Math and Bitmanipulation Math

1.Extended Euclidean Algorithm(Gcd and Hcf)

```
We already know Basic Euclidean Algorithm. Now using the Extended Euclidean Algorithm, given a and b calculate the GCD and integer coefficients \mathbf{x}, \mathbf{y}. Using the same. \mathbf{x} and \mathbf{y} must satisfy the equation \mathbf{a}\mathbf{x} + \mathbf{b}\mathbf{y} = \mathbf{gcd}(\mathbf{a}, \mathbf{b}).

Example 1:

Input:

a = 35
b = 15
Output:
5 1 - 2
Explanation:
gcd(a,b) = 5
35*1 + 15*(-2) = 5

Example 2:

Input:
a = 30
b = 20
Output:
10 1 -1
Explanation:
gcd(30,20) = 10
30*(1) + 20*(-1) = 10
```

```
class Solution {
    static int[] gcd(int a, int b) {
        // code here
        if (b == 0) {
            return new int[]{a, 1, 0};
        }

        // Recursively call gcd(b, a % b)
        int[] vals = gcd(b, a % b);
        int gcd = vals[0];
        int x1 = vals[1];
        int y1 = vals[2];

        // Update x and y using results of recursion
        int x = y1;
        int y = x1 - (a / b) * y1;

        return new int[]{gcd, x, y};
    }
}
```

2. Sieve of Eratosthenes

Given a positive integer \mathbf{n} , calculate and return all prime numbers less than or equal to \mathbf{n} using the Sieve of Eratosthenes algorithm.

A prime number is a natural number greater than 1 that has no positive divisors other than 1 and itself.

Examples:

```
Input: n = 10
Output: 2 3 5 7
Explanation: Prime numbers less than equal to 10 are 2 3 5 and 7.

Input: n = 35
Output: 2 3 5 7 11 13 17 19 23 29 31
Explanation: Prime numbers less than equal to 35 are 2 3 5 7 11 13 17 19 23 29 and 31.
```

Constraints:

```
// User function Template for Java
class Solution {
    static ArrayList<Integer> sieveOfEratosthenes(int n) {
        ArrayList<Integer> primes = new ArrayList<>();
        boolean[] isPrime = new boolean[n + 1];
        for (int i = 2; i \le n; i++) {
            isPrime[i] = true;
        }
        for (int i = 2; i <= Math.sqrt(n); i++) {
            if (isPrime[i]) {
                for (int j = i * i; j <= n; j += i) {
                    isPrime[j] = false;
            }
        }
        for (int i = 2; i \le n; i++) {
            if (isPrime[i]) {
                primes.add(i);
        return primes;
    }
}
```

4.Nth Fibonacci Number

```
Given a non-negative integer n, your task is to find the nth Fibonacci number.
The Fibonacci sequence is a sequence where the next term is the sum of the previous two terms.
The first two terms of the Fibonacci sequence are 0 followed by 1. The Fibonacci sequence: 0, 1, 1,
2, 3, 5, 8, 13, 21
The Fibonacci sequence is defined as follows:
  • F(0) = 0
  • F(1) = 1
  • F(n) = F(n - 1) + F(n - 2) for n > 1
Examples:
 Input: n = 5
 Output: 5
 Explanation: The 5th Fibonacci number is 5.
 Input: n = 0
 Output: 0
 Explanation: The 0th Fibonacci number is 0.
 Input: n = 1
 Output: 1
 Explanation: The 1st Fibonacci number is 1.
```

4.Euler Totient Function

Find the **Euler Totient Function (ETF)** $\Phi(N)$ for an input N. ETF is the count of numbers in $\{1, 2, 3, ..., N\}$ that are relatively prime to N, i.e., the numbers whose GCD (Greatest Common Divisor) with N is 1.

Example 1:

```
Input:

N = 11

Output:

10

Explanation:

From 1 to 11,

1,2,3,4,5,6,7,8,9,10

are relatively prime to 11.
```

Example 2:

```
Input:
N = 16
Output:
8
Explanation:
From 1 to 16
1,3,5,7,9,11,13,15
are relatively prime
to 16.
```

```
26 // User function Template for Java
27 class Solution {
        static long ETF(long N) {
29
            long result = N;
            for (long p = 2; p * p <= N; ++p) {
31 -
                 if (N \% p == 0) {
32 -
                     while (N % p == 0)
                         N /= p;
                     result -= result / p;
                 }
37
            if (N > 1)
                 result -= result / N;
            return result;
41
        }
42
    }
```

Learn About Catalan Number(Im

https://www.geeksforgeeks.org/catalan-numbers/

3.Learn About Modular Arithmetic

https://www.geeksforgeeks.org/modular-arithmetic/

Learn About Prime Factorisation

https://www.geeksforgeeks.org/prime-factorization/

Learn About Chinese remainder theorem

https://www.geeksforgeeks.org/chinese-remainder-theorem/

Bit Manipulation

1.Count set bit in an integer(No. of 1 bit)

```
Given a positive integer n. Your task is to return the count of set bits.

Examples:

Input: n = 6
Output: 2
Explanation: Binary representation is '110', so the count of the set bit is 2.

Input: n = 8
Output: 1
Explanation: Binary representation is '1000', so the count of the set bit is 1.

Input: n = 3
Output: 2
```

```
// User function Template for Java
class Solution {
    static int setBits(int n) {
        // code here
        int countN=Integer.bitCount(n);
        return countN;
    }
}
```

