

Build

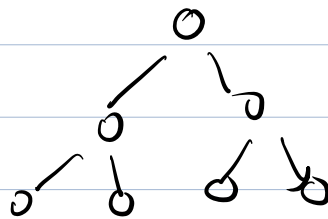
TC: $O(N)$

query (l, r)

TC: $O(\log N)$

update (ith, val)

TC: $O(\log n)$



$\log n$

$1 + 2 + 4 + \dots$

$$S_n = \frac{n(2^n - 1)}{2 - 1}$$

$$= 1 \left(\frac{2^{\log n} - 1}{2 - 1} \right)$$

void build (int idx; l, r)

if (l == r)

{ segtree[idx] = arr[l]
return

mid =

build on left child

build on right child

segtree[idx] = left + right

$$S_n = N - 1$$

$$T(N) = 2 \times T(N/2) + 1$$

$$T(N) = 2 \left[2 \times T\left(\frac{N}{4}\right) + 1 \right] + 1$$

$$T(N) = 4 T\left(\frac{N}{4}\right) + 3$$

$$T(N) = N T\left(\frac{N}{N}\right) + (N-1)$$

$$T(N) = N + N - 1$$

$$T(N) = O(N)$$

sum of intervals.

update $[l, r]$ +1

for($i=1$, $i \leq r$) $i++$)

(
 update i^{th} element = increase it by 1
 update(i^{th} element, val)
)

$$(r-l) \times \log N,$$

$$\downarrow$$

$$N \log N$$

Build

$\rightarrow O(N)$

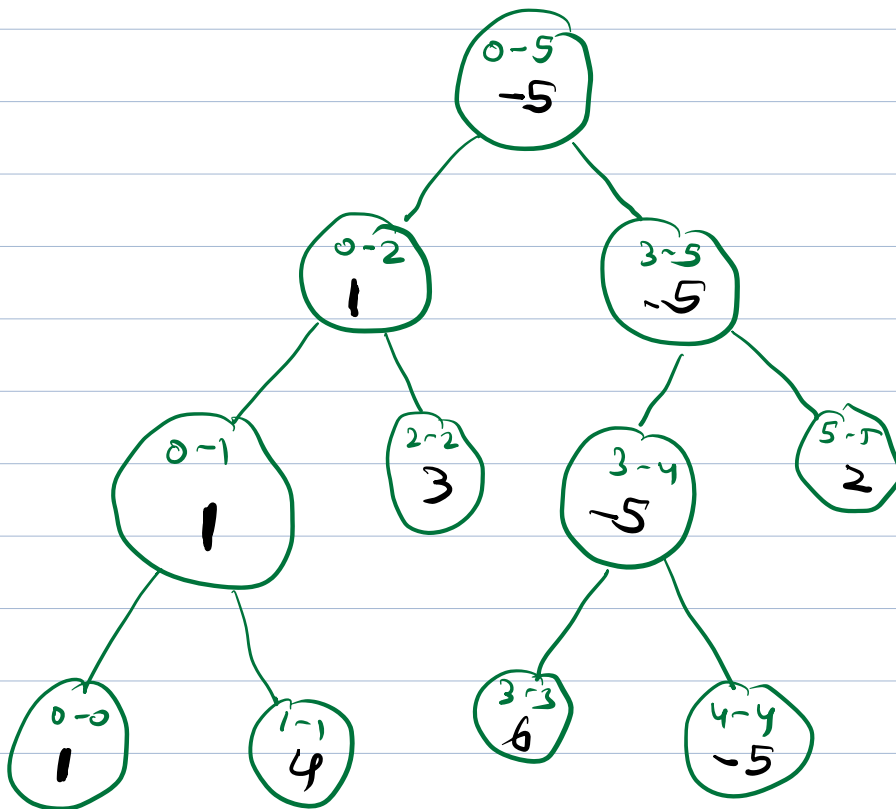
query(l, r)

→ $O(\log N)$

update(l, r, val)

→ $O(N \log N)$

$\{ \overset{0}{1}, \overset{1}{4}, \overset{2}{3}, \overset{3}{6}, \overset{4}{-5}, \overset{5}{2} \}$



Build

→ $O(N)$

query(l, r)

