

# 파이썬으로 배우는 데이터 구조



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# 학습 목표

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트리(Tree) 구조와 용어에 대해 학습하고  
간단하게 구현할 수 있다

# **Data Structures in Python**

## **Chapter 7 - 1**

- Tree Introduction
- Tree Traversals
- Tree Algorithms

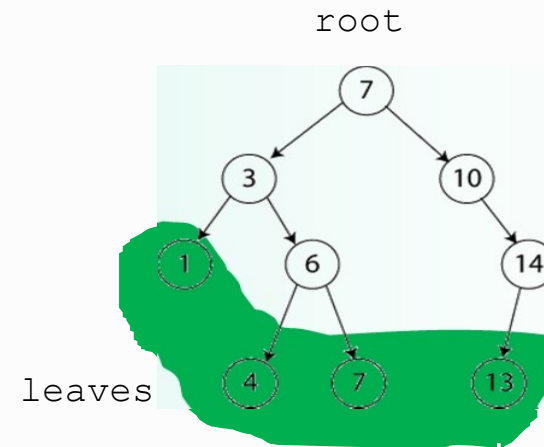
# Agenda & Readings

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- Agenda
  - Tree Terminology
  - Binary Tree Properties
  - Binary Tree and Node Representation
- Reference:
  - Problem Solving with Algorithms and Data Structures
  - Chapter 6 - Tree

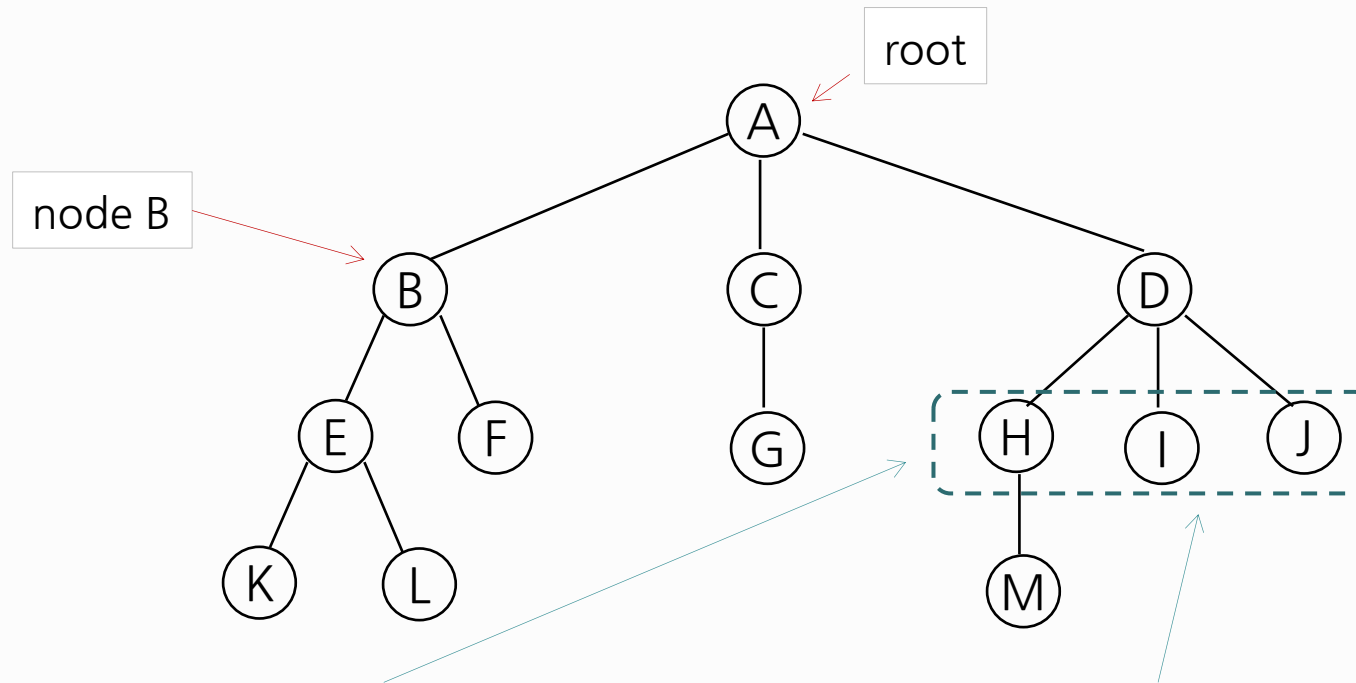
# What is a Tree?

- A non-linear data structure
- An abstraction for a hierarchical structure
- It is defined as a set of points called nodes and a set of lines called edges where an edge connects two distinct nodes.



# Introduction - Terminology

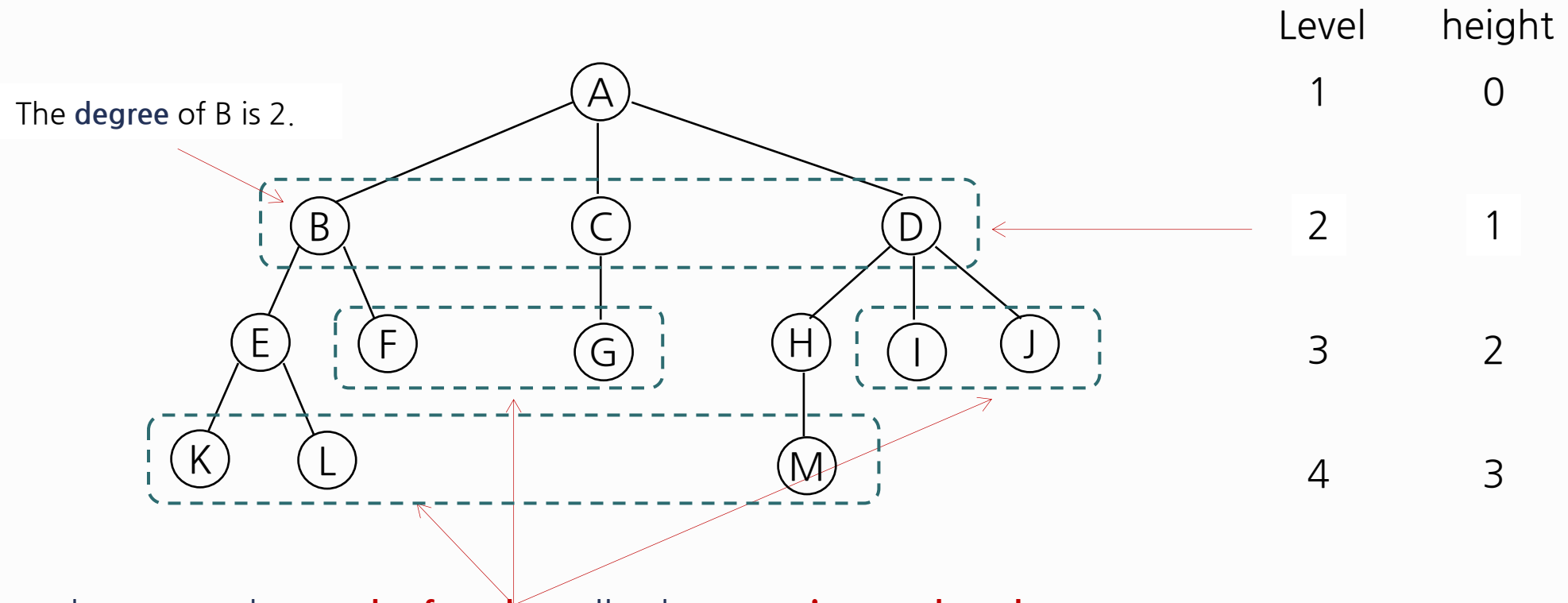
- **A tree data structure:** it is like a linked list that has a **first** node, this node is called as the **root** of the tree.
- **Example.** A **tree** with a root storing the value 'A'



- The **children** of D are H, I, and J; H, I, and J are **siblings**.
- The **parent** of D is A.

