SLICOT Linear Systems Identification Toolbox *

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July 2000

^{*}This document presents research results of the European Community BRITE-EURAM III Thematic Networks Programme NICONET (contract number BRRT-CT97-5040) and is distributed by the Working Group on Software WGS. WGS secretariat: Mrs. Ida Tassens, ESAT - Katholieke Universiteit Leuven, K. Mercierlaan 94, 3001-Leuven-Heverlee, BELGIUM. This report is also available by anonymous ftp from wgs.esat.kuleuven.ac.be in the directory pub/WGS/REPORTS/SLWN2000-4.ps.Z

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Contents

1	Introduction	1
2	SLICOT Identification Toolbox: Algorithmic Outline	2
	2.1 Input-Output Data Processing using Subspace Techniques	2
	2.2 QR Factorization	3
	2.3 Cholesky Factorization	4
	2.4 Computation of System Matrices	6
3	SLICOT Identification Toolbox: Software Components	9
	3.1 Fortran Routines	9
	3.2 Interfaces: Mexfiles and m-files	13
4	Benchmarks for Identification	18
5	Industrial Examples	19
6	Conclusions	23

Abstract

Basic algorithmic, implementation and numerical issues involved in the new, subspacebased linear multivariable system identification toolbox—SLIDENT—incorporated in the freely available Subroutine Library in Control Theory (SLICOT) are described. Reliability, efficiency, and ability to solve large, industrial identification problems received a special consideration in the toolbox development process. Flexibility of usage is enhanced by providing various options: two algorithmic subspace-based approaches (MOESP and N4SID) and their combination, standard or fast techniques for data compression (including abilities to exploit the block-Hankel structure), multiple (possibly connected) data batches processing, availability of both, fully documented drivers and computational routines, the use of structure exploiting algorithms and dedicated linear algebra tools, etc. In addition, optionally, the quality of the intermediate results can be assessed by inspecting the associated condition numbers. The mexfiles developed based on the SLICOT Fortran subroutines enable to combine the efficiency of Fortran calculations with the simplicity of use in a MATLABlike environment. Related m-files are also included in the toolbox, in order to increase the user-friendliness. Extensive comparisons with the available subspace techniques have been made. The numerical results show that SLICOT system identification toolbox—SLIDENT is highly performant.

This report summarizes the achievements and deliverables of Task III.A of the NICONET Project. After a short description of the linear system identification problem and of the available subspace-based techniques to solve it, the numerical algorithms implemented in SLIDENT are surveyed. The associated Fortran routines are then listed and their functional abilities are outlined. The developed interfaces to MATLAB or Scilab, as well as examples of use are presented. Comparisons with the available MATLAB codes are included, illustrating the efficiency and accuracy of the SLIDENT components.

1 Introduction

Consider a linear time invariant (LTI) state space model, described by

$$\begin{aligned}
 x_{k+1} &= Ax_k + Bu_k + w_k, \\
 y_k &= Cx_k + Du_k + v_k,
 \end{aligned} \tag{1}$$

where x_k is the n-dimensional state vector at time k, u_k is the m-dimensional input (control) vector, y_k is the ℓ -dimensional output vector, $\{w_k\}$ and $\{v_k\}$ are state and output disturbance or noise sequences, and A, B, C, and D are real matrices of appropriate dimensions. In system identification problems, the system order, n, and the quadruple of system matrices (A, B, C, D) have to be determined; the only available information is given by an upper bound, s, on n, and by the input and output data sequences, $\{u_k\}$ and $\{y_k\}$, k=1: t (i.e., for k taking integer values from 1 to a given t). In addition, the Kalman gain matrix, as well as the state and output covariance matrices have often to be found.

Two basic subspace-based approaches have been proposed for solving system identification problems: MOESP and N4SID. The main feature of the MOESP class of techniques is the determination of an extended observability matrix of the deterministic part of the model (1). The basic idea was first described in [22]. The main feature of the N4SID class of techniques is the determination of the state sequence of the LTI system to be identified, or of an observer to reconstruct its state sequence, via the intersection of the row spaces of the Hankel-like matrices constructed from "past" and "future" I/O data. The basic idea was described in [12], and extended in [20, 21]. See also the recent survey [3].

Both MOESP and N4SID associated algorithms start by building a block-Hankel-block matrix H, from two concatenated block-Hankel matrices, using the available input-output data, and finding an upper triangular factor from a QR (or RQ) factorization of H. MATLAB¹ and Fortran codes implementing these approaches have been developed, e.g., [19, 23, 8, 18]. The Fortran codes described in [18, 13, 14] include an option to process the input-output data either in a single, or in multiple data "batches", which is important from a practical point of view. These two strategies are referred to as non-sequential, and sequential data processing, respectively.

New codes, implementing the subspace methods for linear system identification, have been written for the SLICOT Library [2], in the framework of the NICONET project. SLICOT is based on the linear algebra package LAPACK [1], and the BLAS collections [10, 6, 5]. The new Fortran routines include algorithmic improvements and refinements over the previous implementations. For instance, it is possible to exploit the particular structure of the block-Hankel-block matrix H, when computing the R factor of its QR factorization, and this could speed-up the calculations by more than an order of magnitude. Special problem structure is exploited whenever feasible. In addition, more flexibility is offered. The SLICOT collection allows to combine the system identification techniques, as proposed in [7]. For instance, it is possible to estimate the matrices A and C using the MOESP approach, and the matrices B and D using either the N4SID approach, or an alternate technique, which resorts again to the input and output records (Subsection 2.4).

In order to increase the user-friendliness, MATLAB interfaces, represented by either mexfiles

¹MATLAB is a trademark of The MathWorks, Inc.

or m-files, are provided. These interfaces have been designed to allow almost the same flexibility as the underlying Fortran routines, but many parameters are represented as options with default values. The MATLAB interfaces have been also used to build similar Scilab [4] interfaces to the SLICOT identification routines.

Basic algorithmic, implementation and numerical issues involved in the SLICOT subspace-based linear multivariable system identification toolbox—SLIDENT—are described in this report. Two numerical techniques for data compression are investigated and compared: the standard QR factorization of the matrix H, and the Cholesky factorization of the associated (inter-)correlation matrix, built exploiting the block-Hankel structure. In addition, a fast QR factorization technique [9], based on the displacement structure, is incorporated in the toolbox. Then, the associated Fortran routines are described and their functional abilities are outlined. The report also summarizes the functions and the structure of the MATLAB/Scilab interface of the toolbox, and presents some examples of use and numerical results obtained using the toolbox components. The comparisons with the available MATLAB identification tools show that SLICOT identification toolbox delivers equally accurate results, but often with much higher efficiency than other toolboxes.

2 SLICOT Identification Toolbox: Algorithmic Outline

2.1 Input-Output Data Processing using Subspace Techniques

The data used by the SLICOT toolbox for subspace identification is basically given as two $t \times m$ and $t \times \ell$ matrices, U and Y, for inputs and outputs, respectively, where t is the number of available measurements. So, the rows in U and Y are all the inputs or outputs measured at a certain time. Specifically,

$$U = \begin{bmatrix} u_1^T \\ u_2^T \\ \vdots \\ u_t^T \end{bmatrix}, \qquad Y = \begin{bmatrix} y_1^T \\ y_2^T \\ \vdots \\ y_t^T \end{bmatrix}.$$

In some applications, the number t is too large, and all data samples cannot be processed as a single batch. For instance, the available memory space could not be sufficient to accommodate for all these data, or the data themselves come in batches. Therefore, provisions have been taken to enable sequential processing of the input/output measurements. In this case, each new data batch is processed by updating a compressed representation of all the data in the previous batches.

