}(B) S_1$	$ B.true = newlabel() \\ B.false = S_1.next = S.next \\ S.code = B.code \mid\mid label(B.true) \mid\mid S_1.code $
$S \rightarrow \text{if } (B) S_1 \text{ else } S_2$	$B.true = newlabel() \\ B.false = newlabel() \\ S_1.next = S_2.next = S.next \\ S.code = B.code \\ label(B.true) S_1.code \\ gen('goto' S.next) \\ label(B.false) S_2.code$
$S \rightarrow $ while $(B) S_1$	$begin = newlabel() \\ B.truc = newlabel() \\ B.false = S.next \\ S_1.next = begin \\ S.code = label(begin) B.code \\ label(B.true) S_1.code \\ gen('goto' begin)$
$S \rightarrow S_1 S_2$	$S_1.next = newlabel()$ $S_2.next = S.next$ $S.code = S_1.code \mid\mid label(S_1.next) \mid\mid S_2.code$

继承属性 S.next: S 的下一条指令

$$P \rightarrow S$$
 $S.next = newlabel()$ $P.code = S.code || label(S.next)$

S.next 为语句 S 指明了"跳出"S 的目标

 $S \rightarrow assign$

S.code = assign.code

表达式语句的翻译,包括数组引用

$$S \rightarrow \mathbf{if} (B) S_1$$

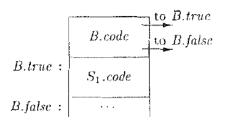
$$\begin{array}{lll} B.true &=& newlabel() \\ B.false &=& S_1.next \\ S.code &=& B.code \mid \mid label(B.true) \mid \mid S_1.code \end{array}$$

31/42

$$S \rightarrow \mathbf{if} (B) S_1$$

$$B.true = \underbrace{newlabel()}_{B.false} = \underbrace{S_1.next}_{S.code} = S.next$$

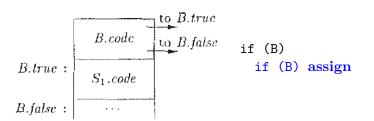
$$S.code = \underbrace{B.code \mid\mid label(B.true) \mid\mid S_1.code}_{label(B.true)}$$



$$S \rightarrow \mathbf{if} (B) S_1$$

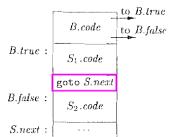
$$B.true = \underbrace{newlabel()}_{B.false} = \underbrace{S_1.next}_{S.code} = S.next$$

$$S.code = \underbrace{B.code}_{||} || \underbrace{label(B.true)}_{||} || S_1.code$$

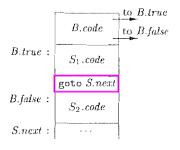


$$S o ext{if } (B) S_1 ext{ else } S_2 \ egin{array}{ll} B. true &= newlabel() \ B. false &= newlabel() \ \hline S_1.next &= S_2.next &= S.next \ S.code &= B.code \ & || label(B.true) || S_1.code \ & || gen('goto' S.next) \ & || label(B.false) || S_2.code \ \hline \end{array}$$

```
S 	o 	ext{if } (B) S_1 	ext{ else } S_2 \ egin{array}{ll} B. true &= newlabel() \ B. false &= newlabel() \ \hline S_1.next &= S_2.next &= S.next \ S.code &= B.code \ & || label(B.true) || S_1.code \ & || gen('goto' S.next) \ & || label(B.false) || S_2.code \ \hline \end{array}
```



```
S 	o 	ext{if } (B) S_1 	ext{ else } S_2 \ | egin{array}{ll} B. true &= newlabel() \ B. false &= newlabel() \ \hline S_1.next &= S_2.next &= S.next \ S.code &= B.code \ & || label(B.true) || S_1.code \ & || gen('goto' S.next) \ & || label(B.false) || S_2.code \ \hline \end{array}
```



if (B) if (B) assign else assign else assign

$$S \rightarrow$$
 while $(B) S_1$

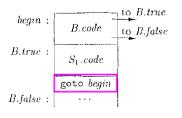
```
begin = newlabel()
B.true = newlabel()
B.false = S.next
S_1.next = begin
S.code = label(begin) || B.code
|| label(B.true) || S_1.code
|| gen('goto' begin)
```

```
S \rightarrow \text{ while } (B) S_1
```

```
\begin{array}{lll} begin &=& newlabel() \\ B.true &=& newlabel() \\ B.false &=& S.next \\ \hline S_1.next &=& begin \\ S.code &=& label(begin) \mid\mid B.code \\ &\mid\mid label(B.true) \mid\mid S_1.code \\ &\mid\mid gen('goto'\ begin) \end{array}
```

```
S \rightarrow while ( B ) S_1
```

```
\begin{array}{ll} begin = newlabel() \\ B.true = newlabel() \\ B.false = S.next \\ \hline S_1.next = begin \\ S.code = label(begin) \mid\mid B.code \\ \mid\mid label(B.true) \mid\mid S_1.code \\ \mid\mid gen('goto'\ begin) \end{array}
```



while (B)
if (B) assign else assign

33 / 42

$$S \rightarrow S_1 S_2$$

$$\begin{array}{ll} S_1.next \\ S_2.next \end{array} = \begin{array}{ll} newlabel() \\ = S.next \\ S.code = S_1.code \mid\mid label(S_1.next) \mid\mid S_2.code \end{array}$$

