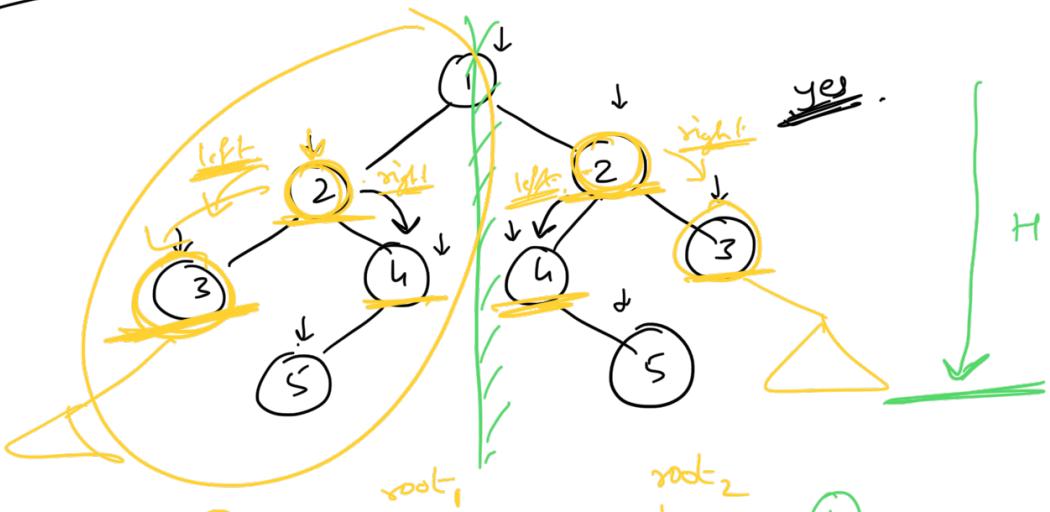


Q.1 Given a B.T., check if its symmetric.

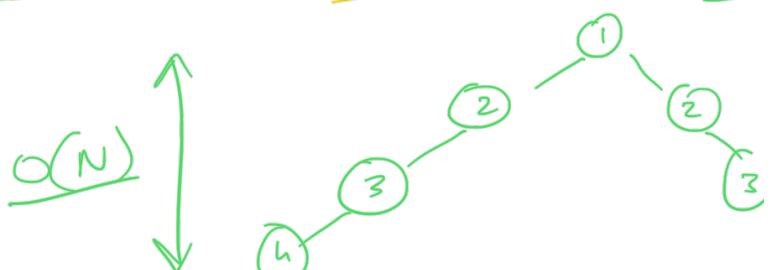


```

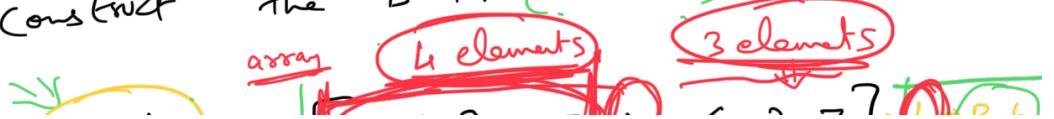
bool check (root1, root2) {
    // base case
    if (root1 == root2 == NULL) return true
    else if (root1 == NULL || root2 == NULL) return false
    if (root1.val != root2.val) return false
    return check (root1.left, root2.right)
        && check (root1.right, root2.left)
}
  
