

香 港 中 文 大 學  
The Chinese University of Hong Kong

版權所有 不得翻印  
Copyright Reserved

二 0 一 四 至 一 五 年 度 上 學 期 科 目 考 試  
Course Examination 1<sup>st</sup> Term, 2014-15

科目編號及名稱  
Course Code & Title : CSCI3130 Formal Languages and Automata Theory

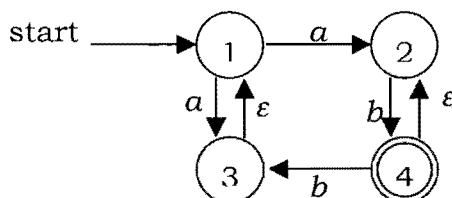
時間  
Time allowed : 2 小時 hours 分鐘 minutes

學號  
Student I.D. No. : 座號  
Seat No. :

**Answer all questions.**

1. **(10 marks)** Let  $n_a(\omega)$  denote the number of occurrences of symbol  $a$  in string  $\omega$ , and  $n_b(\omega)$  the number of occurrences of symbol  $b$  in string  $\omega$ . Is the language  $\{\omega \mid \omega \in \{a, b\}^*, n_a(\omega) = n_b(\omega)\}$  regular? Prove your answer.

2. Consider the following  $\epsilon$ -NFA.



- a. **(10 marks)** Convert the  $\epsilon$ -NFA to a DFA.
- b. **(10 marks)** Give a regular expression that is equivalent to the  $\epsilon$ -NFA.
3. Consider the following language.  
 $\{a^i b^j c^k \mid 0 < j < i, j = i + k\}$   
**(10 marks)** Is the language context free? Prove your answer.
4. **(10 marks)** Prove that the following language is not context free.  
 $\{a^i b^{2i} c^j \mid 0 < j < i\}$

5. A CFG (with  $E$  as its start symbol) and its SLR parsing table are shown below.

- (1)  $E \rightarrow E+T$       (4)  $T \rightarrow F$   
 (2)  $E \rightarrow T$       (5)  $F \rightarrow (E)$   
 (3)  $T \rightarrow T*F$       (6)  $F \rightarrow \mathbf{id}$

STATE	ACTION						GOTO		
	id	+	*	(	)	\$	$E$	$T$	$F$
0	s5			s4			1	2	3
1		s6				acc			
2		r2	s7		r2	r2			
3		r4	r4		r4	r4			
4	s5			s4			8	2	3
5		r6	r6		r6	r6			
6	s5			s4				9	3
7	s5			s4					10
8		s6			s11				
9		r1	s7		r1	r1			
10		r3	r3		r3	r3			
11		r5	r5		r5	r5			

Figure 4.37: Parsing table for expression grammar

(Source: Aho, Lam, Sethi & Ullman, *Compilers: Principles, Techniques, and Tools* (2nd Edition). Addison-Wesley.)

- a. **(10 marks)** Copy the following table to the answer book, and show using the copied table the process of parsing the string ' $\mathbf{id} * (\mathbf{id})$ ' using the parsing table above.

	STACK	SYMBOLS	INPUT	ACTION
(1)			$\mathbf{id} * (\mathbf{id}) \$$	
(2)				

- b. Consider the following CFG, in which  $S$  is the start symbol.

- (1)  $S \rightarrow i B t S e S$   
 (2)  $S \rightarrow a$   
 (3)  $B \rightarrow \mathbf{id}$

- (10 marks)** Give a PDA that accepts the language of this CFG.
- (10 marks)** Define the FIRST and FOLLOW functions for this CFG.
- (10 marks)** Is this grammar LL(1)? Justify your answer.

6. **(10 marks)** Despite the existence of the halting problem, we know that some Turing machines do halt on some particular inputs. Let  $H$  denote the set  $\{\langle T, i \rangle \mid T \text{ is a Turing machine that halts on input } i\}$ . Discuss whether or not the following problem is decidable: given any Turing machine  $X$ , whether or not there exists input  $a$  such that  $\langle X, a \rangle \in H$ .

— End of Paper —