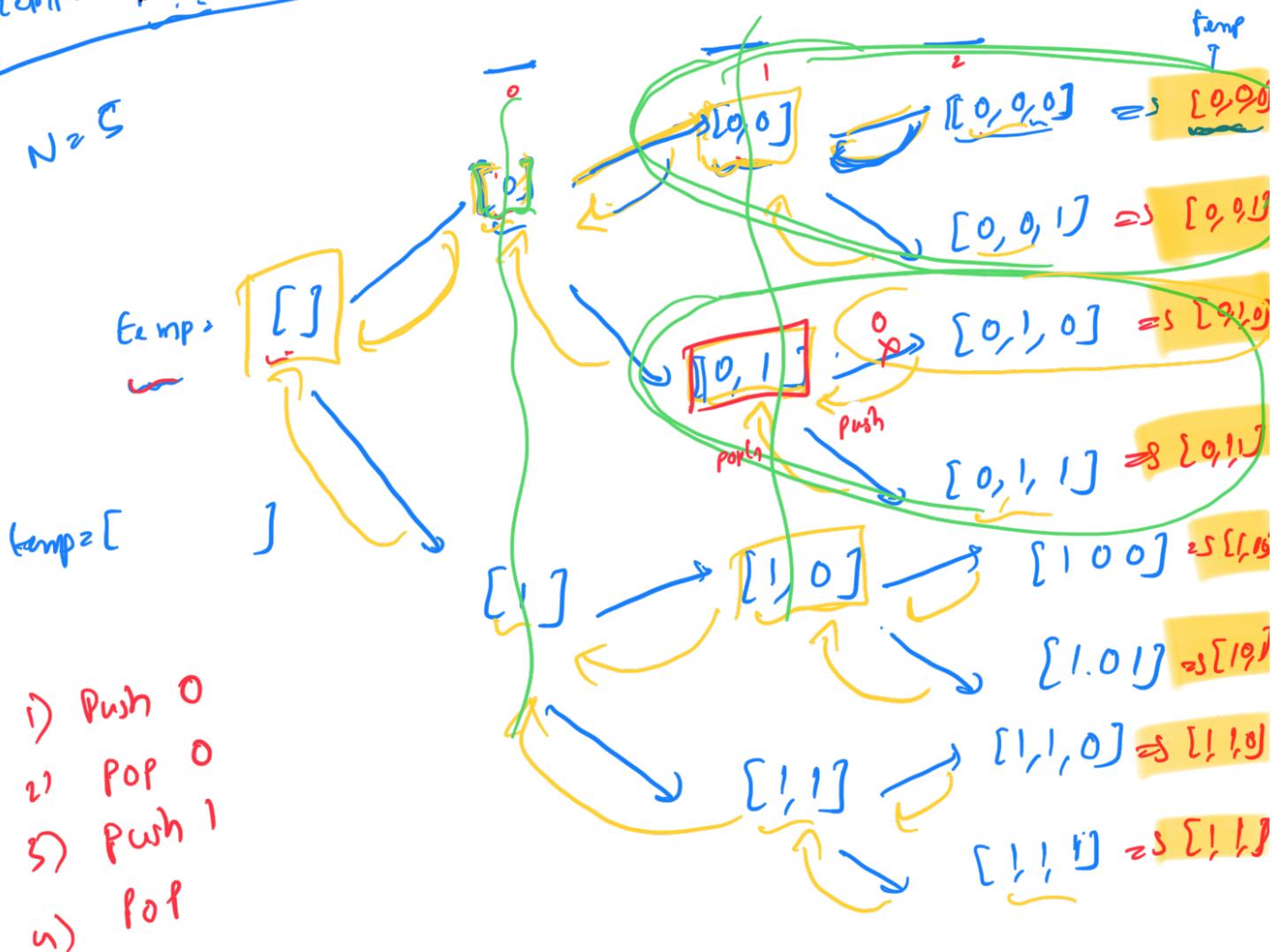
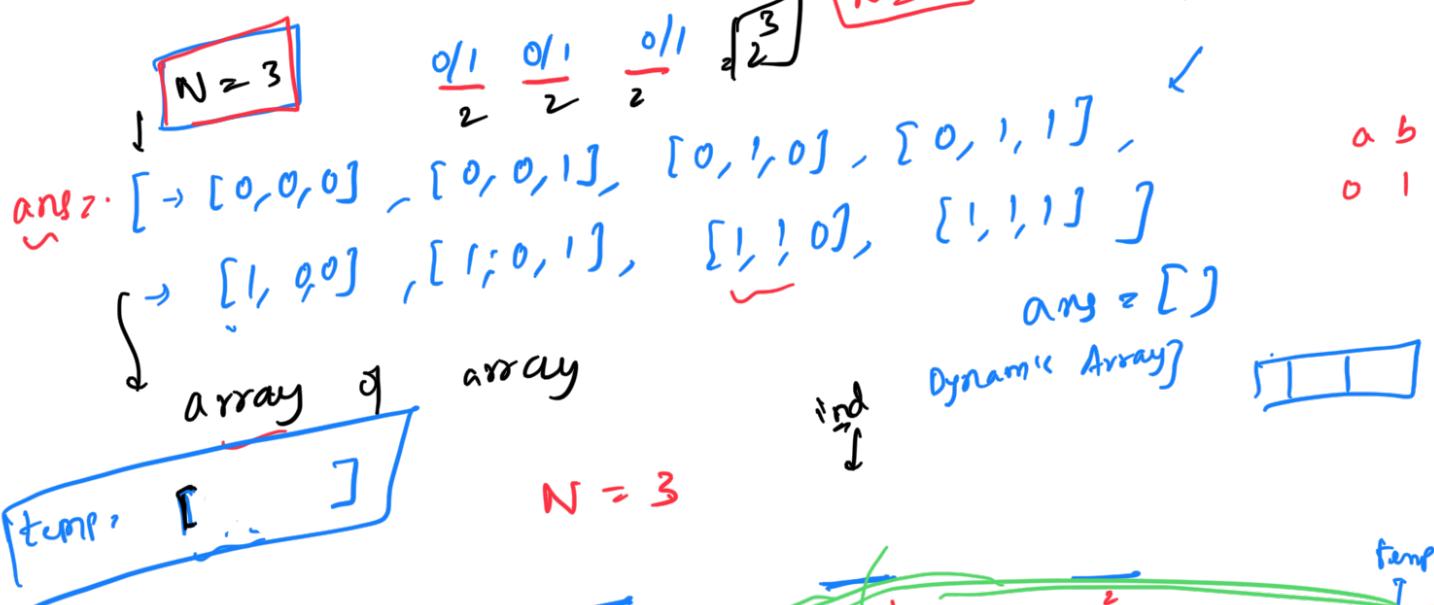


