

학습 목표

힙(Heap)과 BST와의 차이점을 이해하고

힙(Heap)의 시간복잡도를 알 수 있다



Data Structures in Python Chapter 8

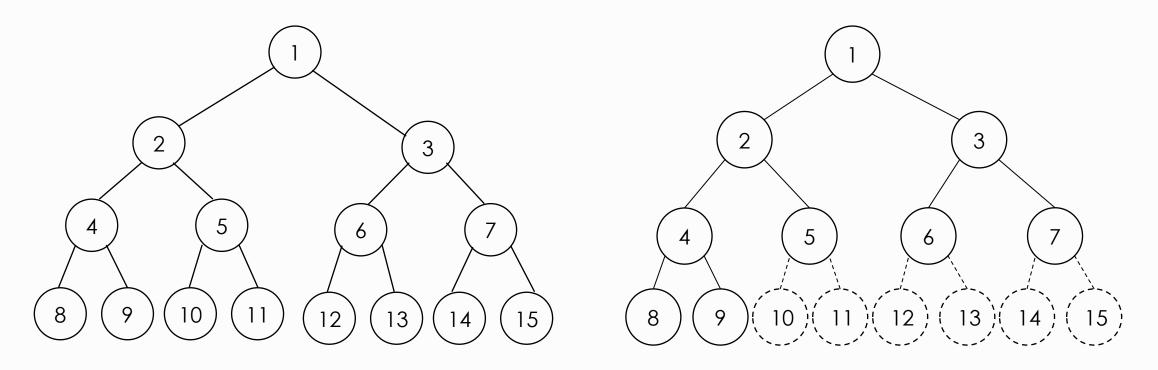
- Heap and Priority Queue
- Heap Coding
- Min/MaxHeap and Heap sort

Agenda & Readings

- Heap and Priority Queue
 - Complete Binary Tee (Review)
 - Heap and Priority Queue
 - Heap ADT
 - Time Complexity
- Reference:
 - Problem Solving with Algorithms and Data Structures

Binary trees - Properties

- **Definition:** A full binary tree of level k is a binary tree having 2^k 1 nodes, $k \ge 0$.
- **Definition**: A binary tree with n nodes and level k is **complete** if and only if its nodes correspond to the nodes numbered from 1 to n in the full binary tree of level k.

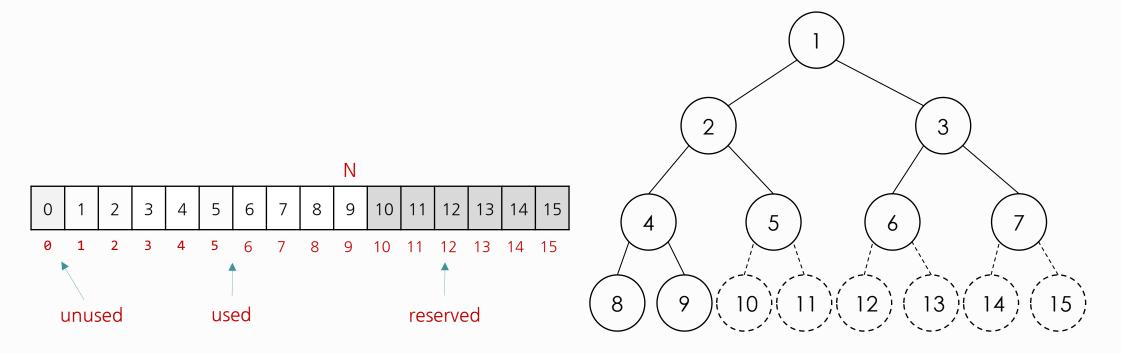


A **full** binary tree

A complete binary tree

Binary trees - Array representation

- A complete binary tree with n nodes, any node index i, $1 \le i \le n$, we have
 - parent(i) is at $\lfloor i/2 \rfloor$ If i = 1, i is at the root and has no parent
 - leftChild(i) is at 2i if 2i <= n. If 2i > n, then i has no left child.
 - rightChild(i) is at 2i+1 if 2i+1 <= n. If 2i+1 > n, then i has no right child.



A complete binary tree

- Heaps are frequently used to implement priority queues.
 - Because it provides an efficient implementation for priority queues.

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- Priority queues.
 - Queues with priorities associated to.
 - Example: A line waiting to be served at a bank and served FIFO except if a senior or a disabled person arrives in the line. They are served first. Seniors and disabled persons have higher priority than others.

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A typical ADT for Priority Queue

- Get the top priority element (min or max)
- Insert an element
- Delete the top priority element
- Decrease the priority of an element

- O(1)
- O(log n)
- O(log n)
- O(log n)

- Challenge: Find the largest M items in a stream of N items.
- Constraints: Not enough memory to store N items.



Order of insert of finding the largest M in a stream of N items

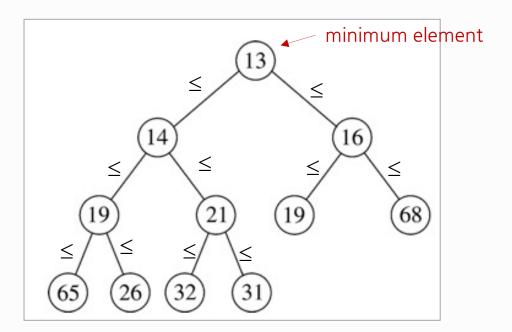
| implementation | insert | delete | min/max |
|-----------------|----------|------------|----------|
| unordered array | 1 | Ν | N |
| ordered array | N | 1 | N huge |
| goal | log N | log N | log N |
| | <u> </u> | • | |
| | AA* | mpossible? | priority |

Binary heap

- Binary heap: array representation of a heap-ordered complete binary tree
- Properties:
 - Heap-ordered: Parent's key no smaller than children's keys. [max-heap]
 - Heap-structure:
 A complete binary tree

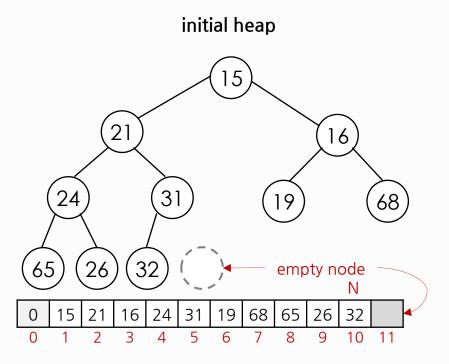
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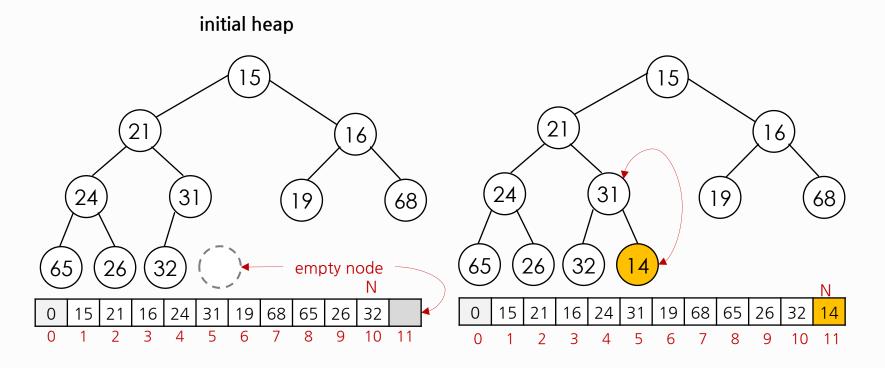


- Duplicates are allowed
- No order implied for elements which do not share ancestor-descendant relationship

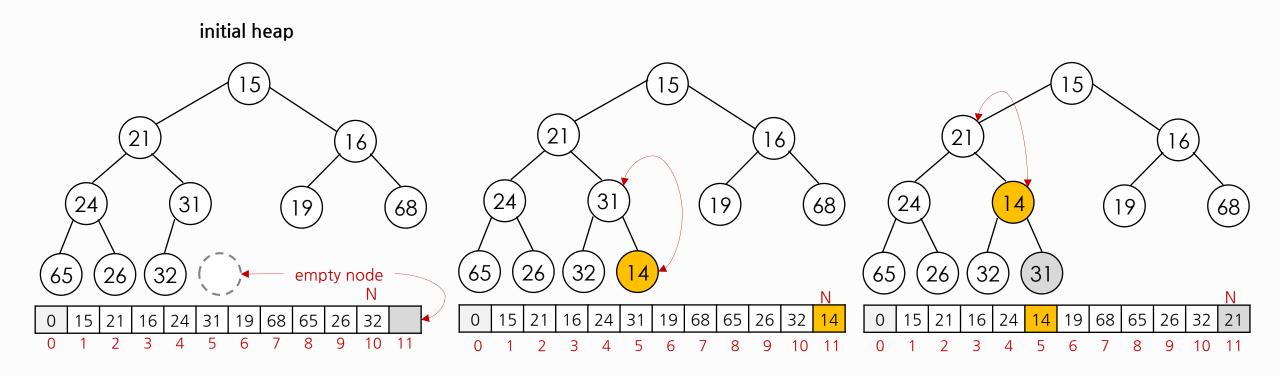
- Insert a new element while maintaining a heap-structure
- Move the element up the heap while not satisfying heap-ordered



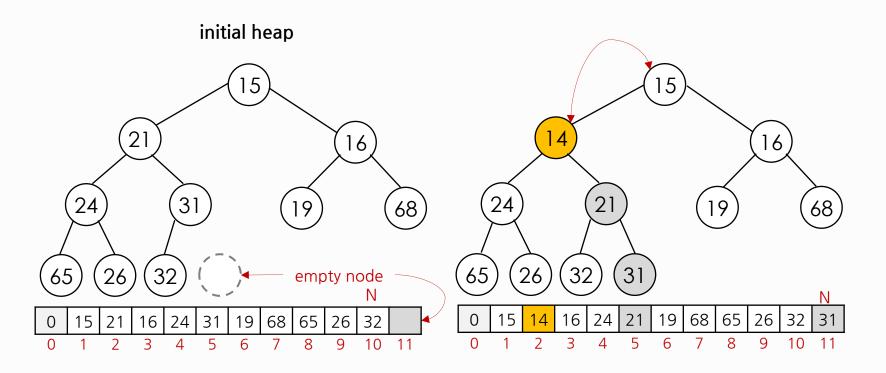
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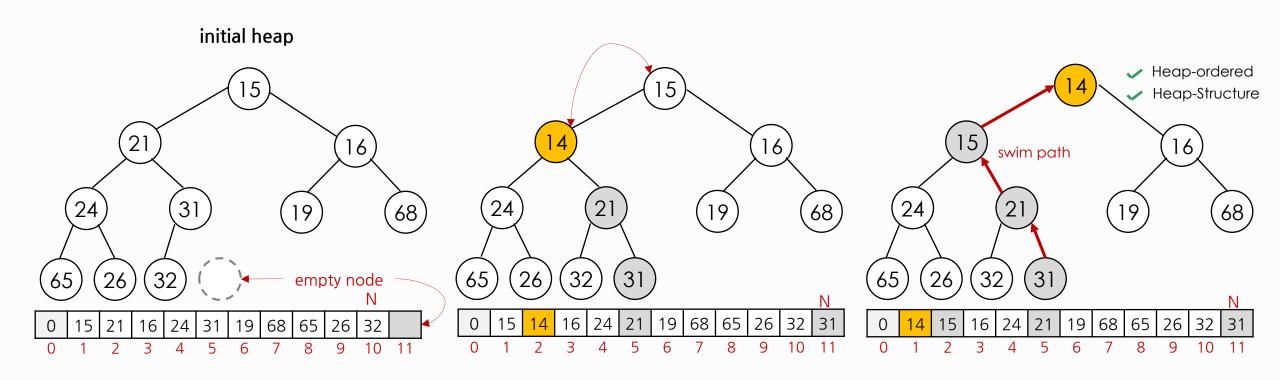
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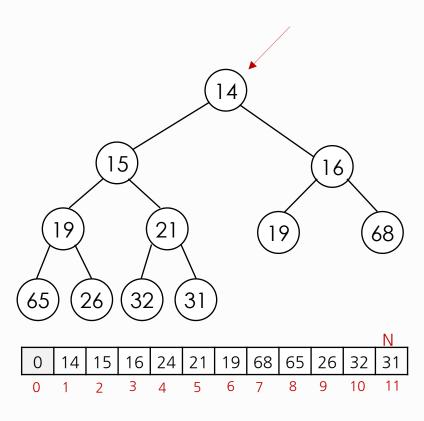


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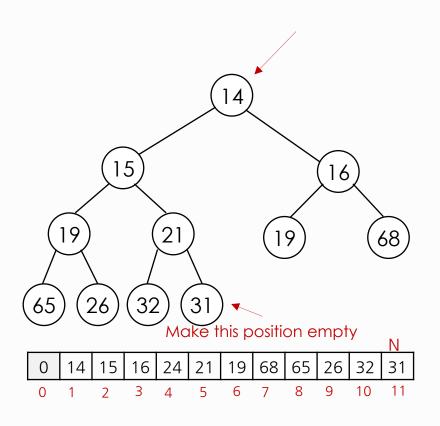


- Swap the root and the last element.
- Heap decreases by one in size.
- Move down (sink) the root while not satisfying heap-ordered.
 - Minimum element is always at the root (by min-heap definition).