```

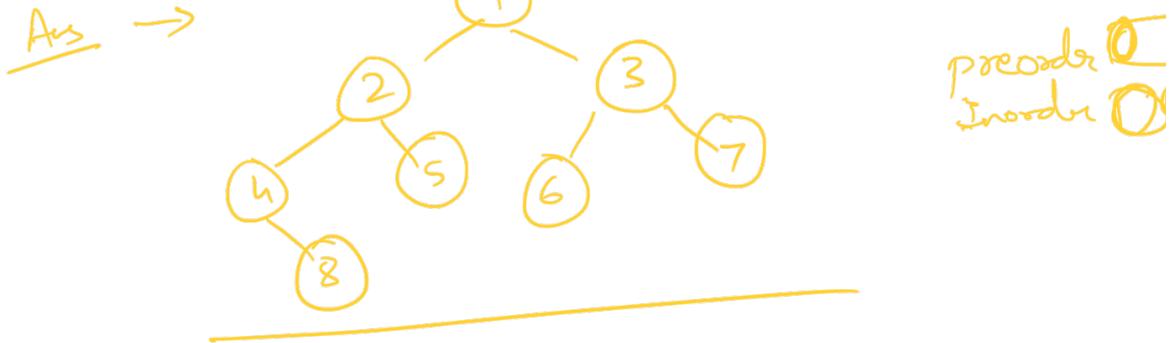
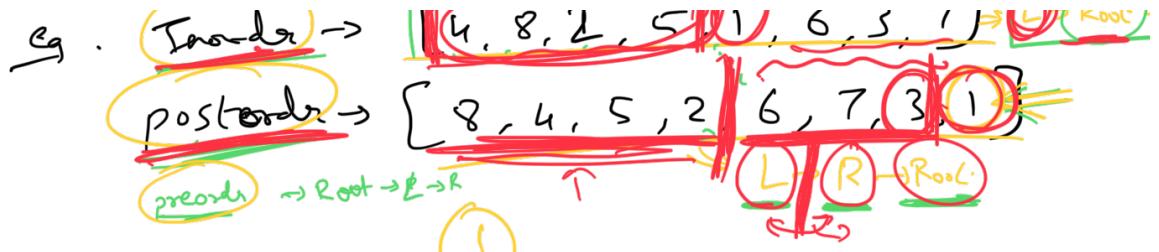
T.C $\Rightarrow O(N)$

S.C $\Rightarrow O(H)$,
worst $O(n)$



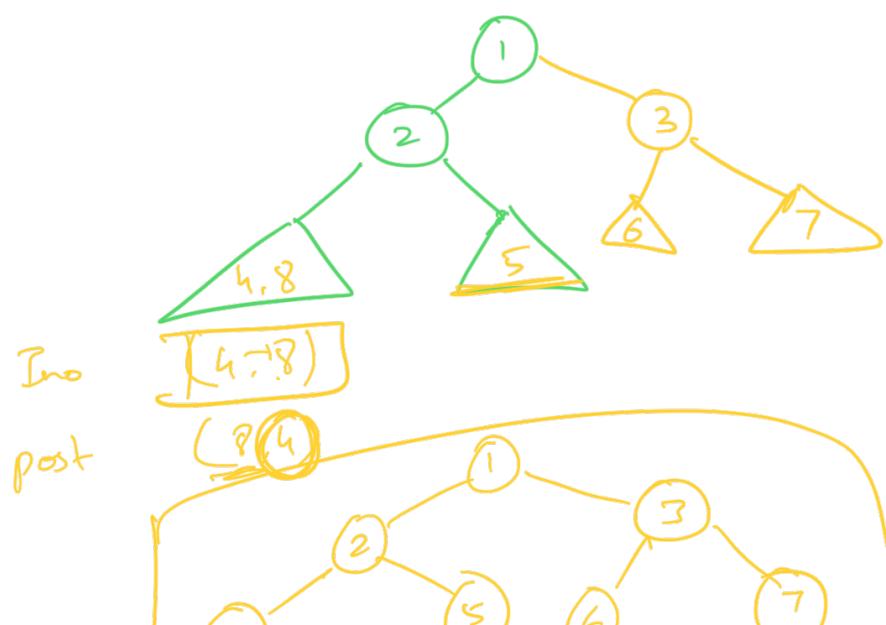
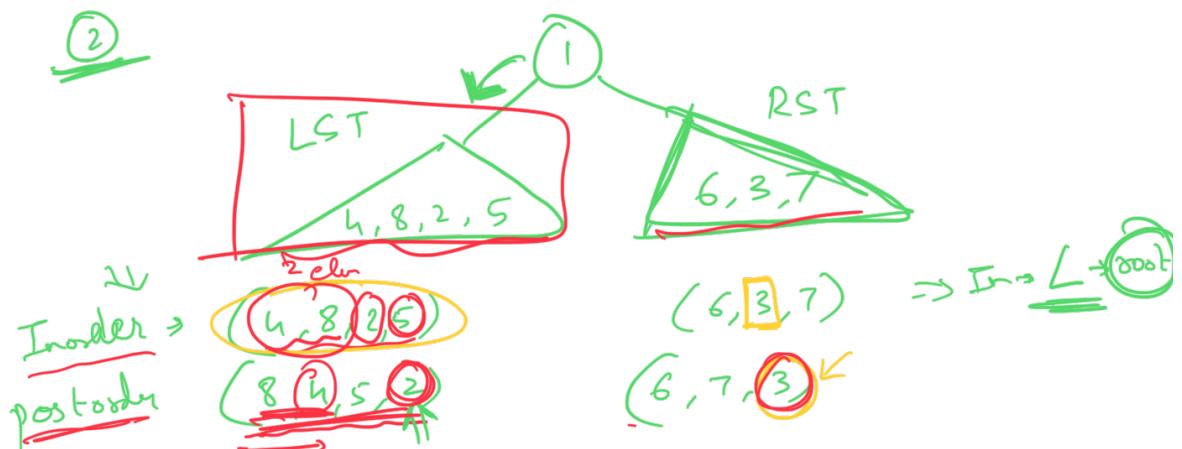
Q.2 Given postorder & inorder traversal, construct the B.T. (distinct)





