Midterm Exam

Instructor: Taewhi Lee April 22, 2013

1. [10pts] Describe two fundamental differences between the relational algebra and SQL.

Relational algebra: set-based (no duplicate tuple), procedural language

SQL: multiset-based (may have duplicate tuples), non-procedural language

2. [5pts] Explain why the following query does not work.

select dept name, id, count(name)

from instructor

group by dept_name;

id cannot be used in the select clause

because id is not included in the group by expression.

3. [10pts] List five integrity constraints and briefly explain each of them.

not null: a specified attribute cannot be null

unique: values for specified attributes must be unique or null

primary key: values for specified attributes must be unique and not null

check: tuples must meet a given predicate

foreign key: a value that appears in one relation for a given set of attributes also appears for a certain set of attributes in another relation

4. a) [5pts] What is the difference between a relation and a view?

A view definition causes the saving of an expression;

the expression is substituted into queries using the view.

b) [5pts] Describe three advantages of using a view that joins two relations, against creating another relation that joins two relations.

-Certain data can be hidden from the view of certain users.

No additional space overhead

Data modifications(insert/delete/update) are needed only for the underlying relations.

5. a) [5pts] Explain two major pitfalls to avoid in designing a database schema.

Redundancy: repeating information may cause data inconsistency

Incompleteness: it is difficult or impossible to model certain aspects of the enterprise

b) [5pts] Describe pros and cons of database normalization.

Pros: we can reduce data redundancy without loss of information.

Cons: more join operations between multiple tables are needed for data retrieval.

c) [5pts] Why do we sometimes need to denormalize a database schema instead of avoiding the pitfalls?

We may want to use non-normalized schema for performance, for example, to reduce join operations between multiple relations.

6. [5pts: 1pt each / 0pt for no answer / -1pt for a wrong answer]
Consider the relation r below. Which of the following FDs does r satisfy?
Choose TRUE or FALSE for each FD.

a)	$A\toD$	FALSE
,	$AB \rightarrow D$	TRUE
,	$C \rightarrow BDE$	TRUE
,		
	$E \rightarrow A$	FALSE
e)	$A \to E$	TRUE

Α	В	С	D	Е
a ₁	b_1	C ₁	d_1	e_1
a ₁	b_2	c_2	d_2	e ₁
a_2	b_1	C ₃	d ₃	e_1
a_2	b_1	C ₄	d_3	e ₁

 C_5

 d_1

 e_1

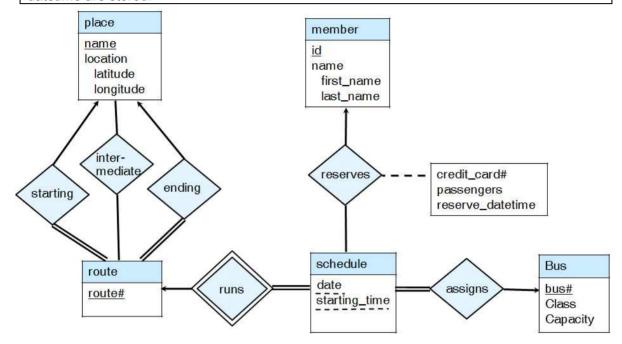
Relation r

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7. a) [10pts] Construct an E-R diagram for the following description.

Design a database for the reservation office of a bus company.

- Each bus has a unique number. We also store its class and capacity.
- Each place has a unique name and location information of the latitude and longitude.
- Routes have a starting place and an ending place; also, some of them have several intermediate places.
- A number of buses are scheduled to a route. A bus is assigned to one schedule; some buses can have multiple schedule. We store the date and starting time of each schedule.
- A member of our company can book a bus by specifying a schedule. We store a unique id, first name, and last name of each member.
- For each reservation, credit card number, the number of passengers, and the reservation datetime are stored.



b) [10pts] Convert your E-R diagram into a relational schema.

place(<u>name</u>, latitude, longitude) route(<u>route#</u>)

starting place(route#, name)

ending_place(route#, name)

intermediate place(route#, name)

schedule(route#, day, starting time)

bus(bus#, class,capacity)

assignment(route#, day, starting time, bus#)

member(id, first name, last name)

reservation(id, route#, day, starting time, credit card#, passengers, reserve datetime)

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8. Let F = \{AB \rightarrow C, B \rightarrow D, CD \rightarrow E, CE \rightarrow GH\}.
Give a derivation sequence on F for the following FDs using only Armstrong's axioms.
a) [5pts] AB \rightarrow E
 AB \rightarrow ABD (augmentation: B \rightarrow D with AB)
 ABD \rightarrow CD (augmentation: AB \rightarrow C with D)
 AB \rightarrow CD (transitivity: AB \rightarrow ABD and ABD \rightarrow CD)
 AB \rightarrow E (transitivity: AB \rightarrow CD and CD \rightarrow E)
b) [5pts] AB \rightarrow G
 AB \rightarrow ABE (augmentation: AB \rightarrow E with AB)
 ABE \rightarrow CE (augmentation: AB \rightarrow C with E)
 AB \rightarrow CE (transitivity: AB \rightarrow ABE and ABE \rightarrow CE)
 AB \rightarrow GH (transitivity: AB \rightarrow CE and CE \rightarrow GH)
 GH \rightarrow G (reflexivity)
 AB \rightarrow G (transitivity: AB \rightarrow GH and GH \rightarrow G)
9. Let F be the set of functional dependencies on relation schema R (A, B, C, D, E, F).
    F = \{A \rightarrow B, C \rightarrow D, AC \rightarrow E, D \rightarrow F\}
a) [5pts] Find all candidate keys for the relation schema R.
 Since AC \rightarrow ABCDEF, AC is a superkey.
 A and C must be in any candidate key since they do not appear on the right of any FD.
 Hence, AC is the only candidate key.
b) [5pts] Give a lossless-join decomposition of R into BCNF.
     R1(A, B), R2(C, D), R3(A, C, E), R4(C, F)
  or R1(A, B), R2(C, D), R3(A, C, E), R4(D, F)
c) [5pts] Give a decomposition of R into 3NF, having a lossless-join and preserving dependencies.
Is your decomposition is in BCNF?
 R1(A, B), R2(C, D), R3(A, C, E), R4(D, F)
 The answer is in BCNF.
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