

학습 목표

이진탐색 트리(BST)의 ADT를 학습하고 이에 맞는 Node클래스를 구현할 수 있다



Data Structures in Python Chapter 7 - 2

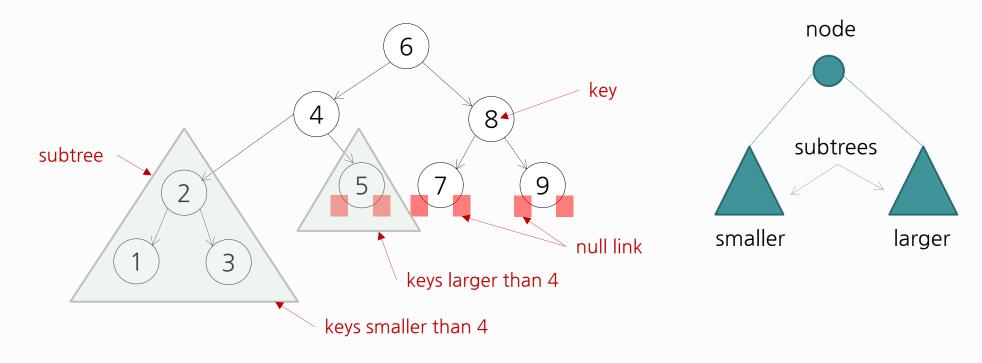
- Binary Search Tree(BST)
- BST Algorithms
- AVL Tree
- AVL Algorithms

Agenda & Readings

- Binary Search Tree
 - Binary Search Tree Properties
 - Binary Search Tree ADT
 - Node Class BST Class Constructors
 - Search & Insert
- Reference:
 - Problem Solving with Algorithms and Data Structures
 - Chapter 6 Tree

Binary Search Tree: Definition

- A binary tree (BT) is a tree data structure in which each node has at most two children, which are referred to as the left child and the right child.
- A binary search tree (BST) is an ordered or sorted binary tree.
 - Each node in a binary search tree has a comparable key (and an associated value) and satisfies the restriction that the key in any node is larger than the keys in all nodes in that node's left subtree and smaller than the keys in all nodes in that node's right subtree.)
 - Equal keys are ruled out.



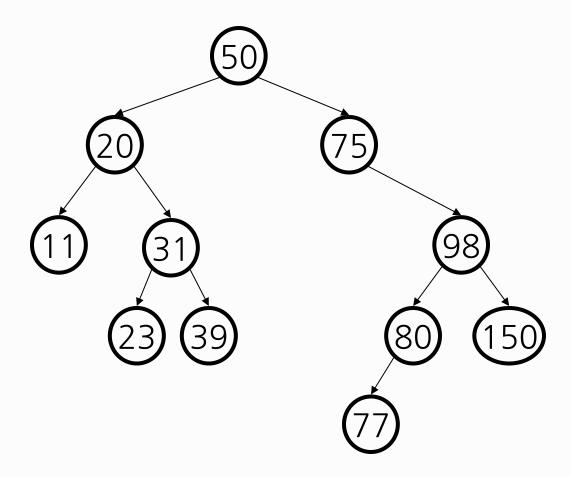
Binary Search Tree: Example

• Insert - Drawing a binary search tree as the following values are being added to an initially empty tree.

www.mathwarehouse.com

Binary Search Tree: Exercise

• Insert - Draw what a binary search tree would look like if the following values were added to an initially empty tree in this order:





Binary Search Tree: ADT

- get() search, contains, size, height, min/max, successor, predecessor, distance
- add() insert, grow
- delete() delete, trim
- BT to BST conversion
- LCA Lowest Common Ancestor

Binary Search Tree:

Node and BST Class Constructors

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Node and BST Class Constructors

```
<__main__.Node object at 0x000002BCE8D22A00>
```

Binary Search Tree:

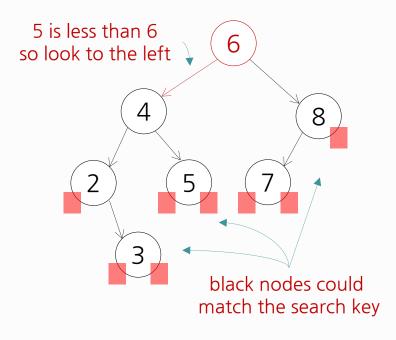
Node and BST Class Constructors

```
class BST:
    def __init__(self, key=None, left=None, right=None):
       if key is None:
           self.root = None
       else:
            self.root = Node(key, left, right)
    def delTree(self):
        self.root = None # gc will do this for us.
    def str (self):
        return self.root.__str__()
    def __iter__(self):
        return self.root. iter ()
    def __len__(self):
        return self.root.__len__()
    def len(self):
        return self. len ()
```

