

Midterm1

● Graded

Student

Jesse Zhang

Total Points

28.5 / 74 pts

Question 1

Multiple choice

12 / 27 pts

1.1 MC1

3 / 3 pts

✓ - 0 pts Correct

- 3 pts Incorrect

1.2 MC2

3 / 3 pts

✓ - 0 pts Correct

- 3 pts Incorrect

1.3 MC3

3 / 3 pts

✓ - 0 pts Correct

- 3 pts Incorrect

1.4 MC4

3 / 3 pts

✓ - 0 pts Correct

- 3 pts Incorrect

1.5 MC5

0 / 3 pts

- 0 pts Correct

✓ - 3 pts Incorrect

1.6 MC6

0 / 3 pts

- 0 pts Correct

✓ - 3 pts Incorrect

1.7 MC7

0 / 3 pts

- 0 pts Correct

✓ - 3 pts Incorrect

1.8 MC8

0 / 3 pts

- 0 pts Correct

✓ - 3 pts Incorrect

1.9 MC9

0 / 3 pts

- 0 pts Correct

✓ - 3 pts Incorrect

Question 2

True/False

11 / 20 pts

2.1 TF1

0 / 2 pts

- 0 pts Correct

✓ - 2 pts Incorrect

2.2 TF2

2 / 2 pts

✓ - 0 pts Correct

- 2 pts Incorrect

2.3 TF3

2 / 2 pts

✓ - 0 pts Correct

- 2 pts Incorrect

2.4 TF4

2 / 2 pts

✓ - 0 pts Correct

- 2 pts Incorrect

2.5 TF5

0 / 2 pts

- 0 pts Correct

✓ - 2 pts Incorrect

2.6 TF6

0 / 2 pts

- 0 pts Correct

✓ - 2 pts Incorrect

2.7 TF7

2 / 2 pts

✓ - 0 pts Correct

- 2 pts Incorrect

2.8 TF8

1 / 2 pts

- 0 pts Correct

✓ - 1 pt Incorrect

2.9 TF9

0 / 2 pts

- 0 pts Correct

✓ - 2 pts Incorrect

2.10 TF10

2 / 2 pts

✓ - 0 pts Correct

- 2 pts Incorrect

Question 3

Recurrences

0 / 12 pts

3.1 (a)

0 / 6 pts

- 0 pts Part 1: Recurrence relation for running time [3 pts/3 pts]

Correctly wrote down the recurrence relation for running time as $T(n) = 64T(n/4) + \Theta(n^3)$.

- 1 pt Part 1: Recurrence relation for running time [2 pts/3 pts]

Could not correctly write down the entire recurrence relation, but the student was on the right track to do so by expressing the running time T as a function of input size n and showing some of the terms in recurrence correctly.

✓ - 3 pts Part 1: Recurrence relation for running time [0 pts/3 pts]

Did not attempt at all or did not express the recurrence relation for running time T as a function of input size n (no mentioning of $T(n)$ at all).

- 0 pts Part 2: Solution for recurrence relation in \mathcal{O} [3 pts/3 pts]

Correctly wrote down the solution of the recurrence relation as $\mathcal{O}(n^3 \log n)$ using Master Theorem.

- 1 pt Part 2: Solution for recurrence relation in \mathcal{O} [2 pts/3 pts]

Did not apply the Master Theorem and attempted to unfold the recurrence relation but could not reach the correct solution. Or found the right values for a,b, and d but could not reach the correct solution.

✓ - 3 pts Part 2: Solution for recurrence relation in \mathcal{O} [0 pts/3 pts]

Did not attempt at all or incorrectly **applied the Master Theorem** to solve the recurrence relation.

- 0 pts Part 1: Recurrence relation for running time [3 pts/3 pts]

Correctly wrote down the recurrence relation for running time as $T(n) = T(19) + T(29) + T(39) + T(n - 1) + 1.9^n$ and *optionally* simplified the expression to $T(n) = T(n - 1) + 1.9^n$.

