

# Optimization Methods (CS1.404), Spring 2024

## Lecture 25

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April 22nd, 2024



# The Gradient Projection Method

- The stationarity condition  $\mathbf{x}^* = P_C(\mathbf{x}^* - s\nabla f(\mathbf{x}^*))$  motivates to solve the optimization problem

$$\begin{aligned} (P) \quad & \min f(\mathbf{x}) \\ & s.t. \mathbf{x} \in C \end{aligned}$$

## Gradient Projection Method

**Input:**  $\epsilon > 0$  (tolerance parameter)

**Initialization:** Pick  $\mathbf{x}_0 \in C$  arbitrarily

**General Steps:** For  $k = 0, 1, 2, \dots$  execute the following steps:

- 1 Pick a stepsize  $t_k$  by a line search procedure.
- 2 Set  $\mathbf{x}_{k+1} = P_C(\mathbf{x}_k - t_k \nabla f(\mathbf{x}_k))$ .
- 3 If  $\|\mathbf{x}_k - \mathbf{x}_{k+1}\| \leq \epsilon$ , then STOP and  $\mathbf{x}_{k+1}$  is the output.

# Sufficient Decrease Lemma for Projected Gradient Descent

## Lemma

Suppose that  $f \in \mathbb{C}^{1,1}(C)$ , where  $C$  is a closed convex set. Then for any  $\mathbf{x} \in C$  and  $t \in (0, 2/L)$ , the following inequality will hold.

$$f(\mathbf{x}) - f(P_C(\mathbf{x} - t\nabla f(\mathbf{x}))) \geq t \left(1 - \frac{Lt}{2}\right) \left\| \frac{1}{t}(\mathbf{x} - P_C(\mathbf{x} - t\nabla f(\mathbf{x}))) \right\|^2.$$

- When  $C = \mathbb{R}^n$ , the obtained inequality is exactly the same as the one obtained for unconstrained case.

- We define gradient mapping as

$$G_M(\mathbf{x}) = M \left[ \mathbf{x} - P_C \left( \mathbf{x} - \frac{1}{M} \nabla f(\mathbf{x}) \right) \right]$$

where  $M > 0$ .

- When  $C = \mathbb{R}^n$ , we see that  $G_M(\mathbf{x}) = \nabla f(\mathbf{x})$ . So, the gradient mapping is an extension of usual gradient operation.
- We see that  $G_M(\mathbf{x}) = \mathbf{0}$  if and only  $\mathbf{x} = P_C \left( \mathbf{x} - \frac{1}{M} \nabla f(\mathbf{x}) \right)$ . This happens if and only if  $\mathbf{x}$  is a stationary point of optimization problem  $(P)$ .
- Thus, we can use  $\|G_M(\mathbf{x})\|$  as an optimality measure.
- Thus, the sufficient decrease property of projected gradient method essentially states that

$$f(\mathbf{x}) - f(P_C(\mathbf{x} - t \nabla f(\mathbf{x}))) \geq t \left( 1 - \frac{Lt}{2} \right) \|G_{\frac{1}{t}}(\mathbf{x})\|^2.$$

# Sufficient decrease of consecutive function values

## Lemma

Consider the optimization problem

$$(P) : \min f(\mathbf{x}) \\ \text{s.t. } \mathbf{x} \in C$$

where  $C$  is a closed convex set and  $f \in \mathbb{C}_L^{1,1}$ . Let  $\mathbf{x}_k$ ,  $k = 0, 1, 2, \dots$  be the sequence generated by the project gradient algorithm for solving  $(P)$  using either constant stepsize  $t_k = \tilde{t} \in (0, \frac{2}{L})$  or by a stepsize chosen using backtracking procedure with parameters  $(s, \alpha, \beta)$  satisfying  $s > 0, \alpha \in (0, 1), \beta \in (0, 1)$ . Then for  $k \geq 0$

$$f(\mathbf{x}_k) - f(\mathbf{x}_{k+1}) \geq M \|G_d(\mathbf{x}_k)\|^2$$

where

$$M = \begin{cases} \tilde{t}(1 - \frac{L\tilde{t}}{2}), & \text{for constant step size} \\ \min\left(s, \frac{2\beta(1-\alpha)}{L}\right), & \text{for Backtracking} \end{cases}$$

and

$$d = \begin{cases} \frac{1}{\tilde{t}}, & \text{for constant step size} \\ \frac{1}{s}, & \text{for Backtracking.} \end{cases}$$