Binary Search Tree: Search

- A recursive algorithm to search for a key in a BST follows immediately from the recursive structure: If the tree is empty, we have a search miss; if the search key is equal to the key at the root, we have a search hit. Otherwise, we search (recursively) in the appropriate subtree.
 - get(root, key), contains(root, key), __getitem__(), __contains__(), contains()

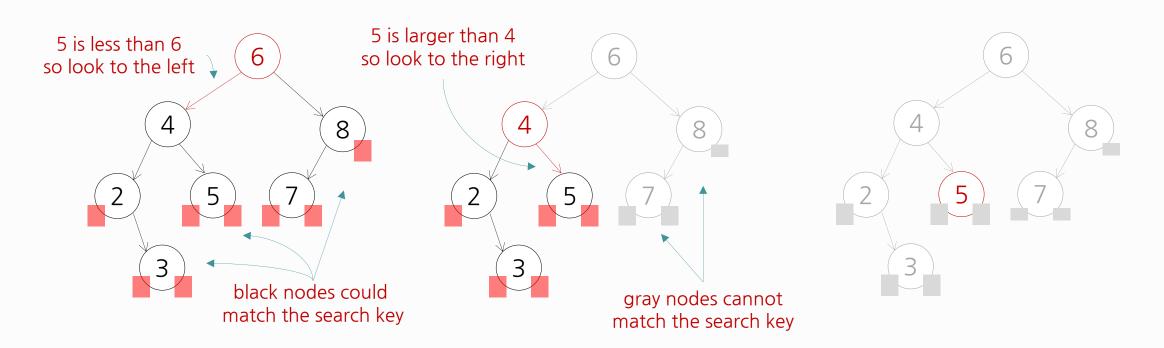
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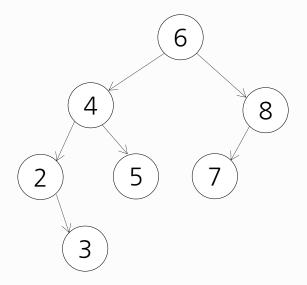


Binary Search Tree: get()

```
class BST:
    def get(self, key):
        return self._get(self.root, key)
    def _get(self, root, key):
        if root is None or root.key == key:
            return root is not None
        elif key < root.key:</pre>
            return _get(root.left, key)
        elif key > root.key:
            return _get(root.right, key)
    def __getitem__(self, key):
        pass
    def __contains__(self, key):
        pass
    def contains(self, key):
        pass
```

Binary Search Tree: get()

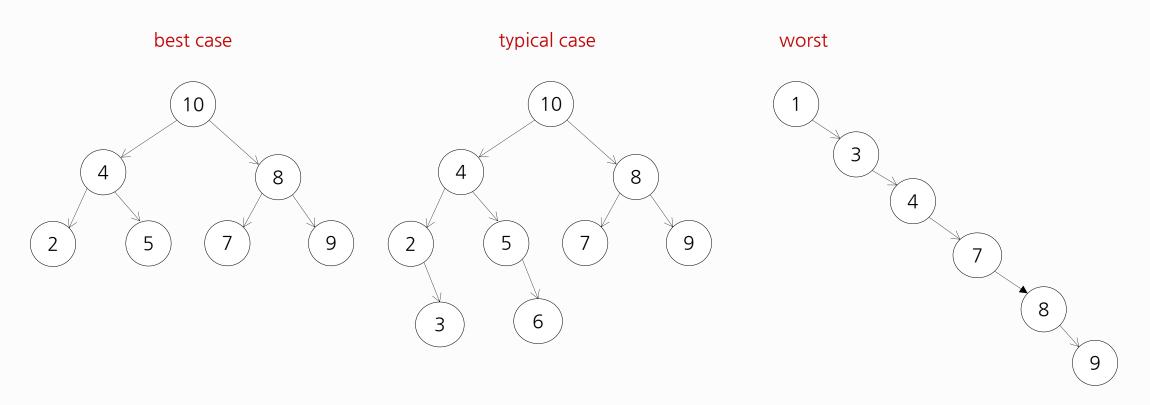
```
def _get(self, root, key):
    if root is None or root.key == key:
        return root is not None
    elif key < root.key:
        return _get(root.left, key)
    elif key > root.key:
        return _get(root.right, key)
```



