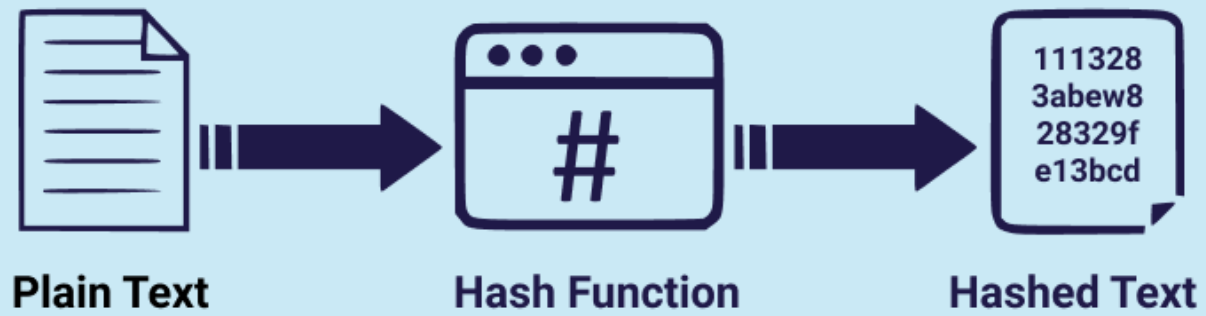


# Data Structures, In Detail

CS 374

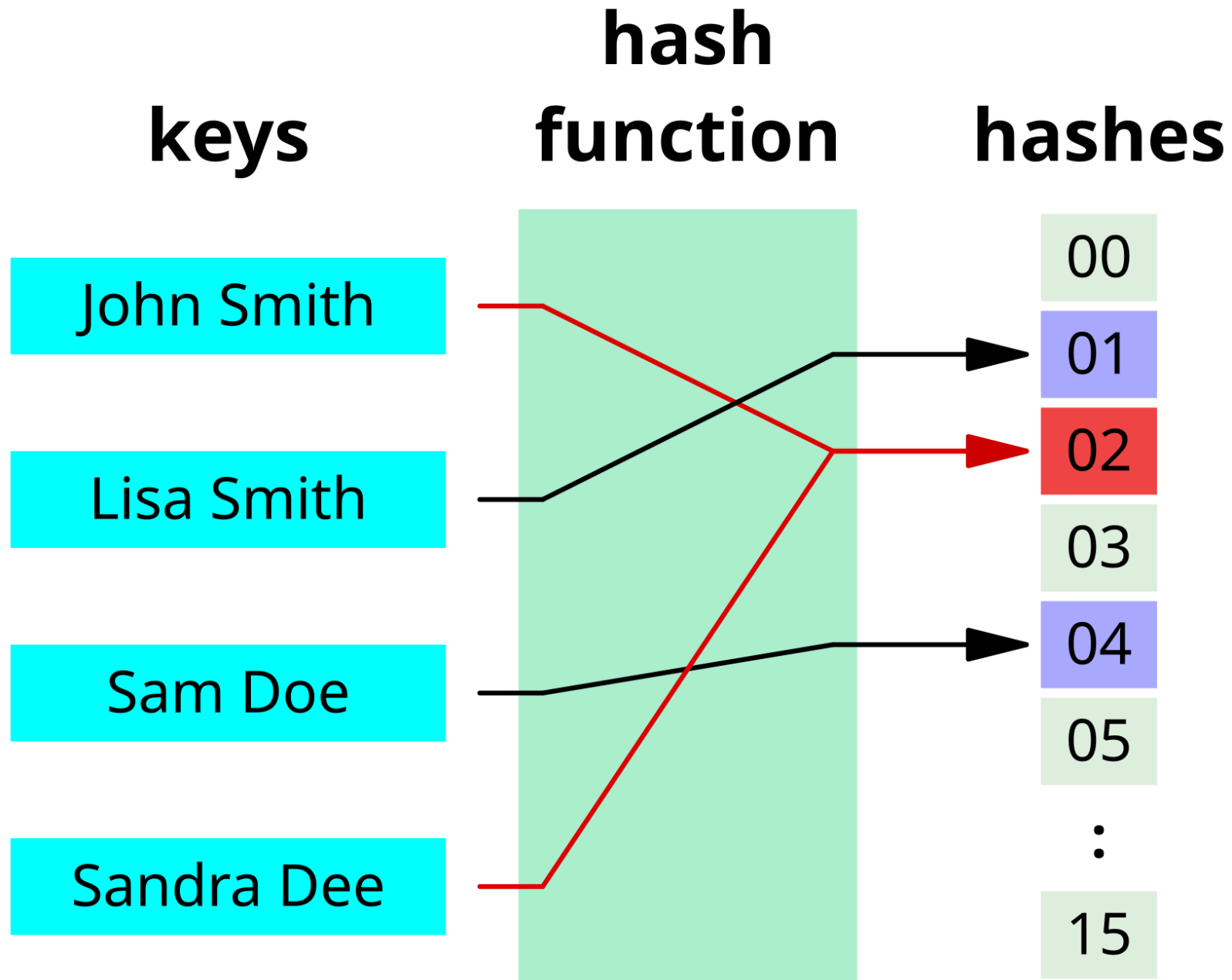
# Hashing/Dictionaries

## Hashing Algorithm



# Hashing

- A *hash* is a function that maps keys to an integer.
- Used in dictionaries, etc. to help you find it
- Also used for cryptography
- Have a function, then mod by  $m$  to stick it into spot  $m$  in the structure
- Collisions are an issue – what if two things have same hash?