# Convergence of the Gradient Projection Method

## Theorem

Consider the optimization problem

$$(P) : \min f(\mathbf{x}) \\ s.t. \mathbf{x} \in C$$

where  $C$  is a closed convex set and  $f \in \mathbb{C}_L^{1,1}$  is bounded below. Let  $\mathbf{x}_k$ ,  $k = 0, 1, 2, \dots$  be the sequence generated by the project gradient algorithm for solving  $(P)$  using either constant stepsize  $t_k = \tilde{t} \in (0, \frac{2}{L})$  or by a stepsize chosen using backtracking procedure with parameters  $(s, \alpha, \beta)$  satisfying  $s > 0, \alpha \in (0, 1), \beta \in (0, 1)$ . Then we have the following:

- The sequence  $\{f(\mathbf{x}_k)\}_{k=0}^{\infty}$  is nonincreasing. In addition,  $f(\mathbf{x}_{k+1}) < f(\mathbf{x}_k)$  unless  $\mathbf{x}_k$  is a stationary point of  $(P)$ .
- $G_d(\mathbf{x}_k) \rightarrow \mathbf{0}$  as  $k \rightarrow \infty$ . where

$$d = \begin{cases} \frac{1}{\tilde{t}}, & \text{for constant step size} \\ \frac{1}{s}, & \text{for Backtracking.} \end{cases}$$

# Equality Constrained Minimization problem

- Consider following convex optimization problem with equality constraints,

$$\min f(\mathbf{x}) \quad \text{s.t.} \quad A\mathbf{x} = \mathbf{b},$$

where  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  is convex and twice continuously differentiable, and  $A \in \mathbb{R}^{p \times n}$  with  $\text{rank}(A) = p < n$ .

- The assumptions on  $A$  mean that there are fewer equality constraints than variables, and that the equality constraints are independent.
- We will assume that an optimal solution  $\mathbf{x}^*$  exists and let  $p^* = \inf\{f(\mathbf{x}) \mid A\mathbf{x} = \mathbf{b}\} = f(\mathbf{x}^*)$ .
- A point  $\mathbf{x}^* \in \text{dom}(f)$  is optimal for if and only if there is a  $\boldsymbol{\mu}^* \in \mathbb{R}^p$  such that

$$A\mathbf{x}^* = \mathbf{b}$$

$$\nabla f(\mathbf{x}^*) + A^\top \boldsymbol{\mu}^* = \mathbf{0}$$

- Above are KKT optimality conditions. There are  $n + p$  equations in the  $n + p$  variables  $\mathbf{x}^*, \boldsymbol{\mu}^*$ .

# Equality Constrained Convex Quadratic Minimization

- Consider the equality constrained convex quadratic minimization problem

$$\begin{aligned} \min f(\mathbf{x}) &= (1/2)\mathbf{x}^\top P\mathbf{x} + \mathbf{q}^\top \mathbf{x} + r \\ \text{s.t. } A\mathbf{x} &= \mathbf{b}, \end{aligned}$$

where  $P \in \mathbb{R}^{n \times n}$  is symmetric positive semi-definite matrix,  $A \in \mathbb{R}^{p \times n}$ ,  $\mathbf{q} \in \mathbb{R}^n$  and  $\mathbf{b} \in \mathbb{R}^p$ .

- $(\mathbf{x}^*, \boldsymbol{\mu}^*)$  is a KKT point of this problem if it satisfies the following conditions.

$$\begin{aligned} A\mathbf{x}^* &= \mathbf{b} \\ P\mathbf{x}^* + \mathbf{q} + A^\top \boldsymbol{\mu}^* &= \mathbf{0} \end{aligned}$$

- We can rewrite the above KKT conditions in the matrix form as follows.

$$\begin{bmatrix} P & A^\top \\ A & \mathbf{0}_{p \times p} \end{bmatrix} \begin{bmatrix} \mathbf{x}^* \\ \boldsymbol{\mu}^* \end{bmatrix} = \begin{bmatrix} -\mathbf{q} \\ \mathbf{b} \end{bmatrix}$$

- There are  $n + p$  linear equations in the  $n + p$  variables  $\mathbf{x}^*, \boldsymbol{\mu}^*$ . The coefficient matrix is called the KKT matrix.



