COMP9334
Capacity Planning for Computer Systems and Networks

Week 2A: Operational Analysis (2).
Workload Characterisation

Last lecture

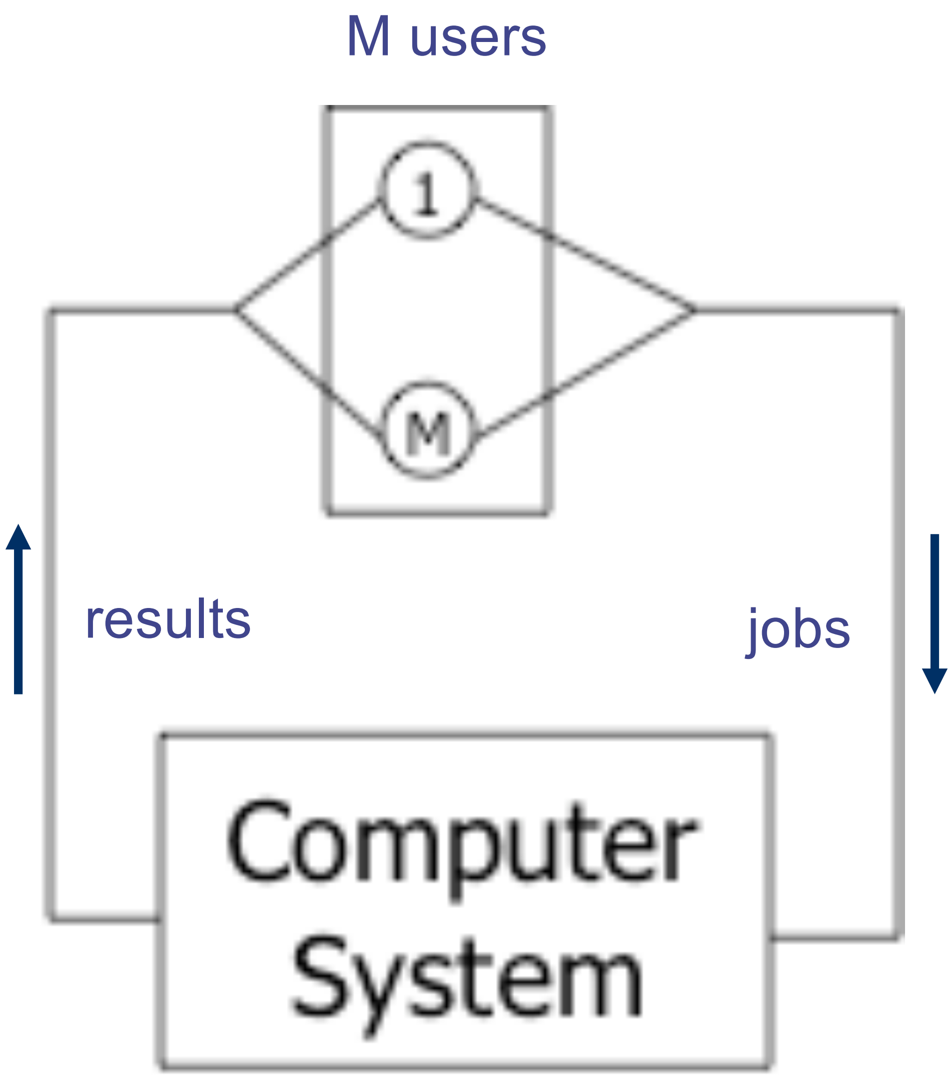
- Modelling a computer system as a queueing network
- Operational analysis on queueing networks
- We have derived these operational laws
  - Utilisation law U(j) = X(j) S(j)
  - Forced flow law X(j) = V(j) X(0)
  - Service demand law D(j) = V(j) S(j) = U(j) / X(0)
  - Little's law N = X R

This lecture

- Operational analysis (Continued)
  - Using operational law for
    - Performance analysis
    - Bottleneck analysis

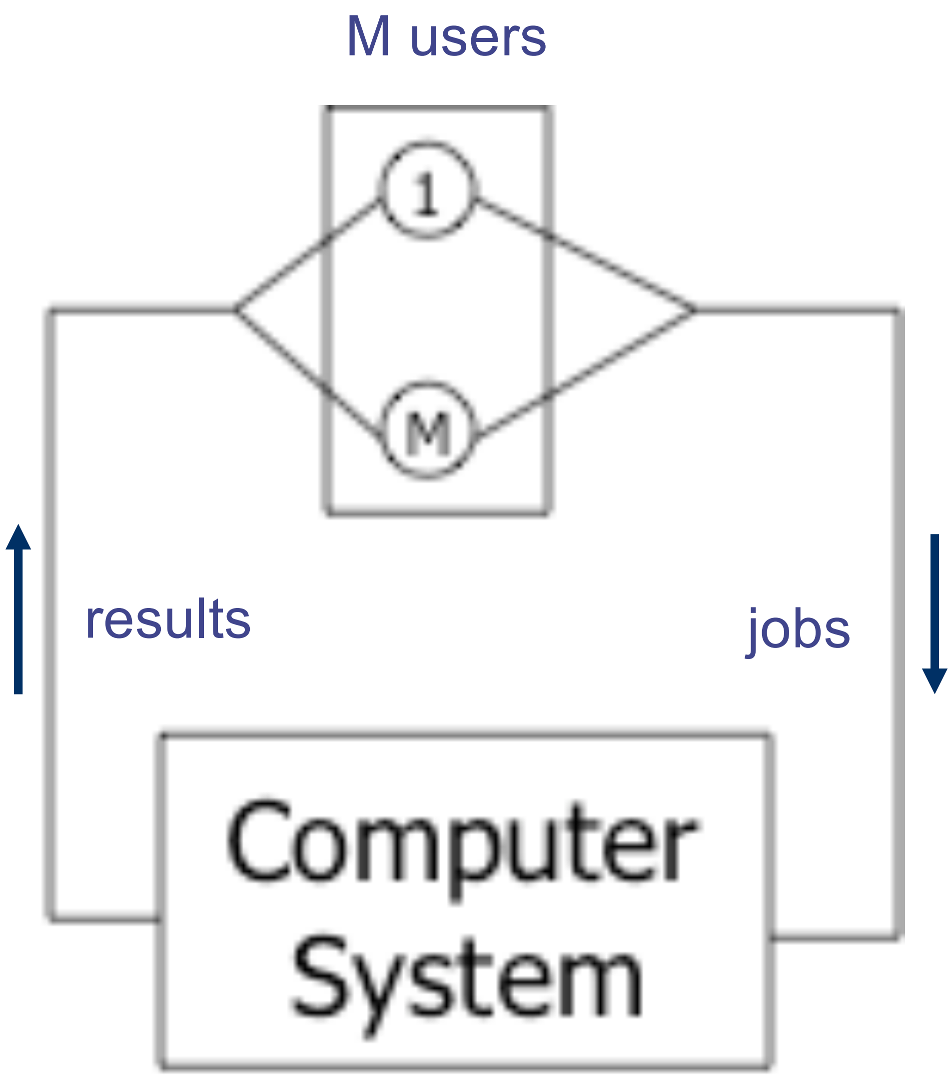
- Workload characterisation
  - Poisson process and its properties

Interactive systems



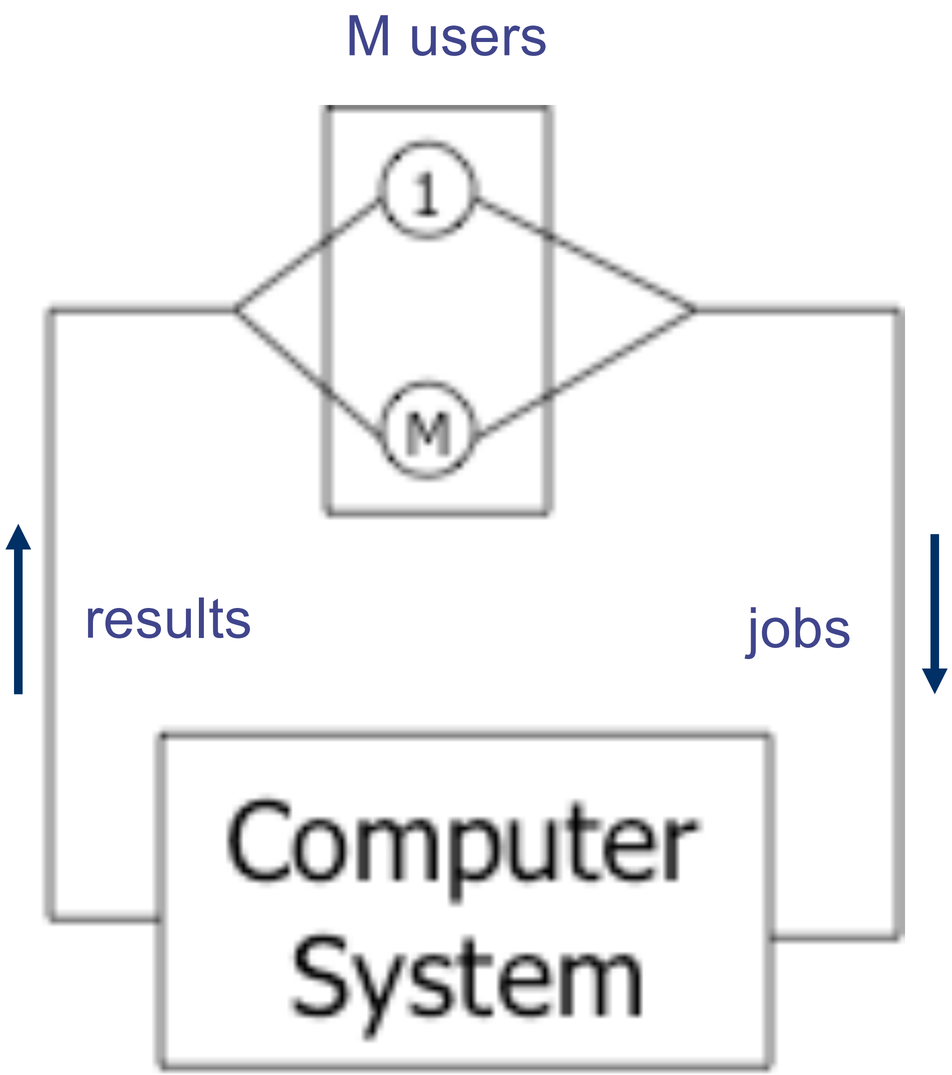
- An interactive system is used to model the interaction between humans (users) and computers
- The system consists of
  - A number of users
  - A computer system

Interactive systems (Cont'd)



