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Standard

Section

Roll No.

Subject

Algorithm for tic tac toe

Step 1 : define print_board(board)

- * Create a 2 D list array to representing the Tic tac toe grid
- * Join the element of each row with " | "
- * print the separate line of dashes " --- " after each row for clarity

Step 2 : define check_winner(board)

```
if row[0] == row[1] == row[2] != " "
    return row[0]
```

→ * This checks each row to see if all 3 cells contain same symbol (x or o) and if empty if winner is found , it returns winning symbol.

→ * if board[0][col] == board[1][col] == board[2][col]
 i = " ";
 return board[0][col].

→ This checks each column for 3 matching symbol , returning winning symbol if a match is found.

→ * if board[0][0] == board[1][1] == board[2][2] != " "
 return board[0][0]

This checks main diagonal for 3 matching symbol & returns if it is found.

→ if board[0][2] == board[1][1] == board[2][0] != " "
 return board[0][2]

This checks anti-diagonal for winning symbol & returns it if found

→ return none

This checks if no more winning condition is met after all checks. It returns none. It indicates there is no winner.

Step 4 :- define tic tac toe()

- * Initialize 3×3 board with empty space
- * Set current player to x
- * Prompt the current player for row and column input ($0, 1, 2$)

Step 5 :- if checks if selected cell is empty
if occupied print an error message & restart it

if $\text{board}[\text{row}][\text{col}] == "$ "
print ("An invalid move")

Step 6 :- calls `check_winner(board)` to check for winning condition

- * if a winner is found, prints the board and announces the winner, then exists the function

* $\text{current_player} = '0'$ if $\text{current_player} = 'x'$
else " x "

→ Switches the current player for the next turn

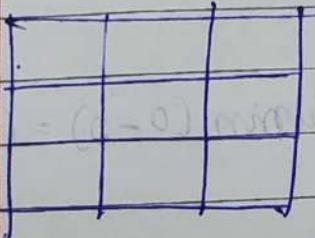
- * After 9 turns into if no winner prints

24/9/2020 the board and declares a draw

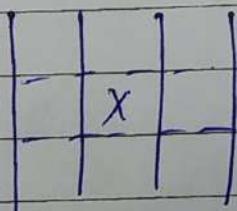
Step 7: if name == "main"
 tic tac toe()

* calls the tic tac toe in main function
& then start the game.

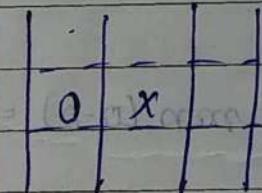
O/P



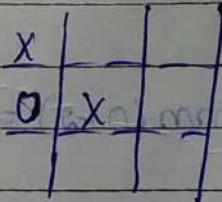
Player X: enter row and column (0-2) = (1, 1)



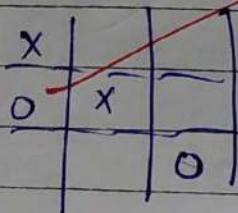
Player O enters row & column (0-2) = (1, 0)



Player X enters row & column (0-2) = (0, 0)



Player O enters row & column (0-2) = (2, 2)



player
enter the x : enter row and column $(0-2) = (0, 2)$

x	-	x
0	x	0

player 0 : enter row & column $(0-2) = (0, 1)$

x	x	x
0	x	0

x is winning

O/P

-	-	-
(0, 1)	-	-

Player x : enter row and column $(0-2) = (0, 1)$

-	x	-
(0, 2)	-	-

player y : enter row and column $(0-2) = (0, 3)$

Invalid move

O/P

-	-	-
-	-	-

player X : Enter row and column (0-2) = (0, 1)

-	X	-
-	-	-
-	-	-

X	X	O
O	O	X
X	O	O

player O : Enter row and column (0-2) = (0, 0)

O	X	-
-	-	-
-	-	-

1	0	1
0	1	0
1	0	0

player X : Enter row and column (0-2) = (0, 2)

O	X	X
-	-	-
-	-	-

player O : Enter row and column (0-2) = (1, 0)

O	X	X
-	O	-
-	-	-

player X : Enter row and column (0-2) = (2, 2)

O	X	X
-	O	O
-	-	X

player O : Enter row and column (0-2) = (1, 2)

O	X	X
-	O	O
-	-	X

player X : Enter row & column (0-2) = (1, 0)

O	X	X
X	O	O
-	-	X

player O : Enter row & column (0-2) = (2, 1)

O	X	X
X	O	O
-	O	X

Player 0, enter row and column (0-2) = (2, 0)

0	x	x
x	0	0
0	0	x

No winner

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9/19/2024

Lab 2: Implement vacuum world cleaner

function REFLEX-VACUUM-AGENT (location, status)
return an action
if status = Dirty then return suck
else if location = A then return right
else if location = B then return left

Algorithm for two quadrants

Step 1:- Initialization

- * Input current room (either A or B)
- * Input the status of the room clean or dirty
- * Initialize a cost to 0

Step 2: Display initial room status

Step 3: cleaning loop

* while either of the room is dirty:

- ① if current room is A and is dirty then clean & increase cost by one
- ② If current room is B and is dirty then clean and increase cost by one
- ③ If current room is not dirty move to next room

Step 4: Display cost

Step 5: Stop

Output:

Enter current room either A or B : A

Is room A dirty ? (Yes : 1 / No : 0) : 0

Is room B dirty ? (Yes : 1 / No : 0) : 0

Initial status of rooms :

Room A : clean

Room B : Dirty

Moving to Room B...

current status:

Room A : clean

Room B : Dirty

Cleaning Room B...

Current status

Room A : clean

Room B : clean

Both rooms are now clean! Total cost: 1

Output:

Enter current room either A or B : A

Is Room A dirty? (Yes:1/No:0) : 1

Is Room B dirty? (Yes:1/No:0) : 0

Initial status of rooms:

Room A : Dirty

Room B : Clean

Cleaning Room A...

Current status:

Room A : clean

Room B : clean

Both rooms are now clean! Total cost: 1

Output:

Enter current room either A or B : A

Is Room A dirty? (Yes: 1 No: 0) : 1

Is Room B dirty? (Yes: N NO: 0) : 1

Initial status of rooms:

Room A : Dirty

Room B : Dirty

Cleaning Room A ...

current status:

Room A : clean

Room B : Dirty

Moving to Room B ...

current status:

Room A : clean

Room B : Dirty

Cleaning Room B ...

current status:

Room A : clean

Room B : clean

Algorithm for Four Quadrants

Step 1: Start

* Accept status of each room (either clean or dirty)

Step 2: Display the initial status of all rooms

Step 3: Cleans rooms

* While all rooms are clean

① If first room is dirty clean and go to next room

② If second room is dirty clean and go to next room

③ If third room is dirty clean and go to next room

④ If fourth room is dirty clean and go to next room

⑤ Increase cost when cleaning and decreases count

Step 4: Show the last status

Step 5: End

Output:

Enter current room (A, B, C, D) : B

Is room A dirty ? (Yes : 1 / No : 0) : 1

Is room B dirty ? (Yes : 1 / No : 0) : 0

Is room C dirty ? (Yes : 1 / No : 0) : 0

Is room D dirty ? (Yes : 1 / No : 0) : 0

Initial status of rooms :

Room A : Dirty Room B : clean Room C : clean

Room D : clean

Moving to Room C ...

Moving to Room D ...

Moving to Room A ...

Cleaning Room A ...

Current status

Room A : clean Room B : clean

Room C : clean Room D : clean

All rooms are now clean : Total cost : 1

Output:

Enter current room (A, B, C, D) : A

Is Room A dirty ? 1

Is Room B dirty ? 1

Is Room C dirty ? 1

Is Room D dirty ? 1

Initial rooms status:

Room A : Dirty Room B : Dirty Room C : Dirty
Room D : Dirty

Cleaning Room A ...

Moving to Room B ...

