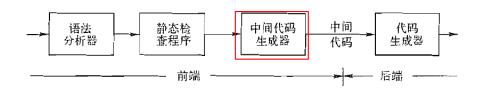
中间代码生成

魏恒峰

hfwei@nju.edu.cn

2020年12月24日





Intermediate Representation (IR)



Intermediate Representation (IR)



精确:不能丢失源程序的信息

独立: 不依赖特定的源语言与目标语言

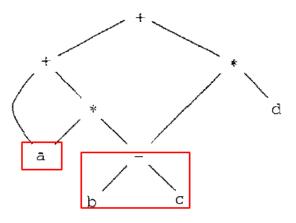
(如,没有复杂的寻址方式)

Intermediate Representation (IR)



图 (抽象语法树)、三地址代码、C 语言

表达式的有向无环图



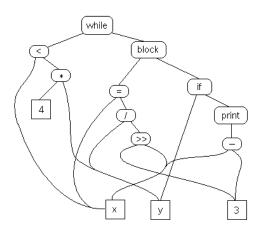
$$a + a * (b - c) + (b - c) * d$$

产生式		语义规则	
1)	$E ightarrow E_1 + T$	$E.node = \frac{\text{new Node}('+', E_1.node, T.node)}{\text{Node}('+', E_1.node, T.node)}$	
2)	$E o E_1 - T$	$E.node = \frac{\text{new }Node('-', E_1.node, T.node)}{}$	
3)	$E \to T$	E.node = T.node	
	$T \rightarrow T_1^*F$	$T.node = $ new $Node('*', T_1.node, F.node)$	
4)	T ightarrow (E)	T.node = E.node	
5)	$T o \mathbf{id}$	$T.node = \frac{\text{new}}{\text{leaf}(\text{id}, \text{id}.entry)}$	
6)	$T \rightarrow \mathbf{num}$	T.node = new $Leaf(num, num.val)$	

产生式		语义规则	
1)	$E ightarrow \overline{E_1 + T}$	$E.node = \frac{\text{new Node}('+', E_1.node, T.node)}{}$	
2)	$E \rightarrow E_1 - T$	$E.node = $ $new Node('-', E_1.node, T.node)$	
3)	$E \to T$	E.node = T.node	
ļ	$T \rightarrow T_1 *F$	$T.node = $ new $Node('*', T_1.node, F.node)$	
4)	T ightarrow (E)	T.node = E.node	
5)	$T o \mathrm{id}$	$T.node = \frac{\mathbf{new}}{\mathbf{new}} Leaf(\mathbf{id}, \mathbf{id}.entry)$	
6)	$T \rightarrow \text{num}$	T.node = new $Leaf(num, num.val)$	

在创建节点之前, 先判断是否已存在 (哈希表)

```
while (x < 4 * y) {
    x = y / 3 >> x;
    if (y) print x - 3;
}
```



Definition (三地址代码 (Three-Address Code (TAC; 3AC))) 每个 **TAC** 指令最多包含三个操作数。

$$x = y \mathbf{op} z \tag{1}$$

$$x = \mathbf{op} \ y \tag{2}$$

$$x = y \tag{3}$$

Definition (三地址代码 (Three-Address Code (TAC; 3AC)))

每个 TAC 指令最多包含三个操作数。

$$x = y \mathbf{op} z \tag{1}$$

 $x = \mathbf{op} \ y \tag{2}$

$$x = y \tag{3}$$

goto L (4)

if x goto L (5)

if False x goto L (6)

 $\textbf{if} \ x \ \textbf{relop} \ y \ \textbf{goto} \ L \qquad (7) \\$

Definition (三地址代码 (Three-Address Code (TAC; 3AC)))

每个 TAC 指令最多包含三个操作数。

		$\mathtt{param}\ x_\mathtt{l}$
		$\mathtt{param}\ x_2$
$\mathbf{param}\;x$	(8)	
$\mathbf{call}\; p, n$	(9)	param x_n
$y = \mathbf{call}\ p, n$	(10)	call p, n
$\mathbf{return}\ y$	(11)	0-11 p,
		$p(x_1, x_2, \dots, x_n)$

Definition (三地址代码 (Three-Address Code (TAC; 3AC))) 每个 **TAC** 指令最多包含三个操作数。

$$x = y[i] \tag{12}$$

$$x[i] = y \tag{13}$$

距离位置 y 处 i 个内存单元

Definition (三地址代码 (Three-Address Code (TAC; 3AC)))

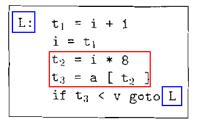
每个 TAC 指令最多包含三个操作数。

$$x = y[i] (12) x = &y (14)$$

$$x[i] = y (13) x = *y$$

距离位置 y 处 i 个内存单元 *x = y (16)

