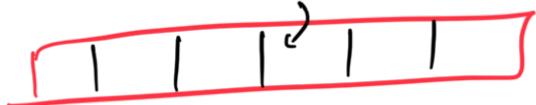


Linked List - I

Bigest advantage of array?

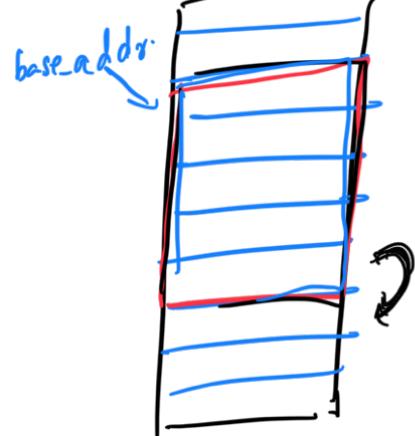


- Fetch the i^{th} index

$O(1)$ Random

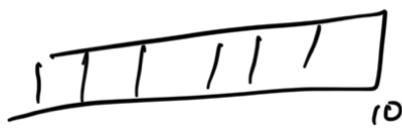
$$\text{addr}(i) = \text{base-addr} + 4 + i$$

$\underline{O(1)}$



$O(1)$

Dynamic Array
Vector / ArrayList



Dynamic Array \Rightarrow Fixed Size array



\Rightarrow



\rightarrow Unused space
 y

99 elmt

Linked List

If j have to
will only

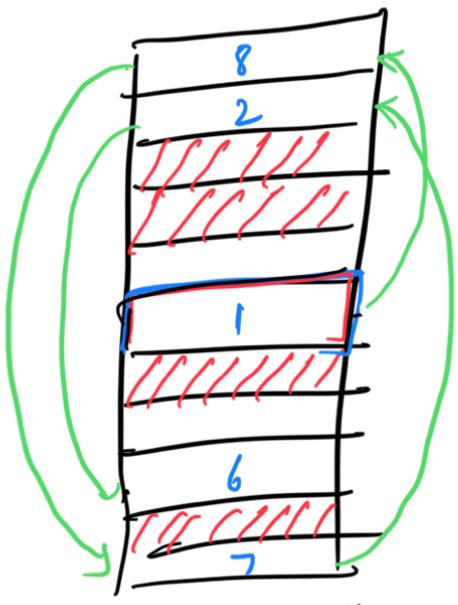
store
allocate 100 elements, 1
space for 100 elmt

$A = [1, 8, 7, 2, 6]$

stored in contiguous locations? Memory

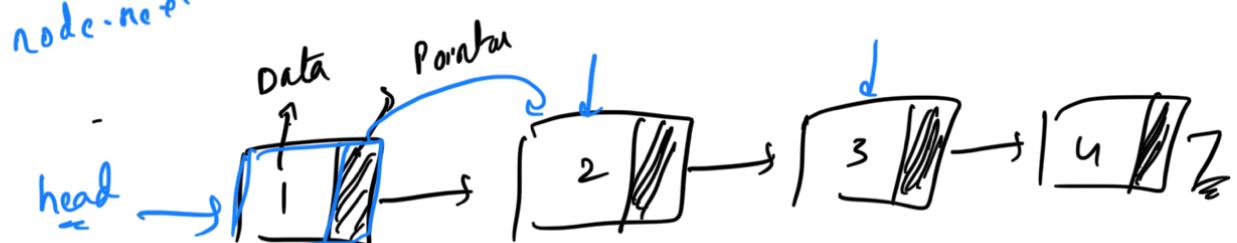
Every node stores the address/reference to the next node.

Linked List: Nodes are linked to each forming a list.



Effective Space Utilization

Extra space for pointers cost 4 bytes



Pointer to the first node of LL: head pointer





if (head == NULL)
Empty Linked List

C++

```
struct Node {
    int data;
    float y;
    string s;
    Node * next;
} Data
// Pointer to next Node
Node obj = new Node();
```

Java / Python

```
class Node {
    int data;
    Node next; // Object of the same class
} References to the object
// Next Pointer.
```

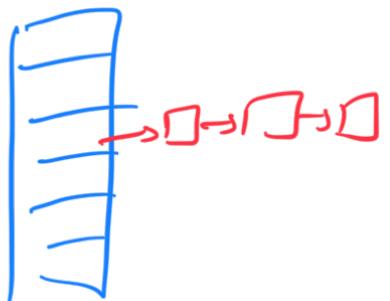
```
string name;
int age;
string address;
char Gender;
vector<string> m;
```