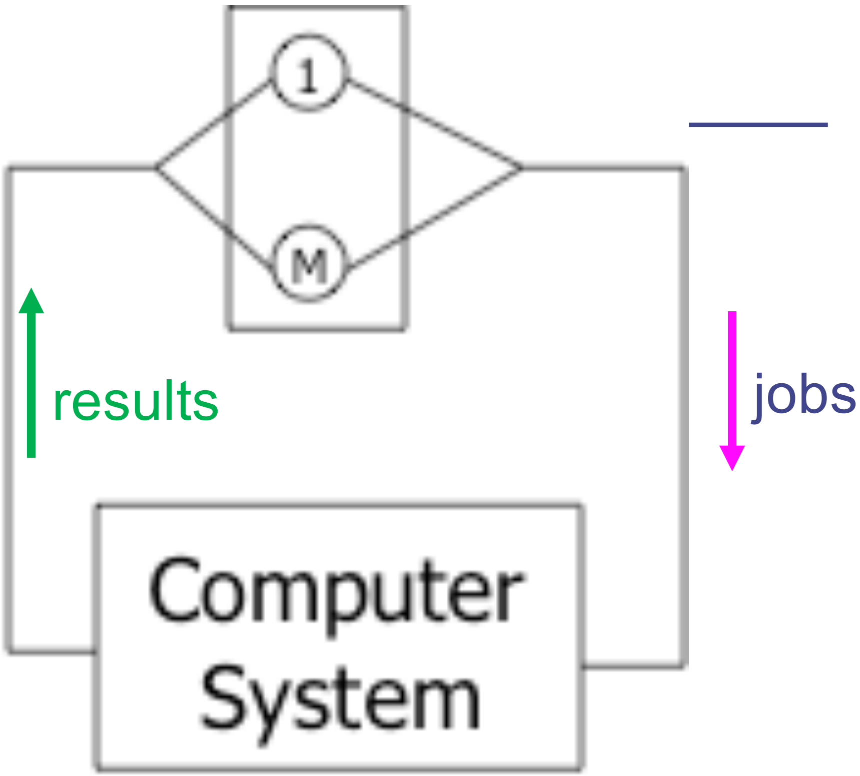
- Interactions
  - Users send jobs to computer systems
  - After finishing processing a job, the computer system returns the result to the user
  - A user, after inspecting the results from the computer system, will send another job to the system

Interactive systems: Modelling assumptions

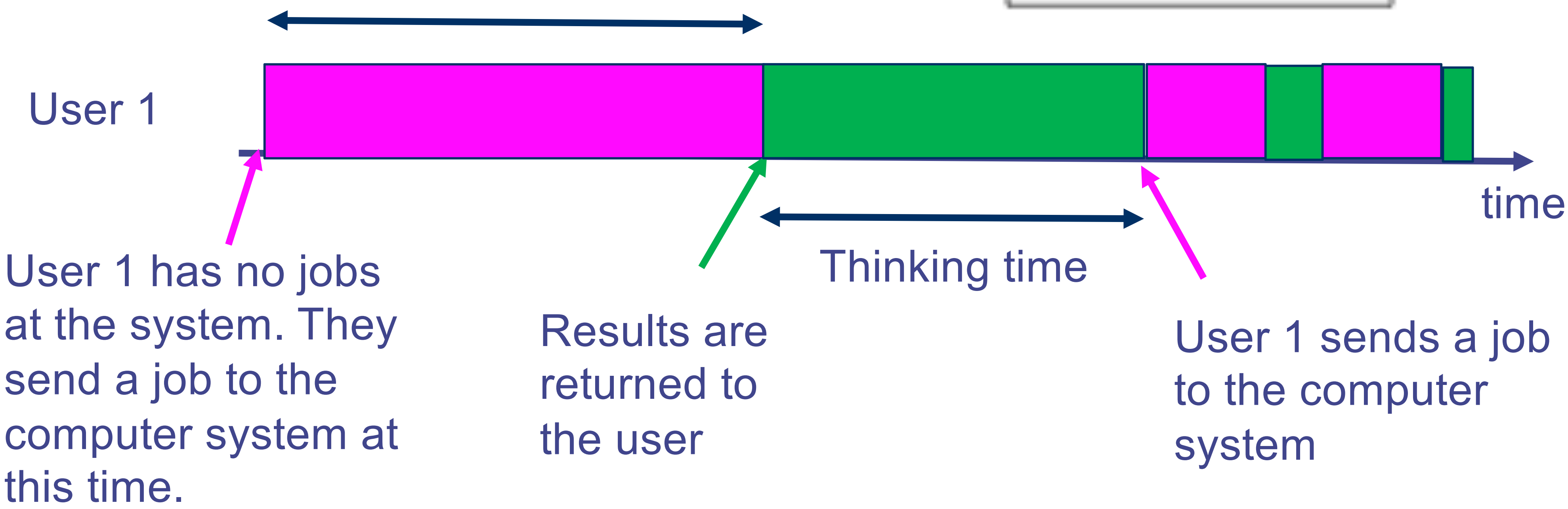


- Analyze interactive systems with specific assumptions
  - Fixed number of users denoted by M
  - Each user can have at most 1 job at the computer system
  - Each user goes through a cycle consisting of
    - Thinking time
    - Waiting for result time

Interactive cycle

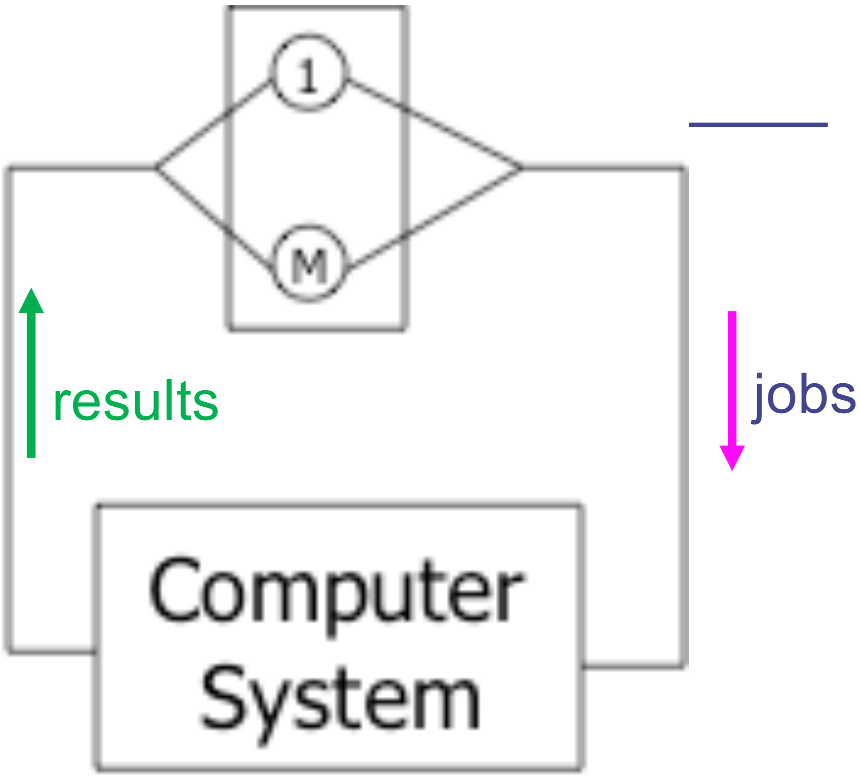


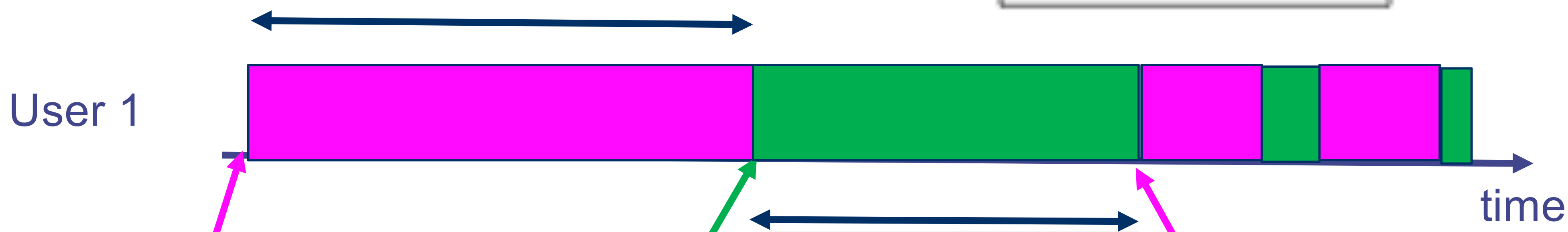
- The time the job from User 1 spends in the computer system
- User 1 waiting for the results



Interactive cycle

- User 1’s perspective: waiting for the result of their job
- Computer system’s perspective: Response time of the job from User 1





User 1 has no jobs at the system. They send a job to the computer system at this time.