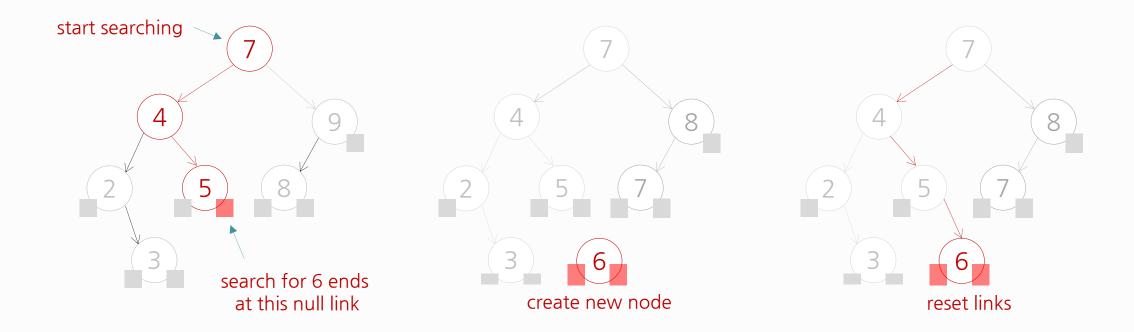
```
def _get_iter(node, key)
   if node is None: return False
   while node is not None:
       if key == node.key: return True;
       if key < node.key:
            node = node.left
       else:
            node = node.right
   return False</pre>
```

Binary Search Tree: Search Analysis

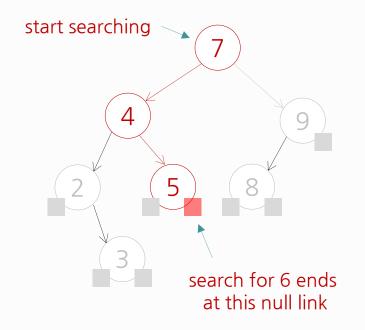
- The running times of algorithms on binary search trees depend on the shapes of the trees, which, in turn, depends on the order in which keys are inserted.
 - It is reasonable, for many applications, to use the following simple model: We assume that the keys are (uniformly) random, or, equivalently, that they are inserted in random order.
 - Insertion and search misses in a BST built from N random keys requires O(h), where h is the height of a BST. The typical case is $\sim 1.39 \log_2 N$ compares on the average.

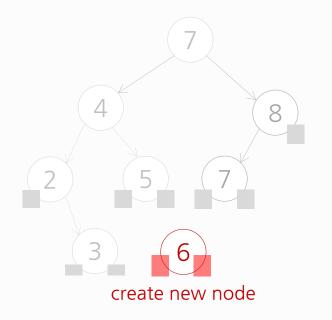


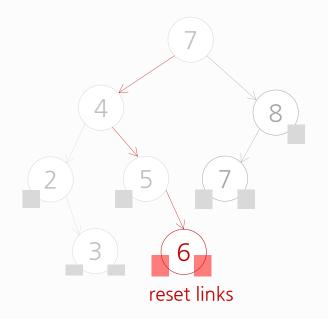
- Insertion operation is very similar to search. Indeed, a search for a key not in the tree ends at a null
 link, and all that we need to do is replace that link with a new node containing the key.
 - add(root, 6)



```
def _add(self, node, key):
    if node is None:
        node = Node(key)
    else:
        if key < node.key:
            node.left = self._add(node.left, key)
        elif key > node.key:
            node.right = self._add(node.right, key)
        else:
            pass
    return node
```

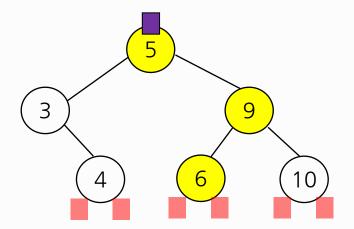






- _add(self, node, key) Insert a node with key
 - Step 1: If the tree is empty, return a new Node(key).
 - Step 2: Pretending to search for key in BST, until locating a None.
 - Step 3: create a new node(key) and link it.

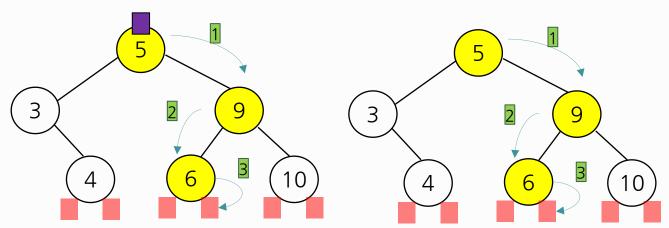
_add(node, 7)



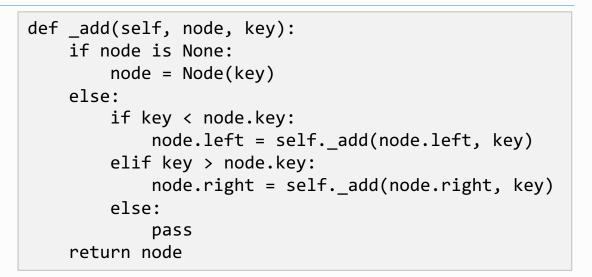
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_add(node, 7) The highlight nodes are compared with key 7.