Node



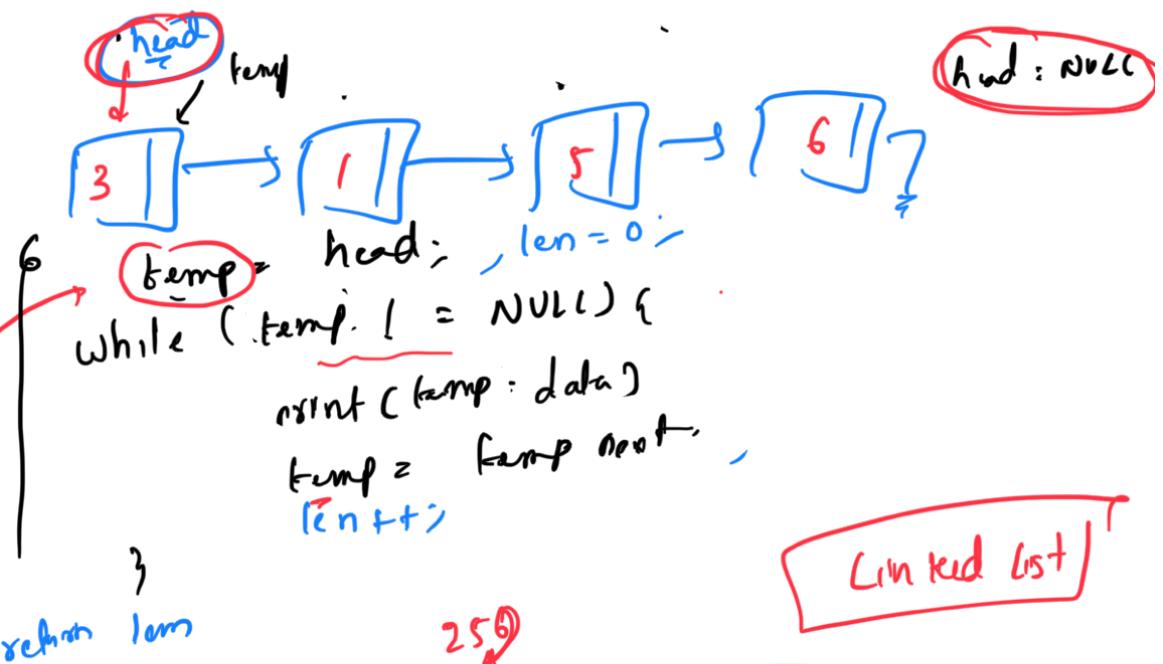
Applications / Uses

- 1) LRU Cache \Rightarrow Doubly LL
- 2) Stack & Queue using LL \times
- 3) Graphs to represent adjacency List \Rightarrow
- 4) Hashing (Chaining)



99% Edge Cases ↗

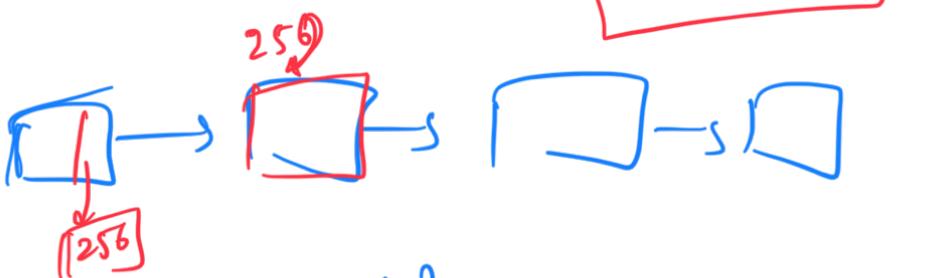
Question: Print a Linked List



3. 1. 5. 6
 $\text{temp} \rightarrow \text{head}; \text{len} = 0;$
 $\text{while } (\text{temp}. \text{l} = \text{NULL}) \{$
 $\quad \text{print}(\text{temp}. \text{data})$
 $\quad \text{temp} = \text{temp}. \text{next},$
 $\quad \text{len}++\}$

return len

Linked List

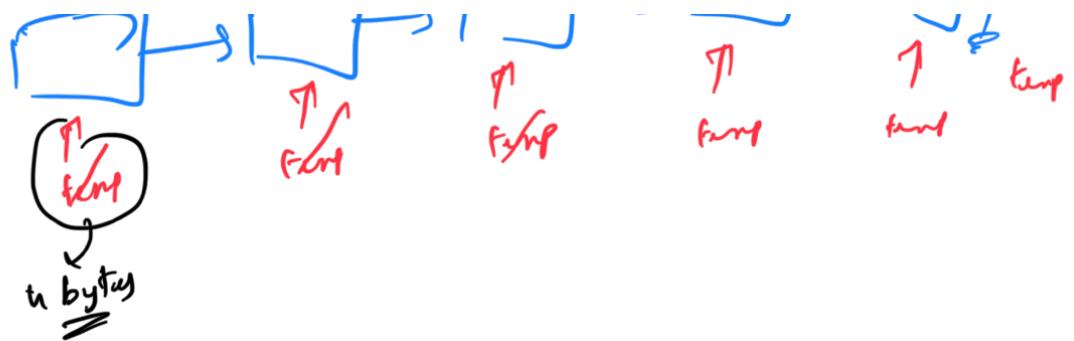


T.C: $O(n)$, S.C: $O(1)$

Question:



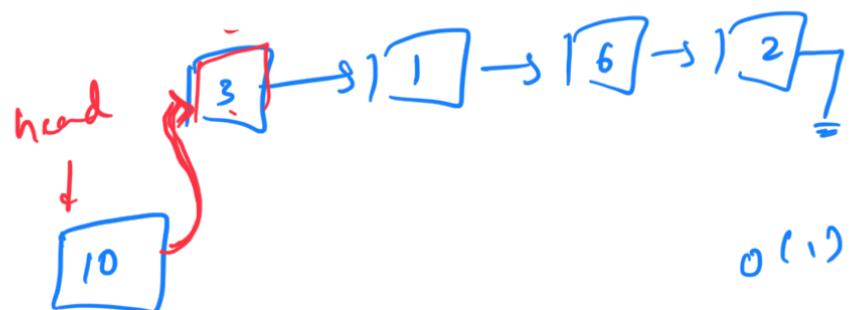
count = 12345
= 5



Operations :

insertion

i) Insert at the start



```
insert At Beginning (head, data){
    newNode = new Node (data);
    newNode . next = head;
    head = newNode;
    return head;
}
```

Delete from the start



head = head.next;

if (head == NULL)
 "none"

O(1)

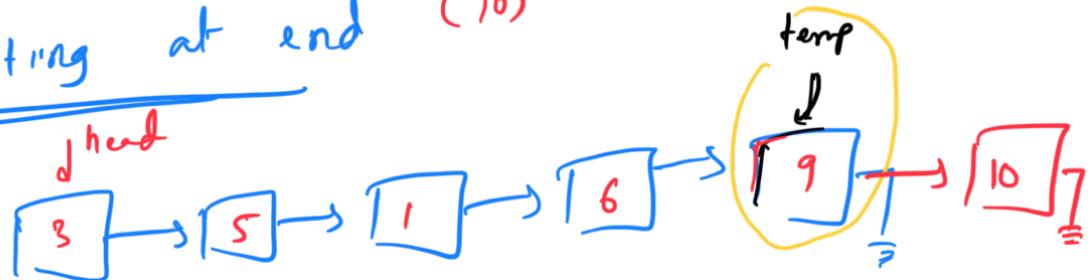
"**FEHLER**

T.C: $O(1)$

```

temp = head;
head = head.next;
free(temp);
    
```

2) Inserting at end (10)



$\text{head} = \text{NULL}$

```

temp = head;
if (temp == NULL) { head = newNode; }
while (temp.next != NULL) {
    temp = temp.next;
}
temp.next = newNode;
    
```

T.C: $O(n)$
S.C: $O(1)$

)