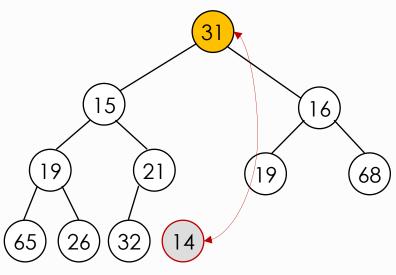
Which position of the node will be empty?



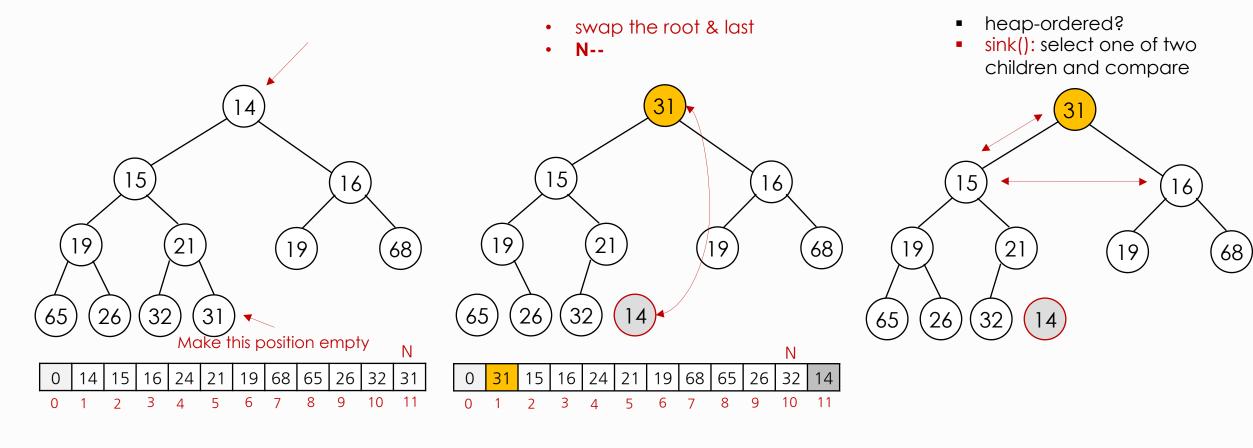
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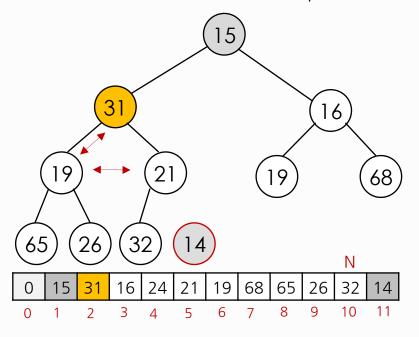
- swap the root & last
- N--



Which position of the node will be empty?

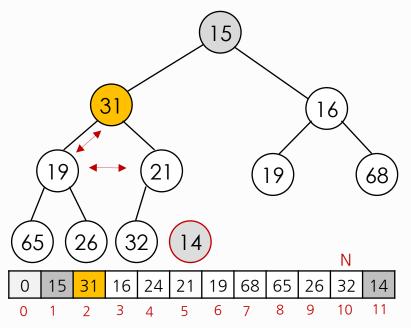


- heap-ordered?
- sink(): select one of two children and compare



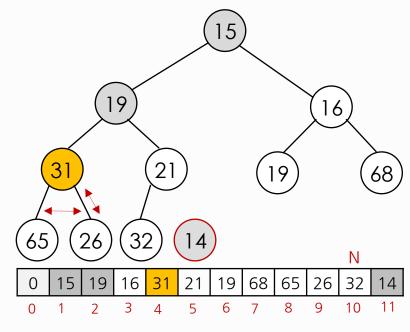
- Is 31 > min(14,16)?
- Yes swap 31 with min(14,16)

- heap-ordered?
- sink(): select one of two children and compare



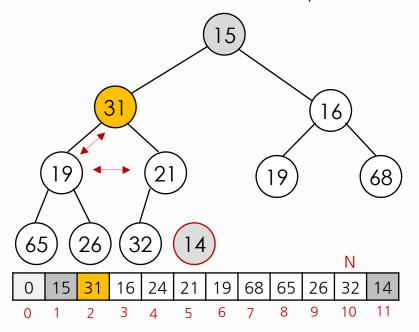
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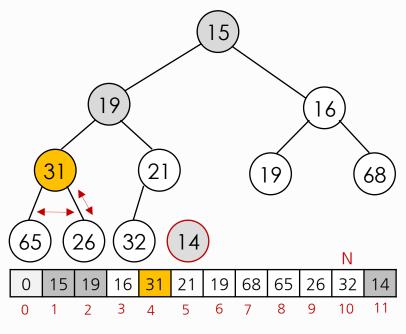
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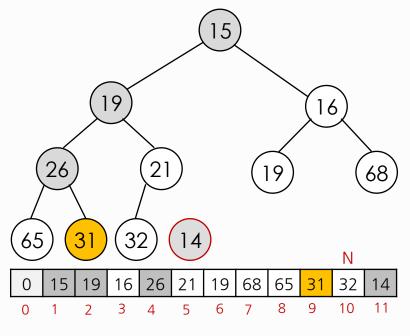
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- Is $31 > \min(19,21)$?
- Yes swap 31 with min(19,21)

- heap-ordered?
- sink(): select one of two children and compare



- Is 31 > min(65,26)?
- Yes swap 31 with min(65,26)
- Heap-ordered
- ✓ Heap-Structure

Binary heap: Time complexity:

- Level of heap is $\lfloor \log_2 N \rfloor$
- insert: O(log N) for each insert
 - In practice, expect less
- delete: O(log N) // deleting root node or any node
- increase/decrease key: O(log N)

| Implementation | Insert | Delete | max |
|-----------------|--------|--------|-----|
| Unordered array | 1 | N | Ν |
| Ordered array | Ν | 1 | 1 |
| Binary heap | log N | log N | 1 |

Mission Completed

학습 정리

- 1) 힙(Heap)은 우선순위 큐(Priority queue)에 주로 사용된다
- 2) **힙(Heap)에서의 삽입과 삭제 작업은 최대 O(**log n) 시간복잡도로 수행할 수 있다





Data Structures in Python

- Heap and Priority Queue
- Heap Coding
- Min/MaxHeap and Heap sort

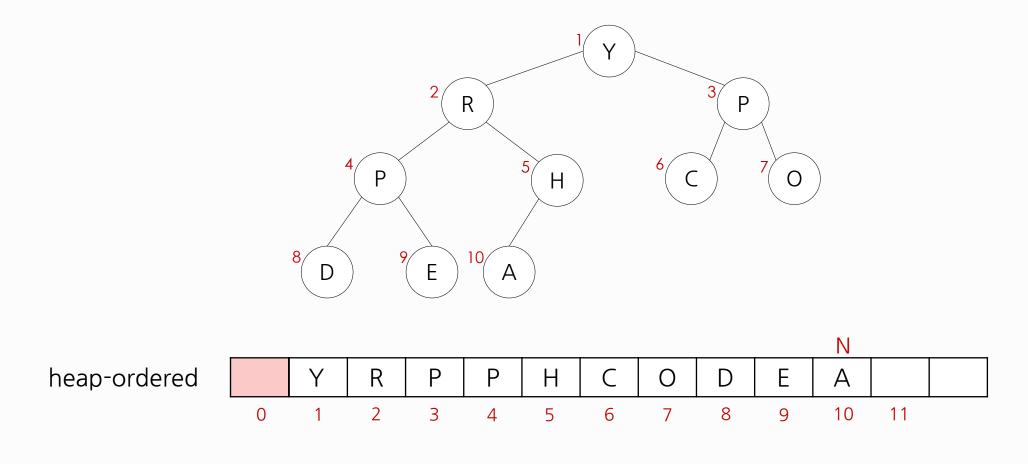
Proof:

https://stackoverflow.com/questions/9755721/how-can-building-a-heap-be-on-time-complexity https://www.insertingwiththeweb.com/data-structures/binary-heap/build-heap-proof/ https://www.quora.com/How-is-the-time-complexity-of-building-a-heap-is-o-n

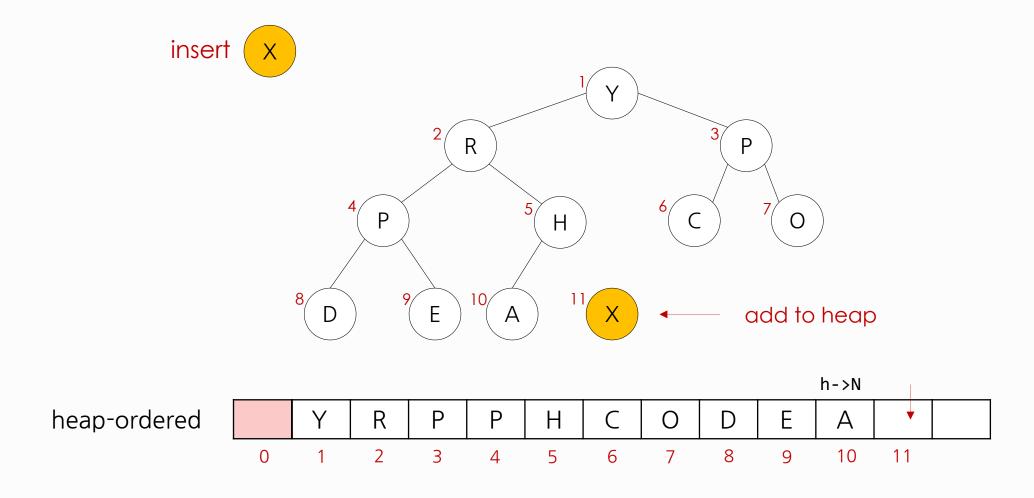
References in Korean:

https://ratsgo.github.io/data%20structure&algorithm/2017/09/27/heapsort/https://zeddios_tistory.com/56

