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1.2. How Did We Get Here?

1.2.1. The Basics

```
#include <bits/stdc++.h>
#define ll long long
#define sz(x) (int)(x).size()
using namespace std;

// mt19937
// rng(chrono::steady_clock::now().time_since_epoch().count());

// uniform_int_distribution<int>(1000,10000)(rng)

int main() {
    ios_base::sync_with_stdio(0);
    cin.tie(0);
    return 0;
}
```

1.2.2. Linux

```
g++ test.cpp
./test
./test < input.txt > output.txt (read/write to file)
```

1.2.3. Macros

Use vectorizations and math optimizations at your own peril. For gcc≥9, there are [[likely]] and [[unlikely]] attributes. Call gcc with -fopt-info-optimized-missed-optall for optimization info.

1.3. Tools

1.3.1. Floating Point Binary Search

```
union di {
    double d;
    ull i;
};
bool check(double);
// binary search in [L, R) with relative error 2^-eps
double binary_search(double L, double R, int eps) {
    di l = {L}, r = {R}, m;
    while (r.i - l.i > 1LL << (52 - eps)) {
        m.i = (l.i + r.i) >> 1;
        if (check(m.d))
            r = m;
else
            l = m;
} return l.d;
}
```

1.3.2. SplitMix64

```
using ull = unsigned long long;
inline ull splitmix64(ull x) {

// change to `static ull x = SEED; `for DRBG
ull z = (x += 0x9E3779B97F4A7C15);

z = (z ^ (z >> 30)) * 0xBF58476D1CE4E5B9;
z = (z ^ (z >> 27)) * 0x94D049BB133111EB;

return z ^ (z >> 31);
}
```

1.3.3. <random>

1.3.4. x86 Stack Hack

```
constexpr size_t size = 200 << 20; // 200MiB
int main() {
   register long rsp asm("rsp");
   char *buf = new char[size];
   asm("movq %0, %%rsp\n" ::"r"(buf + size));
   // do stuff
   asm("movq %0, %%rsp\n" ::"r"(rsp));
   delete[] buf;
}</pre>
```

1.4. Algorithms

1.4.1. Bit Hacks

```
// next permutation of x as a bit sequence
ull next_bits_permutation(ull x) {
  ull c = __builtin_ctzll(x), r = x + (1ULL << c);
  return (r ^ x) >> (c + 2) | r;
}
// iterate over all (proper) subsets of bitset s
void subsets(ull s) {
  for (ull x = s; x;) { --x &= s; /* do stuff */ }
}
```

1.4.2. Aliens Trick

```
// min dp[i] value and its i (smallest one)
pll get_dp(int cost);

l aliens(int k, int l, int r) {
    while (l != r) {
        int m = (l + r) / 2;
        auto [f, s] = get_dp(m);
        if (s == k) return f - m * k;
        if (s < k)
        r = m;
        else
        l = m + 1;
}
return get_dp(l).first - l * k;
}</pre>
```

1.4.3. Hilbert Curve

1.4.4. Infinite Grid Knight Distance

```
1  ll get_dist(ll dx, ll dy) {
    if (++(dx = abs(dx)) > ++(dy = abs(dy))) swap(dx, dy);
    if (dx == 1 && dy == 2) return 3;
    if (dx == 3 && dy == 3) return 4;
    ll lb = max(dy / 2, (dx + dy) / 3);
    return ((dx ^ dy ^ lb) & 1) ? ++lb : lb;
}
```

1.4.5. Poker Hand

```
1 using namespace std;
 3
   struct hand {
      static constexpr auto rk = [] {
        array<int, 256> x{};
auto s = "23456789TJQKACDHS";
        for (int i = 0; i < 17; i++) x[s[i]] = i % 13;
        return x;
 9
      }();
      vector<pair<int, int>> v;
11
      vector<int> cnt, vf, vs;
      int type;
13
      hand(): cnt(4), type(0) {}
      void add_card(char suit, char rank) {
        ++cnt[rk[suit]];
        for (auto δ[f, s] : v)
  if (s == rk[rank]) return ++f, void();
17
        v.emplace_back(1, rk[rank]);
      void process() {
```

```
sort(v.rbegin(), v.rend());
for (auto [f, s] : v) vf.push_back(f), vs.push_back(s);
         bool str = 0, flu = find(all(cnt), 5) != cnt.end();
23
         if ((str = v.size() == 5))
         for (int i = 1; i < 5; i++)
  if (vs[i] != vs[i - 1] + 1) str = 0;
if (vs = vector<int>{12, 3, 2, 1, 0})
  str = 1, vs = {3, 2, 1, 0, -1};
if (str && flu)
25
27
29
            type = 9;
         else if (vf[0] == 4)
31
            type = 8;
33
         else if (vf[0] == 3 && vf[1] == 2)
            type = 7;
         else if (str || flu)
35
            type = 5 + flu;
         else if (vf[0] == 3)
            type = 4;
         else if (vf[0] == 2)
39
            type = 2 + (vf[1] == 2);
41
43
      bool operator<(const hand &b) const {</pre>
45
         return make_tuple(type, vf, vs) <
                  make_tuple(b.type, b.vf, b.vs);
47
   };
```

1.4.6. Longest Increasing Subsequence

```
template <class I> vi lis(const vector<I> &S) {
    if (S.empty()) return {};
    vi prev(sz(S));
    typedef pair<I, int> p;
    vector res;
    rep(i, 0, sz(S)) {
        // change 0 -> i for longest non-decreasing subsequence
        auto it = lower_bound(all(res), p{S[i], 0});
        if (it == res.end())
            res.emplace_back(), it = res.end() - 1;
        *it = {S[i], i};
        prev[i] = it == res.begin() ? 0 : (it - 1)->second;
}
int L = sz(res), cur = res.back().second;
vi ans(L);
while (L--) ans[L] = cur, cur = prev[cur];
return ans;
}
```

1.4.7. Mo's Algorithm on Tree

```
void MoAlgoOnTree() {
        Dfs(0, -1);
        vector<int> euler(tk);
       for (int i = 0; i < n; ++i) {
  euler[tin[i]] = i;</pre>
           euler[tout[i]] = i;
       vector<int> l(q), r(q), qr(q), sp(q, -1);
for (int i = 0; i < q; ++i) {
   if (tin[u[i]] > tin[v[i]]) swap(u[i], v[i]);
           int z = GetLCA(u[i], v[i]);
           sp[i] = z[i];
13
              l[i] = tin[u[i]], r[i] = tin[v[i]];
              l[i] = tout[u[i]], r[i] = tin[v[i]];
           qr[i] = i;
17
       sort(qr.begin(), qr.end(), [δ](int i, int j) {
  if (l[i] / kB == l[j] / kB) return r[i] < r[j];</pre>
           return l[i] / kB < l[j] / kB;
21
23
       vector<bool> used(n);
       // Add(v): add/remove v to/from the path based on used[v] for (int i = 0, tl = 0, tr = -1; i < q; ++i) { while (tl < l[qr[i]]) Add(euler[tl++]);
25
           while (tl > l[qr[i]]) Add(euler[--tl]);
while (tr > r[qr[i]]) Add(euler[tr--]);
27
           while (tr < r[qr[i]]) Add(euler[++tr]);
29
           // add/remove LCA(u, v) if necessary
31
```

2. Data Structures

2.1. GNU PBDS

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
```

```
3 | #include <ext/pb_ds/tree_policy.hpp>
   using namespace __gnu_pbds;
    // most std::map + order_of_key, find_by_order, split, join
   template <typename T, typename U = null_type>
   using ordered_map = tree<T, U, std::less<>, rb_tree_tag,
                                tree_order_statistics_node_update>;
    // useful tags: rb_tree_tag, splay_tree_tag
    template <typename T> struct myhash {
     size_t operator()(T x) const; // splitmix, bswap(x*R), ...
13
   // most of std::unordered_map, but faster (needs good hash)
template <typename T, typename U = null_type>
using hash_table = gp_hash_table<T, U, myhash<T>>;
15
19 // most std::priority_queue + modify, erase, split, join
   using heap = priority_queue<int, std::less<>>;
21 // useful tags: pairing_heap_tag, binary_heap_tag,
                      (rc_)?binomial_heap_tag, thin_heap_tag
```

2.2. Segment Tree (ZKW)

```
1 struct segtree {
       using T = int;
       T f(T a, T b) { return a + b; } // any monoid operation static constexpr T ID = 0; // identity element
       int n:
       vector<T> v;
       segtree(int n_) : n(n_), v(2 * n, ID) {}
       segtree(vector<T> \delta a) : n(a.size()), v(2 * n, ID) {
          copy_n(a.begin(), n, v.begin() + n);
          for (int i = n - 1; i > 0; i--)
v[i] = f(v[i * 2], v[i * 2 + 1]);
11
       void update(int i, T x) {
  for (v[i += n] = x; i /= 2;)
   v[i] = f(v[i * 2], v[i * 2 + 1]);
13
15
17
       T query(int l, int r) {
          T tl = ID, tr = ID;
          for (l += n, r += n; l < r; l /= 2, r /= 2) {
  if (l & 1) tl = f(tl, v[l++]);
19
21
             if (r \delta 1) tr = f(v[--r], tr);
23
          return f(tl, tr);
25 };
```

2.3. Line Container

```
1 struct Line {
      mutable ll k, m, p;
      bool operator<(const Line &o) const { return k < o.k; }</pre>
      bool operator<(ll x) const { return p < x; }</pre>
   };
// add: line y=kx+m, query: maximum y of given x
struct LineContainer : multiset<Line, less<>>> {
 5
      // (for doubles, use inf = 1/.0, div(a,b) = a/b)
      static const ll inf = LLONG_MAX;
      ll div(ll a, ll b) { // floored division
  return a / b - ((a ^ b) < 0 && a % b);</pre>
11
13
      bool isect(iterator x, iterator y) {
        if (y == end()) return x->p = inf, 0;
        if (x->k == y->k)
15
          x->p = x->m > y->m ? inf : -inf;
17
        else
          x->p = div(y->m - x->m, x->k - y->k);
19
        return x->p>=y->p;
21
      void add(ll k, ll m) {
        auto z = insert(\{k, m, \theta\}), y = z++, x = y;
23
        while (isect(y, z)) z = erase(z);
         if (x != begin() && isect(--x, y))
        isect(x, y = erase(y));
while ((y = x) != begin() && (--x)->p >= y->p)
25
27
           isect(x, erase(y));
29
      ll query(ll x) {
        assert(!empty());
        auto l = *lower_bound(x);
31
        return l.k * x + l.m;
33
   };
```

2.4. Li-Chao Tree

```
constexpr ll MAXN = 2e5, INF = 2e18;
struct Line {
```

```
Line(): m(\theta), b(-INF) {}
     Line(ll _m, ll _b) : m(_m), b(_b) {}
     ll operator()(ll x) const { return m * x + b; }
   struct Li Chao {
     Line a[MAXN * 4];
     void insert(Line seg, int l, int r, int v = 1) {
       if (l == r) {
11
         if (seg(l) > a[v](l)) a[v] = seg;
13
         return:
15
       int mid = (l + r) >> 1;
       if (a[v].m > seg.m) swap(a[v], seg);
       if (a[v](mid) < seg(mid)) {
         swap(a[v], seg);
19
         insert(seg, l, mid, v << 1);</pre>
         insert(seg, mid + 1, r, v << 1 | 1);
23
     ll query(int x, int l, int r, int v = 1) {
       if (l == r) return a[v](x);
       int mid = (l + r) >> 1;
       if (x \le mid)
27
         return max(a[v](x), query(x, l, mid, v << 1));
29
         return max(a[v](x), query(x, mid + 1, r, v \ll 1 | 1));
31 };
```

2.5. Heavy-Light Decomposition

struct heavy_light_decomposition {

```
vector<vector<int>> edges;
      vector<int> par, heavy, height, pos, head;
      heavy_light_decomposition(int n)
            n(n), edges(n + 1), par(n + 1), heavy(n + 1),
            height(n + 1), pos(n + 1), head(n + 1) {}
      void add_edge(int x, int y) {
  edges[x].push_back(y);
 9
        edges[y].push_back(x);
      void heavy_dfs(int c, vector<int> &sub) {
        heavy[c] = -1;
13
        sub[c] = 1;
        int max_size = 0;
        for (int i : edges[c]) {
          if (i != par[c]) {
  par[i] = c;
  height[i] = height[c] + 1;
17
19
             heavy_dfs(i, sub);
sub[c] += sub[i];
             if (sub[i] > max_size) {
  heavy[c] = i;
               max_size = sub[i];
25
          }
27
        }
29
      void decompose(int c, int h, int &timer) {
        pos[c] = timer++;
        head[c] = h;
        if (heavy[c] != -1) { decompose(heavy[c], h, timer); }
        for (int i : edges[c]) {
          if (i != par[c] && i != heavy[c]) {
             decompose(i, i, timer);
35
37
        }
      }
39
      void build() {
        int timer = 0;
        vector<int> sub(n + 1);
41
        for (int i = 1; i <= n; ++i) {
43
          if (sub[i] == 0) {
             par[i] = 0;
             height[i] = 0;
45
             heavy_dfs(i, sub);
47
             decompose(i, i, timer);
49
        }
51
      int get_position(int x) { return pos[x]; }
      vector<array<int, 2>> path_queries(int x, int y) {
        vector<array<int, 2>> queries;
while (head[x] != head[y]) {
55
          if (height[head[x]] > height[head[y]]) { swap(x, y); }
          queries.push_back({pos[head[y]], pos[y]});
            = par[head[y]];
```

```
if (height[x] > height[y]) { swap(x, y); }
       queries.push_back({pos[x], pos[y]});
       return queries;
63 };
```

Wavelet Matrix

```
1 #pragma GCC target("popcnt,bmi2")
   #include <immintrin.h>
    // T is unsigned. You might want to compress values first
   template <typename T> struct wavelet_matrix {
      static_assert(is_unsigned_v<T>, "only unsigned T");
      struct bit_vector {
        static constexpr uint W = 64;
        uint n, cnt0;
        vector<ull> bits;
11
        vector<uint> sum;
        bit_vector(uint n_)
             : n(n_), bits(n / W + 1), sum(n / W + 1) {}
13
        void build() {
          for (uint j = 0; j != n / W; ++j)
  sum[j + 1] = sum[j] + _mm_popcnt_u64(bits[j]);
15
          cnt0 = rank0(n):
19
        void set_bit(uint i) { bits[i / W] |= 1ULL << i % W; }</pre>
        bool operator[](uint i) const {
          return !!(bits[i / W] & 1ULL << i % W);
21
        uint rank1(uint i) const {
23
          return sum[i / W]
25
                  _mm_popcnt_u64(_bzhi_u64(bits[i / W], i % W));
27
        uint rank0(uint i) const { return i - rank1(i); }
29
     uint n, lg;
      vector<bit_vector> b;
31
      wavelet_matrix(const vector<T> &a) : n(a.size()) {
        lg
33
         _lg(max(*max_element(a.begin(), a.end()), T(1))) + 1;
        b.assign(lg, n);
        vector<T> cur = a, nxt(n);
35
        for (int h = lg; h--;) {
  for (uint i = 0; i < n; ++i)
37
             if (cur[i] & (T(1) << h)) b[h].set_bit(i);
          b[h].build();
39
          int il = 0, ir = b[h].cnt0;
for (uint i = 0; i < n; ++i)
  nxt[(b[h][i] ? ir : il)++] = cur[i];</pre>
41
43
          swap(cur, nxt);
        }
45
      T operator[](uint i) const {
        T res = 0;
47
        for (int h = lg; h--;)
if (b[h][i])
49
              += b[h].cnt0 - b[h].rank0(i), res |= T(1) << h;
51
            i = b[h].rank0(i);
53
        return res;
      // query k-th smallest (0-based) in a[l, r)
55
      T kth(uint l, uint r, uint k) const {
57
        T res = 0;
        for (int h = lg; h--;) {
          uint tl = b[h].rank0(l), tr = b[h].rank0(r);
if (k >= tr - tl) {
59
            k -= tr - tl:
61
            l += b[h].cnt0 - tl;
            r += b[h].cnt0 - tr;
63
            res |= T(1) << h;
          } else
65
            l = tl, r = tr;
67
        return res;
69
      // count of i in [l, r) with a[i] < u
71
      uint count(uint l, uint r, T u) const {
        if (u >= T(1) << lg) return r - l;
        uint res = 0;
73
        for (int h = lg; h--;) {
          uint tl = b[h].rank0(l), tr = b[h].rank0(r);
          if (u & (T(1) << h)) {
            l += b[h].cnt0 - tl;
r += b[h].cnt0 - tr;
            res += tr - tl;
          } else
             l = tl, r = tr;
```

return res:

```
85 };
   2.7. Link-Cut Tree
   const int MXN = 100005;
   const int MEM = 100005;
   struct Splay {
      static Splay nil, mem[MEM], *pmem;
     Splay *ch[2], *f;
int val, rev, size;
Splay() : val(-1), rev(0), size(0) {
   f = ch[0] = ch[1] = &nil;
      11
13
      bool isr() {
  return f->ch[0] != this && f->ch[1] != this;
15
      int dir() { return f->ch[0] == this ? 0 : 1; }
17
      void setCh(Splay *c, int d) {
19
        ch[d] = c;
        if (c != &nil) c->f = this;
21
        pull();
23
      void push() {
        if (rev) {
          swap(ch[0], ch[1]);
if (ch[0] != &nil) ch[0]->rev ^= 1;
25
          if (ch[1] != &nil) ch[1]->rev ^= 1;
27
          rev = 0:
29
        }
      }
      void pull() {
31
        size = ch[0]->size + ch[1]->size + 1;
if (ch[0] != &nil) ch[0]->f = this;
33
        if (ch[1] != &nil) ch[1]->f = this;
35
    } Splay::nil, Splay::mem[MEM], *Splay::pmem = Splay::mem;
   Splay *nil = &Splay::nil;
37
   void rotate(Splay *x) {
39
      Splay *p = x->f;
int d = x->dir();
      if (!p->isr())
43
        p->f->setCh(x, p->dir());
      else
     x->f = p->f;
p->setCh(x->ch[!d], d);
      x->setCh(p, !d);
     p->pull();
49
      x->pull();
   }
51
   vector<Splay *> splayVec;
   void splay(Splay *x) {
53
      splayVec.clear();
55
      for (Splay *q = x;; q = q -> f) {
        splayVec.push_back(q);
        if (q->isr()) break;
      reverse(begin(splayVec), end(splayVec));
59
      for (auto it : splayVec) it->push();
while (!x->isr()) {
        if (x->f->isr())
63
          rotate(x);
        else if (x-)dir() == x-)f-)dir()
          rotate(x->f), rotate(x);
65
        else
67
          rotate(x), rotate(x);
      }
   }
69
   Splay *access(Splay *x) {
71
      Splay *q = nil;
for (; x != nil; x = x->f) {
        splay(x);
        x->setCh(q, 1);
        q = x;
      return q;
79
   }
    void evert(Splay *x) {
     access(x);
      splay(x);
x->rev ^= 1;
      x->push();
```

```
x->pull();
    void link(Splay *x, Splay *y) {
            evert(x);
 89
        access(x);
        splav(x):
        evert(v);
       x->setCh(y, 1);
 93 }
     void cut(Splay *x, Splay *y) {
 95
       // evert(x);
        access(y);
 97
        splay(y)
       y->push();
 99
       y->ch[\theta] = y->ch[\theta]->f = nil;
101
     int N, Q;
Splay *vt[MXN];
103
    int ask(Splay *x, Splay *y) {
       access(x);
107
        access(y);
        splav(x):
109
        int res = x->f->val;
        if (res == -1) res = x->val;
111
       return res:
113
     int main(int argc, char **argv) {
       scanf("%d%d", &N, &Q);
for (int i = 1; i <= N; i++)
115
117
          vt[i] = new (Splay::pmem++) Splay(i);
        while (Q--) {
119
          char cmd[105];
          int u, v;
scanf("%s", cmd);
if (cmd[1] == 'i') {
    scanf("%d%d", &u, &v);
121
123
            } else if (cmd[0] ==
125
            scanf("%d", &v);
cut(vt[1], vt[v]);
127
          } else {
            etse {
    scanf("%d%d", &u, &v);
    int res = ask(vt[u], vt[v]);
    printf("%d\n", res);
129
131
133
       }
     }
```

3. Graph

3.1. Floyd-Warshall

```
for (int k = 0; k < n; ++k) {
    for (int i = 0; i < n; ++i) {
        for (int j = 0; j < n; ++j) {
            d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
        }
    }
}</pre>
```

3.2. Dijkstra

```
1 struct diikstra {
      using S =
      int; // regular edge weight type, change if needed
using T = long long; // maximum possible distance type,
 3
                                // change if needed
 5
      const T INF = 1e18:
      struct Edge {
        int target;
 9
        S weight;
11
      struct next_shortest {
        int node;
        T dist;
13
        bool operator<(const next_shortest &rhs) const {
           return dist > rhs.dist;
15
17
      vector<vector<Edge>> e;
      dijkstra(int n) : n(n), e(n + 1) {}
void add_directed_edge(int x, int y, S z) {
        e[x].push_back({y, z});
23
      void add_undirected_edge(int x, int y, S z) {
```

```
e[x].push_back({y, z});
        e[y].push_back({x, z});
     vector<T> shortest_path(int src) {
       vector<T> d(n + 1, INF);
29
        d[src] = 0;
       priority_queue<next_shortest> pq;
pq.push({src, 0});
31
        while (!pq.empty()) {
33
         int cur_node = pq.top().node;
          T cur_dist = pq.top().dist;
          pq.pop();
          if (cur_dist > d[cur_node]) continue;
          for (auto i : e[cur_node])
            if (cur_dist + i.weight < d[i.target]) {</pre>
39
              d[i.target] = cur_dist + i.weight;
              pq.push({i.target, d[i.target]});
        return d;
   };
```

3.3. Modeling

```
21
• Maximum/Minimum flow with lower bound / Circulation problem
                                                                                        23
  1. Construct super source S and sink T. 2. For each edge (x,y,l,u), connect x\to y with capacity u-l.
  3. For each vertex v, denote by in(v) the difference between the sum ^{25}
      of incoming lower bounds and the sum of outgoing lower bounds.
  4. If in(v) > 0, connect S \to v with capacity in(v), otherwise, connect
     v \to T with capacity -in(v).
      – To maximize, connect t \to s with capacity \infty (skip this in cir-
        culation problem), and let f be the maximum flow from S to T.
        If f \neq \sum_{v \in V, in(v) > 0} in(v), there's no solution. Otherwise, the
      maximum flow from s to t is the answer.

To minimize, let f be the maximum flow from S to T. Connect
        t \to s with capacity \infty and let the flow from S to T be f'. If _{35}
        f + f' \neq \sum_{v \in V, in(v) > 0} in(v), there's no solution. Otherwise, f'
  is the answer. 5. The solution of each edge e is l_e+f_e, where f_e corresponds to the
     flow of edge e on the graph.
 Construct minimum vertex cover from maximum matching M on
  bipartite graph (X, Y)
1. Redirect every edge: y \to x if (x, y) \in M, x \to y otherwise.

2. DFS from unmatched vertices in X.

3. x \in X is chosen iff x is unvisited.

4. y \in Y is chosen iff y is visited.

• Minimum cost cyclic flow
   1. Consruct super source S and sink T
  2. For each edge (x, y, c), connect x \to y with (cost, cap) = (c, 1) if 47
  c>0, otherwise connect y\to x with (cost,cap)=(-c,1) 3. For each edge with c<0, sum these cost as K, then increase d(y) 49
     by 1, decrease d(x) by 1
  4. For each vertex v with d(v) > 0, connect S \to v with (cost, cap) = 51
      (0, d(v))
  5. For each vertex v with d(v) < 0, connect v \to T with (cost, cap) = 53 };
```

capacity w5. For $v \in G$, connect it with sink $v \to t$ with capacity K + 2T - $(\sum_{e \in E(v)} w(e)) - 2w(v)$ 6. T is a valid answer if the maximum flow f < K|V|

6. Flow from S to T, the answer is the cost of the flow C+K• Maximum density induced subgraph

1. Binary search on answer, suppose we're checking answer T 2. Construct a max flow model, let K be the sum of all weights

3. Connect source $s \to v$, $v \in G$ with capacity K4. For each edge (u, v, w) in G, connect $u \to v$ and $v \to u$ with

• Minimum weight edge cover

(0,-d(v))

1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight

2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v. 3. Find the minimum weight perfect matching on G'.

Project selection problem

1. If $p_v > 0$, create edge (s, v) with capacity p_v ; otherwise, create edge (v,t) with capacity $-p_v$

2. Create edge (u, v) with capacity w with w being the cost of choosing u without choosing v.

3. The mincut is equivalent to the maximum profit of a subset of

projects.
• 0/1 quadratic programming

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})$$
23
25

can be minimized by the mincut of the following graph: 27 1. Create edge (x,t) with capacity c_x and create edge (s,y) with capacity c_y

2. Create edge (x,y) with capacity c_{xy} . 3. Create edge (x,y) and edge (x',y') with capacity $c_{xyx'y'}$.

3.4. Matching/Flows

3.4.1. Dinic's Algorithm

```
struct Dinic {
      struct edge {
         int to, cap, flow, rev;
      static constexpr int MAXN = 1000, MAXF = 1e9;
      vector<edge> v[MAXN];
      int top[MAXN], deep[MAXN], side[MAXN], s, t;
      void make_edge(int s, int t, int cap) {
  v[s].push_back({t, cap, 0, (int)v[t].size()});
  v[t].push_back({s, 0, 0, (int)v[s].size() - 1});
11
      int dfs(int a, int flow) {
   if (a == t || !flow) return flow;
   for (int &i = top[a]; i < v[a].size(); i++) {</pre>
13
           edge &e = v[a][i];
15
           if (deep[a] + 1 == deep[e.to] \&\& e.cap - e.flow) {
              int x = dfs(e.to, min(e.cap - e.flow, flow));
17
              if (x) {
                e.flow += x, v[e.to][e.rev].flow -= x;
19
                 return x;
           }
         deep[a] = -1;
         return 0;
      bool bfs() {
         queue<int> q;
         fill_n(deep, MAXN, 0);
         q.push(s), deep[s] = 1;
         int tmp;
         while (!q.empty()) {
           tmp = q.front(), q.pop();
for (edge e : v[tmp])
   if (!deep[e.to] && e.cap != e.flow)
                deep[e.to] = deep[tmp] + 1, q.push(e.to);
         return deep[t];
      int max_flow(int _s, int _t) {
         s = _s, t = _t;
int flow = 0, tflow;
         while (bfs()) {
           fill_n(top, MAXN, 0);
           while ((tflow = dfs(s, MAXF))) flow += tflow;
         return flow;
      void reset() {
         fill_n(side, MAXN, 0);
         for (auto &i : v) i.clear();
```

3.4.2. Minimum Cost Flow

```
struct MCF {
  struct edge {
    ll to, from, cap, flow, cost, rev;
    * fromE[MAXN]
  vector<edge> v[MAXN];
  ll n, s, t, flows[MAXN], dis[MAXN], pi[MAXN], flowlim;
  void make_edge(int s, int t, ll cap, ll cost) {
     if (!cap) return;
    v[s].pb(edge{t, s, cap, 0LL, cost, v[t].size()});
v[t].pb(edge{s, t, 0LL, 0LL, -cost, v[s].size() - 1});
  bitset<MAXN> vis;
  void dijkstra() {
    vis.reset();
      _gnu_pbds::priority_queue<pair<ll, <mark>int</mark>>> q;
     vector<decltype(q)::point_iterator> its(n);
    q.push({0LL, s})
    while (!q.empty()) {
       int now = q.top().second;
       q.pop();
       if (vis[now]) continue;
       vis[now] = 1;
ll ndis = dis[now] + pi[now];
       for (edge &e : v[now]) {
         if (e.flow == e.cap || vis[e.to]) continue;
         if (dis[e.to] > ndis + e.cost - pi[e.to]) {
           dis[e.to] = ndis + e.cost - pi[e.to];
           flows[e.to] = min(flows[now], e.cap - e.flow);
           fromE[e.to] = &e;
           if (its[e.to] == q.end())
              its[e.to] = q.push({-dis[e.to], e.to});
```

31

13

21

```
q.modify(its[e.to], {-dis[e.to], e.to});
            }
35
          }
        }
37
      bool AP(ll &flow) {
        fill_n(dis, n, INF);
fromE[s] = 0;
39
        dis[s] = 0;
41
        flows[s] = flowlim - flow;
        dijkstra();
43
        if (dis[t] == INF) return false;
45
        flow += flows[t];
        for (edge *e = fromE[t]; e; e = fromE[e->from]) {
  e->flow += flows[t];
47
          v[e->to][e->rev].flow -= flows[t];
49
        for (int i = 0; i < n; i++)
          pi[i] = min(pi[i] + dis[i], INF);
51
        return true;
53
      pll solve(int _s, int _t, ll _flowlim = INF) {
   s = _s, t = _t, flowlim = _flowlim;
}
        pll re;
        while (re.F != flowlim && AP(re.F))
        for (int i = 0; i < n; i++)
for (edge &e : v[i])</pre>
            if (e.flow != 0) re.S += e.flow * e.cost;
61
        re.S /= 2;
63
        return re:
      }
65
      void init(int _n) {
        n = _n;
fill_n(pi, n, 0);
67
        for (int i = 0; i < n; i++) v[i].clear();
69
      void setpi(int s) {
        fill_n(pi, n, INF);
        pi[s] = 0;
        for (ll it = 0, flag = 1, tdis; flag && it < n; it++) {
           flag = 0;
           for (int i = 0; i < n; i++)
             if (pi[i] != INF)
               for (edge \delta e : v[i])
77
                 if (e.cap && (tdis = pi[i] + e.cost) < pi[e.to])
79
                    pi[e.to] = tdis, flag = 1;
81
      }
   };
```

3.4.3. Gomory-Hu Tree

Requires: Dinic's Algorithm

```
int e[MAXN][MAXN];
int p[MAXN];
Dinic D; // original graph
void gomory_hu() {
  fill(p, p + n, 0);
  fill(e[0], e[n], INF);
  for (int s = 1; s < n; s++) {
    int t = p[s];
    Dinic F = D;
    int tmp = F.max_flow(s, t);
  for (int i = 1; i < s; i++)
        e[s][i] = e[i][s] = min(tmp, e[t][i]);
  for (int i = s + 1; i <= n; i++)
        if (p[i] == t && F.side[i]) p[i] = s;
}
</pre>
```

3.4.4. Global Minimum Cut

```
// weights is an adjacency matrix, undirected
   pair<int, vi> getMinCut(vector<vi> &weights) {
     int N = sz(weights);
     vi used(N), cut, best_cut;
     int best_weight = -1;
     for (int phase = N - 1; phase >= 0; phase--) {
       vi w = weights[0], added = used;
       int prev, k = 0;
       rep(i, 0, phase) {
         prev = k;
         k = -1;
13
         rep(j, 1, N) if (!added[j] &&
                           (k == -1 \mid \mid w[j] > w[k])) k = j;
         if (i == phase - 1) {
15
           rep(j, 0, N) weights[prev][j] += weights[k][j];
```

```
rep(j, 0, N) weights[j][prev] = weights[prev][j];
used[k] = true;
19
            cut.push_back(k);
            if (best_weight == -1 || w[k] < best_weight) {
21
              best cut = cut;
              best_weight = w[k];
23
          } else {
25
            rep(j, 0, N) w[j] += weights[k][j];
            added[k] = true;
27
       }
29
     }
     return {best_weight, best_cut};
31 }
```

3.4.5. Bipartite Minimum Cover

Requires: Dinic's Algorithm

```
1 // maximum independent set = all vertices not covered
    // x : [0, n), y : [0, m]
 3 struct Bipartite_vertex_cover {
      Dinic D:
      int n, m, s, t, x[maxn], y[maxn];
void make_edge(int x, int y) { D.make_edge(x, y + n, 1); }
      int matching() {
         int re = D.max_flow(s, t);
         for (int i = 0; i < n; i++)
 9
           for (Dinic::edge &e : D.v[i])
11
              if (e.to != s && e.flow == 1) {
                x[i] = e.to - n, y[e.to - n] = i;
13
                break;
             }
15
         return re;
      // init() and matching() before use
17
      void solve(vector<int> &vx, vector<int> &vy) {
         bitset<maxn * 2 + 10> vis;
        queue<int> q;
for (int i = 0; i < n; i++)
    if (x[i] == -1) q.push(i), vis[i] = 1;</pre>
21
         while (!q.empty())
23
           int now = q.front();
25
           q.pop();
           if (now < n) {
              for (Dinic::edge \delta e : D.v[now])
27
                if (e.to != s && e.to - n != x[now] && !vis[e.to])
29
                  vis[e.to] = 1, q.push(e.to);
           } else {
31
              if (!vis[y[now - n]])
                vis[y[now - n]] = 1, q.push(y[now - n]);
33
35
         for (int i = 0; i < n; i++)
         if (!vis[i]) vx.pb(i);
for (int i = 0; i < m; i++)
37
           if (vis[i + n]) vy.pb(i);
39
      void init(int _n, int _m) {
  n = _n, m = _m, s = n + m, t = s + 1;
  for (int i = 0; i < n; i++)</pre>
41
         x[i] = -1, D.make_edge(s, i, 1);
for (int i = 0; i < m; i++)</pre>
43
45
           y[i] = -1, D.make_edge(i + n, t, 1);
      }
47 };
```

3.4.6. Edmonds' Algorithm