$\text{temp.next} = \text{newNode};$

$\text{newNode} \rightarrow \text{next} = \text{NULL}$

constructor (int data) {
 newNode = new Node();
 newNode.data = data;
 newNode.next = NULL;
 return newNode;
}

Delete the Last Node



```

if (temp == NULL) <-- free(head) ; new
if (head.next == NULL) <-- free(head) ;
while (temp.next != NULL) {
    temp = temp.next ;
}
Create (temp.next) ;
temp.next = NULL ;

```

T.C: $O(n)$
S.C: $O(1)$



insert at End: $O(1)$

S.C: $O(1)$
T.C: $O(1)$

```

insertAtEnd () {
    if (head == NULL) {
        head = tail = newNode;
    }
    tail.next = newNode;
    tail = newNode;      (tail = tail.next)
}

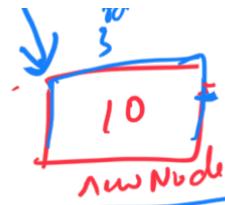
```

T.C: $O(1)$
S.C: $O(1)$



3) Insert at k^{th} Position ($k \neq 0$)
 $k=3$

$\text{temp_next} = \text{newNode}$



$\text{newNode.next} = \text{temp.next}$
 $\text{temp.next} = \text{newNode};$

Insert at K^{th} Position \Rightarrow j iterates fill $(K-1)^{th}$ node.

$$f(K=1)$$

$\text{temp} = \text{NULL}$

0^{th} Node

Edge Case:



1) Empty LL

2) Length 1 LL

3) Length 2/3 LL

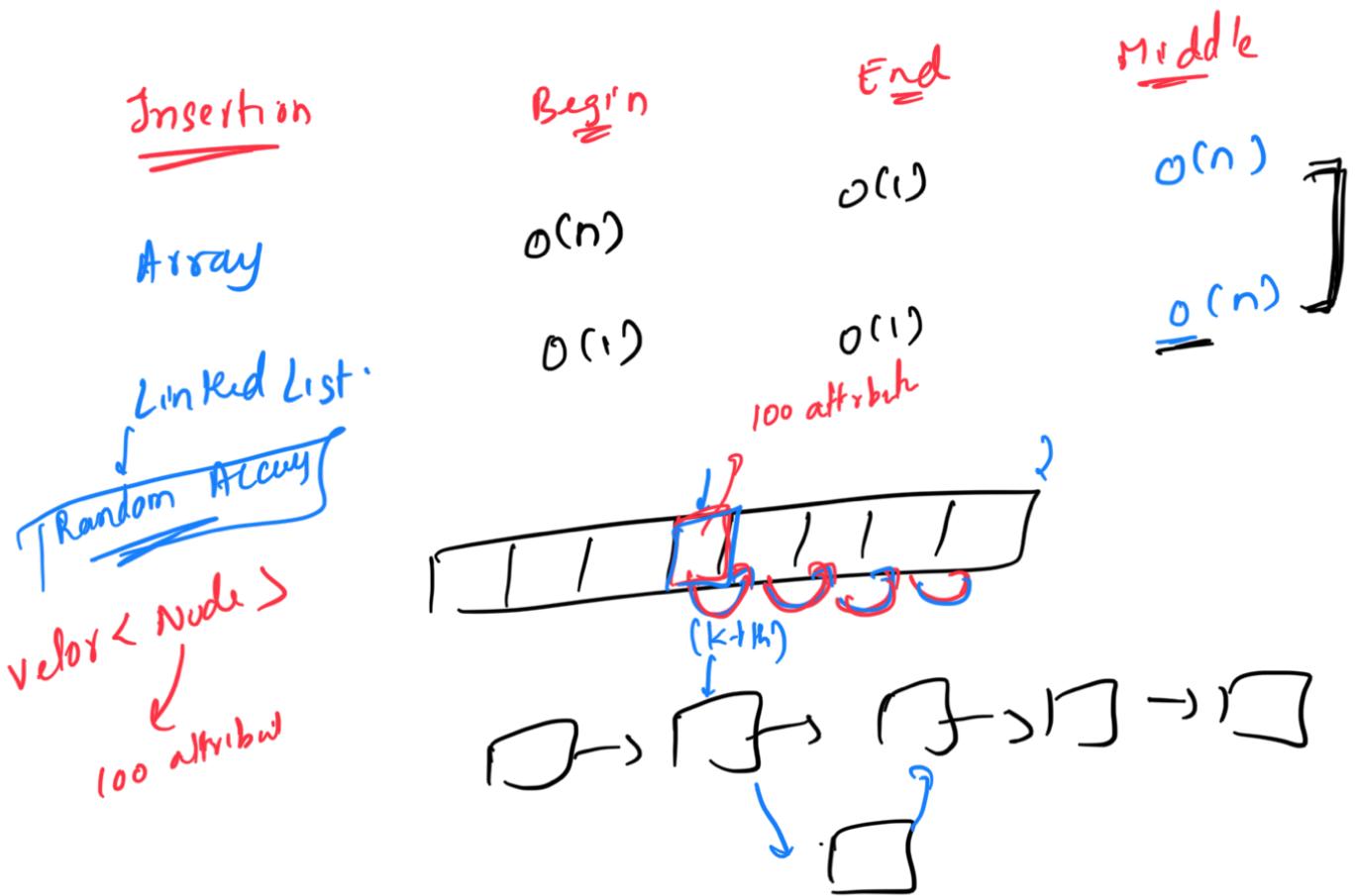
4) Normal Test Case

5) Problem Specific Problem

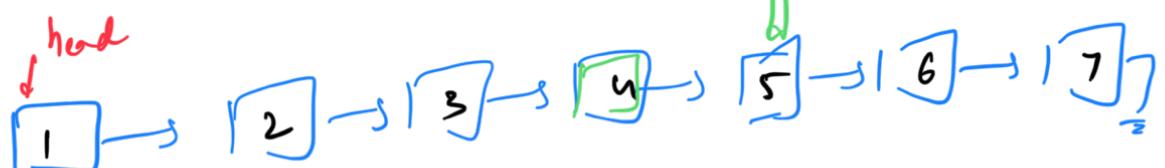
$(K=1)$



```
if (K == 1) {
    newNode.next = head;
    head = newNode;
```