(15)



L:
$$t_1 = i + 1$$

 $i = t_1$
 $t_2 = i * 8$
 $t_3 = a [t_2]$
if $t_3 < v$ goto L

```
100: t_1 = i + 1

101: i = t_1

102: t_2 = i * 8

103: t_3 = a [t_2]

104: if t_3 < v goto 100
```

三地址代码的四元式表示

Definition (四元式 (Quadruple))

一个四元式包含四个字段, 分别为 op、 arg_1 、 arg_2 与 result。

三地址代码的四元式表示

Definition (四元式 (Quadruple))

一个四元式包含四个字段, 分别为 op、 arg_1 、 arg_2 与 result。

$$a + a * (b - c) + (b - c) * d$$

$$t_1 = minus c$$
 $t_2 = b * t_1$
 $t_3 = minus c$
 $t_4 = b * t_3$
 $t_5 = t_2 + t_4$
 $a = t_5$

	о́р	arg ₁	arg_2	result
0	minus	С	ı	tı
1	*	Ъ	t ₁	t_2
2	minus	C	[t ₃
3	*	b	t ₃	t4
4	+	t_2	t4	t ₅
5	=	t_5	1	, a

$$x = y[i]$$
$$x[i] = y$$

$$= [] \qquad y \qquad i \qquad x$$
$$[]= \qquad i \qquad y \qquad x$$

$$x = y[i]$$
$$x[i] = y$$

$$= [] \qquad y \qquad i \qquad x$$
$$[] = \qquad i \qquad y \qquad x$$

$$x = \&y$$
$$x = *y$$
$$*x = y$$

$$x = y[i]$$
$$x[i] = y$$

$$= [] \qquad y \qquad i \qquad x$$

$$[] = \qquad i \qquad y \qquad x$$

$$x = \&y$$
$$x = *y$$
$$*x = y$$

$$=& y & x \\ =* & y & x \\ *= & y & x$$

13 / 61

表达式的中间代码翻译

产生式	语义规则
$S \rightarrow id = E$;	S.code = E.code $gen(top.get(id.lexeme))' = 'E.addr)$
$E \rightarrow E_1 + E_2$	$ E.addr = \mathbf{new} \ Temp() \\ E.code = E_1.code \mid\mid E_2.code \mid\mid \\ gen(E.addr'='E_1.addr'+'E_2.addr) $
- E _i	$E.addr = \mathbf{new} \ Temp() \ E.code = E_1.code \mid \ gen(E.addr'=' 'minus' \ E_1.addr)$
[(E ₁)	$E.addr = E_1.addr \ E.code = E_1.code$
id	E.addr = top.get(id.lexeme) 符号表条目 E.code = ''

综合属性 E.code 与 E.addr

产生式	语义规则	
$S \rightarrow id = E$;	$S.code = E.code \mid gen(top.get(id.lexeme))' = 'E.addr)$	
$E \rightarrow E_1 + E_2$	$ \begin{split} E.addr &= \mathbf{new} \ Temp () \\ E.code &= E_1.code \mid \mid E_2.code \mid \mid \\ gen(E.addr'=' E_1.addr'+' E_2.addr) \end{split} $	
- E _i	$E.addr = \mathbf{new} \ Temp()$ $E.code = E_1.code \mid gen(E.addr'=' 'minus' E_1.addr)$	
[(E ₁)	$E.addr = E_1.addr$ $E.code = E_1.code$	
id	E.addr = top.get(id.lexeme) 符号表条目 E.code = ''	

$$t_1 = minus c$$

 $t_2 = b + t_1$
 $a = t_2$

$$a = b + -c$$

表达式的中间代码翻译 (增量式)

$$S \rightarrow id = E$$
; { $gen(top.get(id.lexeme) '=' E.addr)$; }
 $E \rightarrow E_1 + E_2$ { $E.addr = new Temp()$; $gen(E.addr '=' E_1.addr '+' E_2.addr)$; }

| $-E_1$ { $E.addr = new Temp()$; $gen(E.addr '=' minus' E_1.addr)$; }

| (E_1) { $E.addr = E_1.addr$; }

| id { $E.addr = top.get(id.lexeme)$; }

综合属性 E.addr

数组引用的中间代码翻译

声明: int a[2][3]

数组引用: x = a[1][2]; a[1][2] = x

数组引用的中间代码翻译

声明: int a[2][3]

数组引用: x = a[1][2]; a[1][2] = x

需要计算 a[1][2] 的相对于**数组基地址** a 的**偏移地址**

数组引用的中间代码翻译

int a[2][3]