Cleaning Room B ...

Vaccum is recharging

Moving to Room C ...

Cleaning Room C ...

Moving to Room D ...

Cleaning to Room D ...

Current status :

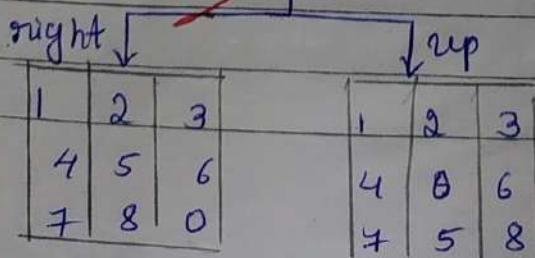
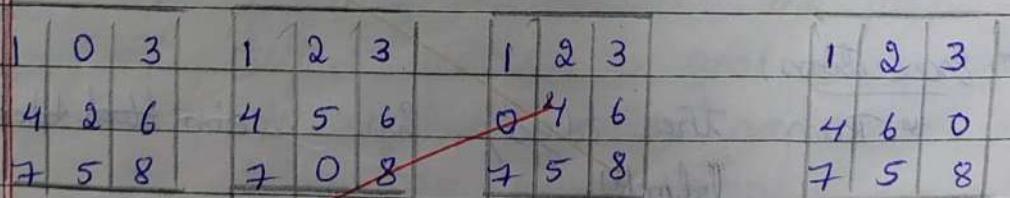
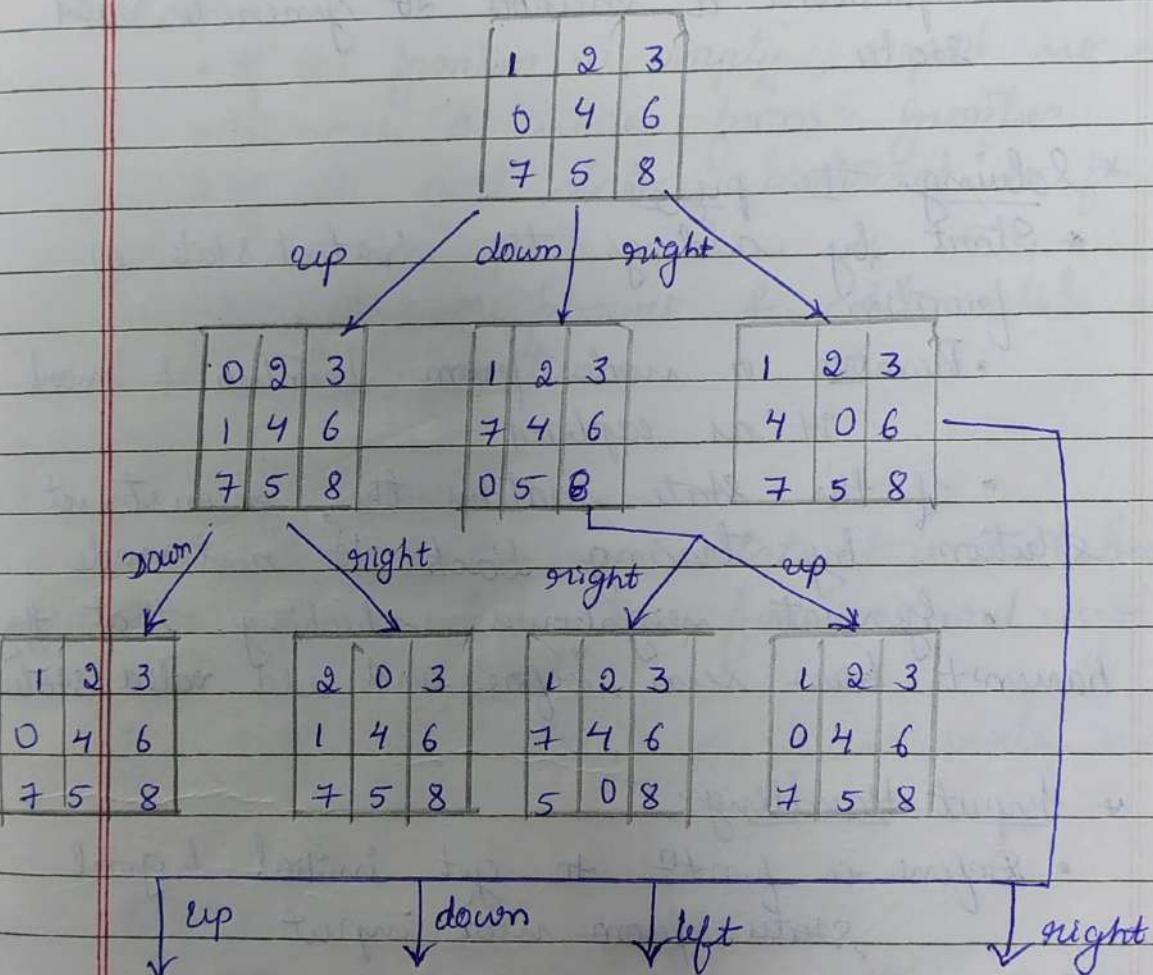
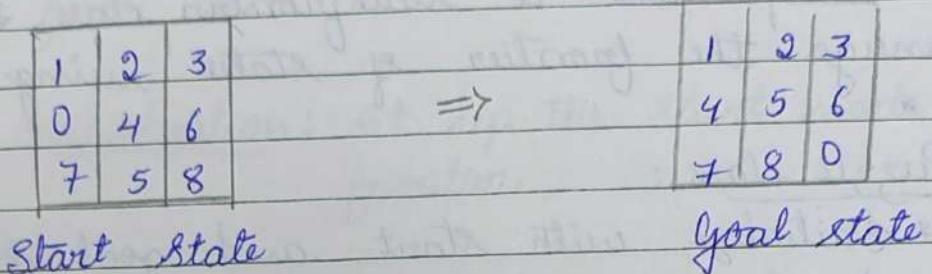
Room A : clean Room B : clean Room C : clean
Room D : clean

All rooms are now clean ! Total cost : 4

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9/10/2024

Solve our 8 puzzle problems using DFA
and BFS

State Space diagram:



Algorithm for 8 puzzle problems by using BFS

* Node Stack frontier :

- * Implement a stack frontier class to manage the frontier of states using stack.
- *

* Puzzle class :

- * Initialize with start and goal states
- * Implement a method to generate valid states

* Solving the puzzle :

- * Start by adding the initial state of frontier
 - Remove a node from frontier & mark it as explored
 - If the state matches the, reconstruct solution by tracing back to root node
 - Generate neighbours, checking that they haven't been seen before and add valid state

* Input handling :

- Define a funtn to get initial & goal states from user input

* Execution :

- * Run the algo, in main block

Algorithm for 8 puzzle problems by using DFS

- * Input: Read the start and goal states from user
- * Initialization: set up the start node and frontier.
- * Loop:
 - If the frontier is empty, report no solution
 - Remove a node from frontier
 - If its goal state, backtrack to find the solution path.
 - Generate neighbours & add valid new states to frontier.
- * Solution process:
 - check if the new state has not been seen before [not in the frontier and not explored].
 - If valid create a new node and add it to frontier.
- * Output:
 - After finding a solution, the print solution method output initial & goal states. the number of explored states and series of moves taken to reach goal.