2.Bit difference

You are given two numbers **a** and **b**. The task is to count the number of bits needed to be flipped to convert a to b.

```
Input: a = 10, b = 20
Output: 4
Explanation:
a = 01010
b = 10100
As we can see, the bits of A that need to be flipped are 01010. If we flip these bits, we get 10100, which is B.

Input: a = 20, b = 25
Output: 3
Explanation:
a = 10100
b = 11001
As we can see, the bits of A that need to be flipped are 10100. If we flip these bits, we get 11001, which is B.
```

3.Count total Set bits

You are given a number ${\bf n}$. Find the total count of set bits for all numbers from 1 to n (both inclusive). **Examples:**Input: n=4Output: 5Explanation: For numbers from 1 to 4. For 1: $0 \ 0 \ 1 = 1$ set bits For 2: $0 \ 1 \ 0 = 1$ set bits For 3: $0 \ 1 \ 1 = 2$ set bits For 4: $1 \ 0 \ 0 = 1$ set bits Therefore, the total set bits is 5.

Input: n=17Output: 35Explanation: From numbers 1 to 17(both inclusive), the total number of set bits is 35.

Expected Time Complexity: 0(logn)
Expected Auxiliary Space: 0(1)

Constraints: $1 \le n \le 10^8$

```
class Solution {

public static int countSetBits(int n){

    if(n=0) return 0;
    // Sum = x * 2^(x-1) + (n-2^x+1) + countSetBits(n-2^x)

    int x = calculateX(n);
    int countBitsTillTwoX = x * (1<<(x-1) ); // x * (2 to power x-1)
    int msb2xtoN = n - (1<<x) + 1;
    int rest = n - (1<<x); // solve for 3 for n = 11 ( 11 - 2 power 3)

    return countBitsTillTwoX + msb2xtoN + countSetBits(rest);

}

public static int calculateX(int n){
    int x = 0;
    while( (1<<x) <= n ){ // 2 raise to power x is less than or equal to n
        x++;
    }
    return x-1;
}
</pre>
```

4.Find Position of set bit

Given a number \mathbf{n} having only one '1' and all other '0's in its binary representation, find the position of the only set bit. If there are 0 or more than 1 set bit the answer should be -1. The position of set bit '1' should be counted starting with 1 from the LSB side in the binary representation of the number.

Examples:

```
Input: n=2
Output: 2
Explanation: 2 is represented as "10" in Binary. As we see there's only one set bit and it's in position 2.
```

```
Input: n=5
Output: -1
Explanation: 5 is represented as "101" in Binary. As we see there's two set bits and thus the output -1.
```

Constraints:

0 <= n <= 10⁸

```
// User function Template for Java
class Solution {
    static int findPosition(int n) {
        // code here
        if(n<=0 || ( n&(n-1))!=0){
            return -1;
        }
        int pos=1;
        while(n > 1){
            n >>=1;
            pos++;
        }
        return pos;
}
```

5.Copy set bit in range

Given two numbers X and Y, and a range [L, R] where $1 \le L \le R \le 32$. You have to copy the set bits of Y' in the range L to R in Y'. Return this modified X.

Note: Range count will be from Right to Left & start from 1.

Example 1:

```
Input:

X = 44, Y = 3

L = 1, R = 5

Output:

47

Explaination:

Binary represenation of 44 and 3 is 101100 and 000011. So in the range 1 to 5 there are two set bits of 3 (1st & 2nd position). If those are set in 44 it will become 101111 which
```

Example 2:

is 47.

```
Input:

X = 16, Y = 2

L = 1, R = 3

Output: 18

Explaination: Binary representation of 16 and 2 is 10000 and 10. If the mentioned conditions are applied then 16 will become 10010 which is 18.
```

```
28
29 class Solution {
        static int setSetBit(int x, int y, int l, int r) {
30 -
31
              for(int i = l; i <= r; i++){
32 -
33 -
                 if((y&(1<<(i-1)))!=0){//Check
                     x = x | (1 << (i-1)); // Set
34
35
                 }
36
             return x;
38
      }
39 }
```

6.Power Set

Given a string **s** of length **n**, find all the **possible non-empty** <u>subsequences</u> of the string **s** in **lexicographically-sorted** order.

Example 1:

```
Input:

s = "abc"

Output:

a ab abc ac b bc c

Explanation:

There are a total 7 number of subsequences possible for the given string, and they are mentioned above in lexicographically sorted order.
```

Example 2:

```
Input:

s = "aa"

Output:
a a aa

Explanation:

There are a total 3 number of subsequences possible for the given string, and they are mentioned above in lexicographically sorted order.
```

```
class Solution {
   public static void solve (String p,int index, List<String> ans, String s){
    if(index==s.length()){
        if(p.length()>0){
            ans.add(p);
         }
        return;
        solve(p,index+1,ans,s);
       p+=s.charAt(index);
        solve(p,index+1,ans,s);
   public List<String> AllPossibleStrings(String s) {
        List<String> ans= new ArrayList<String> ();
        String p ="";
        solve(p,0,ans,s);
        Collections.sort(ans);
        return ans;
```