Question: Find the k^{th} Node from the last



$$N = 7$$

k^{th} Node from last $\Rightarrow (N-k+1)^{\text{th}}$ node from first

$$N - k = 4$$

Approach 1

- 1) Compute length
- 2) Generate till the $(N-k+1)^{\text{th}}$ node

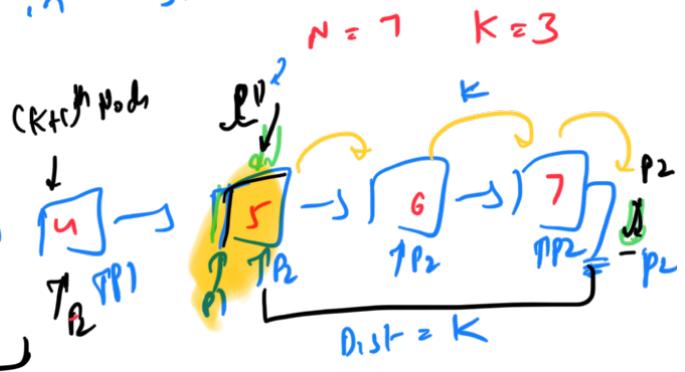
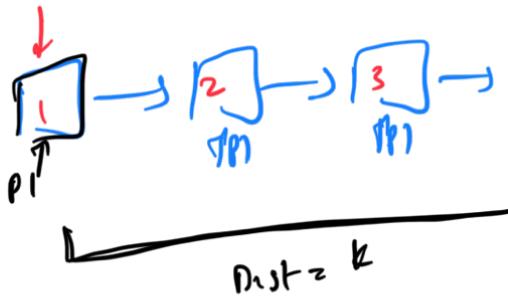
2 iterations

(2)

T.C: $O(n)$
S.C: $O(1)$

Approach 2: Expected in interviews

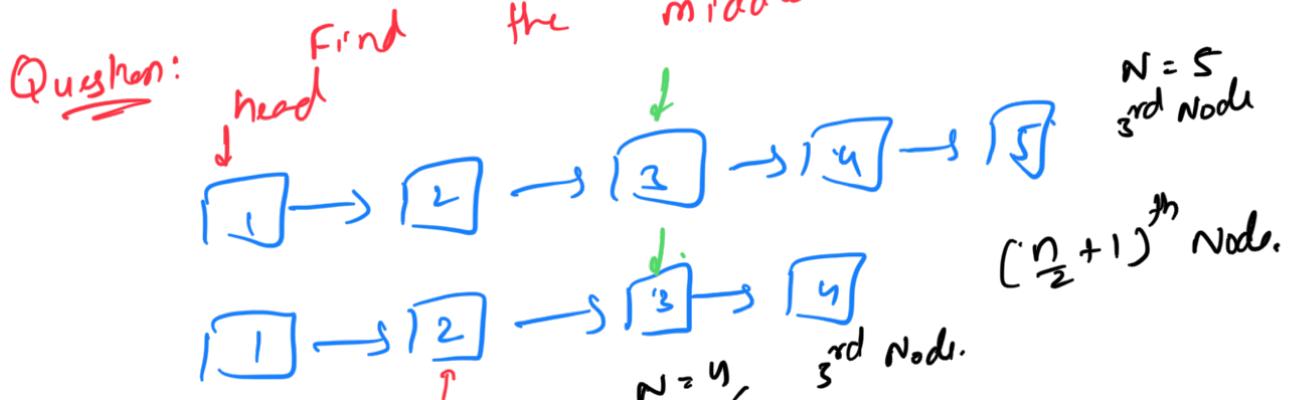
head



$\checkmark p_1 = \text{head}$ ✓
 $\checkmark p_L = (K+1)^{\text{th}}$ Node from head
 1 iteration \leftarrow $p_L = p_1 + (K+1)$
 iteration \leftarrow while ($p_2 \neq \text{NULL}$)
 $\quad\quad\quad p_1 = p_1 \cdot \text{next}$
 $\quad\quad\quad p_2 = p_2 \cdot \text{next}$
 \downarrow
 return p_1 ;

$$\begin{aligned}
 \text{T.C: } & O(K) + O(n) \\
 & O(n) + O(n) \\
 & = O(n)
 \end{aligned}$$

Question: Find the middle node of LL



Approach 1:

Find length of LL = n } 2 iterations

Find

the $(\frac{n}{2} + 1)^{\text{th}}$ node

T.C: $O(n)$

[Approach 2]

Car 1 \Rightarrow v
Soner.

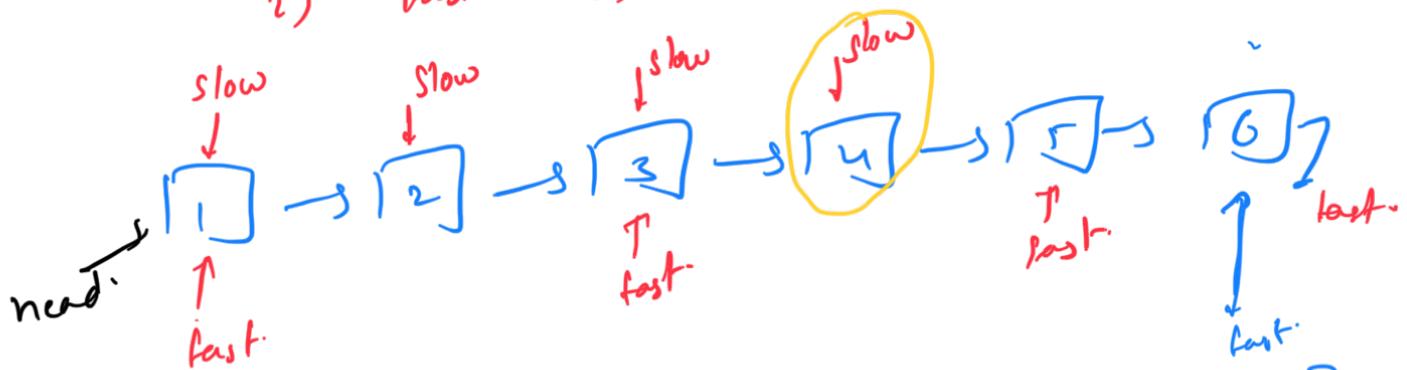
Car 2 \Rightarrow 2v

Car 1 \Rightarrow middle.
↓

Dest
Car

2 Pointers

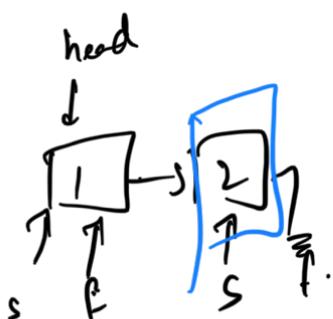
- 1) slow \Rightarrow v speed
- 2) fast- \Rightarrow 2v speed.