→ $O(\log N)$

update(l, r, val)

→ $O(N)$

0 1 2 3 4 5

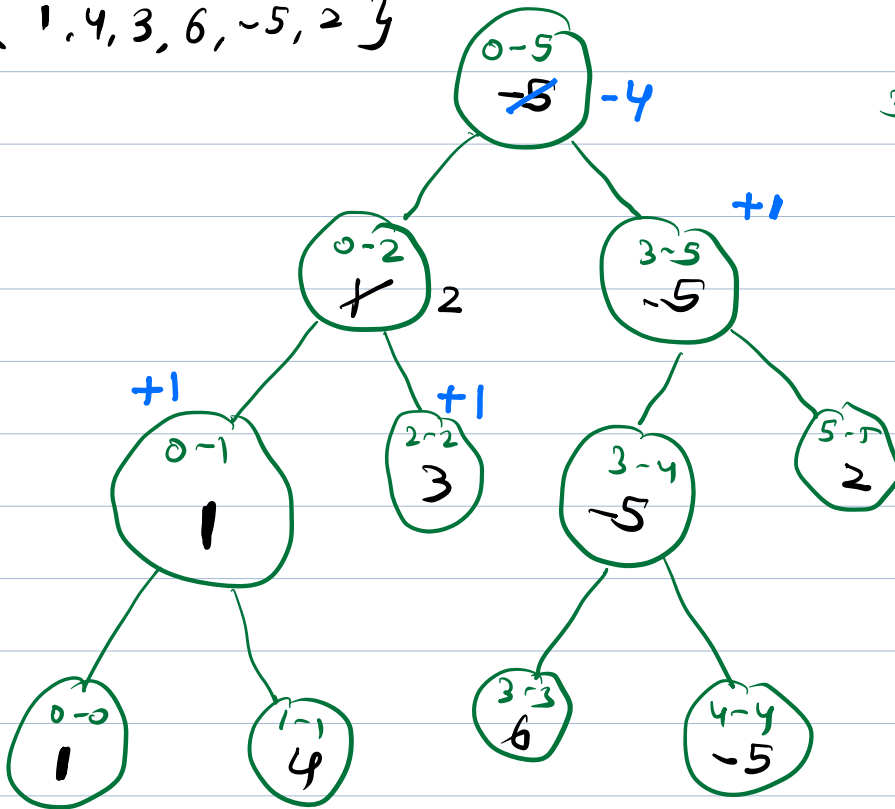
1) update(0-5) +1

↓ 1, 4, 3, 6, -5, 2

2) query (0-5) = -4

3) query (0-2)