Observation:

① Last node in postorder \Rightarrow Root





Q.3 Given postorder & preorder, construct B.T.



\rightarrow pre $\rightarrow 1, 2, 4, 5, 3, 6$

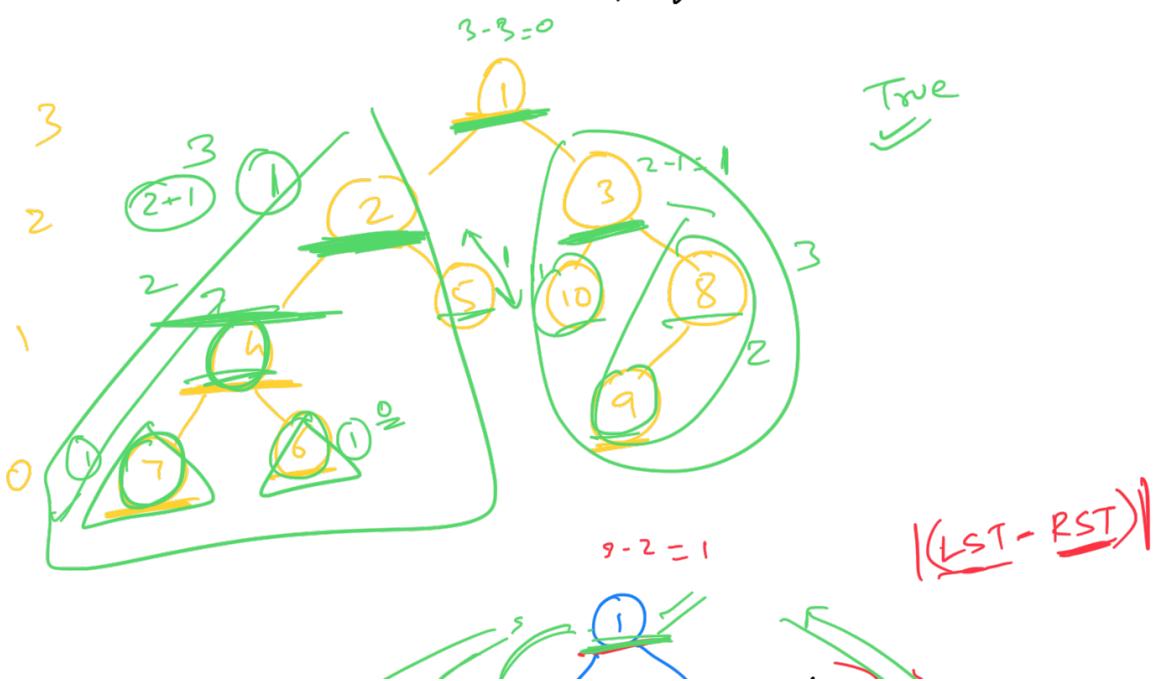
\rightarrow post $\rightarrow 4, 5, 2, 6, 3, 1$