Backtracking - I

Question:

Generate all arrays of size N with 0's / 1's in lexicographical order [Dictionary order]



For i^{th} element :

- 1) Select 0
- 2) Select 1

\rightarrow

```
temp.push(0);
F(ind+1);
temp.pop();
temp.push(1);
E(ind+1);
temp.pop();
```

Parameters :

1) ind (Global variable)

2) temp : size of arrays

3) N : ans \Rightarrow global

temp \Rightarrow global vector
 void allArrays(ind, N) {
 if (ind == N) {
 ans.push(temp);
 return;
 }
 }

// Select 0 // OO

```
temp.push(0);
allArrays(ind+1, N);
temp.pop();
```

// UNDO

// Select 1 // UNDO

```
temp.push(1);
allArrays(ind+1, N);
temp.pop();
```

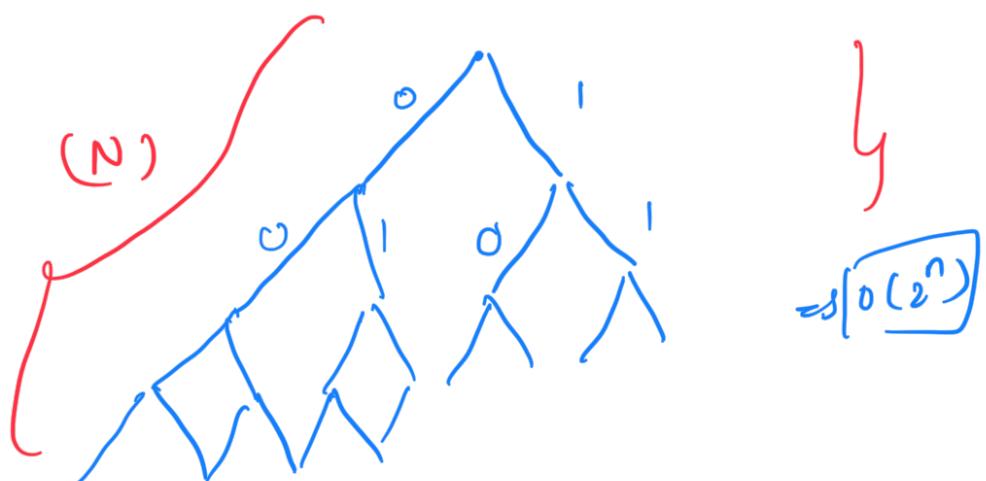
T.C: $O(2^n)$

S.C: $O(N)$

$$O(N) + O(N) = O(N)$$

Recursion stack temp

Depth of tree: (N)



