中间代码生成 (3. Switch/Case 语句与过程调用的翻译)

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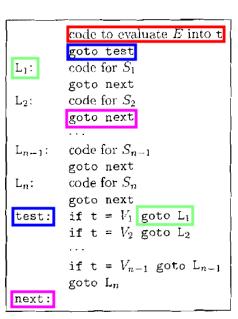
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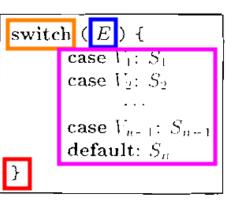
2021年12月24日



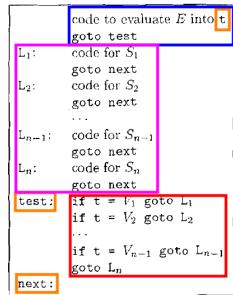
```
switch ( E ) {
    case V_1: S_1
    case V_2: S_2
    ...
    case V_{n-1}: S_{n-1}
    default: S_n
}
```

非 C 语言语义 (break)





 $V_i: L_i$ 队列



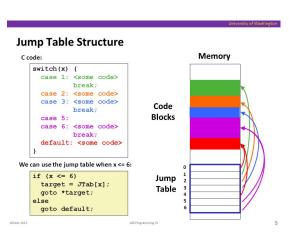
```
code to evaluate E into t
        goto test
L_1:
        code for S_1
       goto next
       code for S_2
L_2:
        goto next
      code for S_{n-1}
L_{n-1}:
       goto next
      S_n
        goto next
test:
        if t = V_1 goto L_1
        if t = V_2 goto L_2
        if t = V_{n-1} goto L_{n-1}
        goto L_n
next:
```

```
case t V_1 \mathbb{L}_1
case t V_2 L_2
case t V_{n-1} L_{n-1}
case ttL,
next:
```

case 三地址代码

```
code to evaluate E into t
         goto test
         code for S_1
L_1:
         goto next
L_2:
         code for S2
         goto next
         code for S_{n-1}
L_{n-1}:
         goto next
L":
         code for S_n
         goto next
test:
         if t = V_1 goto L_1
         if t = V_2 goto L_2
         if t = V_{n-1} goto L_{n-1}
        goto L<sub>n</sub>
next:
```

```
case t V_1 L<sub>1</sub>
case t 1/2 Lo
. . .
case t V_{n-1} L<sub>n-1</sub>
case ttL.
next:
```



Jump Table 优化

语义分析

函数/过程的中间代码翻译

新增文法以支持函数定义与调用

$$D \rightarrow \text{ define } T \text{ id } (F) \{S\}$$

$$F \rightarrow \epsilon \mid T \text{ id }, F$$

$$S \rightarrow \text{ return } E;$$

$$E \rightarrow \text{ id } (A)$$

$$A \rightarrow \epsilon \mid E, A$$

函数定义

$$\begin{array}{ccc} D & \rightarrow & \mathbf{define} \ T \ \mathbf{id} \ (F) \ \{S\} \\ F & \rightarrow & \epsilon \mid T \ \mathbf{id} \ , F \\ S & \rightarrow & \mathbf{return} \ E \ ; \end{array}$$

函数名 id 放入当前符号表, 建立新的符号表, 处理形参 F 与函数体 S

函数调用

$$egin{array}{cccccc} E &
ightarrow & {
m id} & (A) \ A &
ightarrow & \epsilon & E & A \ \end{array}$$

 x_1 param x_2 ... x_n

param x_n call p, n

函数调用

```
S:: = CALL id(Elist) { S. code := Elist. code

| gencode("CALL", id. place, Elist. number)

Elist:: = Elist, ,E { Elist. code := E. code || Elist. code | just |
| gencode("PARAM", E. place);

Elist. number := Elist, number + 1 }

Elist:: = E { Elist. code := E. code || gencode("PARAM", E. place);

Elist. number := 1 }
```

C 语言并未规定参数计算的顺序



计算实参 x_1 的中间代码 $\mathbf{param}\ x_1$ 计算实参 x_2 的中间代码 $\mathbf{param}\ x_2$...

计算实参 x_m 的中间代码 $\mathbf{param}\ x_n$ $\mathbf{call}\ p, n$

计算实参 x_1 的中间代码 计算实参 x_2 的中间代码

计算实参 x_m 的中间代码 $\mathbf{param} x_1$

param x_2

. . .

...

 $\mathbf{param} \ x_n$ $\mathbf{call} \ p, n$

函数调用

```
S:: = CALL id(Elist)
     { Count := 0; S. code := Elist. code;
       while NOT EmptyO(g) do
       begin
         t := HeadO(q):
         S. code := S. code | gencode("PARAM",t);
         DelQ(q); Count := Count + 1
       end:
       S. code := S. code | qencode("CALL", id. place, Cour
Elist: = Elist, ,E { Elist. code := E. code || Elist, code;
                     EnterQ(E. place,q)}
Elist::=E
                   { Elist. code := E. code; CreateQ(q);
                     EnterQ(E.place,q)}
```

集中生成 param 指令, 代码更紧凑

Thank You!



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