```
struct Edmonds {
      int n, T;
      vector<vector<int>> g;
      vector<int> pa, p, used, base;
      Edmonds(int n)
 5
           : n(n), T(\theta), g(n), pa(n, -1), p(n), used(n),
             base(n) {}
      void add(int a, int b) {
9
        g[a].push_back(b);
        g[b].push_back(a);
11
      int getBase(int i) {
13
        while (i != base[i])
          base[i] = base[base[i]], i = base[i];
        return i;
15
17
      vector<int> toJoin;
      void mark_path(int v, int x, int b, vector<int> &path) {
  for (; getBase(v) != b; v = p[x]) {
    p[v] = x, x = pa[v];
19
```

```
toJoin.push_back(v);
           toJoin.push_back(x)
23
           if (!used[x]) used[x] = ++T, path.push_back(x);
25
      bool go(int v) {
27
        for (int x : g[v]) {
           int b, bv = getBase(v), bx = getBase(x);
if (bv == bx) {
29
             continue:
           } else if (used[x]) {
31
             vector<int> path;
33
             toJoin.clear();
             if (used[bx] < used[bv])</pre>
35
               mark_path(v, x, b = bx, path);
37
               mark_path(x, v, b = bv, path);
              for (int z : toJoin) base[getBase(z)] = b;
             for (int z : path)
               if (go(z)) return 1;
           } else if (p[x] == -1) {
41
             p[x] = v;
             p[x] = v,

if (pa[x] == -1) {

for (int y; x != -1; x = v)

y = p[x], v = pa[y], pa[x] = y, pa[y] = x;
43
45
                return 1;
47
             if (!used[pa[x]]) {
               used[pa[x]] = ++T;
49
                if (go(pa[x])) return 1;
51
          }
        }
53
        return 0;
55
      void init_dfs() {
        for (int i = 0; i < n; i++)
used[i] = 0, p[i] = -1, base[i] = i;
59
      bool dfs(int root) {
        used[root] = ++T;
        return go(root);
63
      void match() {
65
        int ans = 0;
        for (int v = 0; v < n; v++)
for (int x : g[v])
if (pa[v] == -1_88 pa[x] == -1) {
67
69
                pa[v] = x, pa[x] = v, ans++;
                break;
             }
71
        init_dfs();
        for (int i = 0; i < n; i++)
        if (pa[i] == -1 88 dfs(i)) ans++, init_dfs(); cout << ans * 2 << "\n";
        for (int i = 0; i < n; i++)
           if (pa[i] > i)
             cout << i + 1 << " " << pa[i] + 1 << "\n";
79
   };
```

3.4.7. Minimum Weight Matching

```
struct Graph {
      static const int MAXN = 105;
      int n, e[MAXN][MAXN];
      int match[MAXN], d[MAXN], onstk[MAXN];
      vector<int> stk;
      void init(int _n) {
        n = n:
        for (int i = 0; i < n; i++)
for (int j = 0; j < n; j++)
             // change to appropriate infinity
// if not complete graph
11
             e[i][j] = 0;
13
      void add_edge(int u, int v, int w) {
        e[u][v] = e[v][u] = w;
15
      bool SPFA(int u) {
17
        if (onstk[u]) return true;
19
        stk.push_back(u);
        onstk[u] = 1;
        for (int v = 0; v < n; v++) {
           if (u != v && match[u] != v && !onstk[v]) {
             int m = match[v];
if (d[m] > d[u] - e[v][m] + e[u][v]) {
23
               d[m] = d[u] - e[v][m] + e[u][v];
onstk[v] = 1;
25
               stk.push_back(v);
```

```
if (SPFA(m)) return true;
29
                stk.pop_back();
                onstk[v] = 0;
31
33
         onstk[u] = 0;
        stk.pop_back();
return false;
35
37
      int solve() {
        for (int i = 0; i < n; i += 2) {
    match[i] = i + 1;
39
41
           match[i + 1] = i;
43
         while (true) {
           int found = 0;
           for (int i = 0; i < n; i++) onstk[i] = d[i] = 0; for (int i = 0; i < n; i++) {
45
47
              stk.clear();
              if (!onstk[i] && SPFA(i)) {
49
                while (stk.size() >= 2) {
                   int u = stk.back();
51
                   stk.pop_back();
int v = stk.back();
53
                  stk.pop_back();
match[u] = v;
55
                  match[v] = u;
57
             }
59
           if (!found) break;
61
         int ret = 0;
63
         for (int i = 0; i < n; i++) ret += e[i][match[i]];
         ret /= 2;
65
         return ret;
67 } graph;
```

3.4.8. Stable Marriage

```
// normal stable marriage problem
    /* input:
 3
    3
    Albert Laura Nancy Marcy
   Brad Marcy Nancy Laura
    Chuck Laura Marcy Nancy
    Laura Chuck Albert Brad
    Marcy Albert Chuck Brad
    Nancy Brad Albert Chuck
    using namespace std;
    const int MAXN = 505;
   int favor[MAXN][MAXN]; // favor[boy_id][rank] = girl_id;
int order[MAXN][MAXN]; // order[girl_id][boy_id] = rank;
int current[MAXN]; // current[boy_id] = rank;
// boy_id will pursue current[boy_id] girl.
17
    int girl_current[MAXN]; // girl[girl_id] = boy_id;
23
   void initialize() {
      for (int i = 0; i < n; i++) {
    current[i] = 0;
25
         girl_current[i] = n;
27
         order[i][n] = n;
29 }
31
   map<string, int> male, female;
    string bname[MAXN], gname[MAXN];
   int fit = 0;
35 void stable_marriage() {
37
       queue<int> que;
       for (int i = 0; i < n; i++) que.push(i);
39
       while (!que.empty()) {
         int boy_id = que.front();
         que.pop();
         int girl_id = favor[boy_id][current[boy_id]];
         current[boy_id]++;
         if (order[girl_id][boy_id] <
   order[girl_id][girl_current[girl_id]]) {
  if (girl_current[girl_id] < n)</pre>
```

```
que.push(girl_current[girl_id]);
          girl_current[girl_id] = boy_id;
        } else {
          que.push(boy_id);
53
     }
55 }
57
   int main() {
     cin >> n:
59
     for (int i = 0; i < n; i++) {
61
        string p, t;
        cin >> p;
male[p] = i;
63
        bname[i] = p;
65
        for (int j = 0; j < n; j++) {
          cin >> t;
          if (!female.count(t)) {
            gname[fit] = t;
69
            female[t] = fit++;
          favor[i][j] = female[t];
71
73
      for (int i = 0; i < n; i++) {
        string p, t;
        for (int j = 0; j < n; j++) {
          order[female[p]][male[t]] = j;
81
     initialize();
85
     stable marriage():
     for (int i = 0; i < n; i++) {
87
        cout << bname[i] <<</pre>
             << gname[favor[i][current[i] - 1]] << endl;</pre>
89
91 }
```

3.4.9. Kuhn-Munkres algorithm

```
// Maximum Weight Perfect Bipartite Matching
    // Detect non-perfect-matching:

    set all edge[i][j] as INF

   // 2. if solve() >= INF, it is not perfect matching.
   typedef long long ll;
   struct KM {
      static const int MAXN = 1050;
      static const ll INF = 1LL << 60;
int n, match[MAXN], vx[MAXN], vy[MAXN];
      ll edge[MAXN][MAXN], lx[MAXN], ly[MAXN], slack[MAXN];
11
      void init(int _n) {
        n = _n;
for (int i = 0; i < n; i++)</pre>
13
          for (int j = 0; j < n; j++) edge[i][j] = 0;
15
      void add_edge(int x, int y, ll w) { edge[x][y] = w; }
17
      bool DFS(int x) {
19
        vx[x] = 1;
        for (int y
                     = 0; y < n; y++) {
          if (vy[y]) continue;
          if (lx[x] + ly[y] > edge[x][y]) {
             slack[y] =
23
             min(slack[y], lx[x] + ly[y] - edge[x][y]);
25
          } else {
             vy[y] = 1;
             if (match[y] == -1 || DFS(match[y])) {
               match[y] = x;
29
               return true:
            }
          }
31
        }
        return false;
33
35
      ll solve() {
        fill(match, match + n,
                                  -1);
37
        fill(lx, lx + n, -INF);
        fill(ly, ly + n, \theta);
        for (int i = 0; i < n; i++)
for (int j = 0; j < n; j++)
39
        lx[i] = max(lx[i], edge[i][j]);
for (int i = 0; i < n; i++) {</pre>
          fill(slack, slack + n, INF);
43
          while (true) {
45
             fill(vx, vx + n, 0);
```

```
fill(vy, vy + n, 0);
47
               if (DFS(i)) break;
49
               for (int j
                              = 0; j < n; j++)
              if (!vy[j]) d = min(d, slack[j]);
for (int j = 0; j < n; j++) {
  if (vx[j]) lx[j] -= d;</pre>
51
                 if (vy[j])
53
                    ly[j] += d;
55
                 else
                    slack[j] -= d;
              }
57
           }
59
         il res = 0;
         for (int i = 0; i < n; i++) {
61
           res += edge[match[i]][i];
63
         return res;
65
    } graph;
```

3.5. Shortest Path Faster Algorithm

```
1 struct SPFA {
      static const int maxn = 1010, INF = 1e9;
      int dis[maxn];
      bitset<maxn> inq, inneg;
      queue<int> q, tq;
vector<pii> v[maxn];
      void make_edge(int s, int t, int w) {
        v[s].emplace_back(t, w);
 9
      void dfs(int a) {
        inneg[a] = 1;
for (pii i : v[a])
11
13
           if (!inneg[i.F]) dfs(i.F);
      bool solve(int n, int s) { // true if have neg-cycle
for (int i = 0; i <= n; i++) dis[i] = INF;</pre>
15
17
         dis[s] = 0, q.push(s);
         for (int i = 0; i < n; i++) {
19
           inq.reset();
           int now;
21
           while (!q.empty()) {
             now = q.front(), q.pop();
23
             for (pii \delta i : v[now]) {
               if (dis[i.F] > dis[now] + i.S) {
  dis[i.F] = dis[now] + i.S;
25
                  if (!inq[i.F]) tq.push(i.F), inq[i.F] = 1;
27
               }
             }
29
           q.swap(tq);
31
        bool re = !q.empty();
33
        inneg.reset();
        while (!q.empty()) {
           if (!inneg[q.front()]) dfs(q.front());
35
           q.pop();
37
        return re;
39
      void reset(int n) {
        for (int i = 0; i <= n; i++) v[i].clear();
43 };
```

3.6. Strongly Connected Components

```
1 struct strongly_connected_components {
      int n. component count:
      vector<vector<int
>> edges, reverse_edges;
 3
      vector<int> component, component_size;
      {\tt strongly\_connected\_components(int n)}
 5
          : n(n), edges(n + 1), reverse_edges(n + 1),
component(n + 1), component_size(n + 1),
component_count(0) {}
 9
      void add_edge(int x, int y) {
        edges[x].push_back(y);
11
        reverse_edges[y].push_back(x);
13
      void dfs(int c, vector<int> &out_order,
                vector<bool> &done) {
        done[c] = true;
        for (int i : edges[c])
          if (!done[i]) dfs(i, out_order, done);
        out_order.push_back(c);
      void dfs_reverse(int c, vector<bool> 8done) {
```

```
component[c] = component_count;
         ++component_size[component_count];
for (int i : reverse_edges[c])
25
           if (!done[i]) dfs_reverse(i, done);
      void build_scc() {
  vector<int> out_order;
  vector<bool> done(n + 1, false),
27
29
         done_reverse(n + 1, false);
for (int i = 1; i <= n; ++i)</pre>
31
           if (!done[i]) dfs(i, out_order, done);
33
         component_count = 0;
         for (int i = out\_order.size() - 1; i >= 0; --i)
            if (!done_reverse[out_order[i]]) {
              ++component_count;
37
              component_size[component_count] = 0;
              dfs_reverse(out_order[i], done_reverse);
39
41 };
```

3.6.1. 2-Satisfiability

23 Requires: Strongly Connected Components

```
// 1 based, vertex in SCC = MAXN * 2
   // (not i) is i + n
  struct two_SAT {
    int n, ans[MAXN];
    SCC S:
    void imply(int a, int b) { S.make_edge(a, b); }
    bool solve(int _n) {
      n = _n;
      11
        ans[i] = (S.scc[i] < S.scc[i + n]);
13
      return true;
    }
15
    void init(int _n) {
      fill_n(ans, n + 1, 0);
19
      S.init(n * 2);
  } SAT;
```

3.7. Biconnected Components

3.7.1. Articulation Points

```
void dfs(int v, int p = -1) {
  visited[v] = true;
      tin[v] = low[v] = timer++;
      int children = 0;
      for (int to : adj[v]) {
        if (to == p) continue;
if (visited[to]) {
           low[v] = min(low[v], tin[to]);
           dfs(to, v);
low[v] = min(low[v], low[to]);
           if (low[to] >= tin[v] \delta\delta p != -1) IS_CUTPOINT(v);
13
      if (p == -1 && children > 1) IS_CUTPOINT(v);
17 }
```

3.7.2. Bridges

```
// if there are multi-edges, then they are not bridges
   void dfs(int v, int p = -1) {
  visited[v] = true;
                                                                            13
      tin[v] = low[v] = timer++;
     for (int to : adj[v]) {
        if (to == p) continue;
        if (visited[to]) {
                                                                            17
          low[v] = min(low[v], tin[to]);
                                                                            19
        } else {
          dfs(to, v);
low[v] = min(low[v], low[to]);
                                                                            21
11
          if (low[to] > tin[v]) IS_BRIDGE(v, to);
                                                                            23
13
     }
15 }
```

3.7.3. Block Cut Tree

9

11

13

15

17

19

21

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29

31

33

35

39

41

43

45

47

49

51

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57

59

5

9

11

15

```
1 struct block_cut_tree {
     int n, edge_id, bct_node_cnt;
vector<vector<array<int, 2>>> edges;
     vector<vector<int>> edges_bct;
     vector<bool> is_split;
     block_cut_tree(int n)
         : n(n), edges(n + 1), edges_bct(n * 3 + 1),
edge_id(0) {}
     void add_edge(int x, int y) {
  edges[x].push_back({y, edge_id});
       edges[y].push_back({x, edge_id});
        ++edge id;
       is_split.push_back(false);
     void dfs_split(int c, int pid, int &timer,
                       vector<int> &tin, vector<int> &low,
                      vector<bool> &done) {
       done[c] = true;
       tin[c] = low[c] = ++timer;
       for (auto i : edges[c])
         if (i[1] != pid) {
            if (done[i[0]])
              low[c] = min(low[c], tin[i[0]]);
            else {
              dfs_split(i[0], i[1], timer, tin, low, done);
low[c] = min(low[c], low[i[0]]);
is_split[i[1]] = tin[c] <= low[i[0]];</pre>
         }
     void dfs_build_bct(int c, int g, int &timer,
                            vector<bool> &done) {
       done[c] = true;
       edges_bct[c].push_back(g + n);
       edges_bct[g + n].push_back(c);
for (auto i : edges[c])
         if (!done[i[0]]) {
            if (is_split[i[1]]) {
              ++timer;
              edges_bct[c].push_back(timer + n);
edges_bct[timer + n].push_back(c);
              dfs_build_bct(i[0], timer, timer, done);
            } else
              dfs_build_bct(i[0], g, timer, done);
         }
     void build_bct() {
       int timer = 0;
       vector<int> tin(n + 1), low(n + 1);
       vector<bool> done_dfs(n + 1), done_build(n + 1);
       bct_node_cnt = 0;
       for (int i = 1; i
                             <= n; ++i)
         if (!done_dfs[i]) {
            dfs_split(i, -1, timer, tin, low, done_dfs);
            ++bct_node_cnt;
            dfs_build_bct(i, bct_node_cnt, bct_node_cnt,
                            done_build);
    }
  };
```

Triconnected Components

```
1 // requires a union-find data structure
  struct ThreeEdgeCC {
     int V, ind;
     vector<int> id, pre, post, low, deg, path;
vector<vector<int>> components;
     UnionFind uf:
     template <class Graph>
     void dfs(const Graph &G, int v, int prev) {
       pre[v] = ++ind;
       for (int w : G[v])
         if (w != v) {
            if (w == prev) {
              prev = -1;
              continue;
            if (pre[w] != -1) {
              if (pre[w] < pre[v]) {
                 low[v] = min(low[v], pre[w]);
              } else {
                 deg[v]--;
                 int &u = path[v];
for (; u != -1 && pre[u] <= pre[w] &&
                         pre[w] <= post[u];) {</pre>
                   uf.join(v, u);
deg[v] += deg[u];
```

25