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10/2024

Lab - 4

15/10/23

For 8 puzzle problem using A* star implementation to calculate $f(n)$ using

@ $g(n)$ = depth of a node

$h(n)$ = Heuristic value



no. of misplaced tiles

$$f(n) = g(n) + h(n)$$

⑥ $g(n)$ = depth of a node

$h(n)$ = Heuristic value



Manhattan distance

$$f(n) = g(n) + h(n)$$

⑦ Number of misplaced tiles

Draw the state space diagram for

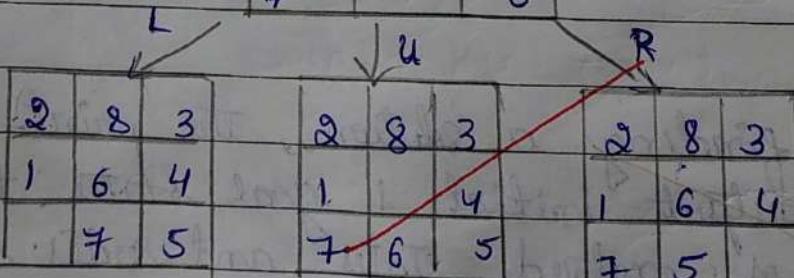
2	8	3
1	6	4
7	.	5

1	2	3
8		4
7	6	5

$$g(n) = 0 \quad h(n) = 4.$$

goal state

9	8	3
1	6	4
7		5



$$g(n) = 1$$

$$h(n) = 5$$

$$f(n) = 1+5 = 6$$

$$g(n) = 1$$

$$h(n) = 3$$

$$f(n) = 4$$

$$g(n) = 1$$

$$h(n) = 5$$

$$f(n) = 6$$

2	8	3
1		4
7	6	5

L

d

R

2	8	3
1	4	
7	6	5

2	8	3
1	6	4
7	5	

2	8	3
1	4	
7	6	5

2	3
1	8
7	6

$$g(n) = 2$$

$$h(n) = 3$$

$$f(n) = 5$$

$$g(n) = 2$$

$$h(n) = 4$$



$$g(n) = 2$$

$$h(n) = 4$$

$$f(n) = 6$$

$$g(n) = 2$$

$$h(n) = 3$$

$$f(n) = 5$$

Don't explore

it is same as

Start State

L	2	8	3
	1	4	
	7	6	5

2	3
1	8
7	6

d

R

U

L

d

on

2	8	3
7	1	4
6	5	

2	8	3
1	4	
7	6	5

2	8	3
1	4	
7	6	5

3	3
1	8
7	6

2	8	3
1	4	
7	6	5

2	3
1	8
7	6

$$g(n) = 3$$

$$h(n) = 4$$

$$f(n) = 7$$

$$g(n) = 3$$

$$h(n) = 3$$

$$f(n) = 6$$

$$g(n) = 3$$

$$h(n) = 2$$

$$f(n) = 5$$

$$g(n) = 3$$

$$h(n) = 5$$

$$f(n) = 8.$$

Don't explore

it already
exist

2	3
1	8
7	6

2	3
1	8
7	6

1	2	3
8	4	
7	6	5

$$g(n) = 4$$

$$A(n) = 3$$

$$f(n) = 7$$

$$g(n) = 4$$

$$A(n) = 1$$

$$f(n) = 5$$

	1	2	3	
	8	4		
	7	6	5	

a) n d.

	2	3		1	2	3		1	2	3	
	1	8	4	8		4		7	8	4	
	7	6	5	7	6	5		7	6	5	

$g(n) = 5$ $g(n) = 5$ $g(n) = 5$
 $h(n) = 2$ $h(n) = 0$ $h(n) = 1$
 $f(n) = 7$ $f(n) = 5$ $f(n) = 6$.

↓

Goal State

Algorithm :

- * Defined node structure: Each node should have
 - * The current state of puzzle
 - * The cost to reach to its position
 - * $f(n)$ means cost of estimation
 - * $g(n)$ means level of each step
 - * $h(n)$ means heuristic values.
- * The 8 puzzle problem using the A* Star Algorithm
 - $f(n) = h(n) + g(n)$
- * Consider first initial state and goal state as level zero
 $g(n) = 0$.
- * Consider the initial state and move the number possible into empty space. (like left, right, up, down direction only)

- * In each step consider the each levels and find heuristic value.
- * Compare the 8 puzzle table with goal state compare @ each cell if it doesn't match consider as 1 and add the 1 How much cells of elements does not match with cell of goal state . calculate the all counts as heuristic value.
- * $f(n)$ is cost estimation of that puzzle it has formula $f(n) = h(n) + g(n)$ followed it which one has minimum value of $h(n)$. [heuristic value] that will be consider in next level , it goes on like this till reaching the goal state
- * If any puzzles' already exist in previous level . then don't explore those puzzle

(b) Manhattan distance

2	8	3
1	6	4
7	5	

1	2	3
8	4	
7	6	5

Initial state

$$g(n) = 0$$

$$h(n)$$

goal state

$$1 \ 2 \ 3 \ 4 \ 5 \ 6 \ 7 \ 8 \\ 1+1+0+0+0+1+0+1=4$$

2	8	3
1	6	4
7	5	

u

n

2	8	3
1	6	4
7	5	

2	8	3
1	6	4
7	6	5

2	8	3
1	6	4
7	5	

$$g(n) = 1$$

$$h(n) = 1+1+1+1+2$$

$$h(n) = 6$$

$$f(n) = 7$$

$$g(n) = 1$$

$$h(n) = 1+1+2$$

$$h(n) = 4$$

$$f(n) = 5$$

$$g(n) = 1$$

$$h(n) = 1+1+1+1+2 = 6$$

$$f(n) = 7$$

2	1	8	3
1			4
7		6	5

2	1	8	3
1			4
7		6	5

2	1	8	3
1			4
7		6	5

L

u

d

n

2	8	3
1	4	
7	6	5

2	8	3
1	8	4
7	6	5

2	8	3
1	8	4
7	6	5

$$g(n) = 2$$

$$h(n) = 2+1+2$$

$$= 5$$

$$f(n) = 7$$

$$g(n) = 2$$

$$h(n) = 1+1+1$$

$$= 3$$

$$f(n) = 5$$

2	8	3
1	6	4
7	5	

2	8	3
1	4	
7	6	5

2	8	3
1	4	
7	6	5

L

u

d

n

$$g(n) = 2$$

$$h(n) = 1+1+2$$

$$= 5$$

$$f(n) = 7$$

$$g(n) = 2$$

$$h(n) = 1+1+2$$

$$= 5$$

$$f(n) = 7$$

2	3
1	8 4
7 6 5	

g1 L d

2 3	2 3	2 8 3
1 8 4	1 8 4	1 4
7 6 5	7 6 5	7 6 5

$$g(n) = 3$$

$$h(n) = 1+1=2$$

$$f(n) = 5$$

$$g(n) = 3$$

$$h(n) = 1+1+1=3$$

$$f(n) = 6$$

$$g(n) = 3$$

$$h(n) = 1+1+2=4$$

$$f(n) = 7$$

2	3
1	8 4
7 6 5	

d

2 3	1 2 3
1 8 4	8 4
7 6 5	7 6 5

$$g(n) = 4$$

$$h(n) = 1+1+1=3$$

$$f(n) = 7$$

$$g(n) = 4$$

$$h(n) = 1$$

$$f(n) = 5$$

1	2	3
8	4	
7	6 5	

d

1 2 3	1 2 3
8 4	7 8 4
7 6 5	6 5

$$g(n) = 5$$

$$h(n) = 0$$

$$f(n) = 5$$

$$g(n) = 5$$

$$h(n) = 1+1=2$$

$$f(n) = 7$$

↑
goal state

Algorithm :

- * Defined node structure each node should have the current state of puzzle

$g(n)$ means cost estimation of estimation

$h(n)$ means heuristic values

$$f(n) = g(n) + h(n)$$

$g(n)$ means level of each fuzz step

- * The 8 puzzle problem using the A* Algorithm

$$f(n) = g(n) + h(n)$$

- * consider first initial state and goal

state as level zero

$$g(n) = 0$$

- * consider the initial state. move the possible nodes into empty space.

(like left, right, up, down direction)

- * In Manhattan distance , calculate heuristic value as comparing moves of each node with go each node of goal state.