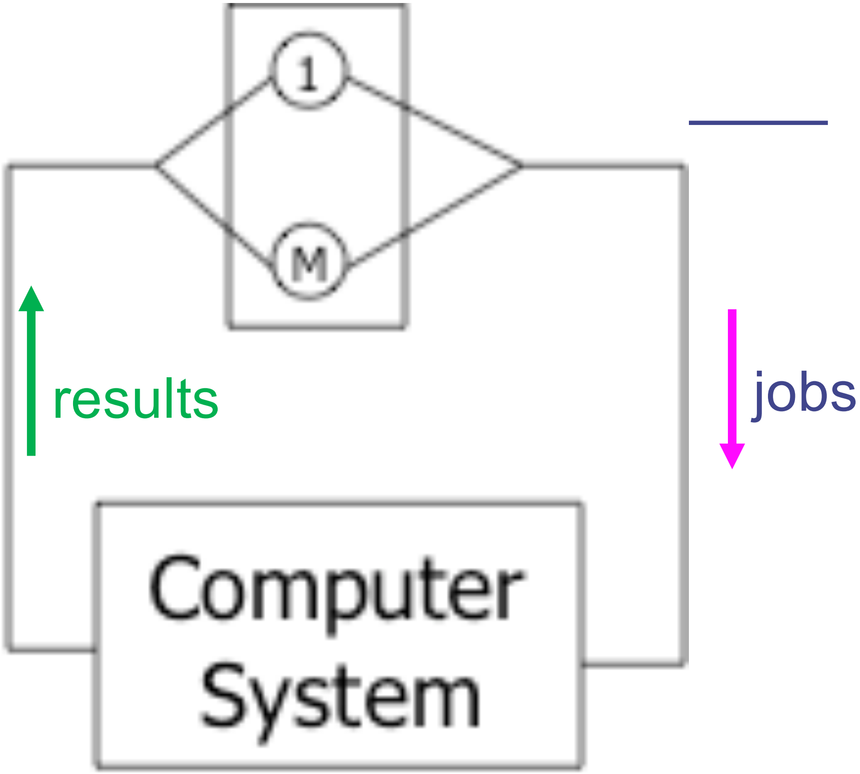
Results are returned to the user

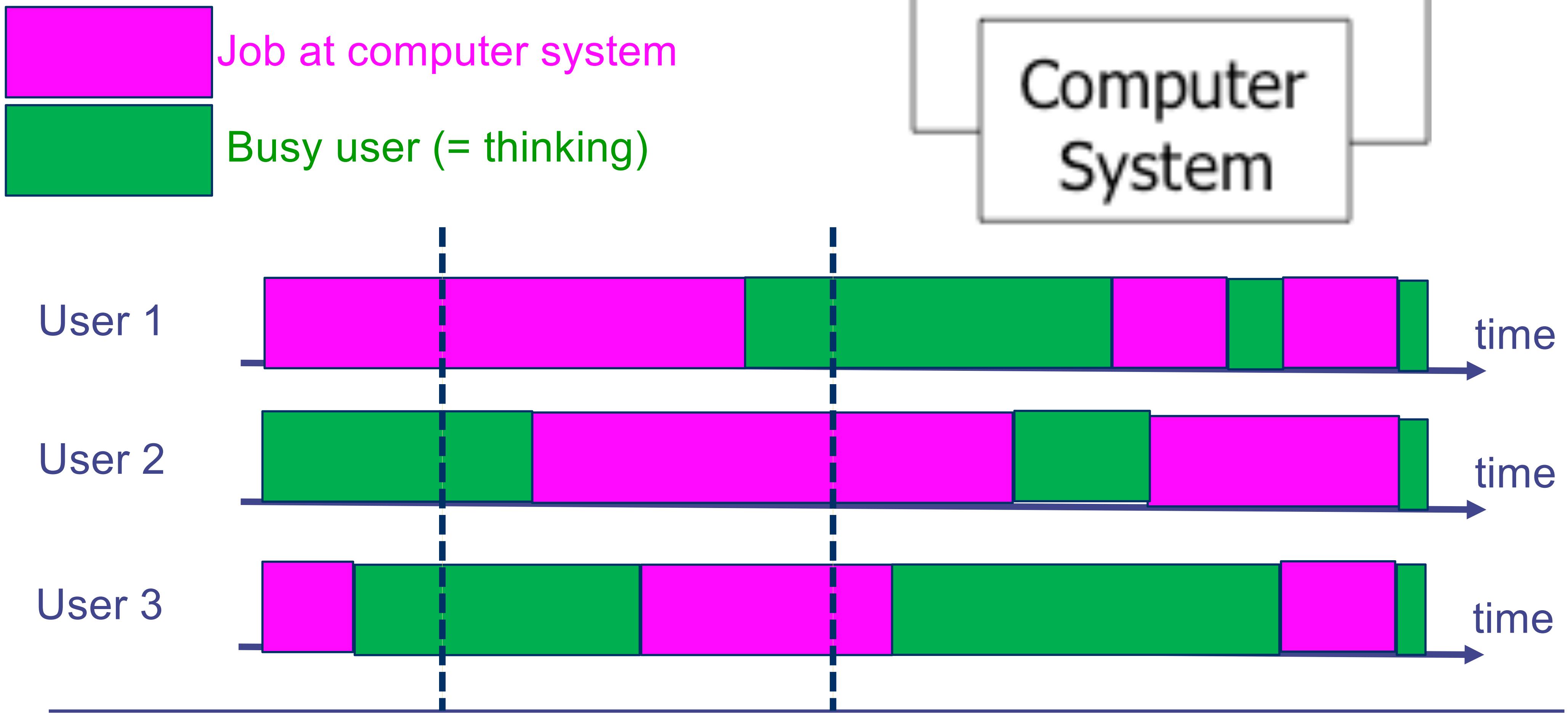
Thinking time = Processing time of the user. That’s why a user is depicted as a circle

User 1 sends a job to the computer system

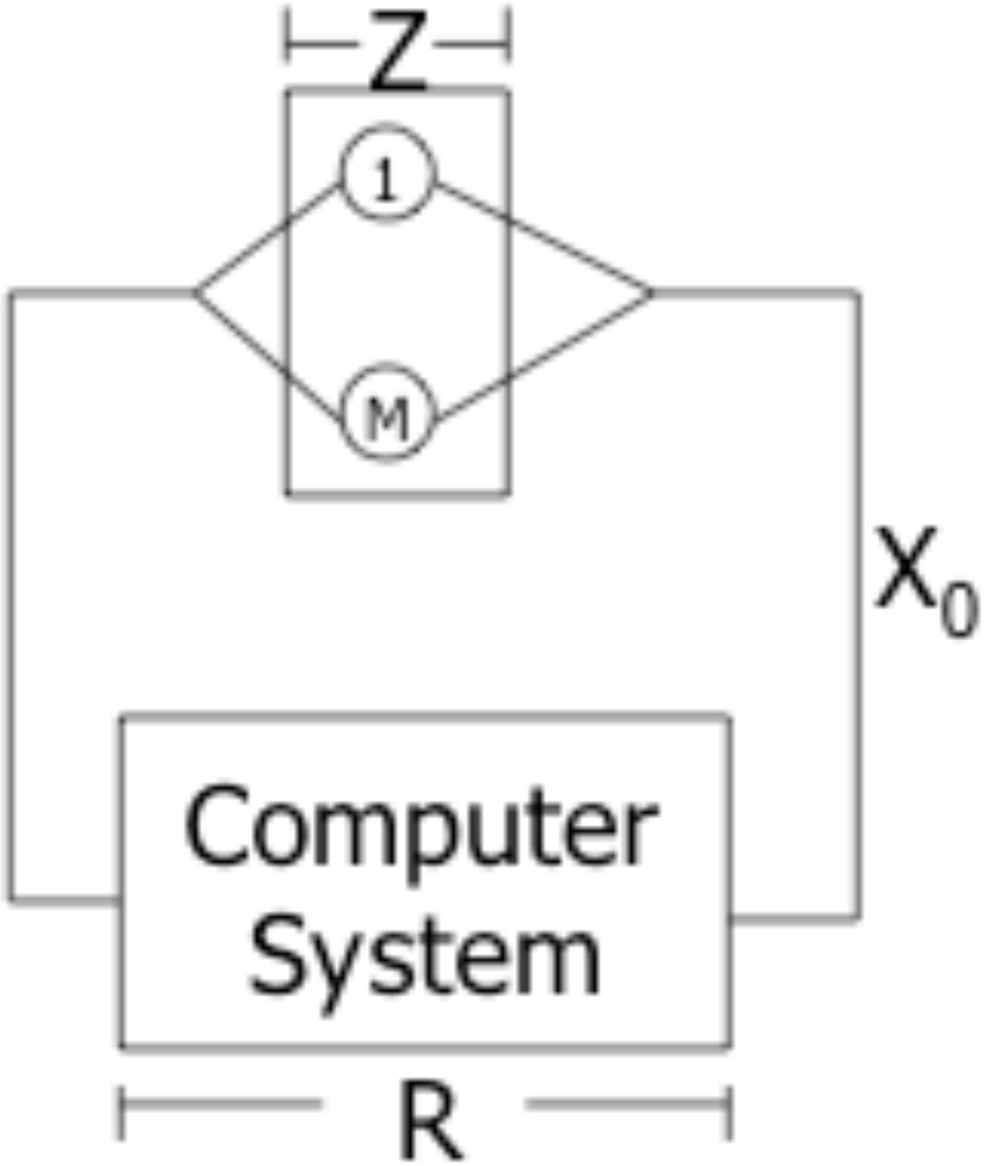
Quiz

- Question: At any time, what is the sum of the number of busy users and the number of jobs at the computer system?