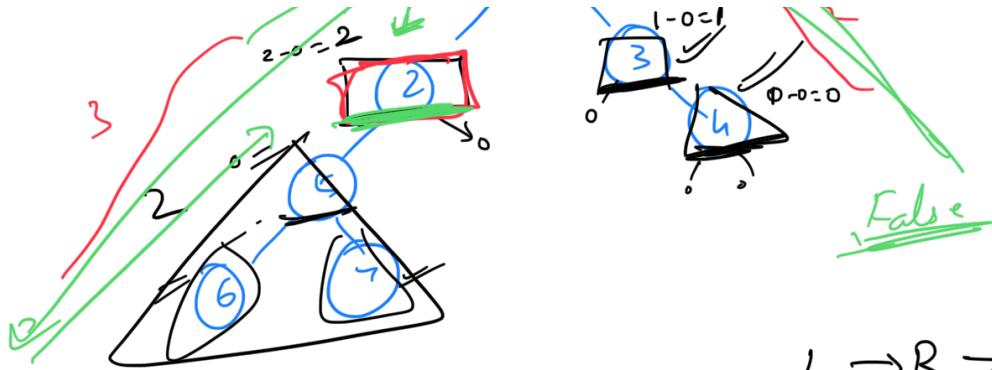
There is no unique B.T. You can constr

Q.4 Given B.T., check if tree is

height balanced.

For all nodes, $\frac{\text{height of LST} - \text{height of RST}}{\text{diff}} \leq 1$





$L \rightarrow R \rightarrow \text{root}$

Brute: For every node,
get height of LST
get height of RST
 \Rightarrow check $|h_{LST} - h_{RST}| \leq 1$ (postorder!)

bool isBalance(root) T.C

{ //Base
 if ($\text{root} == \text{null}$) return True.

$H_{LST} = \text{getHeight}(\text{root.left}) \Rightarrow O(H)$

$H_{RST} = \text{getHeight}(\text{root.right}) \Rightarrow O(H)$

return $(|H_{LST} - H_{RST}| \leq 1) \& \& (\text{isBalance}(\text{root.left})) \& \& \text{isBalance}(\text{root.right}))$

}

T.C $\Rightarrow \underline{O(N * H)}$





$$LST' = \max(\underline{LST}, \underline{RST}) + 1$$



```

pair<int, bool>
Info
{
    // Base
    if (root == NULL) return info(0, true);

    check(root)
}

```

Info LST = check(root.left)

Info RST = check(root.right)

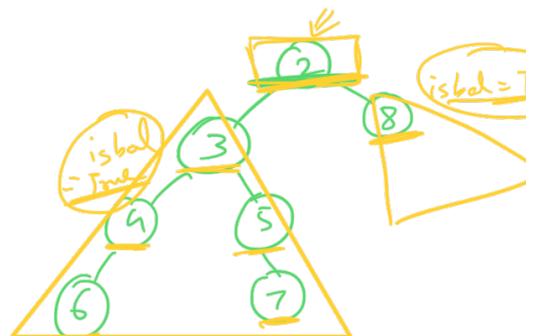
if (left.isbal && right.isbal &&
 $|left.height - right.height| \leq 1$)

return Info(max(left.height, right.height) + 1) true

else return Info(0, false)

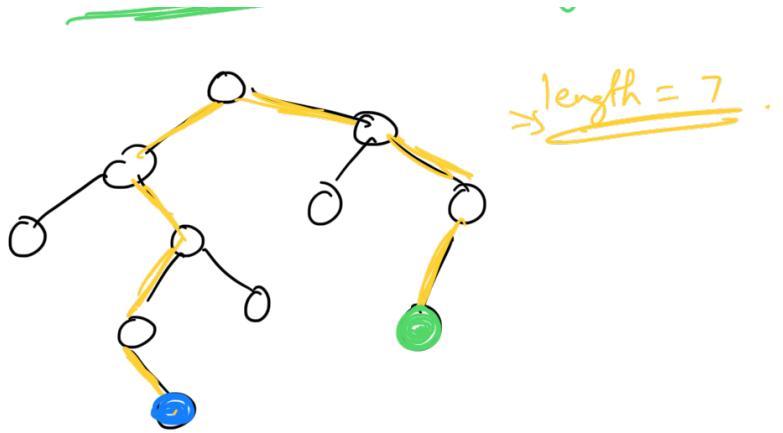
3

T.C $\Rightarrow O(N)$
S.C $\Rightarrow O(H)$

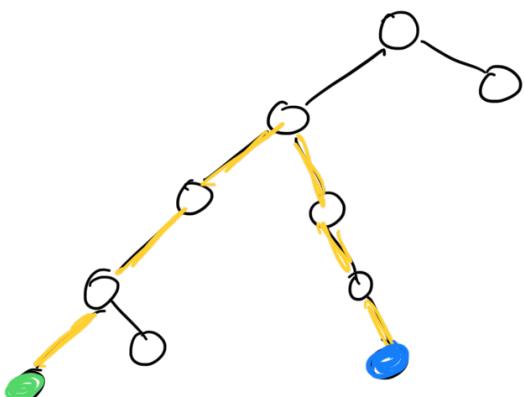


Q.5: Find diameter of the tree

longest path b/w any two nodes.