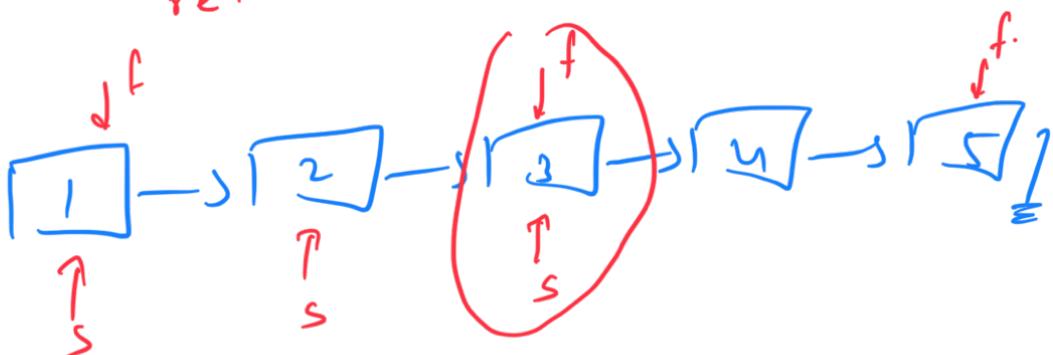
slow = head

fast = head;

while (fast != null) { }
 slow = slow.next;
 fast = fast.next.next;



return slow;



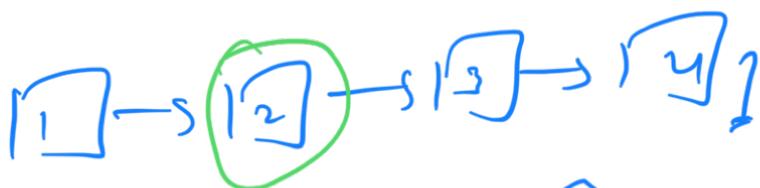
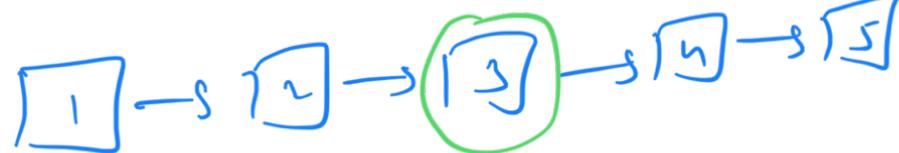
T.C: $O(n)$
 S.C: $O(1)$
 Only 1 iteration

Question:

slow = head
 fast = head.next;

{ HW }

Return the $(\frac{n+1}{2})^{\text{th}}$ Node



fast.next.next.next

\rightarrow fast != NULL
 \rightarrow fast.next != NULL
 \rightarrow fast.next.next != NULL

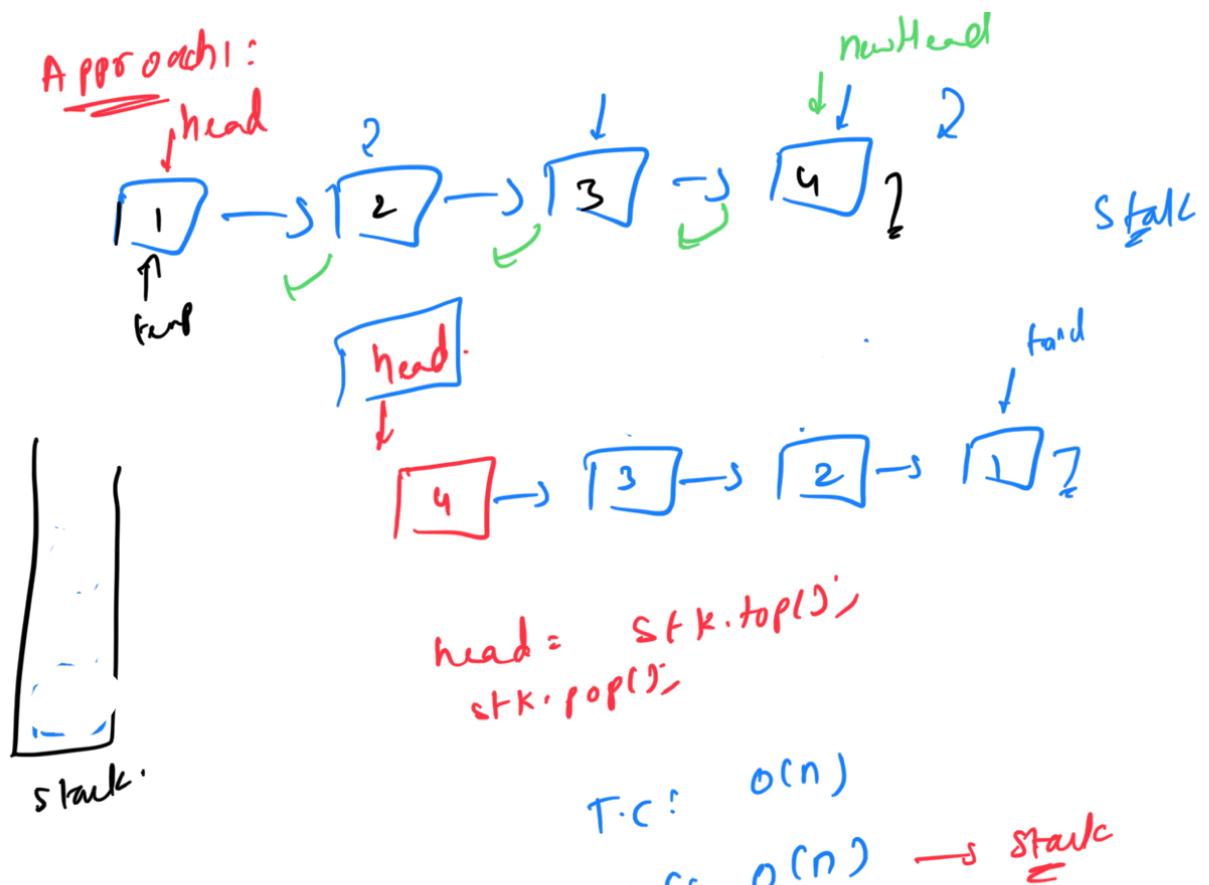
}]

null.next

Question: Reverse a Linked List.