- 1 pt Part 1: Recurrence relation for running time [2 pts/3 pts]

Could not correctly write down the entire recurrence relation, but the student was on the right track to do so by expressing the running time T as a function of input size n and showing some of the terms in recurrence correctly.

✓ - 3 pts Part 1: Recurrence relation for running time [0 pts/3 pts]

Did not attempt at all or did not express the recurrence relation for running time T as a function of input size n (no mentioning of $T(n)$ at all).

- 0 pts Part 2: Solution for recurrence relation in \mathcal{O} [3 pts/3 pts]

Correctly wrote down the solution of the recurrence relation as $\mathcal{O}(1.9^n)$ either by *mentioning HW1 Q5(e)* or unfolding the recurrence.

- 1 pt Part 2: Solution for recurrence relation in \mathcal{O} [2 pts/3 pts]

Did not apply the Master Theorem and attempted to unfold the recurrence relation but could not reach the correct solution.

✓ - 3 pts Part 2: Solution for recurrence relation in \mathcal{O} [0 pts/3 pts]

Did not attempt at all or incorrectly **applied the Master Theorem** to solve the recurrence relation.

Question 4

Algorithms

5.5 / 15 pts

4.1 (a)

- 0 pts Correct

✓ - 2.5 pts Only one of M_L or M_R is incorrect- 5 pts Both M_L and M_R are incorrect/not written

4.2 (b)

3 / 5 pts

- 0 pts Correct

- 1 pt Justification for relation between M_L and M_R to majority of A not provided/incorrect✓ - 2 pts Incorrect answer for how M_L and M_R are related to majority of A, or provided the relation specific to the example not in general- 3 pts Missing answer for how M_L and M_R are related to majority of A

- 2 pts Missing/Incorrect answer for majority of A

4.3 (c)

0 / 5 pts

- 0 pts Correct

✓ - 5 pts wrong algorithm and time complexity

- 4.75 pts only running time correct

- 0.5 pts wrong running time

- 3 pts idea of the algorithm

- 1 pt partially correct algorithm

- 4 pts attempt on algorithm

- 2 pts mostly correct

Midterm 1 - D

Name: Jesse Zhang

Penn State access ID (xyz1234) in the following box:

2fz5311

Student ID number (9XXXXXXXXX):

975557148

TA and/or section time:

MWF 4:35

Instructions:

- Answer all questions. Read them carefully first. Be precise and concise. Handwriting needs to be neat. Box numerical final answers.
- Please clearly write your name and your PSU access ID (i.e., xyz1234) in the box on top of every page.
- Write in only the space provided. You may use the back of the pages only as scratch paper. **Do not write your solutions in the back of pages!**
- **Do not write outside of the black bounding box:** the scanner will not be able to read anything outside of this box.
- We are providing two extra page at the end if you need extra space. Make sure you mention in the space provided that your answer continues there.

Good luck!

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Multiple choice questions (27 points)

For each of the following questions, select the right answer by filling the corresponding grading bubble.

1. If $T(n) = 3T(n/2) + O(n^2)$, which of the following is the tightest upper bound for $T(n)$?

Answer.

- $T(n) = O(n^{\log_3 2})$
- $T(n) = O(n^{\log 3})$
- $T(n) = O(n^2)$
- $T(n) = O(n^2 \log n)$
- None of the above

2. While performing InsertionSort to sort the array $[4, 1, 2, 3, 7, 8, 5, 6]$ in descending order, which of the following is NOT a transition state of the array?

Answer.

- $[4, 2, 1, 3, 7, 8, 5, 6]$
- $[8, 7, 4, 3, 2, 1, 5, 6]$
- $[4, 3, 2, 1, 7, 8, 5, 6]$
- $[7, 4, 3, 2, 1, 8, 5, 6]$
- $[8, 7, 5, 4, 3, 2, 1, 6]$
- None of the above

3. If $T(n) = 9T(n/3) + O(n^2 + n(\log n)^5)$, which of the following is the tightest upper bound for $T(n)$?

Answer.