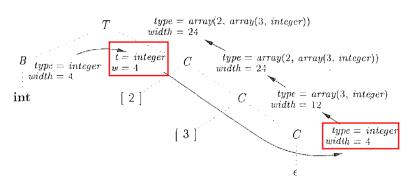


图 6-16 数组类型的语法制导翻译

数组类型声明

int a[2][3]

array(2, array(3, integer))

	类型	宽度
a	array(2, array(3, integer))	24
a[i]	array(3, integer)	12
a[i][j]	integer	4

int a[2][3]

array(2, array(3, integer))

	类型	宽度
a	array(2, array(3, integer))	24
a[i]	array(3, integer)	12
a[i][j]	integer	4

$$addr(a[1][2]) = base + 1 \times 12 + 2 \times 4$$

```
S \rightarrow id = E; { gen(top.get(id.lexeme)' = 'E.addr); }
      L = E:
                  \{ gen(L.array.base' ['L.addr']' '='E.addr); \}
E \rightarrow E_1 + E_2 + E_3 { E.addr = new Temp();
                    gen(E.addr'='E_1.addr'+'E_2.addr);
      id
                 \{E.addr = top.get(id.lexeme);\}
    \mid L \mid
                  \{E.addr = new\ Temp();
                    gen(E.addr'=' L.array.base'[' L.addr']'); }
                  { L.array = top.qet(id.lexeme):
                    L.type = L.array.type.elem;
                    L.addr = new Temp();
                    qen(L.addr'='E.addr'*'L.type.width);
                   \{L.array = L_1.array;
                    L.type = L_1.type.elem;
                    t = new Temp():
                    L.addr = new Temp();
                    qen(t'='E.addr'*'L.type.width);
                    qen(L.addr'='L_1.addr'+'t);
```

综合属性 L.array.base: 数组基地址 (即,数组名)

```
S \rightarrow id = E; { gen(top.get(id.lexeme)' = 'E.addr); }
    | L = E ; { gen(L.array.base' ['L.addr']' '='E.addr);
E \rightarrow E_1 + E_2 + E_2 { E.addr = new Temp();
                      gen(E,addr'='E_1,addr'+'E_2,addr);
       id
                    \{E.addr = top.get(id.lexeme);\}
                    { E.addr = new \ Temp();

gen(E.addr'=' \ L.array.base'[' \ L.addr']'); }
    L
```

综合属性 L.addr: 偏移地址

综合属性 L.array: 数组名 id对应的符号表条目

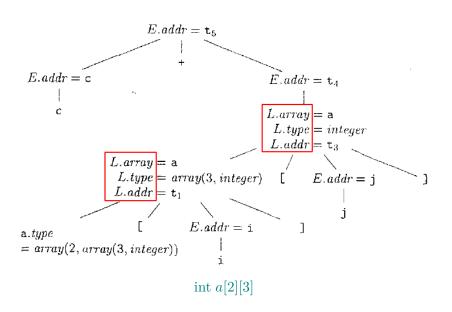
```
L \rightarrow id [E] \{L.array = top.get(id.lexeme);
                    L.type = L.array.type.elem;
                    L.addr = new Temp();
                    qen(L.addr'='E.addr'*'L.type.width); }
   L_1 \ [E] \ \{L.array = L_1.array;
                    L.type = L_1.type.elem:
                    t = \mathbf{new} \ Temp();
                    L.addr = new Temp();
                    qen(t'='E.addr'*'L.type.width);
                    qen(L.addr'='L_1.addr'+'t);
```

综合属性 L.type: (当前) 元素类型

```
L \rightarrow id [E] \{L.array = top.get(id.lexeme);
                    L.type = L.array.type.elem;
                    L.addr = \mathbf{new} \ Temp();
                    qen(L.addr'='E.addr'*'L.type.width); \}
    L_1 [E] \{L.array = L_1.array;
                    L.type = L_1.type, elem;
                    t = \mathbf{new} \ Temp():
                    L.addr = new Temp():
                    gen(t'='E.addr'*'L.type.width);
                    gen(L.addr'='L_1.addr'+'t);
```

综合属性 L.addr: (当前) 偏移地址

```
L \rightarrow id [E] \{L.array = top.get(id.lexeme)\}
                   L.type = L.array.type.elem;
                  L.addr = new Temp();
                   gen(L.addr'='E.addr'*'L.type.width);
   L_1 [E] \{L.array = L_1.array;
                   L.type = L_1.type.elem;
                   t = new Temp();
                  L.addr = new Temp();
                   gen(t'='E.addr'*'L.type.width);
                  gen(L.addr'='L_1.addr'+'t);
```



$$t_1 = i * 12$$
 $t_2 = j * 4$
 $t_3 = t_1 + t_2$
 $t_4 = a [t_3]$
 $t_5 = c + t_4$

int a[2][3]

