

Maths -2

Question:

N doors numbered 1 to N ,

$$N = 10$$

No. of open doors.

(N)

$$P_1 \Rightarrow 1, 2, 3, 4, \dots, 10$$

$$P_2 \Rightarrow 2, 4, 6, 8, 10$$

$$P_3 \Rightarrow 3, 6, 9$$

$$\vdots$$

$$P_N \Rightarrow$$



P_1	o	o	o	o	o	o	o	o	o
P_2	o	c	o	c	o	c	o	c	o
P_3	o	c	c	c	o	d	o	c	c
P_4	o	c	c	o	o	o	o	c	c
P_5	o	c	c	o	c	c	c	c	o
	o	c	c	o	c	c	c	c	o

P4
7

$$\text{Ans} = 3$$

14,9 \angle 10°

brute force:

~~for (i=0; i<n; i++) {
 for ()~~

$$T-C = O(n^y)$$

Efficient Approach

Door [N] how many door people will
foggle this (38) ↗

$$\text{Factors of } 18 = 1, 2, 3, 6, 9, 18$$

All factors of 18

{ 1 F : O
 2 F : C
 3 F : O
 n F : C

18 : 1, 2, 3, 6, 9, 18

Factors will come
 $(1, 18)$ $(2, 9)$ $(3, 6)$

$$(x_1, y_1) = 16$$

in pairs
Even factors

$$\Rightarrow C \in \mathcal{O}^{\text{red}}$$

$$Z \in \mathbb{N}^{2m \times p}$$

$$\Rightarrow \frac{1}{n} \sum_{i=1}^n \left(1 + \frac{1}{i} + \frac{1}{j} + a_i - \right)$$

HP

$$16 = 4^2 \Rightarrow 1, 2, 4, 8, 16$$

[Odd Factors]

$(1, 16) (2, 8) (4,$

 $1^2, 2^2, 3^2, 4^2, 5^2, 6^2, 7^2$

Odd No of factors

$$\begin{array}{|c|c|c|} \hline N=11 & & \\ \hline 1, 4, 9 & \rightarrow & \text{Open} \\ \hline \end{array}$$

No of Perfect squares $\leq N$

Solution:

$$\begin{aligned} N &= 10^0 \Rightarrow 1^0 \\ N &= 11^0 \Rightarrow \left\lceil \text{Floor}(\sqrt{N}) \right\rceil \\ \sqrt{10} &= \frac{10.5}{1} \end{aligned}$$

$$1^2, 2^2, \dots, 8^2, 9^2, 10^2, 100 \Rightarrow 100$$

Question: Rearrange the array elements such that

$$a[i] \rightarrow a[a[i]]$$

$a[a[i]] \in [0, n-1]$

$O(1)$ space
 $O(n)$

$N=4$

	0	1	2	3
$A[0]=3$	3	2	0	1
$A[1]=0$	1	0	3	2
$A[2]=3$	0	3	2	1
$A[3]=2$	2	1	3	0

$a[0] \rightarrow a[a[0]] \rightarrow a[3]$

$a[2] \rightarrow a[a[2]] \rightarrow a[0] \rightarrow 3$

$$a[i][j] \rightarrow a[a[i][j]]$$

$$\rightarrow a[2]$$

$$\rightarrow 0$$

$$a[i][j] \rightarrow a[2][1]$$

$$\rightarrow 2$$

for every i' ,
 $\text{ans}[i][j] = a[i][a[i][j]]$

$$T.C: O(n)$$

$$S.C: O(n)$$

$$a[i][j] \rightarrow a[a[i][j]]$$

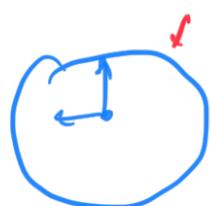
	0	1	2	3	4	$a[i][j] \rightarrow a[a[i][j]]$
$A[0][z]$	4	0	2	1	3	
$A[1][z]$	3	4	2	0	1	
$A[2][z]$	0	1	2	3		\downarrow
$A[3][z]$	1	0	0	1		\downarrow
$A[4][z]$						\downarrow

$$\text{temp} = \underbrace{3}_{1} \underbrace{2}_{O(n)}$$

$O(n)$ extra space

[old val, new val]