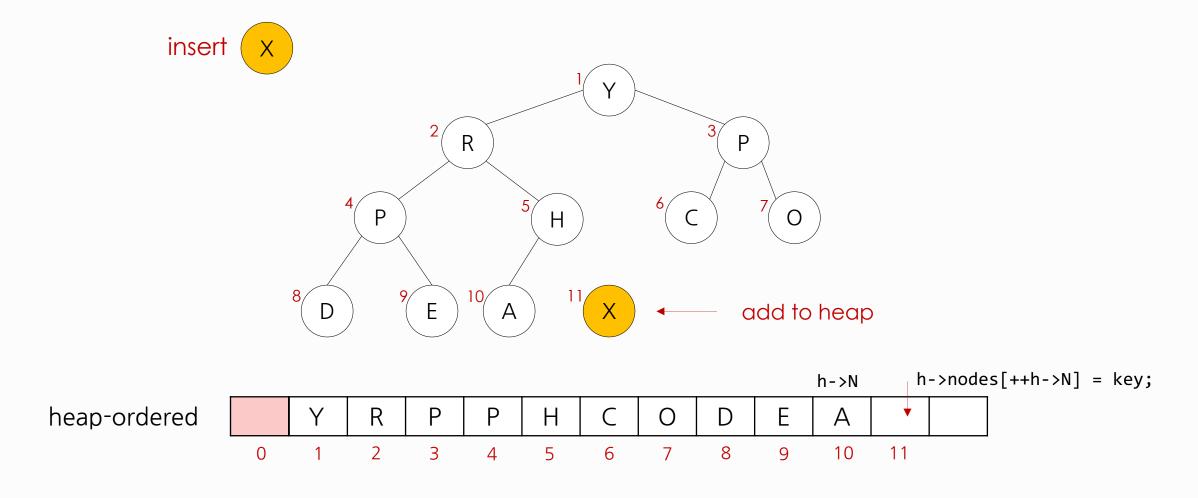
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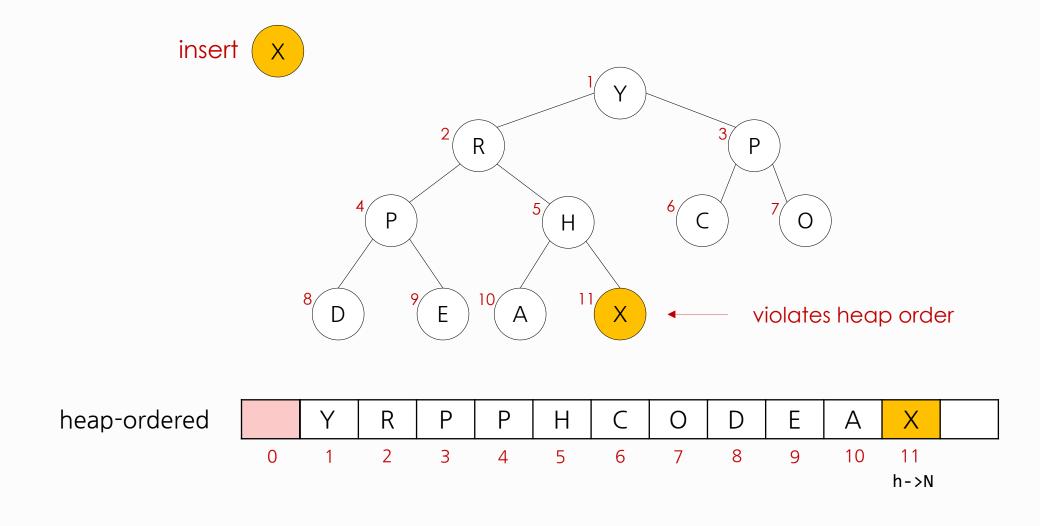
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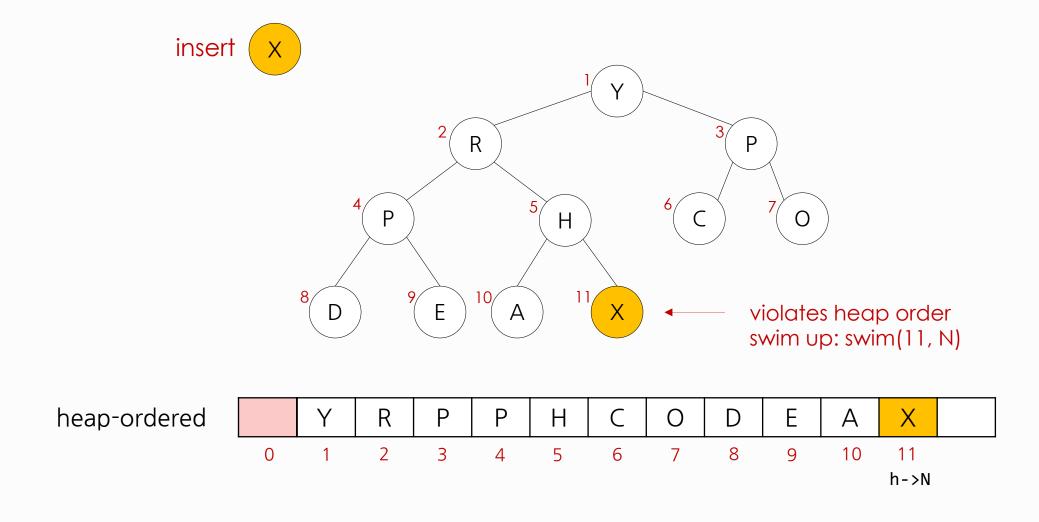
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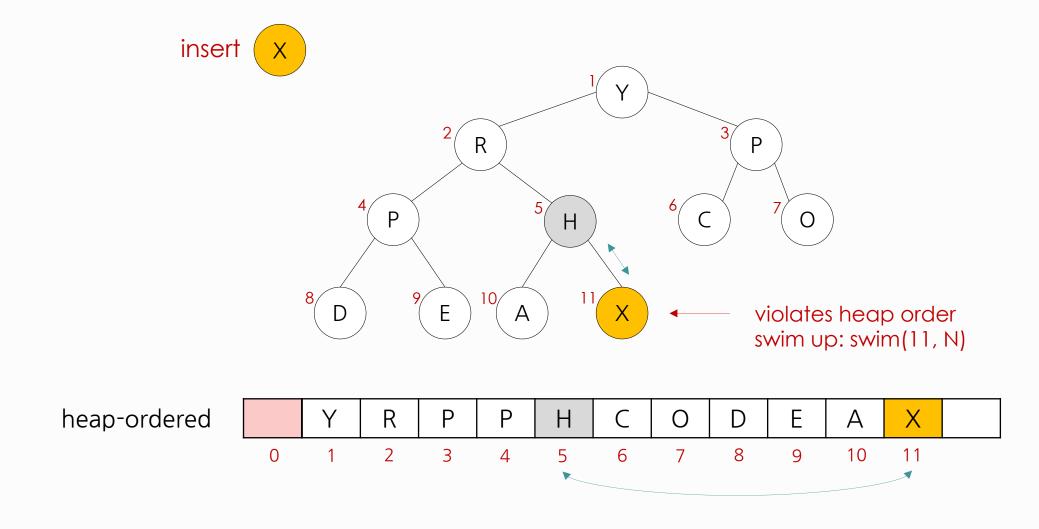
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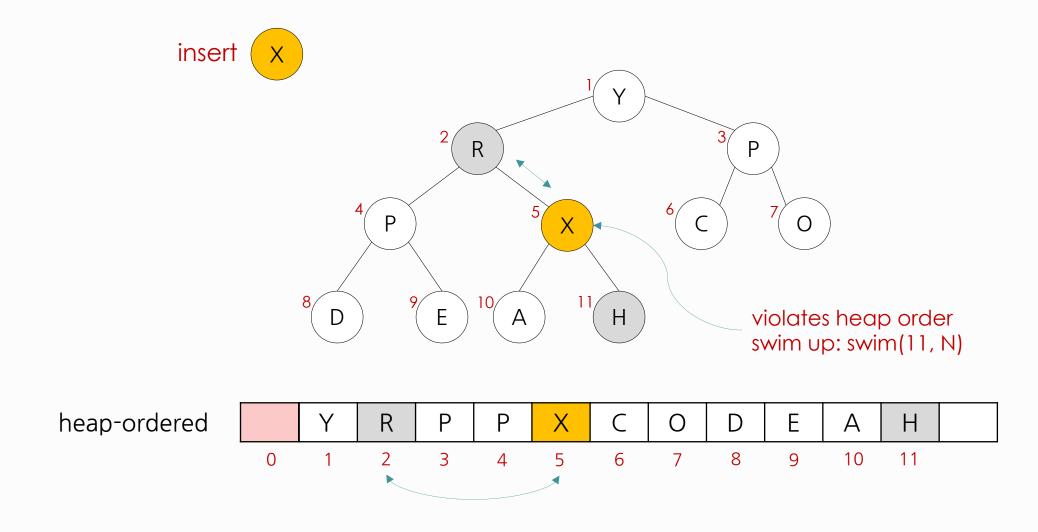
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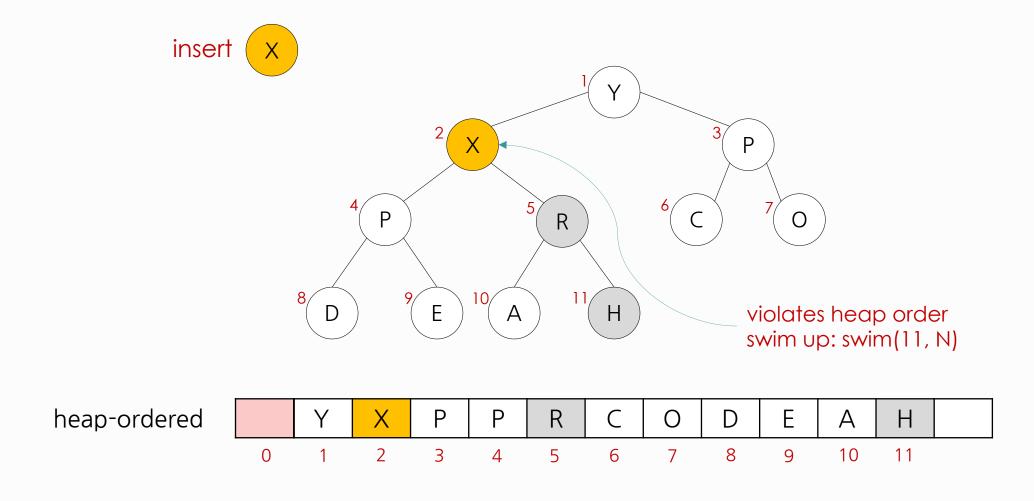
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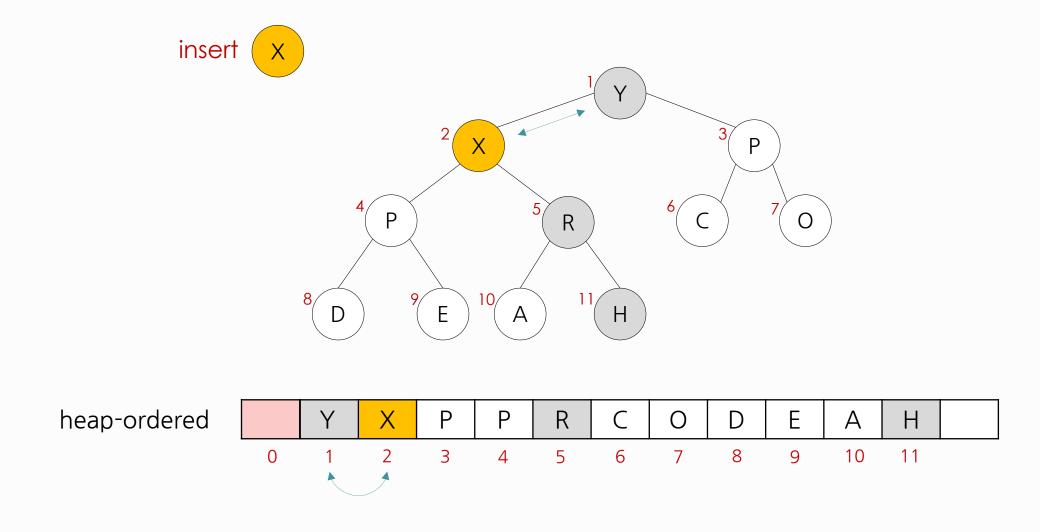
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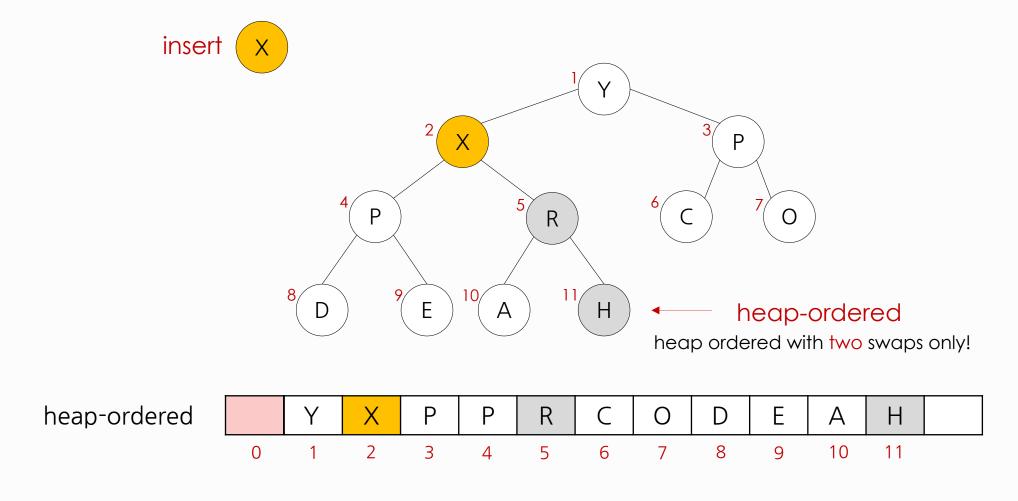
