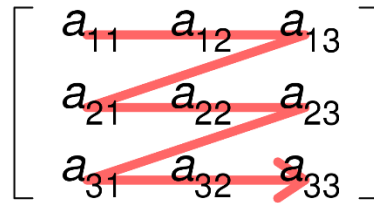
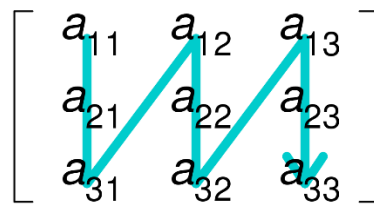


Row Major and Column Major order

Row-major order



Column-major order



Address Calculation in 1D Array:

Memory address calculation Address of an element of an array say $A[I]$ is calculated using the following formula:

$$\text{Address of } A[I] = B + W * (I - LB)$$

Where,

B = Base address

W = Storage Size of one element stored in the array (in byte)

I = Subscript of element whose address is to be found

LB = Lower limit / Lower Bound of subscript, if not specified assume 0 (zero)

Example:

Given the base address of an array $B[1300..1900]$ as 1020 and size of each element is 2 bytes in the memory. Find the address of $B[1700]$.

Solution:

The given values are: $B = 1020$, $LB = 1300$, $W = 2$, $I = 1700$

$$\text{Address of } A[I] = B + W * (I - LB)$$

$$= 1020 + 2 * (1700 - 1300)$$

$$= 1020 + 2 * 400$$

$$= 1020 + 800$$

$$= 1820$$

Address calculation in 2-D array: suppose A is a 2D array with dimensions M1 X M2. So that here total number of elements is equal to M1 X M2. Address should be calculated using one of the two methods as the following.

a. Row-by-Row

$$\text{Location (A [j, k])} = \text{Base (A)} + w [N (j-1) + (k -1)]$$

Here, j and k indicates index value of the element whose address you want to search.

Base (A) is the starting address of the given array. **W** is the memory word required to store given array's element. Here, array is the of size **M(rows) X N(columns)**.

b. Column-by-Column

$$\text{Location (A [j, k])} = \text{Base (A)} + w [M (k-1) + (j -1)]$$

Here, j and k indicates index value of the element whose address you want to search.

Base (A) is the starting address of the given array. **W** is the memory word required to store given array's element. Here, array is the of size **M(rows) X N(columns)**.

Examples:

1. Consider 30 X 4 2D arrays and base address is 200 and 1 word per memory locations. Find out the address of A (15, 3).

$$\text{Base (A),} \quad M=30, \quad N=4, \quad j=15, \quad k=3, \quad w=1$$

a. Row-by-row

$$\text{Location (A [j, k])} = \text{Base (A)} + w [N (j-1) + (k -1)]$$

$$\begin{aligned} \text{Location (A [15, 3])} &= 200 + 1 [4 (15 - 1) + (3 - 1)] \\ &= 200 + 1 [56 + 2] \\ &= 200 + 58 \\ &= 258 \end{aligned}$$

b. Column-by-column

$$\text{Location (A [j, k])} = \text{Base (A)} + w [M (k-1) + (j -1)]$$

$$\begin{aligned} \text{Location (A [15, 3])} &= 200 + 1 [30 (3 - 1) + (15 - 1)] \\ &= 200 + 1 [30 (2) + 14] \\ &= 200 + 74 \\ &= 274 \end{aligned}$$

Address calculation in multidimensional array: Suppose A is a 3D array of size M1 X M2 X M3. So total number if elements are M1 X M2 X M3. In 3D array lower bound (LB) is not always zero (0). It can be any number. Base address of any row of matrix is called effective address.

A [2:11, 3:15] indicates that A is a 2D array with lower bound (LB) of row is 2 and Upper Bound (UB) of row is 11. In the same way lower bound (LB) of column is 3 and upper bound (UB) of column is 15.

