# 算法实验3

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### 问题设定

**0-1背包问题**: 假设有 n 个物品和一个容量为 C 的背包,和物品的重量  $\{w_i\}_{i\in[n]}$ ,以及物品的价值  $\{v_i\}_{i\in[n]}$ ,求解如何选择物品,才能使装进背包的物品价值总和最大。

本次实验我们要求使用以下方法解决0-1背包问题:

- 分治法 (Divide-and-conquer)
- 动态规划法 (Dynamic Programming)
- 贪心算法 (Greedy Method)
- 回溯法 (Back-Tracking Method)
- 分支限界法 (Branch-and-Bound Method)
- 蒙特卡洛算法 (Monte Carlo Method)

## 算法简述

• 分治法 (Divide-and-conquer)

在分治法中,我们对于"是否取用物品"建立一个二叉搜索树,每一层表示一个物品,每个节点的分叉表示选择将这个物品加入背包与否。对该二叉树进行深度优先搜索,每访问一个节点就会记录当前的路径,从而记录当前取用的物品,然后分别递归访问左右子树(若当前访问的物品重量大于背包剩余容量则仅访问右子树),记录得到分别得到的最大价值和对应取用的物品集合,再将这个最优值和最优集合返回。当访问至叶子节点时,将总记录和总价值返回,如此递归,当返回至根节点时,我们就得到了全局最优解。

• 动态规划法 (Dynamic Programming)

在动态规划法中,我维护一个两个二维数组 dp[0..n][0..c] 和 package\_list[0..n][0..c],元素 dp[i][j] 表示包含前 i 个物品,背包容量为 j 的字问题的最优价值,而 package\_list[i][j] 记录了对应的最优序列。当我们对第 [i][j] 个子问题进行访问时,我们进行以下操作:

- 。 若当前物品价值大于背包剩余容量,则该问题的解与第 [i-1][j] 个子问题相同;
- o 否则,研究将该物品装入背包后的总价值,该价值为 value[i] + dp[i 1][j weight[i]],将这个值与第 [i-1][j] 个子问题的最优值作比较,我们就可以得到最优解法是装入当前物品与否,然后进行对应的最优值和最优序列的记录即可。
- 贪心算法 (Greedy Method)

贪心算法原理非常简单,我们将物品按照单位价值进行排序,然后从前往后装进包里,直到装不下,即可获得一个解,但此解未必最优。

• 回溯法 (Back-Tracking Method)

回溯法在分治法的基础上进行了剪枝操作,我们维护一个序列 [rem\_sum[0..n]],其元素表示第 [i] 个物品及其之后的物品价值总和。我们访问到某一个节点时,计算当前已经装入的物品和当前访问物品及其之后物品的价值之和,并将其与目前记录的最优价值进行比较,如果比最有价值 小,那么我们将直接退出访问该节点,如此,我们实现了剪枝,简化了算法的复杂度。

• 分支限界法 (Branch-and-Bound Method)

在分支限界法中,我们维护一个堆 queue ,将我们访问过的节点按照最大上界进行排序,由此,当我们访问到某个节点,其可以达到的最优值比我们预估的最优值小,那么我们直接退出该节点,并且从 queue 中获取最大上界的节点,直接跳跃至该节点进行访问。

• 蒙特卡洛算法 (Monte Carlo Method)

蒙特卡洛算法是一种随机算法,我们直接将物品序列进行随机排列,然后从前往后逐个装入背包, 直到装不下为止。进行多次随机尝试,取其中价值最大的解作为全局解,由此我们可以得到一个 解,但未必是最优的。

### 代码展示

• 分治法 (Divide-and-conquer)

```
def divide_conquer_recursion(i=0, remain=C, optimal_list=None):
   if optimal_list is None:
        optimal_list = []
    if i == n:
        return 0, optimal_list
   if goods[i][0] > remain:
        return divide_conquer_recursion(i + 1, remain, optimal_list) #
cannot putting it in, so we visit next
    else: # can put it in
        not_put_it_in, not_in_list = divide_conquer_recursion(i + 1, remain,
                                                              optimal_list)
 # not putting it in, so we visit next
       optimal_list_copy = optimal_list.copy()
        optimal_list_copy.append(i)
        later_values, in_list = divide_conquer_recursion(i + 1,
                                                         remain - goods[i]
[0], optimal_list_copy) # putting it in
        put_it_in = goods[i][1] + later_values
        if put_it_in > not_put_it_in:
            return put_it_in, in_list
        else:
            return not_put_it_in, not_in_list
```