# Equality Constrained Convex Quadratic Minimization

- When the KKT matrix is nonsingular, there is a unique optimal primal-dual pair  $(\mathbf{x}^*, \boldsymbol{\mu}^*)$ .
- If the KKT matrix is singular, but the KKT system is solvable, any solution yields an optimal pair  $(\mathbf{x}^*, \boldsymbol{\mu}^*)$ .
- If the KKT system is not solvable, the quadratic optimization problem is unbounded below or infeasible.
  - In this case, there exist  $\mathbf{v} \in \mathbb{R}^n$  and  $\mathbf{w} \in \mathbb{R}^p$  such that  $P\mathbf{v} + A^\top \mathbf{w} = \mathbf{0}$ ,  $A\mathbf{v} = \mathbf{0}$ ,  $-\mathbf{q}^\top \mathbf{v} + \mathbf{b}^\top \mathbf{w} > 0$ .
  - Let  $\hat{\mathbf{x}}$  be any feasible point. The point  $\mathbf{x} = \hat{\mathbf{x}} + t\mathbf{v}$  is feasible for all  $t$  and

$$\begin{aligned} f(\hat{\mathbf{x}} + t\mathbf{v}) &= f(\hat{\mathbf{x}}) + t(\mathbf{v}^\top P\hat{\mathbf{x}} + \mathbf{q}^\top \mathbf{v}) + (1/2)t^2 \mathbf{v}^\top P\mathbf{v} \\ &= f(\hat{\mathbf{x}}) + t(-\hat{\mathbf{x}}^\top A^\top \mathbf{w} + \mathbf{q}^\top \mathbf{v}) - (1/2)t^2 \mathbf{w}^\top A\mathbf{v} \\ &= f(\hat{\mathbf{x}}) + t(-\mathbf{b}^\top \mathbf{w} + \mathbf{q}^\top \mathbf{v}) \end{aligned}$$

which decreases without bound as  $t \rightarrow \infty$ .

# Non-Singularity of KKT Matrix

## Claim

KKT matrix is nonsingular if  $P$  is positive definite on the null space of  $A$ . Equivalently,  $A\mathbf{x} = \mathbf{0}, \mathbf{x} \neq \mathbf{0} \Rightarrow \mathbf{x}^\top P\mathbf{x} > 0$ .

## Proof:

- Assume that  $A\mathbf{x} = \mathbf{0}, \mathbf{x} \neq \mathbf{0} \Rightarrow \mathbf{x}^\top P\mathbf{x} > 0$ .
- Let  $(\mathbf{x}, \boldsymbol{\mu})$  be a vector in the null space of the KKT matrix.
- Thus,  $\begin{bmatrix} P & A^\top \\ A & \mathbf{0}_{p \times p} \end{bmatrix} \begin{bmatrix} \mathbf{x} \\ \boldsymbol{\mu} \end{bmatrix} = \mathbf{0}$ .
- This means,  $P\mathbf{x} + A^\top \boldsymbol{\mu} = \mathbf{0}$  and  $A\mathbf{x} = \mathbf{0}$ .
- Thus,  $\mathbf{x}^\top (P\mathbf{x} + A^\top \boldsymbol{\mu}) = 0$ , which implies  $\mathbf{x}^\top P\mathbf{x} + (A\mathbf{x})^\top \boldsymbol{\mu} = 0$ . But using  $A\mathbf{x} = \mathbf{0}$ , we get  $\mathbf{x}^\top P\mathbf{x} = 0$ . But, as per assumption  $P$  is positive definite on the null space of  $A$ , we get  $\mathbf{x} = \mathbf{0}$ .
- Now using the equation  $P\mathbf{x} + A^\top \boldsymbol{\mu} = \mathbf{0}$  and  $\mathbf{x} = \mathbf{0}$ , we get  $A^\top \boldsymbol{\mu} = \mathbf{0}$ . But, this gives  $\boldsymbol{\mu} = \mathbf{0}$  as  $A$  is full rank matrix.
- Thus, the only null space of the KKT matrix is vector  $(\mathbf{x}, \boldsymbol{\mu}) = \mathbf{0}$ . Thus, KKT matrix is nonsingular.