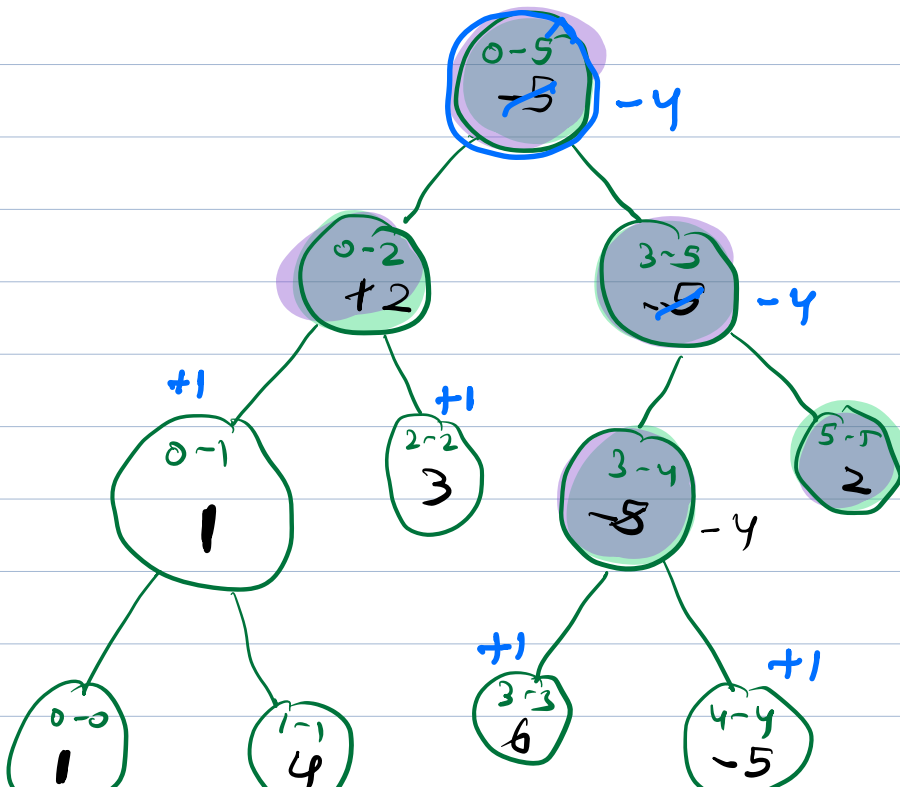
2



query (1-2)

↓
logN.

0-4



update (0-4)

↓
logN.

Build	→	$O(N)$
query(l, r)	→	$O(\log N)$
update(l, r, val)	→	$O(\log N)$

Idea is update node where complete range is overlapping. and store somewhere that its children needed to be updated later.

Lazy propagation

delaying your updates, until they are actually required.

```
int segtree[4N]
int lazy[4N] = {0}
```

Build function ? minimum element

```
void update (idx, start, end, l, r, val)
```

```
if ( lazy[idx] != 0 )
```

```
    segtree[idx] += lazy[idx]
```

```
    if ( start != end )
```

```
    {  
        lazy[2*idx+1] += lazy[idx]  
        lazy[2*idx+2] += lazy[idx]  
    }
```

```
    lazy[idx] = 0
```

```
if ( l > end || r < start ) return
```

```
if ( start ≥ l && end ≤ r )
```

```
    segtree[idx] += val
```

```
    if ( start != end )
```

```
    {  
        lazy[2*idx+1] += val
```

```
        lazy[2*idx+2] += val
```

```
int mid =  $\frac{start + end}{2}$ 
```

```
update ( 2*idx+1, start, mid, l, r, val )
```

```
update ( 2*idx+2, mid+1, end, l, r, val )
```

```
segtree[idx] = min ( segtree[2*idx+1], segtree[2*idx+2] )
```

```
int query (idx, start, end, l, r)
```

```
if (lazy[idx] != 0)
```

```
    segtree[idx] += lazy[idx]
```

```
    if (start != end)
```

```
        { lazy[2*idx+1] += lazy[idx]
```

```
          lazy[2*idx+2] += lazy[idx]
```

```
        lazy[idx] = 0
```

```
if (r < start || l > end) return
```

```
INT_MAX
```

```
if (start >= l || end <= r) return
```

```
    segtree[idx]
```

```
    mid =  $\frac{start + end}{2}$ 
```

```
    query left
```

```
    query right
```

```
    return min(query left, query right)
```

Q. Given an array of N integers.

q queries

Type 1	l, r	(flip sign)
Type 2	(l, r)	sum of elements

$A = \overset{1}{2}, \overset{2}{3}, \overset{3}{4}$

$q=2$	T_1	1	2	{ -2 -3 4 }
	T_2	1	3	-1
	T_1	2	3	{ -2 3 -4 }
	T_1	2	2	{ -2 -3 -4 }
	T_2	2	3	-7

$A[N]$ given

build (0, 0, N-1)

for (i=1 ; i ≤ q ; i++)

{
 if (Type1)
 {
 update (0, 0, N-1, 1, x)
 }
 else
 {
 print (query (0, 0, N-1, 1, x))
 }
 }

void build (idx, start, end)

{
 if (start == end)
 {
 segtree [idx] = A[start]
 return
 }
 int mid = (start + end) / 2

build (2idx+1, start, mid)

build (2idx+2, mid+1, end)

segtree[idx] = segtree[2idx+1] + segtree[2idx+2]

void update (idx, start, end, l, r) 6 = 4 5 -3
 -6 = -4 -5 3

if (lazy[idx] != 0)

segtree[idx] = segtree[idx] * -1

if (start != end)

