

Midterm1

● Graded

Student

Jesse Zhang

Total Points

28.5 / 74 pts

Question 1

Multiple choice

12 / 27 pts

1.1 — MC1

3 / 3 pts

✓ - 0 pts Correct

- 3 pts Incorrect

1.2 — MC2

3 / 3 pts

✓ - 0 pts Correct

- 3 pts Incorrect

1.3 — MC3

3 / 3 pts

✓ - 0 pts Correct

- 3 pts Incorrect

1.4 — MC4

3 / 3 pts

✓ - 0 pts Correct

- 3 pts Incorrect

1.5 — MC5

0 / 3 pts

- 0 pts Correct

✓ - 3 pts Incorrect

1.6 — MC6

0 / 3 pts

- 0 pts Correct

✓ - 3 pts Incorrect

1.7 — MC7

0 / 3 pts

- 0 pts Correct

✓ - 3 pts Incorrect

1.8 — MC8

0 / 3 pts

- 0 pts Correct

✓ - 3 pts Incorrect

1.9 — MC9

0 / 3 pts

- 0 pts Correct

✓ - 3 pts Incorrect

Question 2

True/False

11 / 20 pts

- 2.1 — **TF1** 0 / 2 pts
- 0 pts Correct
- ✓ — 2 pts Incorrect
- 2.2 — **TF2** 2 / 2 pts
- ✓ — 0 pts Correct
- 2 pts Incorrect
- 2.3 — **TF3** 2 / 2 pts
- ✓ — 0 pts Correct
- 2 pts Incorrect
- 2.4 — **TF4** 2 / 2 pts
- ✓ — 0 pts Correct
- 2 pts Incorrect
- 2.5 — **TF5** 0 / 2 pts
- 0 pts Correct
- ✓ — 2 pts Incorrect
- 2.6 — **TF6** 0 / 2 pts
- 0 pts Correct
- ✓ — 2 pts Incorrect
- 2.7 — **TF7** 2 / 2 pts
- ✓ — 0 pts Correct
- 2 pts Incorrect
- 2.8 — **TF8** 1 / 2 pts
- 0 pts Correct
- ✓ — 1 pt Incorrect
- 2.9 — **TF9** 0 / 2 pts
- 0 pts Correct
- ✓ — 2 pts Incorrect

✓ - 0 pts Correct

- 2 pts Incorrect

Question 3

Recurrences

0 / 12 pts

3.1 (a)

0 / 6 pts

- 0 pts Part 1: Recurrence relation for running time [3 pts/3 pts]

Correctly wrote down the recurrence relation for running time as $T(n) = 64T(n/4) + \Theta(n^3)$.

- 1 pt Part 1: Recurrence relation for running time [2 pts/3 pts]

Could not correctly write down the entire recurrence relation, but the student was on the right track to do so by expressing the running time T as a function of input size n and showing some of the terms in recurrence correctly.

✓ - 3 pts Part 1: Recurrence relation for running time [0 pts/3 pts]

Did not attempt at all or did not express the recurrence relation for running time T as a function of input size n (no mentioning of $T(n)$ at all).

- 0 pts Part 2: Solution for recurrence relation in O [3 pts/3 pts]

Correctly wrote down the solution of the recurrence relation as $O(n^3 \log n)$ using Master Theorem.

- 1 pt Part 2: Solution for recurrence relation in O [2 pts/3 pts]

Did not apply the Master Theorem and attempted to unfold the recurrence relation but could not reach the correct solution. Or found the right values for a, b , and d but could not reach the correct solution.

✓ - 3 pts Part 2: Solution for recurrence relation in O [0 pts/3 pts]

Did not attempt at all or incorrectly **applied the Master Theorem** to solve the recurrence relation.

- 0 pts Part 1: Recurrence relation for running time [3 pts/3 pts]

Correctly wrote down the recurrence relation for running time as $T(n) = T(19) + T(29) + T(39) + T(n-1) + 1.9^n$ and *optionally* simplified the expression to $T(n) = T(n-1) + 1.9^n$.

- 1 pt Part 1: Recurrence relation for running time [2 pts/3 pts]

Could not correctly write down the entire recurrence relation, but the student was on the right track to do so by expressing the running time T as a function of input size n and showing some of the terms in recurrence correctly.

