

# 학습 목표

해시(Hash) 테이블과 해시 함수를 이해하고 특징들을 학습한다



# **Data Structures in Python Chapter 6**

- Hash Table
- Collision Resolution
- Double Hashing & Rehashing
- HashMap Coding

# Agenda & Readings

- Agenda
  - Hashing
  - Hash Table
  - Hash Function
- Reference:
  - Problem Solving with Algorithms and Data Structures
  - Chapter 5 Hashing

# Overview

Hashing or Hash Table Data Structure:

Data structures so far

Array of size n	unsorted list	sorted array	Trees BST – average AVL – worst	Heap, Priority Queue	Hashing
insert	find+0(1)	0(n)	O(log n)	O(log n)	
find	0(n)	O(log n)	O(log n)	O(log n)	
remove	find+0(1)	0(n)	O(log n)	O(log n)	

# Overview

 Hashing or Hash Table Data Structure: supports insertion, deletion and search in average case constant time O(1).

Data structures so far

Array of size n	unsorted list	sorted array	Trees BST – average AVL – worst	Heap, Priority Queue	Hashing
insert	find+0(1)	0(n)	O(log n)	O(log n)	0(1)
find	0(n)	O(log n)	O(log n)	O(log n)	O(1)
remove	find+0(1)	0(n)	O(log n)	O(log n)	O(1)

#### Overview

 Hashing or Hash Table Data Structure: supports insertion, deletion and search in average case constant time O(1).

#### Hash table

- It is data structure that stores key-value pairs.
- The key is sent to a hash function that performs arithmetic operations on it.
- The result is called hash value that is the index of the key-value pair in the hash table.

#### Hash function

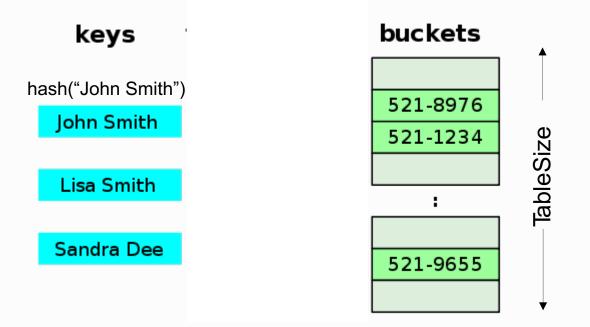
- hash(key) → integer value
- hash("string key") → integer value

Hash table is an array of fixed size elements

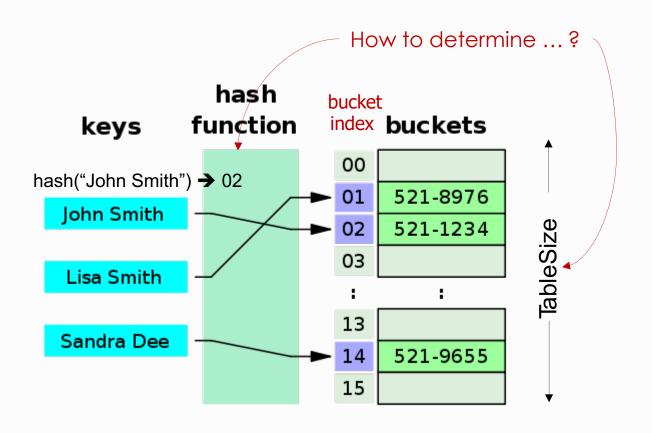
Let us suppose that there are one billion of names and numbers.

• Find, insert, and remove a number by a given name in **O(1)**.

Time Complexity



- Hash table is an array of fixed size elements
- Array elements indexed by a key mapped to a bucket index[0 .. TableSize-1]
- Mapping from key to index using hash(), hash function
  - e.g., hash("John Smith") → 02

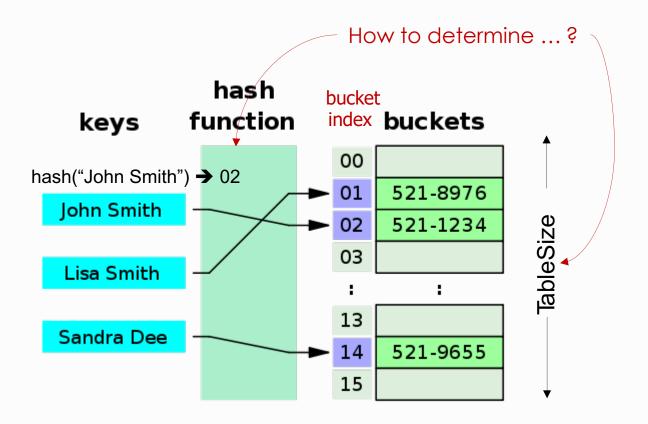


- insert
  - HashTable[hash("John Smith")] = <"John Smith", 521-1234>
- remove
  - HashTable[hash("John Smith")] = None
- find
  - HashTable[hash("John Smith")]
     returns the element hashed for "John Smith"

