# SpotServe: Serving Generative Large Language Models on Preemptible Instances

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## Author

#### **Research Direction:**

- machine learning systems
- data management
- distributed computing

#### Selected Publications:

- SpotServe(ASPLOS'2024)
- SpecInfer(ASPLOS'2024)
- Galvatron(VLDB 2023)







Xupeng Miao <a href="https://hsword.github.io/">https://hsword.github.io/</a>



## Outline

- Background
- Design
- Evaluation
- Thinking

## Background: The Scale Of LLM

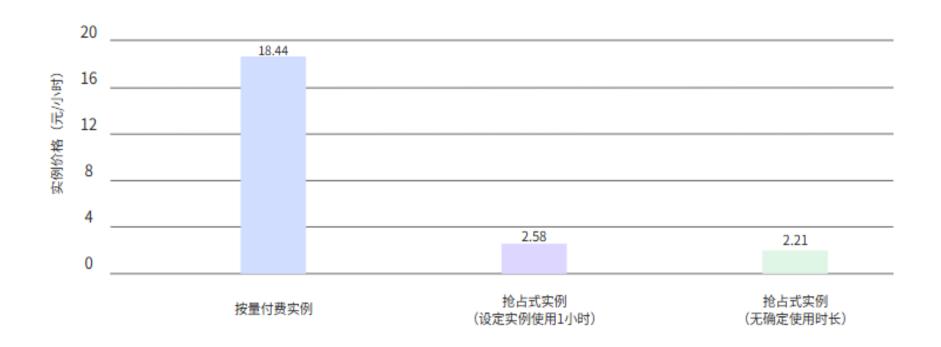
 high computational and memory requirements of LLMs make it challenging to serve them on hardware platforms

model	parameters	device	num	type
GPT-3	175 bilion	Nvidia A100-40GB	16	float

# Background: Spot Instances

#### Modern clouds offer:

- on-demand instances
- spot instances

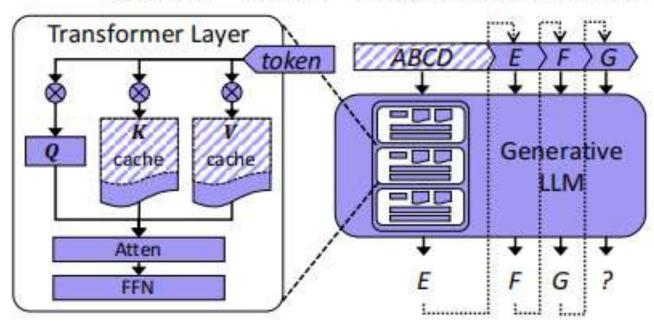


## Background: Generative LLM inference

- $I_{exe}$  is execution latency
- $t_{exe}$  is the LLM's execution time

#### Execution Latency:

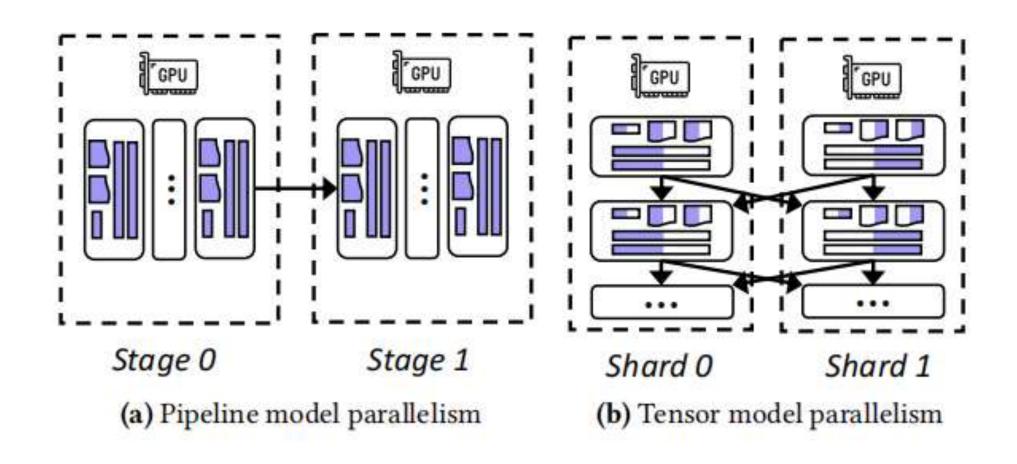
$$l_{exe}(3|4) = t_{exe}(4) + t_{exe}(1) + t_{exe}(1) + t_{exe}(1)$$



$$\begin{split} l_{exe}(S_{out}|S_{in}) &= t_{exe}(S_{in}) + \sum_{i=1}^{S_{out}} t_{exe}(S_{in} + i) \\ &\approx t_{exe}(S_{in}) + S_{out} \times t_{exe}(1) \end{split}$$