For standard non-sequential data processing, the $N \times (m+\ell)p$ block-Hankel-block matrix $H = \begin{bmatrix} U_{1,p,N}^T & Y_{1,p,N}^T \end{bmatrix}$ is constructed, where N denotes the total number of samples that can be used, N = t - p + 1, and $U_{1,p,N}$ and $Y_{1,p,N}$ are block-Hankel matrices defined in terms of the

input and output data, respectively, i.e.,

$$U_{1,p,N} = \begin{bmatrix} u_1 & u_2 & u_3 & \cdots & u_N \\ u_2 & u_3 & u_4 & \cdots & u_{N+1} \\ u_3 & u_4 & u_5 & \cdots & u_{N+2} \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ u_p & u_{p+1} & u_{p+2} & \cdots & u_{N+p-1} \end{bmatrix},$$

and similarly for Y. Each u_i has m entries, and each y_i has ℓ entries. In both MOESP with past inputs and outputs and N4SID approaches, p is taken as 2s, where s is the "number of block rows" considered for both the "past" or the "future" parts. A QR factorization, H=QR, is used for data compression, but the matrix Q is not needed subsequently; then a singular value decomposition (SVD) of a certain matrix reveals the order n of the system as the number of "non-zero" singular values. (For instance, the MOESP approach finds the SVD of the submatrix $R_{ms+1:(2m+\ell)s,(2m+\ell)s+1:2(m+\ell)s}$, while the N4SID approach is more involved, and it first finds an oblique projection.) System matrices are computed from the right singular vectors, and other submatrices of R. The covariance matrices are found using the residuals of a least squares problem. The Kalman gain is obtained by solving a discrete-time algebraic matrix Riccati equation.

Therefore, for both MOESP and N4SID approaches considered, the matrix whose triangular factor R is needed is the $N \times 2(m+\ell)s$ matrix $H = \begin{bmatrix} U_{1,2s,N}^T & Y_{1,2s,N}^T \end{bmatrix}$, where N = t-2s+1. The triangular factor itself is a $2(m+\ell)s \times 2(m+\ell)s$ matrix.

For sequential data processing, the QR factorization is done sequentially, by updating the upper triangular factor R. When all data have been compressed, the system order and system matrices are computed as in the previous case. For convenience, consider that there are only two data batches, denoted as (U_1, Y_1) and (U_2, Y_2) , hence

$$U = \left[\begin{array}{c} U_1 \\ U_2 \end{array} \right], \qquad Y = \left[\begin{array}{c} Y_1 \\ Y_2 \end{array} \right],$$

with the number of rows N_1 and N_2 (possibly the same), and let H_1 and H_2 be the block-Hankel-block matrices, corresponding to (U_1, Y_1) and (U_2, Y_2) , respectively.

The algorithms outlined below are also summarized in [15, 17].

2.2 QR Factorization

The QR-based algorithms first compress the data (U_1, Y_1) . One computes the QR factorization of H_1 , $H_1 = Q_1R_1$, and then the upper triangular factor is updated using a specialized QR factorization for the problem

$$\left[\begin{array}{c} R_1 \\ H_2 \end{array}\right] = QR,\tag{2}$$

where the upper triangular structure of R_1 is exploited. Indeed, QR is a QR factorization of $H = \begin{bmatrix} H_1^T & H_2^T \end{bmatrix}^T$, as, with Q and Q_1 orthogonal matrices, we have

$$H^T H = H_1^T H_1 + H_2^T H_2 = R_1^T R_1 + H_2^T H_2 = R^T R.$$

The Q matrices are not needed in the subspace identification algorithms. The upper triangular factor R_1 is computed by a standard QR factorization, while for (2) one uses $2(m + \ell)s$ Householder transformations of the same order, $1 + N_2$; each Householder transformation i annihilates all the elements of the i-th column of H_2 , modifying the (i,i) element of R_1 . If there are additional data batches, the same procedure is applied repeatedly for each new data batch.

Given the input-output data in a batch j, (U_j, Y_j) , the computational routine has to build the appropriate segment H_j of the global block-Hankel-block matrix, H. This is done using a column-oriented process, exploiting the block-Hankel structure. The first (or single) segment H_1 is built either in the array R, which will store the triangular factor R, or in a workspace array, W, if R is not large enough. The remaining segments, if any, are built in the workspace, because R is anyhow needed for R. Efforts have been made to maximize the speed of the algorithmic process in the memory space provided. For instance, H_1 is constructed in R, if R is large enough. Otherwise, H_1 is constructed in W, its factor R_1 is computed in W and finally moved in R. If the workspace is not large enough to accommodate the entire data batch, then inner-loop sequential processing is additionally used, and the maximal number of rows that can be accommodated in the workspace are compressed in each inner cycle. Usually, the process could be more efficient when processing large data blocks at a time, due to the reduction in the overhead, but the optimal performance is, however, dependent on several factors, such as: use of an optimized BLAS library, cache memory size, block sizes for the underlying LAPACK block algorithms.

One complication appears when processing multiple data batches which are "connected", that is, which cannot be dealt with independently. When building the initial block-Hankel-block matrix H_1 , the last row will have the index N_1-2s+1 , because, due to the Hankel structure, this row will include u_{N_1} and y_{N_1} as the last entries in the U_1 and Y_1 submatrices, respectively. A larger index cannot be used, as the corresponding entries will not be available in U_1 and Y_1 . Therefore, the entries $u_{N_1-2s+2},\ldots,u_{N_1}$ will not appear in the first column of H_1 for the N4SID algorithm. Similarly, the last 2s-1 entries of U_2 will not appear in the first column of H_2 . However, when the data batches are "connected", these last 2s-1 elements in the current data batch should be added in front of the elements of the next data batch, in order to obtain the same triangular factor as when all batches would be processed in a single pass. Connecting the batches has been implemented by storing the $(m+\ell)(2s-1)$ "connection" elements in the workspace at the end of the processing phase of a batch (provided it is not the last), and restoring them at the beginning of the next phase. This way, all data are used and no information is omitted. No saving-restoring operations exist when the data batches are not connected.

2.3 Cholesky Factorization

In order to use the Cholesky factorization algorithm for data compression, one should first build the inter-correlation matrix $G = H^T H$, and then factor G, assuming it is positive definite (which is usually the case in practice, due to various noise components). For the N4SID approach, the

block-Hankel matrix H_u , corresponding to the inputs, is

$$H_{u} = \begin{bmatrix} u_{1}^{T} & u_{2}^{T} & \cdots & u_{2s}^{T} \\ u_{2}^{T} & u_{3}^{T} & \cdots & u_{2s+1}^{T} \\ \vdots & \vdots & \vdots & \vdots \\ u_{N}^{T} & u_{N+1}^{T} & \cdots & u_{N+2s-1}^{T} \end{bmatrix},$$

$$(3)$$

where N=t-2s+1; the block-Hankel matrix H_y , corresponding to the outputs, is written similarly, and $H=\begin{bmatrix} H_u & H_y \end{bmatrix}$. For the MOESP approach, the submatrix built from the last ms columns of H_u (the "future" inputs) comes in front of the submatrix built from the first ms columns of H_u (the "past" inputs part). Similarly, for identification with multiple batches, H_j is defined as H above, but with appropriate values for the inputs and outputs, as well as for the number of samples used.

The algorithm computes the symmetric (block) matrix $G = H^T H$ exploiting the block-Hankel structure. Let G_{uu} , G_{uy} , and G_{yy} denote the subblock-matrices of G corresponding to the input, input-output, and output (inter-)correlations, respectively, hence $G_{uu} = H_u^T H_u$, $G_{uy} = H_u^T H_y$, and $G_{yy} = H_y^T H_y$. Each subblock-matrix G_{uu} , G_{uy} , or G_{yy} , consists of $2s \times 2s$ submatrices of orders $m \times m$, $m \times \ell$, and $\ell \times \ell$, respectively. Denote $G_{uu}^{i,j}$ the (i,j)-th submatrix of G_{uu} , and similarly for G_{uy} and G_{yy} . The nice problem structure can be exploited [15, 17]. The formula for computing $G_{uy}^{i,j}$ can be written as follows

$$G_{uu}^{i,j} = \hat{G}_{uu}^{i,j} + u_i u_i^T + u_{i+1} u_{i+1}^T + \dots + u_{i+N-1} u_{i+N-1}^T,$$

$$\tag{4}$$

where u_i^T is the *i*-th row of U, j = 1:2s, i = 1:j, and $\hat{G}_{uu}^{i,j}$ is a zero matrix if the first (or single) data batch is computed, or it is the matrix $G_{uu}^{i,j}$ computed till the current batch, otherwise. Each matrix $G_{uu}^{j,j}$ is symmetric. The upper triangle of G_{uu} is computed or updated column-wise in the array R, that is, the block matrices are processed in the order $G_{uu}^{1,1}, G_{uu}^{1,2}, G_{uu}^{2,2}, \ldots, G_{uu}^{2s,2s}$. Only the submatrices of the first block-row are fully computed or updated with (4), i.e., (4) is used for i = 1 only. The remaining submatrices are determined exploiting the block-Hankel structure, using the updating formula, which follows from (4),

$$G_{uu}^{i+1,j+1} = \hat{G}_{uu}^{i+1,j+1} - \hat{G}_{uu}^{i,j} + G_{uu}^{i,j} + u_{i+N} u_{i+N}^T - u_i u_i^T,$$

$$(5)$$

applied for j=1:2s-1, and for i=2:j, for a given j; only the upper triangular part is evaluated for i=j. The submatrix $G_{uu}^{1,1}$ is computed using a BLAS 3 operation SYRK, the submatrices $G_{uu}^{1,2:2s}$ are computed using BLAS 3 operations GEMM, and the other submatrices are obtained using BLAS 2 and BLAS 1 updating operations GERR and AXPY, respectively.

