

MISSING NUMBER IN AN ARRAY

GEEKS FOR GEEKS:[MISSING NUMBER](#)

LEETCODE:[MISSING NUMBER](#)

APPROACH-1:

USING THE MATHEMATICAL FORMULA BY COMPUTING $N*(N+1)/2$

this will give sum of n sized elements in array.

and step 2:

i will calculate the sum of array elements using a simple formula and store that answer in the other variable..

when i subtract (1)-(2) i will get the missing number in the array.

```
int MissingNumber(vector<int>& array, int n) {  
  
    int sum=0;  
    for(int i=0;i<n-1;i++){  
        sum=sum+array[i];  
    }  
    int p=n*(n+1)/2;  
    return (p-sum);  
  
}
```

TIME COMPLEXITY = $O(N)$

APPROACH-2:

USING XOR THE DISADVANTAGE FOR THE ABOVE PROCESS IS IF WE PROCEED TO MUCH FOR THE INPUT SIZE N

THE COST OF COMPUTING FOR $N*N+1$ IS VERY HEAVY THAT THE VARIABLE STORING ITS RESULT COULDN'T RESIST AS ITS LIMIT WOULD BE EASILY CROSSED.....

XOR IS LIKE FIRST PERFORM THE XOR FOR N SIZED NATURAL ELEMENTS ...

PERFORM XOR ON ARRAY ELEMENTS...

NOW PERFORM XOR ON THE COMBINED RESULTS AND RETURN THE ANSWER...

NO ADDITIONAL HEADACHE IS OBSERVED HERE..

```
**USING XOR OPERATION **  
THE MAIN THEME HERE IS FIRST PERFORM XOR ON N SIZED NATURAL CONSECUTIVE ELEMENTS.  
NEXT PERFORM XOR ON COMPLETE GIVEN ARRAY.
```

```
class Solution {
public:
    int missingNumber(vector<int>& nums) {
        int n= nums.size();
        int res1=0, res2=0;
        for(int i=0;i<n;i++){
            res1=res1^nums[i];

        }
        for(int i=0;i<=n;i++){
            res2=res2^i;
        }
        return res1^res2;
    }
};
```