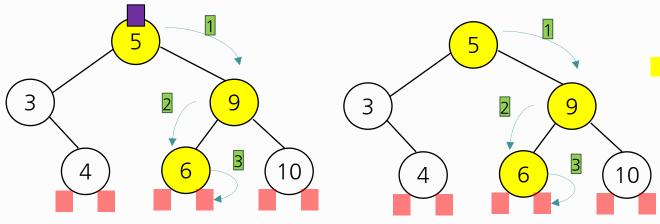


```
after node(6) & key(7) compared, it calls _add(None, 7)
```

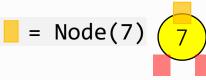


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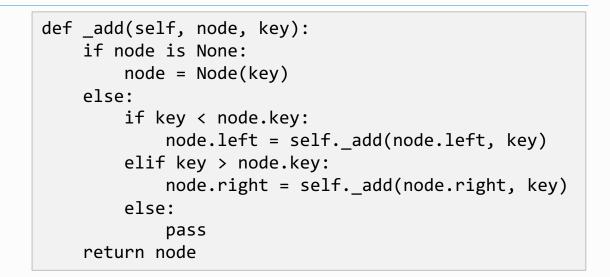
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after node(6) & key(7) compared, it calls **_add(None, 7)**



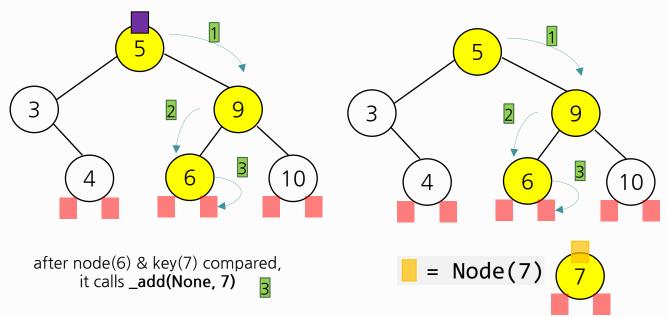
Where does it return to? 6 and 7 are **not** linked.



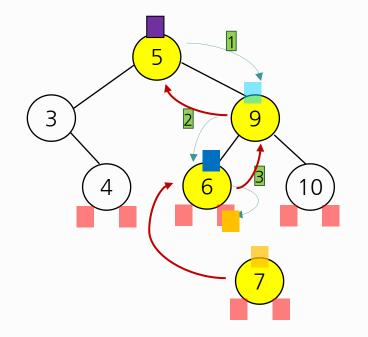
Where does it return to?

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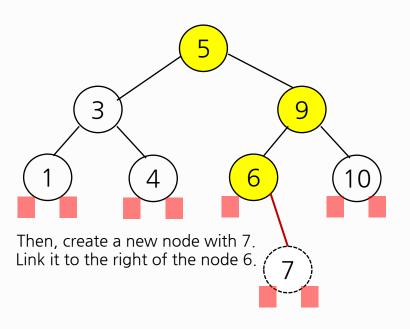
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def _add(self, node, key):
    if node is None:
        node = Node(key)
    else:
        if key < node.key:
            node.left = self._add(node.left, key)
        elif key > node.key:
            node.right = self._add(node.right, key)
        else:
            pass
    return node
```



Questions:

- 1. Explain the differences between the binary tree and binary search tree in this operation.
- 2. To complete inserting **7**, how many times was **_add()** called?
- 3. How many times " if key < node->key" called during this process or inserting 7?
- 4. At the end of this whole process, which **return** will be executed and what is the key value of the node?

```
def _add(self, node, key):
    if node is None:
        node = Node(key)
    else:
        if key < node.key:
            node.left = self._add(node.left, key)
        elif key > node.key:
            node.right = self._add(node.right, key)
        else:
            pass
    return node
```



학습 정리

1) BST의 ADT에는 get(), add(), delete() 등이 있다

2) 균형잡힌 트리에 대한 작업의 시간복잡도는 $O(\log n)$ 이지만, 불균형을 이룬 트리는 최대 O(n)까지 증가할 수 있다