# Example: Equality Constrained Convex Quadratic Minimization

- Consider the following optimization problem

$$\begin{aligned} \min \quad & (x_1 - 2)^2 + 2(x_2 - 1)^2 - 5 \\ \text{s.t.} \quad & x_1 + 4x_2 = 3 \end{aligned}$$

- KKT Conditions give

$$\begin{aligned} 2x_1 + \mu &= 4 \\ 4x_2 + 4\mu &= 4 \\ x_1 + 4x_2 &= 3 \end{aligned}$$

which is equivalent to 
$$\begin{bmatrix} 2 & 0 & 1 \\ 0 & 4 & 4 \\ 1 & 4 & 0 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \mu \end{bmatrix} = \begin{bmatrix} 4 \\ 4 \\ 3 \end{bmatrix}.$$

- The solution of these KKT condition is  $x_1^* = \frac{5}{3}$ ,  $x_2^* = \frac{1}{3}$ ,  $\mu^* = \frac{2}{3}$ .

# Eliminating Equality Constraints

- One general approach to solving the equality constrained problem is to eliminate the equality constraints, and then solve the resulting unconstrained problem using methods for unconstrained minimization.
- We first find a matrix  $F \in \mathbb{R}^{n \times (n-p)}$  and vector  $\hat{\mathbf{x}} \in \mathbb{R}^n$  that parametrize the (affine) feasible set:

$$\{\mathbf{x} \mid A\mathbf{x} = \mathbf{b}\} = \hat{\mathbf{x}} + \{\mathbf{x} \mid A\mathbf{x} = \mathbf{0}\} = \{F\mathbf{z} + \hat{\mathbf{x}} \mid \mathbf{z} \in \mathbb{R}^{n-p}\}.$$

- Here,
  - $\hat{\mathbf{x}}$  is any feasible point. Thus,  $A\hat{\mathbf{x}} = \mathbf{b}$ .
  - $F \in \mathbb{R}^{n \times (n-p)}$  is composed of  $(n-p)$  basis vectors of null space of  $A$ .  
Thus,  $AF = \mathbf{0}_{p \times (n-p)}$ .
- Thus, the reduced unconstrained problem is

$$\min_{\mathbf{z}} f_{\text{new}}(\mathbf{z}) = \min_{\mathbf{z}} f(\hat{\mathbf{x}} + F\mathbf{z})$$

# Eliminating Equality Constraints

$\mathbf{x}^* = \hat{\mathbf{x}} + F\mathbf{z}^*$  fulfils the KKT conditions with  $\boldsymbol{\mu}^* = -(AA^\top)^{-1}A\nabla f(\mathbf{x}^*)$ .

## Proof:

- Clearly,  $\mathbf{x}^* = \hat{\mathbf{x}} + F\mathbf{z}^*$  is primal feasible as  $A\mathbf{x}^* = A\hat{\mathbf{x}} + AF\mathbf{z}^* = \mathbf{b} + \mathbf{0} = \mathbf{b}$ .
- Consider the matrix  $\begin{bmatrix} F^\top \\ A \end{bmatrix}$ . For such a matrix,

$$\begin{aligned} \begin{bmatrix} F^\top \\ A \end{bmatrix} (\nabla f(\mathbf{x}^*) + A^\top \boldsymbol{\mu}^*) &= \begin{bmatrix} F^\top \nabla f(\mathbf{x}^*) + F^\top A^\top \boldsymbol{\mu}^* \\ A \nabla f(\mathbf{x}^*) + AA^\top \boldsymbol{\mu}^* \end{bmatrix} \\ &= \begin{bmatrix} F^\top \nabla f(\mathbf{x}^*) - F^\top A^\top (AA^\top)^{-1} A \nabla f(\mathbf{x}^*) \\ A \nabla f(\mathbf{x}^*) - AA^\top (AA^\top)^{-1} A \nabla f(\mathbf{x}^*) \end{bmatrix} \\ &= \begin{bmatrix} \nabla f_{\text{new}}(\mathbf{z}^*) - (AF)^\top (AA^\top)^{-1} A \nabla f(\mathbf{x}^*) \\ A \nabla f(\mathbf{x}^*) - A \nabla f(\mathbf{x}^*) \end{bmatrix} = \begin{bmatrix} \mathbf{0} \\ \mathbf{0} \end{bmatrix} \end{aligned}$$