Approach 1:



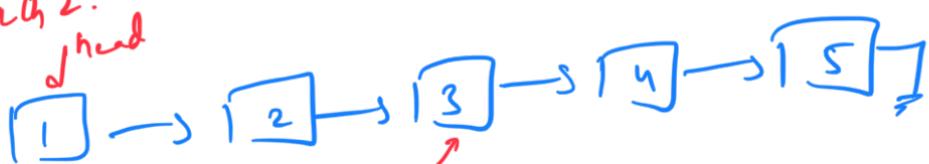
head = stk.top();
stk.pop();

T.C: $O(n)$

S.C: $O(n) \rightarrow \underline{\text{stack}}$

Approach 2:

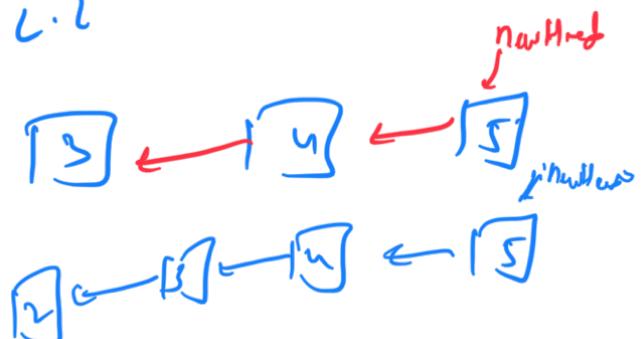
Recurse call stack.



1) Assumption:

reverse(head) reverses and returns the head of the reversed L.L

reverse(3) =
reverse(2) =



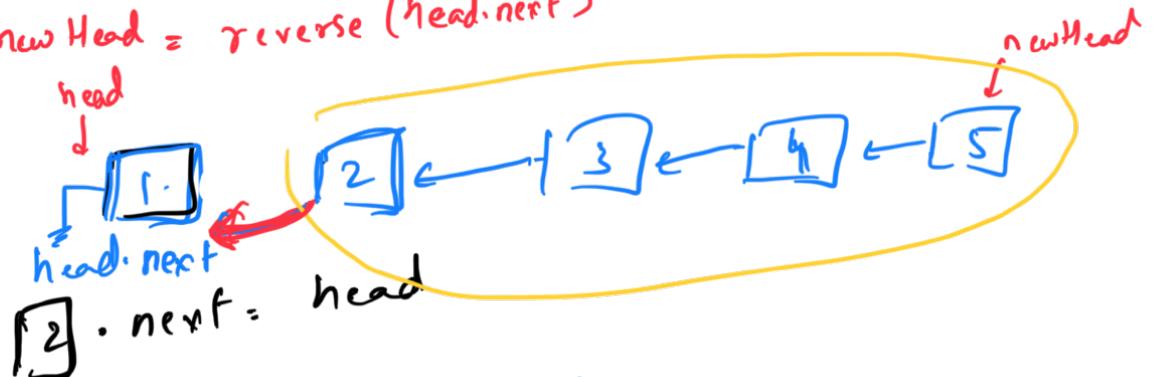
2) Main Logic / Recursive Relation:

How to break smaller ~

a bigger problem
subproblem

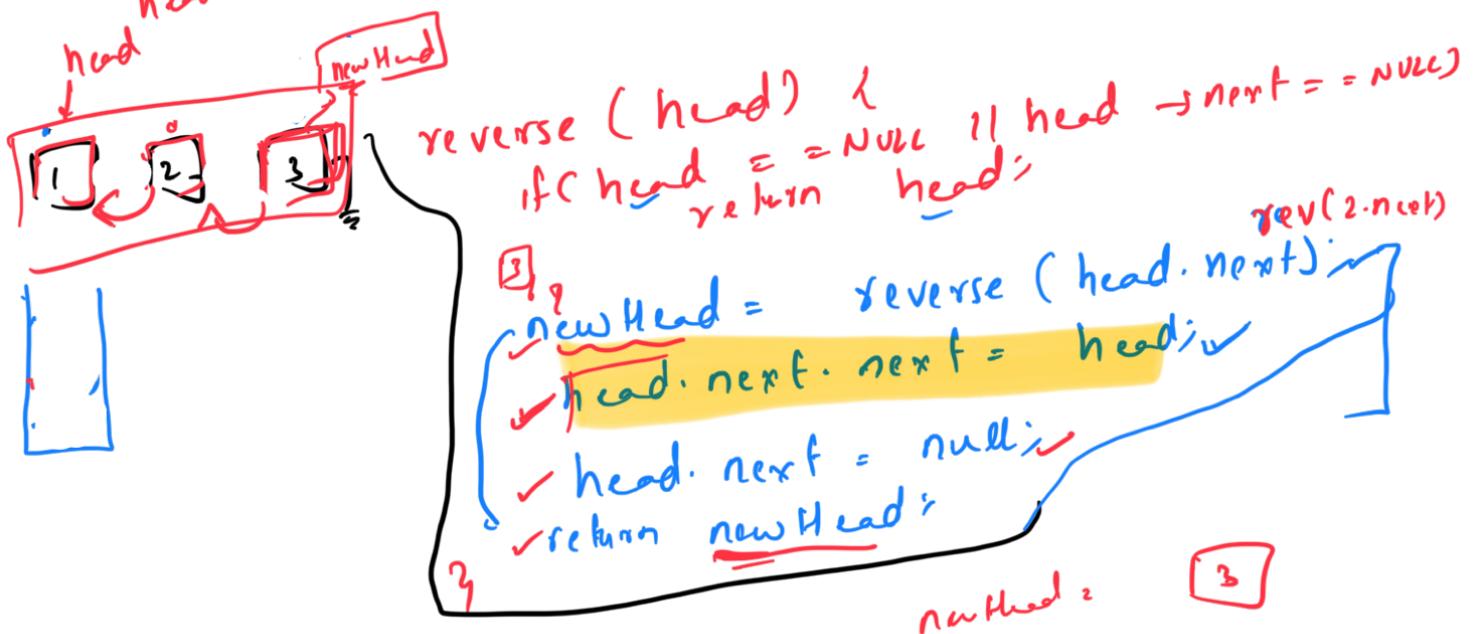


$5 \leftarrow \text{newHead} = \text{reverse}(\text{head}.next)$



$2 \Rightarrow \text{head}.next$

$\text{head} \leftarrow 2$
 $\text{head}.next \leftarrow 3$
 $(\text{head}.next).next \leftarrow \text{head}$
 $\text{head}.next \leftarrow \text{null}$