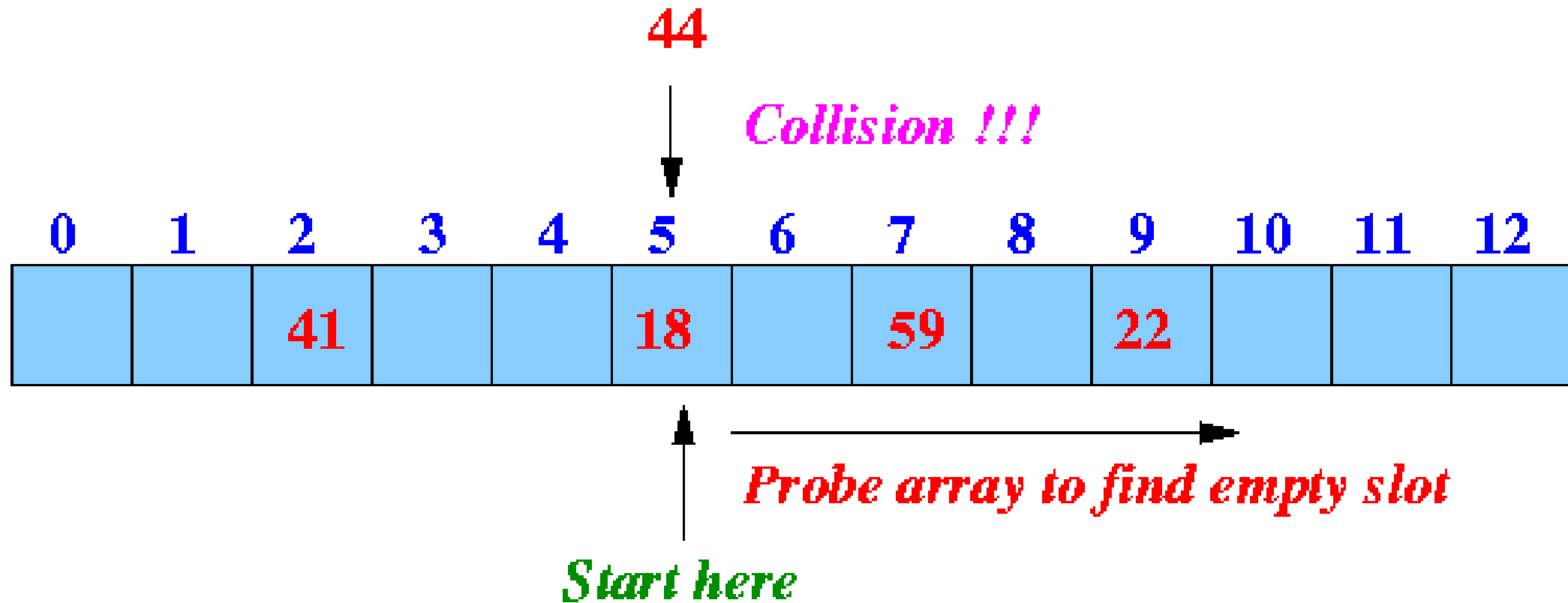
# Example

- 3. Given a hash function of  $f(x) = x + 31 \bmod 11$ , determine *if* there is a collision for items inserted with the keys 60, 9, 45, 3, 70, 5, 90, 93, 39, 77, 70 in that order.
- At which item does the first collision occur?

0	1	2	3	4	5	6	7	8	9
34	5	55	21		2		3	8	13

Figure 3.11: Collision resolution by open addressing and sequential probing, after inserting the first eight Fibonacci numbers in increasing order with  $H(x) = (2x + 1) \bmod 10$ . The red elements have been bumped to the first open slot after the desired location.

# Linear Probing





# Example

- 3. Given a hash function of  $f(x) = x + 31 \bmod 11$ , determine *if* there is a collision for items inserted with the keys 60, 9, 45, 3, 70, 5, 90, 93, 39, 77, 70 in that order.
- What is the result of the array when resolving collisions using linear/sequential probing? Use arrows/different colors to demonstrate when probing occurs.