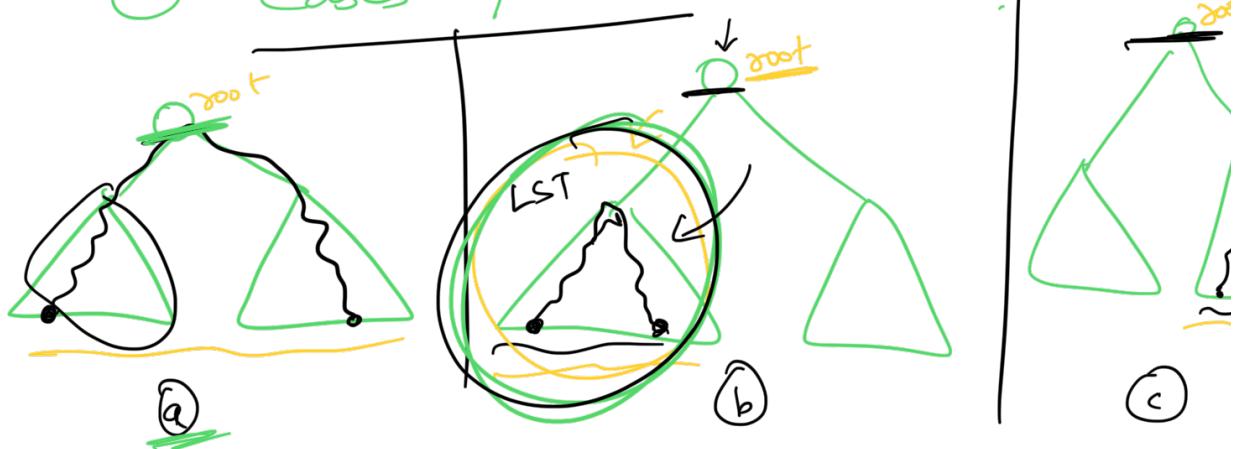


Obs: ~~(left height + right height)~~ wrong
~~max(LST) + max(RST)~~



① 2 nodes are always going to be leaf no.

② Cases possible $\text{getDiameter}(\text{root})$



$$\text{③ } \max \left(\frac{\max(\text{LST}) + \max(\text{RST}),}{\text{diam}(\text{root} \cdot \text{left})} \right)$$

(b) \rightarrow ~~get diam (root)~~
 (c) \rightarrow ~~get Diamet (root · right)~~

int get Diameter (root)

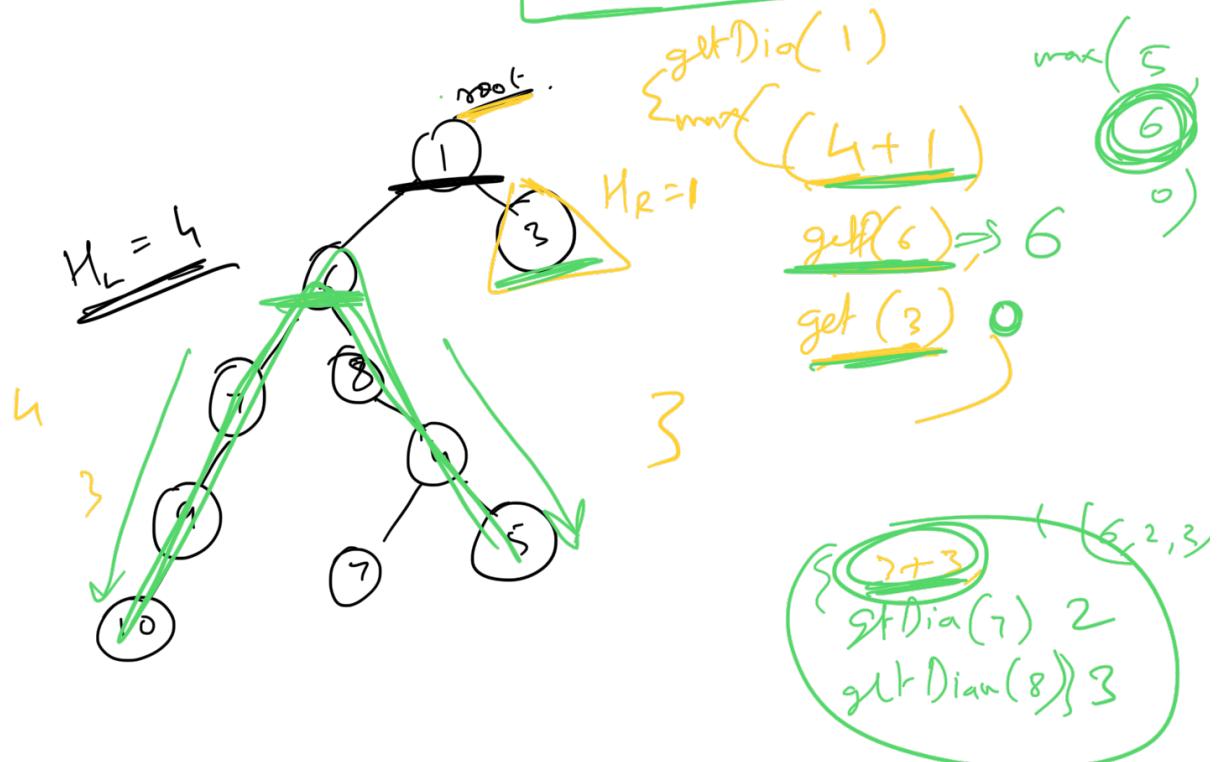
{ // Base if (root == null) return 0;