产生式	语义规则
$P \rightarrow S$	$ \begin{array}{ll} S.next = newlabel() \\ P.code = S.code \mid\mid label(S.next) \end{array} $
$S \rightarrow \text{assign}$	S.code = assign.code
$S \rightarrow \mathbf{if}(B) S_1$	$\begin{array}{lll} B.true &= & newlabel() \\ B.false &= & S_1.next \\ S.code &= & B.code \mid \mid label(B.true) \mid \mid S_1.code \end{array}$
$S \rightarrow \text{if } (B) S_1 \text{ else } S_2$	$B.true = newlabel() \\ B.false = newlabel() \\ [S_1.next = S_2.next] = S.next \\ S.code = B.code \\ label(B.true) S_1.code \\ gen('goto' S.next) \\ label(B.false) S_2.code$
$S \rightarrow \text{ while } (B) S_1$	begin = newlabel()
	$B.true = newlabel()$ $B.false = S.next$ $\boxed{S_1.next} = begin$ $S.code = label(begin) B.code$ $ label(B.true) S_1.code$ $ gen('goto' begin)$
$S \rightarrow S_1 S_2$	

布尔表达式的中间代码翻译

163.0000 1 1 3 1 4 1 Junior		
产生式	语义规则	
$B \rightarrow B_1 \mid \mid B_2$	$B_1.true = B.true$ $B_1.false = newlabel()$	
	$B_2.true = B.true$ $B_2.false = B.false$ $B.code = B_1.code \mid\mid label(B_1.false) \mid\mid B_2.code$	
$B \rightarrow B_1 \&\& B_2$	$B_1.true = newlabel()$ $B_1.false = B.false$ $B_2.true = B.true$ $B_2.false = B.false$ $B.code = B_1.code \mid\mid label(B_1.true) \mid\mid B_2.code$	
$B \rightarrow ! B_1$	$B_1.true = B.false$ $B_1.false = B.true$ $B.code = B_1.code$	
$B ightarrow E_1 \ ext{rel} \ E_2$	$B.code = E_1.code \mid\mid E_2.code$ $\mid\mid gen('if' E_1.addr rel.op E_2.addr 'goto' B.true)$ $\mid\mid gen('goto' B.false)$	
$B ightarrow { m true}$	B.code = gen('goto' B.true)	
$B \rightarrow \text{false}$	B.code = gen('goto' B.false)	

$$B \rightarrow \text{true}$$

$$B.code = gen('goto' B.true)$$

$$B \rightarrow \mathbf{false}$$

$$B.code = gen('goto' B.false)$$

$$B \rightarrow \text{true}$$

$$B.code = gen('goto' B.true)$$

$$B \rightarrow \mathbf{false}$$

$$B.code = gen('goto' B.false)$$

if (true) assign

$$S \rightarrow \mathbf{if} (B) S_1$$

$$\begin{array}{ll} B.true &= newlabel() \\ B.false &= S_1.next \\ S.code &= B.code \mid\mid label(B.true) \mid\mid S_1.code \end{array}$$

if (false) assign

$$B \rightarrow ! B_1$$

$$B_1.true = B.false$$

 $B_1.false = B.true$
 $B.code = B_1.code$

$$B \rightarrow ! B_1$$

$$B_1.true = B.false$$

 $B_1.false = B.true$
 $B.code = B_1.code$

if (!true) assign

$$S \rightarrow \mathbf{if} (B) S_1$$

$$\begin{array}{lll} B.true &= newlabel() \\ B.false &= S_1.next \\ S.code &= B.code \mid\mid label(B.true) \mid\mid S_1.code \end{array}$$

if (!false) assign

if (true || false) assign

$$S \rightarrow \mathbf{if} (B) S_1$$

$$B.true = \underbrace{newlabel()}_{S_1.next} = S.next$$

$$S.code = B.code || label(B.true) || S_1.code$$

if (false || true) assign

$$B \rightarrow B_1 \&\& B_2 \ | \begin{array}{c} B_1.true = newlabel() \\ B_1.false = B.false \\ B_2.true = B.true \\ \hline B_2.false = B.false \\ B.code = B_1.code \mid \mid label(B_1.true) \mid \mid B_2.code \end{array}$$

40/42

$$B \rightarrow B_1 \&\& B_2 \mid B_1.true = newlabel() \ \mid B_1.false = B.false \ \mid B_2.true = B.true \ \mid B_2.false = B.false \ \mid B.code = B_1.code \mid | label(B_1.true) \mid | B_2.code$$

if (true && false) assign

$$S \rightarrow \mathbf{if} (B) S_1$$

$$B.true = newlabel()$$

$$B.false = S_1.next = S.next$$

$$S.code = B.code || label(B.true) || S_1.code$$

if (false && true) assign

$$B \rightarrow E_1 \text{ rel } E_2$$

$$| B.code = E_1.code || E_2.code$$

$$|| gen('if' E_1.addr \text{ rel.}op E_2.addr 'goto' B.true)$$

$$|| gen('goto' B.false)$$

```
if (x < 100 \mid | x > 200 \&\& x != y) x = 0;
```

```
if x < 100 goto L_2
         goto \mathsf{L}_3
 L_3: if x > 200 goto L_4 goto L_1
 L_4: if x != y goto L_2

\begin{array}{ccc}
\text{goto } L_1 \\
L_2: & x = 0
\end{array}
```

Thank You!



Office 926 hfwei@nju.edu.cn