# Chaining

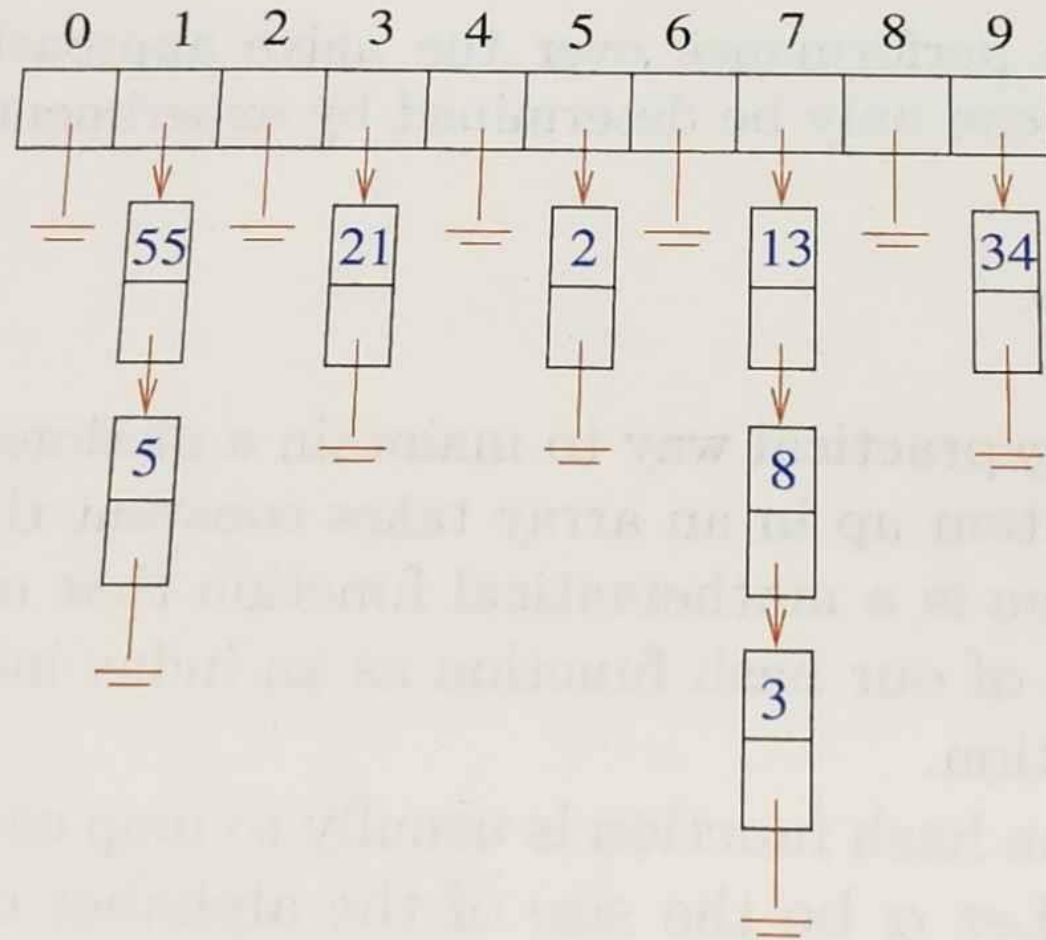


Figure 3.10: Collision resolution by chaining, after hashing the first eight Fibonacci numbers in increasing order, with hash function  $H(x) = (2x + 1) \bmod 10$ . Insertions occur at the head of each list in this figure.

# Example

- 3. Given a hash function of  $f(x) = x + 31 \bmod 11$ , determine *if* there is a collision for items inserted with the keys 60, 9, 45, 3, 70, 5, 90, 93, 39, 77, 70 in that order.
- What is the result of the array when resolving collisions using chaining? Draw a picture.

$$f(x) = x + 31 \bmod 11: 60, 9, 45, 3, 70, 5, 90, 93, 39, 77, 70$$

[illegible]

$$f(x) = x + 31 \bmod 11: 60, 9, 45, 3, 70, 5, 90, 93, 39, 77, 70$$

[illegible]

# More Collision Resolution

- Quadratic hashing
  - If value  $x$  matches to  $y$ ,