The subblock-matrix G_{yy} is computed similarly. Part of the subblock-matrix G_{uy} is got using the formula

$$G_{uy}^{i,j} = \hat{G}_{uy}^{i,j} + u_i y_j^T + u_{i+1} y_{j+1}^T + \dots + u_{i+N-1} y_{j+N-1}^T,$$
(6)

where y_j^T is the j-th row of Y, and either i=1 and j=1:2s, or i=1:2s and j=1, i.e., only the submatrices of the first block-row and block-column are fully computed (or updated) using (6). The remaining ones are determined exploiting the block-Hankel structure, using the updating formula

$$G_{uy}^{i+1,j+1} = \hat{G}_{uy}^{i+1,j+1} - \hat{G}_{uy}^{i,j} + G_{uy}^{i,j} + u_{i+N} y_{j+N}^T - u_i y_j^T, \tag{7}$$

for j = 1:2s - 1, and i = 1:2s - 1. The input-output correlations, G_{uy} , are computed (or updated) column-wise in the array R.

At the final batch for the MOESP approach, the past and future parts of the input submatrix H_u are interchanged after their computation. (Specifically, the upper triangular parts are interchanged, and the (1,2) part is transposed in-situ.) This scheme is simpler and clearer than a directly implemented MOESP calculation.

Finally, the Cholesky factorization algorithm is applied on the computed correlation matrix G. In the rare case when this algorithm fails, the QR factorization is automatically used, for non-sequential processing. This is not possible for sequential processing because not all the data are available for the routine. The calculations could, however, be restarted under user's control.

The SLICOT toolbox includes an option for computing R by a fast QR algorithm, based on displacement rank techniques. The generators of the matrix H^TH are computed and then used to obtain R. Details are given in [9]. The numerical properties and the efficiency of this algorithm are similar to those of the Cholesky factorization algorithm.

2.4 Computation of System Matrices

The other computational steps of the subspace-based identification procedures have been also analyzed for possible improvements, by exploiting any existing structure. For instance, an essential intermediate calculation in the N4SID algorithm is the determination of the so-called "oblique projection". Partition the factor R as $R = [U_p \ U_f \ Y_p \ Y_f]$, where the subscripts p and f stand for "past" and "future" data, respectively, and the four blocks have ms, ms, ℓs , and ℓs columns, respectively. Define $W_p = [U_p \ Y_p]$, and consider the residuals of the two least squares problems giving the oblique projection,

$$r_1 = W_p - U_f X_1, \qquad r_2 = Y_f - U_f X_2,$$

where X_1 and X_2 are the minimum norm least squares solutions of the problems

$$\min \|U_f X - W_p\|_2$$
, $\min \|U_f X - Y_f\|_2$,

respectively. Then, the oblique projection can be computed as $O = r_2^T Q_1 Q_1^T$ [13, 14], where Q_1 consists of the first $k = \text{rank}(r_1)$ columns of the orthogonal matrix in the QR factorization of r_1 . No least squares problems should be actually solved. The advantage is that only orthogonal transformations are used for finding O, and therefore, the problem conditioning is preserved. Both problems have the same coefficient matrix, U_f , and it consists of two $ms \times ms$ submatrices, the second one being upper triangular; this structure is exploited.

A fast algorithm for computing the matrices B and D, using a structure exploiting QR factorization, has been described in [13, 14]. Essentially, the algorithm solves the following problem

$$\begin{bmatrix} Q_{1s} & Q_{1,s-1} & \cdots & Q_{12} & Q_{11} \\ 0 & Q_{1s} & \cdots & Q_{13} & Q_{12} \\ 0 & 0 & \cdots & Q_{14} & Q_{13} \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & \cdots & 0 & Q_{1s} \end{bmatrix} \begin{bmatrix} \Gamma_{-} & 0 \\ 0 & I_{\ell} \end{bmatrix} \begin{bmatrix} B \\ D \end{bmatrix} = \begin{bmatrix} K_{1} \\ K_{2} \\ K_{3} \\ \vdots \\ K_{s} \end{bmatrix},$$

where Q_{1i} is $(\ell s - n) \times \ell$, for i = 1: s, Γ_{-} is $(\ell s - \ell) \times n$ (the leading part of the first n singular vectors), and K_i is $(\ell s - n) \times m$, i = 1: s. The first matrix product in the left-hand side is not computed, but the first matrix is fastly triangularized exploiting its structure. The matrices B and D are then found in two steps. Unfortunately, this fast algorithm finds least squares solutions in special cases only, namely, when the second matrix in the left hand side is square and nonsingular. The results obtained are often sufficiently accurate, especially for the MOESP technique. However, in some applications, the true least squares solutions should be computed using a less efficient algorithm, based on Kronecker products calculations; this algorithm is also available in the toolbox, and it solves a problem having half the size of the corresponding problem solved by the N4SID code. Extensive testing has shown that the use of the total least squares (TLS) solutions, as described in [13, 14], could provide less accurate results; consequently, the TLS strategy has been omitted in the toolbox.

The N4SID algorithm incorporated in the SLICOT toolbox differs from the published algorithm [19, 21] in many aspects and details. However, if few changes are made in the original algorithm, they will essentially produce the same matrices A, B, C, D; there could be differences in the signs of elements, but the systems computed by N4SID and SLICOT are related by a similarity transformation. The changes are:

- 2.4.1 SLICOT routines do not scale the input and output data by $1/\sqrt{t-2s+1}$; this scaling is actually needed only for computing the covariance matrices, when it can be performed in a faster way (on a part of the processed R factor).
- 2.4.2 The first n singular vectors are not scaled by the corresponding singular values.
- 2.4.3 The extended observability matrix Γ is not recomputed from the estimated A and C; extensive testing has shown that this strategy could sometimes slightly improve the accuracy, but equally well, diminish it.

An alternate algorithm is included for the computation of the matrices B and D. An extension and refinement of the method in [18, 24] is used. Specifically, denoting

$$X = \begin{bmatrix} \operatorname{vec}(D^T)^T & \operatorname{vec}(B)^T & x_0^T \end{bmatrix}^T,$$

where $\operatorname{vec}(M)$ is the vector obtained by stacking the columns of the matrix M, then X is the least squares solution of the system $SX = \operatorname{vec}(Y)$, with the matrix $S = \begin{bmatrix} \operatorname{diag}(U) & W \end{bmatrix}$, defined by

where $\operatorname{diag}(U)$ has ℓ block rows and columns. In this formula, y^{ij} are the outputs of the system for zero initial state, computed using the following model, for j=1:m, and for i=1:n,

$$x^{ij}(k+1) = Ax^{ij}(k) + e_i u_j(k), x^{ij}(1) = 0,$$

 $y^{ij}(k) = Cx^{ij}(k),$ (9)

where e_i is the i-th n-dimensional unit vector, Γ is given by

$$\Gamma = \begin{bmatrix} C \\ CA \\ CA^2 \\ \vdots \\ CA^{t-1} \end{bmatrix}, \tag{10}$$

and P is a permutation matrix that groups together the rows of Γ depending on the same row c_i of C, namely

