

香港中文大學
The Chinese University of Hong Kong

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Course Examinations 2003 - 2004

Course Code & Title : CSC 3130 Formal Languages and Automata Theory

Time allowed : 2 hours 0 minutes

Student I.D. No. : Seat No. :

1. Consider the following ambiguous CFG G :
- $$S \rightarrow 0S \mid 1S \mid S0 \mid \varepsilon$$
- (a) Describe in English what the language L generated by G is.
(b) Give a string in L that has two leftmost derivations. Show the two leftmost derivations.
(c) Convert G to a PDA accepting by empty stack.
(d) Convert your answer in (c) to a PDA accepting by final states.
(e) Give an unambiguous grammar for L . Explain briefly why the grammar given is unambiguous.
2. For each of the following pairs of languages, choose the correct relationship from the following four choices:
- Type 1: L_1 is a proper subset of L_2 .
Type 2: L_2 is a proper subset of L_1 .
Type 3: L_1 and L_2 are the same.
Type 4: L_1 and L_2 have no containment relationship, i.e., there are strings in L_1 but not in L_2 and there are strings in L_2 but not in L_1 .
- (a) L_1 is the language of the CFG with productions $S \rightarrow 0S1 \mid 1S0 \mid \varepsilon$.
 L_2 is the language of the regular expression $(0+1)^*$.
(b) L_1 is the language of the CFG with productions $S \rightarrow AS \mid SB \mid \varepsilon$, $A \rightarrow 0$ and $B \rightarrow 1$.
 L_2 is the language of the regular expression 0^*1^* .
(c) L_1 is the language of the regular expression $(0+1)^*11(0+1)^*$.
 L_2 is the language of the regular expression $(0^*1^*11)^*0^*110^*1^*$.
(d) L_1 is the language accepted by the NFA $(\{p, q\}, \{0, 1\}, \delta, q, \{q\})$ with transitions:
 $\delta(q, 0) = \{p\}$, $\delta(q, 1) = \{ \}$, $\delta(p, 0) = \{p, q\}$, $\delta(p, 1) = \{p\}$.
 L_2 is the language of the CFG with productions $S \rightarrow AS \mid \varepsilon$, $A \rightarrow 0B$ and $B \rightarrow 0B \mid 1B \mid 0$.
3. Prove by the Pumping Lemma that the following language is not context-free:
- $$L = \{0^i1^j \mid j = i^3\}$$
4. Give the transition diagram of a Turing machine M that, given two non-zero binary numbers x and y of the same length, will determine whether $x \geq y$:

Input:	Output:
...##101#011##...	...##1#101#011##...
^	^ (since $101_2 \geq 011_2$)
...##0010#0101##...	...##0#0010#0101##...
^	^ (since $0010_2 < 0101_2$)
...##11#11##...	...##1#11#11##...
^	^ (since $11_2 \geq 11_2$)

Notice that the input binary numbers x and y should remain in the tape, the answer is written on the left hand side separated by a blank and the tape head should be pointing to the answer when M accepts. Explain briefly your construction.

5. Consider the following instance of the Modified Post's Correspondence Problem:

Index	A	B
1	10	101
2	101	011
3	110	100

- (a) Explain why the above MPCP instance has no solution.
(b) Give a fourth pair so that there are solutions. Give one solution.
6. (a) Consider the following language:
$$L_{\text{forever}} = \{(k, w) \mid \text{Turing machine } T_k \text{ runs forever on input } w\}$$

Explain why if L_{forever} is recursive, the universal language L_u is also recursive.
- (b) Assume that it is undecidable whether a given context-free grammar is ambiguous. A CFG G is called *super-ambiguous* if there is at least one string w in $L(G)$ where w has at least 10 different leftmost derivations. Explain why it is also undecidable whether a given context-free grammar is super-ambiguous.

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