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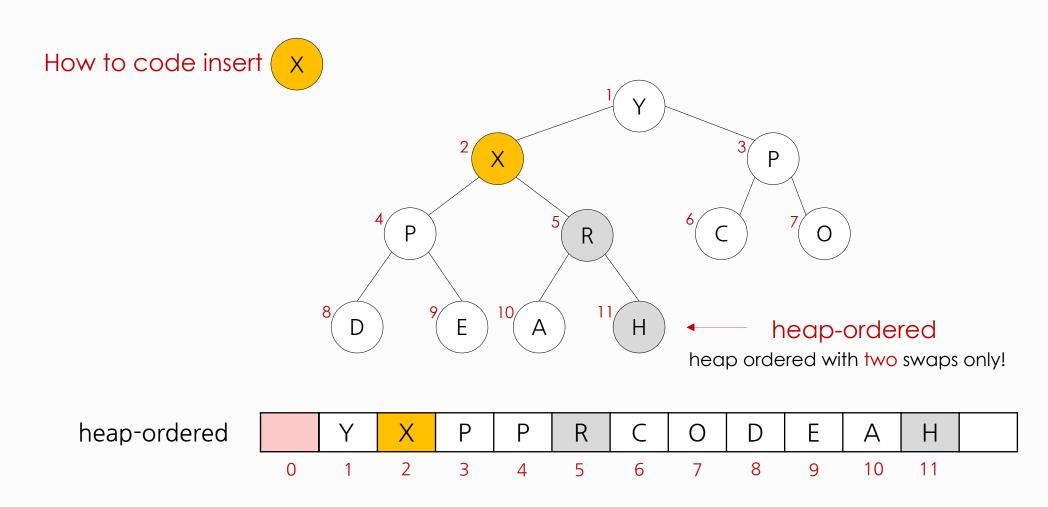
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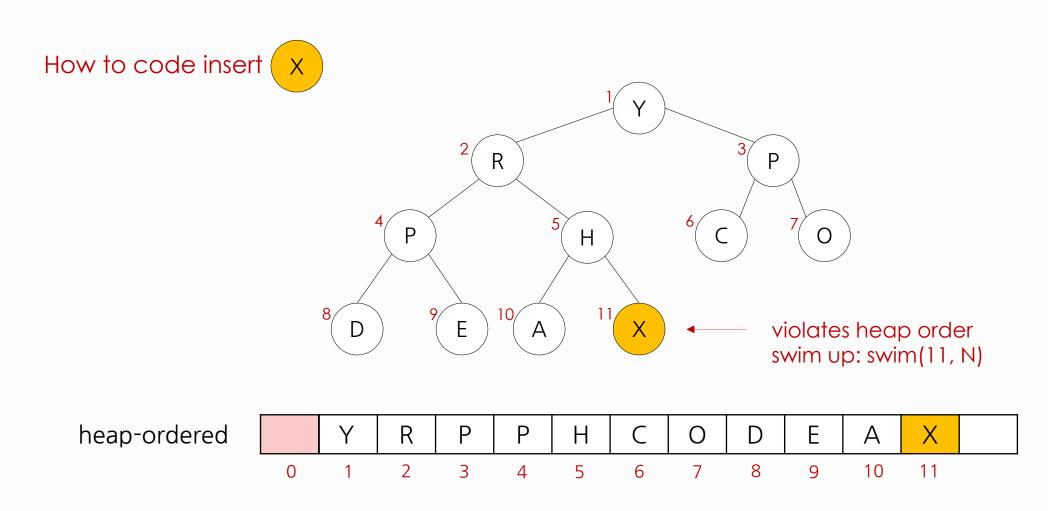
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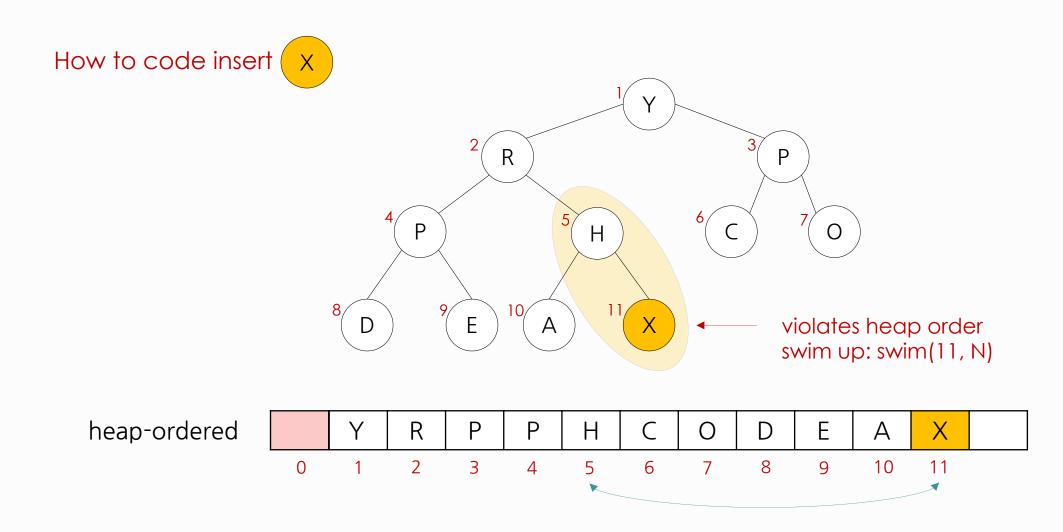
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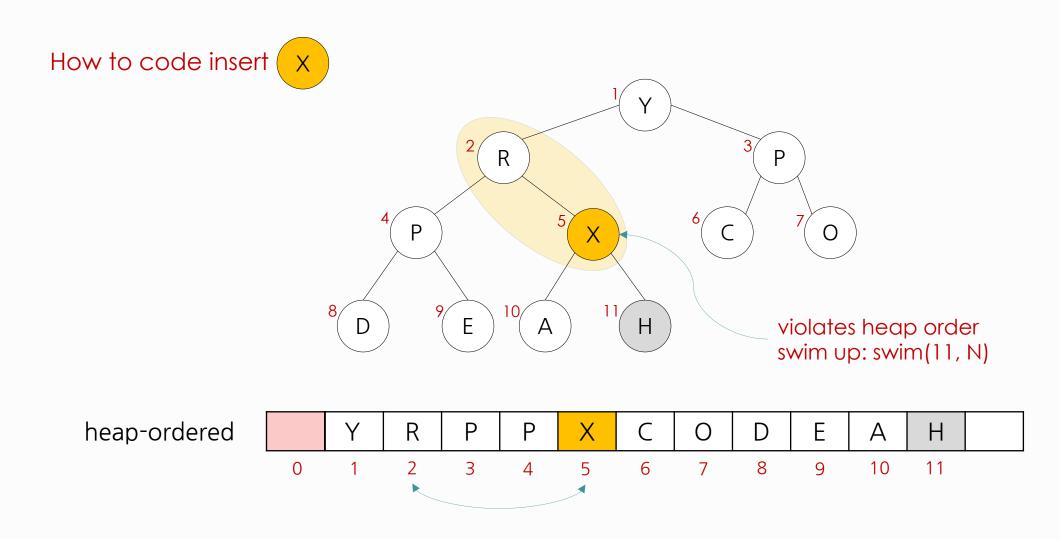
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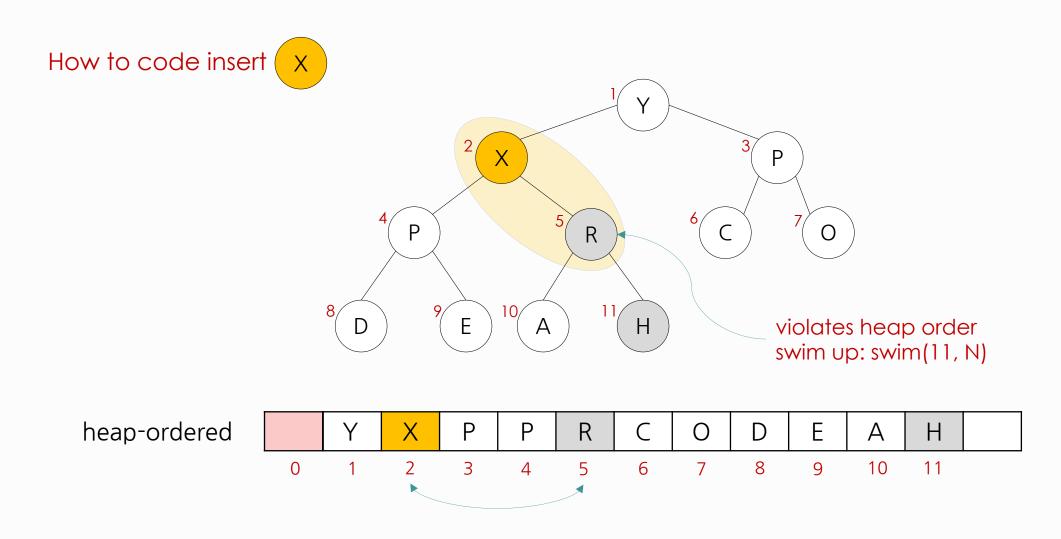
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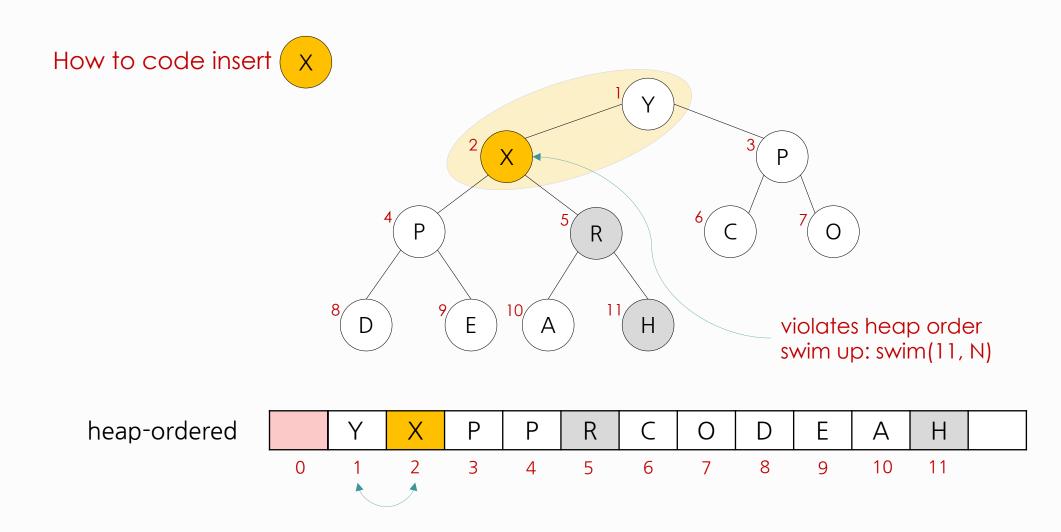
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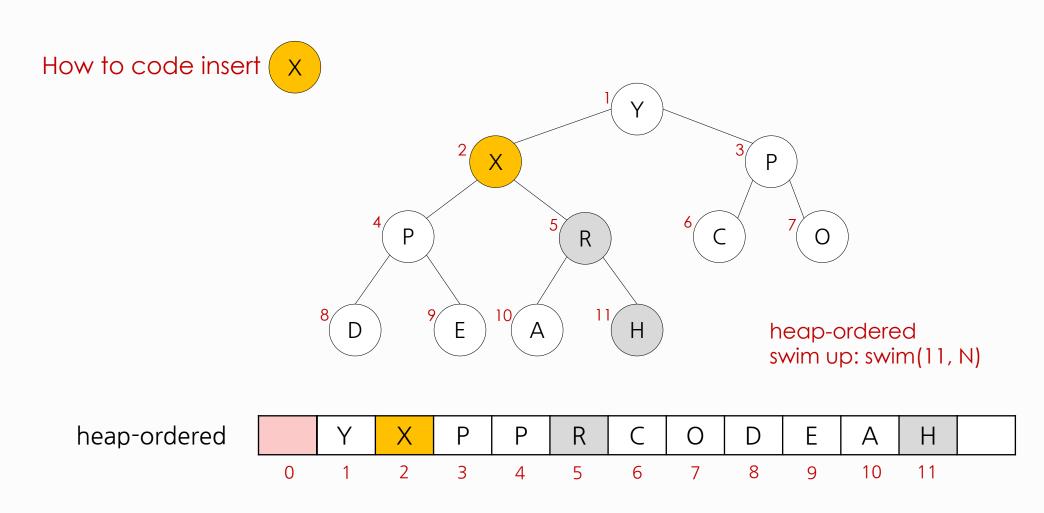
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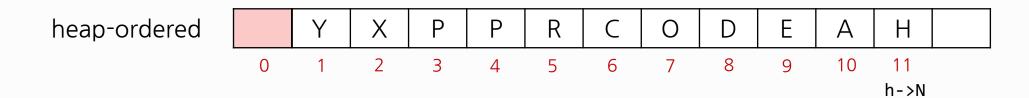