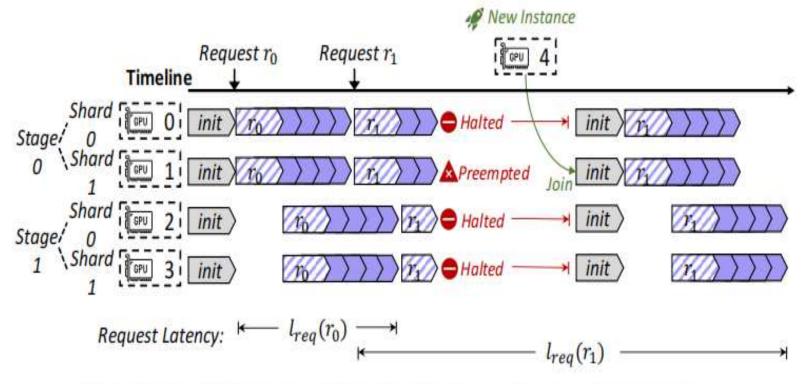
## Background: Distributed Infernece

 Inference pipeline combine pipeline model parallelism and tensor model parallelism



# Background: Preemptible LLM Inference

- One Preempted hang all the other instances
- New instance joins need initialization costs



(b) Distributed LLM inference (P=2, M=2, batch size=1) on preemptible instances

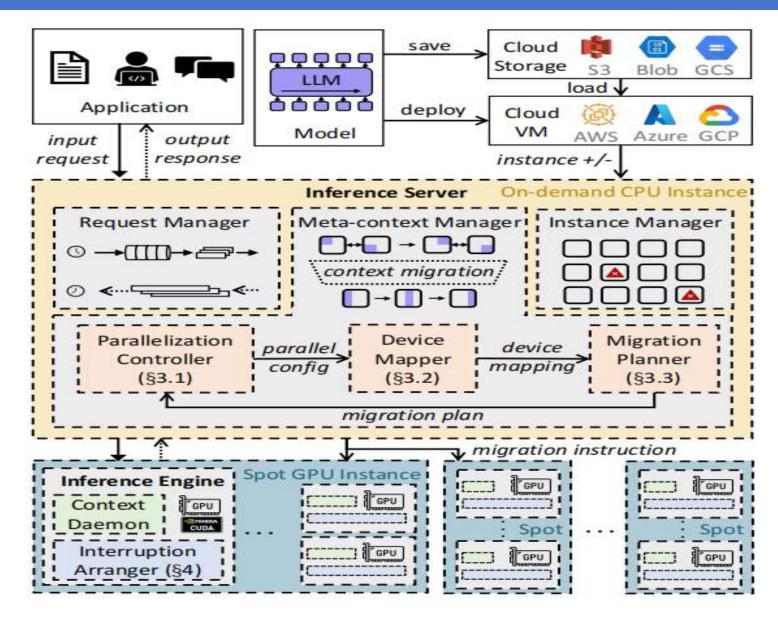
# Design:SpotServe

#### **Inference Server:**

- Request Manager
- Meta-context Manager
- Instance Manager

#### Inference Engine:

- Context Daemon
- Interruption Arranger



# Design:Parallelization Controller

 Parallelization Controller adjusts the parallelization configuration

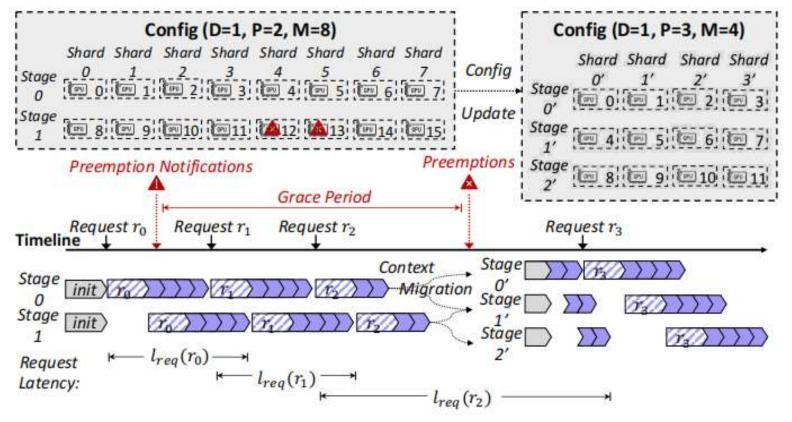
 Algorithm 1 balance the trade-off among throughput latency, and cost

