I/O-efficient Symbolic Model Checking

Steffan Christ Sølvsten, Jaco van de Pol

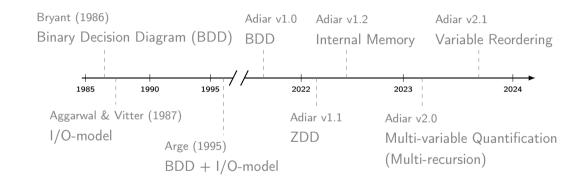
31st of August, 2022

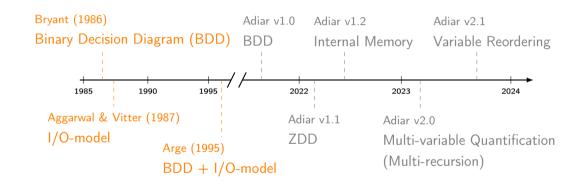


Adiar

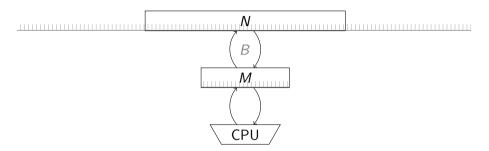
I/O-efficient Decision Diagrams

github.com/ssoelvsten/adiar





Aggarwal and Vitter '87: I/O-model



The I/O-model by Aggarwal and Vitter '87

Aggarwal and Vitter '87: I/O-model

For any realistic values of N, M, and B we have that

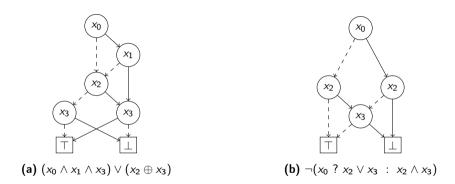
$$N/B < \operatorname{sort}(N) \triangleq N/B \cdot \log_{M/B} N/B \ll N$$
,

Theorem (Aggarwal and Vitter '87) N elements can be sorted in $\Theta(sort(N))$ I/Os.

Theorem (Arge '95)

N elements can be inserted in and extracted from a Priority Queue in $\Theta(sort(N))$ I/Os.

Bryant '86: Binary Decision Diagram

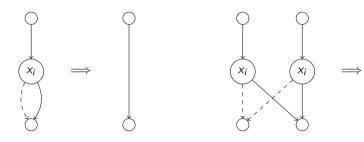


Examples of (Reduced Ordered) Binary Decision Diagrams.

Bryant '86: Binary Decision Diagram

Theorem

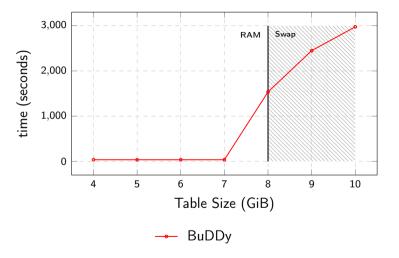
For a fixed variable order, if one exhaustively applies the two rules below, then one obtains the Reduced OBDD, which is a unique canonical form of the function.



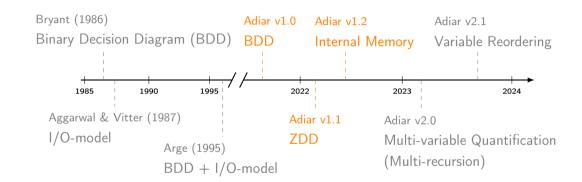
(1) Remove redundant nodes

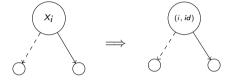
(2) Merge duplicate nodes

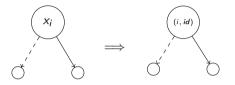
Arge '95 : BDD + I/O-model



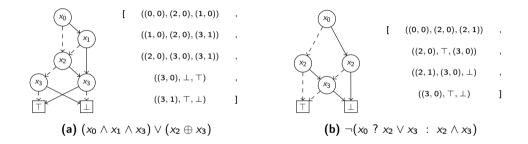
Running time for solving a problem that does not need more than 3 $\,\mathrm{GiB}.$



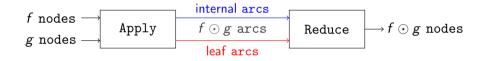


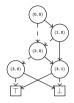


$$(i_1, id_1) < (i_2, id_2) \equiv i_1 < i_2 \lor (i_1 = i_2 \land id_i < id_j)$$



Node-based representation of prior shown BDDs

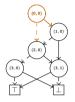




(a) $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$



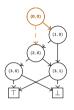
(b) $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$