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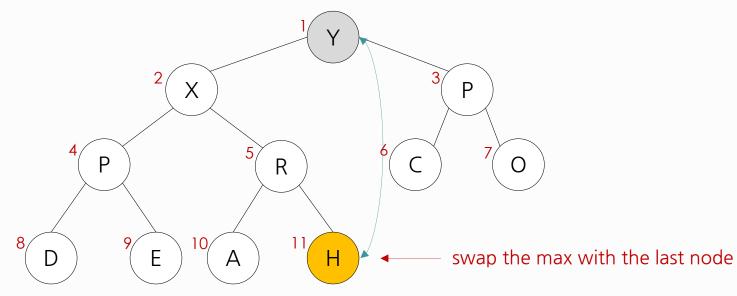
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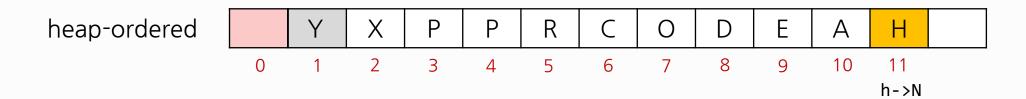
remove the max (root) - pop in priority queue



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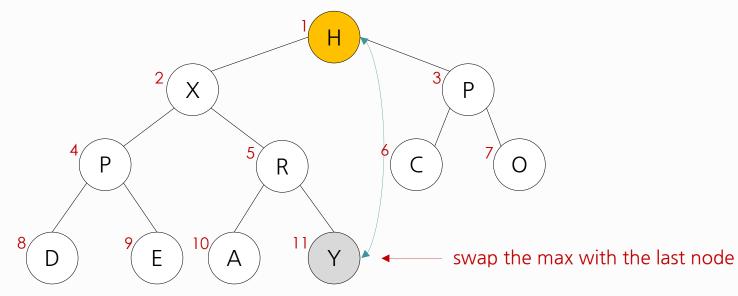
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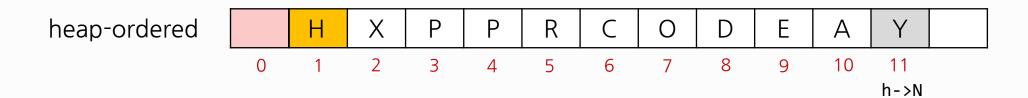




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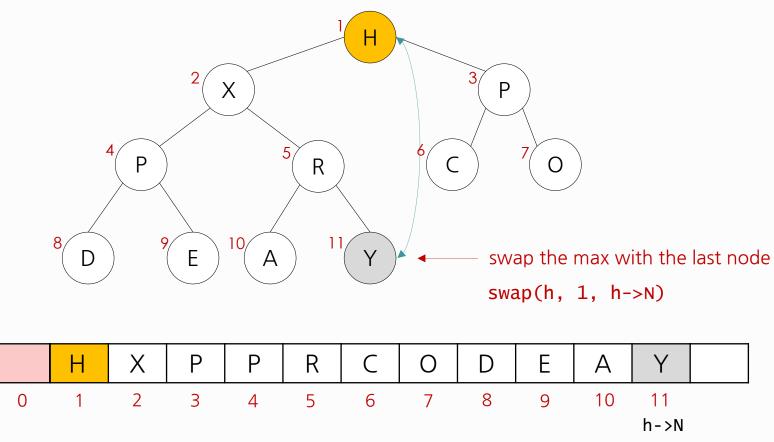




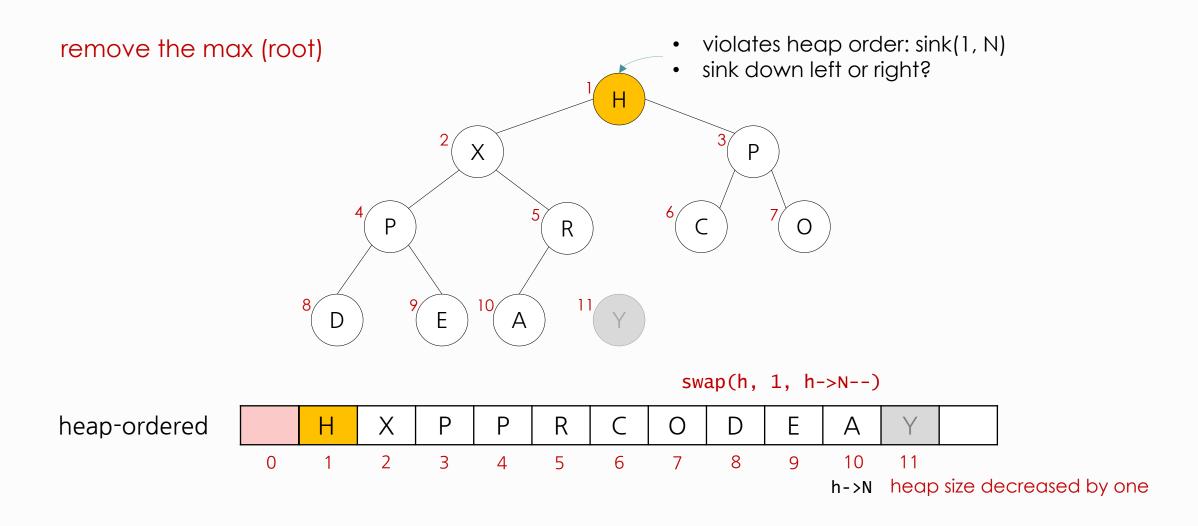
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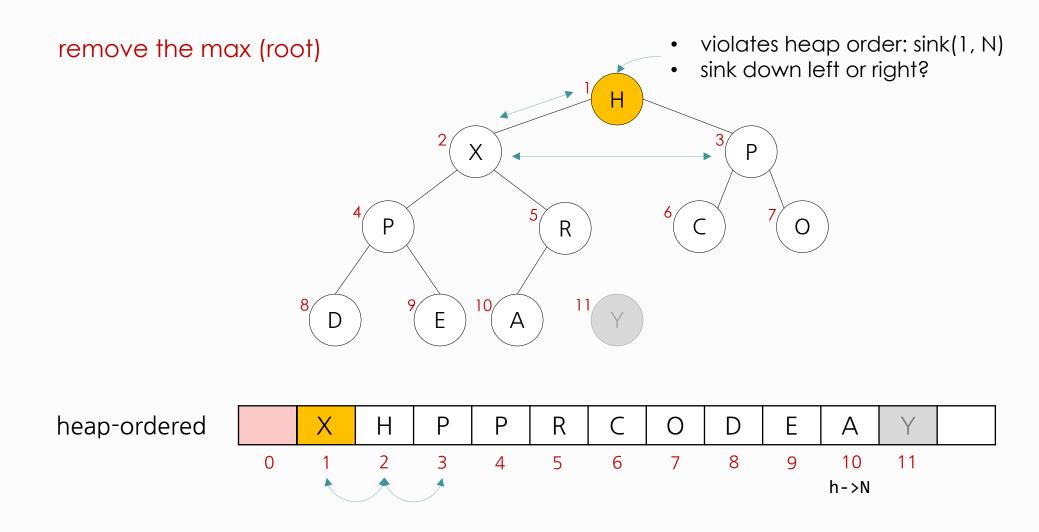
heap-ordered



