

# 학습 목표

Big-O 표기법이 무엇인지 알고 직접 계산할 수 있다



# **Data Structures in Python Chapter 2 - 2**

- Performance Analysis
- Big-O Notation
- Big-O Properties
- Growth Rates
- Growth Rates Examples



그러므로 나의 사랑하는 자들아 너희가 나 있을 때 뿐 아니라 더욱 지금 나 없을 때에도 항상 복종하여 두렵고 떨림으로 너희 구원을 이루라 (Continue to work out your salvation with fear and trembling.) 빌2:12

나는 인애를 원하고 제사를 원하지 아니하며 번제보다 하나님을 아는 것을 원하노라 (호6:6) 하나님은 모든 사람이 구원을 받으며 진리를 아는데에 이르기를 원하시느니라 (딤전2:4)

그런즉 너희가 먹든지 마시든지 무엇을 하든지 다 하나님의 영광을 위하여 하라 (고전10:31)

# Agenda & Reading

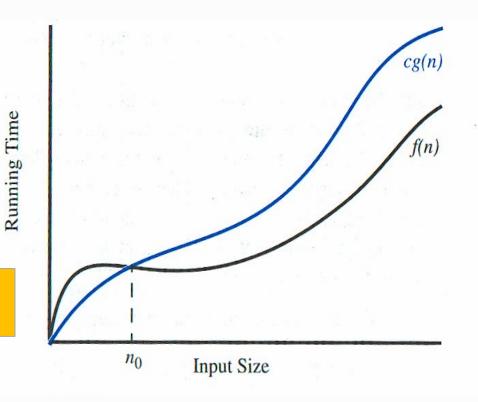
- Big-O Notation
  - Asymptotic Analysis
- Big-O Properties
  - Calculating Big-O

- References:
  - Textbook: Problem Solving with Algorithms and Data Structures
    - Chapter 3. <u>Analysis</u>
  - Textbook: <u>www.github.idebtor/DSpy</u>
    - Chapter 2.1 ~ 3

#### 3 Big-O Definition

- Let f(n) and g(n) be functions that map non-negative integers to real numbers. We say that f(n) is O(g(n)) if there is a real constant c, where c > 0 and an integer constant n, where  $n_0 \ge 1$  such that  $f(n) \le c * g(n)$  for every integer  $n \ge n_0$ .
  - f(n) describe the actual time of the program
  - g(n) is a much simpler function than f(n)
  - With assumptions and approximations, we can use g(n) to describe the complexity i.e., O(g(n))

Big-O Notation is a mathematical formula that best describes an algorithm's performance.

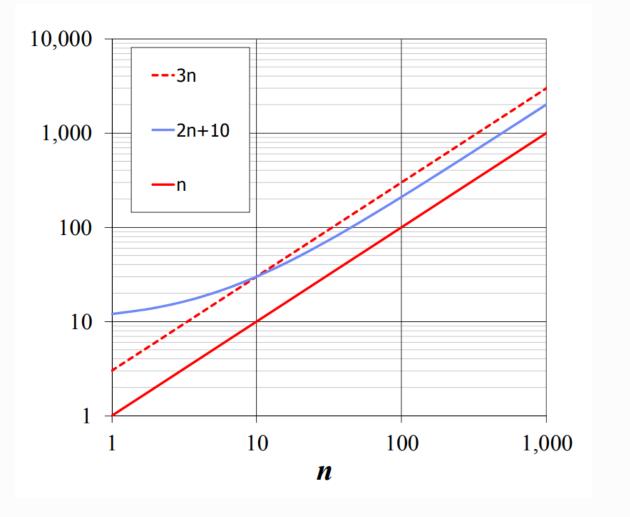


#### 3 Big-O Notation

- We use Big-O notation (capital letter O) to specify the order of complexity of an algorithm.
  - e.g.,  $O(n^2)$ ,  $O(n^3)$ , O(n)
  - If a problem of size n requires time that is directly proportional to n, the problem is O(n) that is, order n.
  - If the time requirement is directly proportional to  $n^2$ , the problem is  $O(n^2)$ , etc.