# Introduction - Terminology

- **Definition.** child, parent, sibling, degree, leaf nodes, level, and internal node



- Zero degree nodes are **leaf nodes**, all others are **internal nodes**.
  - An **internal node** is any node that has at least one non-empty child.
- The **degree** of a node is the number of children.
- The **degree of a tree** is the **maximum of the degree of the nodes** in the tree.

# Introduction - Representation of trees

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- **Exercise.** The tree representing the HTML document below:

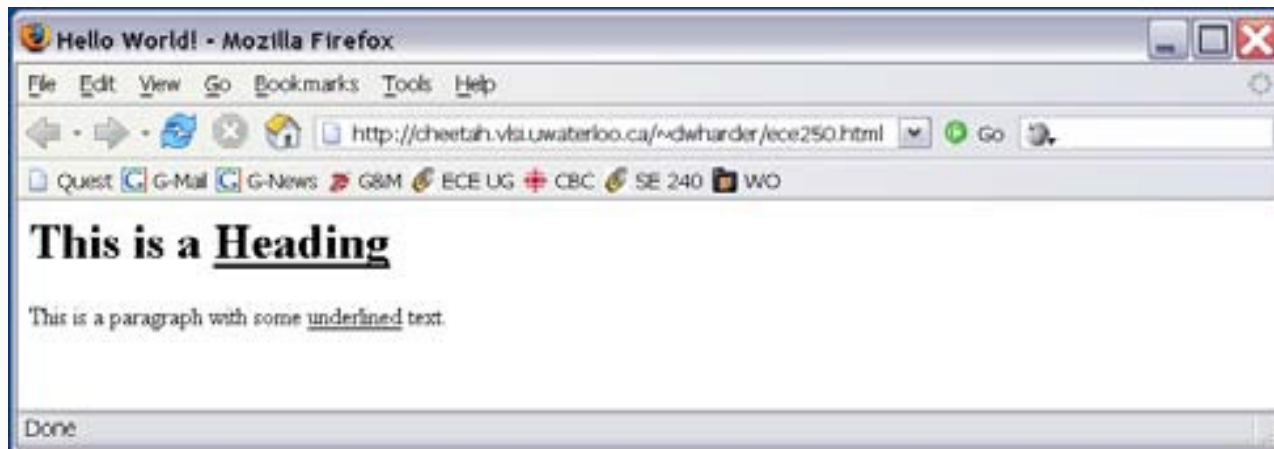
```
<html>
  <head>
    <title>Hello World!</title>
  </head>
  <body>
    <h1>This is a <u>Heading</u></h1>
    <p>This is a paragraph with some <u>underlined</u> text.</p>
  </body>
</html>
```



# Introduction - Representation of trees

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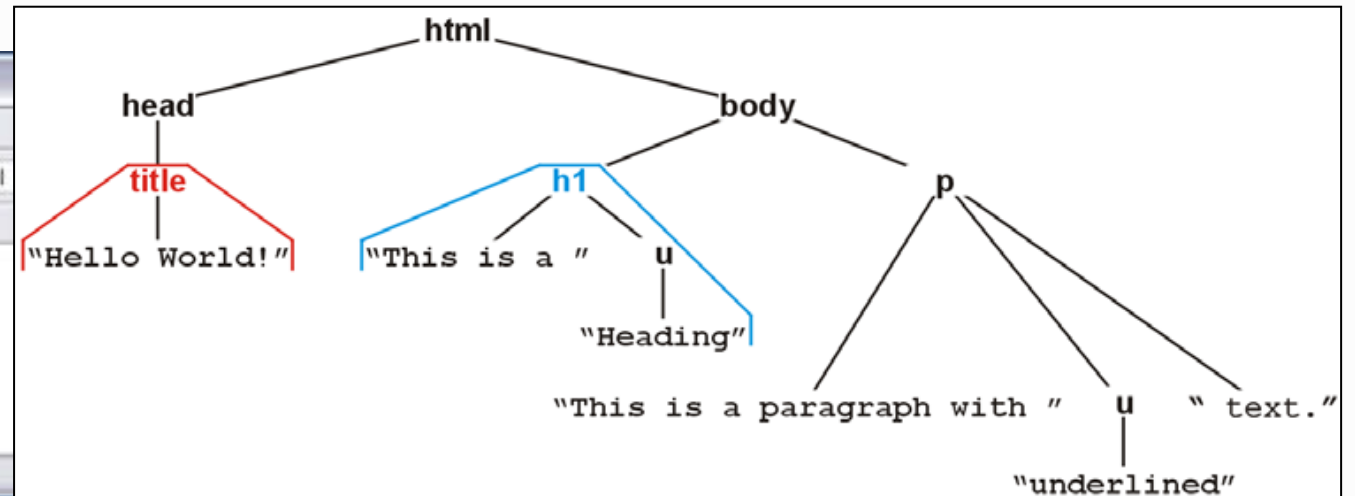
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# Introduction - Representation of trees

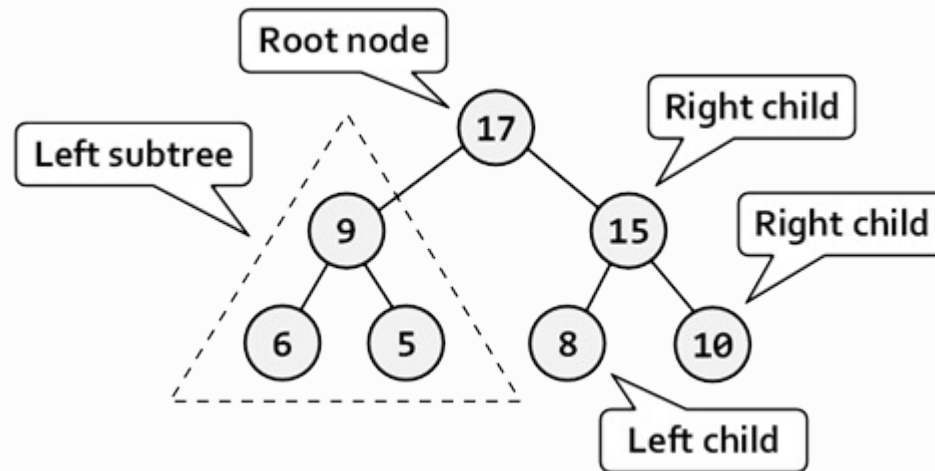
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```