控制流语句与布尔表达式的中间代码翻译

$$S \rightarrow \text{ if } (B) S_1$$

 $S \rightarrow \text{ if } (B) S_1 \text{ else } S_2$
 $S \rightarrow \text{ while } (B) S_1$

控制流语句与布尔表达式的中间代码翻译



产生式	语义规则
$P \rightarrow S$	S.next = newlabel() $P.code = S.code \mid label(S.next)$
$S \rightarrow assign$	S.code = assign.code
$S \rightarrow \mathbf{if}(B) S_1$	$ \begin{array}{ll} B.true &= newlabel() \\ B.false &= S_1.next = S.next \\ S.code &= B.code \mid\mid label(B.true) \mid\mid S_1.code \end{array} $
$S \rightarrow \text{if } (B) S_1 \text{ else } S_2$	$B.true = newlabel() \\ B.false = newlabel() \\ S_1.next = S_2.next = S.next \\ S.code = B.code \\ label(B.true) S_1.code \\ gen('goto' S.next) \\ label(B.false) S_2.code$
$S \rightarrow \text{ while } (B) S_1$	$begin = newlabel() \\ B.true = newlabel() \\ B.false = S.next \\ S_1.next = begin \\ S.code = label(begin) B.code \\ label(B.true) S_1.code \\ gen('goto' begin)$
$S \rightarrow S_1 S_2$	$ \begin{array}{ll} S_1.next &= newlabel() \\ S_2.next &= S.next \\ S.code &= S_1.code \mid\mid label(S_1.next) \mid\mid S_2.code \end{array} $

继承属性 S.next: S 的下一条指令

$$P \rightarrow S$$
 $S.next = newlabel()$ $P.code = S.code || label(S.next)$

S.next 为语句 S 指明了"跳出"S 的目标

 $S \rightarrow assign$

S.code = assign.code

代表了表达式的翻译,包括数组引用

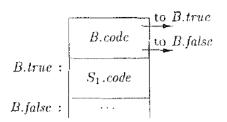
$$S \rightarrow \mathbf{if} (B) S_1$$

$$\begin{array}{lll} B.true &=& newlabel() \\ B.false &=& S_1.next \\ S.code &=& B.code \mid \mid label(B.true) \mid \mid S_1.code \end{array}$$

$$S \rightarrow \mathbf{if} (B) S_1$$

$$B.true = \underbrace{newlabel()}_{B.false} = \underbrace{S_1.next}_{S.code} = S.next$$

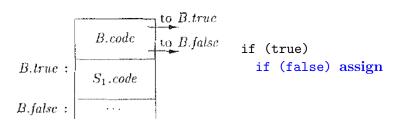
$$S.code = \underbrace{B.code}_{||} || \underbrace{label(B.true)}_{||} || S_1.code$$



$$S \rightarrow \mathbf{if} (B) S_1$$

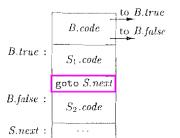
$$B.true = \underbrace{newlabel()}_{B.false} = \underbrace{S_1.next}_{S.code} = S.next$$

$$S.code = \underbrace{B.code}_{||} || \underbrace{label(B.true)}_{||} || S_1.code$$

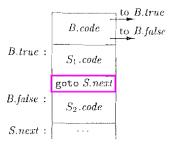


$$S o ext{if } (B) S_1 ext{ else } S_2 \ egin{array}{ll} B. true &= newlabel() \ B. false &= newlabel() \ \hline S_1.next &= S_2.next &= S.next \ S.code &= B.code \ & || label(B.true) || S_1.code \ & || gen('goto' S.next) \ & || label(B.false) || S_2.code \ \hline \end{array}$$

```
S 	o 	ext{if } (B) S_1 	ext{ else } S_2 \ egin{array}{ll} B. true &= newlabel() \ B. false &= newlabel() \ \hline S_1.next &= S_2.next &= S.next \ S.code &= B.code \ &\parallel label(B.true) \parallel S_1.code \ &\parallel gen('goto' \ S.next) \ &\parallel label(B.false) \parallel S_2.code \ \end{array}
```



```
S 	o 	ext{if } (B) S_1 	ext{ else } S_2 \ | egin{array}{ll} B. true &= newlabel() \ B. false &= newlabel() \ \hline S_1.next &= S_2.next &= S.next \ S.code &= B.code \ & || label(B.true) || S_1.code \ & || gen('goto' S.next) \ & || label(B.false) || S_2.code \ \hline \end{array}
```



if (true) if (true) assign else assign else assign

```
S \rightarrow  while (B) S_1
```

```
begin = newlabel() \\ B.true = newlabel() \\ B.false = S.next \\ \hline S_1.next = begin \\ S.code = label(begin) || B.code \\ || label(B.true) || S_1.code \\ || gen('goto' begin)
```