- * we count the each movements as two heuristic value and $f(n)$ calculate the $f(n) = g(n) + h(n)$

- * Consider the minimum value of $h(n)$ to calculate the 8 puzzles problem

- * It goes iteration till it reaches the goal state.
- * If any puzzle already exist in previous level. then we don't need to explore that state.

Output: of misplaced tiles

Initial state:

(2, 8, 3)

(1, 6, 4)

7 0 5

goal state:

1 2 3

8 0 4

7 6 5

Solution found with cost: 5

steps:

2 8 3

1 6 4

7 0 5

2 8 3

1 0 4

7 6 5

2 0 3

1 8 4

7 6 5

0 2 3

1 8 4

7 6 5

.

1 2 3

8 0 4

7 6 5

Output for Manhattan

Enter Initial state

2 8 3

1 6 4

7 0 5

Enter goal state

1 2 3

8 0 4

7 6 5

level 1

2 8 3

1 0 4

7 6 5

level 1

2 8 3

1 6 4

0 7 5

level 1

2 8 3

1 6 4

7 5 0

level : 2

2 8 3

1 6 4

7 5 0

level : 2

2 8 3

1 6 0

7 5 4

level : 2

2 0 3

1 8 4

7 6 5

level : 3

2 3 0

1 8 4

7 6 5

level 4

2 0 3

1 8 4

7 6 5

level 5

1 2 3

8 0 4

7 6 5

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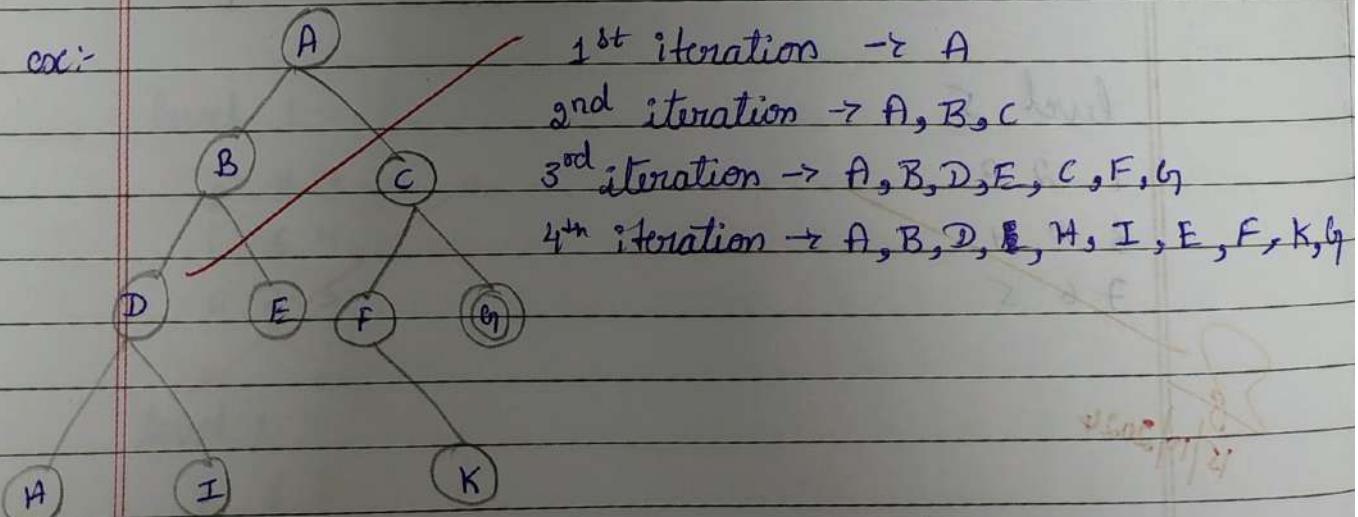
Implement Iterative Deepening Search Algorithm

Algorithm:-

function Iterative-Deeping-Search(problem) returns
 a solution , if failure
 for depth = 0 to ∞ do
 result \leftarrow Depth-limited-Search(problem,
 depth)
 if result \neq cutoff then return result

1. For each child of current node.
2. If it is target node, return
3. If the current maximum depth is reached, return
4. Set the current node to this node & go back to step 1
5. After having gone through all children, go to the next node child of the parent (next sibling)
6. After having gone through all children of start node, increase the maximum depth & go back to 1
7. If we have reached all leaf (bottom) nodes, the goal node doesn't exist.

ex:-



Initial
goal state

1	4	3
7		6
5	8	2

level 0

goal state

1	4	3
7	6	2
5	8	

up left down right

1	3	1	4	3	1	4	3	1	4	3
7	4	6		7	5	7	8	6	7	6
5	8	2	5	8	2	5		2	5	8

left right up down left right

1	3	1	3	4	3	1	4	3	1	4	3
7	4	6	7	4	6	1	7	6	7	6	7
5	8	2	5	8	2	5	8	2	5	2	5

level 2

up down

7	9	1	4	3
7	6	3	7	6
5	8	2	5	8

goal state

exit

22/10/24

Implement Hill climbing search algorithm to solve N-Queens problem

function Hill-climbing(problem) returns a state that is a local maximum

current \leftarrow make-node(problem, initial-state)

loop do

neighbor \leftarrow a highest-valued successor of current
if neighbor.value \leq current.value. then

return current.state

current \leftarrow neighbor

* State: 4 queens on the board. One queen per column.

- val variables : x_0, x_1, x_2, x_3 where x_i is the row position of the queen in column i.
Assume that there is one queen per column.

- Domain for each variable : $x_i \in \{0, 1, 2, 3\} \forall i$.