Length of each dimension is equal to $L_i = UB - LB + 1$

Lower bound (LB) is not always started with zero (0). So, we have to calculate effective address of each value.

$$\text{Effective Address (E}_i\text{)} = K_i - LB$$

Real address of any element is calculated as follows,

Row-by-Row: **Base (A) + W [(E₁L₂ + E₂) L₃ + E₃]**

Column-by-column: **Base (A) + W [(E₃L₂ + E₂) L₁ + E₁]**

Example: Suppose multidimensional array A and B are declared using **A [-2:2, 2:22]** and **B [1:8, -5:5, -10:5]**.

1. Find out the length of each dimension and number of elements in array A and B.
2. Consider B [3, 3, 3] elements in array B find effective address E₁, E₂ and E₃. Also find out the real address of these elements. (In this case the base address of array B is 400 and W is 4)

Answer: For array A

$$\begin{aligned} L_1 &= UB - LB + 1 & L_2 &= UB - LB + 1 \\ &= 2 - (-2) + 1 & &= 22 - 2 + 1 \\ &= 4 + 1 = 5 & &= 21 \end{aligned}$$

For array B

$$\begin{aligned} L_1 &= UB - LB + 1 & L_2 &= UB - LB + 1 & L_3 &= UB - LB + 1 \\ &= 8 - 1 + 1 & &= 5 - (-5) + 1 & &= 5 - (-10) + 1 \\ &= 8 & &= 11 & &= 16 \end{aligned}$$

Total number of elements in array B is: $L_1 \times L_2 \times L_3 = 8 \times 11 \times 16 = 1408$.

Here in given question $K_i = 3, 3, 3$ therefore $K_1=3, k_2 = 3$ and $K_3 = 3$

$$LB_1 = 1, LB_2 = -5, LB_3 = -10$$

$$E_1 = K_1 - LB_1 = 3 - 1 = 2$$

$$E_2 = K_2 - LB_2 = 3 - (-5) = 8$$

$$E_3 = K_3 - LB_3 = 3 - (-10) = 13$$

The real address Row – by- Row:

$$\begin{aligned} &= \text{Base (A)} + W [(E_1 L_2 + E_2) L_3 + E_3] \\ &= 400 + 4 [(2 (11) + 8) 16 + 13] \\ &= 400 + 4 [(22+8) 16 + 13] \\ &= 400 + 4 [(30) 16 + 13] \\ &= 400 + 4 [480 + 13] \\ &= 400 + 4 [493] \\ &= 400 + 1972 \\ &= 2372 \end{aligned}$$

The real address Column – by – column:

$$\begin{aligned} &= \text{Base (A)} + W [(E_3 L_2 + E_2) L_1 + E_1] \\ &= 400 + 4 [(13(11) + 8) 8 + 2] \\ &= 400 + 4 [(143 + 8) 8 + 2] \\ &= 400 + 4 [(151) 8 + 2] \\ &= 400 + 4 [1208 + 2] \\ &= 400 + 4840 \\ &= 5240 \end{aligned}$$

Let **C** be an n-dimensional array.

The length **L_i** of dimension **i** of **C** can be calculated by,

$$L_i = \text{upper bound} - \text{lower bound} + 1$$

For a given subscript **K_i**, the effective index **E_i** of **L_i** can be calculated from,

$$E_i = K_i - \text{lower bound}$$

Then address **C[K₁,K₂,...,K_n]** of element of **C** is,

$Base(C) + w[(((\dots(E_n L_{n-1} + E_{n-1}) L_{n-2}) + \dots + E_3) L_2 + E_2) L_1 + E_1]$, when C is stored in column major, and

$Base(C) + w[(\dots((E_1 L_2 + E_2) L_3 + E_3) L_4 + \dots + E_{n-1}) L_n + E_n]$, when C is stored in row major order.

$Base(C)$ denotes the address of 1st element of C and w denotes the number of words per memory location.

Compute the location of $A[10][10]$ if the base address of $A[-15:20, 10:35]$ is given as 3000 and occupying 4 bytes of memory. Assume row major storage.