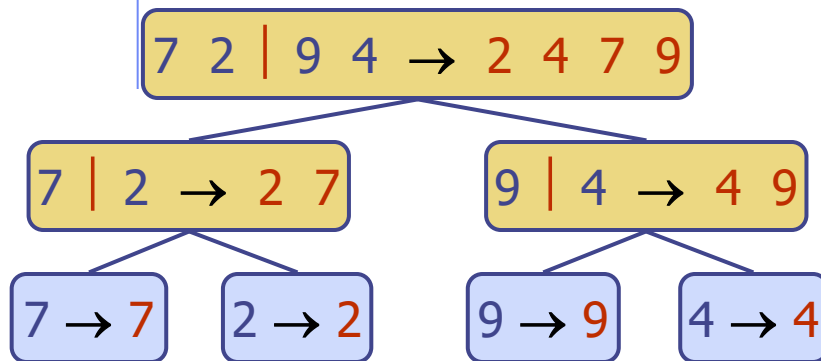


Lesson 4

Recursion: Collapsing Infinity To a Point

Wholeness of the Lesson

Recursive algorithms keep reducing the size of the input instances until a base case is reached, then the solution is computed from the base case up to the solution for the whole problem. Maharishi describes the process of creation as a self-referral process within pure consciousness that unfolds sequentially.



The slide features a minimalist design with thin blue lines. A vertical line on the left and a horizontal line intersect near the top-left corner, with a small blue circle at the intersection. Another horizontal line is positioned below the title. A vertical line on the right and a horizontal line intersect near the bottom-right corner, also with a small blue circle at the intersection. A solid blue horizontal bar is at the very top of the slide.

Recursive Programming

Basic Concepts

- Recognizing Recursion
 - When smaller or simpler instances form sub-constituents of the overall solution
 - E.g., when a function calls itself on smaller subproblem instances to solve the larger, global problem
- Theoretically, any problem that can be solved using iteration (while and for loops) can be solved using recursion (functional style is supported by functional languages)

Recursive Thinking

1. Define the base cases
 - Instance(s) that can be calculated without using recursive calls
2. Decompose the problem into simpler or smaller instances of the original problem
 - A smaller/simpler instance must be moving toward one of the base cases (so the function terminates)
3. Create an induction diagram to determine what to do in addition to the recursive calls

Types of Recursion

- Linear recursion
- Tail recursion
- Multiple recursion
- Mutual recursion
- Nested recursion

Types of Recursion

- Linear recursion

- When a method calls itself only once in the body of the function

Algorithm **sumFirst**(n)

if $n < 0$ then Throw `InvalidInputException`

if $n = 0$ then

 return 0

else

 return $n + \text{sumFirst}(n-1)$

Types of Recursion

- Tail recursion

- A special case of linear recursion in which a method calls itself only once but the call occurs as the last operation executed in the body of the method
- Functional languages optimize tail recursive functions since there is no need to create a new stack frame (activation record)

Algorithm **sumFirst(n)**

```
if n < 0 then Throw InputException
return sumFirstHelper(n, 0)
```

Algorithm **sumFirstHelper(n, s)**

```
if n = 0 then
    return s
else
    return sumFirstHelper(n-1, n+s)
```

Types of Recursion

- **Multiple recursion**
 - When a function calls itself two or more times
- Example is MergeSort and QuickSort (next week)

Algorithm **Fib**(n)

if $n = 0$ then

return 0

else if $n = 1$ then

return 1

else

return **Fib**(n-2) + **Fib**(n-1)

Types of Recursion

- Mutual recursion

- When a group of methods repeatedly call each other until a base case is reached

Algorithm **isEven**(n)

if $n = 0$ then

return true

else

return **isOdd**(n-1)

Algorithm **isOdd**(n)

if $n = 0$ then

return false

else

return **isEven**(n-1)

Types of Recursion

- **Nested recursion**

- When the argument to a recursive call is calculated via another recursive call
- Sometimes called Double Recursion
- Ackermann function is a recursive mathematical function that takes two non-negative integers as inputs and produces a non-negative integer as its output.