Approach 2:



n by n

9 AM | 9 PM ?

$(0-1)$

$21 =$

$$9 \rightarrow (21) \quad \%12 = [0-11]$$

$$1 \times 12 + 9 \quad 12 \text{ mod}$$

$\boxed{11111}$

⑨, $(21) \rightarrow$ AM / PM
 Time of the 12hr slot

$$(21) = 1 \times 12 + 9$$

$$21 \rightarrow 21 \% 12$$

$$\dots 9 \% 12$$

21

$$n \% 12 = 0, 1$$

$\frac{1}{n}$ 21 to 14

(old + n + new val)



Old Val
New Value

Org:

3	2	0	1
---	---	---	---

Arr:

1	0	3	2
---	---	---	---

$$((3 \times 4) + 1) \text{ mod } 0$$

$$13 \mod 4 \rightarrow 1$$

$$3 \rightarrow \lceil N \rceil$$

$$8 \mod 4 \rightarrow 0$$

$$2 \rightarrow \lfloor N \rfloor$$

$$1 \rightarrow \text{New value}$$

$$\text{old} = 3$$

$$\text{new} = 1$$

(N)

[0...n-1]

(N) % n

$$a[0] \rightarrow a[\underline{s}]$$

(2)

$$\begin{cases}
 A_{\text{arr}} = [3, 2, 0, 1] \\
 \text{Step 1} \Rightarrow A[0] = 3 \\
 \text{Step 2} \Rightarrow A[1] = 2 \\
 \text{Step 3} \Rightarrow A[2] = 0
 \end{cases}$$

Step 2 =

$$\begin{array}{ccccccccc}
 & & 3 & & 2 & & 0 & & 1 \\
 & & \downarrow & & \downarrow & & \downarrow & & \downarrow \\
 & & 13 & & 8 & & 3 & & 6
 \end{array}$$

Step 3 =

$$\begin{array}{ccccccccc}
 & & 3 & & 2 & & 0 & & 1 \\
 & & \downarrow & & \downarrow & & \downarrow & & \downarrow \\
 & & 13 & & 8 & & 3 & & 6
 \end{array}$$

$$a[0] \rightarrow a[\underline{s}]$$

$$a[1] \rightarrow a[\underline{\underline{s}}]$$

$$a[0] \rightarrow a[\underline{s}]$$

$$a[1] \rightarrow a[\underline{\underline{s}}]$$

$$a[2] \rightarrow a[\underline{\underline{\underline{s}}}]$$

$$a[0] \rightarrow a[\underline{s}]$$

$$a[1] \rightarrow a[\underline{\underline{s}}]$$

$$a[2] \rightarrow a[\underline{\underline{\underline{s}}}]$$

new value

algo :-

Algo

- 1) Multiply by N
- 2) $\text{arr}[i] += \text{arr}[\lceil \text{arr}[i]/N \rceil] \% N$
- 3) $\text{arr}[i] = \text{arr}[i] \% N$

$A[4] =$	1	0	3	2	
	0	1	2	3	$N=4$
$A =$	3	2	0	$\underline{1}$	
$A =$	$\underline{\underline{12}}$	8	0	$\underline{\underline{9}}$	$a[3]$
$A =$	$\underline{\underline{4}}$				$a[a \% 2]$
					$a[n \% 4]$
					$\underline{\underline{5-5}}$

Question.