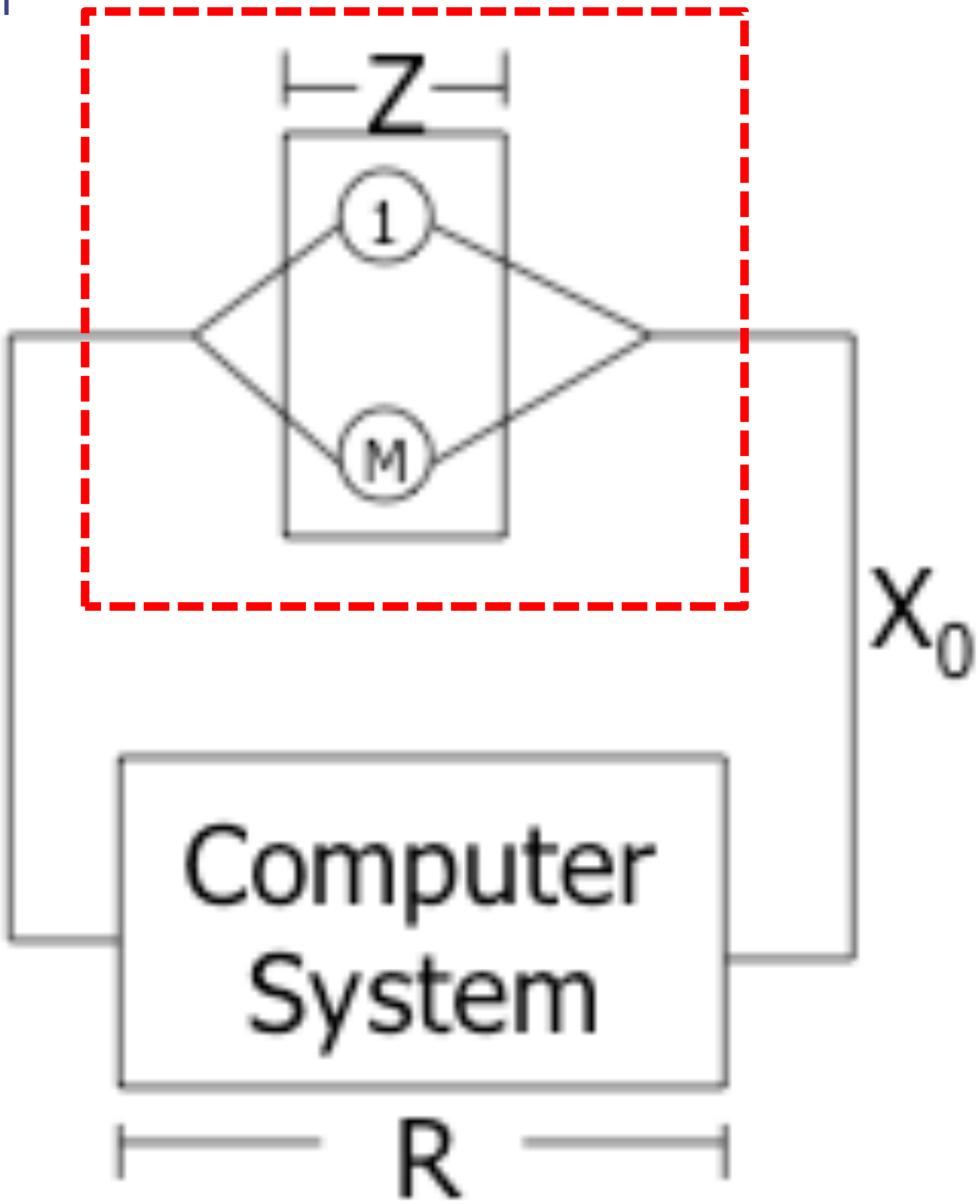


Interactive system: Parameters



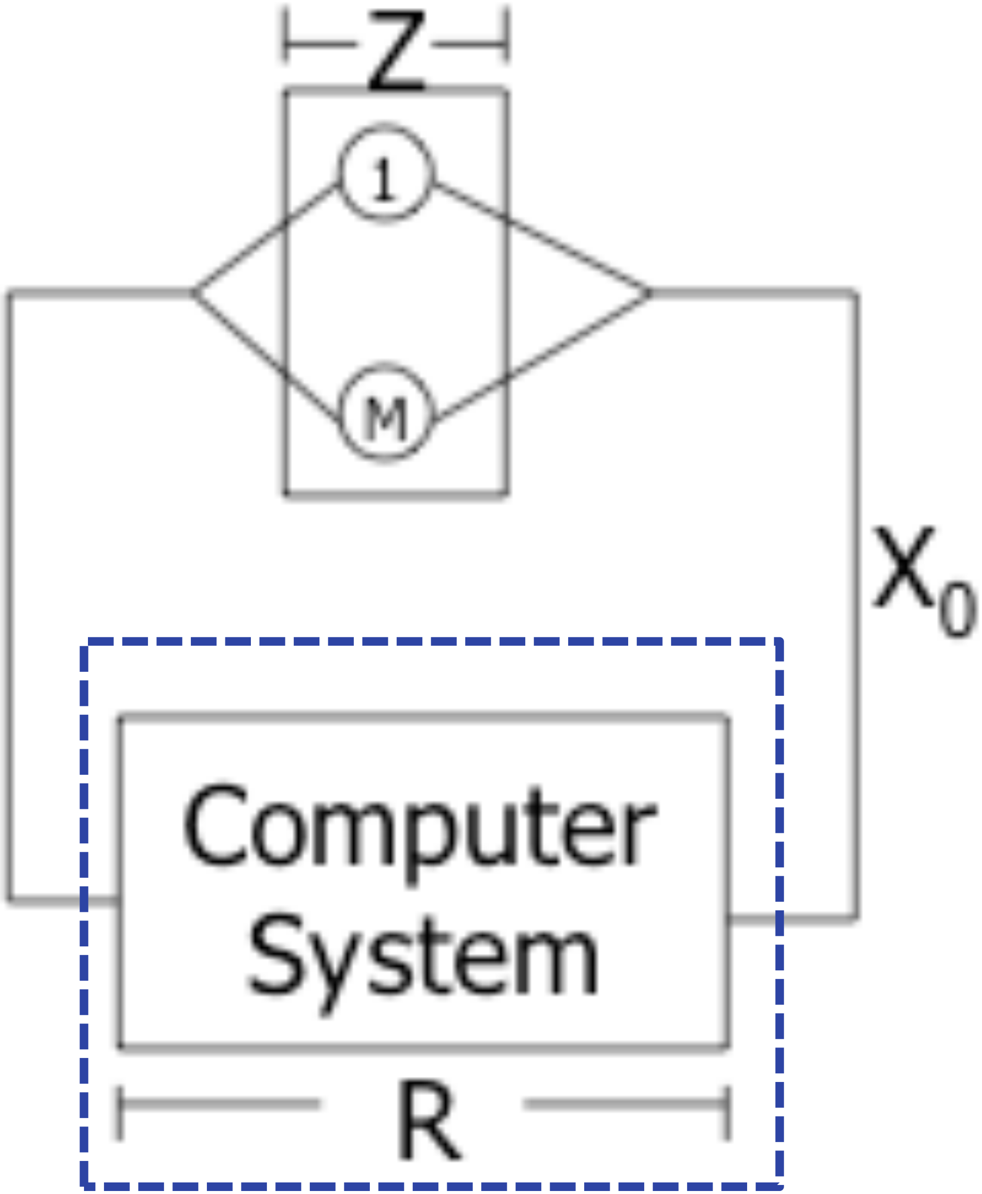
- M interactive users
- Z = mean thinking time
- R = mean response time of the computer system
- X0 = throughput

Analyzing interactive system: Quiz 1



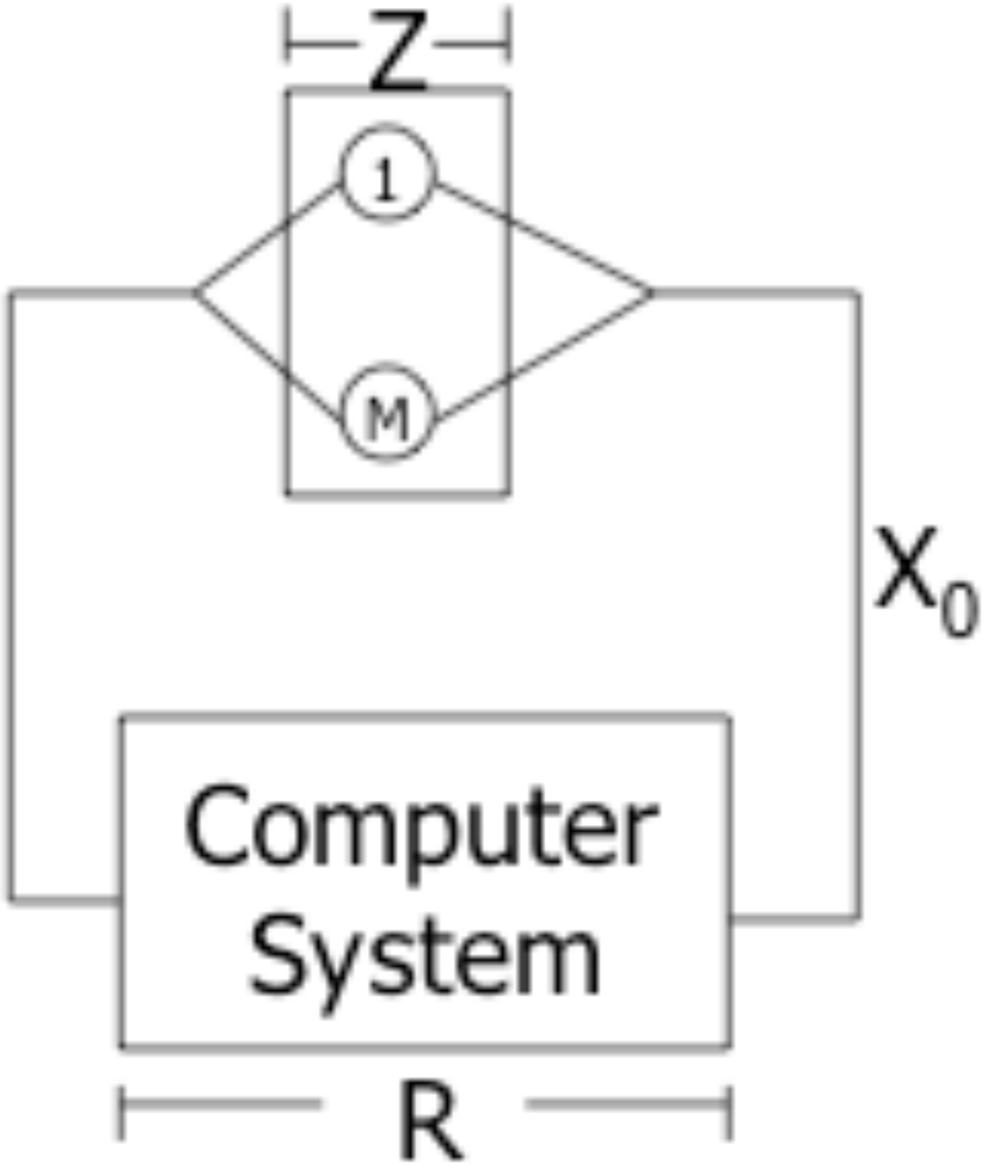
- Mavg = mean # busy users
- Z = mean thinking time
- X0 = throughput
- Apply Little's Law to the red box. What do you get?
  -

Analyzing interactive system: Quiz 2



- Navg = average # jobs in the computer system
- R = mean response time at the computer system
- X0 = throughput
- Apply Little's Law to the computer system (i.e. the blue box), what do you get?
  -

Analyzing interactive system: Quiz 3



- Quiz 1:
- Quiz 2:

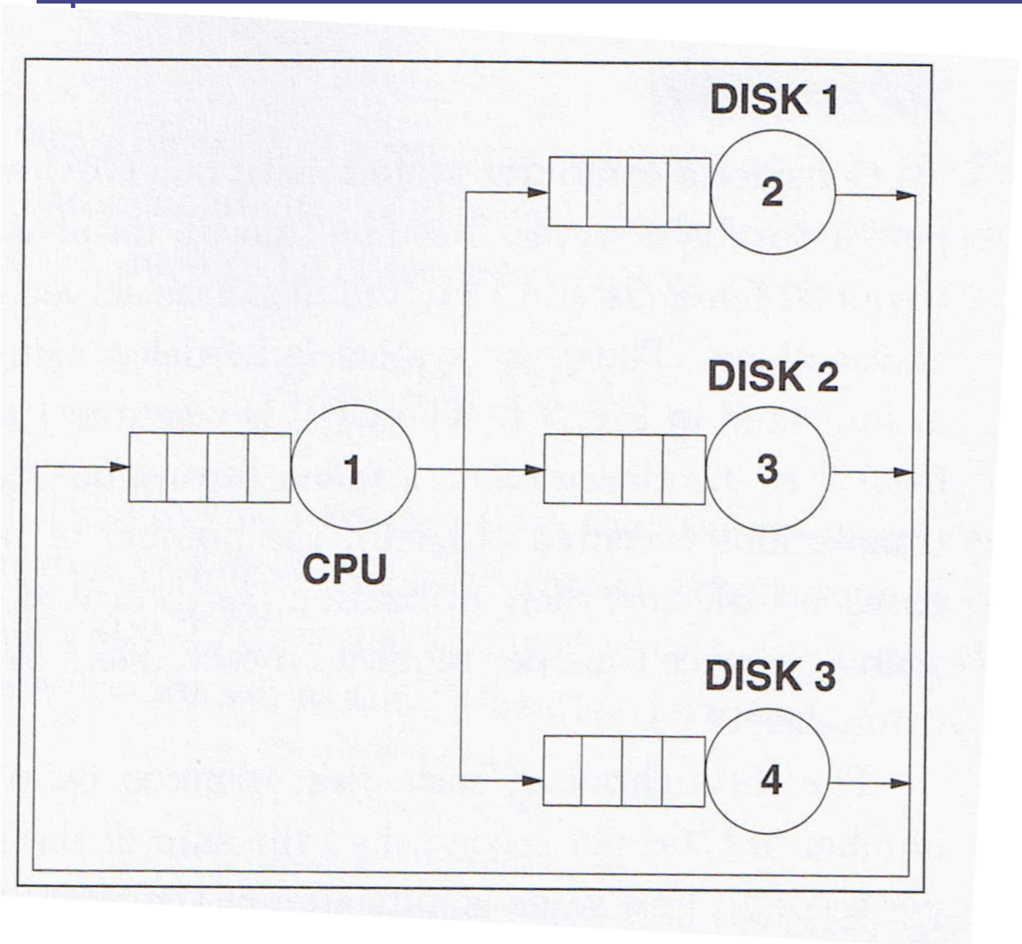
- What is Mavg + Navg?
  -

- Interactive response time law
  -

The operational laws

- These are the operational laws
  - Utilisation law $U(j) = X(j)\ S(j)$
  - Forced flow law $X(j) = V(j)\ X(0)$
  - Service demand law $D(j) = V(j)\ S(j) = U(j)\ /\ X(0)$
  - Little's law $N = X\ R$
  - Interactive response time $M = X(0)\ (R+Z)$
- Applications
  - Mean value analysis (later in the course)
  - Bottleneck analysis
  - Modification analysis

