

DeCent Vote

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Abstract— This application is to provide a web-based decentralized voting application where each and every voter has a fair role to play. It uses blockchain mechanism and prevents double-voting and fraudulent voting. In this application, each voter registers with valid details and then votes to the corresponding party. Each registered voter has a unique voter ID. A set of authentication keys is generated against each voter. When a user votes, it is checked whether it is the first time, or it is already done before and it is ensured that there can be utmost only one vote casted from each voter. After passing through this verification, the vote is then recorded as done. All votes that are accepted require consensus across the network. No single node controls it and every node is an owner and each vote is peer-peer verified. Given that every recorded vote on the blockchain needs consensus on the network and the fact that it is merely impossible to manipulate too many systems at the time, the chances of fraudulent votes are very low. The voting system is 100% transparent, no central authority owns it and the voter's identification remains confidential. Also, people don't have to leave their sofa to cast their votes, but they can do it online at their own convenience. Since this application is de-centralized, it ensures high availability and data security.

It provides a. Anonymity of voters ensures that whoever is casting a vote, they are authorized to do so.
b. Only one vote per person: No one would be able to vote more than once in the same election.
c. Data integrity: Ensures that once voted, it cannot be manipulated.

Keywords—Blockchain Platform, Hyperledger Fabric, JavaScript

INTRODUCTION

With the rise of blockchain technology, the core concept of decentralization has gradually drawn attention. In this context, the main objective of this research is to realize more convenient and secure applications through the use of blockchain technology. Currently, the service industry, such as the financial and banking industry, transmits private information through a trusted third party. However, they are facing many problems and complicated procedures. Since the blockchain technology and smart contract have the characteristics of decentralization, the researchers analyzed the architecture of the existing e-voting systems and found the integration of blockchain and smart contract into the application, which could enhance data verifiability and lower the cost while maintaining the openness and transparency of the voting.

The anonymities of voters, the security of ballot transmission and the verifiability of votes during the billing phase are the most fundamental requirements for voting. The anonymity and security can be achieved by the secret sharing scheme with Paillier's public-key cryptosystem while the verifiability of votes can be realized by taking advantage of the transparency and non-repudiation of blockchain. Voters can calculate the ballots and verify the election results on their own without a trusted third party

Application working

An application is built using the ReactJS and NodeJS. The overall front end is written in ReactJS. The user is asked to Sign up or Log in into the application using his trusted credentials, like driver's license or passport. Upon login, the user is presented to enter the details of his voter ID and to pick his candidate.

Proposed Working on Hyperledger and SmartContracts

Hyperledger Fabric is a permissioned blockchain infrastructure, originally contributed by IBM and Digital Asset, providing a modular architecture with a delineation of roles between the nodes in the infrastructure, execution of Smart Contracts (called "chaincode" in Fabric) and configurable consensus and membership services. A Fabric Network comprises "Peer nodes", which execute chaincode, access ledger data, endorse transactions and interface with applications. "Orderer nodes" which ensure the consistency of the blockchain and deliver the endorsed transactions to the peers of the network, and Membership Service Providers (MSPs), generally implemented as a Certificate Authority, managing X.509 certificates which are used to authenticate member identity and roles.^[13]

Fabric is primarily aimed at integration projects, in which a Distributed Ledger Technology (DLT) is required, offering no user facing services other than an SDK for Node.js, Java and Go. Fabric supports chaincode in Go and JavaScript (via Hyperledger Composer, or natively since v1.1) out-of-the-box, and other languages such as Java by installing appropriate modules. It is therefore potentially more flexible than competitors that only support a closed Smart Contract language.

A smart contract also can be regarded as a secured stored procedure as its execution and codified effects like the transfer of some value between parties are strictly enforced and cannot be manipulated, after a transaction with specific contract details is stored into a blockchain or distributed ledger. That's because the actual execution of contracts is

controlled and audited by the platform, not by any arbitrary server-side programs connecting to the platform.