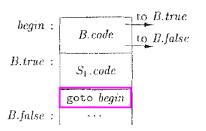
```
S \rightarrow  while (B) S_1
```

```
\begin{array}{ll} begin = newlabel() \\ B.true = newlabel() \\ B.false = S.next \\ \hline S_1.next = begin \\ S.code = label(begin) \mid\mid B.code \\ \mid\mid label(B.true) \mid\mid S_1.code \\ \mid\mid gen('goto'\ begin) \end{array}
```



```
S \rightarrow  while ( B ) S_1
```

```
begin = newlabel()
B.true = newlabel()
B.false = S.next
S_1.next = begin
S.code = label(begin) || B.code
|| label(B.true) || S_1.code
|| gen('goto' begin)
```



while (true)
if (false) assign else assign

$$S \rightarrow S_1 S_2$$

$$S \rightarrow S_1 S_2$$

if (true) assign else assign assign

产生式	语义规则
$P \rightarrow S$	
$S \rightarrow assign$	S.code = assign.code
$S \rightarrow \text{if } (B) S_1$	$\begin{array}{ll} B.true &= newlabel() \\ B.false &= \left \frac{S_1.next}{S_1.next} \right = S.next \\ S.code &= B.code \mid\mid label(B.true) \mid\mid S_1.code \end{array}$
$S \rightarrow \mathbf{if} (B) S_1 $ else S_2	$B.true = newlabel() \\ B.false = newlabel() \\ [S_1.next = S_2.next] = S.next \\ S.code = B.code \\ label(B.true) S_1.code \\ gen('goto' S.next) \\ label(B.false) S_2.code$
$S \rightarrow \text{ while } (B) S_1$	begin = newlabel()
	$B.true = newlabel()$ $B.false = S.next$ $\boxed{S_1.next} = begin$ $S.code = label(begin) B.code$ $ label(B.true) S_1.code$ $ gen('goto' begin)$
$S \rightarrow S_1 S_2$	

布尔表达式的中间代码翻译

产生式	语义规则
$B \rightarrow B_1 \mid \mid B_2$	$B_1.true = B.true$ $B_1.false = newlabel()$
	$B_2.true = B.true$ $B_2.false = B.false$ $B.code = B_1.code \mid\mid label(B_1.false) \mid\mid B_2.code$
$B \rightarrow B_1 \&\& B_2$	$B_1.true = newlabel()$ $B_1.false = B.false$ $B_2.true = B.true$ $B_2.false = B.false$ $B.code = B_1.code \mid\mid label(B_1.true) \mid\mid B_2.code$
$B \rightarrow ! B_1$	$B_1.true = B.false$ $B_1.false = B.true$ $B.code = B_1.code$
$B ightarrow E_1 \ ext{rel} \ E_2$	$B.code = E_1.code \mid\mid E_2.code$ $\mid\mid gen('if' E_1.addr rel.op E_2.addr 'goto' B.true)$ $\mid\mid gen('goto' B.false)$
$B ightarrow { m true}$	B.code = gen('goto' B.true)
$B \rightarrow \text{false}$	B.code = gen('goto' B.false)

$$B \rightarrow \text{true}$$

$$B.code = gen('goto' B.true)$$

$$B \rightarrow \mathbf{false}$$

$$B.code = gen('goto' B.false)$$

$$B \rightarrow \text{true}$$

$$B \rightarrow \text{false}$$

$$B.code = gen('goto' B.true)$$

$$B.code = gen('goto' B.false)$$

if (true) assign

$$S \rightarrow \mathbf{if} (B) S_1$$

$$\begin{array}{ll} B.true &= newlabel() \\ B.false &= S_1.next \\ S.code &= B.code \mid\mid label(B.true) \mid\mid S_1.code \end{array}$$

$$S.code = B.code || label(B.true) || S_1.code$$

if (false) assign

$$B \rightarrow ! B_1$$

$$B_1.true = B.false$$

 $B_1.false = B.true$
 $B.code = B_1.code$

$$B \rightarrow ! B_1$$

$$B_1.true = B.false$$

 $B_1.false = B.true$
 $B.code = B_1.code$

if (!true) assign

$$S \rightarrow \mathbf{if} (B) S_1$$

if (!false) assign

$$B \rightarrow B_1 \mid \mid B_2$$

$$B \rightarrow B_1 \mid \mid B_2$$

$$\begin{vmatrix} B_1.true \\ B_1.false = newlabel() \\ B_2.true \\ B_2.false = B.true \\ B_2.false = B.false \\ B.code = B_1.code \mid \mid label(B_1.false) \mid \mid B_2.code \end{vmatrix}$$