```
u = path[u];
                 }
29
               continue;
31
            dfs(G, w, v);
            if (path[w] == -1 88 deg[w] <= 1) {
  deg[v] += deg[w];
  low[v] = min(low[v], low[w]);</pre>
33
35
               continue:
37
             if (deg[w] == 0) w = path[w];
             if (low[v] > low[w]) {
39
               low[v] = min(low[v], low[w]);
               swap(w, path[v]);
             for (; w = -1; w = path[w]) {
               uf.join(v, w);
               deg[v] += deg[w];
47
          }
        post[v] = ind;
      template <class Graph>
      ThreeEdgeCC(const Graph &G)
          : V(G.size()), ind(-1), id(V, -1), pre(V, -1),
            post(V), low(V, INT_MAX), deg(V, 0), path(V, -1),
53
            uf(V) {
        for (int v = θ; v < V; v++)
if (pre[v] == -1) dfs(G, v, -1);
55
57
        components.reserve(uf.cnt);
        for (int v = 0; v < V; v++)
          if (uf.find(v) == v) {
59
            id[v] = components.size();
             components.emplace_back(1, v);
61
             components.back().reserve(uf.getSize(v));
        for (int v = 0; v < V; v++)
          if (id[v] == -1)
            components[id[v] = id[uf.find(v)]].push_back(v);
   };
```

3.9. Centroid Decomposition

```
void get_center(int now) {
     v[now] = true:
     vtx.push_back(now);
     sz[now] = 1;
     mx[now] = 0;
     for (int u : G[now])
  if (!v[u]) {
          get_center(u);
          mx[now] = max(mx[now], sz[u]);
          sz[now] += sz[u];
11
   void get_dis(int now, int d, int len) {
13
     dis[d][now] = cnt;
      v[now] = true;
      for (auto u : G[now])
        if (!v[u.first]) { get_dis(u, d, len + u.second); }
17
19
   void dfs(int now, int fa, int d) {
     get_center(now);
     int c = -1;
for (int i : vtx) {
21
        if (max(mx[i], (int)vtx.size() - sz[i]) <=</pre>
23
            (int)vtx.size() / 2)
          c = i;
25
        v[i] = false;
     }
27
      get_dis(c, d, 0);
29
      for (int i : vtx) v[i] = false;
     v[c] = true;
31
     vtx.clear();
     dep[c] = d;
33
     p[c] = fa;
      for (auto u : G[c])
        if (u.first != fa && !v[u.first]) {
          dfs(u.first, c, d + 1);
   }
```

3.10. Minimum Mean Cycle

```
1 // d[i][j] == 0 if {i,j} !in E long long d[1003][1003], dp[1003][1003]; 61

pair<long long, long long> MMWC() {
```

```
memset(dp, 0x3f, sizeof(dp));
      for (int i = 1; i \le n; ++i) dp[0][i] = 0;
      for (int i = 1; i <= n; ++i) {
        for (int j = 1; j <= n; ++j) {
  for (int k = 1; k <= n; ++k) {
    dp[i][k] = min(dp[i - 1][j] + d[j][k], dp[i][k]);</pre>
11
        }
13
      15
        long long u = 0, d = 1;
17
        for (int j = n - 1; j >= 0; --j) {
  if ((dp[n][i] - dp[j][i]) * d > u * (n - j)) {
19
             u = dp[n][i] - dp[j][i];
21
          }
23
        if (u * ad < au * d) au = u, ad = d;
25
      long long g =
                        _gcd(au, ad);
      return make_pair(au / g, ad / g);
```

3.11. Directed MST

```
1 template <typename T> struct DMST {
      T g[maxn][maxn], fw[maxn];
      int n, fr[maxn];
      bool vis[maxn], inc[maxn];
      void clear() {
        for (int i = 0; i < maxn; ++i) {
  for (int j = 0; j < maxn; ++j) g[i][j] = inf;</pre>
          vis[i] = inc[i] = false;
 9
11
      void addedge(int u, int v, T w) {
        g[u][v] = min(g[u][v], w);
13
      T operator()(int root, int _n) {
        if (dfs(root) != n) return -1;
17
        T ans = 0:
        while (true) {
          for (int i = 1; i <= n; ++i) fw[i] = inf, fr[i] = i;
          for (int i = 1; i <= n; ++i)
             if (!inc[i]) {
21
               for (int j = 1; j <= n; ++j) {
  if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
23
                   fw[i] = g[j][i];
                    fr[i] = j;
25
27
               }
             }
          int x = -1;
29
           for (int i = 1; i <= n; ++i)
             if (i != root && !inc[i]) {
31
               int j = i, c = 0;
               while (j != root && fr[j] != i && c <= n)
33
               ++c, j = fr[j];
if (j == root || c > n)
                 continue;
37
               else {
39
                 break;
               }
41
          if (!~x) {
43
             for (int i = 1; i \le n; ++i)
               if (i != root δδ !inc[i]) ans += fw[i];
45
             return ans:
          }
47
          int v = x;
          for (int i = 1; i <= n; ++i) vis[i] = false;
          do {
49
             ans += fw[y];
51
             y = fr[y];
             vis[y] = inc[y] = true;
           } while (y != x);
53
          inc[x] = false;
55
           for (int k = 1; k \le n; ++k)
             if (vis[k]) {
               for (int j = 1; j \le n; ++j)
                 if (!vis[j]) {
59
                    if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
                    if (g[j][k] < inf &&
g[j][k] - fw[k] < g[j][x])
g[j][x] = g[j][k] - fw[k];
                 }
```

3.12. Maximum Clique

```
// source: KACTL
    typedef vector<br/>bitset<200>> vb;
    struct Maxclique {
      double limit = 0.025, pk = 0;
      struct Vertex {
        int i, d = 0;
      typedef vector<Vertex> vv;
11
      vector<vi> C;
      vi qmax, q, S, old;
      void init(vv &r) {
         for (auto \delta v : r) v.d = 0;
15
         for (auto &v : r)
17
           for (auto j : r) v.d += e[v.i][j.i];
         sort(all(r),
                        [](auto a, auto b) { return a.d > b.d; });
         int mxD = r[0].d;
19
        rep(i, \theta, sz(r)) r[i].d = min(i, mxD) + 1;
21
      void expand(vv δR, int lev = 1) {
   S[lev] += S[lev - 1] - old[lev];
   old[lev] = S[lev - 1];
23
         while (sz(R))
25
           if (sz(q) + R.back().d <= sz(qmax)) return;</pre>
           q.push_back(R.back().i);
           for (auto v: R)
29
             if (e[R.back().i][v.i]) T.push_back({v.i});
           if (sz(T)) {
              if (S[lev]++
                             / ++pk < limit) init(T);</pre>
              int j = 0, mxk = 1,
mnk = max(sz(qmax) - sz(q) + 1, 1);
33
             C[1].clear(), C[2].clear();
for (auto v : T) {
35
                int k = 1;
37
                auto f = [8](int i) { return e[v.i][i]; };
while (any_of(all(C[k]), f)) k++;
39
                if (k > mxk) mxk = k, C[mxk + 1].clear();
if (k < mnk) T[j++].i = v.i;</pre>
                C[k].push_back(v.i);
              if (j > 0) T[j - 1].d = 0;
             rep(k, mnk, mxk + 1) for (int i : C[k]) T[j].i = i,
                                                               T[j++].d =
           expand(T, lev + 1);
} else if (sz(q) > sz(qmax))
49
             amax = a:
           q.pop_back(), R.pop_back();
51
        }
53
      }
      vi maxClique() {
55
        init(V), expand(V);
        return qmax;
57
      Maxclique(vb conn)
59
           : e(conn), C(sz(e) + 1), S(sz(C)), old(S) {
         rep(i, 0, sz(e)) V.push_back({i});
61
```

3.13. Dominator Tree

```
// idom[n] is the unique node that strictly dominates n but
// does not strictly dominate any other node that strictly
// dominates n. idom[n] = 0 if n is entry or the entry
// cannot reach n.
struct DominatorTree {
   static const int MAXN = 200010;
   int n, s;
   vector<int> g[MAXN], pred[MAXN];
   vector<int> cov[MAXN];
   int dfn[MAXN], nfd[MAXN], ts;
```

```
int par[MAXN];
int sdom[MAXN], idom[MAXN];
int mom[MAXN], mn[MAXN];
inline bool cmp(int u, int v) { return dfn[u] < dfn[v]; }</pre>
int eval(int u) {
  if (mom[u] == u) return u;
  int res = eval(mom[u]);
  if (cmp(sdom[mn[mom[u]]), sdom[mn[u]]))
    mn[u] = mn[mom[u]];
  return mom[u] = res;
void init(int _n, int _s) {
  n = n;
 s = _s;
REP1(i, 1, n) {
  g[i].clear();
    pred[i].clear();
    idom[i] = 0;
void add_edge(int u, int v) {
  g[u].push_back(v);
  pred[v].push_back(u);
void DFS(int u) {
  dfn[u] = ts;
  nfd[ts] = u;
  for (int v : g[u])
    if (dfn[v] == 0) {
      par[v] = u;
      DFS(v);
    }
void build() {
  ts = 0;
  REP1(i, 1, n) {
    dfn[i] = nfd[i] = 0;
    cov[i].clear();
    mom[i] = mn[i] = sdom[i] = i;
  for (int i = ts; i >= 2; i--) {
    int u = nfd[i];
    if (u == 0) continue;
    for (int v : pred[u])
      if (dfn[v]) {
        eval(v):
        if (cmp(sdom[mn[v]], sdom[u]))
          sdom[u] = sdom[mn[v]];
    cov[sdom[u]].push_back(u);
    mom[u] = par[u];
for (int w : cov[par[u]]) {
      eval(w);
      if (cmp(sdom[mn[w]], par[u]))
        idom[w] = mn[w];
        idom[w] = par[u];
    cov[par[u]].clear();
  REP1(i, 2, ts) {
  int u = nfd[i];
    if (u == 0) continue;
    if (idom[u] != sdom[u]) idom[u] = idom[idom[u]];
```

3.14. Manhattan Distance MST

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} dom;

```
int j = it->second;
            Pd = ps[i] - ps[j];
            if (d.y > d.x) break;
19
            edges.push_back({d.y + d.x, i, j});
21
          sweep[-ps[i].y] = i;
        for (P &p : ps)
if (k & 1)
23
25
            p.x = -p.x;
          else
27
            swap(p.x, p.y);
     }
29
     return edges;
```

Math

4.1. Number Theory

4.1.1. Mod Struct

A list of safe primes: 26003, 27767, 28319, 28979, 29243, 29759, 30467 910927547, 919012223, 947326223, 990669467, 1007939579, 1019126699929760389146037459, 975500632317046523, 989312547895528379

```
NTT prime p
                        p-1
                                    primitive root
65537
                        1\ll 16
                                    3
998244353
                        119 \ll 23
                                    3
2748779069441
                        5 \ll 39
                                    3
1945555039024054273 \mid 27 \ll 56
                                    5
```

```
template <typename T> struct M {
       static T MOD; // change to constexpr if already known
       Tv;
       M(T x = 0) \{
          v = (-MOD \le x \&\& x \le MOD) ? x : x \% MOD;
          if (v < 0) v += MOD;
        explicit operator T() const { return v; }
       bool operator==(const M &b) const { return v == b.v; }
       bool operator!=(const M &b) const { return v != b.v; }
       M operator-() { return M(-v); }
11
       M operator+(M b) { return M(v + b.v); }
M operator-(M b) { return M(v - b.v); }
M operator*(M b) { return M(v - b.v); }
M operator*(M b) { return M((__int128)v * b.v % MOD); }
M operator/(M b) { return *this * (b ^ (MOD - 2)); }
15
        // change above implementation to this if MOD is not prime
       M inv() {
          auto [p, _, g] = extgcd(v, MOD);
19
          return assert(g == 1), p;
        friend M operator^(M a, ll b) {
21
          M ans(1);
23
          for (; b; b >>= 1, a *= a)
             if (b \delta 1) ans *= a;
25
          return ans:
       friend M & Soperator+=(M & a, M b) { return a = a + b; } friend M & Soperator-=(M & a, M b) { return a = a - b; } friend M & Soperator*=(M & a, M b) { return a = a * b; } friend M & Soperator*=(M & a, M b) { return a = a * b; }
27
29
       friend M & operator/=(M & a, M b) { return a = a / b; }
31
    };
    using Mod = M<int>;
    template <> int Mod::MOD = 1'000'000'007;
     int &MOD = Mod::MOD;
```

4.1.2. Miller-Rabin

Requires: Mod Struct

```
// checks if Mod::MOD is prime
     bool is_prime() {
        if (MOD < 2 || MOD % 2 == 0) return MOD == 2;
Mod A[] = {2, 7, 61}; // for int values (< 2^31)
// ll: 2, 325, 9375, 28178, 450775, 9780504, 1795265022
int s = __builtin_ctzll(MOD - 1), i;</pre>
         for (Mod a : A) {
   Mod x = a ^ (MOD >> s);
            for (i = 0; i < s \delta\delta (x + 1).v > 2; i++) x *= x;
            if (i && x != -1) return 0;
11
        return 1:
13 }
```

4.1.3. Linear Sieve

```
constexpr ll MAXN = 1000000;
   bitset<MAXN> is_prime;
   vector<ll> primes
   ll mpf[MAXN], phi[MAXN], mu[MAXN];
   void sieve() {
     is prime.set():
      is_prime[1] = 0;
     mu[1] = phi[1] = 1;
for (ll i = 2; i < MAXN; i++) {
        if (is_prime[i]) {
  mpf[i] = i;
11
13
          primes.push_back(i);
          phi[i] = i - 1;
          mu[i] = -1;
15
        for (ll p : primes) {
  if (p > mpf[i] || i * p >= MAXN) break;
17
19
          is\_prime[i * p] = 0;
          mpf[i * p] = p;
21
          mu[i * p] = -mu[i];
          if (i % p == 0)
23
             phi[i * p] = phi[i] * p, mu[i * p] = 0;
             phi[i * p] = phi[i] * (p - 1);
25
27
   }
```

4.1.4. Get Factors

Requires: Linear Sieve

```
vector<ll> fac = {1};
                          while (n > 1) {
  const ll p = mpf[n];
                            vector<ll> cur = {1};
                            while (n \% p == 0) \{
                             n /= p;
                              cur.push_back(cur.back() * p);
                            }
                      9
                            vector<ll> tmp;
                     11
                            for (auto x : fac)
                              for (auto y : cur) tmp.push_back(x * y);
                     13
                            tmp.swap(fac);
                          return fac;
                     15
                        }
```

4.1.5. Binary GCD

```
\ensuremath{//} returns the gcd of non-negative a, b
   ull bin_gcd(ull a, ull b) {
     if (!a || !b) return a + b;
     int s = __builtin_ctzll(a);
                _builtin_ctzll(a ˈ b);
     while (b) {
       if ((b >>= __builtin_ctzll(b)) < a) swap(a, b);</pre>
     return a << s;
11 }
```

4.1.6. Extended GCD

```
1 // returns (p, q, g): p * a + q * b == g == gcd(a, b)
    // g is not guaranteed to be positive when a < 0 or b < 0
3 tuple<ll, ll, ll> extgcd(ll a, ll b) {
    ll s = 1, t = 0, u = 0, v = 1;
    while (b) {
        ll q = a / b;
swap(a -= q * b, b);
         swap(s -= q * t, t);
         swap(u -= q * v, v);
      return {s, u, a};
11
```

4.1.7. Chinese Remainder Theorem

Requires: Extended GCD

```
1 // for 0 <= a < m, 0 <= b < n, returns the smallest x >= 0
   // such that x \% m == a and x \% n == b
3 | ll crt(ll a, ll m, ll b, ll n) {
     if (n > m) swap(a, b), swap(m, n);
    auto [x, y, g] = extgcd(m, n);
assert((a - b) % g == 0); // no solution
x = ((b - a) / g * x) % (n / g) * m + a;
     return x < 0 ? x + m / g * n : x;
```

4.1.8. Baby-Step Giant-Step

Requires: Mod Struct

```
1 // returns x such that a ^ x = b where x \in [l, r)
    ll bsgs(Mod a, Mod b, ll l = θ, ll r = MOD - 1) {
        int m = sqrt(r - l) + 1, i;
        unordered_map<ll, ll> tb;
        Mod d = (a ^ l) / b;
        for (i = θ, d = (a ^ l) / b; i < m; i++, d *= a)
        if (d == 1)
            return l + i;
        else
            tb[(ll)d] = l + i;
        Mod c = Mod(1) / (a ^ m);
        for (i = θ, d = 1; i < m; i++, d *= c)
        if (auto j = tb.find((ll)d); j != tb.end())
            return j->second + i * m;
        return assert(θ), -1; // no solution
    }
}
```

4.1.9. Pollard's Rho

4.1.10. Tonelli-Shanks Algorithm

Requires: Mod Struct