How to determine ...? hash bucket function index buckets keys 00 hash("John Smith") → 02 521-8976 01 John Smith 02 521-1234 **TableSize** 03 Lisa Smith 13 Sandra Dee 14 521-9655 15

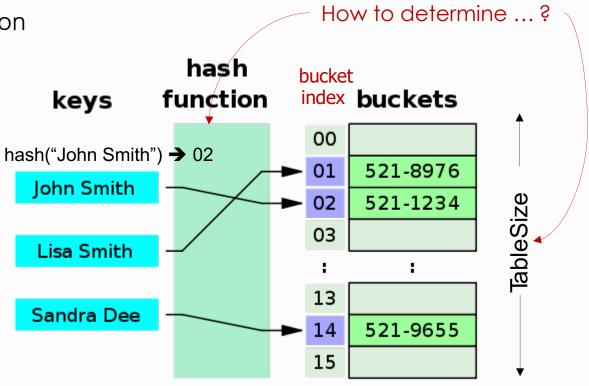
What happens if hash("John Smith") == hash("Joe Blow")? "Collision"

- Factors affecting Hash Table Design
  - Hash function
  - Table size Usually fixed at the start
  - Collision handling schemes Array or Linked List



# **Hash Function**

- It maps an element's key into a valid hash table index
  - hash(key) → hash table index
- Note that this is (slightly) different from saying:
  - hash(string) → int
  - Because the key can be of any type
    - e.g., "hash(int) → int" is also a hash function



# **Hash Function Properties**

- It maps an element's key into a valid hash table index
  - hash(key) → hash table index
- It maps key to integer
  - Constraint: Integer should be between [0, TableSize-1]
- A hash function can result in a many-to-one mapping (causing collision)
  - Collision occurs when hash function maps two or more keys to same array index
- Collisions cannot be avoided but its chances can be reduced using a "good" hash function

#### Hash Function - Effective use of table size

- Simple hash function (assume integer keys)
  - hash(Key) = Key % TableSize
- For random keys, hash() distributes keys evenly over table
  - What if TableSize = 100 and keys are ALL multiples of 10?
  - Better if TableSize is a prime number

# Hash Function Example: String Keys

- Using a very simple function to map strings to integers:
  - Add up character ASCII values (0-255) to produce integer keys
    - e.g., "abcd" = 97 + 98 + 99 + 100 = 394
    - hash("abcd") = 394 % TableSize
- Potential problems:
  - Anagrams will map to the same index
    - hash("abcd") = hash("dbac")
  - Small strings may not use all of table
    - strlen(s) \* 255 < TableSize</li>
  - Time proportional to length of the string

# Hash Function Example: String Keys

- Another approach:
  - Treat first 3 characters of string as base-27 integer(26 letters plus space)
    - e.g., Key =  $s[0] + (27^1 * s[1]) + (27^2 * s[2])$
    - Better than previous approach because ...
- But, potential problems:
- Apple
   Apply
   Appointment
   Apricot

# Hash Function Example: String Keys

- Last approach:
  - Use all N characters of string as an N-digit and base-K number
  - Choose K to be prime number larger than number of different digits (characters)
    - i.e., K = 29, 31, 37
    - If L = Length of string S, then

$$hash(S) = \sum_{i=0}^{L-i} S[L-1-i] * 37^{i} \% TableSize$$
 (1)

- Use Horner's rule to compute hash(S).
- Limit L for long strings
- Potential problems
  - Overflow
  - Larger runtime

```
# a hash function for strings
hash(key, tablesize)
  code = 0
  for x in key:
     code = code * 37 + x
  code %= tablesize
  if code < 0: code += tablesize
  return code</pre>
```

# Summary

- Using a hash table we can, on average (if table large enough and hash function suitable), insert, delete and search for items in constant time - O(1).
- The hash function is the mapping between an item and the slot where the item is stored.
- A collision occurs when an item is mapped to an occupied slot.
- A perfect hash function is able to map m items into a table of size m with no collisions. Perfect hash functions are hard to come by.
- Handling collisions systematically is required collision resolution.

# 학습 정리

1) 해시 테이블을 이용하면, 삽입/삭제/검색 작업을 O(1)으로 해결할 수 있다

2) 해시 함수가 반환하는 해시 값(Value)은 항상 정수이다

3) 해시 테이블을 구현할 때, 충돌을 최소화할 수 있는 해시 함수를 사용해야 한다