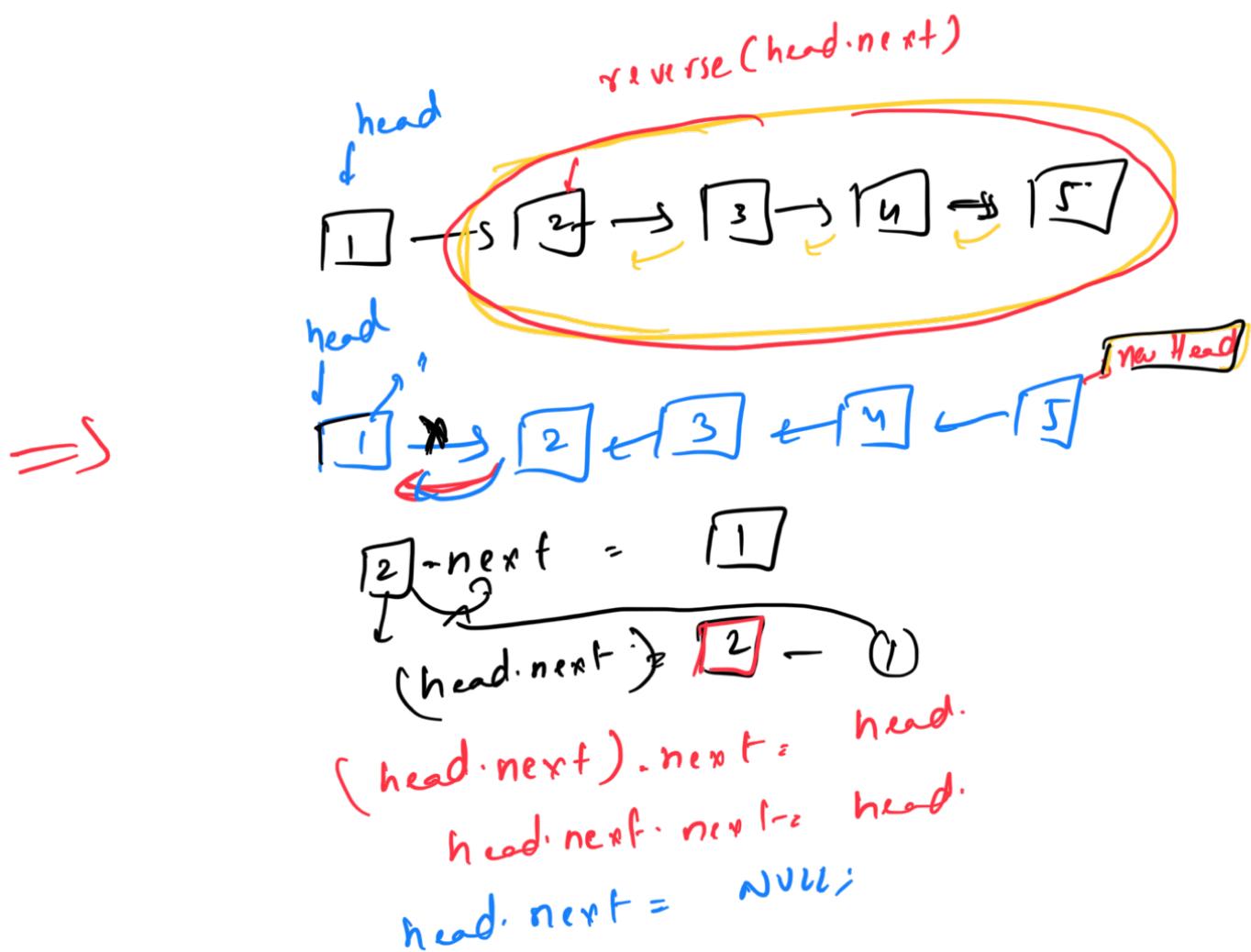


Step 3: Base Case

1) Empty LL
 $\text{if} (\text{head} == \text{null}) \text{return null};$

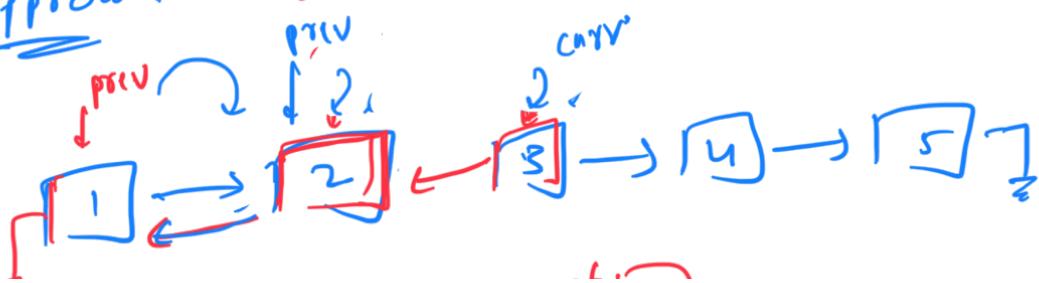


$\text{if} (\text{head}.next == \text{null}) \{$
 $\text{return head};$

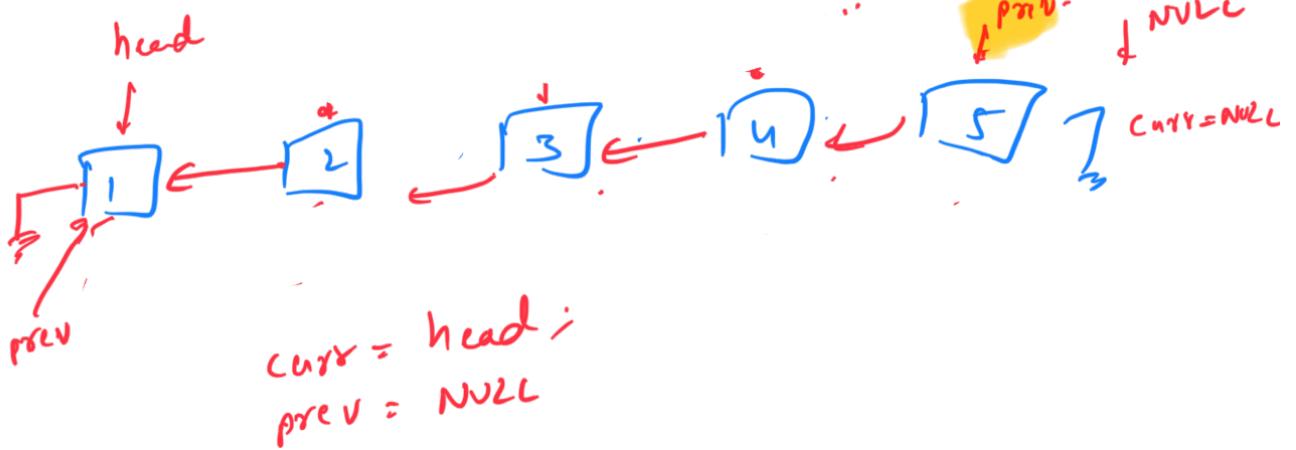


T.C: $O(n)$
 S.C: $O(n)$ [Recursion stack]

Approaches: Iteration.



\Rightarrow
 curr = null
 curr.next = prev;
 curr = next;
 prev = curr
 curr = next;



$\left[\begin{array}{l} \text{next} = \text{curr.next}; \\ \text{curr.next} = \text{prev}; \\ \text{prev} = \text{curr}; \\ \text{curr} = \text{next} \end{array} \right]$

reverse (head) {

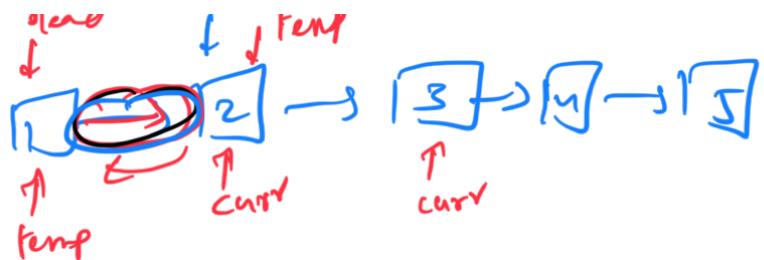
prev = NULL, curr = head;
 while (curr != NULL) {
 next = curr.next;
 curr.next = prev;
 prev = curr;
 curr = next;

return prev;

↑ **(node)** **func**

T.C: $O(n)$
S.C: $O(1)$





```
node = fun->next
temp->next = NULL;
```

Question: Print LL

```
void printLL(head) {
    temp = head;
    while(temp != null) {
        print(temp.data);
        temp = temp.next;
    }
}
```

T.C : O(n), S.C : O(1)

Question: Length of LL

```
int lengthLL(head) {
    temp = head;
    int len = 0;
    while(temp != null) {
        len++;
        temp = temp.next;
    }
    return len;
}
```

Insert at the beginning of LL

```
insertBeginning(head, data) {
    newNode = new Node(data);
    newNode.next = head;
    head = newNode;
    return head;
}
```

Insert at end of LL using tail pointer

```
head = tail = NULL;
insertAtEnd(int val){
    Node* newNode = new Node(val);
    if(head == NULL){
        head = tail = newNode;
    }
    else{
        tail->next = newNode;
        tail = newNode;
    }
}
```

Insert at end of LL

```
Node* insertEnd(Node* head, data)
{
    Node* newNode = new Node(data);
    if(head == NULL)
        head = newNode;

    Node* temp = head;
    while(temp->next != NULL)
        temp = temp->next;

    temp->next = newNode;

    return head;