```
int legendre(Mod a) {
       if (a == 0) return 0;
return (a ^ ((MOD - 1) / 2)) == 1 ? 1 : -1;
    Mod sqrt(Mod a) {
       assert(legendre(a) != -1); // no solution
       ll p = MOD, s = p - 1;
if (a == 0) return 0;
       if (p == 2) return 1;
       if (p % 4 == 3) return a ^ ((p + 1) / 4);
       int r, m;
11
       for (r = 0; !(s & 1); r++) s >>= 1;
       Mod n = 2;
13
       while (legendre(n) != -1) n += 1;
Mod x = a ^ ((s + 1) / 2), b = a ^ s, g = n ^ s;
while (b != 1) {
15
17
          Mod t = b;
          for (m = 0; t != 1; m++) t *= t;
Mod gs = g ^ (1LL << (r - m - 1));
19
          g = gs * gs, x *= gs, b *= g, r = m;
       return x;
23 }
    // to get sqrt(X) modulo p^k, where p is an odd prime: 
// c = x^2 (mod p), c = X^2 (mod p^k), q = p^k-1) 
// X = x^2 * x^2 * x^2 (mod p^k)
```

4.1.11. Chinese Sieve

```
const ll N = 1000000;
   // f, g, h multiplicative, h = f (dirichlet convolution) g
ll pre_g(ll n);
   ll pre_h(ll n);
    // preprocessed prefix sum of f
   ll pre_f[N];
      prefix sum of multiplicative function f
   ll solve_f(ll n) {
      static unordered_map<ll, ll> m;
      if (n < N) return pre_f[n];</pre>
      if (m.count(n)) return m[n];
     ll ans = pre_h(n);
for (ll l = 2, r; l <= n; l = r + 1) {
    r = n / (n / l);</pre>
13
15
        ans -= (pre_g(r) - pre_g(l - 1)) * djs_f(n / l);
17
      return m[n] = ans;
```

4.1.12. Rational Number Binary Search

```
1 struct QQ {
       ll p, q;
       QQ go(QQ b, ll d) \{ return \{p + b.p * d, q + b.q * d\}; \}
 5 bool pred(QQ);
    // returns smallest p/q in [lo, hi] such that
    // pred(p/q) is true, and 0 <= p,q <= N
    QQ frac_bs(ll N) {
       QQ lo{0, 1}, hi{1, 0};
       if (pred(lo)) return lo;
       assert(pred(hi));
bool dir = 1, L = 1, H = 1;
11
       for (; L | | H; dir = !dir) {
13
         for (int t = 0; step = 1;
for (int t = 0; t < 2 && (t ? step /= 2 : step *= 2);)
   if (QQ mid = hi.go(lo, len + step);
      mid.p > N || mid.q > N || dir ^ pred(mid))
15
17
               t++;
19
            else
          len += step;
swap(lo, hi = hi.go(lo, len));
21
         (dir ? L : H) = !!len;
23
       return dir ? hi : lo;
25 }
```

4.1.13. Farey Sequence

```
// returns (e/f), where (a/b, c/d, e/f) are
// three consecutive terms in the order n farey sequence
// to start, call next_farey(n, 0, 1, 1, n)
pll next_farey(ll n, ll a, ll b, ll c, ll d) {
    ll p = (n + b) / d;
    return pll(p * c - a, p * d - b);
}
```

4.2. Combinatorics

4.2.1. Matroid Intersection

This template assumes 2 weighted matroids of the same type, and that removing an element is much more expensive than checking if one can be added. Remember to change the implementation details.

The ground set is $0, 1, \ldots, n-1$, where element *i* has weight w[i]. For the unweighted version, remove weights and change BF/SPFA to BFS.

```
1 constexpr int N = 100;
   constexpr int INF = 1e9;
 3
                             // represents an independent set
     Matroid(bitset<N>); // initialize from an independent set
                            // if adding will break independence
     bool can add(int);
     Matroid remove(int); // removing from the set
   auto matroid_intersection(int n, const vector<int> &w) {
     bitset<N> S;
11
     for (int sz = 1; sz <= n; sz++) {
        Matroid M1(S), M2(S);
13
        vector<vector<pii>>> e(n + 2);
15
        for (int j = 0; j < n; j++)
          if (!S[j]) {
17
            if (M1.can_add(j)) e[n].emplace_back(j, -w[j]);
19
            if (M2.can_add(j)) e[j].emplace_back(n + 1, 0);
        for (int i = 0; i < n; i++)
21
          if (S[i]) {
            Matroid T1 = M1.remove(i), T2 = M2.remove(i);
23
            for (int j = 0; j < n; j++)
25
              if (!S[j]) {
                if (T1.can_add(j)) e[i].emplace_back(j, -w[j]);
27
                 if (T2.can_add(j)) e[j].emplace_back(i, w[i]);
29
          }
31
        vector<pii> dis(n + 2, {INF, 0});
        vector<int> prev(n + 2, -1);
33
        dis[n] = \{0, 0\};
        // change to SPFA for more speed, if necessary
        bool upd = 1;
        while (upd) {
37
          upd = 0:
          for (int u = 0; u < n + 2; u++)
            for (auto [v, c] : e[u]) {
   pii x(dis[u].first + c, dis[u].second + 1);
   if (x < dis[v]) dis[v] = x, prev[v] = u, upd = 1;</pre>
39
```

```
43
                                                                          41
45
        if (dis[n + 1].first < INF)</pre>
                                                                          43
          for (int x = prev[n + 1]; x != n; x = prev[x])
            S.flip(x);
                                                                          45
        else
49
          break;
                                                                          47
51
       // S is the max-weighted independent set with size sz
                                                                          49
53
     return S:
                                                                          51
                                                                          53
```

4.2.2. De Brujin Sequence

```
55
   int res[kN], aux[kN], a[kN], sz;
                                                                            57
   void Rec(int t, int p, int n, int k) {
      if (t > n) {
        if (n \% p == 0)
                                                                            59
          for (int i = 1; i <= p; ++i) res[sz++] = aux[i];
                                                                            61
      } else {
        aux[t] = aux[t - p];
Rec(t + 1, p, n, k);
for (aux[t] = aux[t - p] + 1; aux[t] < k; ++aux[t])</pre>
                                                                            63
                                                                            65
          Rec(t + 1, t, n, k);
11
                                                                            67
13
   int DeBruijn(int k, int n) {
      // return cyclic string of length k^n such that every
                                                                            69
15
      // string of length n using k character appears as a
                                                                            71
      // substring.
     if (k == 1) return res[0] = 0, 1;
                                                                            73
      fill(aux, aux + k * n, \theta);
     return sz = 0, Rec(1, 1, n, k), sz;
                                                                            75
```

4.2.3. Multinomial

```
// ways to permute v[i]
  ll multinomial(vi &v) {
                                                                                     81
     ll c = 1, m = v.empty() ? 1 : v[0];
for (int i = 1; i < v.size(); i++)</pre>
                                                                                     83
       for (int j = 0; i < v[i]; j++) c = c * ++m / (j + 1);
     return c;
                                                                                     85
7 }
```

4.3. Algebra

4.3.1. Formal Power Series

```
template <typename mint>
   struct FormalPowerSeries : vector<mint> {
                                                                   93
     using vector<mint>::vector:
     using FPS = FormalPowerSeries;
     FPS &operator+=(const FPS &r) {
       if (r.size() > this->size()) this->resize(r.size());
       for (int i = 0; i < (int)r.size(); i++)
                                                                   99
         (*this)[i] += r[i];
       return *this;
                                                                  101
11
                                                                  103
     FPS \deltaoperator+=(const mint \deltar) {
       if (this->empty()) this->resize(1);
                                                                  105
       (*this)[0] += r;
15
       return *this;
                                                                  107
17
     109
19
                                                                  111
21
         (*this)[i] -= r[i];
                                                                  113
23
       return *this;
                                                                  115
     FPS & operator -= (const mint &r) {
                                                                  117
27
       if (this->empty()) this->resize(1);
       (*this)[0] -= r;
                                                                  119
29
       return *this;
                                                                  121
31
     FPS &operator*=(const mint &v) {
                                                                  123
       for (int k = 0; k < (int)this -> size(); k++)
         (*this)[k] *= v;
                                                                  125
35
       return *this;
                                                                  127
     FPS &operator/=(const FPS &r) {
                                                                  129
39
       if (this->size() < r.size()) {</pre>
```

```
this->clear():
    return *this;
  int n = this->size() - r.size() + 1;
  if ((int)r.size() <= 64) {
    FPS f(*this), g(r);
    g.shrink();
    mint coeff = g.back().inverse();
for (auto &x : g) x *= coeff;
int deg = (int)f.size() - (int)g.size() + 1;
    int gs = g.size();
    FPS quo(deg);
    for (int i = deg - 1; i >= 0; i--) {
  quo[i] = f[i + gs - 1];
  for (int j = 0; j < gs; j++)</pre>
        f[i + j] -= quo[i] * g[j];
    *this = quo * coeff;
    this->resize(n, mint(0));
    return *this;
  return *this = ((*this).rev().pre(n) * r.rev().inv(n))
                   .pre(n)
                   .rev():
}
FPS &operator%=(const FPS &r) {
  *this -= *this / r * r;
  shrink();
  return *this;
FPS operator+(const FPS &r) const {
 return FPS(*this) += r;
FPS operator+(const mint &v) const {
  return FPS(*this) += v;
FPS operator-(const FPS &r) const {
 return FPS(*this) -= r;
FPS operator-(const mint &v) const {
 return FPS(*this) -= v;
FPS operator*(const FPS &r) const {
 return FPS(*this) *= r;
FPS operator*(const mint &v) const {
  return FPS(*this) *= v;
FPS operator/(const FPS &r) const {
 return FPS(*this) /= r;
FPS operator%(const FPS &r) const {
 return FPS(*this) %= r;
FPS operator-() const {
  FPS ret(this->size());
  for (int i = 0; i < (int)this->size(); i++)
ret[i] = -(*this)[i];
 return ret:
}
void shrink() {
 while (this->size() && this->back() == mint(0))
    this->pop_back();
FPS rev() const {
  FPS ret(*this);
  reverse(begin(ret), end(ret));
  return ret;
FPS dot(FPS r) const {
  FPS ret(min(this->size(), r.size()));
  for (int i = 0; i < (int)ret.size(); i++)
    ret[i] = (*this)[i] * r[i];
  return ret:
FPS pre(int sz) const {
  return FPS(begin(*this),
              begin(*this) + min((int)this->size(), sz));
FPS operator>>(int sz) const {
  if ((int)this->size() <= sz) return {};</pre>
  FPS ret(*this);
  ret.erase(ret.begin(), ret.begin() + sz);
```

77

87

89

91

95

```
return ret:
131
133
        FPS operator<<(int sz) const {</pre>
          FPS ret(*this);
           ret.insert(ret.begin(), sz, mint(0));
135
           return ret;
137
        FPS diff() const {
  const int n = (int)this->size();
139
           FPS ret(max(0, n - 1));
141
           mint one(1), coeff(1);
           for (int i = 1; i < n; i++) {
  ret[i - 1] = (*this)[i] * coeff;</pre>
143
             coeff += one;
145
147
           return ret;
149
        FPS integral() const {
           const int n = (int)this->size();
           FPS ret(n + 1);
153
           ret[0] = mint(0);
           if (n > 0) ret[1] = mint(1);
          auto mod = mint::get_mod();
for (int i = 2; i <= n; i++)
  ret[i] = (-ret[mod % i]) * (mod / i);
for (int i = 0; i < n; i++) ret[i + 1] *= (*this)[i];</pre>
155
157
159
           return ret;
161
        mint eval(mint x) const {
163
           mint r = 0, w = 1;
           for (auto \delta v : *this) r += w * v, w *= x;
165
           return r;
167
        FPS log(int deg = -1) const {
           assert((*this)[0] == mint(1));
           if (deg == -1) deg = (int)this->size();
           return (this->diff() * this->inv(deg))
171
           .pre(deg - 1)
           .integral();
173
175
        FPS pow(int64_t k, int deg = -1) const {
  const int n = (int)this->size();
  if (deg == -1) deg = n;
  for (int i = 0; i < n; i++) {
    if ('ethis)[i] != mint(0)) {
        if ('i this)[i] |= mint(0) }</pre>
177
179
                if (i * k > deg) return FPS(deg, mint(\theta));
181
                mint rev = mint(1) / (*this)[i];
183
                FPS ret =
                (((*this * rev) >> i).log(deg) * k).exp(deg) *
                ((*this)[i].pow(k));
ret = (ret << (i * k)).pre(deg);
185
                if ((int)ret.size() < deg) ret.resize(deg, mint(0));</pre>
187
                return ret;
189
          return FPS(deg, mint(θ));
191
193
        static void *ntt_ptr;
static void set_fft();
195
        FPS &operator*=(const FPS &r);
197
        void ntt()
        void intt();
        void ntt_doubling();
199
        static int ntt_pr();
        FPS inv(int deg = -1) const;
        FPS exp(int deg = -1) const;
     template <typename mint>
     void *FormalPowerSeries<mint>::ntt_ptr = nullptr;
205
```

4.4. Theorems

4.4.1. Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i), \ L_{ij} = -c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

4.4.2. Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

4.4.3. Cayley's Formula

• Given a degree sequence d_1, d_2, \ldots, d_n for each labeled vertices, there are

$$\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$$

spanning trees.

• Let $T_{n,k}$ be the number of *labeled* forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.

4.4.4. Erdős-Gallai Theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1 + d_2 + \ldots + d_n$ is even and

$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k)$$

holds for all $1 \le k \le n$.

4.4.5. Burnside's Lemma

Let X be a set and G be a group that acts on X. For $g \in G$, denote by X^g the elements fixed by g:

$$X^g = \{x \in X \mid gx \in X\}$$

Then

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

5. Randomisation

5.1. Simulated Annealing

```
double get_rand_double() {
     return double(
3
            uniform_int_distribution<int>(0, INF)(rng)) /
5 }
   int main() {
     const int iterations = 10000000;
     double mpl = pow((1e-5) / temp, (double)1 / iterations);
     for (int i = 0; i < iterations; ++i) {
       t *= mpl;
       if (new_score >= score ||
           get_rand_double() <</pre>
13
           exp((new_score - score) / temp)) {
         score = new_score;
         return:
17
     }
19 }
```

6. Numeric

6.1. Barrett Reduction

```
using ull = unsigned long long;
using uL = __uint128_t;
// very fast calculation of a % m
struct reduction {
   const ull m, d;
   explicit reduction(ull m) : m(m), d(((uL)1 << 64) / m) {}
   inline ull operator()(ull a) const {
      ull q = (ull)(((uL)d * a) >> 64);
      return (a -= q * m) >= m ? a - m : a;
   }
};
```

6.2. Long Long Multiplication

```
using ull = unsigned long long;
using ll = long long;
using ld = long double;
// returns a * b % M where a, b < M < 2**63
ull mult(ull a, ull b, ull M) {
    ll ret = a * b - M * ull(ld(a) * ld(b) / ld(M));
    return ret * M * (ret < 0) - M * (ret >= (ll)M);
}
```

6.3. Fast Fourier Transform

```
template <typename T>
void fft_(int n, vector<T> &a, vector<T> &rt, bool inv) {
    vector<int> br(n);
    for (int i = 1; i < n; i++) {
        br[i] = (i & 1) ? br[i - 1] + n / 2 : br[i / 2] / 2;
        if (br[i] > i) swap(a[i], a[br[i]]);

}

for (int len = 2; len <= n; len *= 2)
    for (int i = 0; i < n; i += len)
        for (int j = 0; j < len / 2; j++) {
        int pos = n / len * (inv ? len - j : j);
        T u = a[i + j], v = a[i + j + len / 2] * rt[pos];
        a[i + j] = u + v, a[i + j + len / 2] = u - v;

}

if (T minv = T(1) / T(n); inv)
    for (T &x : a) x *= minv;
}</pre>
```

Requires: Mod Struct

void ntt(vector<Mod> &a, bool inv, Mod primitive_root) {
 int n = a.size();

Mod root = primitive_root ^ (MOD - 1) / n;
 vector<Mod> rt(n + 1, 1);
 for (int i = 0; i < n; i++) rt[i + 1] = rt[i] * root;
 fft_(n, a, rt, inv);

void fft(vector<complex<double>> &a, bool inv) {
 int n = a.size();
 vector<complex<double>> rt(n + 1);
 double arg = acos(-1) * 2 / n;
 for (int i = 0; i <= n; i++)
 rt[i] = {cos(arg * i), sin(arg * i)};
 fft_(n, a, rt, inv);

}

6.4. Fast Walsh-Hadamard Transform