#### Algorithm 1 Adaptive configuration optimizer.

```
1: function ConfigOptimizer(N_t, C_t, \alpha_t)
         if \exists C.\phi(C) \geq \alpha_t and cloud has enough instances for C
    then
             C_{t+1} \leftarrow \operatorname{arg\,min}_{C|\phi(C) \geq \alpha_t} l_{req}(C)
 3:
         else
 4:
 5:
             C_{t+1} \leftarrow \arg\max_{C|N_t} \phi(C)
         \Delta \leftarrow \# \operatorname{Instances}(C_{t+1}) - N_t
         if \Delta > 0 then
              InstanceManager.alloc(\Delta, ondemand_and_spot)
 9:
         else
              InstanceManager.free(-\Delta, ondemand_first)
10:
         ConfigUpdate(C_t, C_{t+1})
11:
```

## Design:Device Mapper

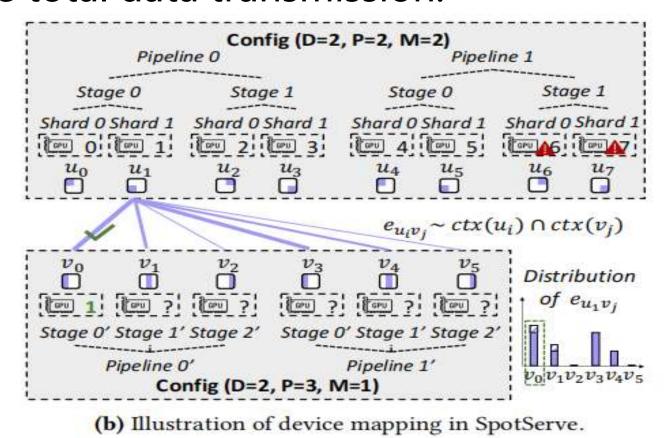
 A context migration mechanism can resume interrupted requests' inference



(a) Illustration of configuration update and context migration in SpotServe.

# Design:Device Mapper

- Device mapping problem → bipartite graph matching task
- Maximally reuses the model parameters and KV cache and minimizes the total data transmission.



## Design:Migration Planner

Algorithm2 is a progressive migration schedule

 Consider the memory usage during the progressive migration

```
Algorithm 2 Workflow of the SpotServe migration planner.
    ▶ Progressive Migration
 1: function MigrationPlanner(context ctx, plan = [])
        plan.append(<migrate, ctx.cache>)
 2:
        O ← Layer migration order from MemOptMigPlanner
 3:
        for layer index i in range(0, #layers) do
 4:
            plan.append(<migrate, ctx.weight[O_i]>)
 5:
            p \leftarrow \text{Get pipeline stage index of layer } O_i
 6:
            if stage p completes all context migration then
 7:
                plan.append(<start, instances of stage p>)
 8:
    ▶ Memory Optimized Migration
 9: function MemOptMigPlanner(maximum buffer size U_{max})
        0 \leftarrow [], X \leftarrow \{\}
10:
        Instance buffer memory usage U \in \{0\}^N
11:
        for layer index i in range(0, #layers) do
12:
            if (migrate, ctx.weight[i]) doesn't exceed U_{max} then
13:
                Update buffer memory usage U
14:
                O.append(i)
15:
            else
16:
                X.add(i)
17:
        while X is not empty do
18:
19:
            x_{opt} \leftarrow
                 \underset{x \in \mathbf{X}}{\arg\min} \max_{0 \le i \le N-1} \{ \mathbf{U}_i \mid (\text{migrate}, \mathsf{ctx.weight}[x]) \}
            O.append(r_{opt})
20:
            X.remove(x_{opt})
21:
```

## Design: JIT Arrangement

#### Utilize grace period

- Each spot GPU instance includes an interruption arranger
- Just-in-time (JIT) arrangement decide when to stop decoding