Given a prime number

$$(p^2 - 1)$$

$$p \leq 10^{1000}$$

$$p \geq 3$$

$$p = "903 \dots \dots "$$

$$p \geq 2$$

check if $\frac{(p^2 - 1)}{24} \% 24 = 0$?

$$p = 5, \quad (25 - 1) \% 24 \rightarrow \text{True}$$

$$p = 7, \quad (49 - 1) \% 24 \rightarrow 0$$

$$B \quad 24 - 1 = (20) \% 24 \\ \dots \dots \dots$$

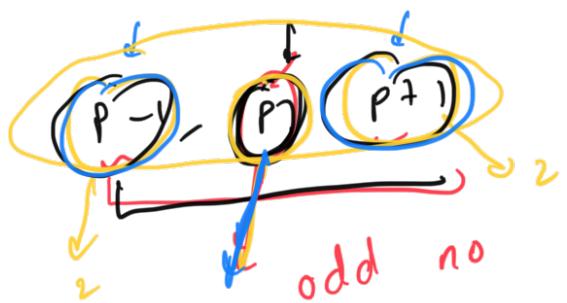
2^u = 2 * 2 * 2 * 5

$$2^u = 2 \times 2 \times 2 \times 5$$

$(p^2 - 1) = (p-1)(p+1)$

$$(a^2 - b^2) = (a-b)(a+b)$$

$a = n+1, b = n+2$



consecutive nos

$u = i's$ $x^{1/3}$ p is not divisible by 3



TRUE

$$(p^2 - 1) \% 2^u = 0$$



14 15 16

%3?

45 46 47

$$\frac{1}{n} \frac{n+1}{2} \frac{n+2}{3}$$

$(1) \frac{2}{2} \frac{3}{3}$

9. 4



13

119.

$$(p^{r-1})^4 \cdot 2^4$$

$$(1^{r-1})$$

$$(1^{(r-1)})$$

$$(1^{(2r-1)})$$

$$(1^{(2r-1)})$$

$$2^{4n} \cdot 5^{r-1} \cdot 10^r$$

Question:

N people \rightarrow standing in a circle RH.Hers

$$\boxed{N = 5}$$



\Rightarrow

$$\boxed{\deg = 3}$$

$N = 9$



$$\boxed{deg = 5}$$

$$\boxed{(3)}$$

$$\boxed{(4)}$$

$$\boxed{(5)}$$

$$\boxed{(6)}$$

$$\boxed{(7)}$$

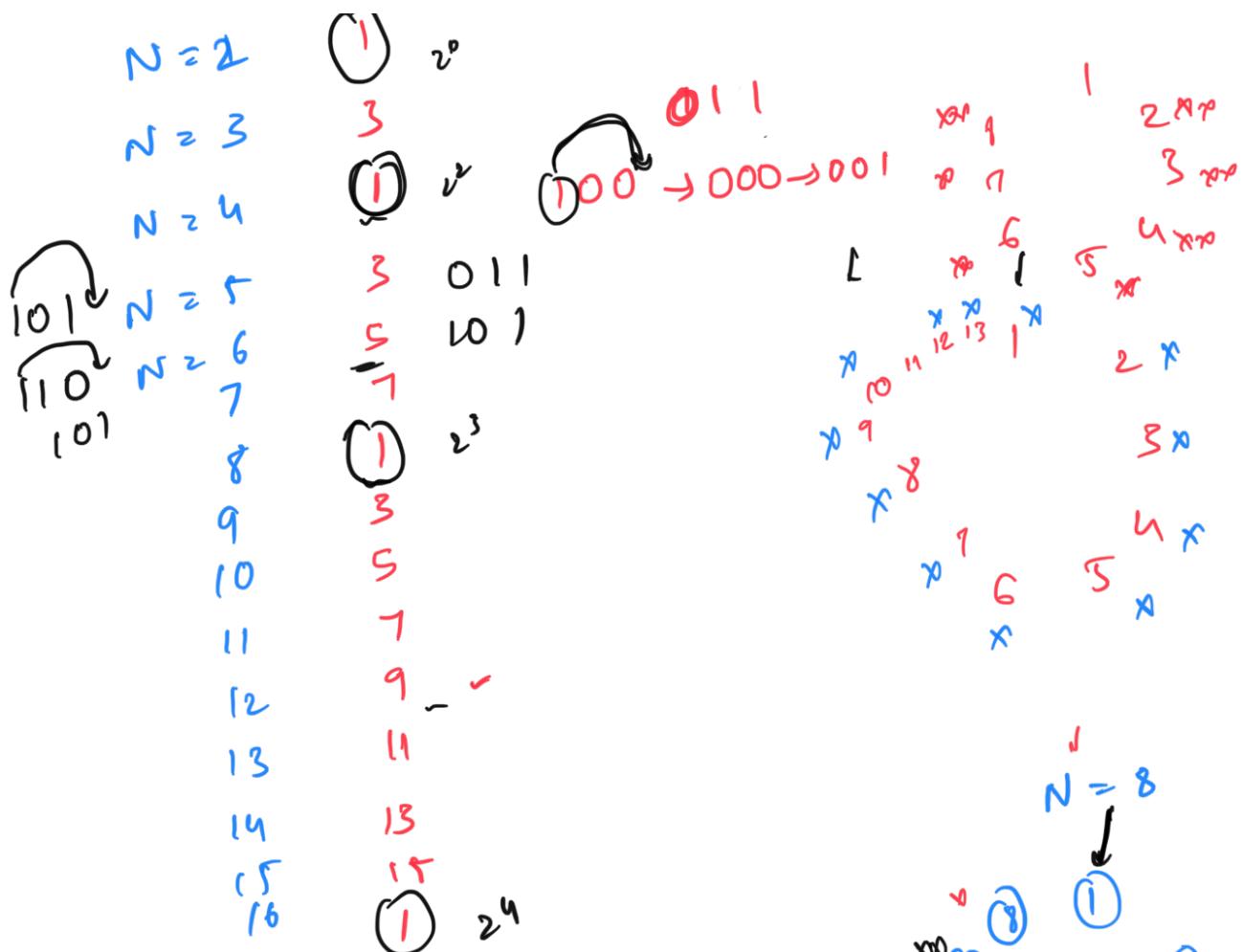
$$\boxed{(8)}$$

$$\boxed{(9)}$$

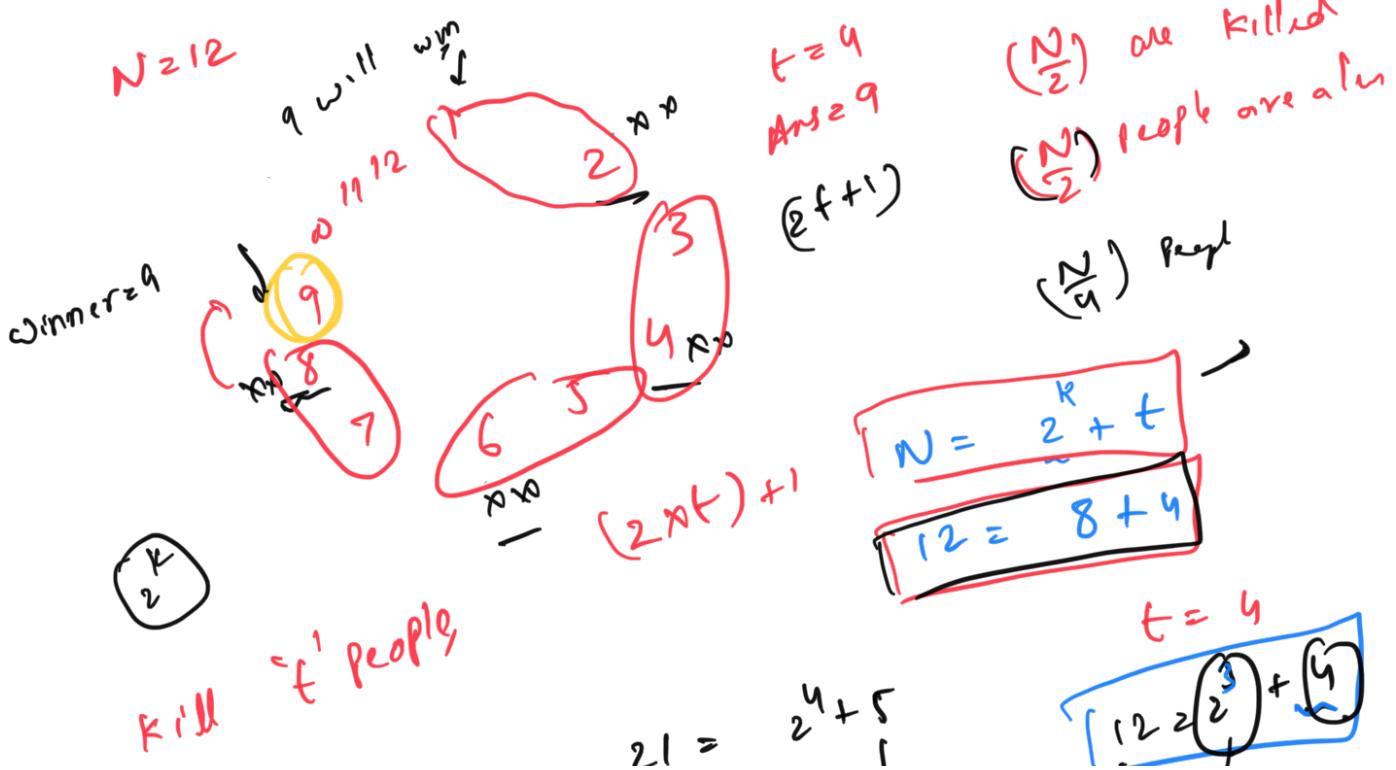
$N = 10^9$

$N = 1$





$$N \rightarrow \frac{N}{2} \rightarrow \frac{N}{4} \rightarrow \frac{N}{8} \dots \textcircled{1}$$



t

$\frac{2^3}{2}$

$$28 = 16 + (2)$$

$2t+1 = 25$

$$28 = 2^4 + 12$$

t

2^4
 t

$(2t+1)$ will be
winner

$(2t+1)$

$$32 = 2^5 + 0$$

$t=0$

$$2^{xt+1} = 0$$

$= 1$

Algo

(1)

$$18 = \rightarrow 10010$$

(F)

$$\rightarrow 10000$$

$$18 = 16 + 2$$

New $2 = 00010$

(N)

$$t = N-2$$

$$16 = 2^4$$

$$\log_2^{16} = 4$$

$$(\log_2^{18})_2 4 \dots$$

$$18 = 2^4 + \dots$$