$$\left[\begin{array}{c}c_j\\c_jA\\c_jA^2\\\vdots\\c_jA^{t-1}\end{array}\right],$$

for $j=1:\ell$. The first block column, $\operatorname{diag}(U)$, is not explicitly constructed, but its structure is exploited. The last block column is evaluated using powers of A with exponents 2^k , but no nonnecessary powers are computed (ℓt is not expanded to a power of 2). No interchanges are applied. A special QR decomposition of the matrix S is computed. Let $U=q\begin{bmatrix} r^T & 0 \end{bmatrix}^T$ be the QR decomposition of U, if m>0, where r is an $m\times m$ upper triangular matrix. Then, $\operatorname{diag}(q^T)$ is applied to W and $\operatorname{vec}(Y)$. The block-rows of S and $\operatorname{vec}(Y)$ are implicitly permuted so that the matrix S becomes

$$\left[\begin{array}{cc} \operatorname{diag}(r) & W_1 \\ 0 & W_2 \end{array}\right],$$

where W_1 has ℓm rows. Then, the QR decomposition of W_2 is computed (sequentially, if m>0) and used to obtain B and x_0 . The intermediate results and the QR decomposition of U are needed to find D. If a triangular factor is too ill conditioned, then singular value decomposition (SVD) is employed. SVD is not generally needed if the input sequence is sufficiently persistently exciting and t is large enough. If the matrix W cannot be stored in the workspace provided, the QR decompositions of W_2 and U are computed sequentially. The calculations simplify if D and/or x_0 are not needed.

The algorithm discussed above for computing the system matrices B and D, implemented in the SLICOT Library routine IB01QD (called by the driver IB01CD), is usually less efficient than the MOESP or N4SID algorithms implemented in SLICOT Library routine IB01PD (called by the driver IB01BD), because a large least squares problem has to be solved. The advantage is that the accuracy is better, since the computed matrices B and D are fitted to the input and output trajectories. However, if matrix A is unstable, the computed matrices B and D could be inaccurate.

When only x_0 should be determined, given the system matrices A, B, C, D, and the inputoutput sequences, a simpler method is implemented in the SLICOT Library routine IB01RD (also called by the driver IB01CD). This is an extension and refinement of the method in [24]. Specifically, the output $y_0(k)$ of the system for zero initial state is computed for k = 1:t using the given model. Then, the following least squares problem is solved for x_0

$$\Gamma x_0 = \begin{bmatrix} y(1) - y_0(1) \\ y(2) - y_0(2) \\ \vdots \\ \vdots \\ y(t) - y_0(t) \end{bmatrix}. \tag{11}$$

The coefficient matrix Γ is evaluated using powers of A with exponents 2^k . The QR decomposition of Γ is computed. If its triangular factor R is too ill conditioned, then singular value decomposition of R is used. If the matrix Γ cannot be stored in the workspace provided, the QR decomposition is computed sequentially. Various particular cases are dealt with by special code portions, in order to increase the efficiency.

Structure exploiting algorithms and specialized linear algebra algorithms, performing all the processing steps of the MOESP and N4SID approaches, are used whenever possible. For instance, for computing the covariance matrices, the residuals of the linear least squares problem MV=N should be computed, where

$$M = \begin{bmatrix} \begin{bmatrix} R_{11} \\ 0 \end{bmatrix} & Y \\ 0 & 0 \end{bmatrix}, \quad N = \begin{bmatrix} R_{1:3\tilde{\ell},4} & X \end{bmatrix}, \tag{12}$$

the matrices M, N, Y and X have dimensions $((2m+\ell)s+\ell, ms+n)$, $((2m+\ell)s+\ell, \ell+n)$, $((2m+\ell)s, n)$, and $((2m+\ell)s+\ell, n)$, respectively, $R_{1:3\widetilde{\ell},4} = R_{1:(2m+\ell)s+\ell,(2m+\ell)s+1:(2m+\ell)s+\ell}$, and $R_{11} = R_{1:ms,1:ms}$ is upper triangular. The structure of this problem is fully exploited. Note also that the two block columns in M have been interchanged, compared to what would directly appear, in order to gain in efficiency. For the N4SID approach, the solution of this system is also needed, as it directly gives the system matrices A and C, and another matrix used for finding B and D.

3 SLICOT Identification Toolbox: Software Components

The SLICOT system identification toolbox consists in 13 user-callable and computational routines, three MATLAB interfaces (mexfiles), and over ten MATLAB m-files. The toolbox software components are also partly summarized in [16].

3.1 Fortran Routines

The Fortran routines are listed in Table 1, together with their function; the first three are driver routines. Together, these routines totalize about 0.4 Mb of source code.

Table 1: SLICOT system identification routines.

Routine	Function
IBO1AD	preprocesses the input-output data for estimating the matrices of a linear
	time-invariant dynamical system and finds an estimate of the system order.
	The input-output data can, optionally, be processed sequentially.
IBO1BD	estimates the system matrices A, C, B , and D , the noise covariance matri-
	ces Q , R_y , and S , and the Kalman gain matrix K of a linear time-invariant state space model, using the processed triangular factor R of the concate-
	nated block-Hankel matrices, provided by SLICOT Library routine IB01AD.
IBO1CD	estimates the initial state and, optionally, the system matrices B and D
	of a linear time-invariant discrete-time system, given the system matrices
	(A, B, C, D), or only the matrix pair (A, C) , respectively, and the input
	and output trajectories of the system.
IBO1MD	computes the upper triangular factor in the QR factorization of the block-
	Hankel matrix built from input-output data for subspace identification.
IBO1MY	computes the upper triangular factor in the QR factorization of the block-
IBO1ND	Hankel matrix, using a fast QR algorithm. finds the singular value decomposition giving the system order, using the
IDOIND	triangular factor of the concatenated block-Hankel matrices.
IBO10D	finds the system order, using the singular value decomposition.
IBO10Y	asks for user's confirmation of the system order found.
IBO1PD	estimates the system matrices and covariance matrices of a linear time-
	invariant state space model.
IBO1PX	computes the matrices B and D of a linear time-invariant state space model using Kronecker products.
IBO1PY	computes the triangular QR factor of a structured matrix and estimates
	the matrices B and D of a linear time-invariant state space model.
IBO1QD	estimates the initial state and the system matrices B and D of a linear time-
	invariant state space model, given the matrix pair (A, C) and the input and output trajectories.
IBO1RD	estimates the initial state of a linear time-invariant state space model, given
	the system matrices (A, B, C, D) and the input and output trajectories.

The calculations are organized so that much flexibility is achieved. For instance, it is possible to compute the A and C matrices via the MOESP approach, and the B and D matrices either via the N4SID approach, or by using the input-output data again (Subsection 2.4).

The design of these subroutines aimed to provide the user the most convenient and flexible way to perform the identification calculations. However, there is no unique driver to estimate the system matrices starting from the input-output records. The calculations have been split in two or three parts.

The first part compresses the input-output data and delivers a processed upper triangular factor R of the block-Hankel-block matrix H (in an array R), the singular values defining the

system order, and an estimate of the system order, n. These calculations are performed by the driver IBO1AD, which calls the computational routines IBO1MD, IBO1ND, IBO10D, and indirectly IBO1MY (by IBO1MD), and IBO10Y (by IBO10D). The numerical values returned in the array R by the driver IBO1AD do not represent the factor R, but an appropriate input for the next calculations. If the user actually needs R in a particular application, the lower-level computational routine IBO1MD should be directly called. The available options of the driver IBO1AD, which can be selected using character strings as actual arguments in the call, 2 are:

METH Specifies the subspace identification method to be used, as follows:

- = 'M': MOESP algorithm with past inputs and outputs;
- = 'N': N4SID algorithm.

ALG Specifies the algorithm for computing the triangular factor R, as follows:

- = 'C': Cholesky algorithm applied to the correlation matrix of the input-output data;
- = 'F': Fast QR algorithm;
- = 'Q': QR algorithm applied to the concatenated block Hankel matrices.

JOBD Specifies whether or not the matrices B and D should later be computed using the MOESP approach, as follows:

- = 'M': the matrices B and D should later be computed using the MOESP approach;
- = 'N': the matrices B and D should not be computed using the MOESP approach.