Question:

Given an array in print all subsets in order

size N ,
lexicographical
 $2 \times 2 \times 2 \dots n = 2^n$

$\boxed{[7, 3, 6]}$

$A =$

subsets =

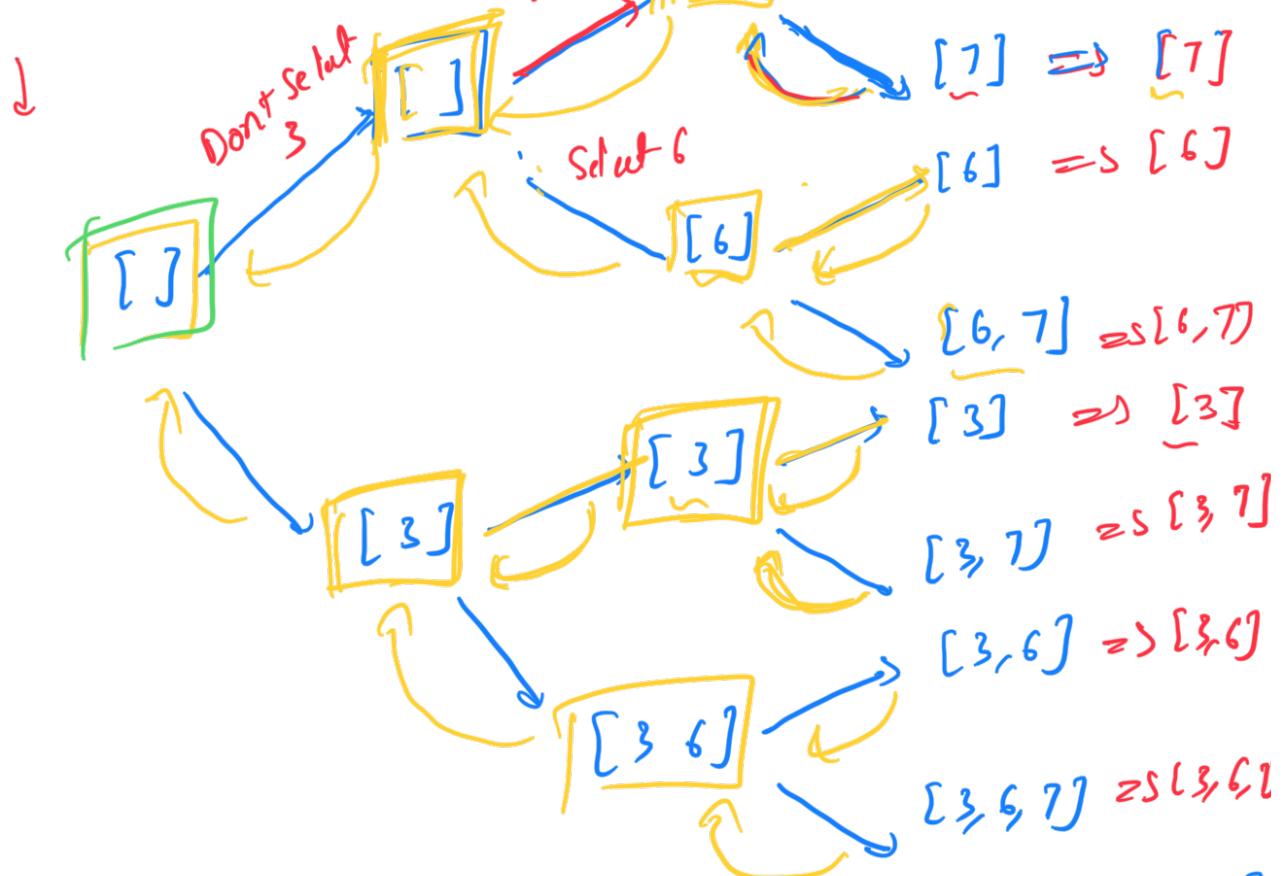
$[7], [3], [6], [7, 3], [7, 6], [3, 6]$
 $[7, 3, 6], [] \Rightarrow 8 = \boxed{2^3}$

Output = $\boxed{[]} \quad \boxed{\underbrace{[5]}_{[6]}} \quad \boxed{\underbrace{[3, 6]}_{[6, 7]}} \quad \boxed{\underbrace{[3, 6, 7]}_{[7]}} \quad \boxed{\underbrace{[3, 7]}_{[]}}$