$$\text{Floor}(\log_2^n)$$

t
 $t+1$

$$\frac{2 \times t}{2t+1} + 1$$

Algo

(1) Remove

MSB

2) Left Shift
3) 1 in USP

Question: N^{th} Magic Number

→ Magic No can be either represented as a unique power of 5 or sum of powers of 5.

$$5, 5^2, 5^3, 5^3 + 5^2, 5^3 + 5^2 + 5^1, 5^3 + 5^2 + 5^1 + 5^0, 625, 630, 650, 655, 250$$

↓
Ascending Order

$N=10$

↓
Binary Rep

$$N=1 \rightarrow 5$$

$$N=2 \rightarrow 5^2$$

$$N=3 \rightarrow 5^3 + 5$$

$$N=4 \rightarrow 5^3 + 0.5^2 + 0.5$$

$$N=5 \rightarrow 5^3 + 5 \rightarrow 5^3 + 0.5^2 + 1.5$$

$$N=6 \rightarrow 5^3 + 5^2 \rightarrow 5^3 + 5^2 + 0.5$$

$$N=7 \rightarrow 5^3 + 5^2 + 5$$

$$N=8 \rightarrow 5^4$$

$$N=9 \rightarrow 5^4 + 5^1$$

$$N=10 \rightarrow 5^4 + 5^2$$

$$N=11 \rightarrow 5^4 + 5^3$$

$$1 \ 0 \ 1 \ 0 \rightarrow 2^{3 \times 1} + 2^{2 \times 0} + 2^1 \times 1 + 2^0 \times 0$$