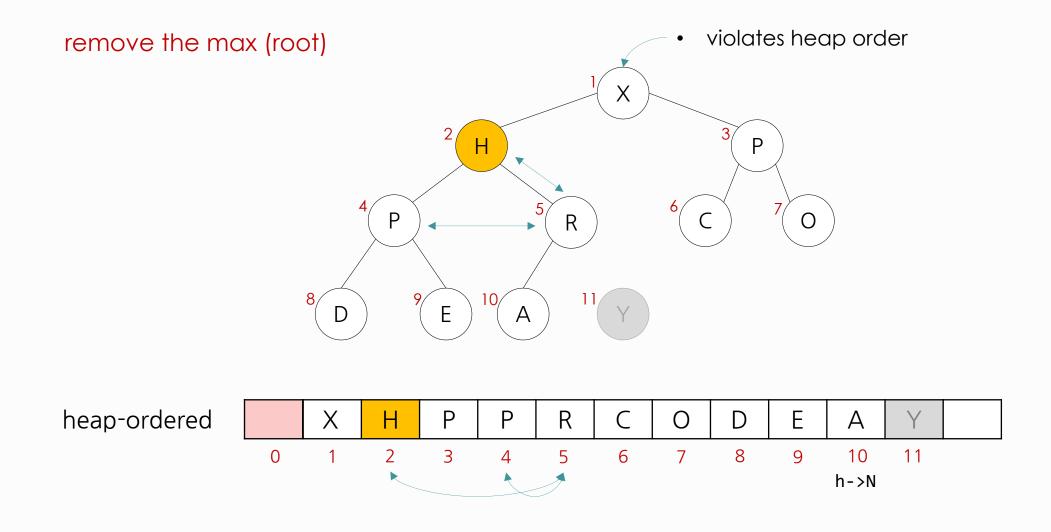
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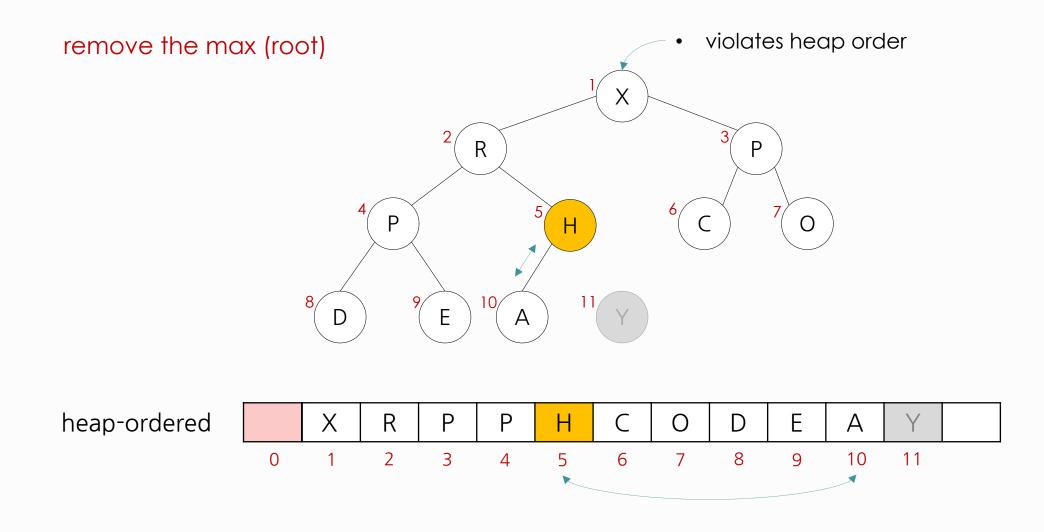
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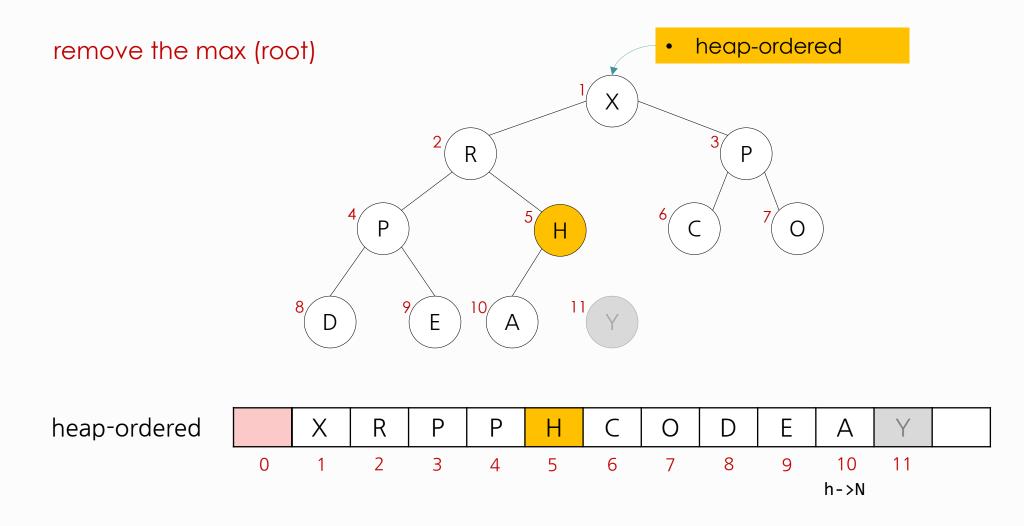
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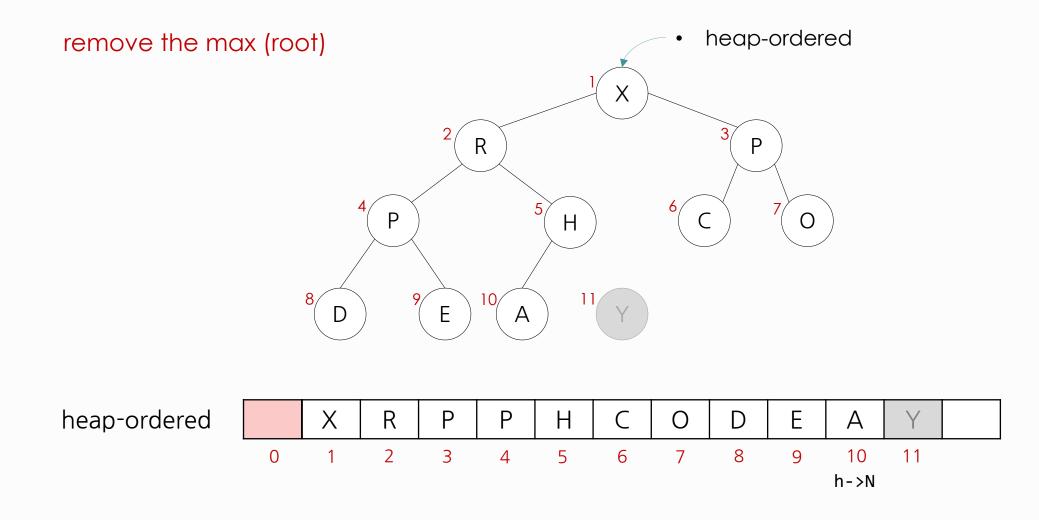
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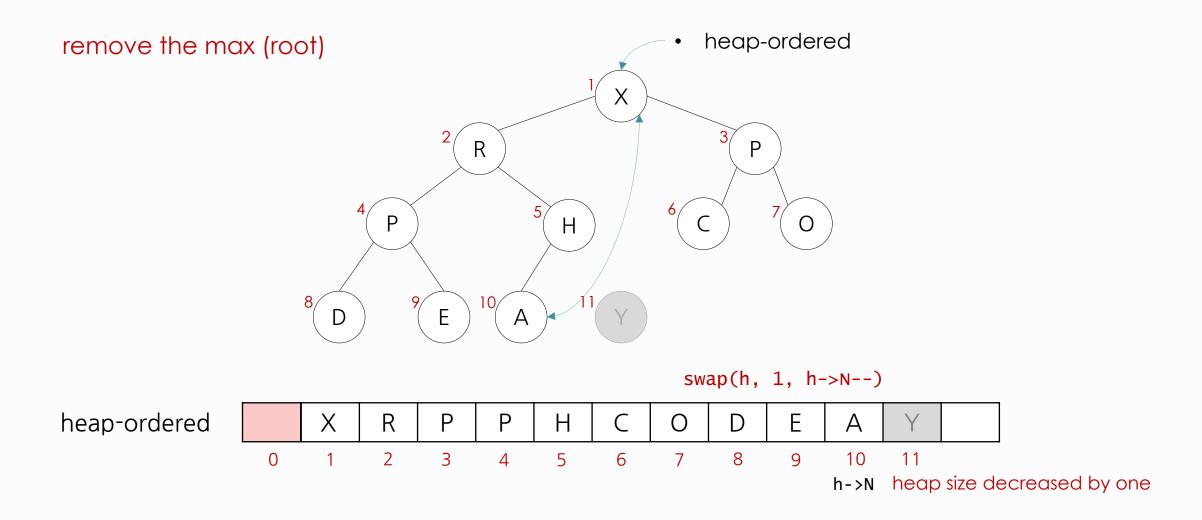
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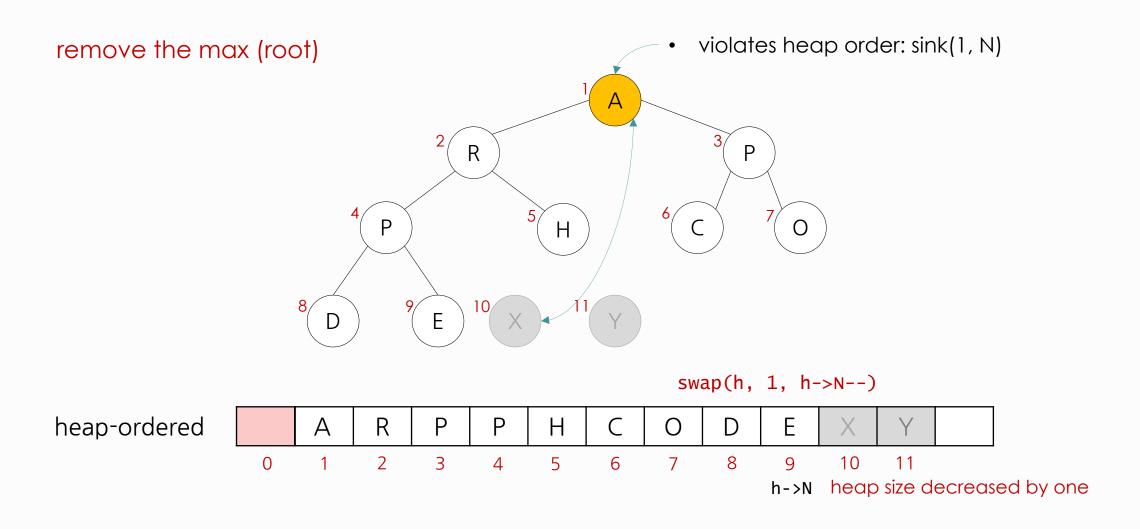
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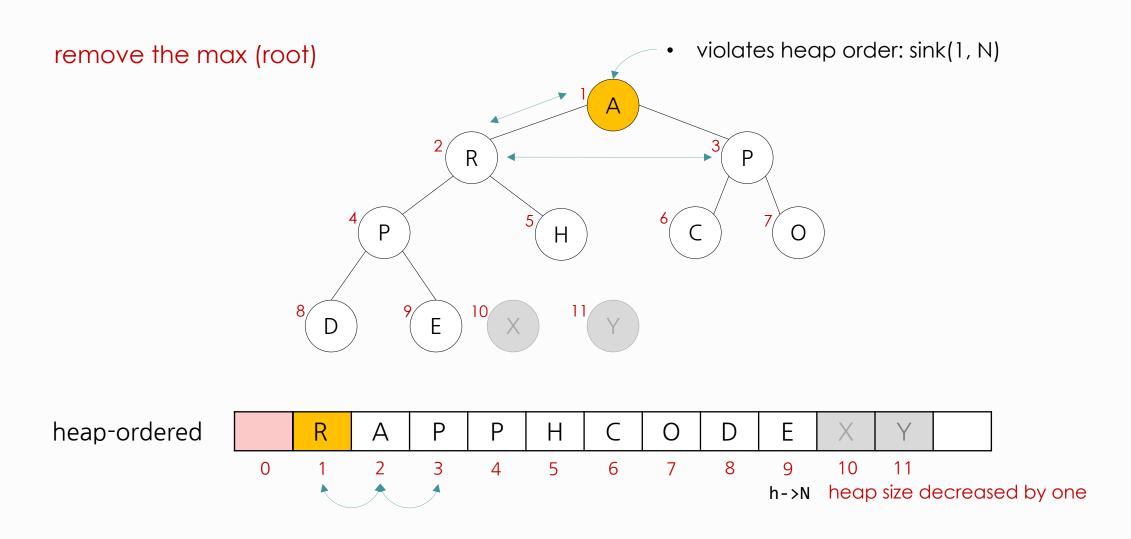
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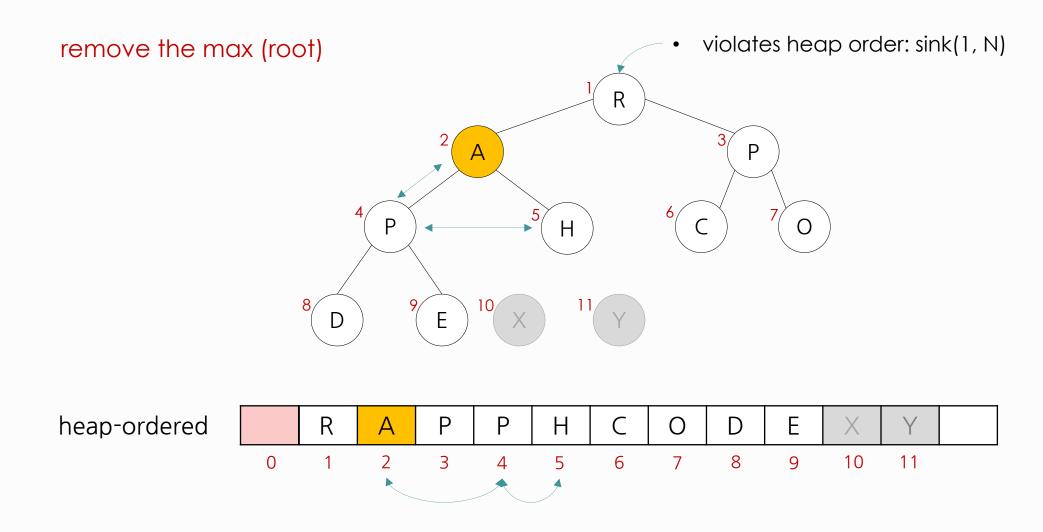
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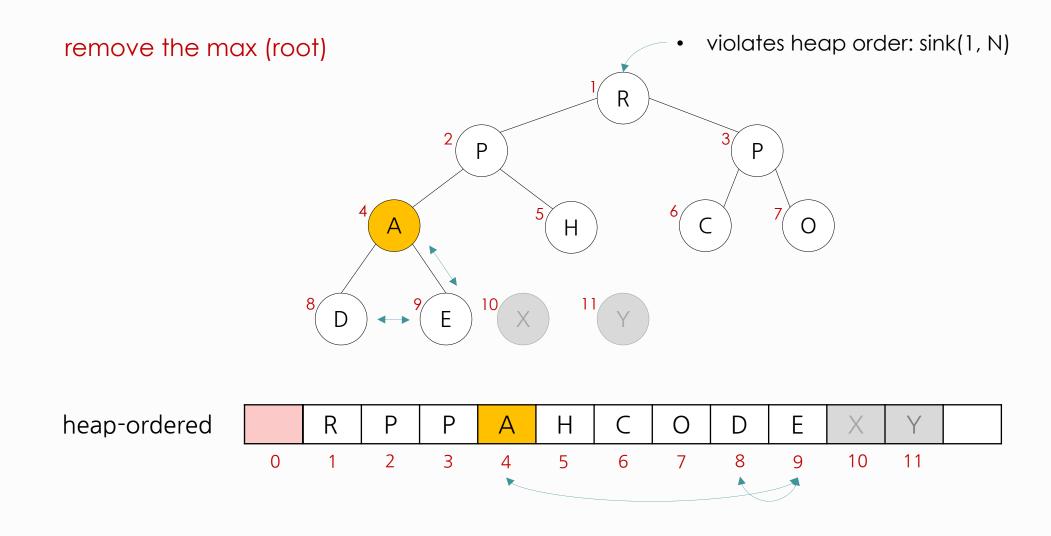
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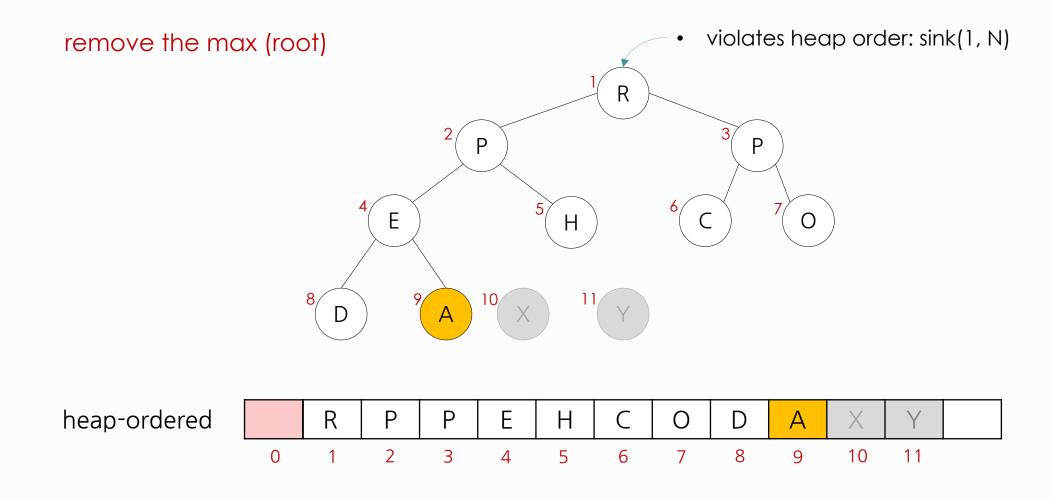