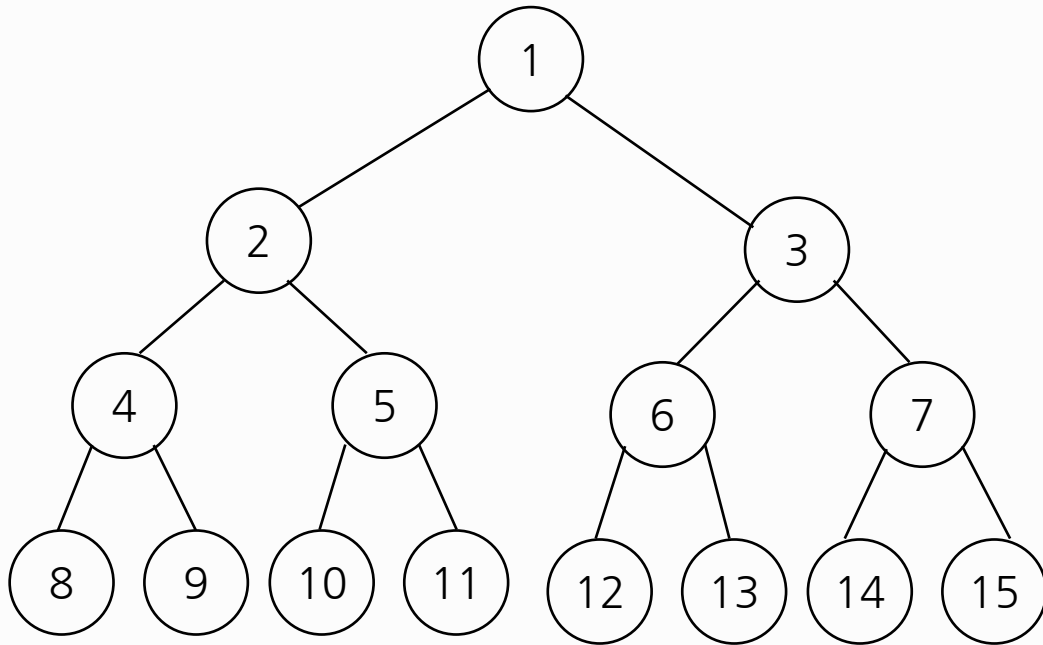
# Binary trees

- **Definition:** A tree such that each node has *exactly* two children.
  - Notice, exactly two children - not up to two children! Because *exactly* two children means a left child **and/or** right child, no middle child.
  - Each child is either empty or another binary tree.
  - Given this constraint, we can label the two children as left and right nodes or subtrees.



# Binary trees - Properties

- **Observation:**
  - Q: Maximum number of nodes in binary trees in each level and all levels?
  - Q: What is the max level  $k$  if there are  $n$  nodes?  $k(n) = ?$

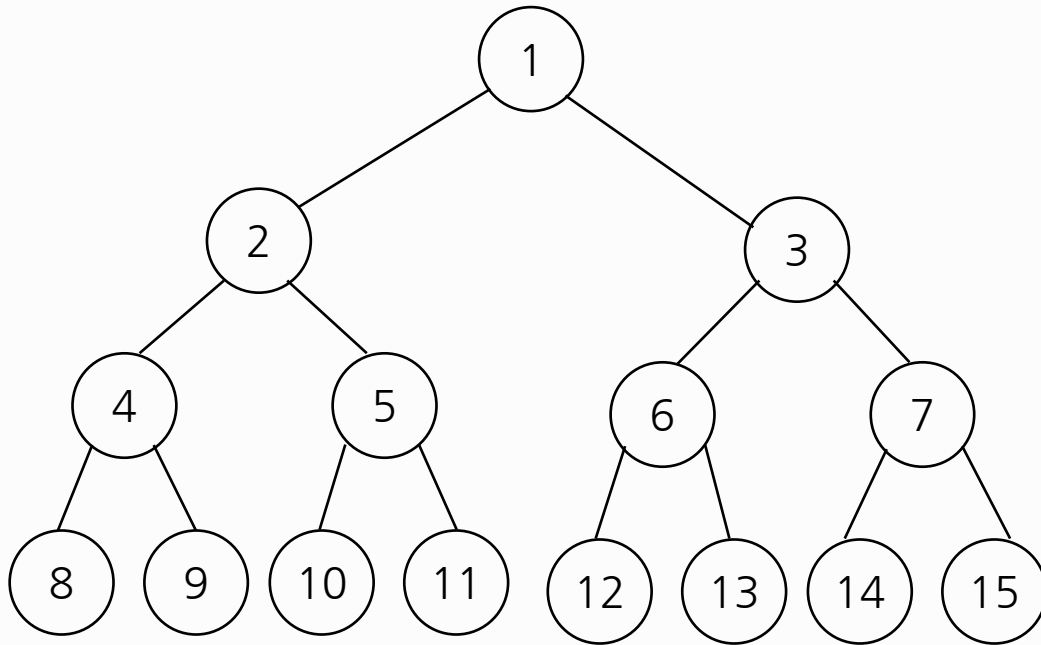


*A full binary tree*

# Binary trees - Properties

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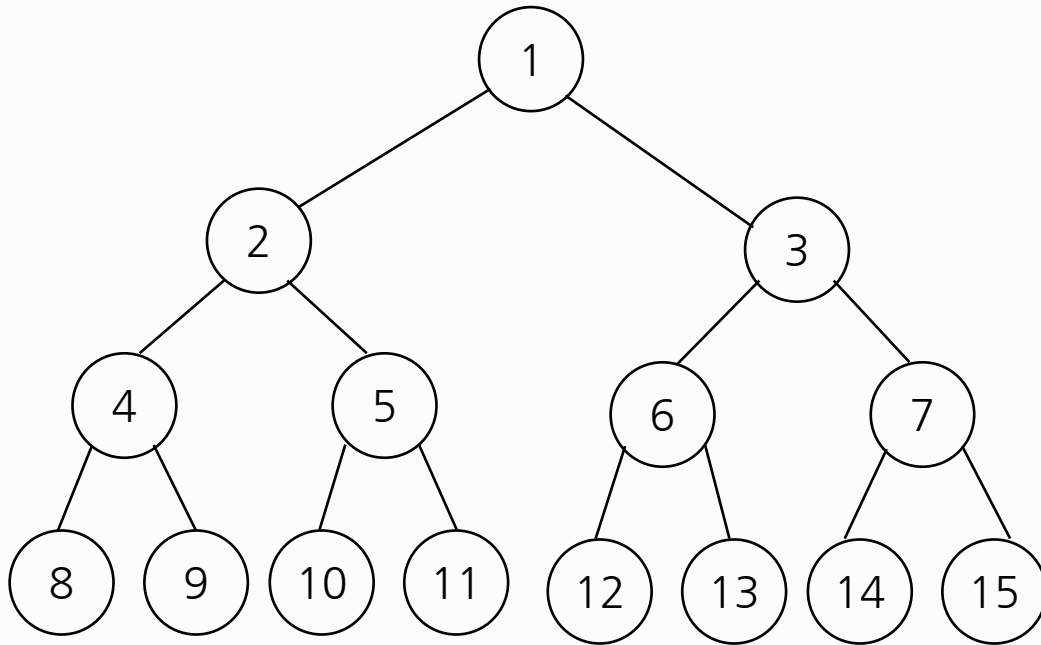
*A full binary tree*