This parameter is not relevant for METH = 'N'.

- BATCH Specifies whether or not sequential data processing is to be used, and, for sequential processing, whether or not the current data block is the first block, an intermediate block, or the last block, as follows:
 - = 'F': the first block in sequential data processing;
 - = 'I': an intermediate block in sequential data processing;
 - = 'L': the last block in sequential data processing;
 - = '0': one block only (non-sequential data processing).

NOTE that when 100 cycles of sequential data processing are completed for BATCH = 'I', a warning is issued, to prevent for an infinite loop.

- CONCT Specifies whether or not the successive data blocks in sequential data processing belong to a single experiment, as follows:
 - = 'C': the current data block is a continuation of the previous data block and/or it will be continued by the next data block;
 - = 'N': there is no connection between the current data block and the previous and/or the next ones.

This parameter is not used if BATCH = 'O'.

- CTRL Specifies whether or not the user's confirmation of the system order estimate is desired, as follows:
 - = 'C': user's confirmation;
 - = 'N': no confirmation.

If CTRL = 'C', a reverse communication routine, IBO10Y, is indirectly called (by SLICOT

²One could use, for instance, 'MOESP' for the parameter METH, to improve the readability of the code.

Library routine IBO10D), and, after inspecting the singular values and system order estimate, n, the user may accept n or set a new value. IBO10Y is not called if CTRL = 'N'.

The reason for including the option JOBD is just to skip some computations when they are not necessary. The lower-level routine IBO10Y is actually a simple template, which the user could call, or easily modify, in order to tailor it for a specific application. The other options are self-explanatory.

The second part uses the results delivered by IBO1AD and computes the system matrices, (A,B,C,D), the covariance matrices (Q,R_y,S) (representing the state, output, and state-output (cross-)covariance matrices, respectively), and the Kalman gain matrix. These calculations are performed by the driver IBO1BD, which calls the computational routine IBO1PD, and this, in turn, calls IBO1PX (for the N4SID approach) and IBO1PY (for the standard MOESP approach). The available options of the driver IBO1BD are:

METH Specifies the subspace identification method to be used, as follows:

- = 'M': MOESP algorithm with past inputs and outputs;
- = 'N': N4SID algorithm;
- = 'C': combined method: MOESP algorithm for finding the matrices A and C, and N4SID algorithm for finding the matrices B and D.

JOB Specifies which matrices should be computed, as follows:

- = 'A': compute all system matrices, A, B, C, and D;
- = 'C': compute the matrices A and C only;
- = 'B': compute the matrix B only;
- = 'D': compute the matrices B and D only.

JOBCK Specifies whether or not the covariance matrices and the Kalman gain matrix are to be computed, as follows:

- = 'C': the covariance matrices only should be computed;
- = 'K': the covariance matrices and the Kalman gain matrix should be computed;
- = 'N': the covariance matrices and the Kalman gain matrix should not be computed.

The ability to compute only part of the results, using the options JOB and JOBCK, is important in some applications. For instance, typical calculations performed by the Leuven Measurement Systems (LMS) International software, for mechanical systems, involve the computation of the matrices A and C for an increasing sequence of system orders; the poles (the eigenvalues of A) are computed for each system order, and stabilization diagrams are used to decide which poles actually correspond to physical poles, and only these poles are retained and used to build the final A and C matrices; then, the corresponding matrices B and D are computed using the approach implemented in the third SLICOT driver, IBO1CD.

The third part uses the matrices A and C delivered by IBO1BD and the input-output trajectories, $\{u_k\}$ and $\{y_k\}$, k=1:t, and estimates the matrices B and D, and/or the initial state of the system, x_0 , denoted as $\mathbf{x}(0)$ below. These calculations are performed by the driver IBO1CD, which calls the computational routine IBO1QD, if both the matrices B, D, and the initial state x_0 are desired, or IBO1RD, if x_0 only is desired. For the second case, the matrices B and D should also be provided at input to the driver. The available options of the driver IBO1CD are:

	Table 2: Driv	er IBO1CD: c	ptions, 1	the data	used, a	and the	returned	results.
--	---------------	--------------	-----------	----------	---------	---------	----------	----------

JOBXO	COMUSE	JOB	Data used	Returned results
Х	С	В	A, C, u, y	x, B
X	C	D	A,C,u,y	x, B, D
N	C	В	A,C,u,y	x = 0, B
N	C	D	A,C,u,y	x = 0, B, D
Х	U	В	A,B,C,u,y	x
X	U	D	A, B, C, D, u, y	x
N	U	*	_	x = 0
Х	N	*	A, C, y	x
N	N	*	_	_

JOBXO Specifies whether or not the initial state should be computed, as follows:

- = 'X': compute the initial state x_0 ;
- = 'N': do not compute the initial state (possibly, because x_0 is known to be zero).

COMUSE Specifies whether the system matrices B and D should be computed or used, as follows:

- = 'C': compute the system matrices B and D, as specified by JOB;
- = 'U': use the system matrices B and D, as specified by JOB;
- = 'N': do not compute/use the matrices B and D.

If JOBXO = 'N' and $COMUSE \iff 'N'$, then x_0 is set to zero.

If JOBXO = 'N' and COMUSE = 'N', then x_0 is neither computed nor set to zero.

- JOB If COMUSE = 'C' or 'U', specifies which of the system matrices B and D should be computed or used, as follows:
 - = 'B': compute/use the matrix B only (D is known to be zero);
 - = 'D': compute/use the matrices B and D.

The value of JOB is irrelevant if COMUSE = 'N' or if JOBXO = 'N' and COMUSE = 'U'.

The combinations of options, the data used, and the returned results, are given in Table 2, where '*' denotes an irrelevant value. For JOBXO = 'N' and COMUSE = 'N', the driver IBO1CD does not perform any calculations, but just returns.

3.2 Interfaces: Mexfiles and m-files

The SLICOT system identification routines can be used from a MATLAB environment using the interfaces described in this section. The general framework of integrating the SLICOT Fortran routines in a MATLAB-like environment is described in [11]. Both mexfiles and m-files are provided in the SLIDENT toolbox. The mexfiles correspond to the Fortran drivers. Their interface is more complicated than that of the m-files, but they provide more flexibility

and generality. The m-files call the mexfiles and are included for user's convenience. On line information about the mexfiles and m-files can be obtained in the usual manner, e.g., by typing help sident, for getting details about the mexfile sident. The same interfaces could also be used in a Scilab environment [4]. The mexfiles are listed in Table 3.

Table 3: SLICOT system identification: mexfile interfaces to MATLAB/Scilab.

Mexfile	Function
order	preprocesses the input-output data for estimating the matrices of a linear
	time-invariant dynamical system and finds an estimate of the system order.
	The input-output data can, optionally, be processed sequentially.
sident	computes a state-space realization (A, B, C, D) , the Kalman predictor gain
	K, and the covariances of a discrete-time system, given the system order
	and the relevant part of the R factor of the concatenated block-Hankel ma-
	trices, using subspace identification techniques (MOESP, N4SID, or their
	combination).
findBD	estimates the initial state x_0 and/or the matrices B and D of a discrete-
	time system, given the system matrices A, C , and possibly B, D , and the
	input and output trajectories $\{u_k\}$, $\{y_k\}$, of the system.

The mexfiles and m-files include several options which could be of interest for the user; all these options have default values, and hence the calling statements could be very simple. One such option is printw, a switch for printing the warning messages; it should be set to 1 to print these messages, or to 0, otherwise (the default).