{
 lazy[2idx+1] = ! lazy[2idx+1]
 lazy[2idx+2] = ! lazy[2idx+2]

lazy[idx] = 0

if (l > end || r < start) return

if (start >= l && end <= r)

segtree[idx] = segtree[idx] * -1

if (start != end)

{
 lazy[2idx+1] = ! lazy[2idx+1]
 lazy[2idx+2] = ! lazy[2idx+2]

return mid = start + end

$$\text{mid} = \frac{\text{start} + \text{end}}{2}$$

update (2idx+1, start, mid, l, r, val)

update (2idx+2, mid+1, end, l, r, val)

segtree[idx] = segtree[2idx+1] + segtree[2idx+2]

int query (idx, start, end, l, r)

if (lazy[idx] != 0)

segtree[idx] = segtree[idx] * -1

if (start != end)

{ lazy[2idx+1] = !lazy[2idx+1]
lazy[2idx+2] = !lazy[2idx+2]

lazy[idx] = 0

if (r < start || l > end) return 0

if (start >= l || end <= r) return
segtree[idx]

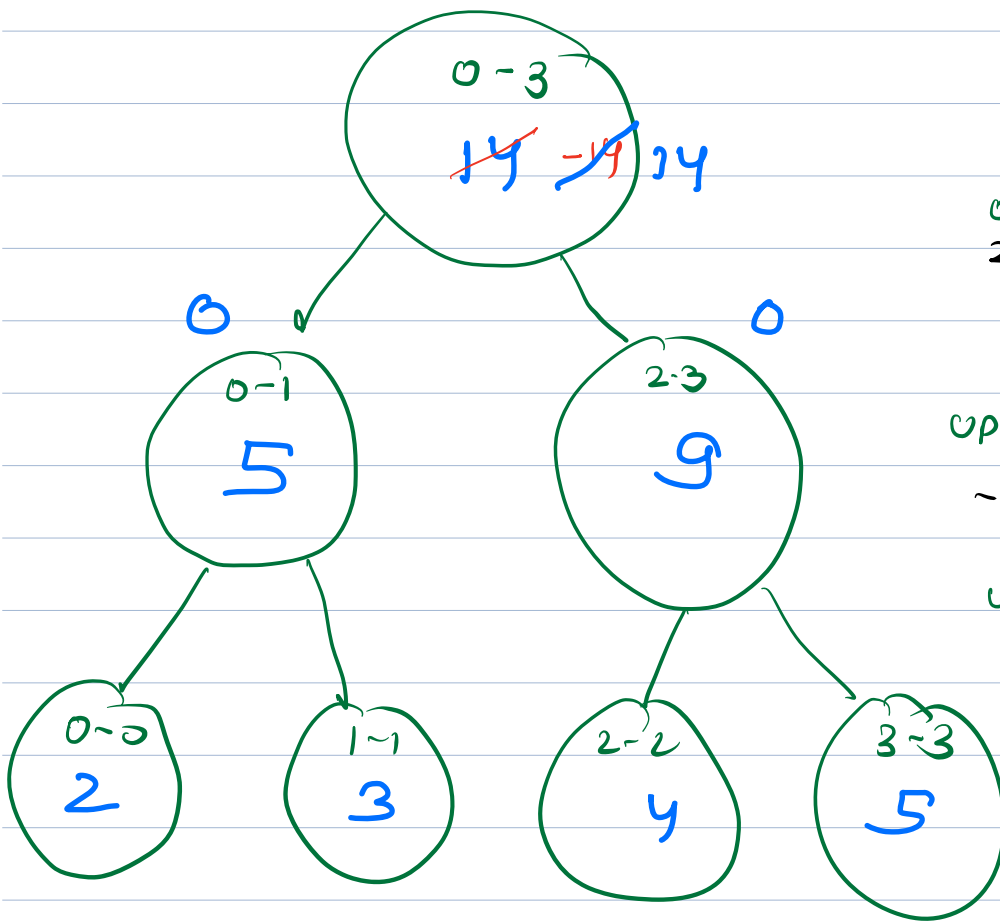
mid = $\frac{\text{start} + \text{end}}{2}$

query left

query right

return query left + query right

Doubts



0 1 2 3
2, 3, 4, 5

update (0-3)

~2, ~3, ~4 ~5

update (0-3)