Edit - Distance

Minimum steps to connent

Word 1 to word 2 Recursive >DP (i) Word = "abe" -> Delete all character Base Case: word 2 = " " » tetur len (nox) (ii) Word 1 = " "
word 2 = "abe" => insent all elements
of mand 2 retur len(mort 2) Operations Insert }
Delete }
minimum
replace Word 1 = abdWord 2 = abd 5 = abd* notequal word 2 [s] → Calculat with which operation steps will equal ward1[i] == word2[i] be minimum. a Don't need any operation. Word2 = abd, equal (n-1, m-1)word2 = acd mword = a (n-1, m-1)word 2 = a (n-1, m-1)word 2 = a (n-1, m-1)both became equal by replace so decrease both paintery. Word 1 = a (n-1, m)word 2 = a (n-1, m)As d'is already also alread an alread (n-1, m)As d'is already checked so check for a in words sol

Word 1 = a b d Insert

ward 2 = a b d \Rightarrow As for (c' of mord 2 me added 'e' in mord 1 so c is checked. 30 now we have to check for remaining ic. left side elements of words (n, m-1): (m-1)det solve (NI, W2, n, m): if n==0 or m==0: return m or nelif w1[n-1] == w2[m-1]: return salne (W1, W2, N-1, M-1) return 1+ min (solne (w1, w2, n-1, m-1) # replace solve(W1, W2, n-1, m) # Delete solve (WI, WZ, n, m-1) # Insert Tabulation s it W1 is empty and we is at longth ? i.e. warac minimum 2 steps needed if wie = "" and w1 = abd minimum 3 steps that is why this cell contain 3

Top-donun dp = [[0]*(m+i) for i inscange (n+i)]# Base case For i in scornge (m+1): 0 th naw 4P[0][j] = j W1="" for i instange (n+1): # Bose case oth calumn dp[i][0] = i W2="" bor i in range (1, n+1): for j in scange (1, m+1): ék W1[i-1] == W2[-1] op[i][j] = op [i-1][j-1] else: dP[i][j] = 1 + min(dp[i-1][j-1] # Replace 4 [i-1][j] # Delete of [i] [j-1] # Insent oceturn dp[-1][-1]