

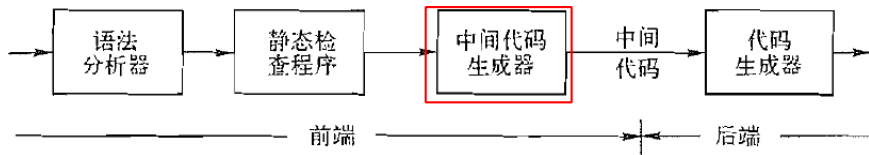
# 中间代码生成

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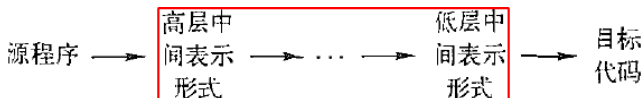
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2020 年 12 月 21 日





## Intermediate Representation (IR)



**精确:** 不能丢失源程序的信息

**独立:** 不依赖特定的源语言与目标语言  
(如, 没有复杂的寻址方式)

## Intermediate Representation (IR)

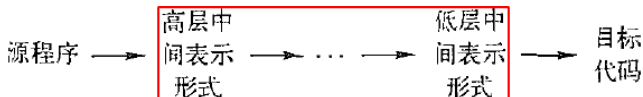
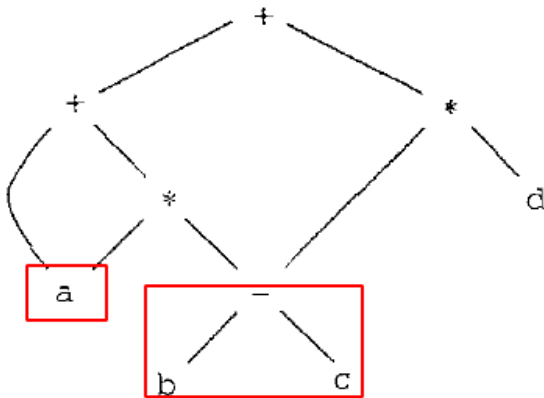


图 (抽象语法树)、**三地址代码**、C 语言

## 表达式的有向无环图

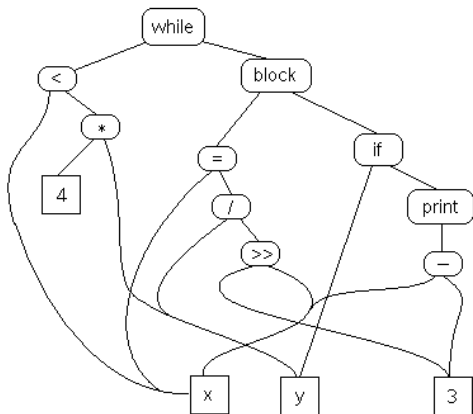


$$a + a * (b - c) + (b - c) * d$$

产生式	语义规则
1) $E \rightarrow E_1 + T$	$E.node = \text{new Node}('+', E_1.node, T.node)$
2) $E \rightarrow E_1 - T$	$E.node = \text{new Node}('-', E_1.node, T.node)$
3) $E \rightarrow T$	$E.node = T.node$
$T \rightarrow T_1 * F$	$T.node = \text{new Node}('*', T_1.node, F.node)$
4) $T \rightarrow ( E )$	$T.node = E.node$
5) $T \rightarrow \text{id}$	$T.node = \text{new Leaf}(\text{id}, \text{id.entry})$
6) $T \rightarrow \text{num}$	$T.node = \text{new Leaf}(\text{num}, \text{num.val})$

在创建节点之前, 先判断是否已存在 (哈希表)

```
while (x < 4 * y) {
    x = y / 3 >> x;
    if (y) print x - 3;
}
```



## Definition (三地址代码 (Three-Address Code (TAC; 3AC)))

每个 **TAC** 指令**最多**包含三个操作数。

$x = y \text{ op } z$  (1)

$x = \text{op } y$  (2)

$x = y$  (3)

**goto**  $L$  (4)

**if**  $x$  **goto**  $L$  (5)

**if False**  $x$  **goto**  $L$  (6)

**if**  $x$  **relop**  $y$  **goto**  $L$  (7)



Definition (三地址代码 (Three-Address Code (TAC; 3AC)))