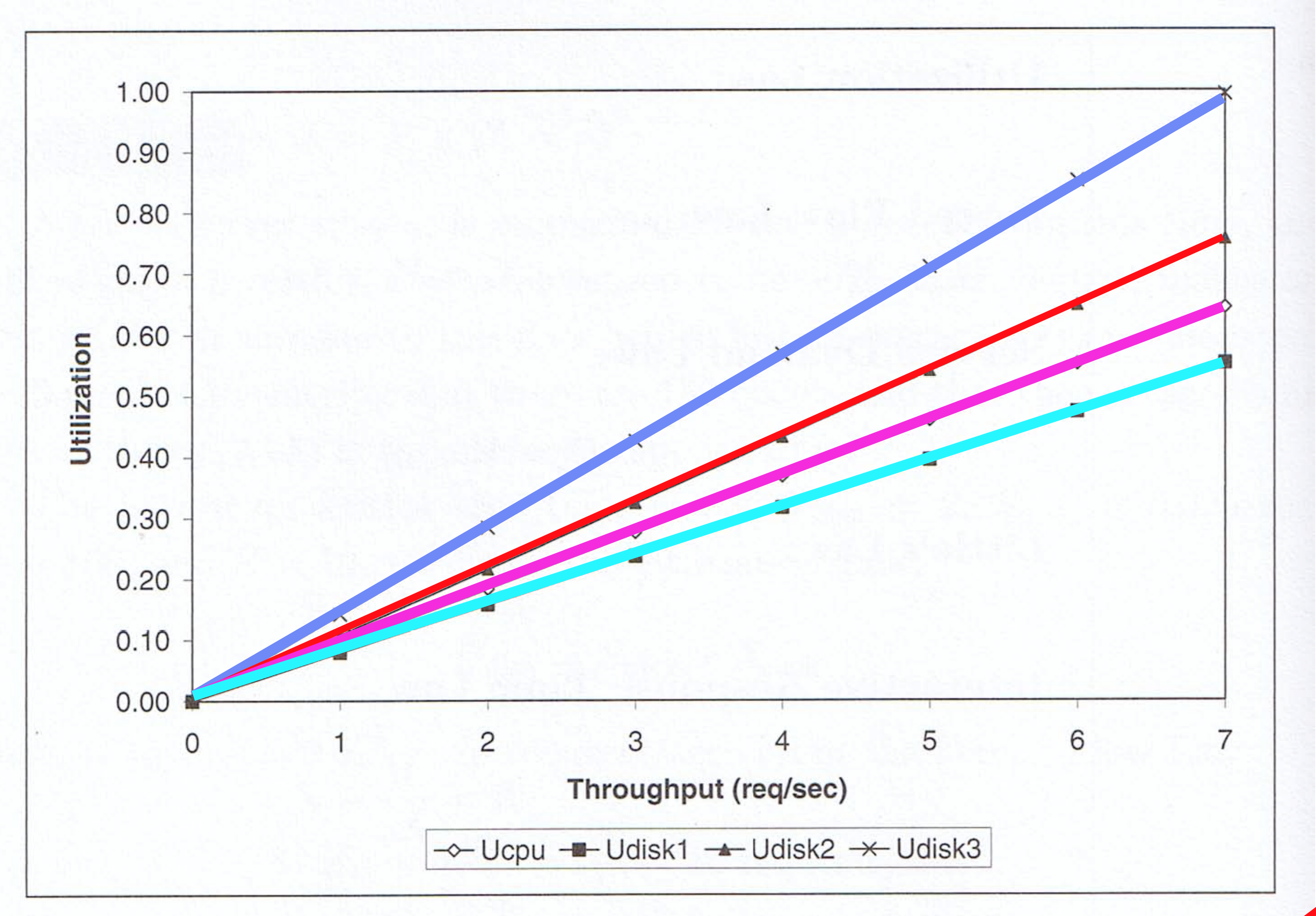
Bottleneck analysis - motivation



	D(j)	Utilisation
Disk 1	79ms	0.30
Disk 2	108ms	0.41
Disk 3	142ms	0.54
CPU	92ms	0.35

Service demand law: D(j) = U(j) / X(0)
==> U(j) = D(j) X(0)
Utilisation increases with increasing throughput and service demand

Utilisation vs. throughput plot U(j) = D(j) X(0)



Disk 3

Disk 2

CPU

Disk 1

What determines this order?

Observation: For all system throughput:
Utilisation of Disk 3 > Utilisation of Disk 2 >
Utilisation of CPU > Utilisation of Disk 1

Bottleneck analysis

- Recall that utilisation is the busy time of a device divided by measurement time
  - What is the maximum value of utilisation?
- Based on the example on the previous slide, which device will reach the maximum utilisation first?

Bottleneck (1)

- Disk 3 has the highest service demand
- It is the bottleneck of the whole system

Operational law: $X(0) = \frac{U(j)}{D(j)}$

Utilisation limit: $U(j) \leq 1$

$$\left.\begin{array}{l} X(0) = \frac{U(j)}{D(j)} \\ U(j) \leq 1 \end{array}\right\} \quad X(0) \leq \frac{1}{D(j)}$$

Bottleneck (2)

$$X(0) \leq \frac{1}{D(j)}$$ Should hold for all K devices in the system

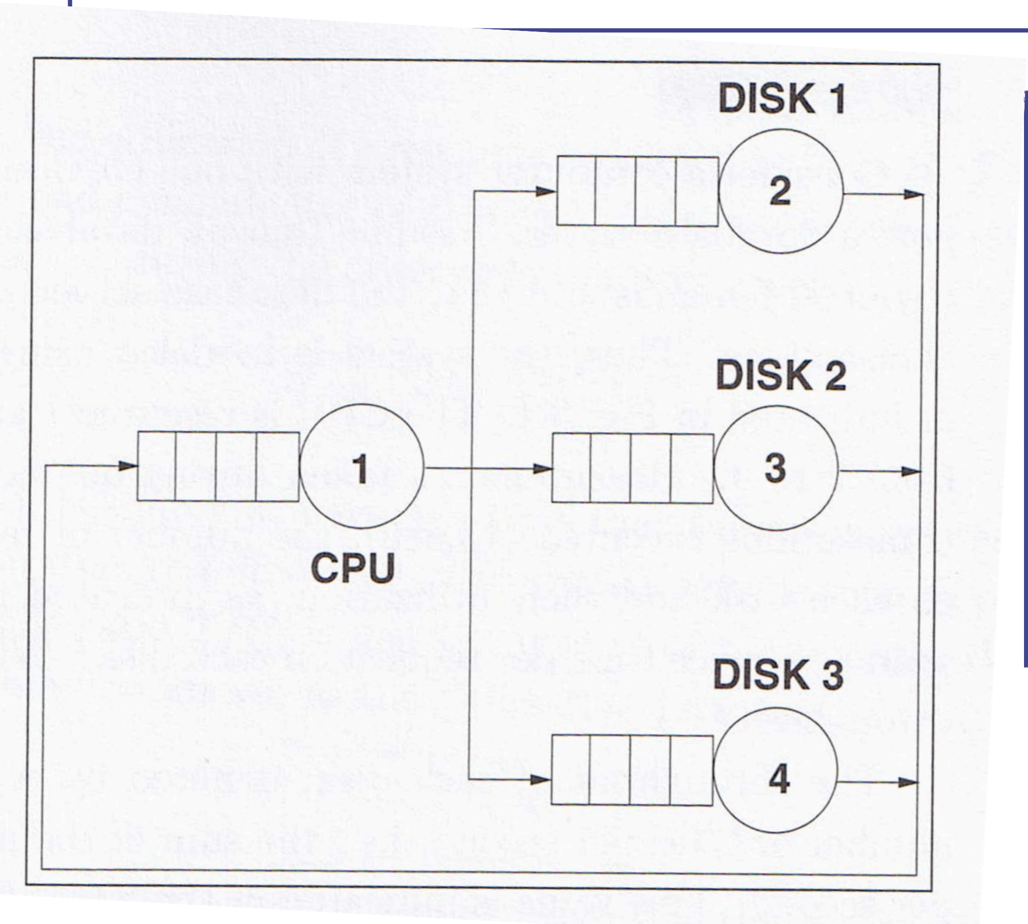
$$i.e. X(0) \leq \frac{1}{D(1)}, \ldots, X(0) \leq \frac{1}{D(K)}$$

$$\Rightarrow X(0) \leq \min \frac{1}{D(j)}$$

$$\Rightarrow X(0) \leq \frac{1}{\max D(j)}$$

Bottleneck throughput is limited by the maximum service demand

Bottleneck exercise



	D(j)	Utilisation
Disk 1	79ms	0.30
Disk 2	108ms	0.41
Disk 3	142ms	0.54
CPU	92ms	0.35

The system throughput is upper bounded by $\frac{1}{0.142}$ = 7.04 jobs/s

If we upgrade Disk 3 by a new disk which is 2 times faster, which device will be the bottleneck after the upgrade? You can assume that service time is inversely proportional to disk speed.