- Since, matrix  $\begin{bmatrix} F^\top \\ A \end{bmatrix}$  has full rank,  $\begin{bmatrix} F^\top \\ A \end{bmatrix} (\nabla f(\mathbf{x}^*) + A^\top \boldsymbol{\mu}^*) = \begin{bmatrix} \mathbf{0} \\ \mathbf{0} \end{bmatrix}$  implies  $\nabla f(\mathbf{x}^*) + A^\top \boldsymbol{\mu}^* = \mathbf{0}$ . Thus, KKT conditions are satisfied by  $\mathbf{x}^* = \hat{\mathbf{x}} + F\mathbf{z}^*$  and  $\boldsymbol{\mu}^* = -(AA^\top)^{-1}A\nabla f(\mathbf{x}^*)$ .

# Newton Method for Equality Constraint Problem

- Here we describe an extension of Newton's method to include linear equality constraint.
- The methods are almost the same except for two differences:
  - the initial point must be feasible  $A\mathbf{x} = \mathbf{b}$ ,
  - the Newton step must be a feasible direction  $A\Delta\mathbf{x}_{nt} = \mathbf{0}$
- $\Delta\mathbf{x}_{nt}$  solves the second order approximation of  $f$  at  $\mathbf{x}$  with variable  $\mathbf{v}$ .
- Let  $\hat{f}(\mathbf{x} + \mathbf{v}) = f(\mathbf{x}) + \nabla f(\mathbf{x})^\top \mathbf{v} + \frac{1}{2} \mathbf{v}^\top \nabla^2 f(\mathbf{x}) \mathbf{v}$ . Then,  $\Delta\mathbf{x}_{nt}$  is the minimizer of the following problem.

$$\begin{aligned} \min \quad & f(\mathbf{x}) + \nabla f(\mathbf{x})^\top \mathbf{v} + \frac{1}{2} \mathbf{v}^\top \nabla^2 f(\mathbf{x}) \mathbf{v} \\ \text{s.t.} \quad & A(\mathbf{x} + \mathbf{v}) = \mathbf{b} \end{aligned}$$

which is same as

$$\begin{aligned} \min \quad & f(\mathbf{x}) + \nabla f(\mathbf{x})^\top \mathbf{v} + \frac{1}{2} \mathbf{v}^\top \nabla^2 f(\mathbf{x}) \mathbf{v} \\ \text{s.t.} \quad & A\mathbf{v} = \mathbf{0} \end{aligned}$$

# Newton Method for Equality Constraint Problem

- $\Delta \mathbf{x}_{nt}$  satisfies KKT Optimality conditions for this optimization problem, which are:

$$\begin{aligned}\nabla f(\mathbf{x}) + \nabla^2 f(\mathbf{x}) \Delta \mathbf{x}_{nt} + A^\top \boldsymbol{\mu} &= \mathbf{0} \\ A \Delta \mathbf{x}_{nt} &= \mathbf{0}\end{aligned}$$

where  $\boldsymbol{\mu}$  is the associated dual variable for the quadratic problem.

- Thus, Newton step  $\Delta \mathbf{x}_{nt}$  at  $\mathbf{x}$  is the solution of following linear system

$$\begin{bmatrix} \nabla^2 f(\mathbf{x}) & A^\top \\ A & \mathbf{0}_{p \times p} \end{bmatrix} \begin{bmatrix} \Delta \mathbf{x}_{nt} \\ \boldsymbol{\mu} \end{bmatrix} = \begin{bmatrix} -\nabla f(\mathbf{x}) \\ \mathbf{0} \end{bmatrix}$$

# Newton Method for Equality Constraint Problem

- **Newton Decrement:** Newton decrement is defined as

$$\lambda(\mathbf{x}) = (\Delta \mathbf{x}_{nt}^\top \nabla f(\mathbf{x}) \Delta \mathbf{x}_{nt})^{\frac{1}{2}}$$

- Newton decrement is related to  $f(\mathbf{x}) - \min_{\mathbf{y}: A\mathbf{y}=\mathbf{b}} \hat{f}(\mathbf{y})$  as:

$$f(\mathbf{x}) - \min_{\mathbf{y}: A\mathbf{y}=\mathbf{b}} \hat{f}(\mathbf{y}) = \frac{1}{2} \lambda(\mathbf{x})^2$$

where  $\hat{f}$  is second order approximation of  $f$  at  $\mathbf{x}$ .