- $T(n) = O(n^2)$
- $T(n) = O(n^2 \log n)$
- $T(n) = O(n(\log n)^5)$
- $T(n) = O(n^2(\log n)^5)$
- None of the above

4. Suppose that $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0$. Which of the following must be true?

Answer.

- $f(n) = O(g(n))$
- $f(n) = \Omega(g(n))$
- $f(n) = \Theta(g(n))$
- None of the above

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5. Consider the Partition function used in QuickSort. If we choose the last element as the pivot, what is the resulting array after calling Partition on [5, 8, 6, 7, 1, 3, 2, 9, 4]?

Answer.

- [1, 2, 3, 4, 5, 6, 7, 8, 9]
- [1, 3, 2, 4, 5, 8, 6, 7, 9]
- [1, 2, 3, 4, 9, 8, 7, 6, 5]
- [1, 3, 2, 4, 5, 8, 6, 9, 7]
- [1, 3, 2, 7, 5, 8, 6, 9, 4]
- None of the above

4 5 8 6 7
4 5 1 6 7 8
1 5 4 6 7 8 3
1 4 8 6 5 7 2
1 4 7 6 5 8 4 6
1 5 3 2 7 8 4 6
1 2 3 5 7 8 4 6 9

6. Run MergeSort to sort the array [5, 2, 9, 4, 6, 1, 10, 7, 3, 11] in ascending order, which two subarrays are the inputs to the final Merge call that produces the fully sorted array?

Answer.

1 5
1 2 5 4 9
1 2 3 4 5 6 7 9 11

- [5, 2, 9, 4, 6] and [1, 10, 7, 3, 11]
- [2, 5, 4, 8, 6] and [4, 10, 3, 7, 11]
- [2, 5, 4, 8, 6, 9] and [9, 10, 11]
- [2, 4, 5, 6, 9] and [1, 3, 7, 10, 11]
- [3, 5, 6, 9, 10] and [1, 2, 4, 7, 11]
- None of the above

7. How many integer multiplications does Strassen's algorithm need to do to multiply two 4×4 matrices of integers?

Answer.

- 7
- 8
- 9
- 49
- 64
- None of the above

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8. Running the Median-of-Medians algorithm with group size 5 to select the 3rd smallest element from the array [8, 2, 4, 13, 6, 1, 7, 10, 14, 15, 9, 6, 11, 3, 12], which number is the first pivot used to partition the array?

Answer.

- 7
- 12
- 9
- 5
- 8
- None of the above

9. Consider running InsertionSort (for sorting in ascending order) on an array with the first $n - \lceil n^{1/3} \rceil$ elements already sorted in ascending order. What is the tightest upper bound of its running time?

Answer.

- $O(n^{1/3})$
- $O(n - (n)^{1/3})$
- $O(n^{4/3})$
- $O(n \log n)$
- $O(n^2)$

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True/False (20 points)

True or false? Fill in the correct bubble. No justification is needed.

T F

- ① $n^{0.1} \log n = O(\log n)$.
- ② $\log n/n = O(n^{1/n})$.
- ③ $(\sqrt{n})^n = \Omega(n^{\sqrt{n}})$.
- ④ $\log(n!) = \Omega(n)$.
- ⑤ If $f(n)$ and $g(n)$ are non-negative functions, then either $f(n) = O(g(n))$ or $g(n) = O(f(n))$.
- ⑥ If $f(n) = O(g(n))$, then there must be some value of $n > 0$ where $f(n) \leq g(n)$.
- ⑦ If $f(n)$ is a nonnegative function, then $f(n+1) = \Theta(f(n))$.
- ⑧ The best-case time complexity of InsertionSort is $O(n^2)$.
- ⑨ The information theory lower bound of the comparison-based insertion problem is $\Omega(n)$.
- ⑩ The median-of-medians selection algorithm runs in $O(n)$ time with group size 7.