每个 **TAC** 指令**最多**包含三个操作数。

`param x` (8)  
`call p, n` (9)  
`y = call p, n` (10)  
`return y` (11)

`param x1`  
`param x2`  
`...`

`param xn`  
`call p, n`

$p(x_1, x_2, \dots, x_n)$

Definition (三地址代码 (Three-Address Code (TAC; 3AC)))

每个 **TAC** 指令**最多**包含三个操作数。

$$x = y[i] \quad (12)$$

$$x[i] = y \quad (13)$$

$$x = \&y \quad (14)$$

$$x = *y \quad (15)$$

$$*x = y \quad (16)$$

距离位置  $y$  处  $i$  个内存单元

do i = i + 1; while (a[i] < v);

```
L:  t1 = i + 1  
    i = t1  
    t2 = i * 8  
    t3 = a [ t2 ]  
    if t3 < v goto L
```

```
100: t1 = i + 1  
101: i = t1  
102: t2 = i * 8  
103: t3 = a [ t2 ]  
104: if t3 < v goto 100
```

### 三地址代码的**四元式**表示

#### Definition (四元式 (Quadruple))

一个四元式包含四个字段, 分别为  $op$ 、 $arg_1$ 、 $arg_2$  与  $result$ 。

$$a + a * (b - c) + (b - c) * d$$

$t_1 = \text{minus } c$

$t_2 = b * t_1$

$t_3 = \text{minus } c$

$t_4 = b * t_3$

$t_5 = t_2 + t_4$

$a = t_5$

	$op$	$arg_1$	$arg_2$	$result$
0	minus	c		$t_1$
1	*	b	$t_1$	$t_2$
2	minus	c		$t_3$
3	*	b	$t_3$	$t_4$
4	+	$t_2$	$t_4$	$t_5$
5	=	$t_5$		a
		...		

$x = y[i]$	$=[$	$y$	$i$	$x$
$x[i] = y$	$[$	$i$	$y$	$x$

$x = \&y$	$=\&$	$y$	$x$
$x = *y$	$=*$	$y$	$x$
$*x = y$	$*=$	$y$	$x$

## 表达式的中间代码翻译

产生式	语义规则
$S \rightarrow id = E ;$	$S.code = E.code \parallel$ $gen(top.get(id.lexeme) '=' E.addr)$
$E \rightarrow E_1 + E_2$	$E.addr = new Temp()$ $E.code = E_1.code \parallel E_2.code \parallel$ $gen(E.addr '=' E_1.addr '+' E_2.addr)$
$  - E_1$	$E.addr = new Temp()$ $E.code = E_1.code \parallel$ $gen(E.addr '=' 'minus' E_1.addr)$
$  ( E_1 )$	$E.addr = E_1.addr$ $E.code = E_1.code$
$  id$	$E.addr = top.get(id.lexeme)$ $E.code = ''$

符号表条目

综合属性  $E.code$  与  $E.addr$

产生式	语义规则
$S \rightarrow id = E ;$	$S.code = E.code \parallel$ $gen(top.get(id.lexeme) '=' E.addr)$
$E \rightarrow E_1 + E_2$	$E.addr = new Temp()$ $E.code = E_1.code \parallel E_2.code \parallel$ $gen(E.addr '=' E_1.addr '+' E_2.addr)$
$  - E_1$	$E.addr = new Temp()$ $E.code = E_1.code \parallel$ $gen(E.addr '=' 'minus' E_1.addr)$
$  ( E_1 )$	$E.addr = E_1.addr$ $E.code = E_1.code$
$  id$	$E.addr = top.get(id.lexeme)$ 符号表条目 $E.code = ''$

$t_1 = \text{minus } c$   
 $t_2 = b + t_1$   
 $a = t_2$

$a = b + -c$

## 表达式的中间代码翻译 (增量式)

$S \rightarrow id = E ;$	$\{ gen( top.get(id.lexeme) \neq E.addr); \}$
$E \rightarrow E_1 + E_2$	$\{ E.addr = new Temp();$ $gen(E.addr \neq E_1.addr '+' E_2.addr); \}$
$  - E_1$	$\{ E.addr = new Temp();$ $gen(E.addr \neq 'minus' E_1.addr); \}$
$  ( E_1 )$	$\{ E.addr = E_1.addr; \}$
$  id$	$\{ E.addr = top.get(id.lexeme); \}$