Level	Node Numbers at Each Level	Total Numbers of Nodes
1	$1 = 2^0$	
2	$2 = 2^1$	
3	$4 = 2^2$	
4	$8 = 2^3$	
.	.	
11	$1024 = 2^{10}$	
.	.	
k		

# Binary trees - Properties

## ■ Observation:

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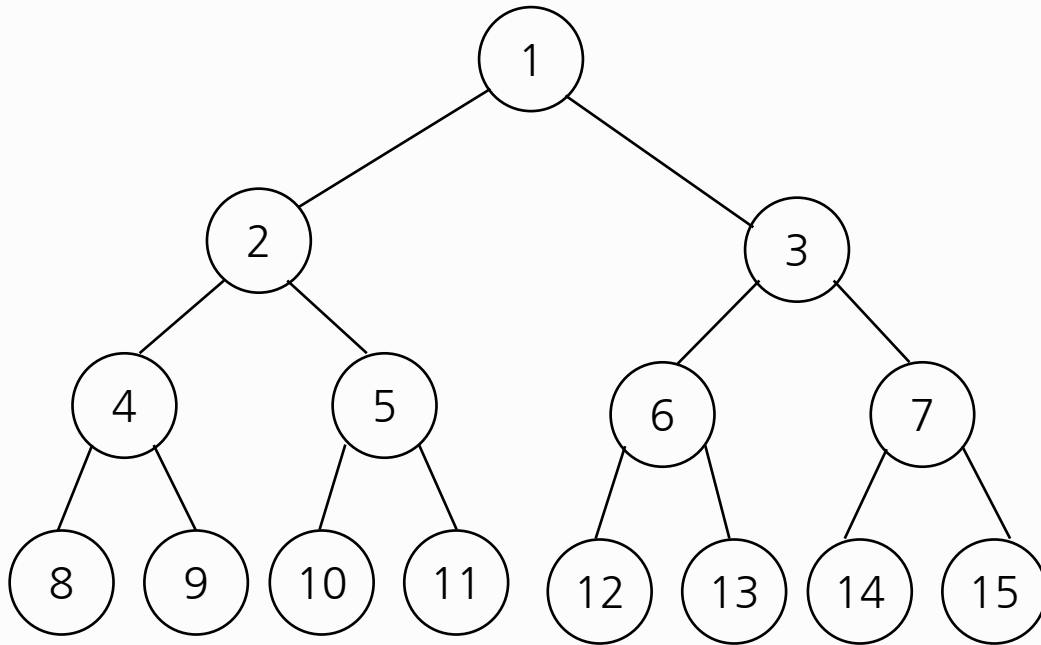
*A full binary tree*

Level	Node Numbers at Each Level	Total Numbers of Nodes
1	$1 = 2^0$	$1 = 2^1 - 1$
2	$2 = 2^1$	$3 = 2^2 - 1$
3	$4 = 2^2$	$7 = 2^3 - 1$
4	$8 = 2^3$	$15 = 2^4 - 1$
.	.	.
11	$1024 = 2^{10}$	$2047 = 2^{11} - 1$
.	.	.
k		

# Binary trees - Properties

## ■ Observation:

- Q: Maximum number of nodes in binary trees in each level and all levels?
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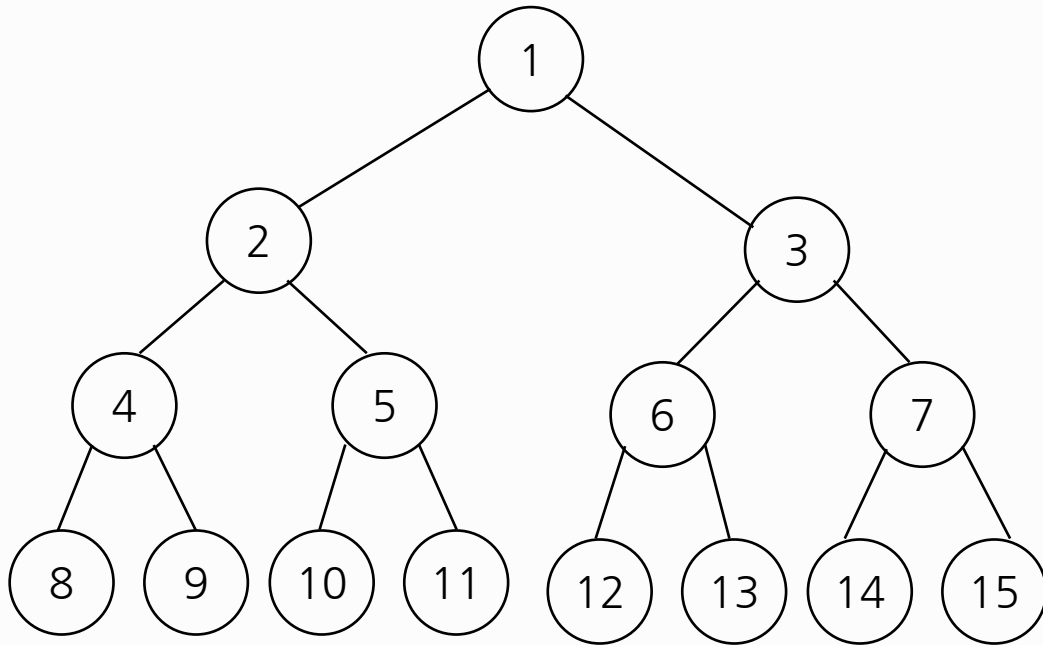


*A full binary tree*

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.	.	.
11	$1024 = 2^{10}$	$2047 = 2^{11} - 1$
.	.	.
k	$2^{k-1}$	$2^k - 1$
h	$2^h$	$2^{h+1} - 1$

# Binary trees - Properties

- **Definition:** *A full binary tree* of level  $k$  is a binary tree having  $2^k - 1$  nodes,  $k \geq 0$ .
- **Definition:** A binary tree with  $n$  nodes and level  $k$  is **complete** iff its nodes correspond to the nodes numbered from 1 to  $n$  in the full binary tree of level  $k$ .