$$B \rightarrow B_1 \mid \mid B_2$$

$$B_1.true = B.true$$

$$B_1.false = newlabel()$$

$$B_2.true = B.true$$

$$B_2.false = B.false$$

$$B.code = B_1.code \mid \mid label(B_1.false) \mid \mid B_2.code$$

if (true || false) assign

$$S \rightarrow \mathbf{if} (B) S_1$$

$$\begin{array}{lll} B.true &=& newlabel() \\ B.false &=& S_1.next \\ S.code &=& B.code \mid\mid label(B.true) \mid\mid S_1.code \end{array}$$

if (false || true) assign



if (true && false) assign

$$S \rightarrow \mathbf{if} (B) S_1$$

$$B.true = \underbrace{newlabel()}_{B.false} = \underbrace{S_1.next}_{S.code} = S.next$$

$$S.code = B.code || label(B.true) || S_1.code$$

if (false && true) assign



 $B \rightarrow E_1 \text{ rel } E_2$ $B.code = E_1.code \mid\mid E_2.code$ $\mid\mid gen('if' E_1.addr \text{ rel.} op E_2.addr 'goto' B.true)$ $\mid\mid gen('goto' B.false)$

if
$$(x < 100 \mid | x > 200 \&\& x != y) x = 0;$$

```
if x < 100 goto L_2
      goto {\tt L}_3
L_3: if x > 200 goto L_4 goto L_1
L_4: if x != y goto L_2
goto L_1
L_2: x = 0
```

布尔表达式的作用: 布尔值 vs. 控制流跳转

$$S \rightarrow \text{id} = E$$
; | if $(E) S$ | while $(E) S \mid S$ | $E \rightarrow E \parallel E \mid E \& \& E \mid E \text{ rel } E \mid E + E \mid (E)$ | id | true | false

布尔表达式的作用: 布尔值 vs. 控制流跳转

$$S \rightarrow \text{id} = E$$
; | if $(E) S$ | while $(E) S \mid S$ | $E \rightarrow E \parallel E \mid E \& \& E \mid E \text{ rel } E \mid E + E \mid (E)$ | id | true | false

函数 jump(t, f): 生成控制流代码

函数 rvalue(): 生成计算布尔值的代码,并将结果存储在临时变量中

产生式	语义规则
$S \rightarrow id = E$;	S.code = E.code $gen(top.get(id.lexeme))' = 'E.addr)$
$E \rightarrow E_1 + E_2$	$E.addr = new Temp()$ $E.code = E_1.code E_2.code $ $gen(E.addr'=' E_1.addr'+' E_2.addr)$
$-E_i$	$E.addr = new \ Temp()$ $E.code = E_1.code \parallel gen(E.addr'=''minus' \ E_1.addr)$
[(E ₁)	$E.addr = E_1.addr$ $E.code = E_1.code$
id	E.addr = top.get(id.lexeme) 符号表条目 E.code = ''

 $E \rightarrow E_1 \&\& E_2$

为 E 生成**跳转代码**, 在**真假出口处**将 true 或 false 存储到临时变量

x = a < b && c < d

```
ifFalse a < b goto L<sub>1</sub>
    ifFalse c < d goto L<sub>1</sub>
    t = true
    goto L<sub>2</sub>
L<sub>1</sub>: t = false
L<sub>2</sub>: x = t
```

$$S \rightarrow \mathbf{if} (B) S_1$$

```
B.true = newlabel()

B.false = S_1.next = S.next

S.code = B.code \mid\mid label(B.true) \mid\mid S_1.code
```

B 还不知道 S.next 的指令地址, 如何跳转?

$$S \rightarrow \mathbf{if} (B) S_1$$

$$\begin{array}{lll} B.true &=& newlabel() \\ B.false &=& S_1.next &=& S.next \\ S.code &=& B.code \mid\mid label(B.true) \mid\mid S_1.code \end{array}$$

B 还不知道 S.next 的指令地址, 如何跳转?

再扫描一遍中间代码,将标号替换成指令(相对)地址

$$S \rightarrow \mathbf{if} (B) S_1$$

$$\begin{array}{lll} B.true &=& newlabel() \\ B.false &=& S_1.next &=& S.next \\ S.code &=& B.code \mid\mid label(B.true) \mid\mid S_1.code \end{array}$$

B 还不知道 S.next 的指令地址, 如何跳转?

再扫描一遍中间代码,将标号替换成指令(相对)地址

可否在生成中间代码的时候就填入指令地址?