Another throughput bound

- Little's law

$$N = R \times X(0) \geq (\sum_{i=1}^{K} D_i) \times X(0)$$

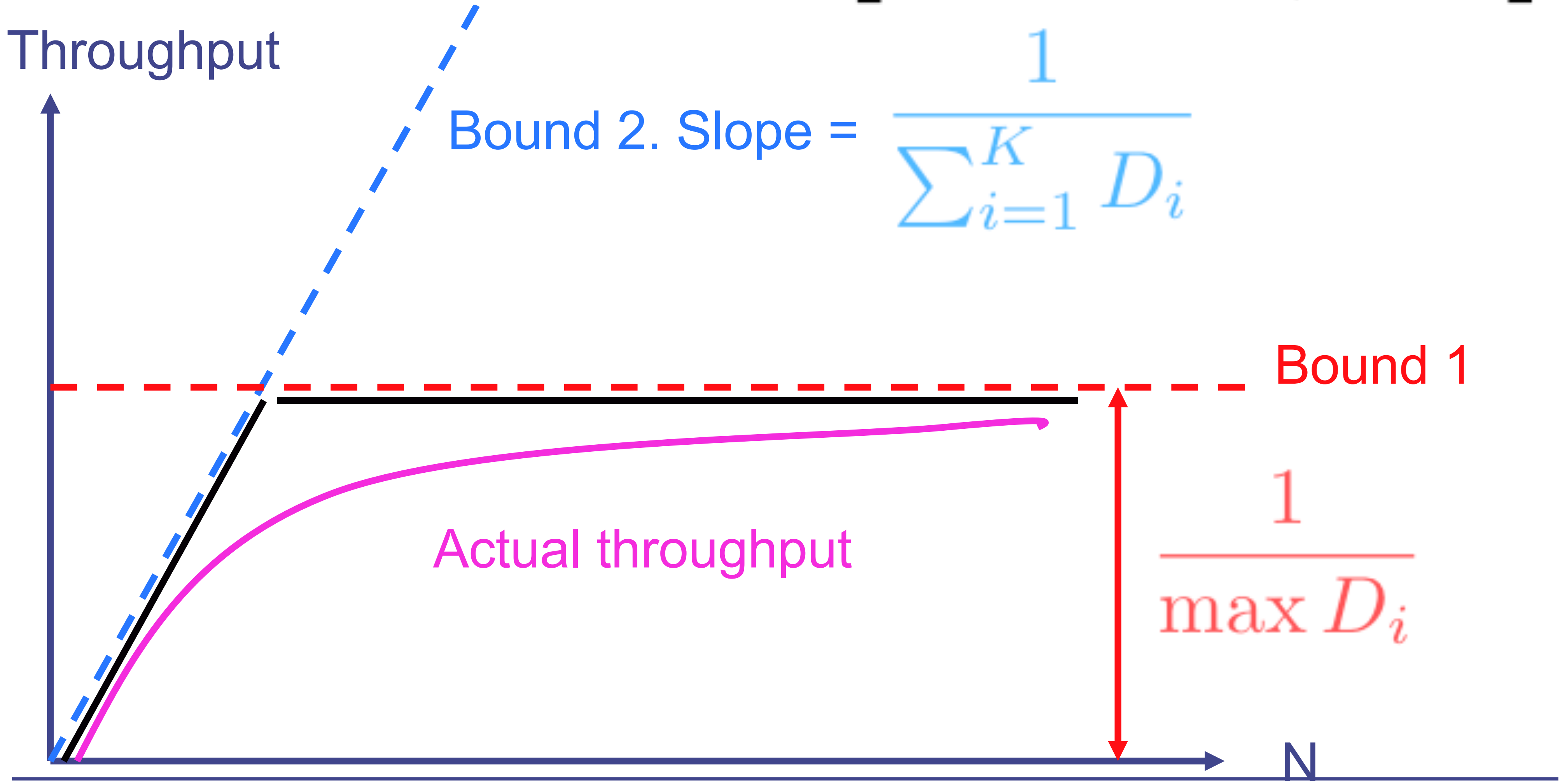
$$\Rightarrow X(0) \leq \frac{N}{\sum_{i=1}^{K} D_i}$$

Previously, we have $X(0) \leq \dfrac{1}{\max D(j)}$

Therefore: $$X(0) \leq \min \left[\frac{1}{\max D_i}, \frac{N}{\sum_{i=1}^{K} D_i} \right]$$

Throughput bounds

$$X(0) \leq \min\left[\frac{1}{\max D_i}, \frac{N}{\sum_{i=1}^{K} D_i}\right]$$



Bottleneck analysis

- Simple to use
  - Needs only utilisation of various components
- Assumes service demand is load independent

Modification analysis (1)

- (Reference: Lazowska Section 5.3.1)
- A company currently has a system (3790) and is considering switching to a new system (8130). The service demands for these two systems are given below:

System	Service demand (seconds)	
	CPU	Disk
3790	4.6	4.0
8130	5.1	1.9

- The company uses the system for interactive application with a think time of 60s.
- Given the same workload, should the company switch to the new system?
- Exercise: Answer this question by using bottleneck analysis. For each system, plot the upper bound of throughput as a function of the number of interactive users.

Modification analysis (2)

Operational analysis

- Operational analysis allows you to bound the system performance but it does NOT allow you to find the throughput and response time of a system
- To order to find the throughput and response time, we need to use queueing analysis
- To order to use queueing analysis, we need to specify the workload

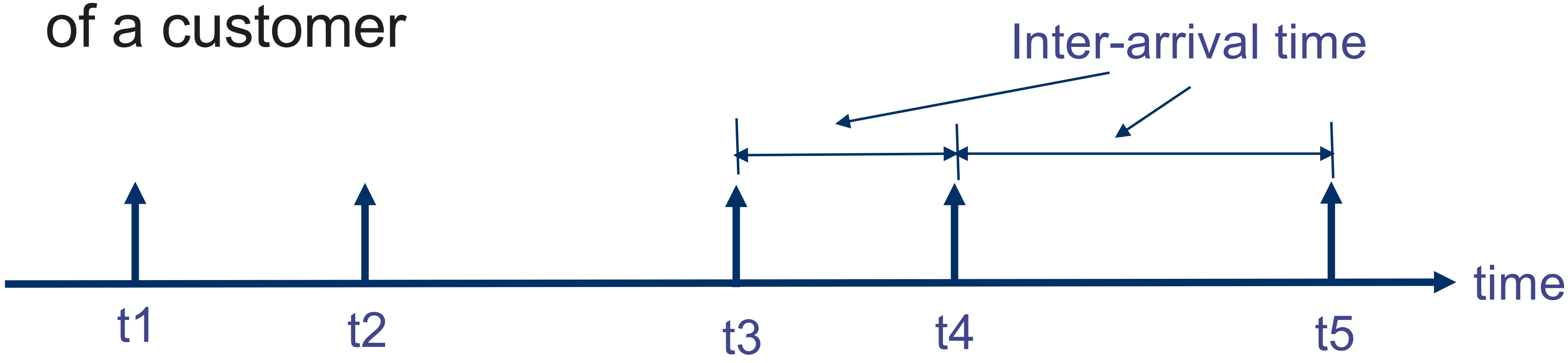
Workload analysis

- Performance depends on workload
  - When we look at the performance bound earlier, the bounds depend on number of users and service demand
  - Queue response time depends on the job arrival probability distribution and job service time distribution
    - Recall from Lecture 1A:
      - Uniform arrival times and uniform processing times result in zero waiting time
      - But non-uniform distributions give non-zero waiting time

- Need to specify workload by using probability distribution.
- We will look at a well-known arrival process called Poisson process today.

Arrival process

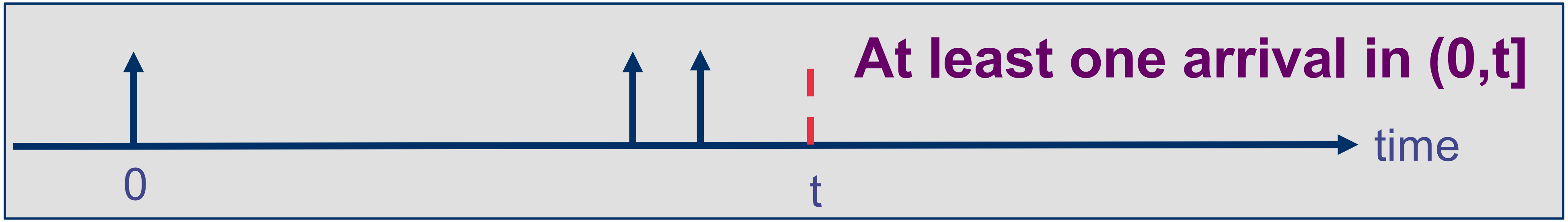
- Each vertical arrow in the time line below depicts the arrival of a customer

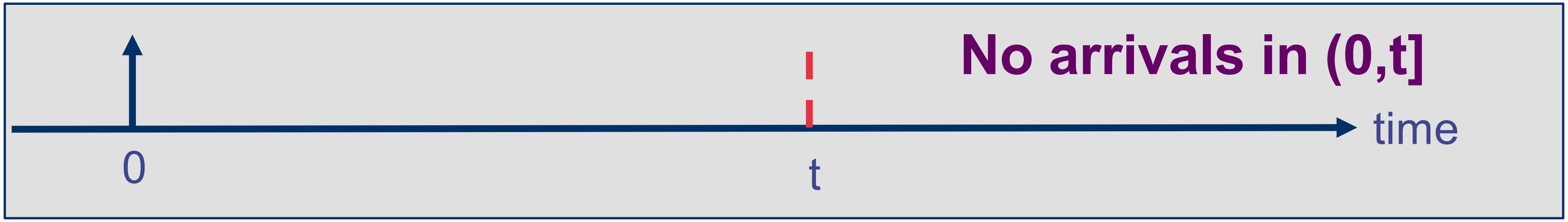


- An arrival can mean
  - A telephone call arriving at a call centre
  - A transaction arriving at a computer system
  - A customer arriving at a checkout counter
  - An HTTP request arriving at a web server
- The inter-arrival time distribution will impact on the response time.
- We will study an inter-arrival distribution that results from a large number of **independent** customers.