~~HL = get height (root · left)~~

~~HR = get height (root · right)~~

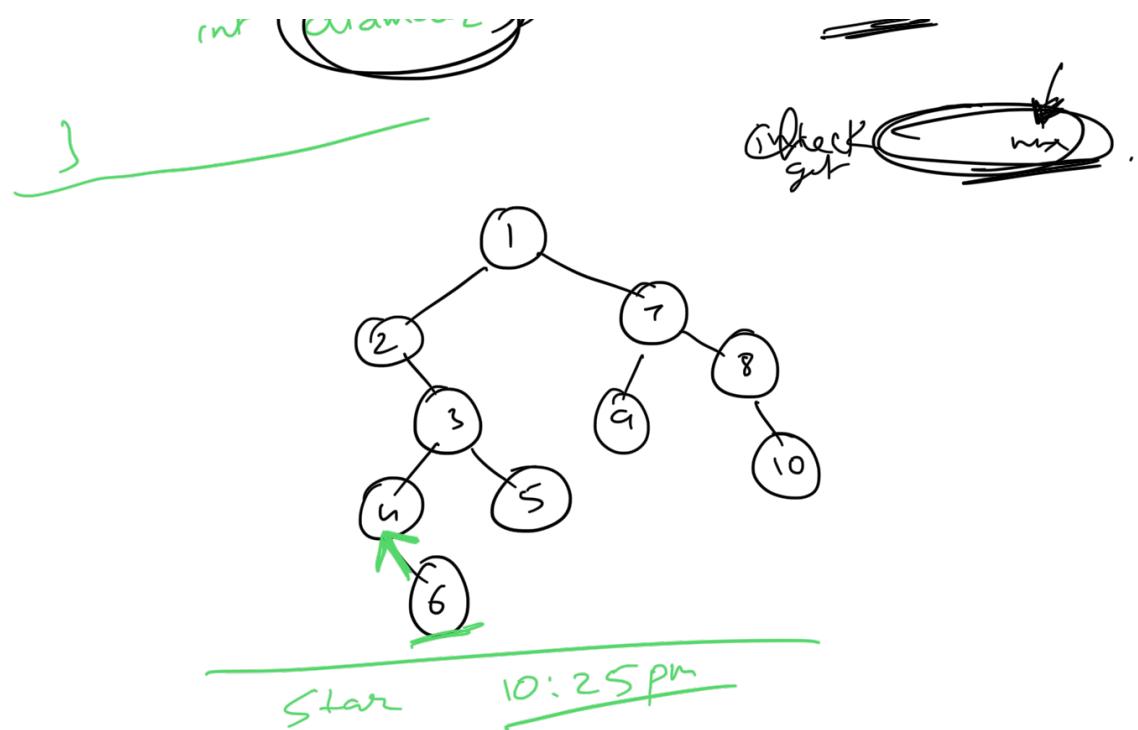
~~return max (HL + HR), getDiam (x · left), getDiam (x · right)~~