- This gives an estimate of  $f(\mathbf{x}) - f^*$  using quadratic approximation.
- The Newton decrement comes up in the line search as well, since the directional derivative of  $f$  in the direction  $\Delta \mathbf{x}_{nt}$  is

$$\frac{\partial}{\partial t} f(\mathbf{x} + t \Delta \mathbf{x}_{nt})|_{t=0} = \nabla f(\mathbf{x})^\top \Delta \mathbf{x}_{nt} = -\Delta \mathbf{x}_{nt}^\top \nabla^2 f(\mathbf{x}) \Delta \mathbf{x}_{nt} = -\lambda(\mathbf{x})^2.$$

Here, we used the Newton step property  $\nabla^2 f(\mathbf{x}) \Delta \mathbf{x}_{nt} = -\nabla f(\mathbf{x})$ .



# Newton Method for Equality Constraint Problem

- We say that  $\mathbf{v} \in \mathbb{R}^n$  is a feasible direction for  $A(\mathbf{x} + \mathbf{v}) = \mathbf{b}$  if  $A\mathbf{v} = \mathbf{0}$ .
- We say that  $\mathbf{v}$  is a descent direction for  $f$  at  $\mathbf{x}$ , if for small  $t > 0$ ,  $f(\mathbf{x} + t\mathbf{v}) < f(\mathbf{x})$ .
- The Newton step  $\Delta\mathbf{x}_{nt}$  is always a feasible descent direction (except when  $\mathbf{x}$  is optimal, in which case  $\Delta\mathbf{x}_{nt} = \mathbf{0}$ ).

# Newton Method for Equality Constraint Problem

## Newton Method for Equality Constraint Problem

**Input:** starting point  $\mathbf{x} \in \text{domain}(f)$  with  $A\mathbf{x} = \mathbf{b}$ , tolerance  $\epsilon > 0$ .  
**while**  $\frac{1}{2}\lambda(\mathbf{x})^2 > \epsilon$  **do**  
    Compute the Newton step  $\Delta\mathbf{x}_{nt}$  and decrement  $\lambda(\mathbf{x})$ .  
    Choose step size  $t$  by backtracking line search.  
    Update  $\mathbf{x} = \mathbf{x} + t\Delta\mathbf{x}_{nt}$ .  
**end while**

- The method is called a feasible descent method, since all the iterates are feasible, with  $f(\mathbf{x}(k+1)) < f(\mathbf{x}(k))$  (unless  $\mathbf{x}(k)$  is optimal).

# Precise Algorithm: Newton Method for Equality Constraint Problem

## Newton Method for Equality Constraint Problem

**Input:** starting point  $\mathbf{x}_0 \in \text{domain}(f)$  with  $A\mathbf{x}_0 = \mathbf{b}$ , tolerance  $\epsilon > 0$ , Maximum number of iterations  $K$ .

**for**  $k = 0, 1, \dots, K$  **do**

Find  $\Delta\mathbf{x}_k$  and  $\mu_k$  as

$$\begin{bmatrix} \Delta\mathbf{x}_k \\ \mu_k \end{bmatrix} = \begin{bmatrix} \nabla^2 f(\mathbf{x}_k) & A^\top \\ A & \mathbf{0}_{p \times p} \end{bmatrix}^{-1} \begin{bmatrix} -\nabla f(\mathbf{x}_k) \\ \mathbf{0} \end{bmatrix}$$

**if**  $\lambda(\mathbf{x}_k) \leq \epsilon$  **then**

return  $\mathbf{x}_k$

**end if**

Choose step size  $t$  by backtracking line search.

Update  $\mathbf{x}_{k+1} = \mathbf{x}_k + t\Delta\mathbf{x}_k$ .

**end for**

# Convergence of Feasible Newton Method for Equality Constraint Problem

- The iterates  $\mathbf{x}(k)$  are the same as those of the Newton algorithm for the eliminated unconstrained problem

$$\hat{f}(\mathbf{z}) := f(\mathbf{x}_0 + F\mathbf{z}), \quad \mathbf{x}(k) = \mathbf{x}_0 + F\mathbf{z}(k)$$

as they fulfil the KKT conditions of the quadratic approximation.

- Thus convergence is the same as in the unconstrained case.