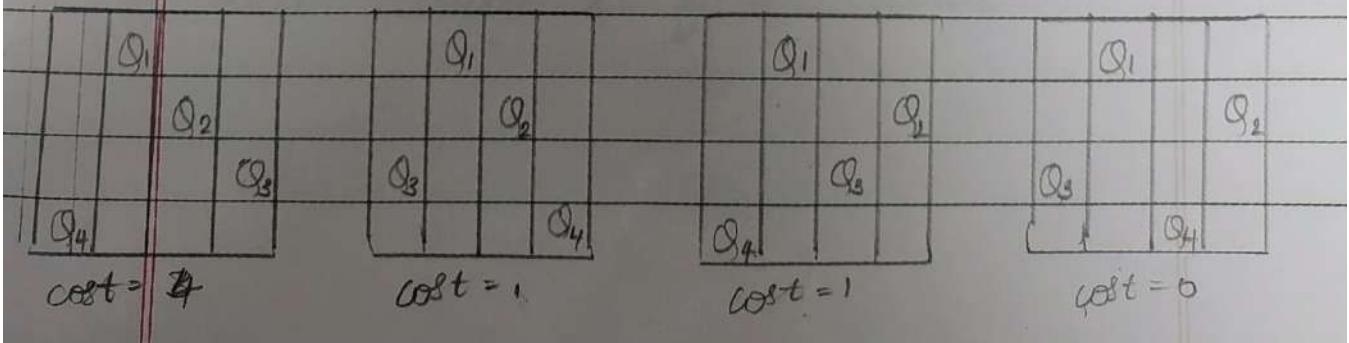
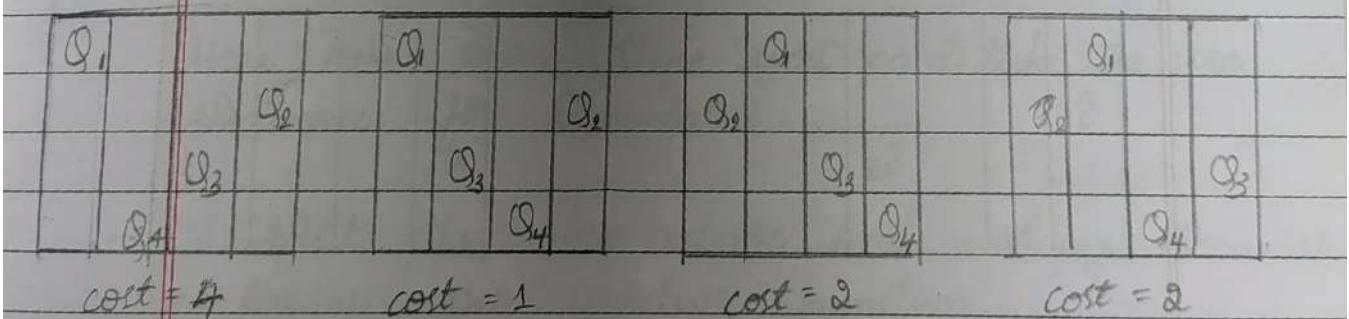
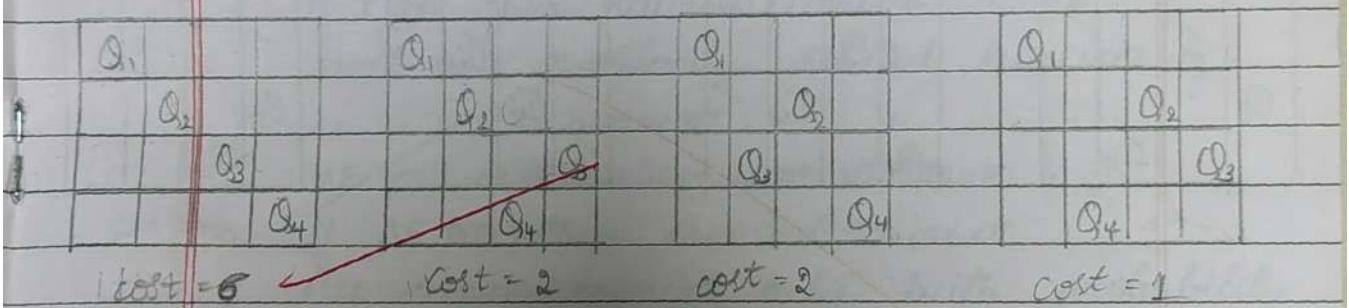
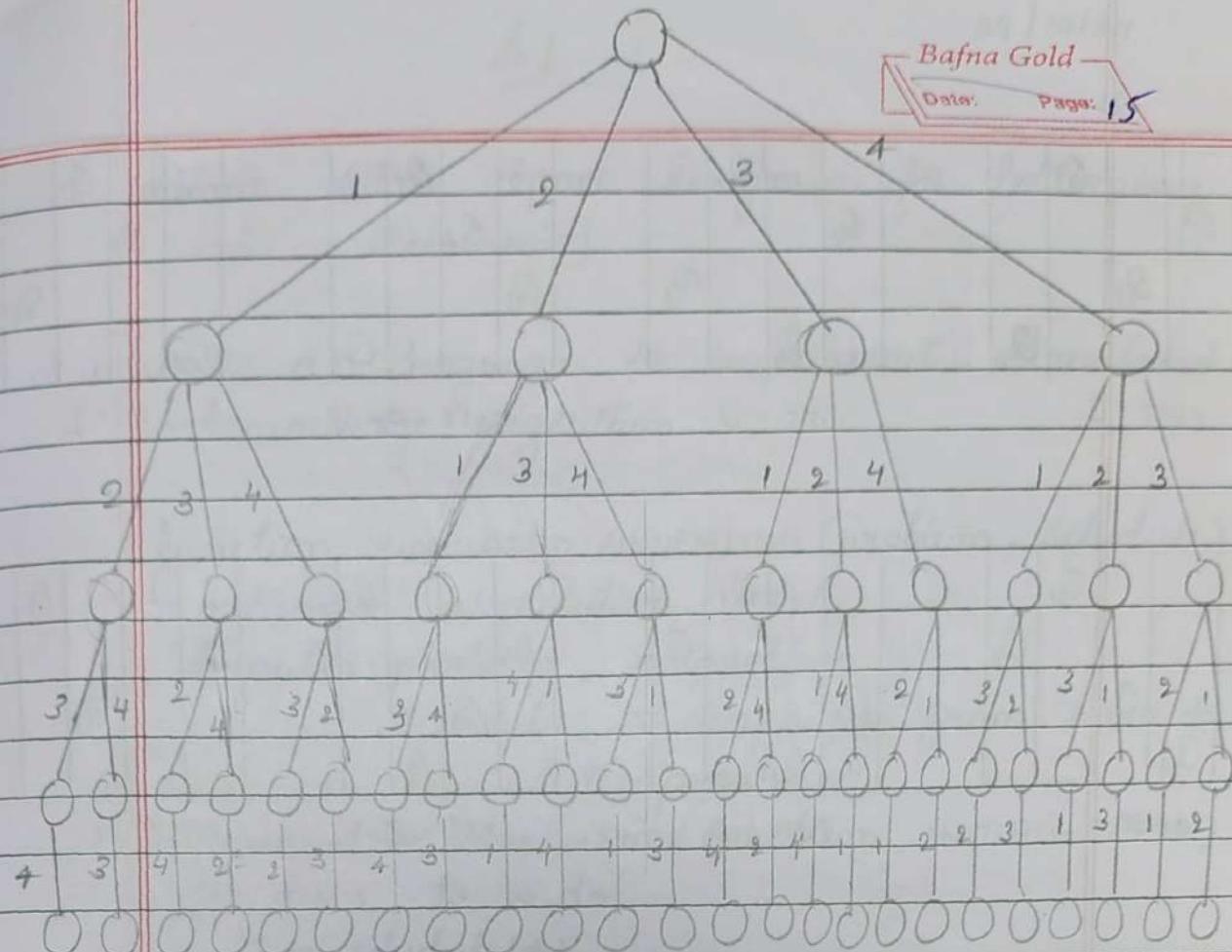
* Initial state : a ~~bad~~ random state

* goal state : 4 queens on the board. No pair of queens are attacking each other.

* Neighbours relation :

Swap the row positions of two queens

* cost function: The number of pairs of queens attacking each other, directly or indirectly



Q_1	Q_1	Q_1	Q_1	Q_1
Q_2	Q_2	Q_2	Q_2	Q_2
Q_3		Q_3	Q_3	Q_3
Q_4		Q_4	Q_4	Q_4

cost = 1

cost = 0

cost = 4

cost = 1

Q_1	Q_1	Q_1	Q_1	Q_1	Q_1
Q_2		Q_2	Q_2	Q_2	Q_2
Q_3		Q_3		Q_3	Q_3
Q_4		Q_4		Q_4	Q_4

cost = 2

cost = 2

cost = 4

cost = 1

Q_1	Q_1	Q_1	Q_1	Q_1	Q_1
Q_2		Q_2	Q_2	Q_2	Q_2
Q_3		Q_3		Q_3	Q_3
Q_4		Q_4		Q_4	Q_4

cost = 1

cost = 2

cost = 6

cost = 2

Best Solution :-

Q_1				Q_1	
Q_2				Q_2	
	Q_3			Q_3	
	Q_4			Q_4	

Draw State Space Diagram for following Scenario

Write a Program to implement Simulated Annealing Algorithm

```

function SIMULATED_ANNEALING(problem, schedule)
    returns a solution state
inputs: problem, a problem
        schedule, a mapping from time to
        "temperature"
current ← MAKE-NODE(problem, INITIAL-STATE)
for t = 1 to ∞ do
    T ← schedule(t)
    if T = 0 then return current
    next ← a randomly selected successor of
          current
    ΔE ← next.VALUE - current.VALUE
    if ΔE > 0 then current ← next
    else current ← next only with probability
           $e^{\Delta E / T}$ 

```

Here, are some steps you can take to use mirose for Simulated

1. Import the mirose and numpy libraries
2. Define the objective function.
3. use the mirose algorithms to get best state, best fitness & fitness curve

The Stimulated Annealing Algorithm can be decomposed in 4 simple steps

1. Start at a random point x .

2. choose a new point x_j on a neighbour $N(x)$.