$K=2$

Magic N°:	\sum of unique powers of 2	Sum of powers of 2
N		
001	$2^0 + 2^1$	$2^0 + 2^1 + 2^0 \times 1$
010	$2^0 + 2^2$	$2^0 + 2^2 + 2^0 \times 1$
011	$2^0 + 2^1 + 2^2$	$2^0 + 2^2 + 2^1 + 2^0 \times 1$
100	$2^0 + 2^3$	$2^0 + 2^3 + 2^0 \times 1$
101	$2^0 + 2^3 + 2^1$	$2^0 + 2^3 + 2^1 + 2^0 \times 1$
110	$2^0 + 2^3 + 2^2$	$2^0 + 2^3 + 2^2 + 2^0 \times 1$
111	$2^0 + 2^3 + 2^2 + 2^1$	$2^0 + 2^3 + 2^2 + 2^1 + 2^0 \times 1$
1000	$2^0 + 2^4$	$2^0 + 2^4 + 2^0 \times 1$

problem \leftarrow

$$10110 = 0 \times 2^0 + 1 \times 2^1 + 1 \times 2^2 + 0 \times 2^3 + 1 \times 2^4 + 1 \times 2^5$$

T.C:

$\rightarrow O(\log_2 n)$

$$N_2 \quad \begin{matrix} 1 & 0 & 1 & 1 & 0 \\ \hline 6 & 5 & 4 & 3 & 2 \end{matrix}$$

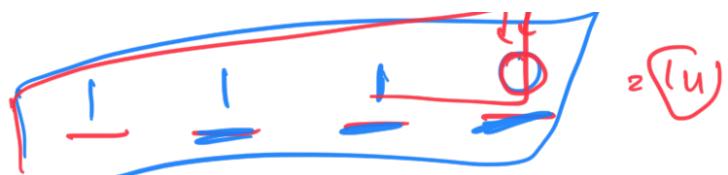
6 bits

%2

(14) =

++ -

(n%2)



$$\begin{aligned} N &= N/2 \\ &= 7 \end{aligned}$$

$\text{ans} = 0$

bit = 1

$1, 5, 5^2 ?$

$7 \text{ in } \\ 7 = 7/2 \quad 5$

$3/2 = 1$

while ($N > 0$) {

LSB = $\underline{\underline{N}}/2;$

ans = ans + LSB * pow(5, bit)

bit++;
 $N = N/2;$

}
(Binary Exponent)
O(log n)

T.C: $\log(n)$

101110

$\rightarrow O(N \cdot 32 \text{ bits})$

18 = 10010

16 = 10000
17 = 10001

$a^b \Rightarrow O(\log b)$

binary exponent

$\log_2^{16} = 4$
 $= 2^4$

10000
11110000
(5)

$\log(17)$
 $= 4.08$

32

$\rightarrow \log(n)$

No- θ bits =

$\text{floor}(\log n)$

$O(\log n)$

$O(32)$
 $\geq O(1)$

$32 \text{ bits} \Rightarrow O(32 \text{ bits})$
 $= 11$

$O(\log h)$

(2×10^9)

$A[0..N-1] = [3, 4, 2, 0, 1]$

$A \rightarrow$

Step 1 → $20 \quad 0 \quad 10 \quad 5 \quad 15$

Step 2 → $20+3 \quad 0+4 \quad 0+2 \quad \sum_{i=0}^{4-1} A[i]$

Step 3 → $\frac{20+3}{5} \quad \frac{0+4}{5} \quad \frac{0+2}{5} \quad \frac{\sum_{i=0}^{4-1} A[i]}{5}$

$$arr[0] = arr[0] + arr[1]/N; \\ arr[1] = arr[1] + arr[2]/N; \\ arr[2] = arr[2] + arr[3]/N; \\ arr[3] = arr[3] + arr[4]/N;$$