✓ - 3 pts Part 1: Recurrence relation for running time [0 pts/3 pts]

Did not attempt at all or did not express the recurrence relation for running time T as a function of input size n (no mentioning of $T(n)$ at all).

- 0 pts Part 2: Solution for recurrence relation in O [3 pts/3 pts]

Correctly wrote down the solution of the recurrence relation as $O(1.9^n)$ either by *mentioning HW1 Q5(e)* or unfolding the recurrence.

- 1 pt Part 2: Solution for recurrence relation in O [2 pts/3 pts]

Did not apply the Master Theorem and attempted to unfold the recurrence relation but could not reach the correct solution.

✓ - 3 pts Part 2: Solution for recurrence relation in O [0 pts/3 pts]

Did not attempt at all or incorrectly **applied the Master Theorem** to solve the recurrence relation.

Question 4

Algorithms

5.5 / 15 pts

4.1 (a)

2.5 / 5 pts

– 0 pts Correct

✓ – 2.5 pts Only one of M_L or M_R is incorrect

– 5 pts Both M_L and M_R are incorrect/not written

4.2 (b)

3 / 5 pts

– 0 pts Correct

– 1 pt Justification for relation between M_L and M_R to majority of A not provided/incorrect

✓ – 2 pts Incorrect answer for how M_L and M_R are related to majority of A, or provided the relation specific to the example not in general

– 3 pts Missing answer for how M_L and M_R are related to majority of A

– 2 pts Missing/Incorrect answer for majority of A

4.3 (c)

0 / 5 pts

– 0 pts Correct

✓ – 5 pts wrong algorithm and time complexity

– 4.75 pts only running time correct

– 0.5 pts wrong running time

– 3 pts idea of the algorithm

– 1 pt partially correct algorithm

– 4 pts attempt on algorithm

– 2 pts mostly correct

Midterm 1 - D

Name: Jesse Zhang

Penn State access ID (xyz1234) in the following box:

2fz 5311

Student ID number (9XXXXXXXXX):

975557148

TA and/or section time:

MWF 4:35

Instructions:

- Answer all questions. Read them carefully first. Be precise and concise. Handwriting needs to be neat. Box numerical final answers.
- Please clearly write your name and your PSU access ID (i.e., xyz1234) in the box on top of every page.
- Write in only the space provided. You may use the back of the pages only as scratch paper. Do not write your solutions in the back of pages!
- Do not write outside of the black bounding box: the scanner will not be able to read anything outside of this box.
- We are providing two extra page at the end if you need extra space. Make sure you mention in the space provided that your answer continues there.

Good luck!

Name: Jesse Zhang

PSU Access ID (xyz1234): 2fc5311

Multiple choice questions (27 points)

For each of the following questions, select the right answer by filling the corresponding grading bubble.

1. If $T(n) = 3T(n/2) + O(n^2)$, which of the following is the tightest upper bound for $T(n)$?

Answer.

- ☐ $T(n) = O(n^{\log_3 2})$
☐ $T(n) = O(n^{\log 3})$
☒ $T(n) = O(n^2)$
☐ $T(n) = O(n^2 \log n)$
☐ None of the above

2. While performing InsertionSort to sort the array $[4, 1, 2, 3, 7, 8, 5, 6]$ in descending order, which of the following is NOT a transition state of the array?

Answer.

- ☐ $[4, 2, 1, 3, 7, 8, 5, 6]$
☐ $[8, 7, 4, 3, 2, 1, 5, 6]$
☐ $[4, 3, 2, 1, 7, 8, 5, 6]$
☐ $[7, 4, 3, 2, 1, 8, 5, 6]$
☐ $[8, 7, 5, 4, 3, 2, 1, 6]$
☒ None of the above

3. If $T(n) = 9T(n/3) + O(n^2 + n(\log n)^5)$, which of the following is the tightest upper bound for $T(n)$?

Answer.

- ☐ $T(n) = O(n^2)$
☒ $T(n) = O(n^2 \log n)$
☐ $T(n) = O(n(\log n)^5)$
☐ $T(n) = O(n^2(\log n)^5)$
☐ None of the above

4. Suppose that $\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0$. Which of the following must be true?

Answer.