3. Decide whether or not to move to new point x_j . The decision will be made based on the probability function $P(x, x_j, T)$

$$P(x, x_j, T) = \begin{cases} 1 & \text{Si } F(x_j) \geq F(x) \\ e^{\frac{F(x_j) - F(x)}{T}} & \text{Si } F(x_j) < F(x) \end{cases}$$

4. Reduce T

Python code

```
import mbrose.hive as mbrose
import numpy as np
```

```
def queensmax(position):
    queennoattacking = 0
    for i in range(len(position)-1):
        noattack = 0
        for j in range(i+1, len(position)):
            if (position[j] != position[i]) and
               (position[j] != position[i-j]) and
               (position[j] != position[i-(j-i)]):
                noattack += 1
        if noattack == len(position)-1:
            queennoattacking += 1
    return queennoattacking
```

try :

user_input = input("Enter the initial position as 8 comma-separated integers: ")

initialpos = np.array([int(x) for x in userinput.split(',')])

if len(initialpos) != 8 or any(x < 0 or x >= 8 for x in initialpos):

raise ValueError("Please enter exactly 8 integer between 0 & 7.")

except ValueError as e:

print(e)

exit()

objective = mbrose.customFitness(Queenmax)

problem = mbrose.DiscreteOpt(length = 8, fitness_fn

= objective, maximize=True, maxval = 8)

T = mbrose.ExpDecay()

result = mbrose.simulated_annealing(problem =

problem, schedule = T, max_attempts = 500,

max_iter = 5000, init_state = initialpos)

beststate = result[0]

best_fitness = result[1]

~~print('The best position found is : ', beststate)~~

~~print('The no. of queens that are not attacking each other is : ', best_fitness)~~

print("\nBest State Diagram: ")

board = [[':' for _ in range(8)] for _ in range(8)]

for row, col in enumerate(beststate):

board[row][col] = 'Q'

for now in board :

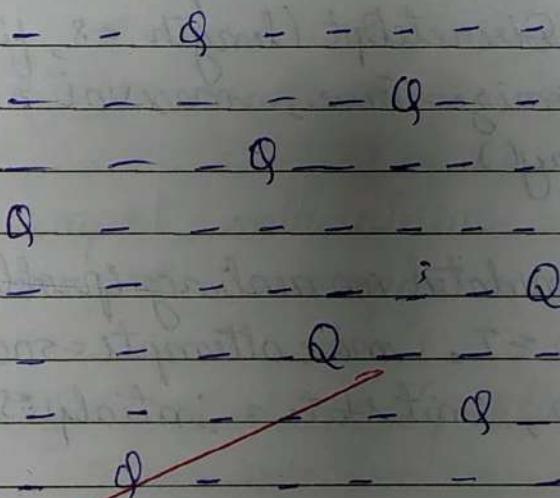
- print ('', join(now))

output

Enter the initial position as 8 comma-separated integers : 1, 2, 4, 6, 7, 5, 3, 0.

The best position found is : [3 7 0 2 5 1 6 4]
The no. of queens that are not attacking each other is : 7.0

Best State Diagram :



SF
22/10/2024

Implementation of truth table enumeration algorithm for deciding propositional

TT - Truth table

PL - True? return true if sentence holds within a model.

The π variable model represents a partial model - an assignment to some of symbols.

The keyword "and" is used here as a logical operator on its two arguments, returning true or false.

Algorithm:-

function $\Pi\text{-ENTAILS?}(KB, \alpha)$ returns true or false

inputs: KB , the knowledge base, a sentence in propositional logic α , the query, a sentence in propositional logic

$\text{symbols} \leftarrow$ a list of propositional symbol in KB and α

return $\Pi\text{-CHECK-ALL}(KB, \alpha, \text{symbols}, \emptyset)$

function $\Pi\text{-CHECK-ALL}(KB, \alpha, \text{symbols}, \text{model})$

returns true or false

if $\text{EMPTY?}(\text{symbols})$ then

if $\text{PL-TRUE?}(KB, \text{model})$ then return PL-True
(α , model)

else return true || when KB is false,
always return true

else do

$P \leftarrow \text{FIRST}(\text{symbols})$

$\text{rest} \leftarrow \text{REST}(\text{symbols})$

$\text{return } (\pi\text{-CHECK-ALL}(KB, } \alpha, \text{rest, model } v \\ \langle P = \text{true} \rangle)$

and

$\pi\text{-CHECK-ALL}(KB, } \alpha, \text{rest, model } v \\ \langle P = \text{false} \rangle)$

Example

$$\alpha = A \vee B$$

$$KB = (A \vee C) \wedge (B \vee \neg C)$$

$$KB \models \alpha$$

A	B	C	$A \vee C$	$B \vee \neg C$	$\neg B$
false	false	false	false	true	
false	false	true	true	false	
false	true	false	false	true	
false	true	true	true	true	
true	false	false	true	false	
true	false	true	true	false	
true	true	false	true	true	
true	true	true	true	true	

KB	α
false	false
false	false
true	true
true	true
false	true
true	true
true	true

Output

A	B	C	KB	Alpha	KB entails Alpha?
false	false	false	false	false	NO
false	false	true	false	false	NO
false	false	false	false	true	NO
false	false	true	true	true	Yes
true	false	false	true	true	Yes
true	false	true	false	true	NO
true	true	false	true	true	Yes
true	true	true	true	true	Yes

Does KB entail alpha? True

S8
18/11/2024

Implement Unification in first Order logic

Algorithm : unify(Ψ_1, Ψ_2)

Step 1 : If Ψ_1 or Ψ_2 is a variable or constant, then :

@ If Ψ_1 or Ψ_2 are identical, then return Ψ_1 .

⑥ Else if Ψ is a variable,

a. Then if Ψ_1 occurs in Ψ_2 , then return FAILURE

b. Else return $\{(\Psi_2 / \Psi_1)\}$.

c. Else if Ψ_2 is a variable,

a. If Ψ_2 occurs in Ψ_1 , then return FAILURE

b. Else return $\{(\Psi_1 / \Psi_2)\}$.

d. Else return FAILURE

Step 2 : If the initial Predicate symbol in Ψ_1 & Ψ_2 are not same, then return FAILURE.

Step 3 : If Ψ_1 & Ψ_2 have a different number of arguments, then return FAILURE.

Step 4 : Set Substitution set (SUBST) to NIL.

Step 5 : For i=1 to number of elements in Ψ ,

⑥ Call unify function with the i^{th} element of Ψ_1 and i^{th} element of Ψ_2 , & put results into S .

⑥ If $S = \text{failure}$ then returns FAILURE

⑥ If $S \neq \text{NIL}$ then do,

a. Apply S to the remainder of

both L1 & L2

b. $\text{SUBST} = \text{APPEND}(S, \text{SUBST})$

Step 6: Return SUBST

* example 1:

$$p(x, f(y)) - ①$$

$$p(a, f(g(x))) - ②$$

① & ② are identical if x is replaced with a
in ①
 a/x

$$p(a, f(y)) - ①$$

if y is replaced with $g(x)$

$$p(a, f(g(x))) - ①$$

Now, ① & ② are equal so it is given as output "TRUE"

* example 2:

~~$Q(a, g(x), a), f(y)) - ①$~~

~~$Q(a, g(f(b)), a), x) - ②$~~

if x is replaced with $f(b)$ in ①. $f(b)/x$

$$Q(a, g(f(b)), a), f(y)) - ①$$

if $f(y)$ is replaced with x in ①,
 ~~$Q(a, g(f(b)), a), x) - ①$~~ $x/f(y)$

now, ① & ② are equal

Cannot able to replace with x in ①
because the ^{same} variable x didn't
replaced twice

so, it is not unified. It is
FAILURE

* example 3 =

$$\Psi_1 = p(b, x, f(g(z))) \quad \text{--- ①}$$

$$\Psi_2 = p(z, f(y), f(y)) \quad \text{--- ②}$$

If b^2 is replaced by b in ②

$$\Psi_2 = p(b, f(y), f(y)) \quad \text{--- ②}$$

If x is replaced with $f(y)$ in ①

$$\Psi_1 = p(b, f(y), f(g(z))) \quad \text{--- ①}$$

If y is replaced with $g(z)$ in ②

$$\Psi_2 = p(b, f(y), f(g(z))) \quad \text{--- ②}$$

Now, ① & ② are equal so, it is unified

* example 4 =

$$\Psi_1 = p(f(a), g(c))$$

$$\Psi_2 = p(x, x)$$

x cannot replaced twice so it
is failure.