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Recurrences (12 points)

Each of the following scenarios outlines a divide-and-conquer algorithm. In each case, write down the appropriate recurrence relation for the running time as a function of the input size n and give its solution. Giving your solution in big- O notation. You need not give a full derivation of your solution, but you should indicate how you arrived at it (e.g., by using the Master Theorem). You may assume that n is of some special form (e.g., a power of two or of some other number), and that the recurrence has a convenient base case with cost $\Theta(1)$.

- (a) An input of size n is broken down into 64 subproblems, each of size $n/4$. The time taken to construct the subproblems, and to combine their solutions, is $\Theta(n^3)$.

$$\frac{n}{4}$$

let $n=1$

$$64 \times \frac{1}{4} = 13.5$$

so the answer is $\Theta(n^3)$

- (b) An input of size n is broken down into four subproblems of sizes 19, 29, 39, and $n - 1$ respectively. The time taken to construct the subproblems, and to combine their solutions, is 1.9^n .

because it is always $\Theta(n \log n)$

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Algorithms. (15 points)

2/2

Suppose the instructor is considering several possible grading curves for the exam, and you are asked to collect students' opinions. Let an array $A[1..n]$ represent the choices of n students. A *majority* choice is one that is selected by more than $n/2$ students. For example, $A = [a, a, b, a, c, c, a]$ has a majority choice a because it appears in $4 > 7/2$ positions; on the other hand, $A = [1, 4, 3, 1, 3, 3]$ does not have a majority choice because no choice appears in more than $6/2 = 3$ positions.

You would like to design a divide-and-conquer algorithm to find the majority choice (if it exists):

- (a) Consider the following array $A[1..16] = [5, 2, 2, 5, 5, 2, 5, 2 | 3, 5, 3, 5, 5, 5, 2, 5]$. Divide it into two equal-sized halves. Let M_L and M_R be the majority choices of the left and right halves of the array, respectively. Do M_L and M_R exist? If so, what are they?

$$M_L = 3 \quad M_R = 5$$

M_{left} didn't $\Rightarrow [3, 5, 3, 5, 5, 5, 2, 5]$

exit because no choice more than

- (b) In the example of part (a), does the array A have a majority choice? In general, how is the majority choice of A related to M_L and M_R ? Justify your answer.

Yes, A majority choice is $5 > \frac{16}{2}$ position

Only 9 times 5 appear and there is 16 position, so only the side have more than 4 times about 5 have M_R or M_L

- (c) Based on your answer for part (b), describe a general divide-and-conquer algorithm that finds a majority choice of an array or return NULL if it doesn't exist. What is the running time of your algorithm?

3

3 3 5 5 5 5 5 5 2 2 2 2

Should be NULL

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$$\begin{array}{r} 45867 \\ \times 45867 \\ \hline 31400 \\ 18320 \\ \hline 45867 \\ 45867 \\ \hline 458671 \\ 1234578469 \\ 1234578469 \\ \hline 15327846 \\ 15327846 \\ \hline 15367842 \\ 15367842 \\ \hline 458671 \\ 458671 \\ \hline 4586783 \\ 15467843 \\ 15367843 \\ \hline 451678 \\ 451678 \\ \hline 458678 \\ 458678 \\ \hline 8477 \\ 8477 \\ \hline \end{array}$$

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$$\begin{array}{r}
 492 \\
 4213785^6 \\
 -17 \quad 123 \\
 \hline
 1 \\
 17 \quad 123 \\
 \hline
 177 \\
 2943 \\
 -23491 \\
 \hline
 21493 \\
 -21 \quad 0 \\
 \hline
 0 \\
 124937 \\
 -124937 \\
 \hline
 0 \\
 1249376 \\
 -1249376 \\
 \hline
 0 \\
 109373 \\
 -109373 \\
 \hline
 0 \\
 41 \\
 -41 \\
 \hline
 0 \\
 109373 \\
 -109373 \\
 \hline
 0 \\
 41 \\
 -41 \\
 \hline
 0 \\
 109373 \\
 -109373 \\
 \hline
 0 \\
 1249376 \\
 -1249376 \\
 \hline
 0
 \end{array}$$