$\boxed{\frac{n}{2}}$

$ans = []$

$A = [3, 6, 7]$
 $i = 0$
 $\text{temp} = []$
 $i = 1$
 $i = 2$
 $i = 3$
 $i = 4$
 $i = 5$
 $i = 6$
 $i = 7$



Ans → $\{ [], [1], [6], [5, 7], [3], [3, 7], [3, 6], [3, 6, 7] \}$

\downarrow

sort(ans);

Parameters:

- 1) ind
- 2) N
- 3) temp

```

> allSubsets(ind, N, temp) {
    if(ind == N) {
        ans.push_back(temp);
        return;
    }
    // Don't include A[ind]
    allSubsets(ind+1, N, temp);
    // Include A[ind]
    temp.push_back(A[ind]);
    allSubsets(ind+1, N, temp);
    temp.pop_back();
}
    returns;
}
    
```

ans = hashset
 ↓
 output of {

main() {

// Sort the given

sort(A);

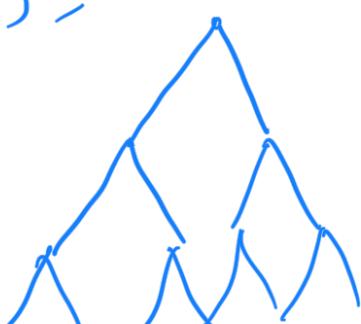
allSubsets(0, A.size() - 1);

sort(ans);

ans = hashset();

f.c.: $O(2^n)$

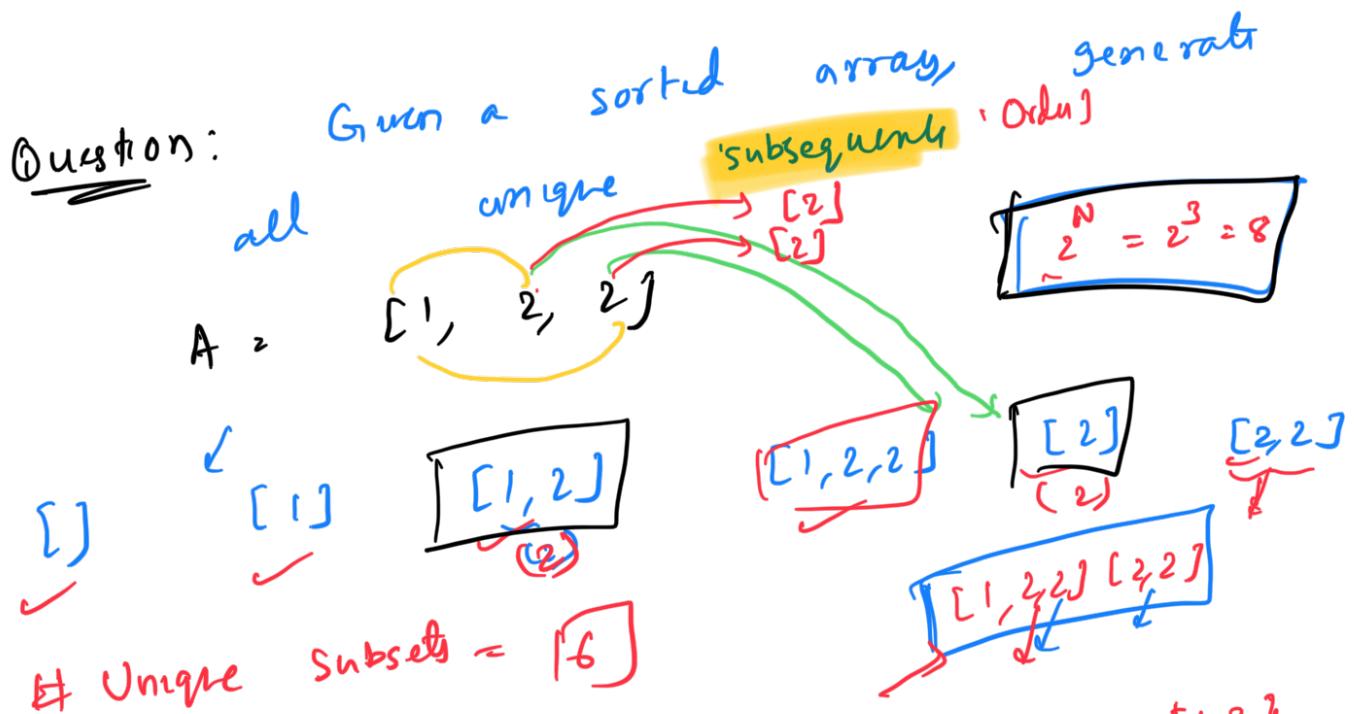
s: $O(n)$



1th position

$A = [3, 4, 5, 7, 7, 8] \nearrow \nearrow \nearrow$

$$\begin{aligned} a &= n! \\ n! &= n + k - 1 + (n-1)! \end{aligned}$$



Solution1:

Hashset

\times

$$\begin{aligned} \{1, 2\} \\ \{2, 1\} \\ \{1, 2\} = \{2, 1\} \end{aligned}$$

Solution2:

→ Can we do without using hashset

→ Don't even generate these duplicates

$A = [1, 1, 2, 3]$

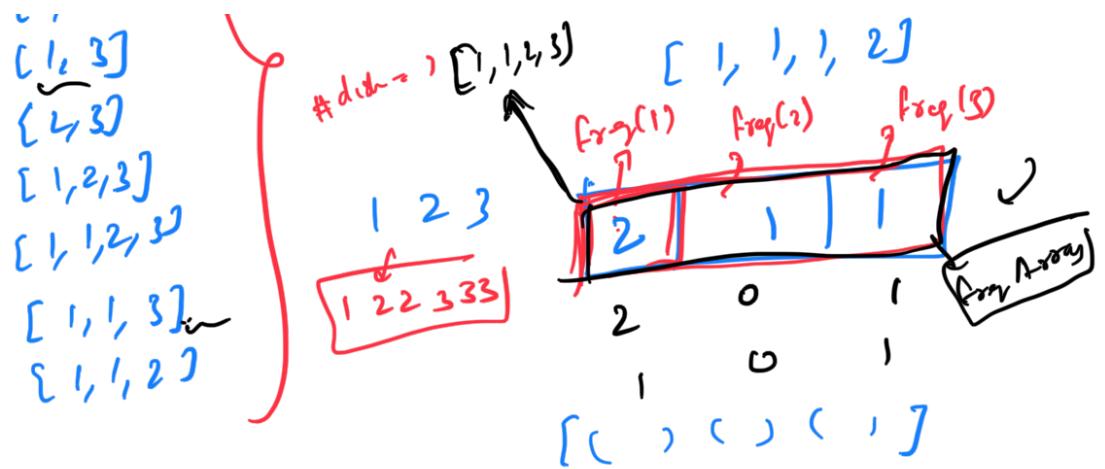
Unique Subsets = 12

$\begin{bmatrix} [] \\ [1] \\ [2] \\ [3] \\ [1, 2] \end{bmatrix}$

→ All these 12 subsequences have to be in sorted order

Distinct: $\begin{bmatrix} 1, 2, 3 \\ 3, 1, 0 \end{bmatrix} \quad (3)$

$$N = 4 \\ 2^4 = 16$$



Size of Array:
No. of distinct unique elements

$$N = 6$$

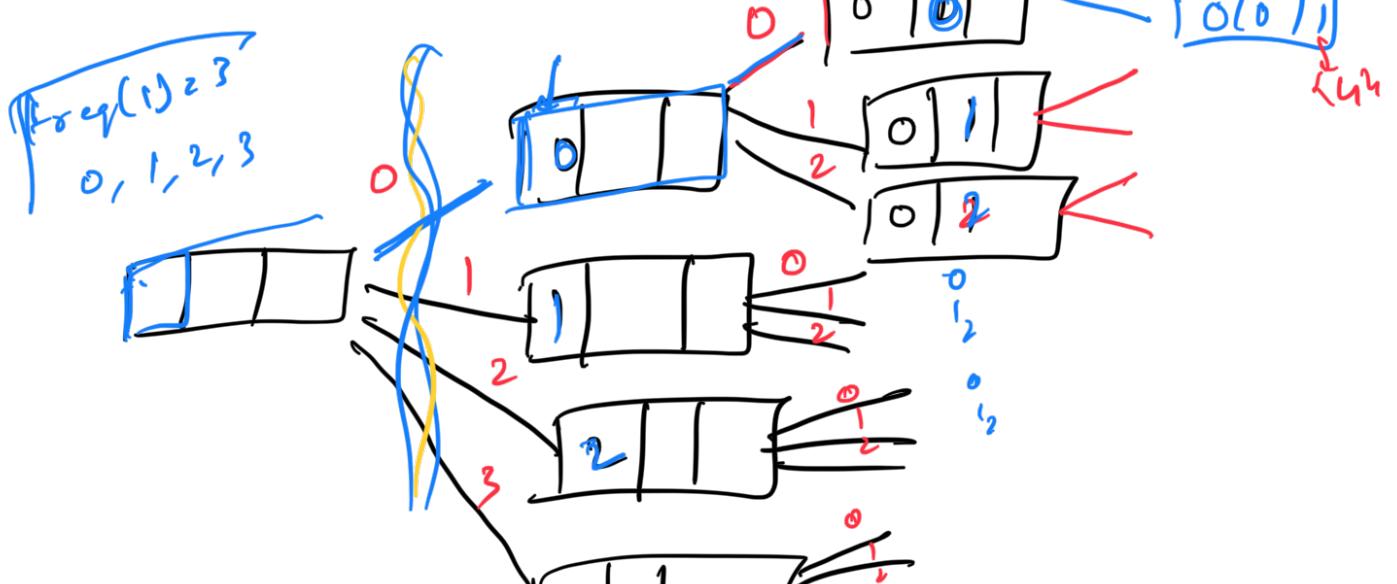
→ F.size()