综合属性  $E.addr$



## 数组引用的中间代码翻译

声明 :  $\text{int } a[2][3]$

数组引用 :  $x = a[1][2]; a[1][2] = x$

需要计算  $a[1][2]$  的相对于**数组基地址**  $a$  的**偏移地址**

## 数组引用的中间代码翻译

`int a[2][3]`

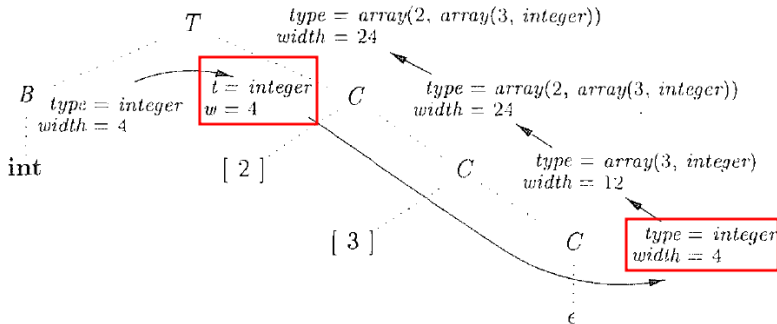


图 6-16 数组类型的语法制导翻译

## 数组类型声明

`int a[2][3]`

`array(2, array(3, integer))`

	元素类型	元素宽度
<code>a[ ]</code>	<code>array(3, integer)</code>	12
<code>a[ ][ ]</code>	<code>integer</code>	4

$$addr(a[1][2]) = base + 1 \times 12 + 2 \times 4$$

```

S → id = E ; { gen( top.get(id.lexeme) '=' E.addr); }

| L = E ; { gen(L.array.base '[' L.addr ']' '=' E.addr); }

E → E1 + E2 { E.addr = new Temp();
                gen(E.addr '=' E1.addr '+' E2.addr); }

| id { E.addr = top.get(id.lexeme); }

| L { E.addr = new Temp();
      gen(E.addr '=' L.array.base '[' L.addr ']'); }

L → id [ E ] { L.array = top.get(id.lexeme);
               L.type = L.array.type.elem;
               L.addr = new Temp();
               gen(L.addr '=' E.addr '*' L.type.width); }

| L1 [ E ] { L.array = L1.array;
               L.type = L1.type.elem;
               t = new Temp();
               L.addr = new Temp();
               gen(t '=' E.addr '*' L.type.width);
               gen(L.addr '=' L1.addr '+' t); }

```

int a[2][3]

### 综合属性 $L.array.base$ : 数组基地址 (即, 数组名)

$S \rightarrow id = E ; \quad \{ gen(top.get(id.lexeme) \neq E.addr); \}$

$\quad | \quad L = E ; \quad \{ gen(L.array.base '[' L.addr ']' \neq E.addr); \}$

$E \rightarrow E_1 + E_2 \quad \{ E.addr = new Temp();$   
 $\quad \quad \quad gen(E.addr \neq E_1.addr '+' E_2.addr); \}$

$\quad | \quad id \quad \quad \quad \{ E.addr = top.get(id.lexeme); \}$

$\quad | \quad L \quad \quad \quad \{ E.addr = new Temp();$   
 $\quad \quad \quad gen(E.addr \neq L.array.base '[' L.addr ']); \}$

### 综合属性 $L.addr$ : 偏移地址

## 综合属性 $L.array$ : 数组名 $id$ 对应的符号表条目

```
 $L \rightarrow id [ E ] \quad \{ L.array = top.get(id.lexeme);$   
     $L.type = L.array.type.elem;$   
     $L.addr = new Temp();$   
     $gen(L.addr '=' E.addr '*' L.type.width); \}$   
  
|  $L_1 [ E ] \quad \{ L.array = L_1.array;$   
     $L.type = L_1.type.elem;$   
     $t = new Temp();$   
     $L.addr = new Temp();$   
     $gen(t '=' E.addr '*' L.type.width);$   
     $gen(L.addr '=' L_1.addr '+' t); \}$ 
```