回填 (Backpatch) 技术

回填 (Backpatch) 技术



子节点挖坑、祖先节点填坑

针对布尔表达式的回填技术

```
1) B \rightarrow B_1 \parallel M B_2
                               { backpatch(B_1,falselist,M.instr);
                                  B.truelist = merge(B_1.truelist, B_2.truelist):
                                  B.falselist = B_2.falselist; 
     B \rightarrow B_1 \&\& M B_2
                                  backpatch(B_1.truelist.M.instr):
                                  B.truelist = B_{\uparrow}.truelist;
                                  B.falselist = merge(B_1.falselist, B_2.falselist); 
                               \{B.truelist = B_1.falselist;
                                  B.falselist = B_1.truelist;
     B \rightarrow (B_{\perp})
                               \{B.truelist = B_1.truelist;
                                  B.falselist = B_1.falselist; 
    B \to E_1 \text{ rel } E_2 { B.truelist = makelist(nextinstr):
                                  B.falsclist = makelist(nextinstr + 1);
                                  gen('if' E1.addr rel.op E2.addr 'goto _'):
                                  gen('goto _'): }
     B \to \mathbf{true}
                               \{ B.truelist = makelist(nextinstr); \}
                                 gen('goto _'); }
     B \to \mathbf{false}
                               \{ B.falselist = makelist(nextinstr): \}
                                 gen('goto _'); }
                               \{ M.instr = nextinstr. \}
```

2020 年 12 月 24 日

综合属性 B.truelist 保存 需要跳转到 B.true 的指令地址

6) $B \rightarrow \text{true}$ { $B.truelist = makelist(nextinstr); \\ <math>gen('goto \ _');$ }

7) $B \rightarrow \text{false}$ { $B.falselist = makelist \ nextinstr); \\ <math>gen('goto \ _');$ }

综合属性 B.falselist 保存 需要跳转到 B.false 的指令地址

综合属性 B.truelist 保存 需要跳转到 B.true 的指令地址

6)
$$B \rightarrow \text{true}$$
 { $B.truelist = makelist(nextinstr); } gen('goto _'); }$

7)
$$B \rightarrow \text{false}$$
 { $B.\text{falselist} = makelist next instr); } gen('goto _'); }$

综合属性 B.falselist 保存 需要跳转到 B.false 的指令地址

$$B o ext{true}$$
 $B.code = gen('goto' B.true)$ $B o ext{false}$ $B.code = gen('goto' B.false)$

50 / 61

```
5) B \to E_1 \text{ rel } E_2
```

```
B.truelist = makelist(nextinstr):
B.falselist = makelist(nextinstr + 1);
gen('if' E<sub>1</sub>.addr rel.op E<sub>2</sub>.addr 'goto _'):
gen('goto _'): }
```

$$B \rightarrow E_1 \text{ rel } E_2$$
 | $B.code = E_1.code \parallel E_2.code$ | $||gen('if' E_1.addr \text{ rel.op } E_2.addr 'goto' B.true$ | $||gen('goto' B.false)$

$$3) \quad B \to 1 B_1$$

4)
$$B \rightarrow (B_1)$$

$$B \rightarrow ! B_1$$

{
$$B.truelist = B_1.falselist;$$

 $B.falselist = B_1.truelist;$ }
{ $B.truelist = B_1.truelist;$
 $B.falselist = B_1.falselist;$ }

$$B_1.true = B.false$$

 $B_1.false = B.true$
 $B.code = B_1.code$

```
2) B \rightarrow B_1 \&\& M B_2 = \{\begin{array}{c} backpatch(B_1.truelist, M.instr); \\ B.truelist = B_2.truelist; \\ B.falselist = merge(B_1.falselist, B_2.falselist); \} \end{array}
```

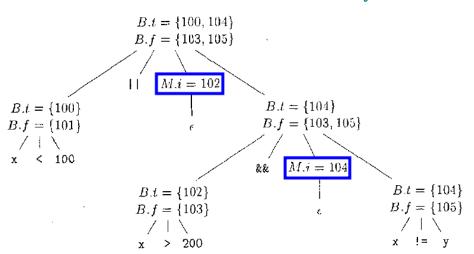
8)
$$M \to \epsilon$$
 { $M.instr = nextinstr.$ }

$$B \rightarrow B_1 \&\& B_2$$
 $| B_1.true = newlabel() | B_1.false = B.false | B_2.true = B.true | B_2.false | B.false | B.false | B.code = B_1.code || label(B_1.true) || B_2.code || B_2.code || B_2.true || B_2.code || B_2.true || B_2.code || B_2.true || B_2$

1)
$$B \rightarrow B_1 \parallel M B_2 = \{ \begin{array}{ll} backpatch(B_1, falselist, M.instr); \\ B.truelist = merge(B_1, truelist, B_2, truelist); \\ B.falselist = B_2, falselist; \} \end{array}$$

