

## Homework Week 2

### Text Questions

Monica Quaintance

mjq2102@columbia.edu

#### 1. Weiss Exercise 3.1

You are given a list,  $L$ , and another list,  $P$ , containing integers sorted in ascending order. The operation  $\text{printLots}(L,P)$  will print the elements in  $L$  that are in positions specified by  $P$ . For instance, if  $P = 1, 3, 4, 6$ , the elements in positions 1, 3, 4, and 6 in  $L$  are printed. Write the procedure  $\text{printLots}(L,P)$ . You may use only the public Collections API container operations. What is the running time of your procedure?

```
import java.util.ArrayList;
import java.util.ListIterator;

public class collectionPrinter<T> {

    ArrayList<T> fullList;
    ArrayList<Integer> targetList;
    ArrayList<T> resultList = new ArrayList<T>();

    public collectionPrinter(){
    }

    public ArrayList<T> printLots(ArrayList<T> L, ArrayList<Integer> P) {
        fullList = L;
        targetList = P;

        ListIterator<Integer> iterTarget = targetList.listIterator();
        ListIterator<T> iterFull = fullList.listIterator();

        int i = 0;
        int targetIndex = iterTarget.next();
        while (iterFull.hasNext() && i<= targetIndex)
        {
            T element = iterFull.next();
            if (i == targetIndex) {
                resultList.add(element);
                targetIndex = iterTarget.next();
            }
        }
    }
}
```

```

        else if (iterTarget.hasNext() == false)
            return resultList;

        else {
        }
        i++;
    }
    return resultList;
}
}

```

You must check at least  $j$  elements, where  $j$  is the value of the final element in  $P$ . Running time is  $O(N)$ .

## 2. Weiss Exercise 3.2

*Swap two adjacent elements by adjusting only the links (and not the data) using:*

*a. Singly linked lists.*

List A:  $[w] \rightarrow [x] \rightarrow [y] \rightarrow [z]$

becomes:  $[w] \rightarrow [y] \rightarrow [x] \rightarrow [z]$

- $w.\text{next} = y$  ( $x.\text{next}$ )
- $x.\text{next} = z$  ( $y.\text{next}$ )
- $y.\text{next} = x$

*b. Doubly linked lists.*

List B:  $[w] \leftrightarrow [x] \leftrightarrow [y] \leftrightarrow [z]$

becomes:  $[w] \leftrightarrow [y] \leftrightarrow [x] \leftrightarrow [z]$

- $x.\text{next} = z$  ( $y.\text{next}$ )
- $w.\text{next} = y$  ( $x.\text{next}$ )
- $y.\text{next} = x$
- $z.\text{prev} = x$
- $x.\text{prev} = y$
- $y.\text{prev} = w$

## 3. Weiss, Exercise 3.24

*Write routines to implement two stacks using only one array. Your stack routines should not declare an overflow unless every slot in the array is used.*

```

import java.util.EmptyStackException;

public class doubleStack
{
    // default stack of size 16
    String[] stackArray = (String[]) new String[16];
    private int leftTop = -1;
    private int rightTop = stackArray.length;

    public doubleStack() {
    }

    // specified stack size
    public doubleStack(int arraySize) {
        this.stackArray = (String[]) new String[arraySize];
        this.leftTop = -1;
        this.rightTop = stackArray.length;
    }

    // prints left stack (for debugging)
    public String printLeft()
    {
        if (leftTop == -1)
            throw new EmptyStackException();
        else {
            StringBuilder leftPrint = new StringBuilder();
            for (int i = 0; i < leftTop + 1; i++)
                leftPrint.append(stackArray[i]);
            return leftPrint.toString();
        }
    }

    // prints right stack (for debugging)
    public String printRight()
    {
        if (rightTop == stackArray.length)
            throw new EmptyStackException();
        else {
            StringBuilder rightPrint = new StringBuilder();
            for (int i = stackArray.length - 1; i > rightTop - 1; i--)
                rightPrint.append(stackArray[i]);
            return rightPrint.toString();
        }
    }
}

```

```

// push onto left side of array
public void leftPush(String newItem) {
    if (leftTop >= rightTop - 1)
        throw new IndexOutOfBoundsException("Both stacks are full");
    else {
        stackArray[leftTop+1] = newItem;
        leftTop++;
    }
}

// pop from left side of array
public String leftPop() {
    if (leftTop == -1)
        throw new EmptyStackException();
    else {
        String popItem = stackArray[leftTop];
        leftTop--;
        return popItem;
    }
}

// push onto right side of array
public void rightPush(String newItem) {
    if (leftTop >= rightTop-1)
        throw new IndexOutOfBoundsException("Both stacks are full");
    else {
        stackArray[rightTop-1] = newItem;
        rightTop--;
    }
}

// pop from right side of array
public String rightPop() {
    if (rightTop == stackArray.length)
        throw new EmptyStackException();
    else {
        String popItem = stackArray[rightTop];
        rightTop++;
        return popItem;
    }
}
}

```

**Weiss, Exercise 4.5 (this is asking for a proof)**

*Show that the maximum number of nodes in a binary tree of height  $h$  is  $2^{h+1}-1$ .*

For max # of nodes, assume tree is always complete. Number of nodes at height  $h = N(h)$ ,  $N(h) = 2^{h+1} - 1$ .

**Step One:** Base Case:  $N(0) = 1$

$2^{0+1} - 1 = 1$ , Base Case is true.

**Step Two:** Assume for all values of  $h$  up to  $h$ ,  $2^{h+1} - 1 = N(h)$

**Step Three:** Using assertion in 2:

$$N(h) = 2[N(h-1) - N(h-2)] + N(h-1)$$

$$N(h) = 2 * N(h-1) - 2 * N(h-2) + N(h-1)$$

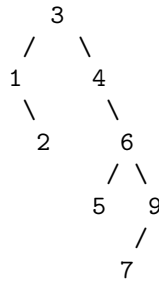
$$N(h) = 2[2^{h-1+1} - 1] - 2[2^{h-2+1} - 1] + [2^{h-1+1} - 1]$$

$$N(h) = 2^{h+1} - 2 - 2^h + 2 + 2^h - 1$$

$$N(h) = 2^{h+1} - 1$$

**Weiss, Exercise 4.9 (using a full deletion)**

*a. Show the result of inserting 3, 1, 4, 6, 9, 2, 5, 7 into an initially empty binary search tree.*



*b. Show the result of deleting the root.*