## 综合属性 $L.type$ : (当前) 元素类型

```
 $L \rightarrow id [ E ] \quad \{ L.array = top.get(id.lexeme);$   
     $L.type = L.array.type.elem;$   
     $L.addr = new Temp();$   
     $gen(L.addr '=' E.addr '*' L.type.width); \}$   
  
 $| L_1 [ E ] \quad \{ L.array = L_1.array;$   
     $L.type = L_1.type.elem;$   
     $t = new Temp();$   
     $L.addr = new Temp();$   
     $gen(t '=' E.addr '*' L.type.width);$   
     $gen(L.addr '=' L_1.addr '+' t); \}$ 
```

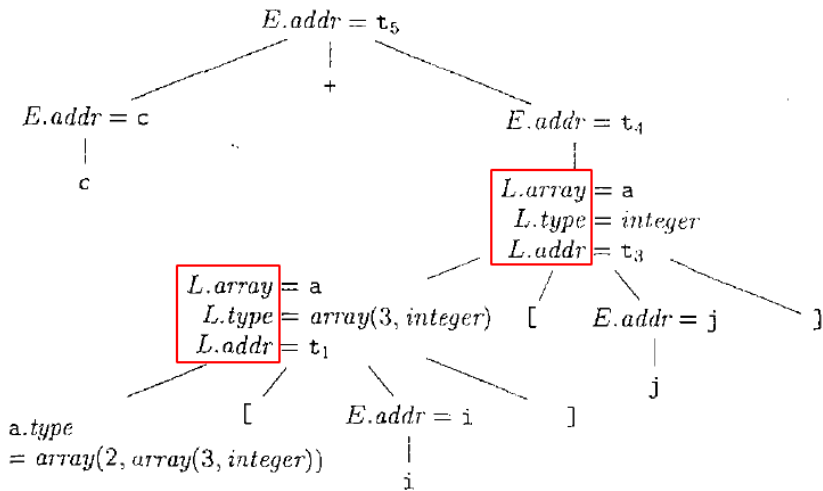
## 综合属性 $L.addr$ : (当前) 偏移地址

```
 $L \rightarrow id [ E ] \quad \{ \begin{array}{l} L.array = top.get(id.lexeme); \\ L.type = L.array.type.elem; \\ L.addr = new Temp(); \\ gen(L.addr '=' E.addr '*' L.type.width); \end{array} \}$ 
```

```
 $| \quad L_1 [ E ] \quad \{ \begin{array}{l} L.array = L_1.array; \\ L.type = L_1.type.elem; \\ t = new Temp(); \\ L.addr = new Temp(); \\ gen(t '=' E.addr '*' L.type.width); \\ gen(L.addr '=' L_1.addr '+' t); \end{array} \}$ 
```





`int a[2][3]`

```
t1 = i * 12  
t2 = j * 4  
t3 = t1 + t2  
t4 = a [ t3 ]  
t5 = c + t4
```

`int a[2][3]`

## 控制流语句的中间代码翻译

$S \rightarrow \text{if } (B) S_1$

$S \rightarrow \text{if } (B) S_1 \text{ else } S_2$

$S \rightarrow \text{while } (B) S_1$

产生式	语义规则
$P \rightarrow S$	$S.next = newlabel()$ $P.code = S.code \parallel label(S.next)$
$S \rightarrow assign$	$S.code = assign.code$
$S \rightarrow if ( B ) S_1$	$B.true = newlabel()$ $B.false = S_1.next = S.next$ $S.code = B.code \parallel label(B.true) \parallel S_1.code$
$S \rightarrow if ( B ) S_1 else S_2$	$B.true = newlabel()$ $B.false = newlabel()$ $S_1.next = S_2.next = S.next$ $S.code = B.code$ $\parallel label(B.true) \parallel S_1.code$ $\parallel gen('goto' S.next)$ $\parallel label(B.false) \parallel S_2.code$
$S \rightarrow while ( B ) S_1$	$begin = newlabel()$ $B.true = newlabel()$ $B.false = S.next$ $S_1.next = begin$ $S.code = label(begin) \parallel B.code$ $\parallel label(B.true) \parallel S_1.code$ $\parallel gen('goto' begin)$
$S \rightarrow S_1 S_2$	$S_1.next = newlabel()$ $S_2.next = S.next$ $S.code = S_1.code \parallel label(S_1.next) \parallel S_2.code$