In 2018, a US Senate report said: "While smart contracts might sound new, the concept is rooted in basic contract law. Usually, the judicial system adjudicates contractual disputes and enforces terms, but it is also common to have another arbitration method, especially for international transactions. With smart contracts, a program enforces the contract built into the code." By implementing the Decree on Development of Digital Economy, Belarus has become the first-ever country to legalize smart contracts. Belarusian lawyer Denis Aleinikov is considered to be the author of a smart contract legal concept introduced by the decree

Security Concerns

A smart contract is "a computerized transaction protocol that executes the terms of a contract". A blockchain-based smart contract is visible to all users of said blockchain. However, this leads to a situation where bugs, including security holes, are visible to all yet may not be quickly fixed.

Such an attack, difficult to fix quickly, was successfully executed on The DAO in June 2016, draining US\$50 million in Ether while developers attempted to come to a solution that would gain consensus. The DAO program had a time delay in place before the hacker could remove the funds; a hard fork of the Ethereum software was done to claw back the funds from the attacker before the time limit expired.

Issues in Ethereum smart contracts, in particular, include ambiguities and easy-but-insecure constructs in its contract language Solidity, compiler bugs, Ethereum Virtual Machine bugs, attacks on the blockchain network, the immutability of bugs and that there is no central source documenting known vulnerabilities, attacks and problematic constructs.

Distributed Ledger Technology for Decentralized applications

The distributed ledger database is spread across several nodes (devices) on a peer-to-peer network, where each replicate and saves an identical copy of the ledger and updates itself independently. The primary advantage is the lack of central authority. When a ledger update happens, each node constructs the new transaction, and then the nodes vote by consensus algorithm on which copy is correct. Once a consensus has been determined, all the other nodes update themselves with the new, correct copy of the ledger. Security is accomplished through cryptographic keys and signatures.

PROPOSED BLOCKCHAIN-BASED E-VOTING SYSTEM

There are seven roles in the system: (1) Voters(V_i) with qualification for voting; (2) Registration Server (RS) which verifies the voter's identity and provides eligible voters with voting certificate, $Cert_{\delta V_i P}$; (3) Authentication Server (AS) verifies the certificate derives from RS which issued V_i 's $Cert_{\delta V_i P}$; (4) Voting Website (VWeb), the voting site of the system which is under the control of the electoral authorities; (5) Recording Center (RC) stores $Cert_{\delta V_i P}$ and ballot signatures when V_i is voting; (6) Distributed Data Servers

(DDS) store the encrypted coordinates of points of the selected number when V_i is voting; (7) Smart Contract (SC), a dynamic space to replace the functionality of the traditional bulletin board. It can count the ballots to enhance the credibility and reliability of the election. In the system, we assume that there are n_1 voters, n_2 candidates and 5 distributed data servers. Moreover, all the transmission procedures are executed via https connection.

3.1 Initial Phase Before the protocol, AS and RS have to generate their RSA-based public/private key pair ($e; N_P = \delta P d; N$ and $(e_0; N_0 P = d_0; N_0 \delta P$, respectively, where $\delta P d; N, d_0; N_0 \delta P$ are the signing key and $(e; N_P, (e_0; N_0 P$ are the public key used for signature verification. h denotes the hash function (e.g., SHA-256 or SHA-3). On the other hand, RC has its Paillier-based encryption and decryption key pair $(pk_{RC}; sk_{RC})$:

3.2 Registration Phase This phase includes two major procedures and is done off-line before the election: 306 J.-H. Hsiao et al. Step1: User code generation, $V_i; 1 \leq i \leq n_1$. It proceeds by each V_i as follows: • Pick a random number $t \in \mathbb{Z} \setminus N$ and generate its unique user code PID_i where $PID_i = \frac{1}{4} h_{\delta SSN V_i} ktP$: Here SSN is the social security number. Finally, send PID_i to RS for verification. Step2: Verification of V_i 's identity, $1 \leq i \leq n_1$, RS proceeds as follows: • Accept PID_i and issue voting certificate $Cert_{\delta V_i P} = \frac{1}{4} PID_i; Sig_{\delta f g} \delta \delta P PID_i$ to V_i if PID_i is correct. The certificate is the signature of PID_i signed by RS. • Publish the PID_i of eligible voters onto the bulletin board. Here the bulletin board is implemented by the smart contract (SC).