  - Instead of moving  $y+1$ , move  $y+1^2$ .
  - Then, if  $y+1^2$  is full, try  $y+2^2$
  - Chances are, you'll find an empty spot faster.
- Double hashing
  - Have a separate hash function that you run when there are collisions
- And more!

# How many things are in there?

- Hash table has fixed size
  - Like if implemented with arrays
- So, need to keep track of "how full" it is
- Load factor:  $\alpha = \frac{n}{m}$ 
  - $n$  = size of data (num. entries)
  - $m$  = number of buckets (capacity)

# Load factor

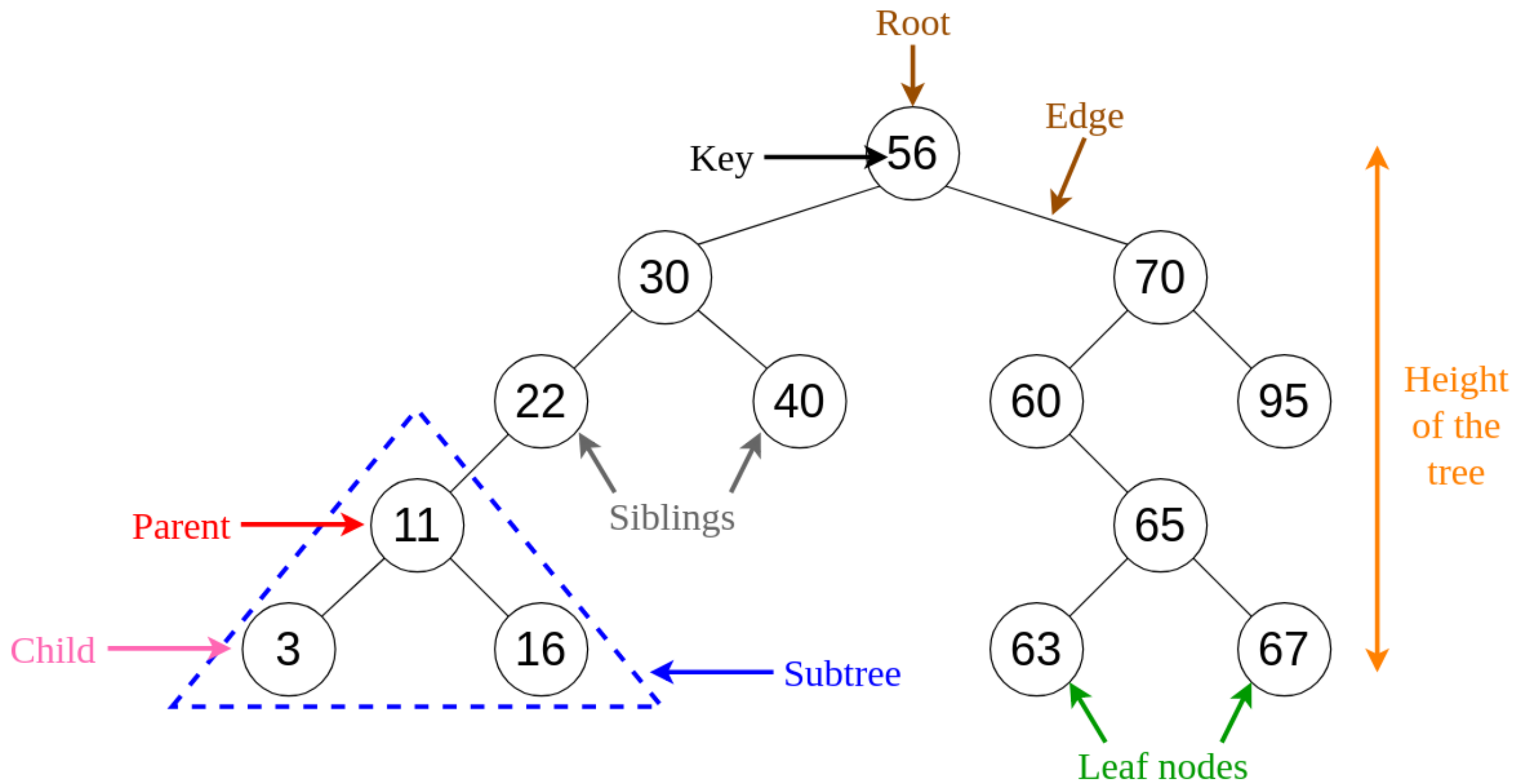
- You decide how full you want it to get
  - Set some  $\alpha_{max}$  threshold
- Resize, *then rehash* when it gets too big ( $\alpha_{max}$ ), or when gets too small ( $\alpha_{max}/4$  ??)
- Also may depend on method of collision resolution