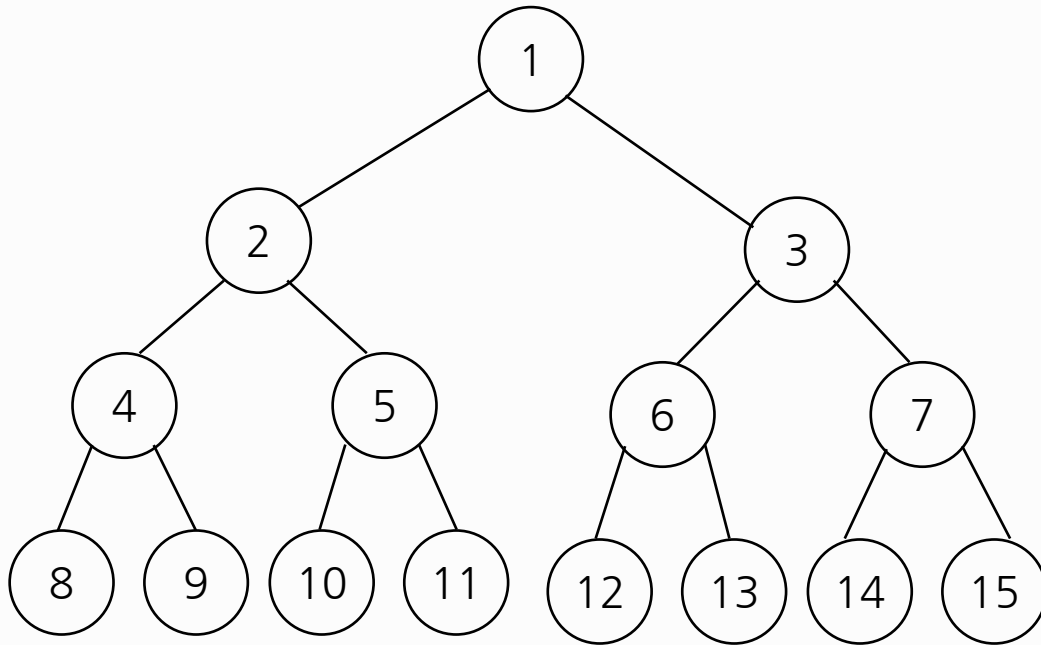


*A full binary tree*

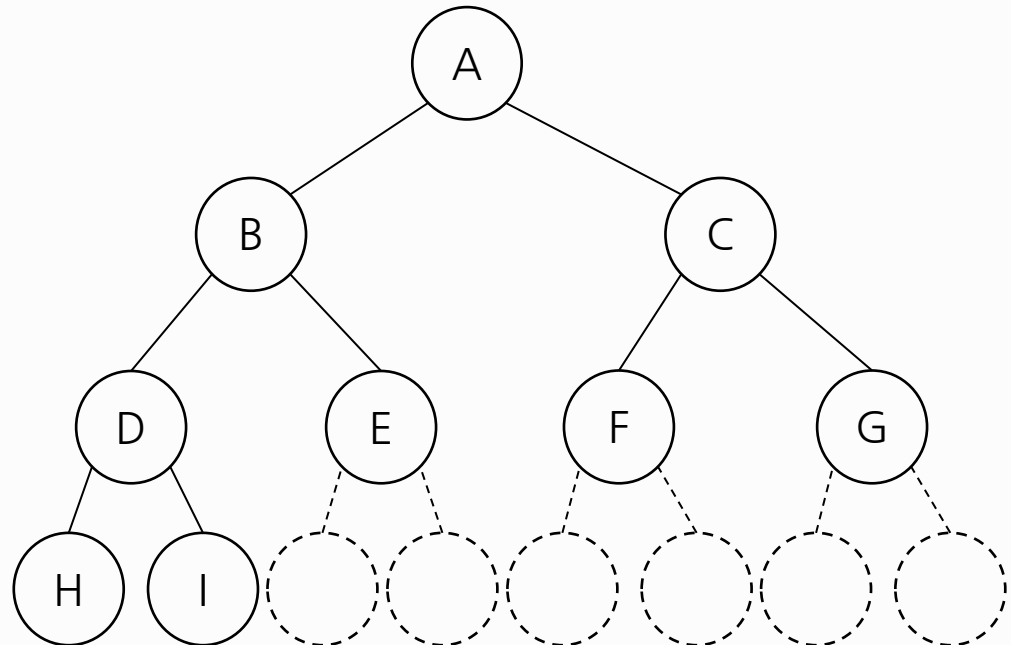


# Binary trees - Properties

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- **Definition:** A binary tree with  $n$  nodes and level  $k$  is **complete** iff its nodes correspond to the nodes numbered from 1 to  $n$  in the full binary tree of level  $k$ .



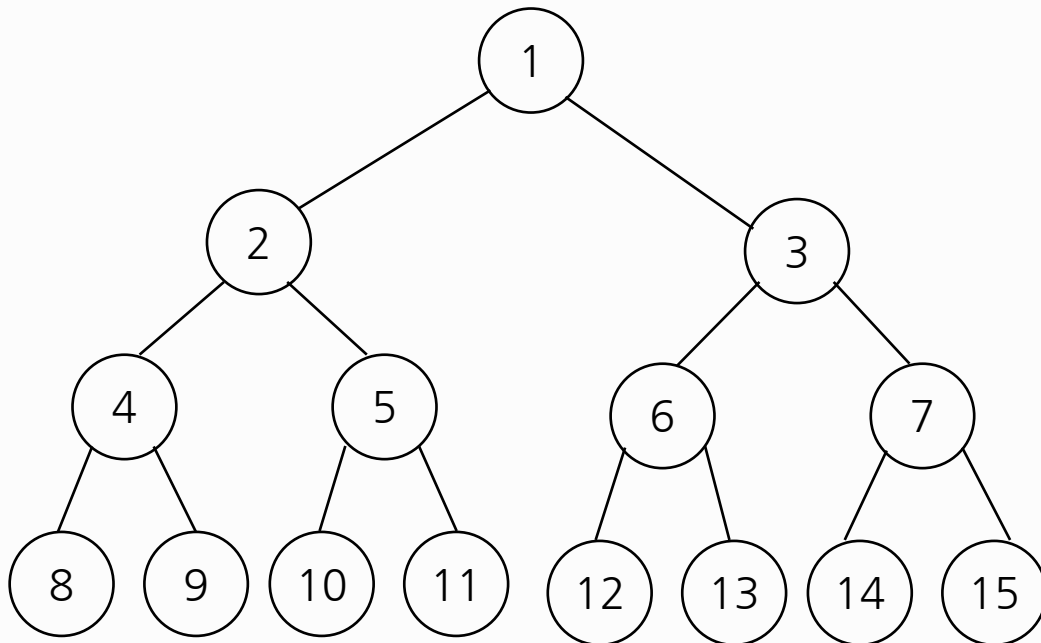
*A full binary tree*



*A complete binary tree*

# Binary trees - Array representation

- **Q:** Let's suppose that you have a **complete binary tree** in an array, how can we locate node  $x$ 's parent or child?
- A **complete** binary tree with  $n$  nodes, any node index  $i$ ,  $1 \leq i \leq n$ , we have
  - $\text{parent}(i)$  is at  $\lfloor i/2 \rfloor$  If  $i = 1$ ,  $i$  is at the root and has no parent
  - $\text{leftChild}(i)$  is at  $2i$  if  $2i \leq n$ . If  $2i > n$ , then  $i$  has no left child.
  - $\text{rightChild}(i)$  is at  $2i+1$  if  $2i+1 \leq n$ . If  $2i+1 > n$ , then  $i$  has no right child.