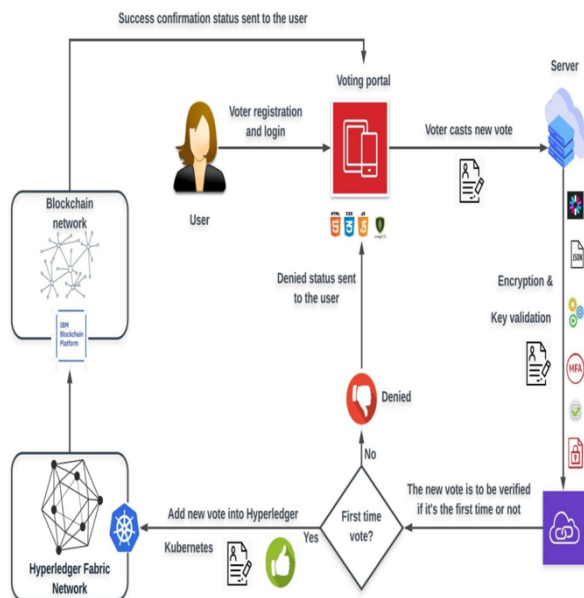
3.3 Voting Phase In which follows are the procedures for $V_i; 1 \leq i \leq n_1$, to obtain ballot signature and personal key pair from AS, then he/she can use the ballot signature for voting. Step1: V_i provides $Cert_{\delta V_i P}$ to AS and asks for a ballot signature. Step2: AS generates V_0 is Paillier-based public/private key pair $(pk_{V_i}; sk_{V_i} P$ and sends it back to V_i if V_0 is $Cert_{\delta V_i P}$ is correct. Step3: Assume that V_i wants to vote to the k -th candidate, $k \in \{1, \dots, n_2\}; V_i$ proceeds as follows: • After receiving $(pk_{V_i}; sk_{V_i} P$, compute $h_{\delta Pk}$ and generate $Ep_{k V_i} = \frac{1}{4} h_{\delta Pk} Pk$. • Pick n_2 random numbers $r_j \in \mathbb{Z} \setminus N$; and generate $c_j = \frac{1}{4} m_j r_j; 1 \leq j \leq n_2$. Here m_j means the ballot which is correspond to the j -th candidate. • Send $Ep_{k V_i} = \frac{1}{4} h_{\delta Pk} Pk$ and $c_j; 1 \leq j \leq n_2$; to AS. Step4: After receiving $Ep_{k V_i} = \frac{1}{4} h_{\delta Pk} Pk$ and c_j from $V_i; 1 \leq i \leq n_1; 1 \leq j \leq n_2$; AS proceeds as follows: • First, check each m_j from c_j , to avoid signing on incorrect or unrelated documents and compute the hash value of each c_j , (i.e., $h_{\delta c_j P}$). Second, sign each to generate the RSA signatures, $h_{\delta c_j} \square d$; by using key d . Finally, encrypt the signatures and the hash values of c_j , (i.e., $h_{\delta k_j P}$) using V_i 's public key pk_{V_i} to obtain $X_j = \frac{1}{4} Ep_{k V_i} h_{\delta c_j} d$ and $Ep_{k V_i} h_{\delta k_j P}$. • Pick n_2 random numbers, $k_j \in \mathbb{Z} \setminus N; 1 \leq j \leq n_2$ and compute $\delta Ep_{k V_i} = \frac{1}{4} Pk h_{\delta Pk} k_j$ and $\delta Ep_{k V_i} h_{\delta k_j P} P k_j$. Notice that according to the Paillier's additive homomorphic property, $Ep_{k V_i} = \frac{1}{4} h_{\delta Pk} Pk j = \frac{1}{4} Ep_{k V_i} = \frac{1}{4} h_{\delta Pk} Pk k_j$ and $Ep_{k V_i} h_{\delta k_j P} = \frac{1}{4} Ep_{k V_i} h_{\delta k_j P} k_j$. • Compute $M_{\delta P} = \frac{1}{4} Ep_{k V_i} k_j h_{\delta Pk} k h_{\delta k_j P} \square P h_{\delta c_j} \square d$; and sent $M_{\delta P} = \frac{1}{4} i j$ to V_i . Step5: $V_i; 1 \leq i \leq n_1$; proceeds as follows: Decentralized E-Voting Systems Based on the Blockchain Technology 307 • Obtain $M_{\delta P} = \frac{1}{4} i j; 1 \leq j \leq n_2$; then use his private key sk_{V_i} to decrypt $M_{\delta P} = \frac{1}{4} i j$, and to get n_2 ciphertexts and verify the k -th ciphertext by using AS's