```
Requires: Mod Struct 17

void fwht(vector<Mod> &a, bool inv) {
    int n = a.size();
    for (int d = 1; d < n; d <<= 1)
        for (int m = 0; m < n; m++)
        if (!(m & d)) {
            inv ? a[m] -= a[m | d] : a[m] += a[m | d]; // AND
            inv ? a[m | d] -= a[m] : a[m | d] += a[m]; // OR
            Mod x = a[m], y = a[m | d]; // XOR
        a[m] = x + y, a[m | d] = x - y; // XOR
        }

if (Mod iv = Mod(1) / n; inv) // XOR
        for (Mod &i : a) i *= iv; // XOR

29
```

6.5. Subset Convolution

13 }

```
Requires: Mod Struct
    #pragma GCC target("popcnt")
    #include <immintrin.h>
    void fwht(int n, vector<vector<Mod>> &a, bool inv) {
      for (int h = 0; h < n; h++)
for (int i = 0; i < (1 << n); i++)
            if (!(i & (1 << h)))
               for (int k = 0; k <= n; k++)
inv ? a[i | (1 << h)][k] -= a[i][k]
: a[i | (1 << h)][k] += a[i][k];
    // c[k] = sum(popcnt(i \& j) == sz \&\& i | j == k) a[i] * b[j]
    vector<Mod> subset_convolution(int n, int sz, const vector<Mod> &a_
13
                                             const vector<Mod> &b_) {
      int len = n + sz + 1, N = 1 << n;</pre>
       vector<vector<Mod>> a(1 << n, vector<Mod>(len, \theta)), b = a;
17
       for (int i = 0; i < N; i++)
a[i][_mm_popent_u64(i)] = a_[i],
19
         b[i][_mm_popcnt_u64(i)] = b_[i];
       fwht(n, a, 0), fwht(n, b, 0);
for (int i = 0; i < N; i++) {</pre>
21
         vector<Mod> tmp(len);
         for (int j = 0; j < len; j++)
for (int k = 0; k <= j; k++)
25
               tmp[j] += a[i][k] * b[i][j - k];
       fwht(n, a, 1)
      vector<Mod> c(N);
for (int i = 0; i < N; i++)
         c[i] = a[i][_mm_popcnt_u64(i) + sz];
33
      return c;
```

6.6. Linear Recurrences

6.6.1. Berlekamp-Massey Algorithm

```
template <typename T>
vector<T> berlekamp_massey(const vector<T> &s) {
    int n = s.size(), l = 0, m = 1;
    vector<T> r(n), p(n);
    r[0] = p[0] = 1;
    T b = 1, d = 0;
    for (int i = 0; i < n; i++, m++, d = 0) {
        for (int j = 0; j <= l; j++) d += r[j] * s[i - j];
        if ((d /= b) == 0) continue; // change if T is float auto t = r;
        for (int j = m; j < n; j++) r[j] -= d * p[j - m];
        if (l * 2 <= i) l = i + 1 - l, b *= d, m = 0, p = t;
}
return r.resize(l + 1), reverse(r.begin(), r.end()), r;
}</pre>
```

6.6.2. Linear Recurrence Calculation

```
1 template <typename T> struct lin_rec {
        using poly = vector<T>;
        poly mul(poly a, poly b, poly m) {
           int n = m.size();
           poly r(n);
          poty r(n);
for (int i = n - 1; i >= 0; i--) {
    r.insert(r.begin(), 0), r.pop_back();
    T c = r[n - 1] + a[n - 1] * b[i];
    // c /= m[n - 1]; if m is not monic
    for (int j = 0; j < n; j++)
        r[j] += a[j] * b[i] - c * m[j];
}</pre>
11
13
          return r;
       poly pow(poly p, ll k, poly m) {
  poly r(m.size());
15
           r[0] = 1;
           for (; k; k >>= 1, p = mul(p, p, m))
              if (k \& 1) r = mul(r, p, m);
19
21
        T calc(poly t, poly r, ll k) {
           int n = r.size();
           poly p(n);
           p[1] = 1:
           poly q = pow(p, k, r);
27
           T ans = 0:
           for (int i = 0; i < n; i++) ans += t[i] * q[i];
29
           return ans:
31 };
```

6.7. Matrices

6.7.1. Determinant

Requires: Mod Struct

```
1 Mod det(vector<vector<Mod>> a) {
      int n = a.size();
      Mod ans = 1;
      for (int i = 0; i < n; i++) {
         int b = i;
        for (int j = i + 1; j < n; j++)
if (a[j][i] != 0) {
              b = j;
              break:
         if (i != b) swap(a[i], a[b]), ans = -ans;
11
         ans *= a[i][i];
         if (ans == 0) return 0;
for (int j = i + 1; j < n; j++) {
   Mod v = a[j][i] / a[i][i];</pre>
15
           if (v != 0)
              for (int k = i + 1; k < n; k++)
17
                a[j][k] -= v * a[i][k];
19
21
      return ans;
```

```
double det(vector<vector<double>> a) {
   int n = a.size();

double ans = 1;
   for (int i = 0; i < n; i++) {
   int b = i;
   for (int j = i + 1; j < n; j++)
   if (fabs(a[j][i]) > fabs(a[b][i])) b = j;
```

```
if (i != b) swap(a[i], a[b]), ans = -ans;
         ans *= a[i][i];
         if (ans == 0) return 0;
        for (int j = i + 1; j < n; j++)
  double v = a[j][i] / a[i][i];</pre>
11
           if (v != 0)
13
             for (int k = i + 1; k < n; k++)
                a[j][k] -= v * a[i][k];
15
      }
17
      return ans;
19 }
```

6.7.2. Inverse

```
// Returns rank.
    // Result is stored in A unless singular (rank < n).</pre>
   // For prime powers, repeatedly set
    // A^{-1} = A^{-1} (2I - A*A^{-1}) \pmod{p^k}
   // where A^{-1} starts as the inverse of A mod p,
    // and k is doubled in each step.
   int matInv(vector<vector<double>> &A) {
 9
     int n = sz(A);
      vi col(n);
11
      vector<vector<double>> tmp(n, vector<double>(n));
      rep(i, 0, n) tmp[i][i] = 1, col[i] = i;
13
      rep(i, 0, n) {
        int r = i, c = i;
15
        rep(j, i, n)
        rep(k, i, n) if (fabs(A[j][k]) > fabs(A[r][c])) r = j,
17
19
        if (fabs(A[r][c]) < 1e-12) return i;
        A[i].swap(A[r])
21
        tmp[i].swap(tmp[r]);
        rep(j, 0, n) swap(A[j][i], A[j][c]),
        swap(tmp[j][i], tmp[j][c]);
        swap(col[i], col[c]);
25
        double v = A[i][i];
        rep(j, i + 1, n) {
           double f = A[j][i] / v;
          A[j][i] = 0;
29
          rep(k, i + 1, n) A[j][k] -= f * A[i][k];
          rep(k, 0, n) tmp[j][k] -= f * tmp[i][k];
31
        rep(j, i + 1, n) A[i][j] /= v;
rep(j, θ, n) tmp[i][j] /= v;
A[i][i] = 1;
33
35
37
      for (int i = n - 1; i > 0; --i) rep(j, 0, i) {
    double v = A[j][i];
39
          rep(k, 0, n) tmp[j][k] -= v * tmp[i][k];
41
43
      rep(i, 0, n) rep(j, 0, n) A[col[i]][col[j]] = tmp[i][j];
45
   }
   int matInv_mod(vector<vector<ll>>> &A) {
      int n = sz(A);
      vi col(n);
49
      vector<vector<ll>> tmp(n, vector<ll>(n));
      rep(i, \theta, n) tmp[i][i] = 1, col[i] = i;
      rep(i, 0, n) {
   int r = i, c = i;
53
        rep(j, i, n) rep(k, i, n) if (A[j][k]) {
55
          r = j;
c = k;
57
          goto found;
59
        return i;
      found:
61
        A[i].swap(A[r]);
        tmp[i].swap(tmp[r]);
63
        rep(j, 0, n) swap(A[j][i], A[j][c]),
swap(tmp[j][i], tmp[j][c]);
        swap(col[i], col[c]);
        ll v = modpow(A[i][i], mod - 2);
        rep(j, i + 1, n) {
    ll f = A[j][i] * v % mod;
          A[j][i] = 0;
          rep(k, i + 1, n) A[j][k] =
(A[j][k] - f * A[i][k]) % mod;
          rep(k, 0, n) tmp[j][k] = (tmp[j][k] - f * tmp[i][k]) % mod;
73
75
```

```
rep(j, i + 1, n) A[i][j] = A[i][j] * v % mod;
        rep(j, 0, n) tmp[i][j] = tmp[i][j] * v % mod;
A[i][i] = 1;
77
79
           (int i = n - 1; i > 0; --i) rep(j, 0, i) {
ll v = A[j][i];
rep(b ^ )
81
      for (int i = n -
           rep(k, 0, n) tmp[j][k] = (tmp[j][k] - v * tmp[i][k]) % mod;
83
85
      rep(i, 0, n) rep(j, 0, n) A[col[i]][col[j]] =
87
      tmp[i][j] \% mod + (tmp[i][j] < 0 ? mod : 0);
89
      return n;
```

6.7.3. Characteristic Polynomial

```
1 // calculate det(a - xI)
     template <typename T>
    vector<T> CharacteristicPolynomial(vector<vector<T>> a) {
        int N = a.size();
        for (int j = 0; j < N - 2; j++) {
  for (int i = j + 1; i < N; i++) {
    if (a[i][j] != 0) {
                 swap(a[j + 1], a[i]);
for (int k = 0; k < N; k++)
11
                    swap(a[k][j + 1], a[k][i]);
                 break;
             }
13
          if (a[j + 1][j] != 0) {
  T inv = T(1) / a[j + 1][j];
  for (int i = j + 2; i < N; i++) {</pre>
15
17
                 if (a[i][j] == 0) continue;
19
                 T coe = inv * a[i][j];
                 for (int l = j; l < N; l++)
a[i][l] -= coe * a[j + 1][l];
21
                 for (int k = 0; k < N; k++)
                    a[k][j + 1] += coe * a[k][i];
23
25
          }
27
        vector<vector<T>>> p(N + 1);
        p[0] = {T(1)};
for (int i = 1; i <= N; i++) {
29
          p[i].resize(i + 1);

for (int j = 0; j < i; j++) {

   p[i][j + 1] -= p[i - 1][j];

   p[i][j] += p[i - 1][j] * a[i - 1][i - 1];
31
33
35
           T x = 1;
           for (int m = 1; m < i; m++) {
    x *= -a[i - m][i - m - 1];
37
39
              T coe = x * a[i - m - 1][i - 1];
              for (int j = 0; j < i - m; j++)
p[i][j] += coe * p[i - m - 1][j];
41
          }
43
        return p[N];
45 }
```

6.7.4. Solve Linear Equation

```
1 typedef vector<double> vd;
   const double eps = 1e-12;
    // solves for x: A * x = b
   int solveLinear(vector<vd> &A, vd &b, vd &x) {
      int n = sz(A), m = sz(x), rank = 0, br, bc;
     if (n) assert(sz(A[0]) == m);
      vi col(m);
9
      iota(all(col), 0);
      rep(i, 0, n) {
11
        double v, bv = 0;
13
        rep(r, i, n) rep(c, i, m) if ((v = fabs(A[r][c])) > bv)
15
        bc = c, bv = v;
        if (bv <= eps) {
17
          rep(j, i, n) if (fabs(b[j]) > eps) return -1;
        swap(A[i], A[br]);
        swap(b[i], b[br]);
       swap(col[i], col[bc]);
rep(j, 0, n) swap(A[j][i], A[j][bc]);
bv = 1 / A[i][i];
23
```

```
rep(j, i + 1, n) {
  double fac = A[j][i] * bv;
         b[j] -= fac * b[i];
                                                                       55
         rep(k, i + 1, m) A[j][k] -= fac * A[i][k];
                                                                       57
       rank++;
31
                                                                       59
     33
                                                                       61
35
                                                                       63
       x[col[i]] = b[i];
37
       rep(j, \theta, i) b[j] -= A[j][i] * b[i];
                                                                       65
     return rank; // (multiple solutions if rank < m)</pre>
39
                                                                       67
   }
                                                                       69
```

6.8. Polynomial Interpolation

```
// returns a, such that a[0]x^0 + a[1]x^1 + a[2]x^2 + ...
   // passes through the given points
   typedef vector<double> vd;
   vd interpolate(vd x, vd y, int n) {
     vd res(n), temp(n);
     rep(k, 0, n - 1) rep(i, k + 1, n) y[i] = (y[i] - y[k]) / (x[i] - x[k]);
     double last = 0;
      temp[0] = 1;
     rep(k, 0, n) rep(i, 0, n) {
  res[i] += y[k] * temp[i];
        swap(last, temp[i]);
        temp[i] -= last * x[k];
13
      return res;
```

6.9. Simplex Algorithm

```
// Two-phase simplex algorithm for solving linear programs
   // of the form
   //
   //
           maximize
                         c^T x
 5
   //
           subject to
                         Ax <= b
                         x >= 0
   // INPUT: A -- an m x n matrix
             b -- an m-dimensional vector
 9
              c -- an n-dimensional vector
              x -- a vector where the optimal solution will be
   // OUTPUT: value of the optimal solution (infinity if
15
   // unbounded
   11
               above, nan if infeasible)
17
   // To use this code, create an LPSolver object with A, b,
19
   // and c as arguments. Then, call Solve(x).
   typedef long double ld;
   typedef vector<ld> vd;
   typedef vector<vd> vvd;
23
   typedef vector<int> vi;
   const ld EPS = 1e-9;
27
   struct LPSolver {
29
     int m, n;
     vi B. N:
     vvd D:
31
33
     LPSolver(const vvd &A, const vd &b, const vd &c)
          : m(b.size()), n(c.size()), N(n + 1), B(m),
35
            D(m + 2, vd(n + 2)) {
        for (int i = 0; i < m; i++)
37
          for (int j = 0; j < n; j++) D[i][j] = A[i][j];
        for (int i = 0; i < m; i++) {
          B[i] = n + i;
39
          D[i][n] = -1;
41
          D[i][n + 1] = b[i];
        for (int j = 0; j < n; j++) {
  N[j] = j;
  D[m][j] = -c[j];</pre>
45
        N[n] = -1;
       D[m + 1][n] = 1;
49
     void Pivot(int r, int s) {
  double inv = 1.0 / D[r][s];
```

```
for (int i = 0; i < m + 2; i++)
            if (i != r)
               for (int j = 0; j < n + 2; j++)
                 if (j != s) D[i][j] -= D[r][j] * D[i][s] * inv;
          for (int j = 0; j < n + 2; j++)
            if (j != s) D[r][j] *= inv;
          for (int i = 0; i < m + 2; i++)
            if (i != r) D[i][s] *= -inv;
          D[r][s] = inv;
          swap(B[r], N[s]);
       bool Simplex(int phase) {
          int x = phase == 1 ? m + 1 : m;
          while (true) {
            int s = -1;
            for (int j = 0; j <= n; j++) {
  if (phase == 2 && N[j] == -1) continue;
  if (s == -1 || D[x][j] < D[x][s] ||</pre>
                   D[x][j] == D[x][s] && N[j] < N[s])
            if (D[x][s] > -EPS) return true;
            int r = -1;
for (int i = 0; i < m; i++) {
   if (D[i][s] < EPS) continue;</pre>
               if (r == -1 \mid | D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s] \mid |
                    (D[i][n + 1] / D[i][s]) ==
(D[r][n + 1] / D[r][s]) &&
                    B[i] < B[r]
            if (r == -1) return false;
            Pivot(r, s);
       }
       ld Solve(vd &x) {
          for (int i = 1; i < m; i++)
  if (D[i][n + 1] < D[r][n + 1]) r = i;
          if (D[r][n + 1] < -EPS) {
            Pivot(r, n);
            if (!Simplex(1) \mid | D[m + 1][n + 1] < -EPS)
            return -numeric_limits<ld>::infinity();
for (int i = 0; i < m; i++)
               if (B[i] == -1) {
                 int s = -1;
                 for (int j = 0; j <= n; j++)
if (s == -1 || D[i][j] < D[i][s] ||
                        D[i][j] == D[i][s] && N[j] < N[s])
                 Pivot(i, s);
          if (!Simplex(2)) return numeric_limits<ld>::infinity();
          x = vd(n);
          for (int i = 0; i < m; i++)
            if (B[i] < n) \times [B[i]] = D[i][n + 1];
          return D[m][n + 1];
       }
115 };
    int main() {
       const int m = 4;
const int n = 3;
       vvd A(m);
       vd b(_b, _b + m);
       vd\ c(\bar{c}, \bar{c}, \bar{c} + n);
for (int i = 0; i < m; i++) A[i] = vd(\bar{A}[i], \bar{A}[i] + n);
       LPSolver solver(A, b, c);
       vd x;
       ld value = solver.Solve(x);
       cerr << "VALUE: " << value << endl; // VALUE: 1.29032</pre>
       cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];</pre>
       cerr << endl:
       return 0:
     }
```

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123

125

127

129

131

133

135

137

139

7. Geometry

7.1. Point

