

Johns Hopkins Engineering

625.464 Computational Statistics

Squeezed Rejection Sampling

Module 4 Lecture 4D



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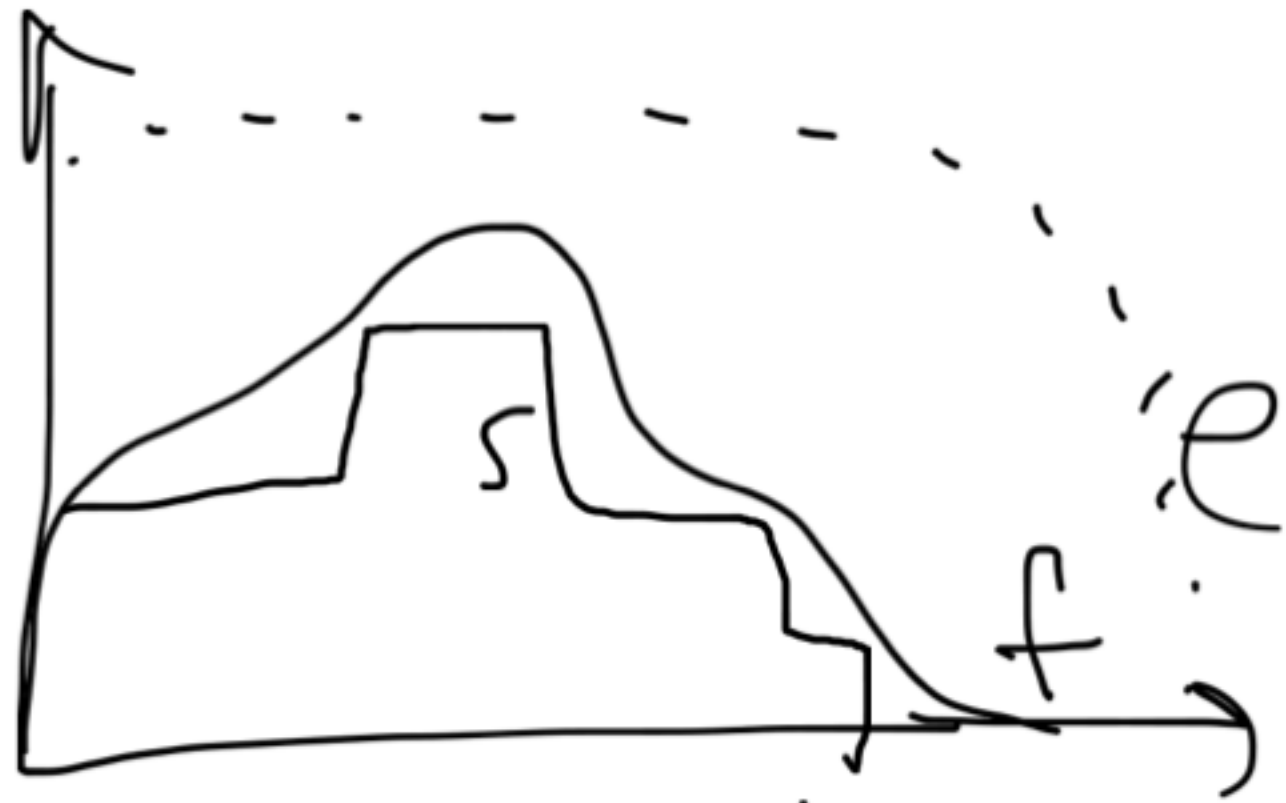
Why Squeezed Rejection Sampling?

Ordinary rejection sampling

- ① Sample $y \sim q$
- ② Sample $u \sim U(0,1)$
- ③ If $u > f(y)/e(y)$ reject
- ④ o.w. keep.

Squeezed Rejection Sampling

The Basic Idea



- ① $S(x) \leq f(x)$
- ② easy to compute

- ④ check if $u \leq \frac{f(y)}{e(y)}$
if so keep $x = y$
→ go to 6

- Squeezed rejection
- ① Sample $y \sim q$
 - ② Sample $u \sim u(0,1)$
 - ③ If $u \leq \frac{S(y)}{e(y)}$ Keep y
Set $y = x \rightarrow$ go to ⑥

- ⑤ Reject & Start over

- ⑥ go to 1 until done

Pseudo code for Squeezed Rejection Sampling

Pseudo code for Squeezed Rejection Sampling.

Suppose you want to generate n random numbers from distribution f using S.R.S. w/ S, g and $e = g/\alpha$.

Pseudocode

```
- Initialize constants  $n, \alpha$ 
- Initialize variables  $y[n] = \alpha$ ; accepted = 0
- While (accepted < n) {
     $x = \text{r.n. from } g$  (sample  $g$ )
     $u = \text{r.n. from } U(0,1)$  (sample  $U(0,1)$ )
     $S = S(x)$  (eval  $S$  at  $x$ )
     $e = g(x)/\alpha$  (eval  $e$  at  $x$ )

    If ( $u \leq S/e$ ) {
        accepted = accepted + 1
         $y[\text{accepted}] = x$ 
    }
    Else {
         $f = f(x)$ 
        If ( $u \leq f/e$ ) {
            accepted = accepted + 1
             $y[\text{accepted}] = x$ 
        }
    }
}
Return  $y$ 
```