Algorithm **A**(n, s)

if $n = 0$ then

 return $s + 1$

else if $s = 0$ then

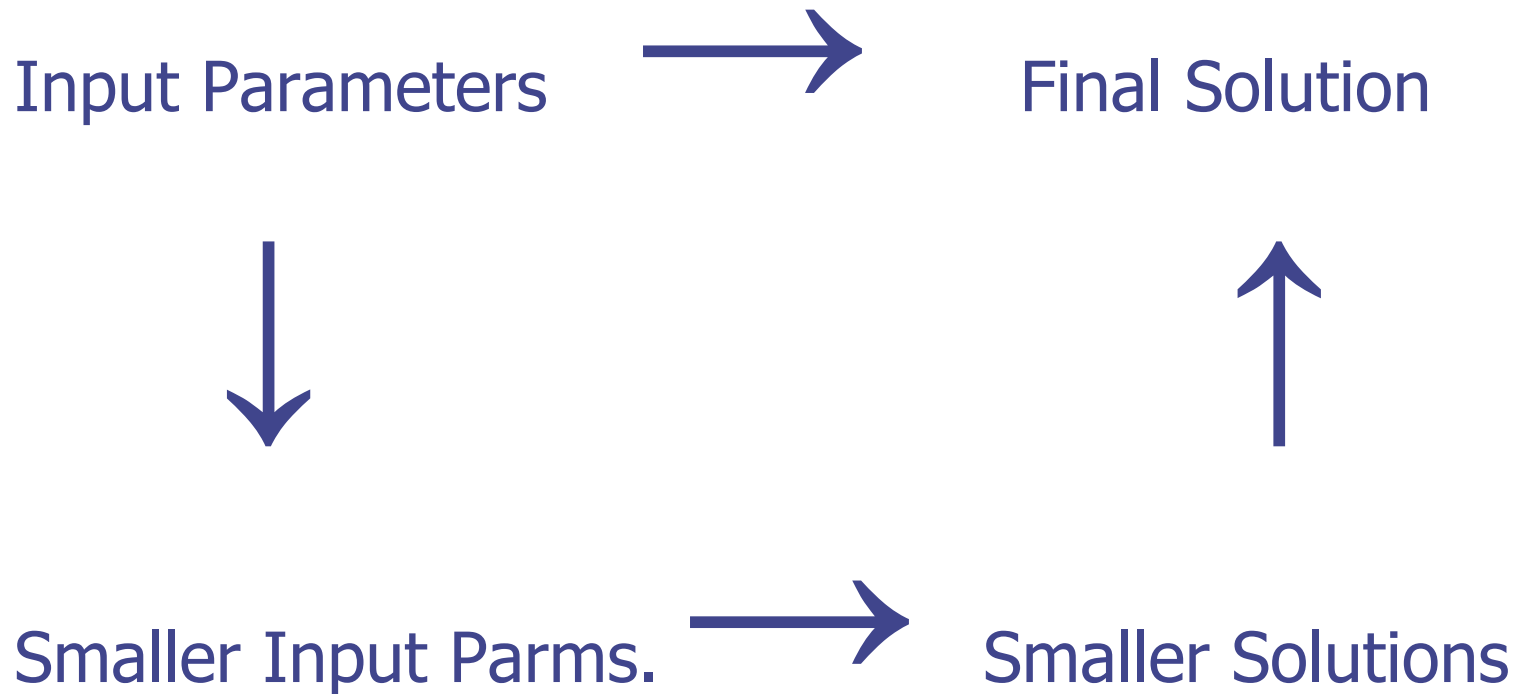
 return $A(n-1, 1)$

else

// $\{n > 0 \text{ and } s > 0\}$

 return $A(n-1, A(n, s-1))$

Recursive Thinking (AKA Subgoal Induction)



Exercises

1. Write a pseudo code function, *isEven*(n) to recursively determine whether a natural number, n, is an even number.
2. Write a pseudo code function, *power*(x, k), that computes x^k . Can you do this in $\log k$ time?

Exercise on List

- Generic methods:
 - integer `size()`
 - boolean `isEmpty()`
 - objectIterator `elements()`
- Accessor methods:
 - position `first()`
 - position `last()`
 - position `before(p)`
 - position `after(p)`
- Query methods:
 - boolean `isFirst(p)`
 - boolean `isLast(p)`
- Update methods:
 - `swapElements(p, q)`
 - `replaceElement(p, o)`
 - `insertFirst(o)`
 - `insertLast(o)`
 - `insertBefore(p, o)`
 - `insertAfter(p, o)`
 - `remove(p)`

Exercise:

- Write a recursive method to calculate the sum of the integers in a list of integers

