Propositional Probability III: Compilation[1]

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- **Goal:** to evaluate terms from PROP as efficiently as possible

- Last time we saw $(\pi, (\varphi, w)) \Downarrow^e v$. We observed that it is efficient (i.e., polynomial in $|\varphi|$) if φ has specific forms: for instance, if it is a simple conjunction of variables. This time, we will see some more sophisticated evaluation strategies

1   Memoization

- Consider the formula $(a \vee b) \wedge (c \vee d)$ and some w.

- Let's see a sketch of a derivation tree (where we elide w and π for now):

$$(\star)\ \cfrac{\cfrac{\cdots}{c \vee d \Downarrow^e v_3} \qquad \text{false} \Downarrow^e 0}{\cfrac{(b \wedge (c \vee d)) \Downarrow^e v_1}{((a \vee b) \wedge (c \vee d)) \Downarrow^e w(a)v_1 + w(\overline{a})v_2}} \qquad \cfrac{\cdots}{(c \vee d)) \Downarrow^e v_2}\ (\star)$$

- **Observe**: both of the subtrees marked by $(\star)$ are *identical*. We are wasting effort by re-deriving both. We should *reuse* these derivations rather than recomputing them.

- Now we will define a new reduction rule that memoizes previously computed results

- Shape of new relation: $(\rho, \mathsf{p}) \Downarrow^m (v, \rho')$, where ρ is a memoization table that maps PROP programs to probabilities:

$$(\pi, \rho, (\mathtt{true}, w)) \Downarrow^m (1, \rho) \qquad (\pi, \rho, (\mathtt{false}, w)) \Downarrow^m (0, \rho)$$

$$(\textsc{Memo})\ \cfrac{\rho(\varphi) = v}{(\pi, \rho, (\varphi, w)) \Downarrow^e (v, \rho)}$$

$$(\textsc{Split})\ \cfrac{\varphi \notin \rho \qquad (\pi, \rho, (\varphi[\mathtt{true}/x], w)) \Downarrow^m (\rho', v_1)}{(x :: \pi, \rho, (\varphi[\mathtt{false}/x] \mapsto v_1), (\varphi[\mathtt{false}/x], w)) \Downarrow^m (\rho'', v_2)}{(x :: \pi, \rho, (\varphi, w)) \Downarrow^m (\rho'' \cup (\varphi[\mathtt{false}/x] \mapsto v_2), w(x)v_1 + w(\overline{x})v_2)}$$

So far we've considered *tree-like* proof-systems (also called *Gentzen sequent calculus [Gentzen, 1969]*): these are systems of deductions that describe trees, and cannot share sub-derivations. There are more powerful DAG-like proof-systems that permit the reusing of subderivations. If you are interested in learning more about the power of various proof systems, check out Beame and Pitassi [2001]

Analysis of memoization

- What kinds of PROP programs does our new memoized reduction scale well on?

- For the purposes of analysis, it can be useful to give a *cost-annotated* version of our rules that counts how many derivations were produced during runtime. This is fairly easy to do: we augment our relation with an integer n that accumulates the total runtime, so our new relation is $(\pi, \rho, (\varphi, w)) \Downarrow (v, \rho', n)$:

$$(\pi, \rho, (\texttt{true}, w)) \Downarrow^m (1, \rho, 1) \qquad (\pi, \rho, (\texttt{false}, w)) \Downarrow^m (0, \rho, 1)$$

$$(\textsc{Memo}) \ \frac{\rho(\varphi) = v}{(\pi, \rho, (\varphi, w)) \Downarrow^e (v, \rho, 1)}$$

$$(\textsc{Split}) \ \frac{\varphi \notin \rho \qquad (\pi, \rho, (\varphi[\texttt{true}/x], w)) \Downarrow^m (\rho', v_1, n_1) \quad (\pi, \rho' \cup (\varphi[\texttt{true}/x] \mapsto v_1), (\varphi[\texttt{false}/x], w)) \Downarrow^m (\rho'', v_2, n_2)}{(x :: \pi, \rho, (\varphi, w)) \Downarrow^m (\rho'' \cup (\varphi[\texttt{false}/x] \mapsto v_2), w(x)v_1 + w(\overline{x})v_2, n_1 + n_2 + 1)}$$

- Complexity in the *elimination-width*, similar to variable elimination (won't get into that now...)