8)
$$M \to \epsilon$$
 { $M.instr = nextinstr$, }

$x < 100 \mid \mid x > 200 \&\& x != y$



```
100: if x < 100 goto _
101: goto _
102: if x > 200 goto 104
103: goto _
104: if x != y goto _
105: goto _
```

a) 将 104 回填到指令 102 中之后

```
100: if x < 100 goto _

101: goto 102

102: if x > 200 goto 104

103: goto _

104: if x != y goto _

105: goto _
```

b) 将 102 回填到指令 101 中之后

$$S \rightarrow \text{if } (B) S \mid \text{if } (B) S \text{ else } S \mid \text{ while } (B) S \mid \{L\} \mid A;$$

 $L \rightarrow L S \mid S$

1)
$$S \to if(B) M S_1 \{ backpatch(B.truelist, M.instr); \\ S.nextlist = merge(B.falselist, S_1.nextlist); \}$$

6)
$$M \to \epsilon$$
 { $M.instr = nextinstr$, }

1)
$$S \to if(B) M S_1 \{ backpatch B.truelist, M.instr);$$

 $S.nextlist = merge(B.falselist, S_1.nextlist); \}$

6)
$$M \to \epsilon$$
 { $M.instr = nextinstr$, }

$$S \rightarrow \mathbf{if}(B) S_1$$

$$B.true = \underbrace{newlabel()}_{B.false} = \underbrace{S_1.next}_{S.code} = S.next$$

$$S.code = B.code || label(B.true) || S_1.code$$

```
S \rightarrow \mathbf{if}(B) M_1 S_1 N \mathbf{else} M_2 S_2
\{ \begin{array}{l} backpatch \ B.truelist, \ M_1.instr); \\ backpatch \ B.falselist, \ M_2.instr); \\ temp = merge(S_1.nextlist, \ N.nextlist); \\ S.nextlist = merge(temp, \ S_2.nextlist); \} \end{array}
```

```
    6) M → ε { M.instr = nextinstr, }
    7) N → ε { N.nextlist = makelist(nextinstr); gen('goto _'); }
```

```
S \rightarrow \mathbf{if}(B) M_1 S_1 N \text{ else } M_2 S_2
                                     { backpatch B.truelist, M_1.instr); backpatch B.falselist, M_2.instr);}
                                         \underline{temp} = \underline{merge}(S_1.nextlist, N.nextlist);
                                        S.nextlist = merge(temp, S_2.nextlist); 
            6) M \to \epsilon
                                                  \{ M.instr = nextinstr, \}
            7) N \to \epsilon
                                                  \{ N.nextlist = makelist(nextinstr); \}
                                                     gen('goto _'); }
                 S \rightarrow \text{if } (B) S_1 \text{ else } S_2
B.true = newlabel()
B.false = newlabel()
S_1.next = S_2.next = S.next
S.code = B.code
|| label(B.true) || S_1.code
|| gen('goto' S.next) || S_1.code
                                                                      || label(B.false) || S_2.code
```

```
3) S \rightarrow \text{ while } M_1 \ (B) \ M_2 S_1 \ \{ \begin{array}{ll} backpatch \ (S_1.nextlist, \ M_1.instr); \\ backpatch \ B.truelist, \ M_2.instr); \\ S.nextlist \ = B.falselist; \\ gen('goto' \ M_1.instr); \ \} \end{array}
```

6)
$$M \rightarrow \epsilon$$
 { $M.instr = nextinstr$, }

```
3) S \rightarrow \text{ while } M_1 (B) M_2 S_1
                                                                                   \{\begin{array}{ll} backpatch & S_1.nextlist, & M_1.instr); \\ backpatch & B.truelist, & M_2.instr); \\ S.nextlist & = & B.falselist; \\ gen('goto' & M_1.instr); & \} \end{array}
                                                                                                    \{ M.instr = nextinstr, \}
           6) M \to \epsilon
                                                                                         \begin{array}{lll} begin &= newlabel() \\ B.truc &= newlabel() \\ B.false &= S.next \\ \hline S_1.next &= begin \\ S.code &= label(begin) \mid\mid B.code \\ &\mid\mid label(B.true) \mid\mid S_1.code \\ &\mid\mid gen('goto'\ begin) \end{array}
                     S \rightarrow  while (B) S_1
```

60/61

4)
$$S \rightarrow \{L\}$$

$$\{ S.nextlist = L.nextlist; \}$$

5)
$$S \rightarrow A$$
;

$$\{S.nextlist = null;\}$$

6)
$$M \to \epsilon$$

$$\{ M.instr = nextinstr, \}$$

8)
$$L \rightarrow L_1 M S$$

{
$$backpatch(L_1.nextlist, M.instr);$$

 $L.nextlist = S.nextlist;$ }

9)
$$L \rightarrow S$$

$$\{L.nextlist = S.nextlist;\}$$

Thank You!



Office 926 hfwei@nju.edu.cn