Describing arrivals probabilistically

- Assume a customer arrives at time 0

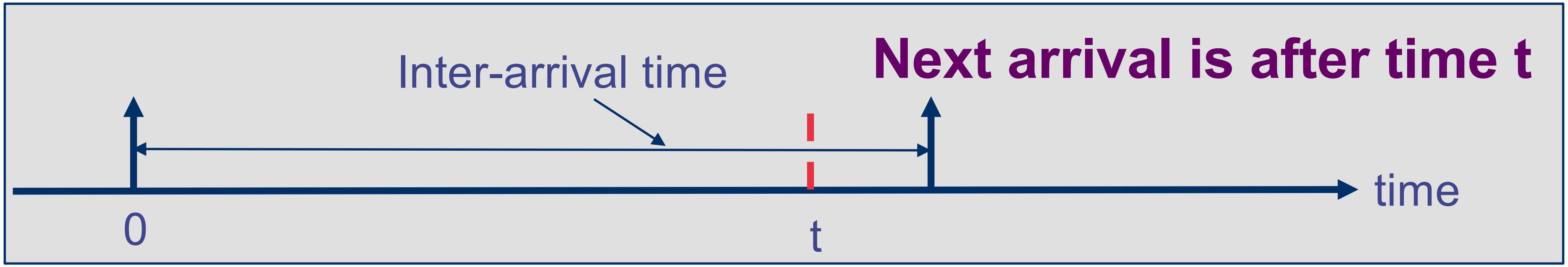


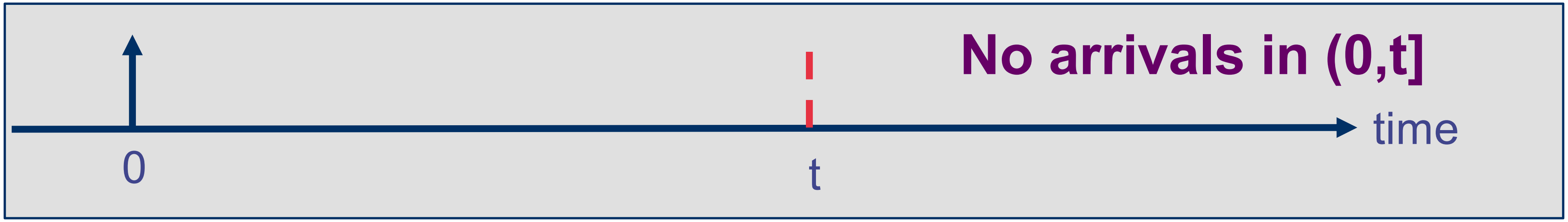


- Quiz: What is the relation between the following two probabilities?
  - Prob[at least one arrival in (0,t]]
  - Prob[no arrivals in (0,t]]
- Answer:
- Moral: ”No arrivals” is not boring, it tells you something

Inter-arrival probability

- Assume a customer arrives at time 0



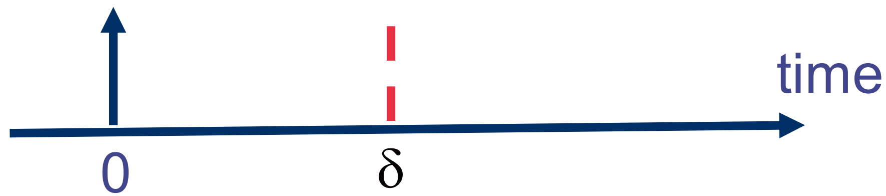


- Quiz: What is the relation between the following two probabilities?
  - Prob[Inter-arrival time is >= t]
  - Prob[no arrivals in (0,t]]
- Answer:
- Next step: Find Prob[no arrivals in (0,t]] for independent customers

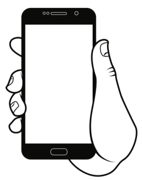
Many independent arrivals (1)

- Problem set up:
  - An arrival at time 0
  - A large pool of N independent customers
  - Behaviour of each customer: Within a small time interval of δ, a customer sends a request (or arrives) with a probability of $p\delta$
    - p is a constant

- Quiz: If there are 2 (= N) customers, what is the probability that both of them do not send any request in the time interval δ
  - Answer:



Customer 1



Prob[Send] = $p\delta$

Prob[Not send] = $1 - p\delta$

Customer 2

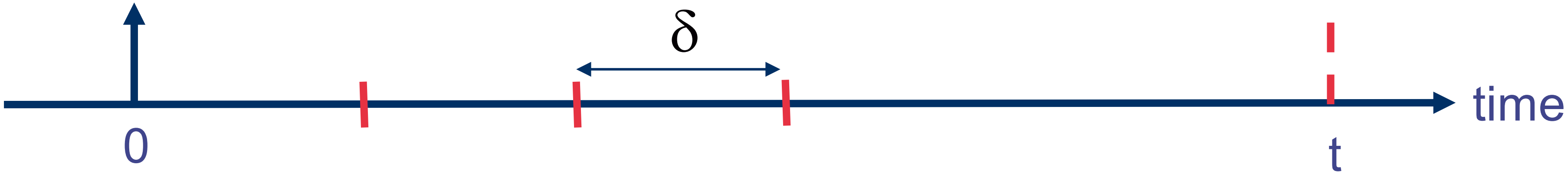
Prob[Send] = $p\delta$

Prob[Not send] = $1 - p\delta$

 Picture credit: openclipart.org

Many independent arrivals (2)

- Aim: Want to find the probability of no arrivals in (0,t]
- Divide the time t into intervals of width δ



- No arrival in (0,t] = no arrival in each interval δ from N users
- Probability of no arrival in δ =
- There are t / δ intervals
- Probability of no arrival in (0,t] is

$$(1 - Np\delta)^{\frac{t}{\delta}} \rightarrow e^{-Npt} \text{ as } \delta \rightarrow 0$$

Exponential inter-arrival time

- We have showed

  Probability(no arrival in (0,t]) = $\exp(-Npt)$

- Probability(inter-arrival time > t) = $\exp(-Npt)$

- This means

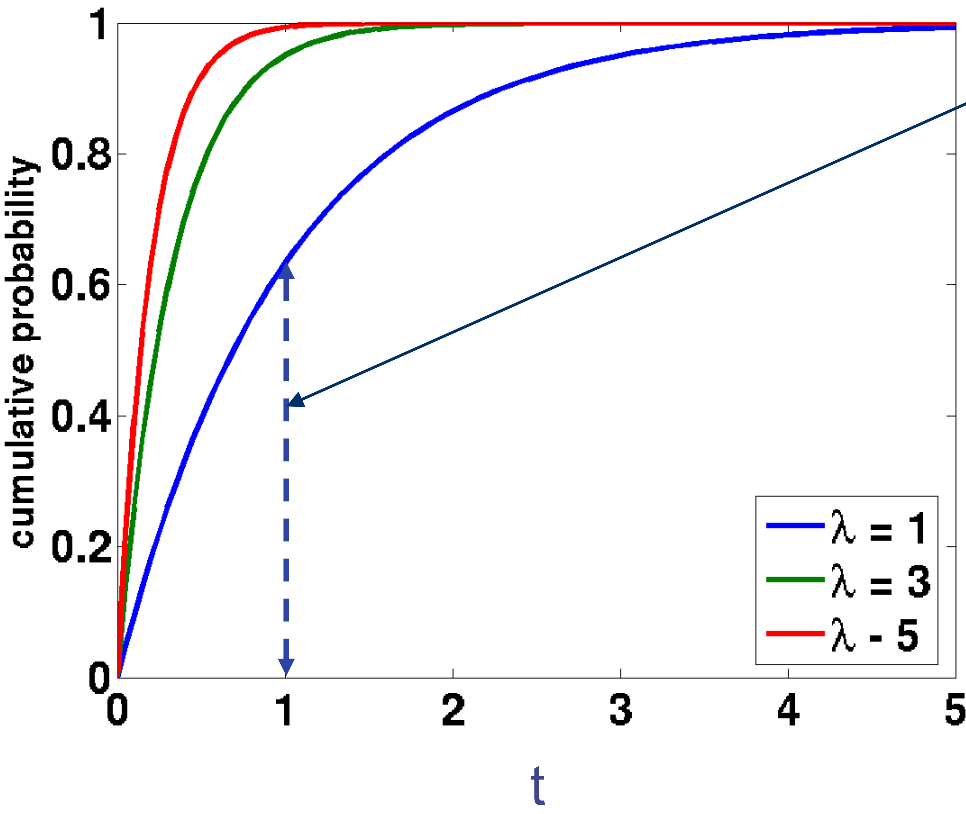
  Probability(inter-arrival time ≤ t) = $1 - \exp(-Npt)$

- What this shows is the inter-arrival time distribution for independent arrival is exponentially distributed

- Define: $\lambda = Np$
  - λ is the mean arrival rate of customers

Exponential distribution - cumulative distribution

- Cumulative distribution of inter-arrival time with customer arrival rate λ
  - Prob(inter-arrival time $\leq$ t) = 1 - exp(- λ t)



Prob that at least a customer will arrive before t = 1 for λ = 1

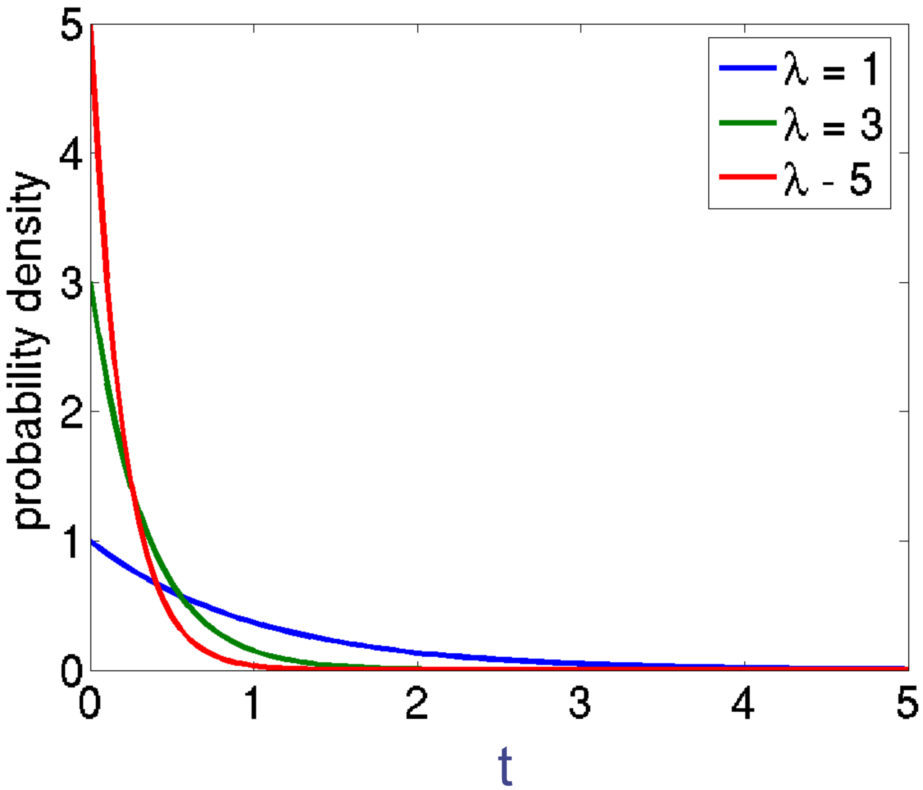
$\wedge$

Prob that at least a customer will arrive before t = 1 for λ = 3

Exponential distribution

- A continuous random variable is exponentially distributed with rate λ if it has probability density function

$$f(t) = \begin{cases} \lambda \exp(-\lambda t) & \text{for } t \geq 0 \\ 0 & \text{for } t < 0 \end{cases}$$