$a[0]/5$

$$a[1] = a[1] + a[a[0]/N]/N \\ + a[0]/N \\ 2a[1]$$

$$a[2] = a[2] + a[a[1]/5]/5$$

$$= (0 + a[2]/5)$$

$$10 + 2$$

Array = $O(n)$ space
 $O(n^2)$ space

$$a[3] = a[3] + a[a[2]/5]/5$$

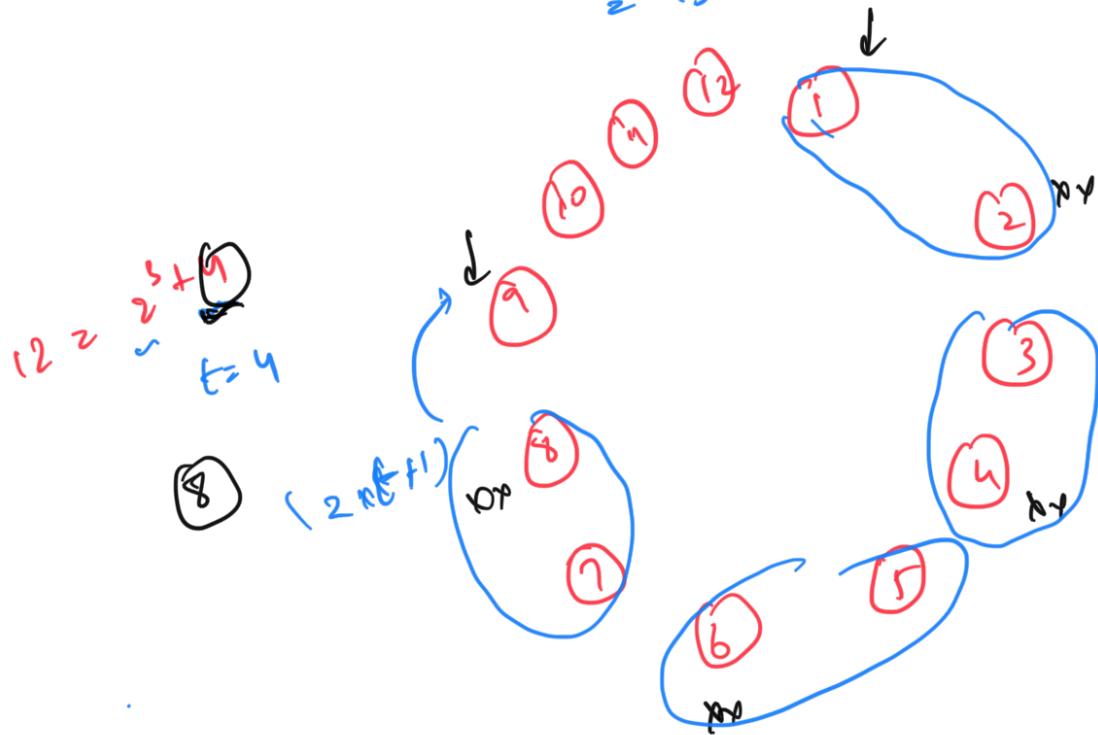
$$= 5 + a[4]/5$$

$$= 5 + 0/5$$

$$= 5 + 0$$

$$\dots = 1 - a[4] + a[a[3]/5]/5$$

$$\begin{aligned}
 & a \Delta M Z = u \cdot v \\
 & = 15 + a \{ (5/\Gamma) / 5 \} \\
 & = 15 + a \{ 3 \} / \Gamma \\
 & = 15 + 5 / \Gamma \\
 & = 15 + 1
 \end{aligned}$$



$18 =$ Even Factors \Rightarrow $\{$

$$\begin{aligned}
 18 = & \\
 (1, 18) & \\
 (3, 6) & \\
 (2, 9) &
 \end{aligned}$$

$$\begin{aligned}
 16 = & \\
 (1, 16) & \\
 (2, 8) & \\
 (4, 4) & \\
 (1, 16) &
 \end{aligned}$$

1, 2, 4, 8, 16

$$1^2, 2^2, 3^2, 4^2$$

$$N \neq 100$$

1, 2, 2, 3, 2, 4, 5, 5, 6, 2, 7, 2, 8, 2, 9, 2, 10

Nth Magic Number

```
ans = 0;
int bit = 1;
while(N > 0) {
    ans += N%2 * pow(5, bit);
    N = N/2;
    bit++;
}
```

Rearrange the array

```
for(int i = 0; i<n; i++)
    arr[i] = arr[i] * n;
for(int i = 0; i<n; i++)
    arr[i] += arr[arr[i]/n]/n;
for(int i = 0; i<n; i++)
    arr[i] = arr[i] % n;
```