$$S_t = arg \max_{0 \le S \le S_{out}} \{l_{exe}(S|C_t) < T^- - T_{mig}\}$$
  

$$\cap T_{mig} < l_{exe}(S_t|C_t)$$

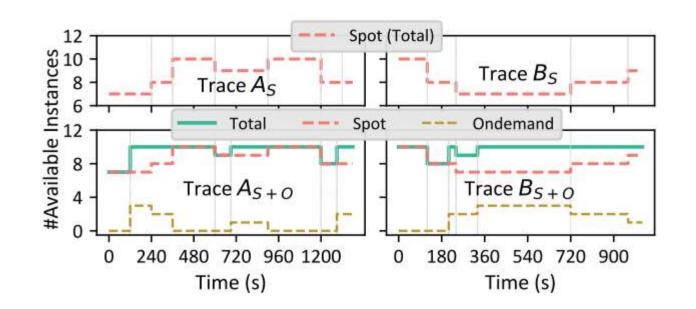
# Evaluation: Experiment Setup

#### Baseline:

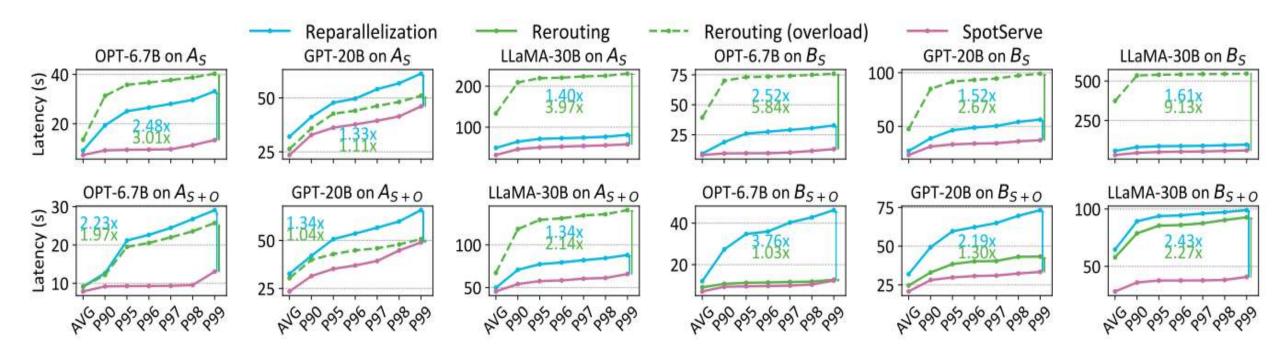
- Request rerouting
- Model reparallelization

#### Models:

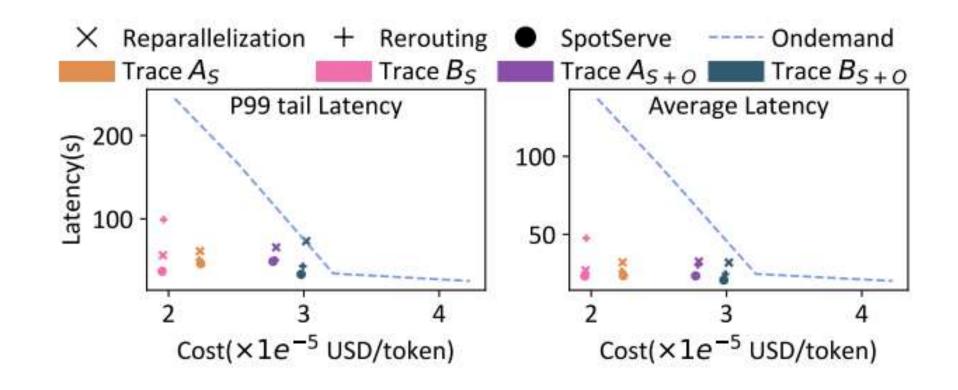
- OPT-6.7B
- GPT-20B
- LLaMA-30B



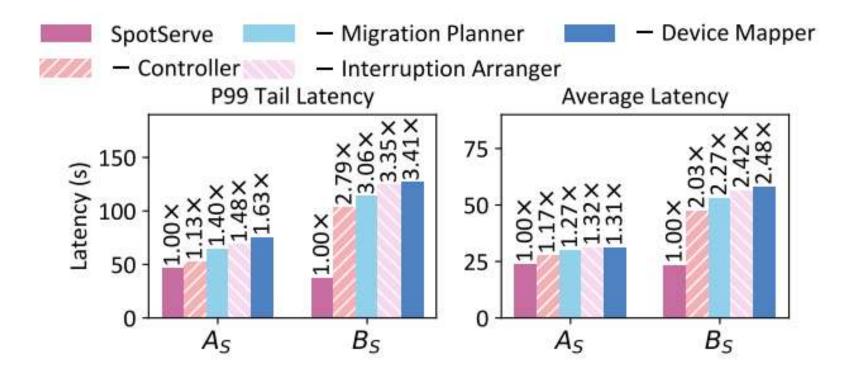
## Evaluation



## Evaluation



## Evaluation



# Thinking

- 文章有什么问题
- 该系统的并行策略中考虑的都是相同类型的GPU实例,并行策略上也是将 其按照同构资源处理
- 2. 在模型映射部分,只将模型参数的可复用程度作为边权的考虑,没有考虑设备间数据传输的能力差异
- 该方法能不能应用到别的场景
   对于智能家居的分布式计算来说,对于一个有处理任务的设备来说,其他当前状态下空闲的设备的资源对它来说视作可抢占资源,利用目前可用资源配置出最优的并行推理流水线

# Thanks