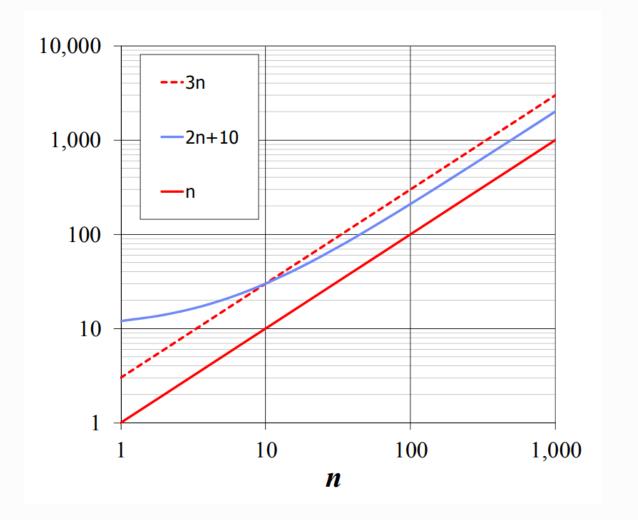
Given functions f(n) and g(n), we say that f(n) is O(g(n)) if there are positive constants, c, and  $n_0$  such that  $f(n) \le c * g(n)$  for every integer  $n \ge n_0$ .

- Example: T(n) = 2n + 10 T(n) is O(n)
- Question:



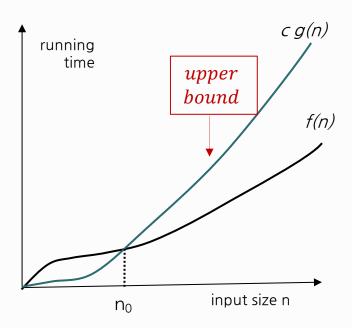
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- Example: T(n) = 2n + 10 T(n) is O(n)
- Question:
  - $n_0$
  - (
  - g(n)
  - $f(n) \leq c * g(n)$
  - f(n) is O(g(n))



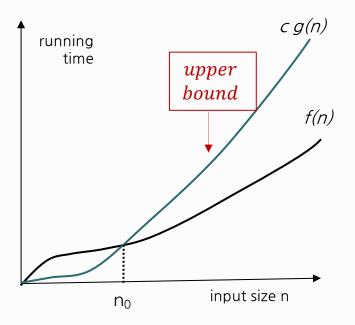
• Find c and  $n_0$  to justify that the function 7n + 5 is O(n).

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We must find c and n_0 such that 7n + 5 \le c n for <math>n \ge n_0
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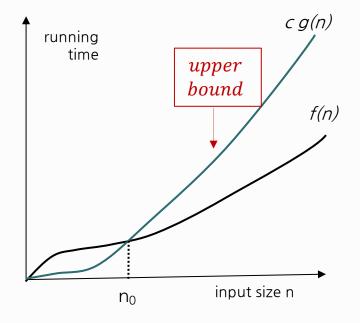
• Find c and  $n_0$  to justify that the function 7n + 5 is O(n).

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We must find c and n_0 such that 7n + 5 \le c n for n \ge n_0 7n + 5 \le 7 n + n 7n + 5 \le 8 n for n \ge n_0 = 5 Therefore, 7n + 5 \le c n for c = 8 and n_0 = 5, g(n) = n and O(n)
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• Find c and  $n_0$  to justify that the function 7n + 5 is O(n).

```
We must find c and n_0 such that 7n + 5 \le c n \qquad for n \ge n_07n + 5 \le 7n + n7n + 5 \le 8n \qquad for n \ge n_0 = 5Therefore, 7n + 5 \le c n for c = 8 and n_0 = 5, f(n) is O(n)
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7n + 5 \le c n  for n \ge n_0

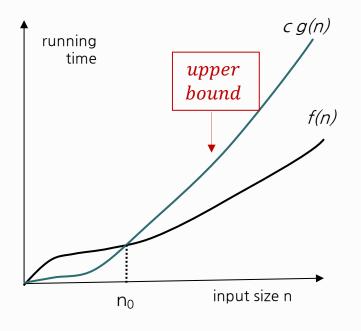
7n + 5 \le 12 n  for n \ge n_0 = 1

Therefore, 7n + 5 \le c n for c = 12 and n_0 = 1

g(n) = n, f(n) is O(n)
```