Output :

Enter the first term : [`'p', 'b', '?x', ['f', ['g', '?z']]`]

Enter the second term : [`'p', '?z', ['f', '?y'], ['f', '?y']`]

unification successfull !

Substitution : $?z : b, ?x : [f, ?y], ?y : [g, ?z]$

Enter the first term : [`'p', ['f', ['a']], ['g', ['?y']]`]

Enter the second term : [`'p', '?x', '?x'`]

unification failed : Predicate symbols
don't match : $g := f$

Ex 12/11

Routing Information Protocol (RIP)

Topology :-

Create KB consisting of FOL statements & prove the given query using forward reasoning

Algorithm:-

function FOL-FC-Ask(KB, α) returns a substitution β

input: KB, the knowledge base, a set of first order definite clauses α , the query, an atomic sentence.

local variable: new, the new sentence inferred on each iteration

repeat until new is empty

new $\leftarrow \emptyset$

for each rule in KB do

$(P_1 \wedge \dots \wedge P_n \Rightarrow q) \leftarrow \text{STANDARDIZE-VARIABLES}$
(rule.)

if some $P'_1 \wedge \dots \wedge P'_n$ in KB

$q' \leftarrow \text{SUBST}(P'_n, q)$

if q' does not unify with some sentence already in KB in new then add q' to new

$\phi \leftarrow \text{UNIFY}(q', \alpha)$

if ϕ is not fail then return ϕ

add new to KB

return false

Output :-

Enter your FOL statements :

- > American(p) \wedge weapon(q) \wedge sells(p, q, n) \wedge
- Hostile(n) \Rightarrow Criminal(p)
- > $\exists x \text{ owns}(A, x) \wedge \text{missile}(x)$
- > $\text{owns}(A, T1)$
- > $\text{missile}(T1)$
- > $\forall x \text{ missile}(x) \wedge \text{owns}(A, x) \Rightarrow \text{sells}(\text{Robert}, x, A)$
- > $\text{missile}(x) \Rightarrow \text{weapon}(x)$
- > $\forall x \text{ Enemy}(x, \text{America}) \Rightarrow \text{Hostile}(x)$
- > American(Robert)
- > Enemy(A, America)
- > done

Enter the query to prove : Criminal(Robert)

Proven : Criminal(Robert)

Example :-

It is a crime for an American to sell weapons to hostile nations

Let's say p, q and n are variable
American

Country a has some missiles.

so

~~Existential instantiation, introducing a new constant T1:~~

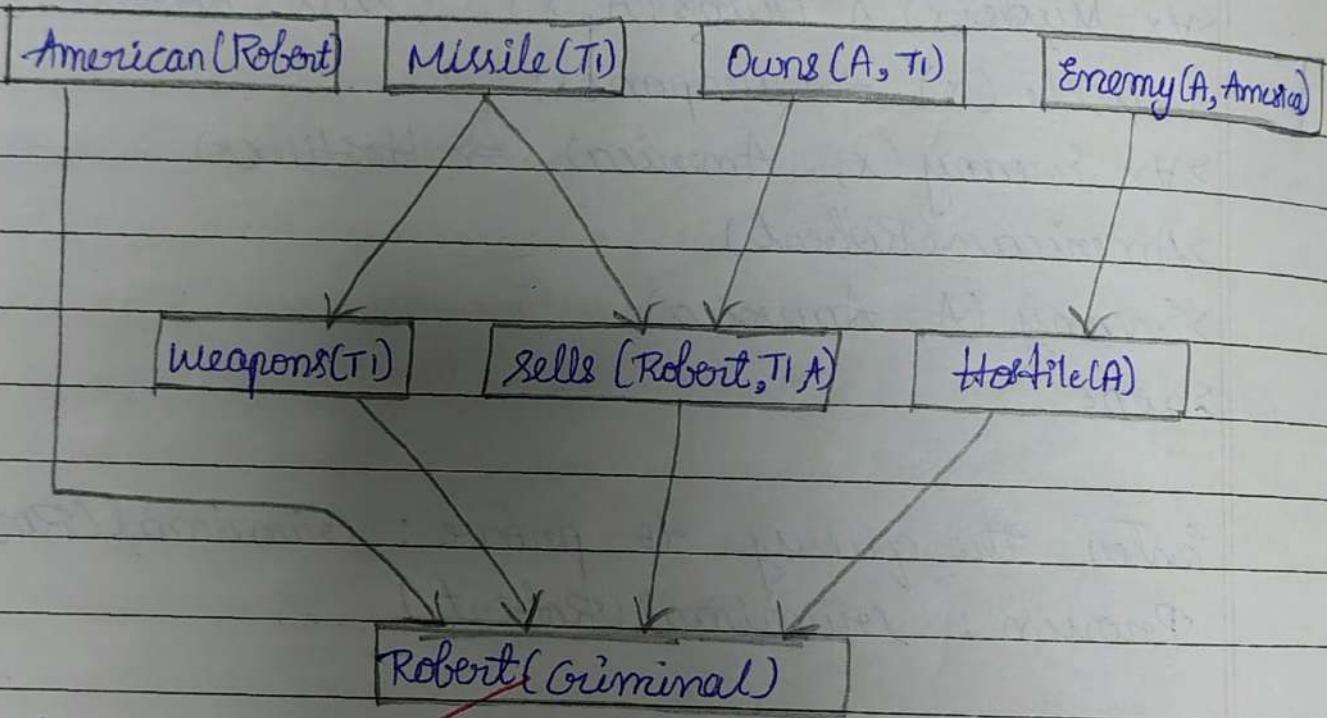
All of the missiles were sold to country A by Robert

missiles are weapons

Enemy of America is known as hostile

Robert is an American

The country A, an enemy of America.



American(p) 1 weapon(q) 1 sells(p, q, x) 1 Hostile(x)
=> Criminal

Se
st
planar

Convert a given first order logic statement into conjunctive normal form (CNF) to Resolution

Basic steps for providing a conclusion
of given premises

Premises, ..., Premise
(all expressed in FOL):

1. Convert all sentences to CNF
2. Negate conclusion & convert result to CNF
3. Add negated conclusion & to the premises clauses
4. Repeat until contradiction or no progress is made:
 - @ Select 2 clauses (call them parent clauses)
 - ② Resolve them together, performing all required unification
 - ③ If resolvent is empty clause, a contradiction has been found (i.e., \perp follows from premises).
 - ④ If not, add resolvent to premises

~~If we succeed in step 4, we have proved the conclusion~~

Give the KB of premises:

- a. John likes all kind of food
- b. Apple and vegetables are food
- c. Anything anyone eats and not killed is food
- d. Anil eats everything that Aril eats peanuts & still alive
- e. Harry eats everything that Anil eats
- f. Anyone who is alive implies not killed
- g. Anyone who is not killed implies alive

Prove by resolution that:

Representation in FOL

- a. $\forall x : \text{food}(x) \rightarrow \text{likes}(\text{John}, x)$
- b. $\text{food}(\text{Apple}) \wedge \text{food}(\text{vegetables})$
- c. $\forall x \forall y : \text{eats}(x, y) \wedge \neg \text{killed}(x) \rightarrow \text{food}(y)$
- d. $\text{eats}(\text{Anil}, \text{Peanuts}) \wedge \text{alive}(\text{Anil})$
- e. $\forall x : \text{eats}(\text{Anil}, x) \rightarrow \text{eats}(\text{Harry}, x)$
- f. $\forall x : \neg \text{killed}(x) \rightarrow \text{alive}(x)$
- g. $\forall x : \text{alive}(x) \rightarrow \neg \text{killed}(x)$
- h. $\text{likes}(\text{John}, \text{Peanuts})$

Eliminate implication \Rightarrow with $\neg \alpha \beta$

- a. $\forall x \neg \text{food}(x) \vee \text{likes}(\text{John}, x)$
- b. ~~$\text{food}(\text{Apple}) \wedge \text{food}(\text{vegetables})$~~
- c. ~~$\forall x \forall y \neg [\text{eats}(x, y) \wedge \neg \text{killed}(x)] \vee \text{food}(y)$~~
- d. ~~$\text{eats}(\text{Anil}, \text{Peanuts}) \wedge \text{alive}(\text{Anil})$~~
- e. $\forall x : \text{eats}(\text{Anil}, x) \rightarrow \text{eats}(\text{Harry}, x)$
- f. $\forall x : \neg \text{killed}(x) \rightarrow \text{alive}(x)$
- g. $\forall x : \text{alive}(x) \rightarrow \neg \text{killed}(x)$
- h. ~~$\text{likes}(\text{John}, \text{Peanuts})$~~

More negation (\neg) onwards & rewrite

- $\forall x \neg \text{food}(x) \vee \text{likes}(\text{John}, x)$
- $\text{food}(\text{Apple}) \wedge \text{food}(\text{vegetables})$
- $\forall x \forall y \neg \text{eats}(x, y) \vee \text{killed}(x) \vee \text{food}(y)$
- $\text{eats}(\text{Anil}, \text{Peanuts}) \wedge \text{alive}(\text{Anil})$
- $\forall x \neg \text{eats}(\text{Anil}, x) \vee \text{eats}(\text{Harry}, x)$
- $\forall x \neg \text{killed}(x) \wedge \text{alive}(x)$
- $\forall x \neg \text{alive}(x) \vee \neg \text{killed}(x)$
- $\text{likes}(\text{John}, \text{Peanuts})$

Rename variable & Standardize variable

- $\forall x \neg \text{food}(x) \vee \text{likes}(\text{John}, x)$
- $\text{food}(\text{Apple}) \wedge \text{food}(\text{vegetables})$
- $\forall y \forall z \neg \text{eats}(y, z) \vee \text{killed}(y) \vee \text{food}(z)$
- $\text{eats}(\text{Anil}, \text{Peanuts}) \wedge \text{alive}(\text{Anil})$
- $\forall w \neg \text{eats}(\text{Anil}, w) \vee \text{eats}(\text{Harry}, w)$
- $\forall g \neg \text{killed}(g) \wedge \text{alive}(g)$
- $\forall k \neg \text{alive}(k) \vee \neg \text{killed}(k)$
- $\text{likes}(\text{John}, \text{Peanuts})$

Drop universal Quantifiers

- $\neg \text{food}(x) \vee \text{likes}(\text{John}, x)$
- ~~$\text{food}(\text{Apple})$~~
- ~~$\text{food}(\text{vegetable})$~~
- ~~$\neg \text{eats}(y, z) \vee \text{killed}(y) \vee \text{food}(z)$~~
- ~~$\text{eats}(\text{Anil}, \text{Peanuts})$~~
- ~~$\text{alive}(\text{Anil})$~~
- $\neg \text{eats}(\text{Anil}, w) \vee \text{eats}(\text{Harry}, w)$
- $\neg \text{killed}(g) \vee \text{alive}(g)$
- $\neg \text{alive}(k) \vee \neg \text{killed}(k)$
- $\text{likes}(\text{John}, \text{Peanuts})$

Proof by Resolution

$\neg \text{likes}(\text{john}, \text{Peanuts})$

$\neg \text{food}(x) \vee \text{likes}(\text{john}, x)$
 $\neg \text{peanuts} \vee x = \text{john}$

$\neg \text{food}(\text{Peanuts})$

$\neg \text{eats}(y, z) \vee \text{killed}(y)$

$\neg \text{food}(z)$

$\neg \text{Peanuts} \vee z = y$

$\neg \text{eats}(y, \text{Peanuts}) \vee \text{killed}(y)$

$\text{eats}(\text{Anil}, \text{Peanuts})$
 $\neg \text{Anil} \vee y = z$

$\text{killed}(\text{Anil})$

$\neg \text{alive}(k) \vee \neg$

$\text{killer}(k)$

$\neg \text{alive}(\text{Anil}) \quad \text{alive}(\text{Anil})$

L3

Hence proved.

Output :-

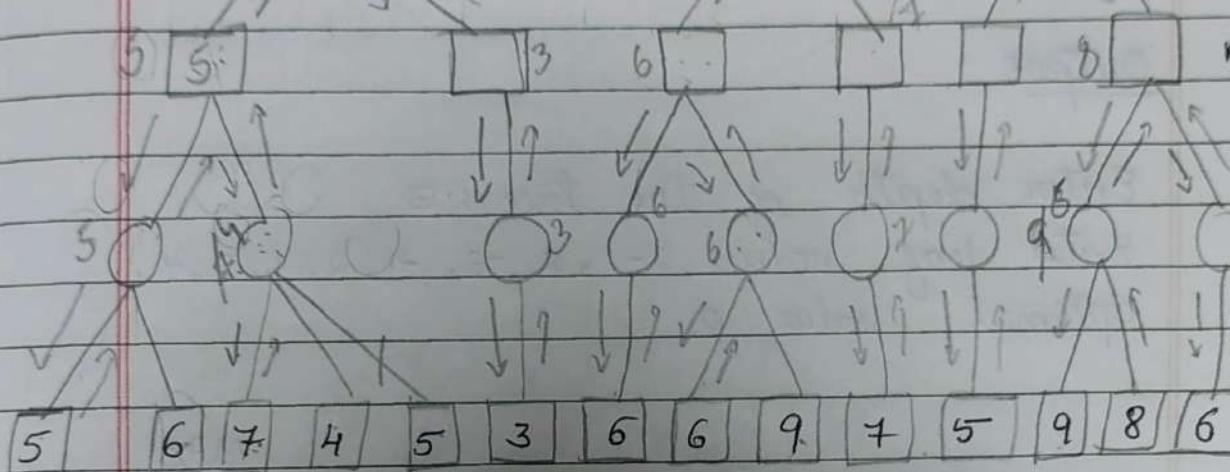
Derived Fact : food (Peanuts)

Derived Fact : likes (john, Peanuts)

Proven : Likes (john, Peanuts)

Simple 2020

Alpha Beta pruning Algorithm

max α min β Max α min β max α min β 

$$\alpha >= \beta$$

Algorithm :-

function MINIMAX-DECISION (state) return an action
 returns any max a \in ACTIONS(s)
 MAX-VALUE (RESULT (s, a))

function MAX-VALUE (s) returns a utility value
 if TERMINAL-TEST (s) then return
 UTILITY (s)

$$v \leftarrow -\infty$$

for each a in ACTION (s) do

$$v \leftarrow \max (v, \text{MIN-VALUE} (\text{RESULT} (s, a)))$$

return v

function MIN-VALUE (state) return a utility value
if TERMINAL-TEST (state) then return
UTILITY (state)

$$v \leftarrow \infty$$

for each a in ACTIONS (state) do
 $v \leftarrow \min(v, \text{MAX-VALUE}(\text{RESULT}(s, a)))$
return v

Output

Enter depth of the tree : 3

Enter leaf values : -1, 8, -3, -1, 2, 1, 3, 4

Optimal value : 2