```
template <typename T> struct P {
       T \times y;
P(T \times Y = 0, T \times Y = 0) : x(x), y(y) {}
       bool operator<(const P &p) const {
          return tie(x, y) < tie(p.x, p.y);
       bool operator==(const P &p) const {
         return tie(x, y) == tie(p.x, p.y);
       P operator-() const { return {-x, -y}; }
       P operator+(P p) const { return {x + p.x, y + p.y}; }
P operator-(P p) const { return {x - p.x, y - p.y}; }
P operator*(T d) const { return {x * d, y * d}; }
P operator/(T d) const { return {x / d, y / d}; }
11
13
       T dist2() const { return x * x + y * y;
       double len() const { return sqrt(dist2()); }
       P unit() const { return *this / len(); }
friend T dot(P a, P b) { return a.x * b.x + a.y * b.y; }
       friend T cross(P a, P b) { return a.x * b.y - a.y * b.x; }
       friend T cross(P a, P b, P o) {
          return cross(a - o, b - o);
21
23 }:
    using pt = P<ll>;
```

7.1.1. Quarternion

```
constexpr double PI = 3.141592653589793;
    constexpr double EPS = 1e-7;
   struct Q {
      using T = double;
      T x, y, z, r;

Q(T r = 0) : x(0), y(0), z(0), r(r) {}

Q(T x, T y, T z, T r = 0) : x(x), y(y), z(z), r(r) {}

friend bool operator==(const Q &a, const Q &b) {
        return (a - b).abs2() <= EPS;
      friend bool operator!=(const Q &a, const Q &b) {
11
        return !(a == b);
13
      Q operator-() { return Q(-x, -y, -z, -r); } Q operator+(const Q &b) const {
15
        return Q(x + b.x, y + b.y, z + b.z, r + b.r);
17
      Q operator-(const Q &b) const {
        return Q(x - b.x, y - b.y, z - b.z, r - b.r);
19
21
      Q operator*(const T &t) const {
        return Q(x * t, y * t, z * t, r * t);
23
      Q operator*(const Q &b) const {
25
        return Q(r * b.x + x * b.r + y * b.z - z * b.y,
                   r * b.y - x * b.z + y * b.r + z * b.x,
27
                   r * b.z + x * b.y - y * b.x + z * b.r
                   r * b.r - x * b.x - y * b.y - z * b.z);
29
      Q operator/(const Q &b) const { return *this * b.inv(); }
      T abs2() const { return r * r + x * x + y * y + z * z; }
      T len() const { return sqrt(abs2()); }
      Q conj() const { return Q(-x, -y, -z, r); }
Q unit() const { return *this * (1.0 / len()); }
Q inv() const { return conj() * (1.0 / abs2()); }
35
      friend T dot(Q a, Q b) {
        return a.x * b.x + a.y * b.y + a.z * b.z;
37
39
      friend Q cross(Q a, Q b) {
        return Q(a.y * b.z - a.z * b.y, a.z * b.x - a.x * b.z,
a.x * b.y - a.y * b.x);
41
43
      friend Q rotation_around(Q axis, T angle) {
        return axis.unit() * sin(angle / 2) + cos(angle / 2);
45
      Q rotated_around(Q axis, T angle) {
        Q u = rotation_around(axis, angle);
47
        return u * *thīs / u;
49
      friend Q rotation_between(Q a, Q b) {
51
        a = a.unit(), b = b.unit();
        if (a == -b) {
           // degenerate case
53
           Q ortho = abs(a.y) > EPS ? cross(a, Q(1, 0, 0))
55
                                          cross(a, Q(0, 1, 0));
           return rotation_around(ortho, PI);
57
        return (a * (a + b)).conj();
      }
59
   };
```

7.1.2. Spherical Coordinates

```
1 struct car_p {
     double x, y, z;
 3
   }:
   struct sph_p {
5
     double r, theta, phi;
   sph_p conv(car_p p) {
     double r = sqrt(p.x * p.x + p.y * p.y + p.z * p.z);
double theta = asin(p.y / r);
9
11
     double phi = atan2(p.y, p.x);
     return {r, theta, phi};
13 }
   car_p conv(sph_p p) {
     double x = p.r * cos(p.theta) * sin(p.phi);
     double y = p.r * cos(p.theta) * cos(p.phi);
     double z = p.r * sin(p.theta);
     return {x, y, z};
19 }
```

7.2. Segments

```
// for non-collinear ABCD, if segments AB and CD intersect
bool intersects(pt a, pt b, pt c, pt d) {
   if (cross(b, c, a) * cross(b, d, a) > 0) return false;
   if (cross(d, a, c) * cross(d, b, c) > 0) return false;
   return true;
}

// the intersection point of lines AB and CD
pt intersect(pt a, pt b, pt c, pt d) {
   auto x = cross(b, c, a), y = cross(b, d, a);
   if (x == y) {
      // if (abs(x, y) < 1e-8) {
      // is parallel
   } else {
      return d * (x / (x - y)) - c * (y / (x - y));
   }
}</pre>
```

7.3. Convex Hull

7.3.1. 3D Hull

```
1 typedef Point3D<double> P3;
   struct PR {
      void ins(int x) { (a == -1 ? a : b) = x; }
void rem(int x) { (a == x ? a : b) = -1; }
      int cnt() { return (a != -1) + (b != -1); }
      int a, b;
   };
    struct F {
     P3 q;
11
      int a, b, c;
13 }:
   vector<F> hull3d(const vector<P3> &A) {
      assert(sz(A) >= 4);
      vector<vector<PR>>> E(sz(A), vector<PR>(sz(A), {-1, -1}));
17
    #define E(x, y) E[f.x][f.y]
vector<F> FS;
19
      auto mf = [8](int i, int j, int k, int l) {
  P3 q = (A[j] - A[i]).cross((A[k] - A[i]));
21
         if (q.dot(A[l]) > q.dot(A[i])) q = q * -1;
        F f{q, i, j, k};
E(a, b).ins(k);
         E(a, c).ins(j);
25
         E(b, c).ins(i)
         FS.push_back(f);
      rep(i, 0, 4) rep(j, i + 1, 4) rep(k, j + 1, 4)
```

```
mf(i, j, k, 6 - i - j - k);
31
      rep(i, 4, sz(A)) {
        rep(j, 0, sz(FS)) {
    F f = FS[j];
33
          if (f.q.dot(A[i]) > f.q.dot(A[f.a])) {
35
            E(a, b).rem(f.c);
            E(a, c).rem(f.b);
37
            E(b, c).rem(f.a);
swap(FS[j--], FS.back());
39
            FS.pop_back();
          }
41
        ŀ
43
        int nw = sz(FS);
        rep(j, 0, nw) {
          F f = FS[j];
   #define C(a, b, c
     if (E(a, b).cnt() != 2) mf(f.a, f.b, i, f.c);
          C(a, b, c);
          C(a, c, b);
          C(b, c, a);
51
      for (F &it : FS)
53
        if ((A[it.b] - A[it.a])
            .cross(A[it.c] - A[it.a])
55
             .dot(it.q) <= 0)
          swap(it.c, it.b);
     return FS;
59 };
```

7.4. Angular Sort

7.5. Convex Polygon Minkowski Sum

```
// O(n) convex polygon minkowski sum
   // must be sorted and counterclockwise
   vector<pt> minkowski_sum(vector<pt> p, vector<pt> q) {
     auto diff = [](vector<pt> δc)
       auto rcmp = [](pt a, pt b) {
         return pt{a.y, a.x} < pt{b.y, b.x};
       rotate(c.begin(), min_element(ALL(c), rcmp), c.end());
       c.push_back(c[0]);
       vector<pt> ret;
       for (int i = 1; i < c.size(); i++)
11
         ret.push_back(c[i] - c[i - 1]);
13
       return ret;
     };
     auto dp = diff(p), dq = diff(q);
15
     pt cur = p[\theta] + q[\theta];
vector<pt> d(dp.size() + dq.size()), ret = {cur};
     // include angle_cmp from angular-sort.cpp
     merge(ALL(dp), ALL(dq), d.begin(), angle_cmp);
19
     // optional: make ret strictly convex (UB if degenerate)
     int now = 0;
     for (int i = 1; i < d.size(); i++) {
       if (cross(d[i], d[now]) == 0)
23
         d[now] = d[now] + d[i];
       else
         d[++now] = d[i];
     d.resize(now + 1);
     // end optional part
     for (pt v : d) ret.push_back(cur = cur + v);
31
     return ret.pop_back(), ret;
```

7.6. Point In Polygon

```
bool on_segment(pt a, pt b, pt p) {
    return cross(a, b, p) == 0 && dot((p - a), (p - b)) <= 0;
}

// p can be any polygon, but this is O(n)

bool inside(const vector<pt> &p, pt a) {
    int cnt = 0, n = p.size();
    for (int i = 0; i < n; i++) {
        pt l = p[i], r = p[(i + 1) % n];
    // change to return 0; for strict version</pre>
```

```
if (on_segment(l, r, a)) return 1;
cnt ^= ((a.y < l.y) - (a.y < r.y)) * cross(l, r, a) > 0;
}
return cnt;
}
```

7.6.1. Convex Version

```
1 // no preprocessing version
    // p must be a strict convex hull, counterclockwise
    // if point is inside or on border
    bool is_inside(const vector<pt> &c, pt p) {
      int n = c.size(), l = 1, r = n - 1;
if (cross(c[0], c[1], p) < 0) return false;
if (cross(c[n - 1], c[0], p) < 0) return false;
while (l < r - 1) {</pre>
         int \dot{m} = (l + r) / 2
         T a = cross(c[\theta], c[m], p);
11
         if (a > 0)
         l = m;
else if (a < 0)
13
           r = m;
15
         else
           return dot(c[\theta] - p, c[m] - p) \ll \theta;
17
       if (l == r)
19
         return dot(c[\theta] - p, c[l] - p) \ll \theta;
21
         return cross(c[l], c[r], p) >= 0;
    // with preprocessing version
   vector<pt> vecs;
    pt center;
    // p must be a strict convex hull, counterclockwise
    // BEWARE OF OVERFLOWS!!
   void preprocess(vector<pt> p) {
      for (auto &v : p) v = v * 3;

center = p[0] + p[1] + p[2];

center.x /= 3, center.y /= 3;

for (auto &v : p) v = v - center;
31
33
       vecs = (angular_sort(p), p);
35
    bool intersect_strict(pt a, pt b, pt c, pt d) {
  if (cross(b, c, a) * cross(b, d, a) > 0) return false;
37
       if (cross(d, a, c) * cross(d, b, c) >= 0) return false;
39
       return true;
    // if point is inside or on border
    bool query(pt p) {
43
      p = p * 3 - center;
       auto pr = upper_bound(ALL(vecs), p, angle_cmp);
       if (pr == vecs.end()) pr = vecs.begin();
       auto pl = (pr == vecs.begin()) ? vecs.back() : *(pr - 1);
       return !intersect_strict({0, 0}, p, pl, *pr);
```

7.6.2. Offline Multiple Points Version

Requires: Point, GNU PBDS

```
using Double = __float128;
using Point = pt<Double, Double>;
   int n, m;
   vector<Point> poly;
   vector<Point> query;
   vector<int> ans;
   struct Segment {
     Point a, b;
11
     int id;
13 vector<Segment> segs;
15
   Double Xnow;
   inline Double get_y(const Segment &u, Double xnow = Xnow) {
17
     const Point &a = u.a;
     const Point &b = u.b;
     return (a.y * (b.x - xnow) + b.y * (xnow - a.x)) /
19
             (b.x - a.x);
   bool operator<(Segment u, Segment v) {</pre>
     Double yu = get_y(u);
     Double yv = get_y(v);
if (yu != yv) return yu < yv;
     return u.id < v.id;
   ordered_map<Segment> st;
   struct Event {
```

```
int type; // +1 insert seg, -1 remove seg, 0 query
      Double x, y;
    bool operator<(Event a, Event b) {</pre>
35
      if (a.x != b.x) return a.x < b.x;</pre>
      if (a.type != b.type) return a.type < b.type;</pre>
      return a.y < b.y;</pre>
    }
39
    vector<Event> events:
 41
    void solve() {
      set<Double> xs;
 43
      set<Point> ps;
for (int i = 0; i < n; i++) {</pre>
 45
        xs.insert(poly[i].x);
 47
        ps.insert(poly[i]);
 49
      for (int i = 0; i < n; i++) {
         Segment s\{poly[i], poly[(i + 1) % n], i\};
 51
         if (s.a.x > s.b.x ||
             (s.a.x == s.b.x && s.a.y > s.b.y)) {
           swap(s.a, s.b);
 55
         segs.push_back(s);
         if (s.a.x != s.b.x) {
           events.push_back({+1, s.a.x + 0.2, s.a.y, i});
           events.push_back({-1, s.b.x - 0.2, s.b.y, i});
 59
61
       for (int i = 0; i < m; i++) {
        events.push_back({0, query[i].x, query[i].y, i});
63
       sort(events.begin(), events.end());
 65
       int cnt = 0;
       for (Event e : events) {
 67
         int i = e.id;
         Xnow = e.x;
 69
         if (e.type == 0) {
           Double x = e.x;
           Double y = e.y;
Segment tmp = {{x - 1, y}, {x + 1, y}, -1};
           auto it = st.lower_bound(tmp);
           if (ps.count(query[i]) > 0) {
            ans[i] = 0:
           } else if (xs.count(x) > 0) {
  ans[i] = -2;
           } else if (it != st.end() &&
                       get_y(*it) == get_y(tmp)) {
 81
             ans[i] = \bar{0};
           } else if (it != st.begin() &&
 83
                       get_y(*prev(it)) == get_y(tmp)) {
             ans[i] = 0;
 85
           } else {
             int rk = st.order_of_key(tmp);
             if (rk % 2 == 1) {
 89
               ans[i] = 1;
             } else {
               ans[i] = -1;
             }
         } else if (e.type == 1) {
           st.insert(segs[i]);
           assert((int)st.size() == ++cnt);
         } else if (e.type == -1) {
           st.erase(segs[i]);
           assert((int)st.size() == --cnt);
99
101
      }
```

7.7. Closest Pair

```
vector<pll> p; // sort by x first!
bool cmpy(const pll &a, const pll &b) const {
    return a.y < b.y;
}

ll sq(ll x) { return x * x; }
// returns (minimum dist)^2 in [l, r)

ll solve(int l, int r) {
    if (r - l <= 1) return 1e18;
    int m = (l + r) / 2;
    ll mid = p[m].x, d = min(solve(l, m), solve(m, r));
    auto pb = p.begin();
    inplace_merge(pb + l, pb + m, pb + r, cmpy);
    vector<pll> s;
    for (int i = l; i < r; i++)
        if (sq(p[i].x - mid) < d) s.push_back(p[i]);</pre>
```

7.8. Minimum Enclosing Circle

```
1 typedef Point<double> P;
     double ccRadius(const P &A, const P &B, const P &C) {
  return (B - A).dist() * (C - B).dist() * (A - C).dist() /
      abs((B - A).cross(C - A)) / 2;
 5
    P ccCenter(const P &A, const P &B, const P &C) {
   P b = C - A, c = B - A;
   return A + (b * c.dist2() - c * b.dist2()).perp() /
 9
                        b.cross(c) / 2;
11
    pair<P, double> mec(vector<P> ps) {
        shuffle(all(ps), mt19937(time(0)));
13
        P o = ps[0];
        double r = 0, EPS = 1 + 1e-8;
        rep(i, \theta, sz(ps)) if ((o - ps[i]).dist() > r * EPS) {
          o = ps[i], r = 0;
rep(j, 0, i) if ((o - ps[j]).dist() > r * EPS) {
17
             o = (ps[i] + ps[j]) / 2;
             r = (o - ps[i]).dist();
rep(k, 0, j) if ((o - ps[k]).dist() > r * EPS) {
o = ccCenter(ps[i], ps[j], ps[k]);
19
21
                r = (o - ps[i]).dist();
23
25
       }
        return {o, r};
27 }
```

7.9. Delaunay Triangulation