Algorithm `sum(L)`

Hint: you also need a helper function with argument Position `p`

Algorithm `sumHelper(L, p)`

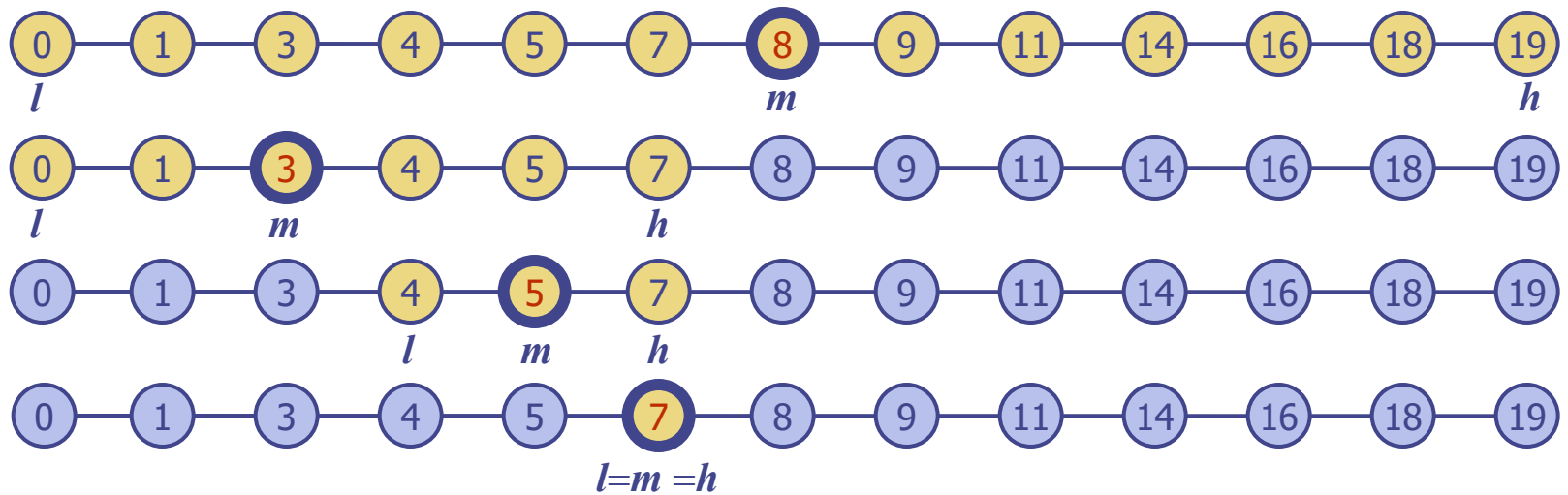
Algorithm Design Strategy

- Prune and Search
 - Also known as Decrease and Conquer
 - Problem solving technique
 - Examples:
 - **binary search**
 - quick select (randomized prune and search)
- Randomized Algorithm (later)
 - Quick Select

Binary Search

<https://yongdanielliang.github.io/animation/web/BinarySearchNew.html>

- Binary search performs operation **findElement(k)** on an array or array-based sequence, sorted by key
 - similar to the high-low game
 - at each step, the number of candidate items is halved
 - terminates after $O(\log n)$ steps
- Example: **findElement(7)**



Binary Search Algorithm (iterative)

Algorithm BinarySearch(S, k):

Input: A sorted array S storing n items

Output: An element of S with value k .

low \leftarrow 0

high \leftarrow S.length - 1

while low \leq high do

 mid \leftarrow floor((low + high)/2)

 if $k = S[\text{mid}]$ then

 return $S[\text{mid}]$ // Found k

 else if $k < S[\text{mid}]$ then

 high \leftarrow mid - 1 // search on left

 else

 low \leftarrow mid + 1 // search on right

return **NO_SUCH_KEY**

Binary Search Algorithm (one key comparison per iteration)

Algorithm BinarySearch(S, k):

Input: A sorted array S storing n items

Output: An element of S with value k .

low \leftarrow 0

high \leftarrow S.size() - 1

mid \leftarrow 0

while low < high do

 mid \leftarrow floor((low + high)/2)

 if $k < S[\text{mid}]$ then

 high \leftarrow mid-1

 else

 low \leftarrow mid

if S.size() > 0 /\ $k = S[\text{mid}]$ then //(/\ - And)

 return S[mid]