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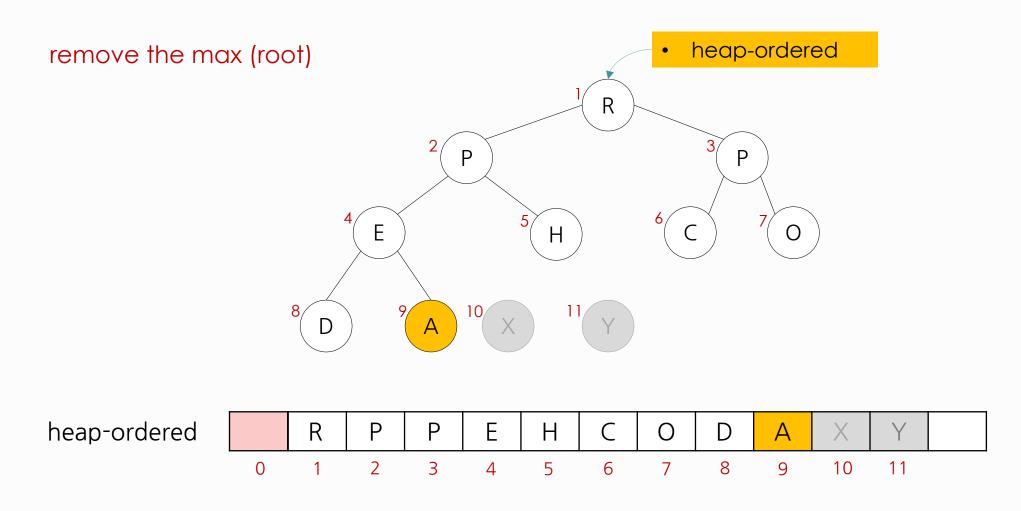
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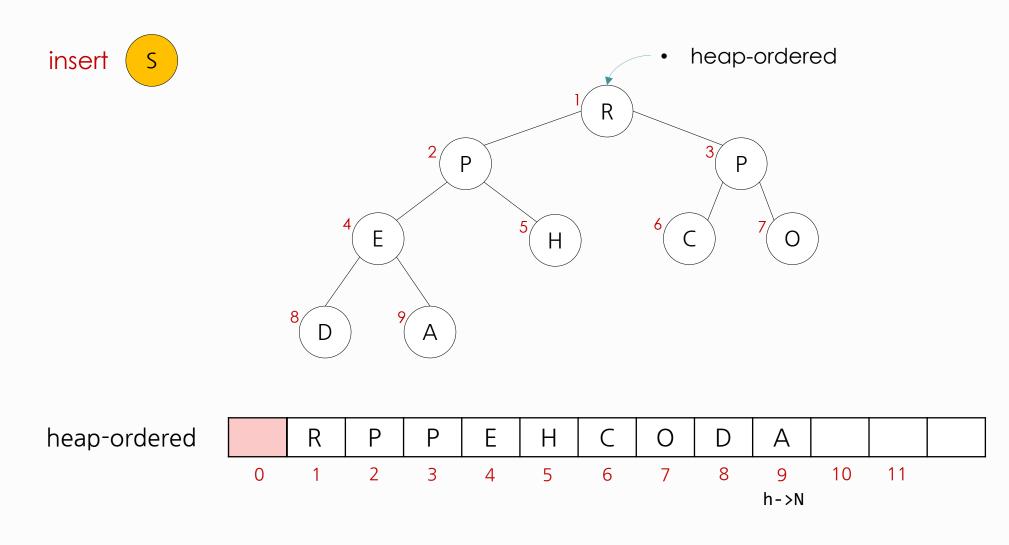
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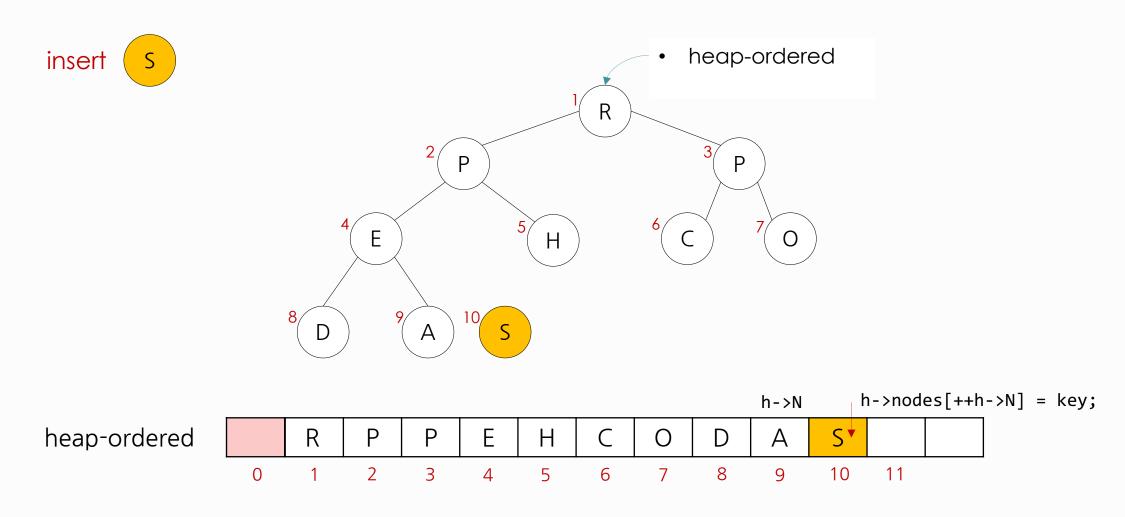
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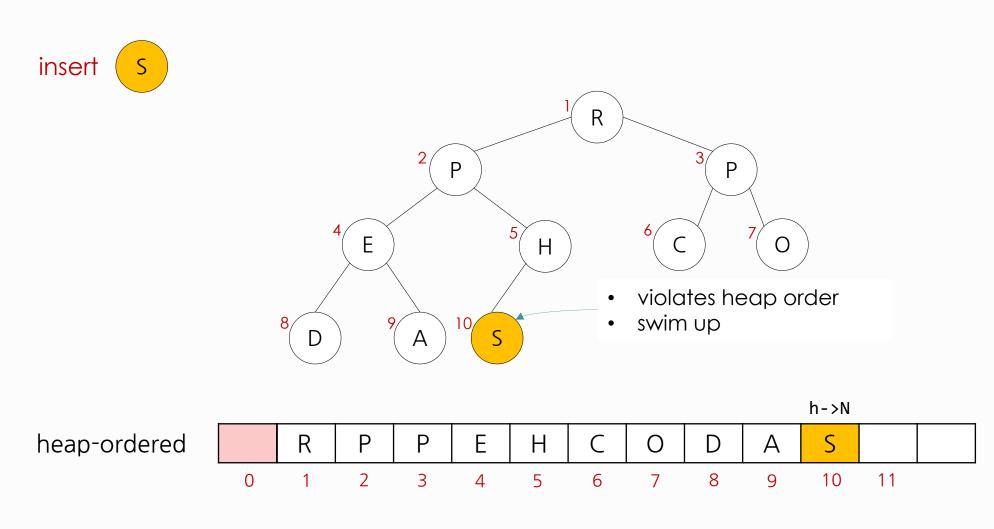
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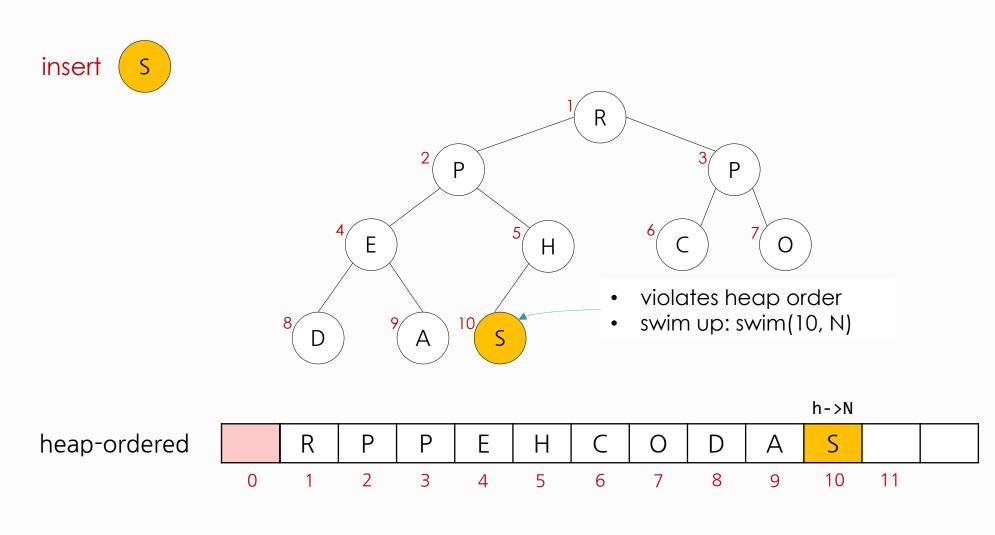
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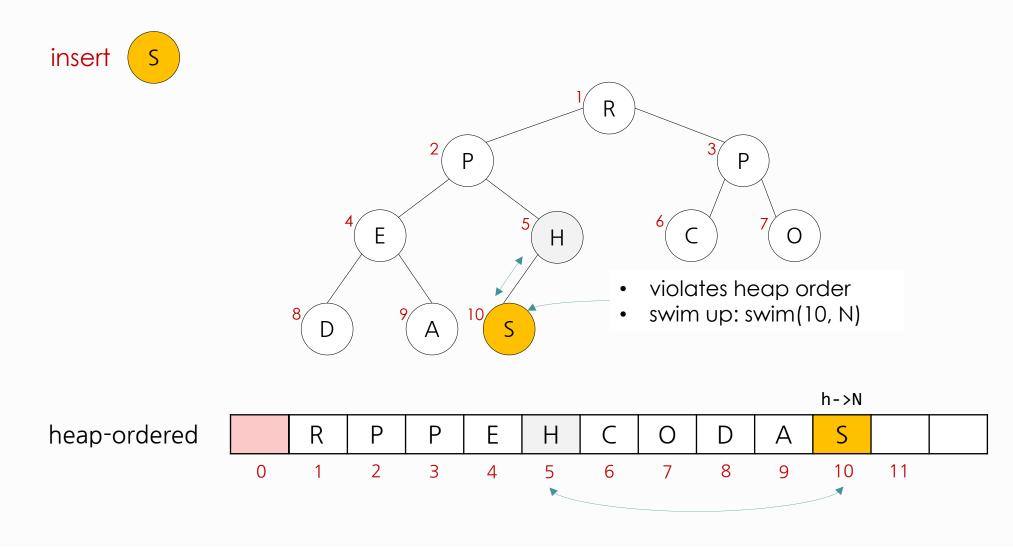
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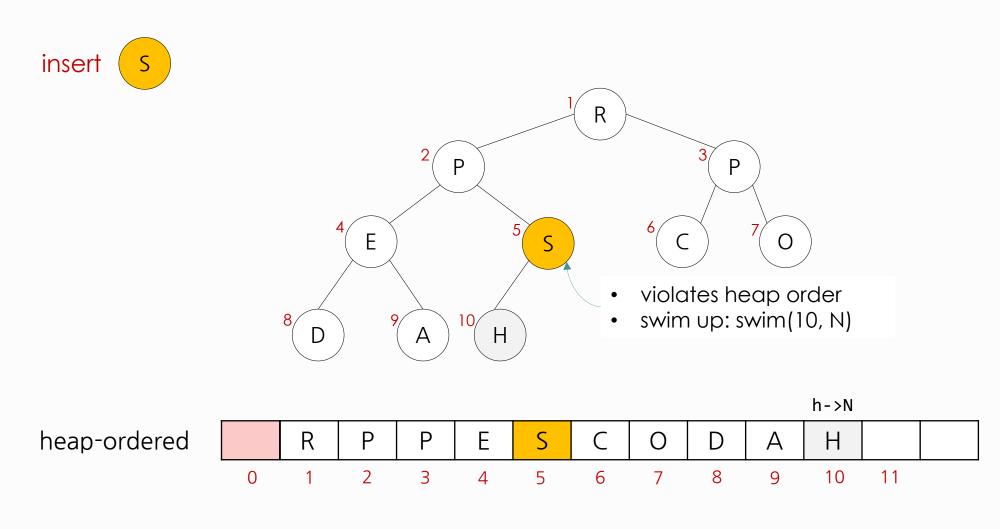
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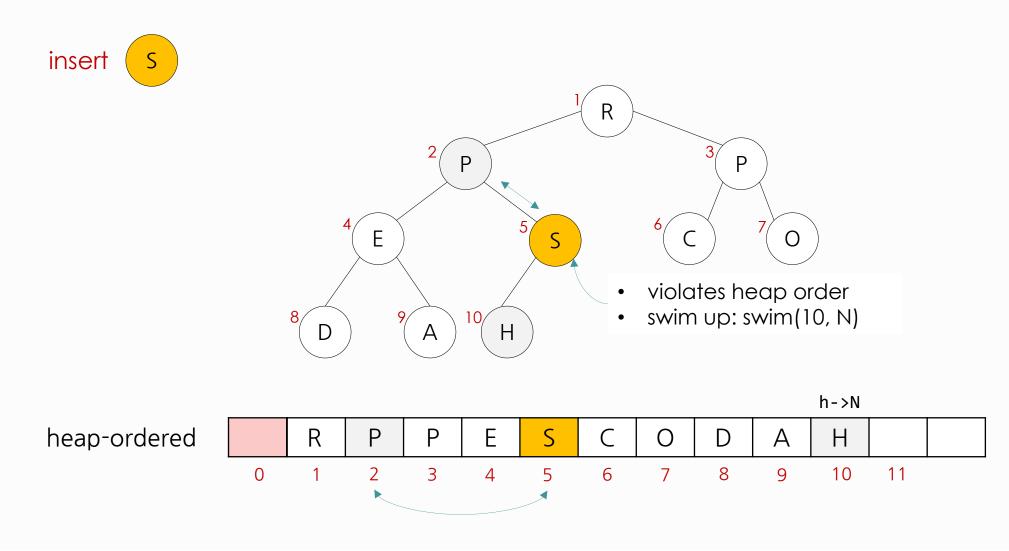
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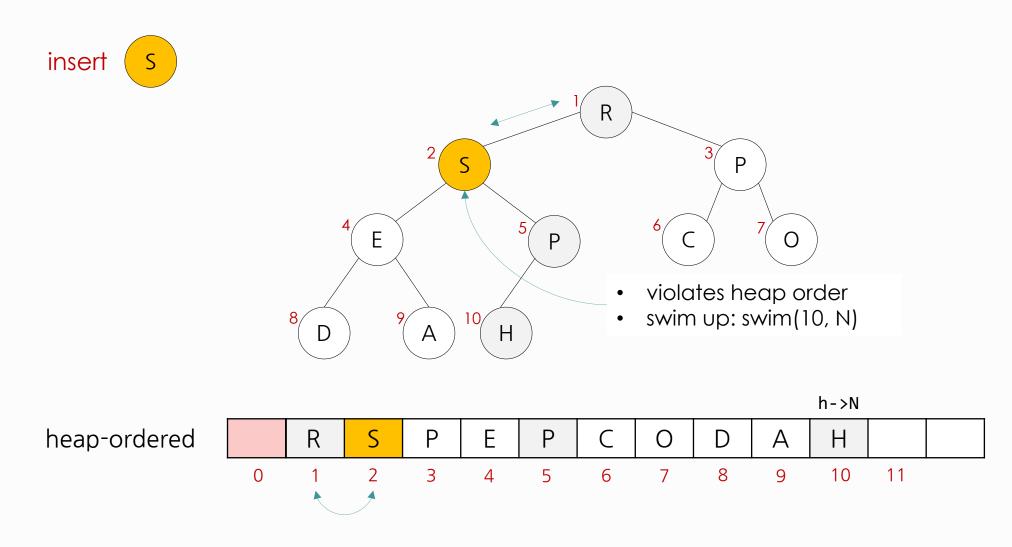