public key to get only the k th ballot signatures; $h = c \oplus p_k \oplus d$. After voting, the candidate number k which chosen by V_i will be divided into k plaintext coordinates $PC = \{i; k \cdot \frac{1}{4} \oplus p_k x_k; y_k; 1 \oplus n_1; 1 \oplus k \cdot 5 \text{ via } 3 \oplus p; 5 \text{ secret sharing scheme which can be recovered from 3-out-of-5 plaintext coordinates. VWeb will use } V_i\text{'s public key } pk_{V_i} \text{ to encrypt the data and store it together with } PID_i \text{ in DDS. After DDS receives the coordinates, it will use } RC\text{'s public key } pk_{RC} \text{ to encrypt the coordinates and ultimately announce the coordinates and } PID_i; 1 \oplus n_1; \text{ by the SC for } V_i \text{ to check whether the correctness of counting.}$

3.4 Billing Phase After voting, an event will be sent via SC to notify all $V_i; 1 \oplus n_1$, the recording of ballots is ready. The procedures are as follows: Step1: AS proceeds as follows: • Publish V_i 's private key sk_{V_i} ; and the random number $r_j; 1 \oplus j \oplus n_2$; selected by voting phase onto SC. Step2: RC publishes its own private key sk_{RC} onto SC. Step3: SC computes and proceeds as follows: • First, decrypt $CC = \{i; k \cdot \frac{1}{4} \oplus p_k RC x_k; E_{pk_{V_i}} \delta y_k p; 1 \oplus k \cdot 5 \text{ and recover the plaintext coordinates } PC = \{i; k \cdot \frac{1}{4} \oplus p_k x_k; y_k; 1 \oplus k \cdot 5; \text{ by using } sk_{V_i} \text{ . Second, recover } k \text{ from plaintext coordinates by } 3 \oplus p; 5 \text{ secret sharing scheme and verify the correctness of the ballot signatures and check the value of } k \text{ to see whether it's consistent with the information on the ballot signatures by } mj.$

3.5 Security Analysis The e-voting systems proposed in this paper meet the following security requirements: Voter qualification, protecting the anonymity of voter's identity, non-repeatable, ballot eligibility and ballot verifiability. Due to the page limitation, we will not discuss this in detail.

The architecture diagram of the application is as follows.



With its distributed ledger, smart contracts, and repudiation

capabilities, blockchain is revolutionizing the way organizations do business, and the election industry is no exception. This developer code pattern shows you how to implement a web-based blockchain app using the IBM Blockchain Platform to facilitate voting and help ensure the prevention of double-voting.

Technology Stack

Programming languages & Libraries: JavaScript
Technologies: Blockchain Platform, Hyperledger Fabric,
Web Technologies: HTML5, CSS3, JavaScript, Vue, Node.JS
Miscellaneous: SOA, Agile, Git, Trello.

ACKNOWLEDGMENT

The authors gratefully acknowledge the contribution of Horea Porutiu for providing the guidelines on the e-vote platform. Also, we would like to thank Rakesh Ranjan and the IBM Design Team for their inputs on this topic which also helped us in research.

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