- ☒ $f(n) = O(g(n))$
☐ $f(n) = \Omega(g(n))$
☐ $f(n) = \Theta(g(n))$
☐ None of the above

Name: Jesse Zhang

PSU Access ID: 2f2531

5. Consider the Partition function used in QuickSort. If we choose the last element as the pivot, what is the resulting array after calling Partition on [5, 8, 6, 7, 1, 3, 2, 9, 4]?

Answer.

- ☐ [1, 2, 3, 4, 5, 6, 7, 8, 9]
☐ [1, 3, 2, 4, 5, 8, 6, 7, 9]
☐ [1, 2, 3, 4, 9, 8, 7, 6, 5]
☐ [1, 3, 2, 4, 5, 8, 6, 9, 7]
☐ [1, 3, 2, 7, 5, 8, 6, 9, 4]
☒ None of the above

4 5 8 6 7 1
4 5 1 6 7 8
4 1 8 6 3 1 5 4 6 7 8 9
1 4 8 6 5 7 1 5 3 6 7 8 4 2
1 4 7 6 5 8 1 5 3 2 7 8 4 6
1 2 3 5 7 8 4 6 9

6. Run MergeSort to sort the array [5, 2, 9, 4, 6, 1, 10, 7, 3, 11] in ascending order, which two subarrays are the inputs to the final Merge call that produces the fully sorted array?

Answer.

- ☐ [5, 2, 9, 4, 6] and [1, 10, 7, 3, 11]
☐ [2, 5, 4, 8, 6] and [2, 10, 3, 7, 11]
☒ [2, 4, 4, 8, 6, 9] and [3, 7, 10, 11]
☐ [2, 4, 5, 6, 9] and [1, 3, 7, 10, 11]
☐ [3, 5, 6, 9, 10] and [1, 2, 4, 7, 11]
☐ None of the above

1 5
1 2 5 4 9
1 2 3 4 5 6 7 9 11

7. How many integer multiplications does Strassen's algorithm need to do to multiply two 4×4 matrices of integers?

Answer.

- ☐ 7
☐ 8
☐ 9
☐ 49
☒ 64
☐ None of the above

Name: Jesse Zhang

PSU Access ID: zt5311

8. Running the Median-of-Medians algorithm with group size 5 to select the 3rd smallest element from the array [8, 2, 4, 13, 6, 1, 7, 10, 14, 15, 9, 6, 11, 3, 12], which number is the first pivot used to partition the array?

Answer.

- ☐ 7
☐ 12
☐ 9
☒ 5
☐ 8
☐ None of the above

9. Consider running InsertionSort (for sorting in ascending order) on an array with the first $n - \lceil n^{1/3} \rceil$ elements already sorted in ascending order. What is the tightest upper bound of its running time?

Answer.

- ☐ $O(n^{1/3})$
☐ $O(n - (n)^{1/3})$
☐ $O(n^{4/3})$
☐ $O(n \log n)$
☒ $O(n^2)$

Name: ~~zfsz~~ Jesse Zhang

PSU Access ID (xyz1234): zfsz11

True/False (20 points)

True or false? Fill in the correct bubble. No justification is needed.

T F

☒ ☐ $n^{0.1} \log n = O(\log n)$.

☒ ☐ $\log n / n = O(n^{1/n})$.

☒ ☐ $(\sqrt{n})^n = \Omega(n^{\sqrt{n}})$.

☒ ☐ $\log(n!) = \Omega(n)$.

☒ ☐ If $f(n)$ and $g(n)$ are non-negative functions, then either $f(n) = O(g(n))$ or $g(n) = O(f(n))$.

☒ ☐ If $f(n) = O(g(n))$, then there must be some value of $n > 0$ where $f(n) \leq g(n)$.

☐ ☒ If $f(n)$ is a nonnegative function, then $f(n+1) = \Theta(f(n))$.

☐ ☒ The best-case time complexity of InsertionSort is $O(n^2)$.

☒ ☐ The information theory lower bound of the comparison-based insertion problem is $\Omega(n)$.

☒ ☐ The median-of-medians selection algorithm runs in $O(n)$ time with group size 7.