2   Tractable runtimes

- **Problem**: evaluating probabilities is hard for *arbitrary formulae*. What if we consider a *restricted class* of formulae for which there exists a runtime semantics that *guaranteed to be efficient in $|\varphi|$*?

  Definition 1 (Tractable runtime). *Let $p \Downarrow (v, n)$ be a cost-annotated big-step reduction relation for some program p written in language L. We call $\Downarrow$ a **tractable runtime** if n is polynomial in $(|\varphi|)$ for any program p written in L.*

- We need to define a notion of size for PROP programs. Intuitively, the size simply counts how many syntactic constructs are utilized:

$$|\texttt{true}| = 1 \qquad |\texttt{false}| = 1 \qquad \frac{|\varphi| = n}{|\neg\varphi| = 1 + n} \qquad |x| = 1$$

$$\frac{|\alpha| = n_1 \qquad |\beta| = n_2}{|\alpha \wedge \beta| = 1 + |\alpha| + |\beta|}$$

And so on similarly for other connectives.

- Observe: neither $\Downarrow^e$ nor $\Downarrow^m$ are tractable runtimes.
  Exercise: find a family of programs that witnesses this intractability for each language.

- Can we come up with a suitable restriction to PROP that guarantees that $\Downarrow^e$ is a tractable runtime?

- **Idea**, our first tractable runtime: restrict propositional terms to be "pre-split" (or *Shannon-expanded*). Call this language PROP$_S$:

```
1  φ_S, α_S, β_S ::= true_S | false_S | (x_S∧α_S) ∨ (¬x_S∧β_S)
2  w_S ::= [x_S ↦ θ_x, ⋯]
3  p_S ::= (φ_S, w_S)
```

- We define a cost-annotated enumeration relation $p_S \Downarrow^e (v, n)$ in a manner identical to PROP for PROP$_S$

Theorem 1. *Assume $p_S \Downarrow^e (v, n)$. Then, $n \leq 1 + |p_S|$.*

Proof. By structural induction on (φ, w), inducting on φ. Base cases are easy: $(\mathsf{true}, w) \Downarrow (1, 1) \leq 1 + 1$.

Now let's handle $((x \wedge \varphi_1) \vee (\neg x \wedge \varphi_2))$. Assume by induction that $\varphi[\mathsf{true}/x] \Downarrow (v_1, n_1)$, $\varphi[\mathsf{false}/x] \Downarrow (v_2, n_2)$, and $n_1 \leq 1 + |\varphi_1|$ and $n_2 \leq 1 + |\varphi_2|$.

Then, by applying the (SPLIT) rule:

$$\frac{\varphi[\mathsf{true}/x] \Downarrow (v_1, n_1) \qquad \varphi[\mathsf{false}/x] \Downarrow (v_2, n_2)}{((x \wedge \varphi_1) \vee (\neg x \wedge \varphi_2)) \Downarrow^e (w(x)v_1 + w(\bar{x})v_2, 1 + n_1 + n_2)}$$

Now we are done if we can conclude that $n_1 + n_2 \leq |\varphi|$. Continuing:

$$
\begin{aligned}
n_1 + n_2 &\leq 1 + |\varphi[\mathsf{true}/x]| + 1 + |\varphi[\mathsf{false}/x]| && \text{by I.H.} && (1) \\
&\leq 1 + |\varphi_1[\mathsf{true}/x]| + 1 + |\varphi_2[\mathsf{false}/x]| && (\star) && (2) \\
&\leq 1 + |\varphi_1| + 1 + |\varphi_2| && (\star\star) && (3) \\
&\leq |(x \wedge \varphi_1) \vee (\neg x \wedge \varphi_2)| && (\dagger) && (4)
\end{aligned}
$$

where $(\star)$ follows from $\varphi[\mathsf{true}/x] = (x \wedge \varphi_1) \vee (\neg x \wedge \varphi_2)[\mathsf{true}/x] = \varphi_1[\mathsf{true}/x]$;

$(\star\star)$ follows from a simple lemma: substitution can only reduce the size of the formula, i.e. $|\varphi[v/x]| \leq |\varphi|$ for any variable x and value v.