Wow! Can we use this to all binary trees?  
**Why not?**

## Problem remains:

The problem with storing an arbitrary binary tree using an array is the inefficiency *in memory usage*.

## Binary trees - Properties

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(1) The maximum number of **nodes on level k** of a binary tree is

$k \geq 1$

(2) The maximum number of **nodes in a binary tree of level k** is

$k \geq 1$

(3) The maximum level of a **complete binary tree** with **n** nodes is

$[x]$  is the smallest integer  $\geq x$ .

## Binary trees - Properties

---

(1) The maximum number of **nodes on level  $k$**  of a binary tree is

$$2^{k-1}, \quad k \geq 1$$

(2) The maximum number of **nodes in a binary tree of level  $k$**  is

$$2^k - 1, \quad k \geq 1$$

(3) The maximum level of a **complete binary tree** with  **$n$**  nodes is

$$k(n) = \lceil \log_2 (n + 1) \rceil, \quad \lceil x \rceil \text{ is the smallest integer } \geq x.$$

$$n = 2^k - 1$$

$$n + 1 = 2^k$$

$$\log(n + 1) = \log 2^k$$

$$\log(n + 1) = k$$

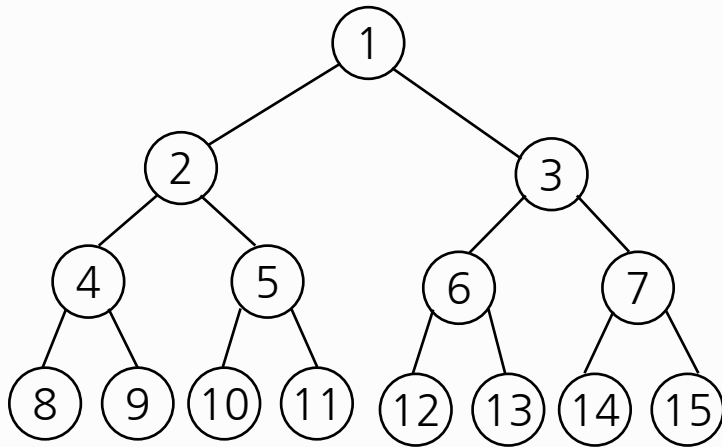
$$k(n) = \lceil \log(n + 1) \rceil$$

$$k(n) = \lceil \log(n) \rceil + 1$$

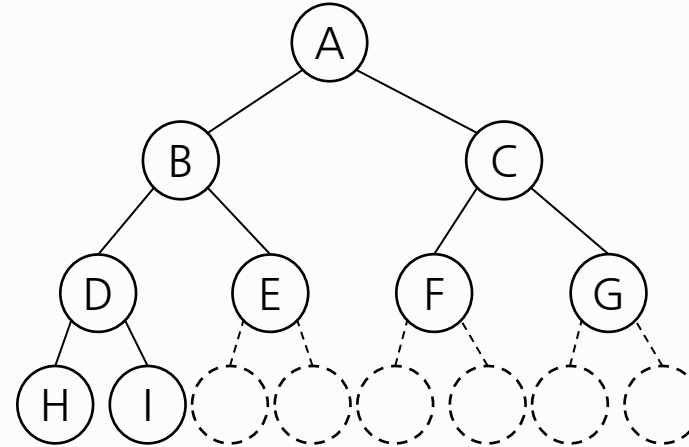
since  $k$  is an integer, and includes  
the max level of complete binary tree.

# Binary trees - Properties

- **Observation:** The max level of a full binary tree of  $n$  nodes is  $k = \text{floor}(\log_2(n)) + 1$  :
  - Many operations with trees have a run time that goes with the max level of some path within the tree;
  - If we have a full binary tree (or something *close* to it), we know that those operations **run in  $O(\log n)$** .



*A full binary tree*



*A complete binary tree*

# Binary trees - Linked representation

- **Node** and **Tree** representations:

```
class Node:
    def __init__(self, key):
        self.key = key
        self.left = None
        self.right = None

    def insertLeft(self, key):
        ...
```

```
root = Node('A')
print(root)
```

key	
left	right

root

'A'	
None	None

# Binary trees - Linked representation

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```

```
root = Node('A')
print(root)
```

```
<__main__.Node object at
0x0000015985006A00>
```

key	
left	right

6A00

root

'A'	
None	None

# Binary trees - Linked representation

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```
class Node:
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    def insertLeft(self, key):
        ...
```

- **Create 'B' and link with left of 'A'**

```
root = Node('A')
node = Node('B')
```

key	
left	right

6A00  
root

'A'	
None	None

7A00  
node

'B'	
None	None



# Binary trees - Linked representation

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```

- Create 'B' and link with left of 'A'.

```
root = Node('A')
node = Node('B')
root.left = node
```

key	
left	right

6A00  
root

'A'	
None	None

7A00  
node

'B'	
None	None

# Binary trees - Linked representation

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```
root = Node('A')
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root.left = node
```

- Simplify the code above.

key	
left	right

6A00  
root

'A'	
7A00	None

7A00  
node

'B'	
None	None

# Binary trees - Linked representation

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```
class Node:
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        ...
```

- Create 'B' and link with left of 'A'.

```
root = Node('A')
node = Node('B')
root.left = node
```

- Simplify the code and diagram.