T.C $\Rightarrow O(N * H)$

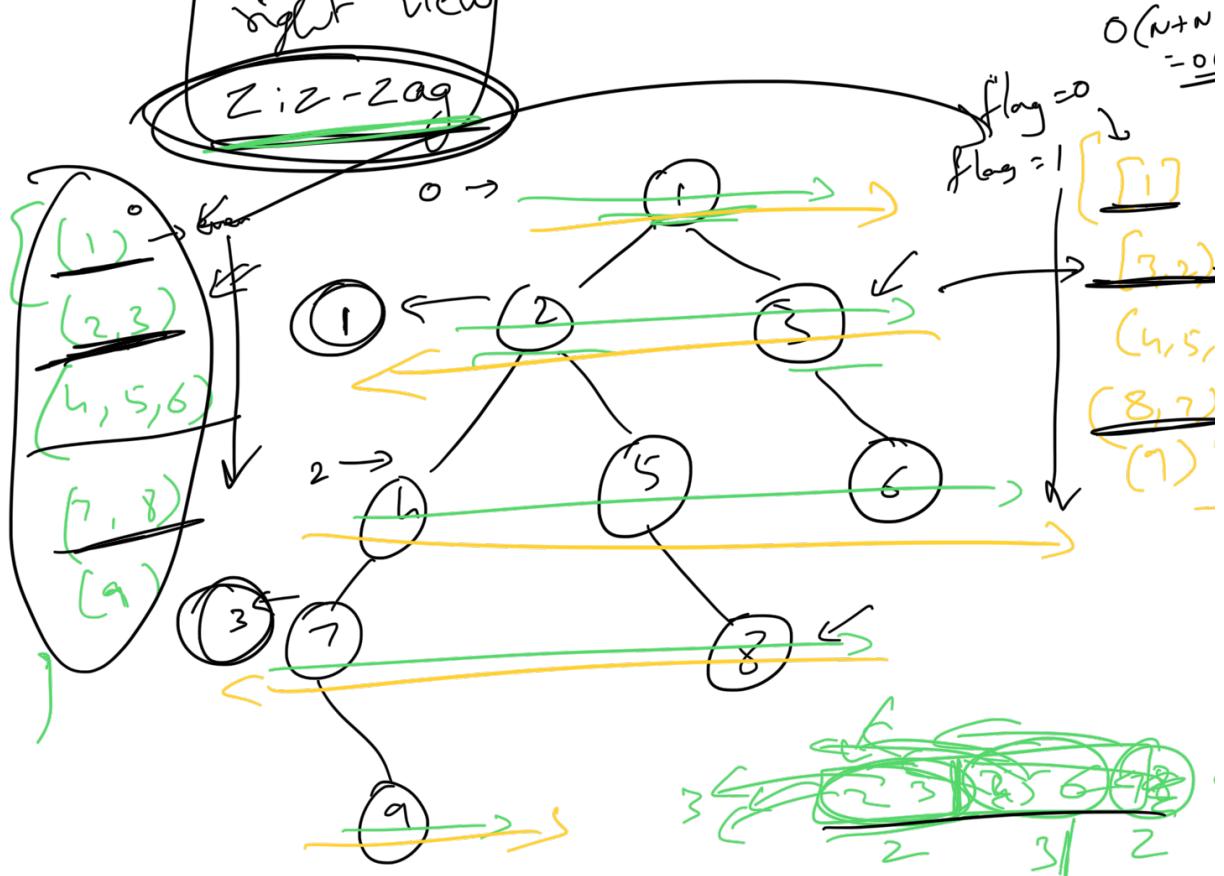


stack Info
 int height, O(.)
 - L: diameter

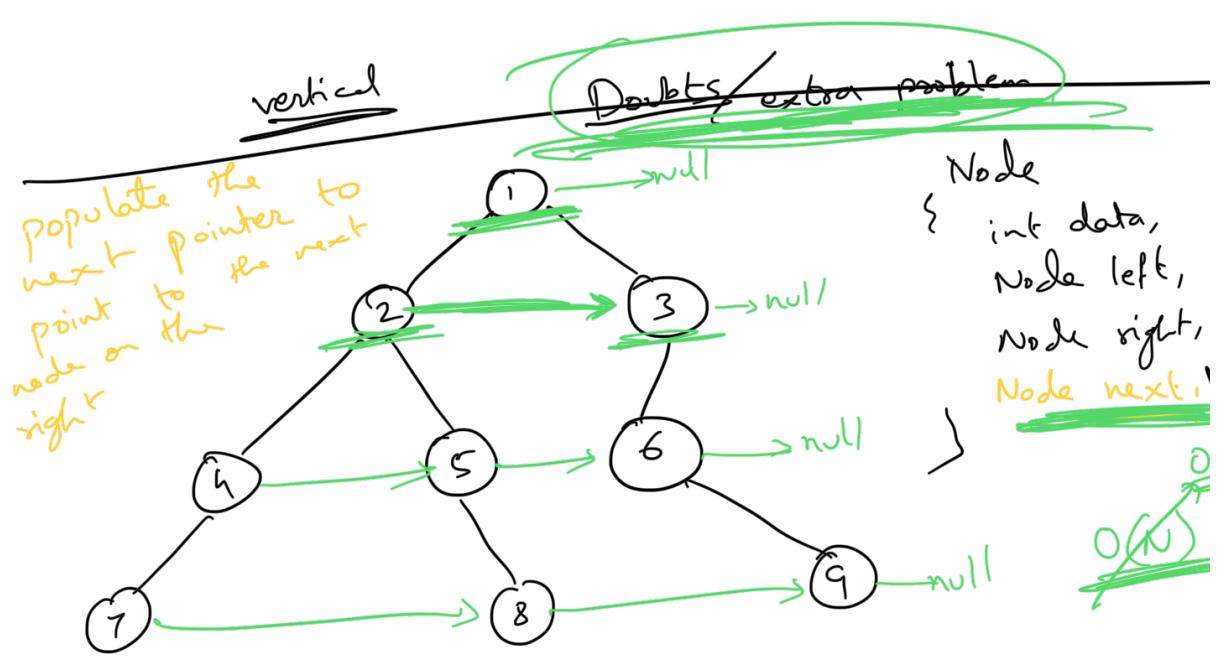
int
 2 int



Q.6 Level order traversal