# Hash Table Operations

	Hash table (expected)	Hash table (worst case)
Search( $L, k$ )	$O(n/m)$	$O(n)$
Insert( $L, x$ )	$O(1)$	$O(1)$
Delete( $L, x$ )	$O(1)$	$O(1)$
Successor( $L, x$ )	$O(n + m)$	$O(n + m)$
Predecessor( $L, x$ )	$O(n + m)$	$O(n + m)$
Minimum( $L$ )	$O(n + m)$	$O(n + m)$
Maximum( $L$ )	$O(n + m)$	$O(n + m)$

# Binary Search Trees



Insert the following items into a BST:

60, 9, 45, 3, 70, 5, 90, 93, 39, 77, 70

Tree A

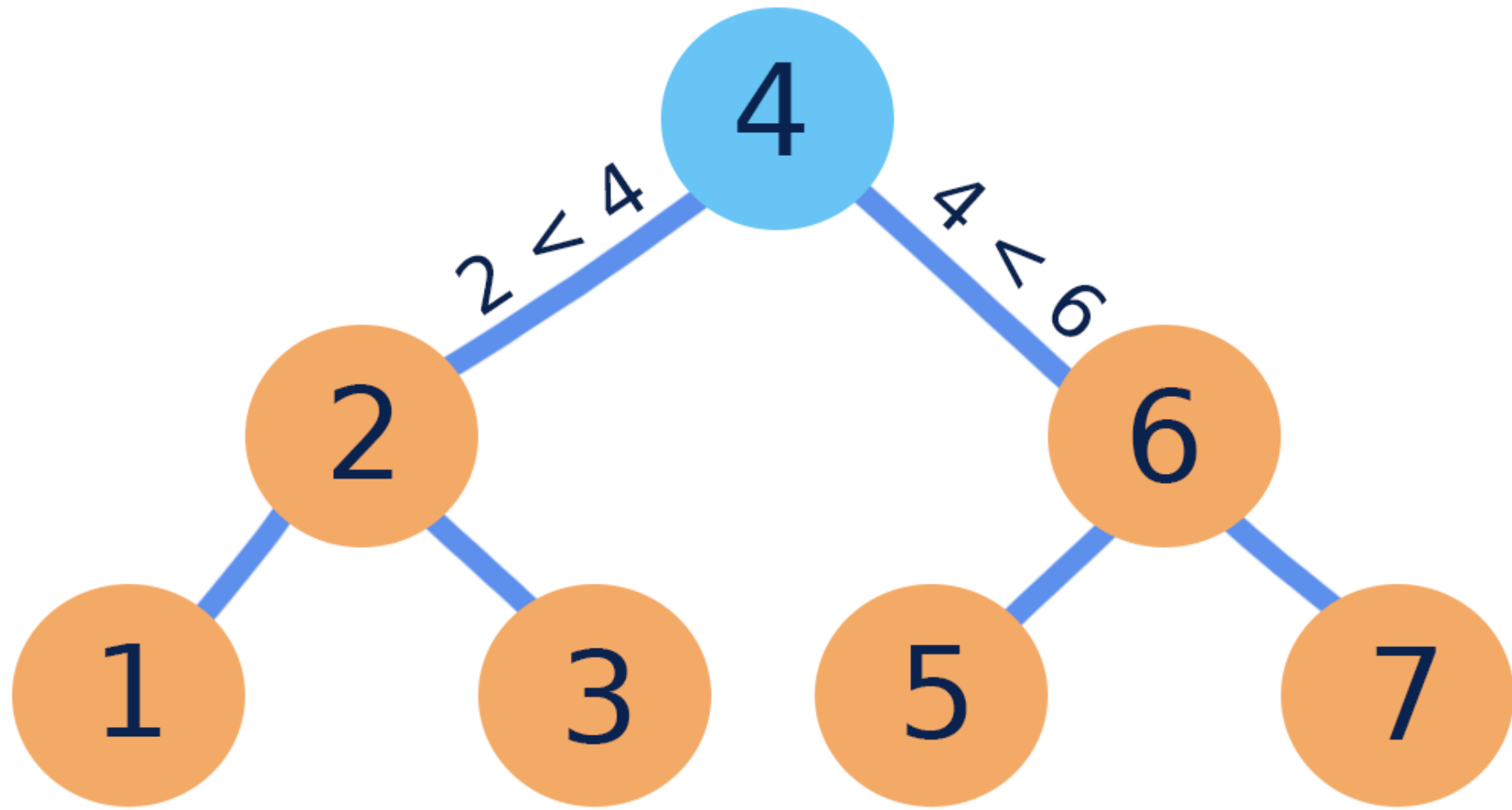
Insert the following items into a BST:

58 47 84 39 55 14 98 53 79 42 87 3 90 1

Tree B

# BST Terms

- Complete
  - All levels full
- Balanced
  - One side not more than other



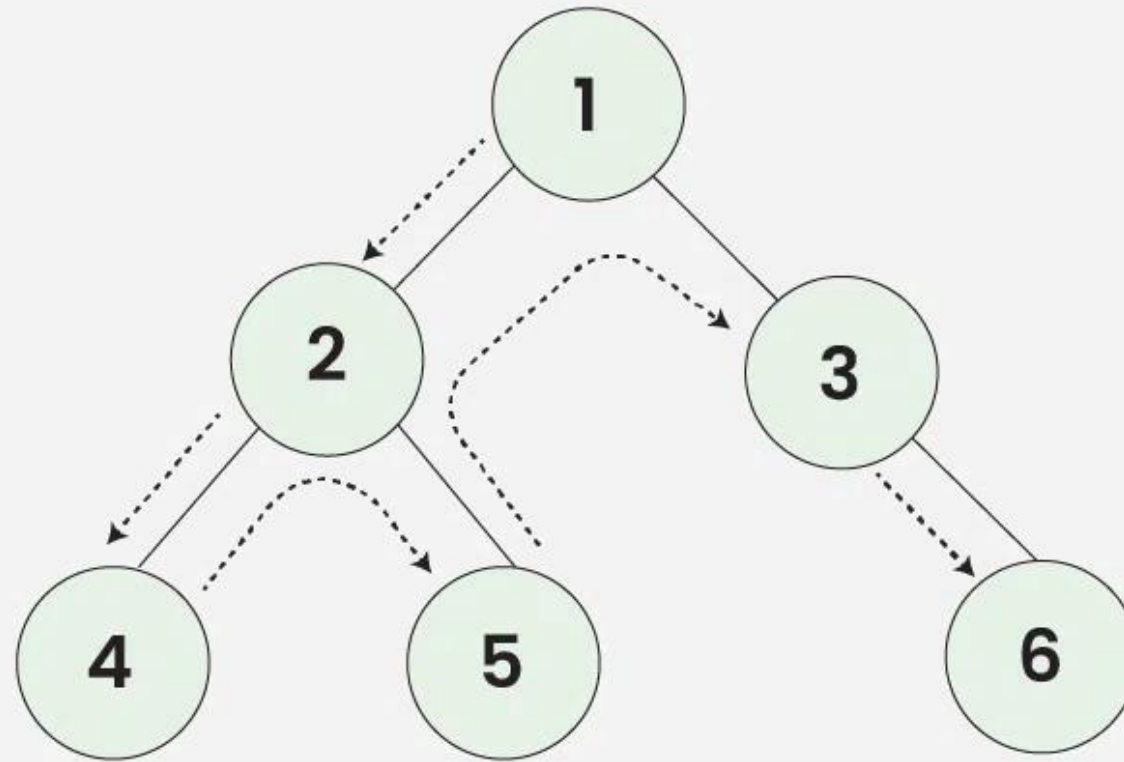
In Order Traversal: 1 2 3 4 5 6 7

# InOrder Traversal of tree A



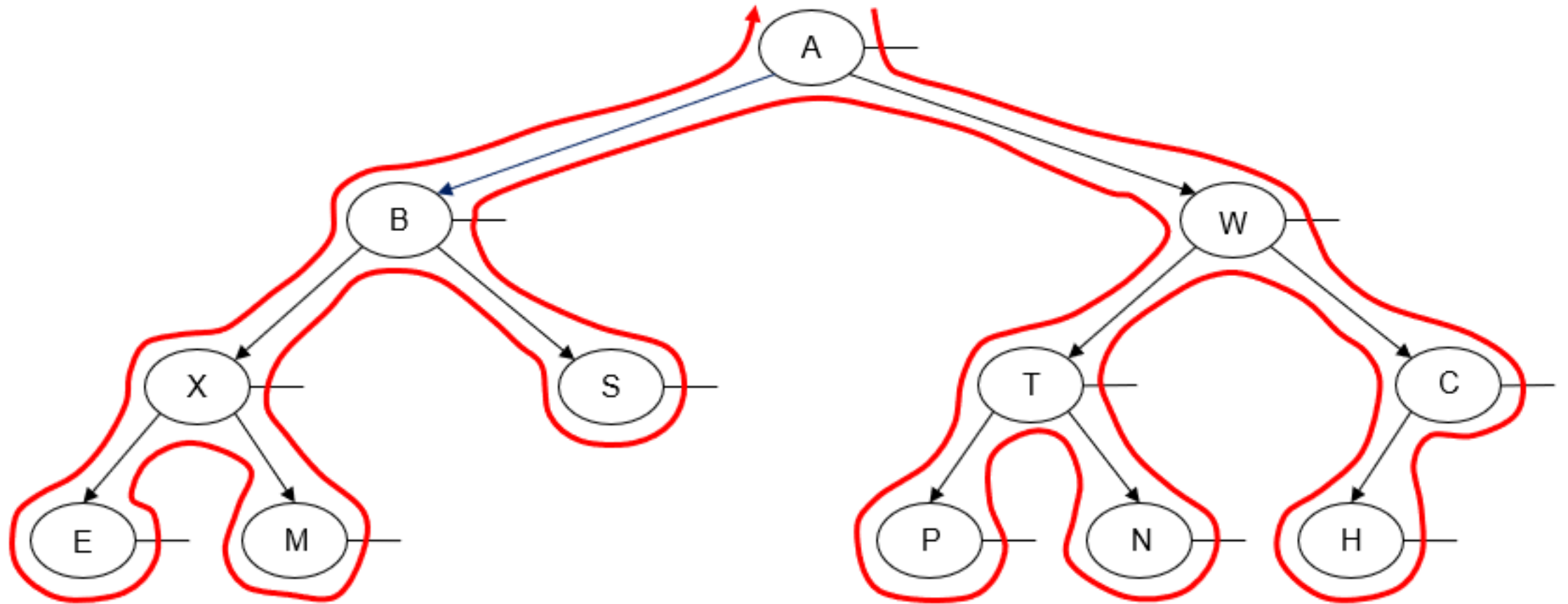


# Preorder Traversal of Binary Tree



Preorder Traversal:  $1 \rightarrow 2 \rightarrow 4 \rightarrow 5 \rightarrow 3 \rightarrow 6$

# PreOrder Traversal of tree A



- Post-order traversal: E M X S B P N T H C W A
- (this isn't a BST, just a BT)

# PostOrder Traversal of tree A

# BST Operations

- Search –  $O(h)$ , where  $h$ = height of tree
- Find min/max –  $O(h)$
- Traversal –  $O(n)$
- Insertion –  $O(h)$
- Deletion –  $O(h)$

# Warm-up

Meet up with your teams. On the whiteboard:

## Problem A.

- Using *mod* 6,
- Create a hash table for 10, 20, 7, 11, 17, 14, 2, 4, 6, 15, 3
- Equation:
  - $f(x) = 18x^6 - x^4 + 4x^3 + 21x + 2$
- Rehash when  $\alpha > 0.6$ . What do you choose for new mod? Why?
- Use your choice of collision resolution

## Problem B.

- Create a binary search tree for 20, 10, 7, 11, 17, 14, 2, 4, 6, 15, 3