继承属性  $S.next$  :  $S$  的下一条指令

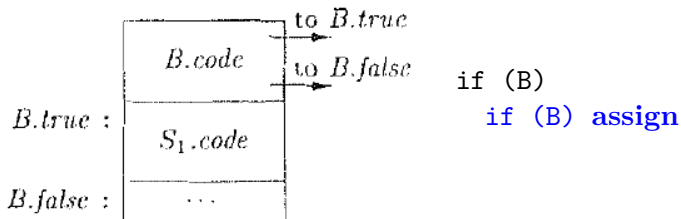
$$P \rightarrow S \quad \left| \begin{array}{l} S.next = newlabel() \\ P.code = S.code || label(S.next) \end{array} \right.$$

$S.next$  为语句  $S$  指明了“跳出”  $S$  的目标

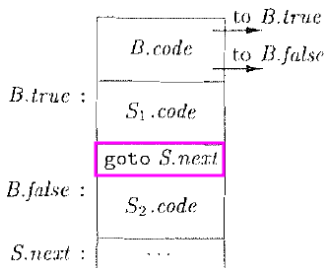
$$S \rightarrow \text{assign} \quad | \quad S.\text{code} = \text{assign}.\text{code}$$

表达式语句的翻译, 包括数组引用

$$S \rightarrow \text{if} ( B ) S_1$$

$$\left\{ \begin{array}{l} B.true = \text{newlabel}() \\ B.false = S_1.next = S.next \\ S.code = B.code || \text{label}(B.true) || S_1.code \end{array} \right.$$


$S \rightarrow \text{if} ( B ) S_1 \text{ else } S_2$	$B.true = \text{newlabel}()$ $B.false = \text{newlabel}()$ <span style="border: 1px solid blue; padding: 2px;"><math>S_1.next = S_2.next = S.next</math></span> $S.code = B.code$  $   \text{label}(B.true)    S_1.code$ $   \text{gen}('goto' S.next)$ $   \text{label}(B.false)    S_2.code$
---	---



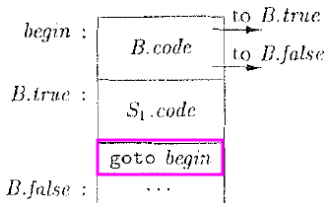
```

if (B)
    if (B) assign else assign
else
    assign
  
```



$S \rightarrow \text{while } (B) S_1$

```
begin = newlabel()
B.true = newlabel()
B.false = S.next
S1.next = begin
S.code = label(begin) || B.code
        || label(B.true) || S1.code
        || gen('goto' begin)
```



```
while (B)
    if (B) assign else assign
```

$S \rightarrow S_1 S_2$

$S_1.next = newlabel()$   
 $S_2.next = S.next$   
 $S.code = S_1.code || label(S_1.next) || S_2.code$

产生式	语义规则
$P \rightarrow S$	$S.next = newlabel()$ $P.code = S.code    label(S.next)$
$S \rightarrow assign$	$S.code = assign.code$
$S \rightarrow if ( B ) S_1$	$B.true = newlabel()$ $B.false = S_1.next = S.next$ $S.code = B.code    label(B.true)    S_1.code$
$S \rightarrow if ( B ) S_1 else S_2$	$B.true = newlabel()$ $B.false = newlabel()$ $S_1.next = S_2.next = S.next$ $S.code = B.code$ $   label(B.true)    S_1.code$ $   gen('goto' S.next)$ $   label(B.false)    S_2.code$
$S \rightarrow while ( B ) S_1$	$begin = newlabel()$ $B.true = newlabel()$ $B.false = S.next$ $S_1.next = begin$ $S.code = label(begin)    B.code$ $   label(B.true)    S_1.code$ $   gen('goto' begin)$
$S \rightarrow S_1 S_2$	$S_1.next = newlabel()$ $S_2.next = S.next$ $S.code = S_1.code    label(S_1.next)    S_2.code$