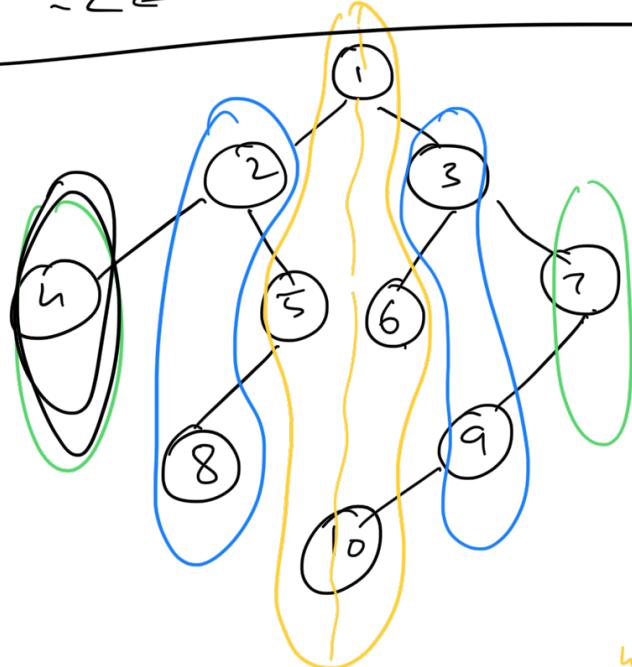
✓ +



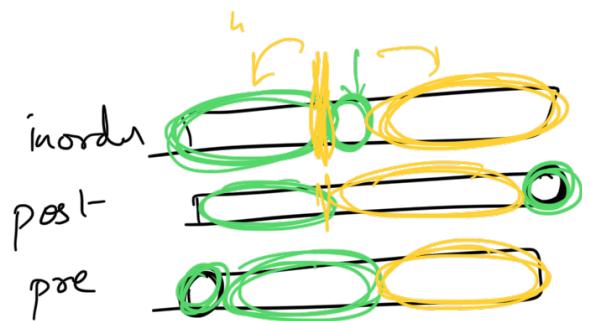
Method 1: queue

Method 2 → O(1) space

-2 ← -1 ← 0 → +1 → 2



(1)
(2, 8)
(1, 5, 6, 10)
(3, 9)
(7) \



— 1 —