Name: Jesse Zhang

PSU Access ID: zfs311

Recurrences (12 points)

Each of the following scenarios outlines a divide-and-conquer algorithm. In each case, write down the appropriate recurrence relation for the running time as a function of the input size n and give its solution. Giving your solution in big- O notation. You need not give a full derivation of your solution, but you should indicate how you arrived at it (e.g., by using the Master Theorem). You may assume that n is of some special form (e.g., a power of two or of some other number), and that the recurrence has a convenient base case with cost $\Theta(1)$.

- (a) An input of size n is broken down into 64 subproblems, each of size $n/4$. The time taken to construct the subproblems, and to combine their solutions, is $\Theta(n^3)$.

$$r \frac{n}{4}$$
$$n \text{ at } n=1$$
$$64 \times \frac{1}{4} = 13.5$$

so the complexity is $\Theta(n^3)$

- (b) An input of size n is broken down into four subproblems of sizes 19, 29, 39, and $n-1$ respectively. The time taken to construct the subproblems, and to combine their solutions, is 1.9^n .

because it is always defined $\Theta(\log n)$

Name: Jesse Zhang

PSU Access ID: 2425311

Algorithms. (15 points)

2/1

Suppose the instructor is considering several possible grading curves for the exam, and you are asked to collect students' opinions. Let an array $A[1..n]$ represent the choices of n students. A *majority* choice is one that is selected by more than $n/2$ students. For example, $A = [a, a, b, a, c, a]$ has a majority choice a because it appears in $4 > 7/2$ positions; on the other hand, $A = [1, 4, 3, 1, 3, 3]$ does not have a majority choice because no choice appears in more than $6/2 = 3$ positions.

You would like to design a divide-and-conquer algorithm to find the majority choice (if it exists):

- (a) Consider the following array $A[1..16] = [5, 2, 2, 5, 5, 2, 5, 2, 3, 5, 3, 5, 5, 5, 2, 5]$. Divide it into two equal-sized halves. Let M_L and M_R be the majority choices of the left and right halves of the array, respectively. Do M_L and M_R exist? If so, what are they?

$$M_L = 8 \quad M_R = 8$$

only the part contain more than 4 times are exists
5 and

$$M_{\text{left didn't}} = [3, 5, 3, 5, 5, 5, 2, 5]$$

exit because no choice more than

- (b) In the example of part (a), does the array A have a majority choice? In general, how is the majority choice of A related to M_L and M_R ? Justify your answer.

Yes, A majority choice is 5 > 16/2 position

only 9 times 5 appear and there is 16 position, so only the side have more than 4 times about 5 have M_R or M_L

- (c) Based on your answer for part (b), describe a general divide-and-conquer algorithm that finds a majority choice of an array or return NULL if it doesn't exist. What is the running time of your algorithm?

3

3 3 5 5 5 5 5 5 2 2 2 2

should be NULL

Name:

Jesse Zhang

PSU Access ID (xyz1234):

2f25311

457681

45867

15367842

15327846

123578469

45867

1235 458671

3

45

4/64
14
12

45

45867

45867

8 478

451678

1546783

1536784

Name: ~~2f2531~~ Jesse Zhang PSU Access ID (xyz1234): 2f2531

$4 \ 2 \ 1 \ 3 \ 7 \ 8 \ 5^6$
 $1 \ 3 \ 12$
 $1 \ 2 \ 5 \ 3 \ 7^4$
 $4 \ 15$
 $a=3 \ b=2 \ d=2$
 $109 \ 2 \ 3 \ d=2$
 $2 \ 8 \ 5$
 $4 \ 2 \ 1 \ 3 \ 7 \ 8 \ 5^6$
 $1 \ 3 \ 12$
 $4 \ 32 \ 7 \ 7$
 $7 \ 11$
 11
 $4 \ 4 \ 4 \ 4$
 2
 $4 \ 9 \ 2$
 $2 \ 9 \ 4 \ 3$
 $2 \ 3 \ 4 \ 9 \ 1$
 $2 \ 1 \ 4 \ 9 \ 3$
 $9 \ 1 \ 2 \ 4 \ 9 \ 3 \ 7$
 $12 \ 4 \ 9 \ 3 \ 7 \ 6$
 $d=5$
 9