## 布尔表达式的中间代码翻译

产生式	语义规则
$B \rightarrow B_1 \    \ B_2$	$B_1.true = B.true$ $B_1.false = newlabel()$ $B_2.true = B.true$ $B_2.false = B.false$ $B.code = B_1.code \    \ label(B_1.false) \    \ B_2.code$
$B \rightarrow B_1 \ \&\& \ B_2$	$B_1.true = newlabel()$ $B_1.false = B.false$ $B_2.true = B.true$ $B_2.false = B.false$ $B.code = B_1.code \    \ label(B_1.true) \    \ B_2.code$
$B \rightarrow ! B_1$	$B_1.true = B.false$ $B_1.false = B.true$ $B.code = B_1.code$
$B \rightarrow E_1 \ rel \ E_2$	$B.code = E_1.code \    \ E_2.code$ $\quad \    \ gen('if' \ E_1.addr \ rel \ op \ E_2.addr \ 'goto' \ B.true)$ $\quad \    \ gen('goto' \ B.false)$
$B \rightarrow true$	$B.code = gen('goto' \ B.true)$
$B \rightarrow false$	$B.code = gen('goto' \ B.false)$

$B \rightarrow \text{true}$        $B.\text{code} = \text{gen}(\text{'goto' } B.\text{true})$

$B \rightarrow \text{false}$        $B.\text{code} = \text{gen}(\text{'goto' } B.\text{false})$

if (true) assign

$S \rightarrow \text{if}(B) S_1$        $\begin{cases} B.\text{true} = \text{newlabel}() \\ B.\text{false} = S_1.\text{next} = S.\text{next} \\ S.\text{code} = B.\text{code} || \text{label}(B.\text{true}) || S_1.\text{code} \end{cases}$

if (false) assign

$$B \rightarrow ! B_1$$
$$\left| \begin{array}{l} B_1.true = B.false \\ B_1.false = B.true \\ B.code = B_1.code \end{array} \right.$$

if (!true) assign

$$S \rightarrow \text{if} ( B ) S_1$$
$$\left| \begin{array}{l} B.true = \text{newlabel}() \\ B.false = S_1.next = S.next \\ S.code = B.code || \text{label}(B.true) || S_1.code \end{array} \right.$$

if (!false) assign

## 短路求值

$$B \rightarrow B_1 \parallel B_2 \quad \left| \begin{array}{l} B_1.true = B.true \\ B_1.false = newlabel() \\ B_2.true = B.true \\ B_2.false = B.false \\ B.code = B_1.code \parallel label(B_1.false) \parallel B_2.code \end{array} \right.$$

if (true || false) assign

$$S \rightarrow \text{if}(B) S_1 \quad \left| \begin{array}{l} B.true = newlabel() \\ B.false = S_1.next = S.next \\ S.code = B.code \parallel label(B.true) \parallel S_1.code \end{array} \right.$$

if (false || true) assign

## 短路求值

$$B \rightarrow B_1 \ \&\& \ B_2 \quad \left| \begin{array}{l} B_1.true = newlabel() \\ B_1.false = B.false \\ B_2.true = B.true \\ B_2.false = B.false \\ B.code = B_1.code \ || \ label(B_1.true) \ || \ B_2.code \end{array} \right.$$

if (true && false) assign

$$S \rightarrow \text{if} ( B ) S_1 \quad \left| \begin{array}{l} B.true = newlabel() \\ B.false = S_1.next = S.next \\ S.code = B.code \ || \ label(B.true) \ || \ S_1.code \end{array} \right.$$

if (false && true) assign



$$B \rightarrow E_1 \text{ rel } E_2 \quad \left| \quad \begin{array}{l} B.\text{code} = E_1.\text{code} \parallel E_2.\text{code} \\ \parallel \text{gen('if' } E_1.\text{addr rel.op } E_2.\text{addr 'goto' } B.\text{true}) \\ \parallel \text{gen('goto' } B.\text{false}) \end{array} \right.$$

```
if (x < 100 || x > 200 && x != y) x = 0;
```

```
        if x < 100 goto L2  
        goto L3  
L3:    if x > 200 goto L4  
        goto L1  
L4:    if x != y goto L2  
        goto L1  
L2:    x = 0  
L1:
```

Thank  
You!



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