$(\dagger)$ follows from the definition of $|\cdot|$. $\qquad\square$

Semantics-preserving translations

- **Idea**: we can give a semantics-preserving translation from an intractable language into a tractable one

- Why would we want to do that?

- – Abstraction: just like we don't want all programmers to program in assembly, it's useful to identify good target languages
- – Modular reasoning: We can separate our efficiency argument into 2 stages: proving the target efficient, and proving the translation efficient (or characterizing the efficiency of the translation)

- Compilation from PROP to PROP$_S$: define a relation $\rightsquigarrow$ that translates each term in each language while preserving semantics:

Figure 1: Compilation from PROP to PROP$_S$.

$$\text{true} \rightsquigarrow \text{true}_S \qquad \text{false} \rightsquigarrow \text{false}_S \qquad x \rightsquigarrow x_S$$

$$\frac{x \text{ free in } \varphi \qquad \varphi[\text{true}/x] \rightsquigarrow \varphi_1 \qquad \varphi[\text{false}/x] \rightsquigarrow \varphi_2}{\varphi \rightsquigarrow (x_S \wedge \varphi_1) \vee (\neg x_S \wedge \varphi_2)}$$

- We will want to prove this compilation correct by proving that it preserves denotation:

Definition 2 (Semantics preserving compilation). *Let L_1 and L_2 be languages with semantic interpretations $[\![\cdot]\!]$ with the same semantic domain. A compilation $\rightsquigarrow$ from language L_1 to language L_2 is **semantics-preserving** if, for any program p_1 written in L_1 where $p_1 \rightsquigarrow p_2$, it is the case that $[\![p_1]\!] = [\![p_2]\!]$.*

**Theorem 2.** *The compilation rules in Figure 1 are semantics-preserving.*

- In a seminal paper Felleisen [1990] distinguished between different kinds of programming languages based on their supported syntactic features. He compared languages on the basis of *expressivity*: whether or not it was possible to express all programs in one language in the other.

- Here we consider a refinement of this notion to *efficient expressiveness*, which captures whether or not an *efficient* (i.e., polynomial-time) translation exists between two languages:

Definition 3 (Efficient expressiveness). *A program p_1 written in language L_1 is **efficiently expressible** as a program p_2 written in language L_2 if there exists a polynomial-time semantics-preserving translation $p_1 \rightsquigarrow p_2$. We say language L_1 is **efficiently expressible** as L_2, written $L_1 \sqsubseteq L_2$, if all programs in L_1 are efficiently expressible as programs in L_2.*

- Observe: PROP$_S$ is efficiently expressible as PROP, but not vice versa

3   Binary decision diagrams

- **Problem**: PROP_S is not very expressive

- **Goal:** More expressive tractable languages

- What can we improve about PROP_S? Observe that there may be *repetitious sub-syntactic terms*, and we'd like to exploit those, just like we observed in our memoized reduction scheme

- How can we represent repetitious subterms syntactically? By changing our syntax to permit *directed acyclic graphs*!

- Syntax of **binary decision diagram** (BDD): a rooted binary directed acyclic graph with two kinds of nodes:

  - *Choice nodes*, written $\text{Ite}(x, l, h)$, where x is a propositional variable (called the top variable), l is an edge called the *low edge*, and h is an edge called the *high edge*

  - *Terminal nodes*, which are labeled `true` or `false`.

- The syntax of BDD is given as graphical notation, for example:
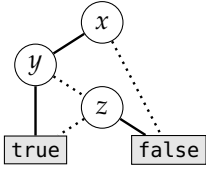
  Solid edges are high edges, dotted edges are low edges

- Semantics of BDD:

$$\left[\!\!\left[\Gamma \vdash \boxed{\texttt{true}}\right]\!\!\right] = [\![\Gamma]\!]$$

$$\left[\!\!\left[\Gamma \vdash \boxed{\texttt{false}}\right]\!\!\right] = \varnothing$$

 $= \{I \in [\![y]\!] \mid I(x) = \texttt{true}\} \cup \{I \in [\![z]\!] \mid I(x) = \texttt{false}\}$

References

Paul Beame and Toniann Pitassi. Propositional proof complexity: Past, present. *Future*, pages 42–70, 2001.

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Gerhard Gentzen. Investigations into logical deduction. *The collected papers of Gerhard Gentzen*, pages 68–131, 1969.