动态规划法 (Dynamic Programming)

```
def dynamic_method():
    package_list = [[[] for j in range(C + 1)] for i in range(n + 1)]
    # dp[i][j] means the max value we can get if we put the first i goods in
the bag with capacity j
    dp = [[0 for j in range(C + 1)] for i in range(n + 1)]

for i in range(1, n + 1): # i is the number of goods
    for j in range(1, C + 1): # j is the capacity
        if goods[i - 1][0] > j: # if the weight of the current good is
larger than the capacity
        package_list[i][j] = package_list[i - 1][j].copy() # we can
not put the good in
```

```
dp[i][j] = dp[i - 1][j] # we can only take the previous one
           else: # if the weight of the current good is smaller than the
capacity
               put_it_i = goods[i - 1][1] + dp[i - 1][j - goods[i - 1][0]]
# we put the good in
               not_put_it_in = dp[i - 1][j] # we don't put the good in
               if put_it_in > not_put_it_in: # if we can put the good in
                   package_list[i][j] = package_list[i - 1][j - goods[i -
1][0]].copy() # we put the good in
                   package_list[i][j].append(i - 1) # we put the good in
                   dp[i][j] = put_it_in # we put the good in
               else: # if we can't put the good in
                   package_list[i][j] = package_list[i - 1][j] # we don't
put the good in
                   dp[i][j] = not_put_it_in # we don't put the good in
   return dp[n][C], package_list[n][C] # return the optimal value and the
optimal list
```

• 贪心算法 (Greedy Method)

```
def greedy_method():
   optimal_list = []
   goods_copy = goods.copy() # copy the goods list
   goods_copy.sort(key=take_unit_value, reverse=True) # sort the goods by
the unit value
   capacity = C # the capacity is the same as the bag
   value_all = 0 # the value of all the goods taken into the bag
   for good in goods_copy: # for each good
       if capacity >= good[0]: # if the capacity is larger than the weight
of the good
           capacity -= good[0] # reduce the capacity
           value_all += good[1] # add the value of the good
           optimal_list.append(good[2]) # add the index of the good to the
optimal list
   return value_all, optimal_list # return the optimal value and the
optimal list
```

回溯法(Back-Tracking Method)

```
def back_tracking():
    goods_copy = goods.copy() # copy the goods list
    goods_copy.sort(key=take_unit_value, reverse=True) # sort the goods by
the unit value
    rem_sum = [0] # the sum of the remaining value of the goods

for i in range(n-1, -1, -1): # for each good
    rem_sum.insert(0, rem_sum[0] + goods_copy[i][1]) # add the weight
of the good to the sum
```

```
# level is the level of the recursion,
    # cur_weight is the weight of the bag,
    # cur_val is the value of the bag,
    # optimal is the optimal value
    def dfs(level, cur_weight, cur_val, optimal, package_list):
        if level == n:
            return max(optimal, cur_val), package_list
       good = goods_copy[level]
       flag = 0
       package_list_copy = package_list.copy()
       if cur_weight + good[0] <= C: # if the weight of the bag is larger</pre>
than the weight of the good
            package_list_copy.append(good[2]) # add the index of the good
to the optimal list
            optimal, left_list = dfs(level + 1, cur_weight + good[0],
                                     cur_val + good[1], optimal,
package_list_copy) # we can put the good in
            left_optimal = optimal
            optimal_list = left_list
            flag = 1
       # if the value of the bag is larger than the sum of the remaining
value of the goods
       if cur_val + rem_sum[level + 1] > optimal:
            # we can not put the good in
            optimal, right_list = dfs(level + 1, cur_weight, cur_val,
optimal, package_list)
            if flag == 0:
                return optimal, right_list
            flag = 2
       if flag == 0:
            return optimal, package_list
        elif flag == 1:
            return optimal, left_list
        else:
            if left_optimal >= optimal:
                return left_optimal, left_list
            else:
                return optimal, right_list
    return dfs(0, 0, 0, 0, [])
```