```
1 typedef Point<ll> P;
   typedef struct Quad *Q;
   typedef __int128_t lll; // (can be ll if coords are < 2e4)</pre>
   P arb(LLONG_MAX, LLONG_MAX); // not equal to any other point
   struct Quad {
     bool mark;
     Q o, rot;
     P p;
P F() { return r()->p; }
     Q r() { return rot->rot; }
     Q prev() { return rot->o->rot; }
     Q next() { return r()->prev(); }
13
15
   17
19
             p.cross(c, a) * B >
21
23 Q makeEdge(P orig, P dest) {
     Q q[] = \{new Quad\{0, 0, 0, orig\}, new Quad\{0, 0, 0, arb\}\}
     new Quad\{0, 0, 0, \text{ dest}\}, new Quad\{0, 0, 0, \text{ arb}\}; rep(i, 0, 4) q[i]->o = q[-i \& 3],
                   q[i] -> rot = q[(i + 1) & 3];
27
     return *q;
29 }
   void splice(Q a, Q b) {
     swap(a->o->rot->o, b->o->rot->o);
     swap(a->o, b->o);
33
   Q connect(Q a, Q b) {
     Q q = makeEdge(a->F(), b->p);
35
     splice(q, a->next());
     splice(q->r(), b);
37
     return q:
39 }
41
   pair<Q, Q> rec(const vector<P> &s) {
     if (sz(s) <= 3) {
43
       Q a = makeEdge(s[0], s[1])
       b = makeEdge(s[1], s.back());
if (sz(s) == 2) return {a, a->r()};
        splice(a->r(), b);
       auto side = s[0].cross(s[1], s[2]);
       Q c = side ? connect(b, a) : 0;
       return {side < 0 ? c \rightarrow r() : a, side < 0 ? c : b \rightarrow r()};
49
```

```
e->F(), e->p
    #define valid(e) (e->F().cross(H(base)) > 0)
      Q A, B, ra, rb;
int half = sz(s) /
      tie(ra, A) = rec({all(s) - half});
      tie(B, rb) = rec({sz(s) - half + all(s)});
while ((B->p.cross(H(A)) < 0 && (A = A->next())) ||
              (A->p.cross(H(B)) > 0 & (B = B->r()->o)))
 59
61
      Q base = connect(B->r(), A);
      if (A->p == ra->p) ra = base->r();
      if (B->p == rb->p) rb = base;
63
 65
    #define DEL(e, init, dir)
      Q e = init->dir;
      if (valid(e))
 67
         while (circ(e->dir->F(), H(base), e->F())) {
          Q t = e->dir;
 69
           splice(e, e->prev());
           splice(e->r(), e->r()->prev());
 73
      base = connect(RC, base->r());
         else
          base = connect(base->r(), LC->r());
 81
 83
      return {ra, rb};
    }
85
    // returns [A_0, B_0, C_0, A_1, B_1, \dots] // where A_i, B_i, C_i are counter-clockwise triangles
87
    vector<P> triangulate(vector<P> pts) {
      sort(all(pts));
 89
      assert(unique(all(pts)) == pts.end());
      if (sz(pts) < 2) return {};
      Q e = rec(pts).first;
      vector<Q> q = {e};
      int qi = 0;
      while (e->o->F().cross(e->F(), e->p) < 0) e = e->o;
 95
    #define ADD
97
        Q c = e;
        do {
99
          c->mark = 1:
          pts.push_back(c->p);
101
          q.push_back(c->r());
103
          c = c->next();
        } while (c != e);
105
      ADD;
107
      pts.clear();
      while (qi < sz(q))
        if (!(e = q[qi++])->mark) ADD;
109
      return pts;
```

7.9.1. Slower Version

```
template <class P, class F>
   void delaunay(vector<P> &ps, F trifun) {
     if (sz(ps) == 3) {
        int d = (ps[\theta].cross(ps[1], ps[2]) < \theta);
       trifun(0, 1 + d, 2 - d);
     vector<P3> p3;
for (P p : ps) p3.emplace_back(p.x, p.y, p.dist2());
     if (sz(ps) > 3)
        for (auto t : hull3d(p3))
11
          if ((p3[t.b] - p3[t.a])
              .cross(p3[t.c] - p3[t.a])
               .dot(P3(0, 0, 1)) < 0)
13
            trifun(t.a, t.c, t.b);
15 }
```

7.10. Half Plane Intersection

```
21
struct Line {
  Point P;
                                                                   25
  bool operator<(const Line &b) const {</pre>
    return atan2(v.y, v.x) < atan2(b.v.y, b.v.x);
                                                                   27
};
bool OnLeft(const Line &L, const Point &p) {
  return Cross(L.v, p - L.P) > 0;
```

```
Point GetIntersection(Line a, Line b) {
      Vector u = a.P - b.P;
      Double t = Cross(b.v, u) / Cross(a.v, b.v);
      return a.P + a.v * t;
15
   int HalfplaneIntersection(Line *L, int n, Point *poly) {
17
     sort(L, L + n);
     int first, last;
Point *p = new Point[n];
19
      Line *q = new Line[n];
21
      q[first = last = 0] = L[0];
     for (int i = 1; i < n; i++) {
  while (first < last && !OnLeft(L[i], p[last - 1]))</pre>
23
25
          last-
        while (first < last \delta\delta !OnLeft(L[i], p[first])) first++;
        q[++last] = L[i];
27
        if (fabs(Cross(q[last].v, q[last - 1].v)) < EPS) {</pre>
29
          last-
          if (OnLeft(q[last], L[i].P)) q[last] = L[i];
31
        if (first < last)
          p[last - 1] = GetIntersection(q[last - 1], q[last]);
33
35
      while (first < last δδ !OnLeft(q[first], p[last - 1]))
      if (last - first <= 1) return 0;
37
      p[last] = GetIntersection(q[last], q[first]);
39
      for (int i = first; i <= last; i++) poly[m++] = p[i];</pre>
41
      return m:
43 }
```

8. Strings

8.1. Knuth-Morris-Pratt Algorithm

```
vector<int> pi(const string &s) {
         vector<int> p(s.size());
for (int i = 1; i < s.size(); i++) {</pre>
            int g = p[i - 1];
           while (g \, \&\& \, s[i] \, != \, s[g]) \, g = \, p[g - 1];

p[i] = g + (s[i] == \, s[g]);
        return p:
 9
     }
     vector<int> match(const string &s, const string &pat) {
  vector<int> p = pi(pat + '\0' + s), res;
  for (int i = p.size() - s.size(); i < p.size(); i++)</pre>
11
            if (p[i] == pat.size())
13
               res.push_back(i - 2 * pat.size());
15
         return res;
```

8.2. Aho-Corasick Automaton

```
struct Aho_Corasick {
     static const int maxc = 26, maxn = 4e5;
     struct NODES {
       int Next[maxc], fail, ans;
     NODES T[maxn];
     int top, qtop, q[maxn];
     int get node(const int &fail) {
       fill_n(T[top].Next, maxc, 0);
       T[top].fail = fail;
       T[top].ans = 0;
11
       return top++;
13
     int insert(const string &s) {
15
       int ptr = 1;
       for (char c : s) { // change char id
17
         if (!T[ptr].Next[c]) T[ptr].Next[c] = get_node(ptr);
19
         ptr = T[ptr].Next[c];
       return ptr;
     } // return ans_last_place
     void build_fail(int ptr) {
23
       int tmp;
       for (int i = 0; i < maxc; i++)
         if (T[ptr].Next[i]) {
           tmp = T[ptr].fail;
while (tmp != 1 && !T[tmp].Next[i])
              tmp = T[tmp].fail;
            if (T[tmp].Next[i] != T[ptr].Next[i])
              if (T[tmp].Next[i]) tmp = T[tmp].Next[i];
```

```
T[T[ptr].Next[i]].fail = tmp;
33
            q[qtop++] = T[ptr].Next[i];
35
     void AC_auto(const string &s) {
       int ptr = 1;
for (char c : s) {
          while (ptr != 1 && !T[ptr].Next[c]) ptr = T[ptr].fail;
39
          if (T[ptr].Next[c]) {
            ptr = T[ptr].Next[c];
41
            T[ptr].ans++;
         }
43
       }
     }
45
     void Solve(string &s) {
47
        for (char &c : s) // change char id
        for (int i = 0; i < qtop; i++) build_fail(q[i]);
49
        AC_auto(s);
        for (int i = qtop - 1; i > -1; i--)
51
          T[T[q[i]].fail].ans += T[q[i]].ans;
     void reset() {
       qtop = top = q[0] = 1;
       get_node(1);
   } AC;
59
   // usage example
   string s, S;
int n, t, ans_place[50000];
61
   int main() {
     Tie cin >> t:
     while (t--) {
65
       AC.reset();
        cin >> S >> n;
        for (int i = 0; i < n; i++) {
          cin >> s
          ans_place[i] = AC.insert(s);
69
       AC.Solve(S);
        for (int i = 0; i < n; i++)
73
          cout << AC.T[ans_place[i]].ans << '\n';</pre>
75 }
```

8.3. Suffix Array

```
// sa[i]: starting index of suffix at rank i
    // 0-indexed, sa[0] = n (empty string)
// lcp[i]: lcp of sa[i] and sa[i - 1], lcp[0] = 0
    struct SuffixArray {
       vector<int> sa, lcp;
       rank(n);
          sa = lcp = y, iota(all(sa), 0);
for (int j = 0, p = 0; p < n;
    j = max(1, j * 2), lim = p) {</pre>
11
13
             p = j, iota(all(y), n - j);
for (int i = 0; i < n; i++)
if (sa[i] >= j) y[p++] = sa[i] - j;
17
              fill(all(ws), 0);
             for (int i = 0; i < n; i++) ws[x[i]]++;

for (int i = 1; i < lim; i++) ws[i] += ws[i - 1];

for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];

swap(x, y), p = 1, x[sa[0]] = 0;

for (int i = 1; i < n; i++)
19
21
                a = sa[i - 1], b = sa[i]
23
                25
27
29
           for (int i = 1; i < n; i++) rank[sa[i]] = i;
          for (int i = 0, j; i < n - 1; lcp[rank[i++]] = k)
  for (k &&-, j = sa[rank[i] - 1];</pre>
31
                    s[i + k] == s[j + k]; k++)
33
35 };
```

8.4. Suffix Tree

```
int max_len, in, times;
      } * root, *last, reg[maxn * 2];
      int top;
 9
      Node *get_node(int _max) {
         Node *re = &reg[top++];
         re->in = 0, re->times = 1;
         re->max_len = _max, re->green = 0;
         for (int i = 0; i < maxc; i++) re->edge[i] = 0;
13
15
      void insert(const char c) { // c in range [0, maxc)
17
         Node *p = last;
         last = get_node(p->max_len + 1);
        while (p && !p->edge[c])
p->edge[c] = last, p = p->green;
19
         if (!p)
21
           last->green = root;
23
         else {
           Node *pot_green = p->edge[c];
25
           if ((pot_green->max_len) == (p->max_len + 1))
              last->green = pot_green;
27
              Node *wish = get_node(p->max_len + 1);
29
              wish->times = 0;
              while (p && p->edge[c] == pot_green)
              p->edge[c] = wish, p = p->green;
for (int i = 0; i < maxc; i++)
  wish->edge[i] = pot_green->edge[i];
33
             wish->green = pot_green->green;
pot_green->green = wish;
35
              last->green = wish;
37
        }
39
      Node *q[maxn * 2];
      int ql, qr;
41
      void get_times(Node *p) {
    ql = 0, qr = -1, reg[0].in = 1;
    for (int i = 1; i < top; i++) reg[i].green->in++;
    for (int i = 0; i < top; i++)</pre>
43
45
           if (!reg[i].in) q[++qr] = &reg[i];
47
         while (ql \ll qr) {
           q[ql]->green->times += q[ql]->times;
49
           if (!(--q[ql]->green->in)) q[++qr] = q[ql]->green;
51
      }
53
      void build(const string δs) {
         top = 0;
         root = last = get_node(0);
for (char c : s) insert(c - 'a'); // change char id
55
57
        get_times(root);
59
       // call build before solve
      int solve(const string &s) {
61
         Node *p = root;
         for (char c : s)
63
           if (!(p = p -> edge[c - 'a'])) // change char id
             return 0:
65
        return p->times;
      }
67 };
```

8.5. Z Value

```
int z[n];
void zval(string s) {
    // z[i] => longest common prefix of s and s[i:], i > 0
    int n = s.size();
    z[0] = 0;
    for (int b = 0, i = 1; i < n; i++) {
        if (z[b] + b <= i)
            z[i] = 0;
    else
        z[i] = min(z[i - b], z[b] + b - i);
    white (s[i + z[i]] == s[z[i]]) z[i]++;
    if (i + z[i] > b + z[b]) b = i;
}
```

8.6. Manacher's Algorithm

```
int z[n];
void manacher(string s) {
    // z[i] => longest odd palindrome centered at i is
    // s[i] => longest odd palindrome centered at i is
    // s[i] - z[i] ... i + z[i]]

// to get all palindromes (including even length),
    // insert a '#' between each s[i] and s[i + 1]

int n = s.size();
    z[0] = 0;
```

```
for (int b = 0, i = 1; i < n; i++) {
    if (z[b] + b >= i)
        z[i] = min(z[2 * b - i], b + z[b] - i);
    else
    z[i] = 0;
    while (i + z[i] + 1 < n && i - z[i] - 1 >= 0 &&
        s[i + z[i] + 1] == s[i - z[i] - 1])
    z[i]++;
    if (z[i] + i > z[b] + b) b = i;
}
```

8.7. Minimum Rotation

```
int min_rotation(string s) {
    int a = 0, n = s.size();
    s += s;
    for (int b = 0; b < n; b++) {
        for (int k = 0; k < n; k++) {
            if (a + k == b || s[a + k] < s[b + k]) {
                b += max(0, k - 1);
                break;
        }
        if (s[a + k] > s[b + k]) {
            a = b;
                break;
        }
    }
}
return a;
}
```

8.8. Palindromic Tree

```
struct palindromic_tree {
      struct node {
         int next[26], fail, len;
        int cnt,
num; // cnt: appear times, num: number of pal. suf.
node(int l = 0) : fail(0), len(l), cnt(0), num(0) {
           for (int i = 0; i < 26; ++i) next[i] = 0;
        }
 9
      vector<node> St;
11
      vector<char> s;
      int last, n;
13
      palindromic_tree() : St(2), last(1), n(0) {
        St[0].fail = 1, St[1].len = -1, s.pb(-1);
15
      inline void clear() {
        St.clear(), s.clear(), last = 1, n = 0;
17
         St.pb(0), St.pb(-1);
        St[0].fail = 1, s.pb(-1);
19
      inline int get_fail(int x) {
  while (s[n - St[x].len - 1] != s[n]) x = St[x].fail;
21
23
        return x;
      inline void add(int c) {
  s.push_back(c -= 'a'), ++n;
  int cur = get_fail(last);
25
27
         if (!St[cur].next[c]) {
           int now = SZ(St);
           St.pb(St[cur].len + 2);
31
           St[now].fail = St[get_fail(St[cur].fail)].next[c];
           St[cur].next[c] = now;
33
           St[now].num = St[St[now].fail].num + 1;
        last = St[cur].next[c], ++St[last].cnt;
35
      inline void count() { // counting cnt
         auto i = St.rbegin();
         for (; i != St.rend(); ++i) {
39
           St[i->fail].cnt += i->cnt;
41
      inline int size() { // The number of diff. pal.
  return SZ(St) - 2;
43
45
    };
```

9. Debug List

```
    Pre-submit:

            Did you make a typo when copying a template?

    Test more cases if unsure.

            Write a naive solution and check small cases.

    Submit the correct file.
```

```
General Debugging:
        Read the whole problem again.
      - Have a teammate read the problem.
        Have a teammate read your code.
        - Explain you solution to them (or a rubber duck).
        Print the code and its output / debug output.
13
      - Go to the toilet.
     Wrong Answer:
15
       · Any possible overflows?
         > __int128` ?
Try `-f+~~
17
                -ftrapv` or `#pragma GCC optimize("trapv")`
      - Floating point errors?
- > `long double` ?
19
        - turn off math optimizations
21
      - check for `==`, `>=`, `acos(1.000000001)`, etc.
- Did you forget to sort or unique?
23
      - Generate large and worst "corner" cases.
- Check your `m` / `n`, `i` / `j` and `x` / `y`.
25
      - Are everything initialized or reset properly?
      - Are you sure about the STL thing you are using?
27
         Read cppreference (should be available).
29
      - Print everything and run it on pen and paper.
31
   - Time Limit Exceeded:
      - Calculate your time complexity again.
33
      - Does the program actually end?
         Check for `while(q.size())` etc.
35
        Test the largest cases locally.
      - Did you do unnecessary stuff?
        - e.g. pass vectors by value
37
        - e.g. `memset` for every test case
39
      - Is your constant factor reasonable?
41
     Runtime Error:
       Check memory usage.
43
        - Forget to clear or destroy stuff?
        - > `vector::shrink_to_fit()
        Stack overflow?
        Bad pointer / array access?
- Try `-fsanitize=address`
        Division by zero? NaN's?
```