

Monday, Nov 03, 2025

- Huffman Properties.** Prove the following: if some character occurs with frequency more than  $2/5$ , then there is guaranteed to be a codeword of length 1.

- Huffman Encoding.** Let  $n \geq 2$  and label symbols  $s_1, \dots, s_n$  with frequencies

$$f_i = 2^{n-i} \quad (i = 1, \dots, n),$$

so the frequencies (in descending order) are  $2^{n-1}, 2^{n-2}, \dots, 2, 1$ . Construct the Huffman code for these frequencies and determine the codeword lengths  $L(s_i)$  for all  $i$ .

- Worst Case for Greedy Set Cover.** Let  $n$  be a power of 2. Show that there exists an instance of the set cover problem such that: 1) there are  $n$  elements in the base set; 2) the optimal solution uses only two sets; and 3) the greedy algorithm picks at least  $\log n$  sets.

#### 4. Weighted Set Cover and the Greedy Algorithm

Let us consider the weighted version of the Set Cover problem.

Suppose the universe is

$$U = \{1, 2, 3, 4, 5, 6, 7, 8\}.$$

The following sets and weights are given:

$$\begin{aligned} T_1 &= \{1, 2, 3, 4\}, & w(T_1) &= 9, \\ T_2 &= \{5, 6\}, & w(T_2) &= 7, \\ T_3 &= \{7\}, & w(T_3) &= 5, \\ A &= \{1, 3, 5, 7\}, & w(A) &= 10, \\ B &= \{2, 4, 6, 8\}, & w(B) &= 10. \end{aligned}$$

- What can be a possible greedy strategy for solving this weighted set cover problem?
- Apply the proposed greedy strategy step by step to find the sets selected by the greedy algorithm and the total cost of the cover it produces.
- Determine the optimal solution and its cost. Compare it with the greedy solution.