The calling sequences for the three mexfiles are shown below.

```
[R,n(,sval)(,rcnd)] = order(meth,alg,jobd,batch,conct,s,Y(,U,tol, printw,ldwork,R));
[(A,C)(,B(,D))(,K(,Q,Ry,S))(,rcnd)] = sident(meth,job,s,n,l,R(,tol,t, A,C,printw));
[(x0)(,B(,D))(,V)(,rcnd)] = findBD(jobx0,comuse(,job),A(,B),C(,D),Y (,U,tol,printw,ldwork));
```

where the parameters put inside the brackets are optional, and their compulsory grouping is also indicated. Most of these parameters have clear meaning, or details have been already given in Subsection 3.1. The input data argument U is optional in the mexfiles and m-files which could involve it. The parameter ldwork in the mexfiles order and findBD allows the user to specify other sizes than the default ones for working arrays. Usually, the SLICOT mexfiles are using exactly the sizes needed for the arrays, but providing larger sizes for working arrays could improve the efficiency, e.g., by enabling the calls of the LAPACK block algorithms. The influence of the workspace length, when Fortran, not optimized BLAS routines are used, is shown in Figure 1 in Section 5. The parameter rcnd returns the reciprocal condition numbers, and, possibly, an error bound for the Riccati equation solution (needed for computing the Kalman gain); these accuracy indicators are useful in assessing the reliability of the computed results.

The basic MATLAB m-files which call the mexfiles are presented in Table 4. The *system* object defined in the MATLAB Control Toolbox is used, whenever possible.

Table 4: SLICOT system identification: m-file interfaces.

m-file	Function
order	help function for order mexfile.
sident	help function for sident mexfile.
findBD	help function for findBD mexfile.
findR	preprocesses the input-output data for estimating the matrices of a linear
	time-invariant dynamical system, using Cholesky or (fast) QR factorization
	and subspace identification techniques (MOESP or N4SID), and estimates
	the system order.
findABCD	finds the system matrices and the Kalman gain of a discrete-time system,
	given the system order and the relevant part of the R factor of the concate-
C: 14 G	nated block-Hankel matrices, using MOESP, N4SID, or their combination.
findAC	finds the system matrices A and C of a discrete-time system, given the system order and the relevant part of the R factor of the concatenated
	block-Hankel matrices, using MOESP or N4SID.
findBDK	finds the system matrices B and D and the Kalman gain of a discrete-time
	system, given the system order, the matrices A and C , and the relevant part
	of the R factor of the concatenated block-Hankel matrices, using MOESP
	or N4SID.
findxOBD	estimates the initial state and/or the matrices B and D of a discrete-time
	linear system, given the (estimated) system matrices A, C , and a set of
	input/output data.
inistate	estimates the initial state of a discrete-time system, given the (estimated)
	system matrices, and a set of input/output data.

The calling sequences for the computational m-files (other than the three help files associated to the mexfiles) are listed below:

```
[R,n(,sval,rcnd)] = findR(s,Y(,U,meth,alg,jobd,tol,printw));
[sys(,K,Q,Ry,S,rcnd)] = findABCD(s,n,l,R(,meth,nsmpl,tol,printw));
[A,C(,rcnd)] = findAC(s,n,l,R(,meth,tol,printw));
[B(,D,K,Q,Ry,S,rcnd)] = findBDK(s,n,l,R,A,C(,meth,job,nsmpl,tol,printw));
[x0,B,D(,V,rcnd)] = findxOBD(A,C,Y(,U,withx0,withd,tol,printw));
[x0(,V,rcnd)] = inistate(sys,Y(,U,tol,printw));
```

The input parameter tol is a vector with two elements in case of findR, and a scalar, otherwise. tol(1) (or tol) is the tolerance for estimating the rank of matrices. If tol(1) > 0, the given value of tol(1) is used as a lower bound for the reciprocal condition number. tol(2) is the tolerance for estimating the system order. If tol(2) \geq 0, the estimate is indicated by the index of the last singular value greater than or equal to tol(2). (Singular values less than

tol(2) are considered as zero.) When tol(2) = 0, then s*eps*sval(1) is used instead tol(2), where sval(1) is the maximal singular value, and eps is the relative machine precision. When tol(2) < 0, the estimate is indicated by the index of the singular value that has the largest logarithmic gap to its successor. Default values are: tol(1) = prod(size(matrix))*eps; tol(2) = -1.

Shorter calls are, of course, possible, e.g.,

```
[sys,rcnd] = findABCD(s,n,1,R);
[B,D,rcnd] = findBDK(s,n,1,R,A,C);
[B,D] = findxOBD(A,C,Y,U,O);
x0 = inistate(A,C,Y);
x0 = inistate(A,B,C,Y,U);
```

The parameter sys is a discrete-time ss MATLAB system object, consisting in a state-space realization sys = (A, B, C, D).

Example 3.1 Assuming u, y already available, the simple MATLAB session shown below

produced the following results, for some given $\{u\}$ and $\{y\}$ data sequences,

Estimated system matrices

```
A =
  9.3206e-01
               2.9657e-01
                           3.5373e-02
                                        1.5266e-02
  -1.9087e-01
               9.2076e-01 -1.6224e-02
                                        5.1043e-02
  -1.3172e-02
               2.0301e-02
                           8.5844e-01 -4.5612e-01
  -1.5919e-03 -2.6322e-02
                           2.4633e-01
                                        7.6515e-01
B =
   1.2245e+00
              8.9248e-01
```

```
7.7178e-01 4.5673e-01
-4.1280e-01 4.1149e-01
2.3463e-01 -1.2672e-01

C =
2.2571e-01 -1.8491e-01 -3.5805e-01 -3.3156e-01
2.1009e-01 -2.6185e-01 2.7019e-01 3.1076e-01

D =
5.4624e-02 1.9648e-03
2.7838e-03 3.7158e-02
```

Relative output error 1-norm and reciprocal condition numbers

```
err =
    1.3073e-06

rcnd =
    9.6651e-01
    1.0000e+00
```

4.9869e-03 3.9396e-01

Four additional method-oriented m-files, briefly described in Table 5, have been included in the SLIDENT toolbox.

Table 5: SLICOT system identification: method-oriented m-file interfaces.

m-file	Function
slmoesp	finds the system matrices and the Kalman gain of a discrete-time system,
	using MOESP subspace identification technique.
sln4sid	finds the system matrices and the Kalman gain of a discrete-time system,
	using N4SID subspace identification technique.
slmoen4	finds the system matrices and the Kalman gain of a discrete-time system,
	using combined MOESP and N4SID subspace identification techniques: the
	matrices A and C are found via MOESP, and B and D , via N4SID.
slmoesm	finds the system matrices, the Kalman gain, and the initial state of a
	discrete-time system, using combined MOESP subspace identification and
	system simulation techniques: the matrices A and C are found via MOESP,
	B and D via simulation.

The calling sequences for these method-oriented m-files are listed below.

```
[sys(,K,rcnd,R)] = slmoesp(s,Y,U,n,alg,tol,printw);
[sys(,K,rcnd,R)] = sln4sid(s,Y,U,n,alg,tol,printw);
[sys(,K,rcnd,R)] = slmoesn(s,Y,U,n,alg,tol,printw);
[sys(,K,rcnd,x0,R)] = slmoesm(s,Y,U,n,alg,tol,printw);
```

If n=0, or n=[], or n is omitted from the input parameters, the user is prompted to provide its value, after inspecting the singular values, shown as a bar plot. If n<0, then n is determined automatically, according to tol(2). The optional output parameter R returns the processed upper triangular factor R of the block-Hankel-block matrix H, built from the input-output data. It can be used for fast identification of systems of various orders, using, for instance, the following commands:

```
[sys,K,rcnd,R] = sln4sid(s,Y,U,n0,alg,tol,printw);
for n = n0+1 : min( n0+nf, s-1 )
    [sys,K,rcnd] = sln4sid(s,Y,U,n,R,tol,printw);
end
```

The data values for Y and U are not used inside the loop (only the size of Y is needed), but R replaces alg. Clearly, the systems of various orders (from n0+1 to min(n0+nf, s-1)), should be used inside the loop. The elements rcnd(1) and rcnd(2) are set to 1 when sln4sid is called with R instead of alg.

Shorter calls are, of course, possible, e.g.,

```
[sys,K] = slmoesp(s,Y);
[sys,K] = slmoesp(s,Y,[],n);
sys = slmoesp(s,Y,U);
```

and similarly for the other m-files. The first two calls estimate the matrices A, C, and K of a stochastic system (with no inputs), for an order n found automatically, or prespecified, respectively.