• Find  $\boldsymbol{c}$  and  $\boldsymbol{n_0}$  to justify that the function  $\boldsymbol{f}(\boldsymbol{n}) = 27\boldsymbol{n}^2 + 16\boldsymbol{n}$  is  $\boldsymbol{O}(\boldsymbol{n}^2)$ .

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We must find c and n_0 such that For 16n \le n^2 27n^2 + 16n \le 27n^2 + n^2 27n^2 + 16n \le 28n^2 \qquad for \ n \ge n_0 = 16 Hence, c = 28 and n_0 = 16, Therefore, g(n) = n^2, f(n) is O(n^2).
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27n^2+16n is \textbf{\textit{O}}(n^2), we must find \textbf{\textit{c}} and \textbf{\textit{n}}_0 such that 27n^2+16n \leq 43n^2 27n^2+16n \leq 43n^2 for n \geq \textbf{\textit{n}}_0=1 Hence, c=43 and \textbf{\textit{n}}_0=1, Therefore, \textbf{\textit{g}}(\textbf{\textit{n}})=\textbf{\textit{n}}^2, \textbf{\textit{f}}(\textbf{\textit{n}}) is \textbf{\textit{O}}(n^2).
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- Suppose an algorithm requires
  - T(n) = 7n-2 operations to solve a problem of size n

$$7n-2 \le 7 * n \text{ for all } n_0 \ge 1$$
  
i.e.,  $c = 7$ ,  $n_0 = 1$ 

•  $T(n) = n^2 - 3 * n + 10$  operations to solve a problem of size n

$$n^2 - 3 * n + 10 < 3 * n^2$$
 for all  $n_0 \ge 2$   
i.e.,  $c = 3$ ,  $n_0 = 2$ 

•  $T(n) = 3n^3 + 20n^2 + 5$  operations to solve a problem of size n

$$3n^3 + 20n^2 + 5 < 4 * n^3$$
 for all  $n_0 \ge 21$  i.e.,  $c = 4$ ,  $n_0 = 21$   $O(n^3)$ 

 $f(n) \le c * g(n)$  for every integer  $n \ge n_0$ 

- 1) 3n + 2 =
- 2) 3n + 3 =
- 3) 100n + 6 =
- 4)  $10n^2 + 4n + 2 =$
- 5)  $6 * 2^n + n^2 =$
- 6) 3n + 3 =
- 7)  $10n^2 + 4n + 2 =$
- (3)  $3n + 2 \neq 0$  (1) as 3n + 2 is **not**  $\leq c$  for any c and all  $n, n \geq n_0$ .
- (3) 9)  $10n^2 + 4n + 2 \neq O(n)$

#### Summary

- Big-O Notation is a mathematical formula that best describes an algorithm's performance.
- Big-O notation is often called the asymptotic notation (점근적 표기법) since it uses so-called the asymptotic analysis (점근적 분석) approach.
- Normally we assume worst-case analysis, unless told otherwise.
- In some cases, it may need to consider the best, worst and/or average performance of an algorithm

# 학습 정리

1) Big-O(빅 오)은 알고리즘의 수행능력을 잘 나타내는 수학적인 표기법이다

2) Big-O를 계산할 때 주어진 함수들에서 가장 근접한 함수를 찾는 것이 좋다