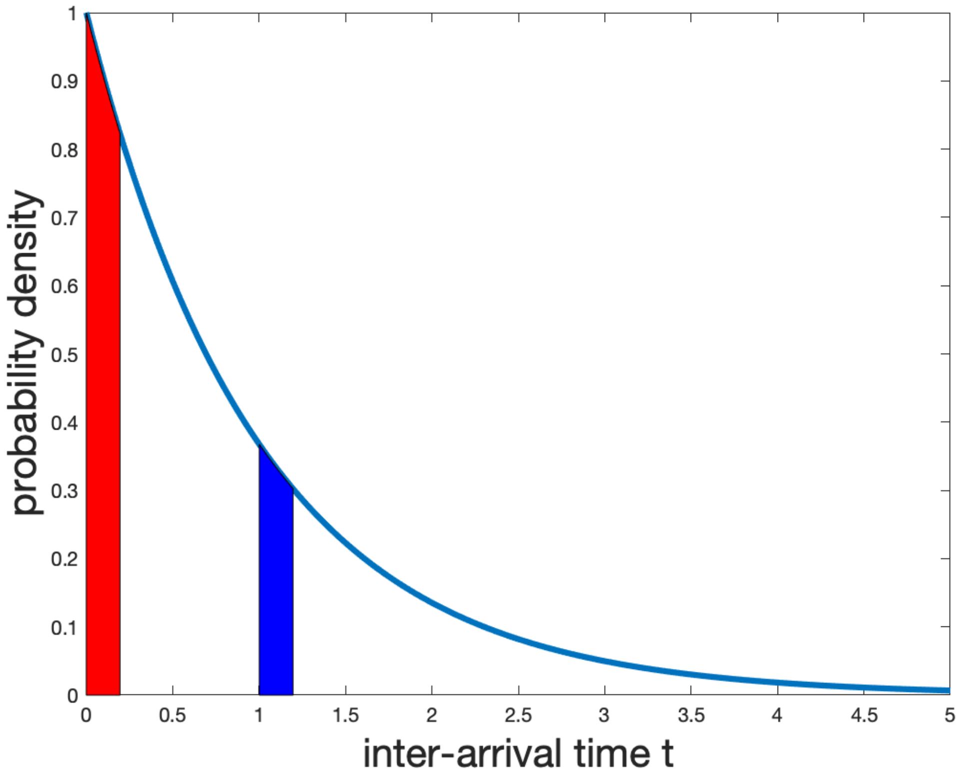
Reminder:

Probability density function $\xrightarrow{\text{integration}}$ Cumulative probability

Probability density function $\xleftarrow[\text{differentiation}]{}$ Cumulative probability

$$1 - \exp(-\lambda t)$$

Probability density function (PDF)



Reminder: PDF $f(t)$

Probability($t \leq T \leq t + \delta t$)

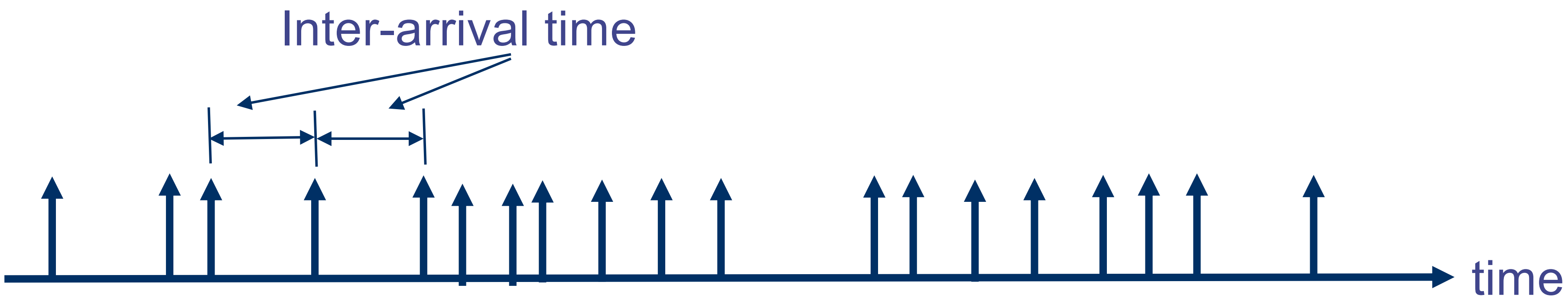
$= f(t)\ \delta t$

Red area = probability that inter-arrival time is in the interval [0,0.2]

Blue area = probability that the inter-arrival is in the interval [1,1.2]

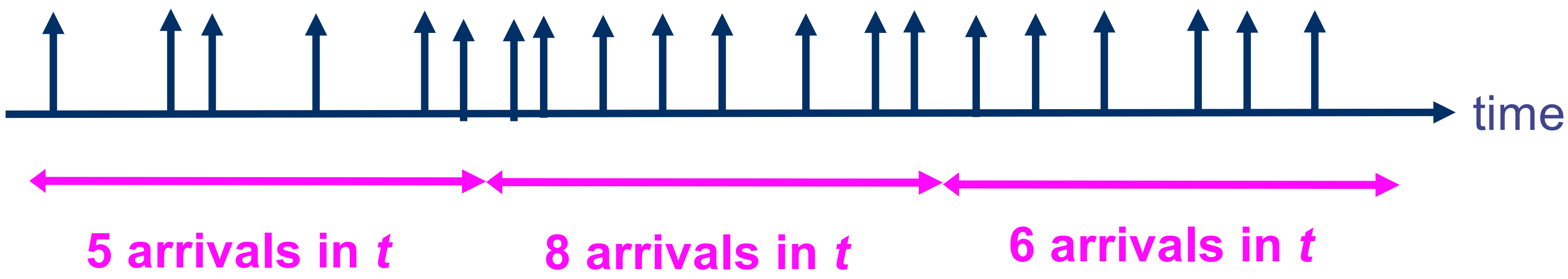
Two different methods to describe arrivals

Method 1: Continuous probability distribution of inter-arrival time



Two different methods to describe arrivals

Method 2: Use a fixed time interval (say *t*), and count the number of arrivals within *t*.

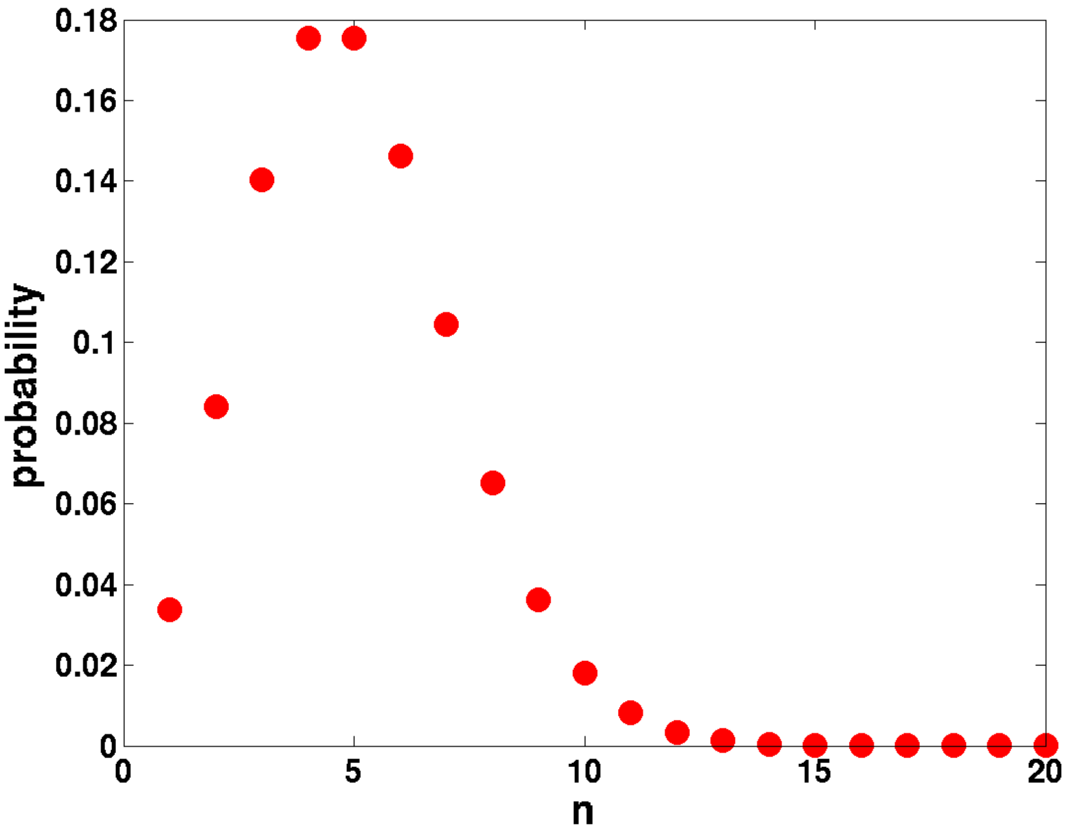


- The number of arrivals in *t* is random
- The number of arrivals must be a non-negative integer
- We need a discrete probability distribution:
  - Prob[#arrivals in t = 0]
  - Prob[#arrivals in t = 1]
  - etc.

Poisson process (1)

- Definition: An arrival process is Poisson with parameter λ if the probability that *n* customer arrive in any time interval *t* is

$$\frac{(\lambda t)^n e^{-\lambda t}}{n!}$$



Example:

Example:

λ= 5 and t = 1

Note: Poisson is a discrete probability distribution.

Poisson process (2)

- Theorem: An exponential inter-arrival time distribution with parameter λ gives rise to a Poisson arrival process with parameter λ

- How can you prove this theorem?
  - A possible method is to divide an interval t into small time intervals of width δ. A finite δ will give a binomial distribution and with $\delta \rightarrow 0$, we get a Poisson distribution.

Customer arriving rate

- Given a Poisson process with parameter λ, we know that the probability of n customers arriving in a time interval of t is given by:

$$\frac{(\lambda t)^n e^{-\lambda t}}{n!}$$

- What is the mean number of customers arriving in a time interval of t?

$$\sum_{n=0}^{\infty} n \frac{(\lambda t)^n e^{-\lambda t}}{n!} = \lambda t$$

- That's why λ is called the arrival rate.

Customer inter-arrival time

- You can also show that if the inter-arrival time distribution is exponential with parameter λ, then the mean inter-arrival time is $1/\lambda$
- Quite nicely, we have

  Mean arrival rate = 1 / mean inter-arrival time

Application of Poisson process

- Poisson process has been used to model the arrival of telephone calls to a telephone exchange successfully
- Queueing networks with Poisson arrival is tractable
  - We will see that in the next few weeks.
- Beware that not all arrival processes are Poisson! Many arrival processes we see in the Internet today are not Poisson. We will see that later.

References

- Operational analysis
  - Lazowska et al, Quantitative System Performance, Prentice Hall, 1984. (Classic text on performance analysis. Now out of print but can be download from http://www.cs.washington.edu/homes/lazowska/qsp/
    - Chapters 3 and 5 (For Chapter 5, up to Section 5.3 only)
  - Alternative 1: You can read Menasce et al, “Performance by design”, Chapter 3. Note that Menasce doesn’t cover certain aspects of performance bounds. So, you will also need to read Sections 5.1-5.3 of Lazowska.
  - Alternative 2: You can read Harcol-Balter, Chapters 6 and 7. The treatment is more rigorous. You can gross over the discussion mentioning ergodicity.

- Poisson process: Harcol-Balter Chapter 11