}
```

Insert at kth position

```
Node* insertAtPosition(Node* head, int pos, int data){  
    int num_nodes = countNodes(head);  
    if(pos > num_nodes + 1 || pos < 1)  
        return NULL;      // Invalid Input  
  
    Node* newNode = createNode(data);  
    if(head == NULL){  
        head = newNode;  
        return head;  
    }  
  
    if(pos == 1){  
        newNode ->next = head  
        head = newNode;  
        return head;  
    }  
  
    curr = head;  
    int count = 1;  
    while(curr && count != pos-1){  
        curr = curr->next;  
        count++;  
    }  
    newNode->next = curr->next;  
    curr->next = newNode;  
  
    return head;  
}
```

Reverse a LL using stacks

```
reverseUsingStacks(Node* head) {  
    stack<Node*> s;  
    Node* temp = head;  
    while(temp!=NULL) {  
        s.push(temp);  
        temp=temp->next;  
    }  
  
    head = s.top();  
    Node* tail = s.top();  
    s.pop();  
  
    while(!s.empty()) {  
        tail->next = s.top();  
        tail = s.top();  
        s.pop();  
    }  
    tail->next = NULL;  
    return head;  
}
```

Reverse a LL using recursion

```
Node* reverse(Node* head) {  
    if(head == NULL) return head;  
    if(head->next == NULL) return head;  
  
    Node* nH = reverse(head->next);  
    head->next->next = head;  
    head->next = NULL;  
    return nH;  
}
```

Reverse a LL using iterative method

```
Node* reverse(Node* head) {
    if(head == NULL || head->next == NULL)
        return head;

    Node* curr = head, prev = NULL, next = NULL;
    while(curr != NULL) {
        next = curr->next;
        curr->next = prev;
        prev = curr;
        curr = next;
    }
    head = prev;
    return head;
}
```

Middle Node

```
Node* middleNode(Node* head) {
    Node* slow = head, *fast = head;
    while(fast != NULL && fast->next !=NULL) {
        slow = slow->next;
        fast = fast->next->next;
    }
    return slow;
}
```