```
root = Node('A')
root.left = Node('B')
```

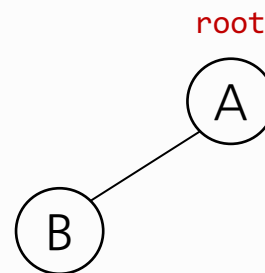
key	
left	right

6A00  
root

'A'	
7A00	None

7A00  
node

'B'	
None	None



# Binary trees - Linked representation

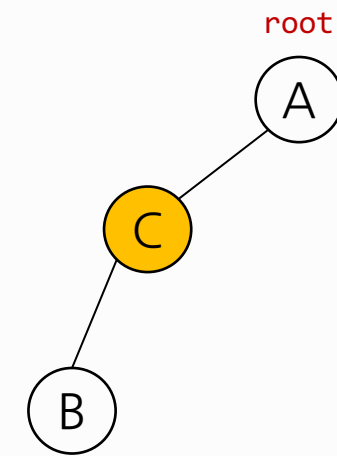
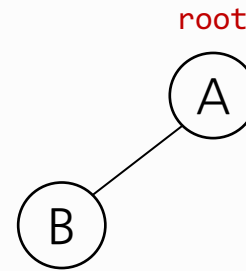
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- Code **insertLeft()** and insert 'B' & 'C'.

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root = Node('A')
root.insertLeft('B')
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# Binary trees - Linked representation

- **Node** and **Tree** representations:

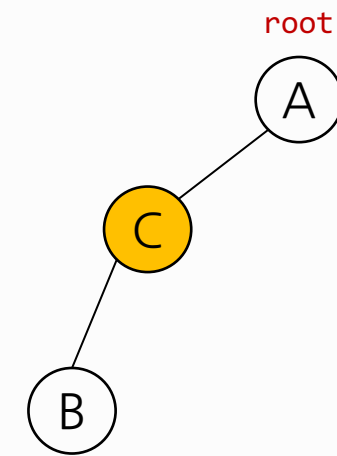
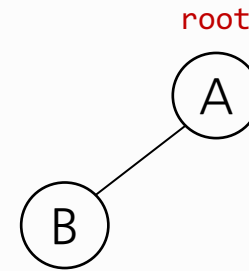
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```
def insertLeft(self, key):
    if self.left == None:
        self.left = Node(key)
    else:
        ...
```



# Binary trees - Linked representation

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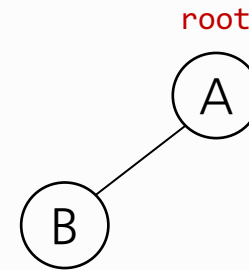
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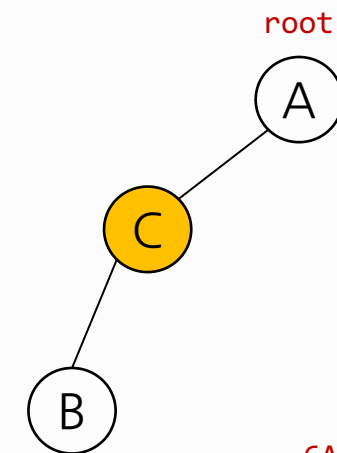


6A00  
root

'A'	
7A00	None

7A00

'B'	
None	None



6A00  
root

'A'	
7CFF	None

7CFF

'C'	
7A00	None

7A00

'B'	
None	None

# Binary trees - Linked representation

- **Node** and **Tree** representations:

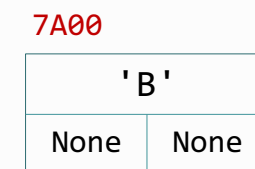
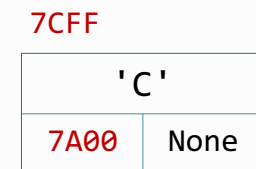
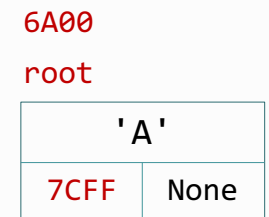
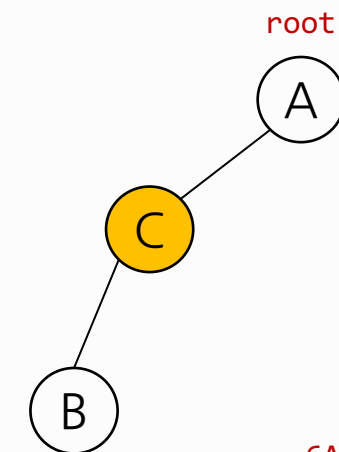
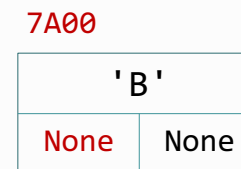
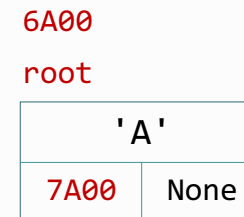
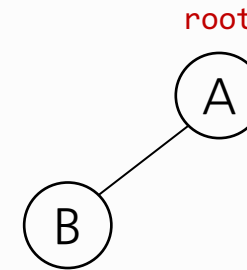
```
class Node:
    def __init__(self, key):
        self.key = key
        self.left = None
        self.right = None

    def insertLeft(self, key):
        ...
```

- Code insertLeft() and insert 'B' & 'C'.

```
root = Node('A')
root.insertLeft('B')
root.insertLeft('C')
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```
def insertLeft(self, key):
    if self.left == None:
        self.left = Node(key)
    else:
        node = Node(key)
        node.left = self.left
        self.left = node
```

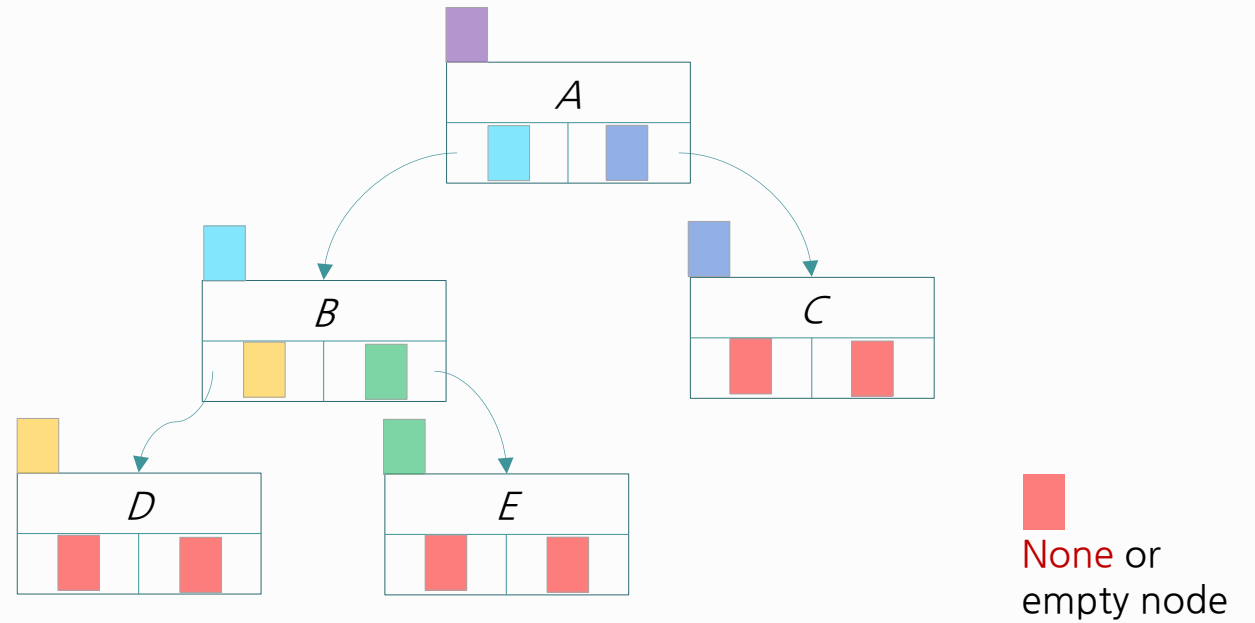


# Binary trees - Linked representation

- **Node** and **Tree** representations:

```
class Node:
    def __init__(self, key):
        self.key = key
        self.left = None
        self.right = None

    def insertLeft(self, key):
        ...
```



- Q. Is this node structure good enough?
  - Not easy to find its parent node. A parent field could be added if necessary.