• 分支限界法 (Branch-and-Bound Method)

```
def branch_and_bound_method():
    goods_copy = goods.copy() # copy the goods list
    goods_copy.sort(key=take_unit_value, reverse=True) # sort the goods by
the unit value

    optimal = 0
    queue = [(0.0, 0, 0, 0, [])]
    heapq.heapify(queue) # min-root heap
```

```
return_list = []
   while queue: # while the queue is not empty
        item = heapq.heappop(queue) # get the item with the minimum value
       upper_bound, index, cur_weight, cur_val, cur_list = item[:] # get
the upper bound, the index of the good, the weight of the bag, and the value
of the bag
       upper_bound = - upper_bound # get the upper bound
       if optimal < cur_val: # if the upper bound is larger than the
optimal value
           optimal = cur_val
            return_list = cur_list
        if index == n or upper_bound < optimal: # if the index is the last</pre>
good or the upper bound is smaller than the optimal value
           continue # we can not put the good in
       weight, value, identity = goods_copy[index] # get the weight and
value of the good
       upper_bound = cur_val + (C - cur_weight) * value / weight # get the
upper bound
       if cur_weight + weight <= C: # if the weight of the bag is larger
than the weight of the good
           cur_list_copy = cur_list.copy() # copy the optimal list
           cur_list_copy.append(identity) # add the index of the good to
the optimal list
           heapq.heappush(queue, (-upper_bound, index + 1, cur_weight +
weight,
                                  cur_val + value, cur_list_copy)) # put
the good in
       heapq.heappush(queue, (-upper_bound, index + 1, cur_weight, cur_val,
cur_list)) # not put the good in
    return optimal, return_list
```

• 蒙特卡洛算法 (Monte Carlo Method)

### 测试结果

对于不同 n 和 C, 进行如下尝试:

```
number of goods: 20
Content of package: 950
Max weight: 200
max_value: 100
divide_conquer took 0.4070272445678711 seconds
Optimal value: 732
optimal list: [0, 1, 3, 5, 6, 10, 11, 14, 15, 16, 18]
dynamic_method took 0.024004220962524414 seconds
Optimal value: 732
optimal list: [0, 1, 3, 5, 6, 10, 11, 14, 15, 16, 18]
greedy_method took 0.0 seconds
Optimal value: 732
optimal list: [5, 11, 3, 0, 10, 18, 1, 6, 16, 15, 14]
back_tracking took 0.00400233268737793 seconds
Optimal value: 732
Optimal list: [5, 11, 3, 0, 10, 18, 1, 6, 16, 15, 14]
branch_and_bound_method took 0.0010020732879638672 seconds
Optimal value: 732
Optimal list: [5, 11, 3, 0, 10, 18, 1, 6, 16, 15, 14]
monte_carlo_method took 2.0482213497161865 seconds
Optimal value: 715
Optimal list: [16, 6, 11, 3, 10, 15, 14, 5, 1, 18]
```

```
number of goods: 10
Content of package: 950
Max weight: 200
max_value: 100
divide_conquer took 0.0010004043579101562 seconds
Optimal value: 468
Optimal list: [0, 1, 2, 3, 4, 6, 7, 8, 9]

dynamic_method took 0.013999700546264648 seconds
Optimal value: 468
Optimal list: [0, 1, 2, 3, 4, 6, 7, 8, 9]

greedy_method took 0.0 seconds
```

Optimal value: 468

optimal list: [0, 4, 3, 6, 2, 1, 8, 9, 7]

back\_tracking took 0.0 seconds

Optimal value: 468

optimal list: [0, 4, 3, 6, 2, 1, 8, 9, 7]

branch\_and\_bound\_method took 0.0009989738464355469 seconds

Optimal value: 468

Optimal list: [0, 4, 3, 6, 2, 1, 8, 9, 7]

monte\_carlo\_method took 1.4659490585327148 seconds

Optimal value: 468

optimal list: [7, 9, 0, 8, 2, 1, 4, 6, 3]