- Create a binary search tree for 17, 10, 7, 11, 20, 14, 2, 4, 6, 15, 3
- These are the same numbers: What's the difference?

# Heaps

# Heaps

- Min-heap: smaller = at the top
- Max-heap: larger = at the top
- Complete binary tree
  - All levels are filled...
  - ...except maybe the last one
  - Last level gets filled Left to Right.



# Heap – Insertion

- Insert latest item in bottommost, leftmost position.
- "Bubble up" if necessary.
  - Swap the item with its parent until heap property is not violated
  - Also called "heapify", sometimes "heapify up"

Example: Insert the following numbers into a min-heap.

18, 12, 3, 7, 15, 4, 11, 16, 5, 14.

Then draw the array.



# Heap – Deletion

- Remove the item at the root
- Move the "last" thing in the heap to the root
- "Bubble down" if necessary.
  - Check the children.
  - Swap the item with the smaller of the children
  - Continue swapping item down, so the smaller thing gets swapped up.
  - Also called "heapify", sometimes "heapify down"

Remove 2 things from the preceding heap.



# Heaps – that's great, but how does it code???

- Heaps use binary trees
- Stored in an array with level order
  - Index 0=minimum
  - For a given index  $i$ ,
    - Parent node is located at  $\frac{i-1}{2}$
    - Left child is located at  $2i + 1$
    - Right child is located at  $2i + 2$

# Heap – Search

- Ehhh not really
- Used to find the \*min\* element
- So don't worry about it!



# What do we use heaps for?

- Priority queues!
- Min-heap removal always gives the minimum element.

# Heap Complexity

Draw the max-heap for the following numbers,  
then remove 3 elements.

3, 14, 4, 8, 15, 16, 9, 12, 20

# Cool-Down: Solve my to-do list

Task	Deadline	Dr. Roscoe's Priority
Grade Algorithms HW 1	Tuesday	1
Knit hat	March 31	5
Read a research paper	Friday	2
Post the Stats homework	5pm today	2
Submit Jan Term proposal	Sunday 11:59pm	1
Post Algorithms video	5pm today	1
Write a section of research paper	5pm today	1

- Discuss on what variable or combination of variables you will optimize.
- Describe on the board how you will construct a priority queue, and use a heap to return the order of tasks I should complete, until the queue is empty.

These slides are work-in-progress; more will be added