The same numerical results as those obtained using findR and findABCD in Example 3.1 can be obtained using the following MATLAB code

4 Benchmarks for Identification

This topic is fully described in another NICONET report, entitled *Benchmarks for Identification*, available through the Niconet web page (Reports link), or by anonymous ftp from wgs.esat.kuleuven.ac.be, file pub/WGS/REPORTS/nic1999-19.ps.Z. This report also describes some rules for benchmarking in system identification. A valuable, publicly available collection of test data is the Database for Identification of Systems (DAISY), which is accessible from the website wgs.esat.kuleuven.ac.be/sista/daisy. This collection contains several datasets:

process industry systems, mechanical systems, biomedical systems, environmental systems, thermic systems, simulators, and time series. Some of the test data from DAISY have been used for testing the identification toolbox SLIDENT (see Section 5).

5 Industrial Examples

Besides the industrial examples from the DAISY data collection, several input-output sequences from the LMS International have been used to test the components of the SLIDENT toolbox. An effort is underway at LMS to incorporate the SLICOT routines for identification into their interactive software. The results obtained on some of the LMS industrial data will be presented in another Niconet report.

The rest of this section summarizes typical results obtained using the subspace-based techniques available in the SLIDENT toolbox. The data used is publicly available on the DAISY internet site:

http://www.esat.kuleuven.ac.be/sista/daisy

8

9

Ball and beam

Wall temperature

in order to increase the accessibility and the reproducibility. Specifically, Table 6 gives a summary description of the applications considered in this report, indicating the number of inputs m, outputs ℓ , block rows s, and data samples used t.

#	Application	m	ℓ	s	t
1	Glass tubes	2	2	20	1401
2	Labo Dryer	1	1	15	1000
3	Glass Oven	3	6	10	1247
4	Mechanical flutter	1	1	20	1024
5	Flexible robot arm	1	1	20	1024
6	Evaporator	3	3	10	6305
7	CD player arm	2	2	15	2048

1

2

1

20

20

1000

1680

Table 6: Summary description of applications.

Only one data batch has been used in all the tests whose results are given below. The numerical results have been obtained on either an IBM PC computer at 500 MHz, with 128 Mb memory, using Digital Visual Fortran V5.0, or on a Sun 4 sparc Ultra-2 computer, using operating system OS 5.6, and Sun WorkShop Compiler FORTRAN 77 5.0. In either cases, MATLAB 5.3.0.10183 (R11) has been employed.

In [15], results obtained on the PC computer for the data compressing step using the QR factorization and the fast Cholesky factorization algorithms (for the MOESP approach) are presented and compared with the MATLAB codes. Additional results are given in [9], using the fast algorithm exploiting the displacement structure. Both fast algorithms could be used for all

applications, except Application 4, for which the computed correlation matrix was not positive definite. Compared to the standard QR factorization, the gain in efficiency was remarkable (usually, more than 50 times faster), and the accuracy was comparable. On the other hand, MATLAB calculations of the correlations appeared as rather inefficient and they are not included in the comparison. The results in [15] have also shown that MATLAB QR algorithm was faster than SLICOT calculations. The reason discovered later was the bad use of the computer cache memory in the SLICOT tests, correlated to the use of Fortran BLAS libraries, rather than of optimized BLAS. By reducing the size of the provided workspace, and implicitly exploiting the abilities of SLICOT routines to process the data sequentially (in either outer on inner cycles), the efficiency of SLICOT calculations has been improved. Timing results for various workspace sizes (on the PC) are shown in Figure 1. The notation "min" and "full" denote that the minimum amount and the recommended (assuming a large enough cache size) amount, respectively, of workspace needed for each application has been specified as the workspace size, while "full/c" means that the "recommended" workspace size has been divided by the scalar c. The value "full/24" provided a much better use of the cache memory for the most time consuming applications, and enabled to slightly outperform the MATLAB calculations. Since the efficiency of some calculations could depend on the size of the workspace provided, two of the mexfiles in the SLIDENT toolbox, order and findBD, include the parameter ldwork—the length of the working array. A default value is used when ldwork is not specified, or if it is less than the minimal workspace needed.

Table 7: Relative output errors using the QR or Cholesky factorization algorithms.

#		Relative output errors				
	$\operatorname{SLmoesp}$	OMOESP	MOESP,	N4SID,		
			SLmoesm	SLn4sid		
1	5.39e-01	1.02e+01	4.86e-01	6.09e-01		
2	3.30e-02	7.33e-02	1.50e-02	2.16e-02		
3	6.03e-01	6.42e-01	4.96e-01	6.21e-01		
4	2.24e+02	6.49e + 10	2.37e-01	2.94e-01		
5	1.14e-01	7.38e + 04	3.60e-02	4.51e-02		
6	5.50e-01	5.67e-01	4.89e-01	5.53e-01		
7	1.13e + 02	8.06e + 04	1.76e-01	3.45e-01		
8	3.61e-01	2.21e+01	2.54e-01	2.78e-01		
9	4.23e-01	4.20e-01	1.38e-01	1.38e-01		

Table 7 presents comparative results, obtained on the PC, for the accuracy of the calculations using SLICOT algorithms, as well as the available MATLAB codes. Standard QR factorization has been used in all these calculations, but exactly the same accuracy results have been obtained using the Cholesky factorization algorithm, except for Application # 4, for which the Cholesky algorithm failed, and the QR algorithm was automatically called. The input and output trajectories for Application # 4 are plotted in the Figure 2 and Figure 3, respectively. The reported relative output error has been computed with the following MATLAB formula

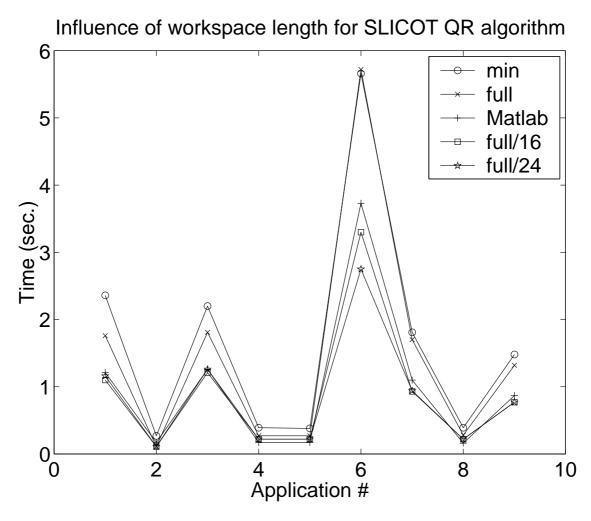


Figure 1: Influence of workspace length for the SLICOT QR factorization algorithm with non-optimized BLAS (MOESP approach).

err = norm(y - ye,1)/norm(y,1);

where ye denotes the estimated output of the system, evaluated using the estimated system matrices, A, B, C, and D, the given input trajectory, $\{u_k\}$, and an estimate of the initial state of the system, found, in turn, using A, B, C, D, and the given trajectories, $\{u_k\}$ and $\{y_k\}$. In the tables and figures, SLmoesp and SLn4sid denote the SLICOT implementations of the standard MOESP and N4SID approaches, OMOESP denotes a previous MATLAB code using TLS [23], MOESP denotes a new MATLAB code which computes the B and D matrices by solving a large least squares problem built using the matrix pair (A, C) and the input and output trajectories of the system [8], and N4SID denotes the robust MATLAB code in [21]. As already shown (Subsection 2.4), the SLICOT toolbox also contains the algorithm in the new MOESP code (called here SLmoesm, corresponding to the m-file slmoesm, or to the combination of the computational routines IBO1PD and IBO1QD, called by the driver routines IBO1BD and IBO1CD,

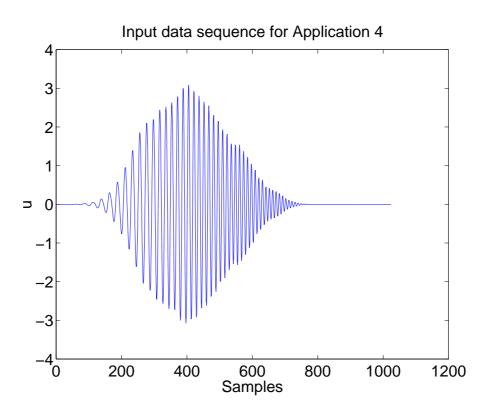


Figure 2: Input data sequence for Application # 4.

respectively); this algorithm gave the same accuracy as MOESP. The applications 4 and 7 are difficult for standard MOESP approach, because ill-conditioned matrices appear during the computation of B and D. However, there is no difficulty in solving these applications using the other methods, implemented in the m-files sln4sid, slmoen4, or slmoesm. Figure 4 plots the trajectories of the output, $\{y\}$, estimated output, $\{y_e\}$, and filtered output, $\{y_f\}$ (with Kalman gain also included in the estimated output calculations), when slmoesm was used for Application # 4. The filtered output is practically indistinguishable from the output.