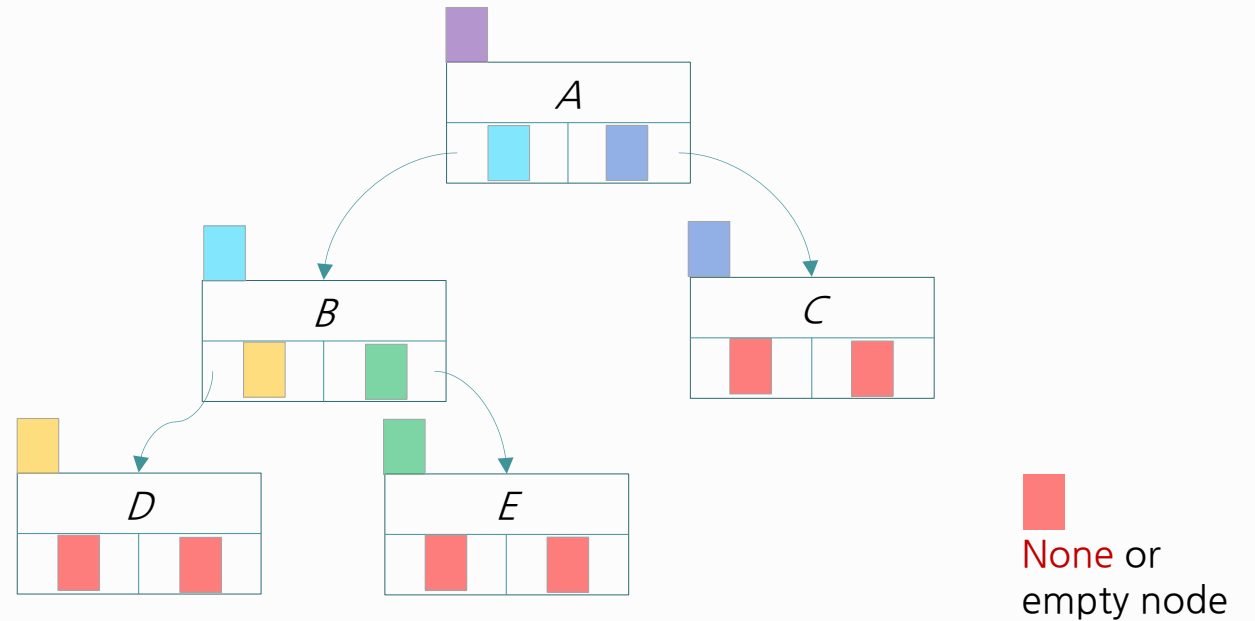


# Binary trees - Linked representation

- **Node** and **Tree** representations:

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class Node:
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        self.right = None

    def insertLeft(self, key):
        ...
```



- Q. Is this node structure good enough?
  - Not easy to find its parent node. A parent field could be added if necessary.
- Q. It is similar to a doubly-linked list(DLL). What is different?
  - One head, but many tails. **None** points empty node conceptually.

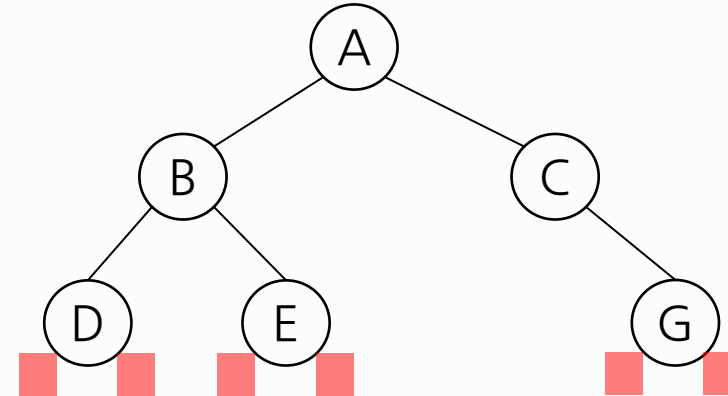
# Binary trees - Exercise 1

- Build a tree shown below using insertLeft() and insertRight().

```
class Node:
    def __init__(self, key):
        self.key = key
        self.left = None
        self.right = None

    def insertLeft(self, key):
        ...
    def insertRight(self, key):
        ...

if __name__ == '__main__':
    root = Node('A')
    root.insertLeft('B')
    root.insertRight('C')
    # your code here
```



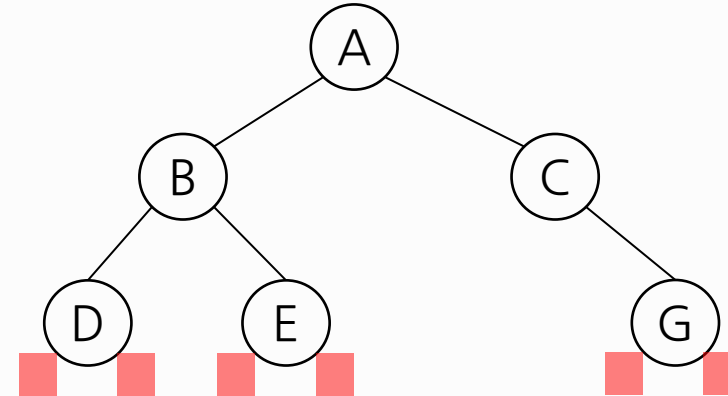
## Binary trees - Exercise 2

- Extend Exercise 1 such that it exactly reproduces the output using the as shown below.

```
class Node:
    def __init__(self, key):
        self.key = key
        self.left = None
        self.right = None

    def insertLeft(self, key):
        ...
    def insertRight(self, key):
        ...

if __name__ == '__main__':
    ...
    for node in [a, b, c, d, e, g]:
        if node.key:
            print(...)
    ...
```



```
A: (B, C)
B: (D, E)
C: (None, G)
D: (None, None)
E: (None, None)
G: (None, None)
```

# 학습 정리

- 1) 트리(Tree)는 비선형 자료 구조이다
- 2) 완전이진트리(Complete binary tree)에 대한 시간복잡도는  $O(\log n)$ 으로 빠른 알고리즘이다

# 파이썬으로 배우는 데이터 구조

수고했습니다  
곧 다음 시간에  
다시 뵙겠습니다