Generate all possible frequency arrays;

Diagram illustrating the execution of a for loop with index $i=5$. The array A is shown with indices 0, 1, 2, 3, 4, 5. The element at index 5 is highlighted with a red box and labeled '3'. The label 'for(j=0) j<5' is written below the array.

$$\text{log}(1) = 3$$

0, 1, 2, 3



13) ↘

```

generate (ind, F, temp) {
    if (ind == F.size()) {
        // Generate the subsequence using curr
        for (j=0; j <= (F[ind].freq); j++) {
            temp.push(j);
            generate (ind+1, F, temp);
            temp.pop();
        }
    }
    return;
}

```

X HashSet

Question: Given an array, generate all
the permutations

A = [4, 7, 8]

	[4, 7, 8]	[7, 4, 8]	[7, 8, 4]	[8, 4, 7]
	[4, 8, 7]			[8, 7, 4]

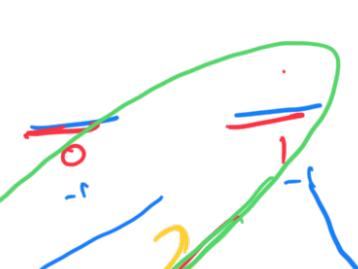
N!

$$\frac{3}{4/18} \cdot \frac{2}{7} \cdot \frac{1}{8} = \boxed{N!}$$

$N \times (N-1) \times (N-2) \dots$

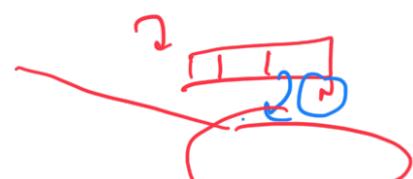
temp = [-1, -1, -1]