The bar charts in Figure 5 illustrate the timing results on the Sun machine for computing system matrices using the QR factorization algorithm. Figure 6 presents the timing results when SLICOT routines used the Cholesky factorization algorithm. Since the bars in Figure 6 for SLICOT calculations have heights almost negligible compared to the MATLAB codes (except for Application 4), Table 8 presents the timing results numerically. When successful, the Cholesky factorization algorithm enabled to significantly increase the efficiency, while preserving the same accuracy as for the QR factorization algorithm. The speed-up factors vary between 10 and 20 when comparing to the SLICOT QR factorization algorithm, and between 15 and 40 when comparing to MATLAB codes. For a PC machine, the gain in efficiency was even larger. Figure 7 illustrates the timing results on the PC machine for computing the system matrices and Kalman gain using the Cholesky factorization algorithm, for all method-oriented m-files. For some applications and methods, the CPU time was below the measurable time (0.05 sec.), and in such cases the corresponding bars could not be plotted. It is apparent from Figure 7 that slmoesp and

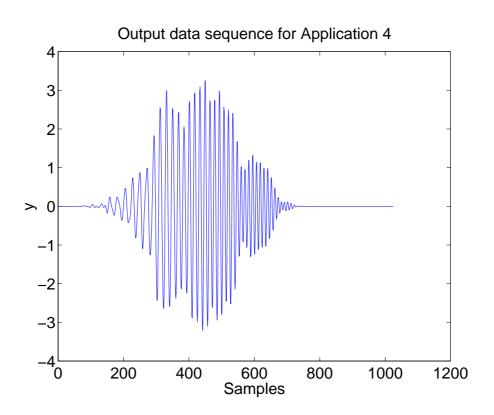


Figure 3: Output data sequence for Application # 4.

slmoen4 were the most efficient methods, while slmoesm was less efficient for some applications, like the applications # 1, 3, 6, and 7.

6 Conclusions

Algorithmic, implementation and numerical details concerning subspace-based techniques for linear multivariable system identification have been described and compared. The techniques are implemented in the new system identification toolbox for the SLICOT Library—SLIDENT. This toolbox includes interfaces (mexfiles and m-files) to the MATLAB and Scilab environments, which improve the user-friendliness of the collection. The results obtained show that the fast algorithmic variants included in the toolbox can frequently be used, and they are significantly more efficient than the standard QR factorization and the MATLAB codes.

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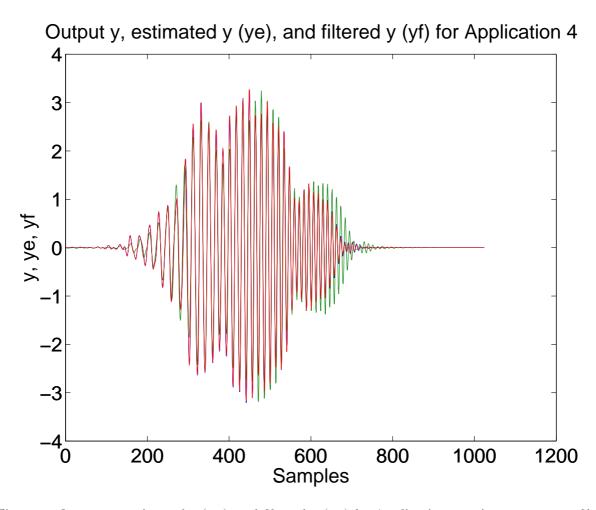


Figure 4: Output y, estimated y (ye), and filtered y (yf) for Application 4, using slmoesm m-file.

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Table 8: CPU time (sec.) for computing the system matrices using Cholesky factorization algorithm in the SLICOT-based calculations.

#			Time		
	$\operatorname{SLmoesp}$	OMOESP	MOESP	SLn4sid	N4SID
1	0.04	0.76	1.32	0.06	0.92
2	0.01	0.38	0.33	0.03	0.41
3	0.27	3.46	4.10	0.36	3.48
4	0.37	0.71	0.55	0.37	0.68
5	0.03	0.72	0.51	0.04	0.69
6	0.32	8.23	10.29	0.37	8.13
7	0.10	2.83	3.37	0.16	2.80
8	0.03	0.69	0.48	0.04	0.62
9	0.08	2.13	1.75	0.13	2.10

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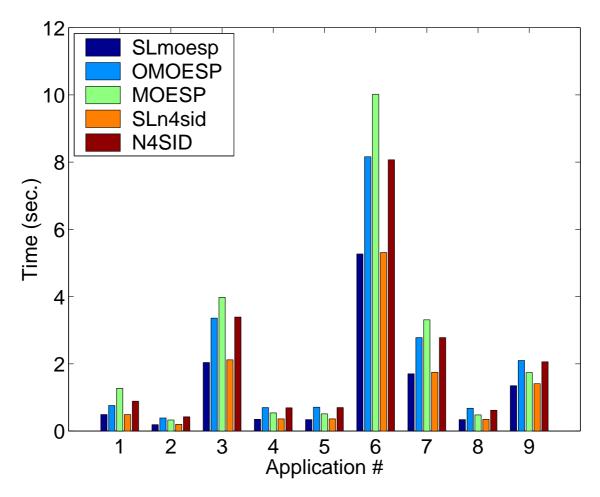


Figure 5: CPU time (sec.) for computing the system matrices using the QR factorization algorithm.

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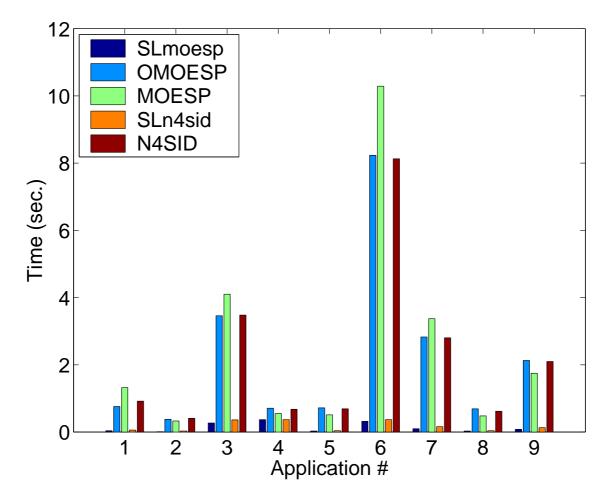


Figure 6: CPU time (sec.) for computing the system matrices using Cholesky factorization algorithm in the SLICOT-based calculations.

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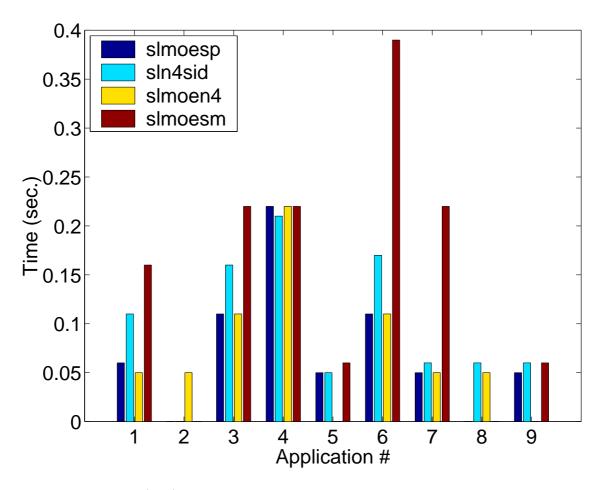


Figure 7: CPU time (sec.) for computing the system matrices and Kalman gain using the Cholesky factorization algorithm and the method-oriented m-files.

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