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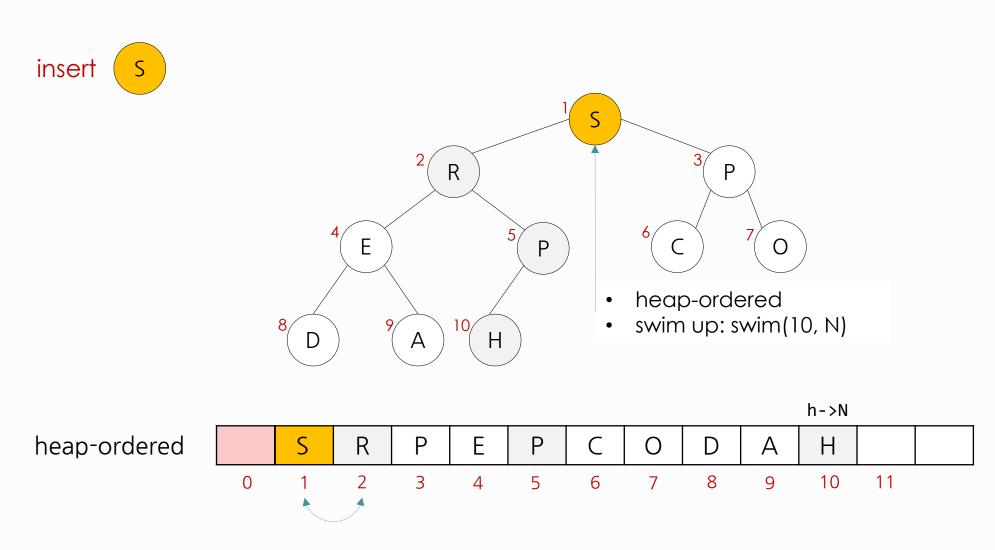
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Binary heap operations time complexity with N items:

- Level of heap is $\lfloor \log_2 N \rfloor$
- insert: O(log N) for each insert
 - In practice, expect less
- delete: O(log N) // deleting root node or any node
- increase/decrease key: O(log N)
- Heapify(): O(N)
- Heapsort(): O(N log N)
- Because O(N) heapify + O(log N) delete = O(N log N)

Proof:

- https://stackoverflow.com/questions/9755721/how-can-building-a-heap-be-on-time-complexity
- https://www.insertingwiththeweb.com/data-structures/binary-heap/build-heap-proof/
- https://www.guora.com/How-is-the-time-complexity-of-building-a-heap-is-o-n

References in Korean:

- https://ratsgo.github.io/data%20structure&algorithm/2017/09/27/heapsort/
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