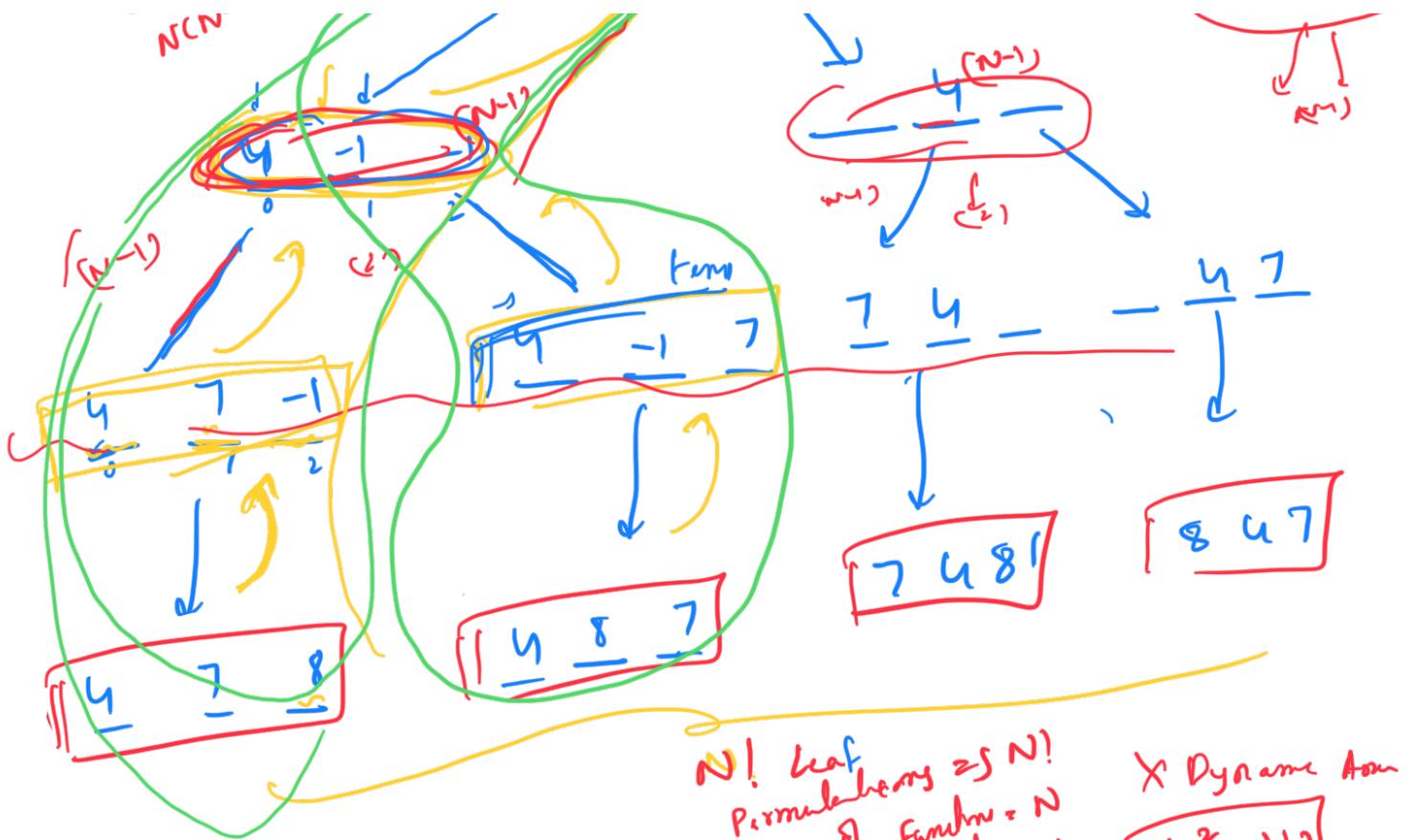
$(N + N(N-1)) \dots$



-1 2

temp = [-1, -1, -1]





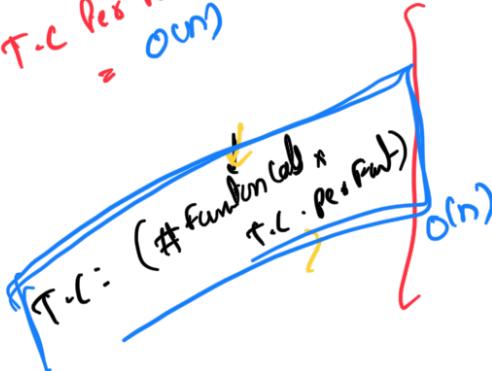
Parameters:

- 1) ind
 - 2) N \Rightarrow size of Array
 - 3) temp \rightarrow vector/Array List
- $[-1, -1, -1]$
 \downarrow
 4

$N!$ leaf
 permute leafs $\geq N!$
 No. of Function = N
 T.C. per function = N

X Dynamic Array
 $p(N^2 \times N!)$

T.C per function
 $= O(n!)$



```

void permute(ind, N, temp) {
    if (ind == N) {
        ans.push_back(temp);
        return;
    }
    for (j = 0; j < N; j++) {
        if (temp[j] == -1) {
            temp[j] = A[ind];
            permute(ind + 1, N, temp); // O(n)
            temp[j] = -1; // UNP
        }
    }
}
  
```

compute the unique permutation

Extension:

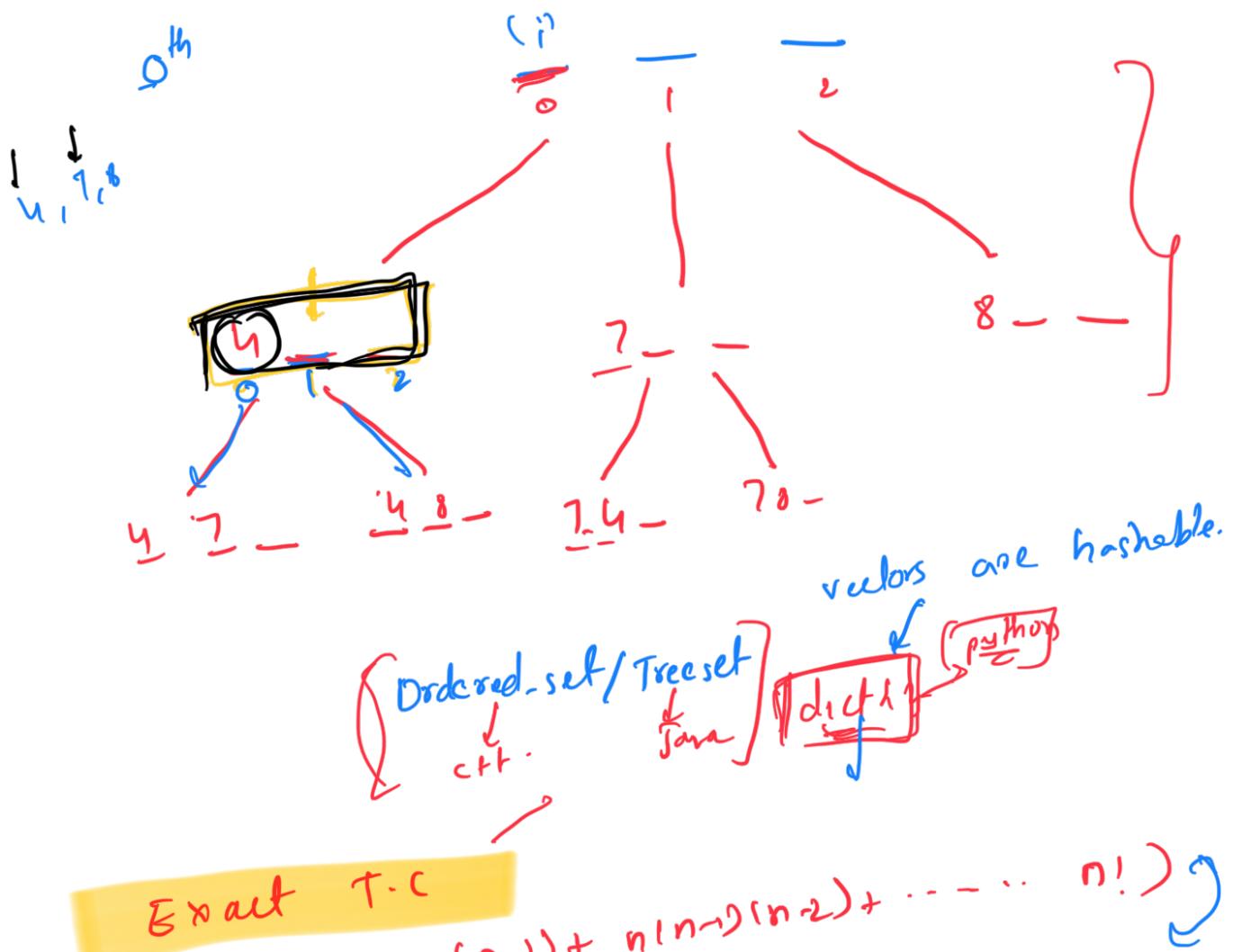
solution 1:

HASH SET

Solution 2:

UTKARSH's Approach

$$A = [4, 7, 8]$$

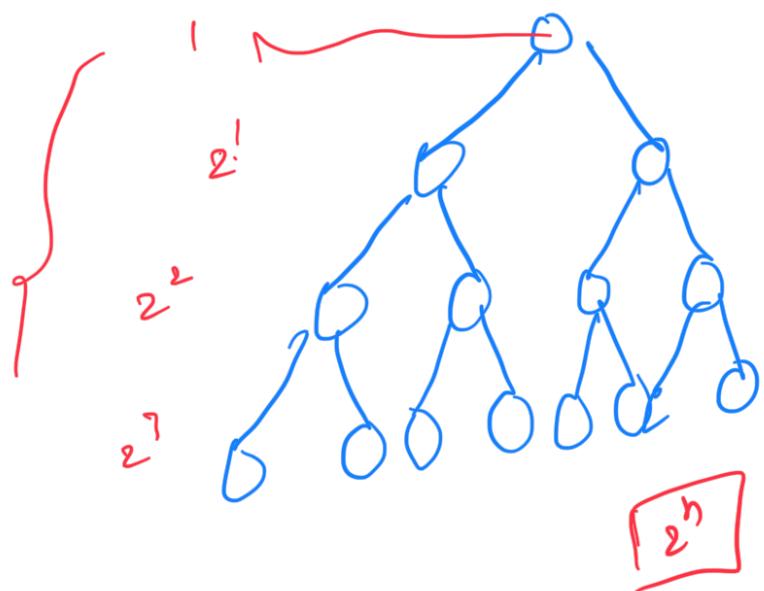


functions = $(n + n(n-1) + n(n-1)(n-2) + \dots + n!)$

computation = $O(n)$

T.C Per tree

Upper bound:



j