

Lecture Section:

Monday, Nov 10, 2025

Student Name:

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1. (2 pts.) What are the key components of formulating and solving a problem using Dynamic Programming? Select all that apply.

- (a) Defining a recurrence.
- (b) Solving the base case.
- (c) Defining a greedy heuristic.
- (d) Defining subproblems with optimal substructure.

2. (2 pts.) The space complexity of the dynamic programming solution for the finding the Longest Increasing Subsequence (LIS) is:

- (a) $O(n)$ - can be done in linear space
- (b) $O(n^2)$ - need space for all $O(n^2)$ subproblems
- (c) $O(1)$ - no space overhead needed, we only incur a cost in time
- (d) $O(n \log n)$

3. (2 pts.) Consider the two strings $x := ACTGGACTT$ and $y := AGTCGTTT$. What is the edit distance $d(x, y)$?

- (a) 3
- (b) 4
- (c) 5
- (d) 6

4. (2 pts.) In class, you learned that the minimum edit distance problem can be solved in polynomial

time using Dynamic Programming. Suppose there are two strings of equal length n . What is the time complexity of determining the **optimal alignment** of the two sequences (not just the minimum edit distance)?

- (a) $O(n^2)$ - same runtime as the edit distance problem
- (b) $O(n^3)$ - Need an additional factor of $O(n)$ to compare the best alignment for each subproblem.
- (c) $O(n^4)$ - Need an additional factor of $O(n^2)$ to compare the best alignment for each subproblem for both strings.
- (d) Cannot be computed in polynomial runtime.

5. (2 pts.) Consider the sequence $S = (4, 6, 2, 3, 8, 1, 5)$. Recall that we defined the following recurrence for finding the length of the Longest Increasing Subsequence ending at a_j :

$$L(j) = \begin{cases} 1 + \max_{i < j, a_i < a_j}(L(i)) \\ 1, \text{ if no such } i \end{cases}$$

What is $L(4)$ while trying to find the length of the LIS in S ? Assume that S is 1-indexed.

- (a) 2
- (b) 3
- (c) 4
- (d) 5