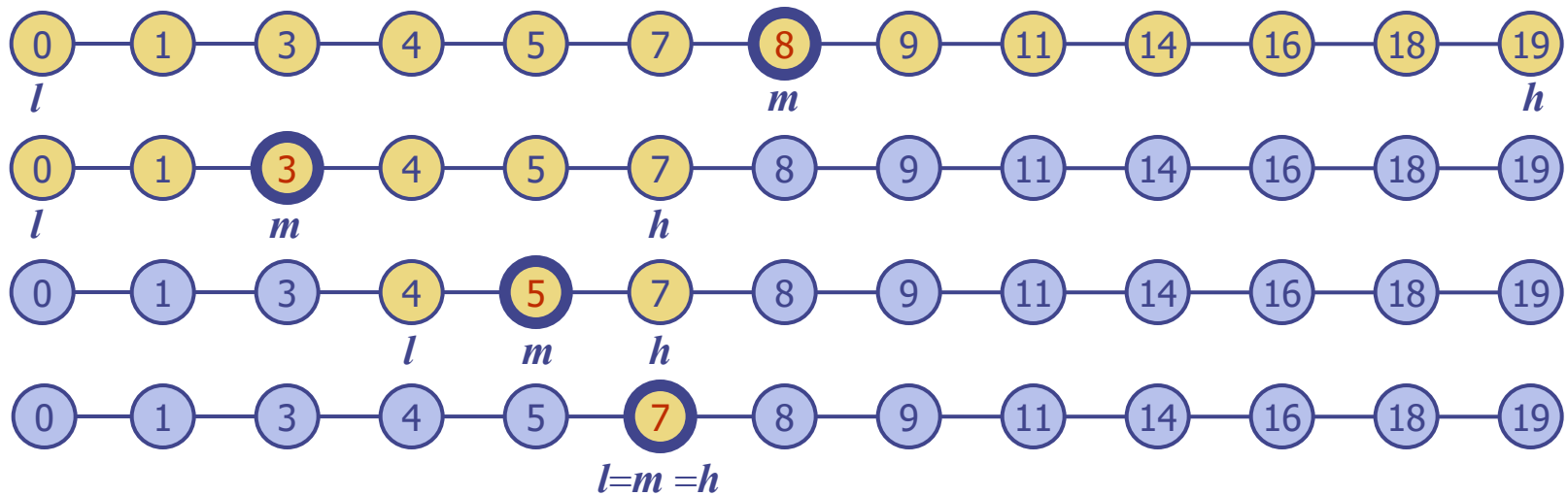
else

 return **NO_SUCH_KEY**

Recursive Binary Search

Binary Search

- Binary search performs operation **findElement**(S, k) on an array or array-based sequence S, sorted by key (requires efficient random access)
 - similar to the high-low game
 - at each step, the number of candidate items (search space) is halved
 - terminates after $O(\log n)$ steps
- Example: **findElement**(S, 7)



Binary Search Algorithm (recursive)

Algorithm BinarySearchRec($S, k, low, high$):

Input: A sorted array S storing n items

Output: An element of S with value k between indices low & $high$.

if $low > high$ then

 return **NO_SUCH_KEY**

else

$mid \leftarrow \text{floor}((low + high)/2)$

 if $k = S[mid]$ then

 return $S[mid]$

 else if $k < S[mid]$ then

 return BinarySearchRec($S, k, low, mid-1$)

 else

 return BinarySearchRec($S, k, mid + 1, high$)

In-Class Practice

- Write a pseudo code function, *sum*(n), to recursively sum the first n natural numbers but divide the problem in half and make two recursive calls.

Main Point

1. Any iterative algorithm can be computed using recursion, i.e., a function calling itself. In fact, the meaning of while- and for-loops are defined using recursive functions in programming language semantics (Denotational Semantics). Recursive algorithms keep reducing the size of the inputs instances until a base case is reached, then the solution is computed from the base case up to the solution for the whole problem.

Science of Consciousness: Maharishi describes the process of creation as a self-referral process that unfolds sequentially. The dynamism of the unified field seems chaotic when studied at the macroscopic level, yet it is a field of perfect order, responsible for the order and balance in creation.

Connecting the Parts of Knowledge With the Wholeness of Knowledge

Recursion

1. A recursive algorithm calls itself repeatedly until it reaches the base case, then the results of these calls are combined to compute the final result.
2. There are different categories of recursive algorithms depending on how many and where the recursive calls occur, i.e., linear, tail, multiple, mutual, and nested (double) recursion.
3. *Transcendental Consciousness* is the self-referral field of *infinite correlation*, from which all of creation emerges, where “an impulse anywhere is an impulse everywhere.”
4. *Impulses within the Transcendental field*. The dynamic natural laws reside in this field, where the self-referral laws of nature govern all of creation.
5. *Wholeness moving within itself*. In Unity Consciousness, the field of action effortlessly unfolds as the play of one’s own Self, one’s own pure consciousness.