(a) $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$



(b) $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$



Priority Queue: Qapp:1:

 $[(0,0) \xrightarrow{\top} ((1,0),(2,1))$

 $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$

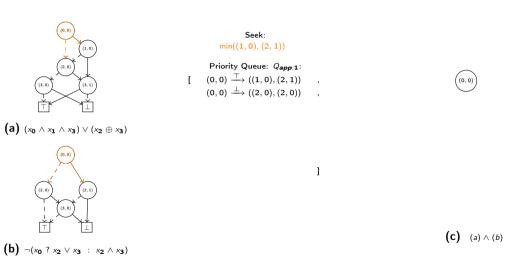
(0,0)

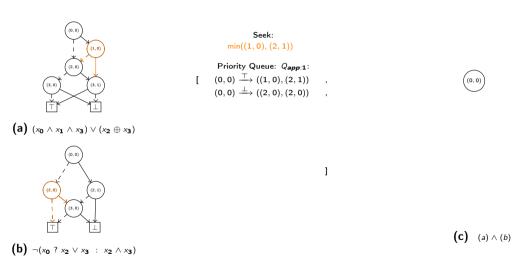
(a) $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$

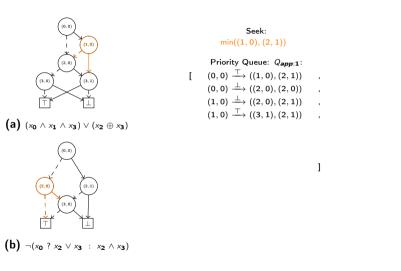


(b) $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$

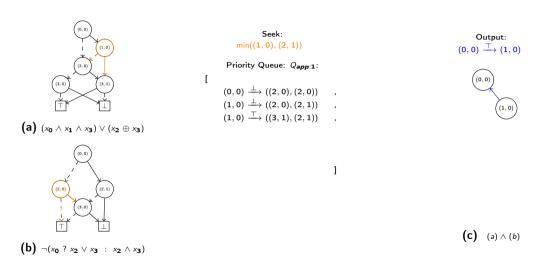
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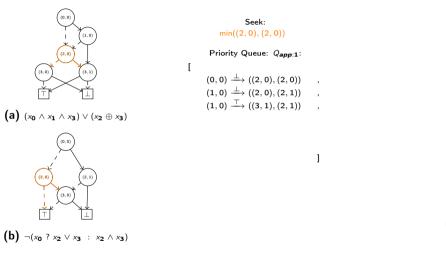






(1,0)

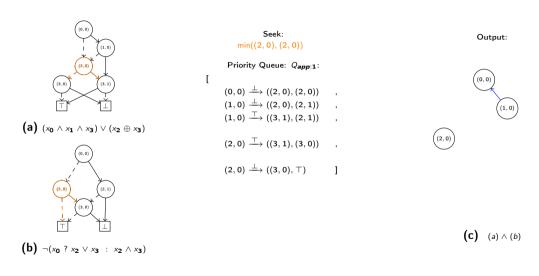


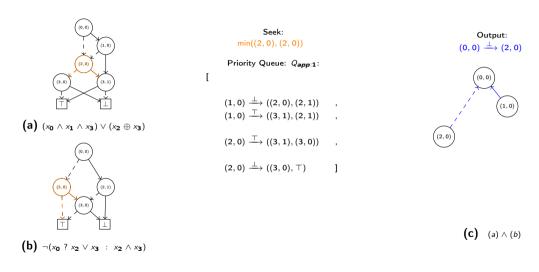


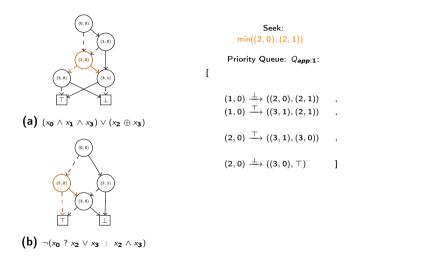
Output:



(c) (a) ∧ (b)



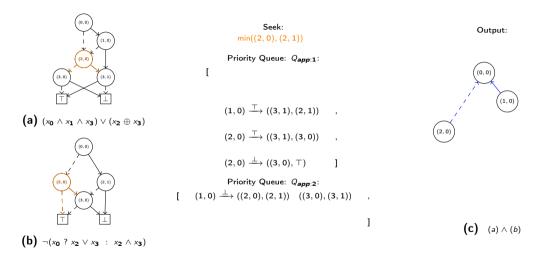


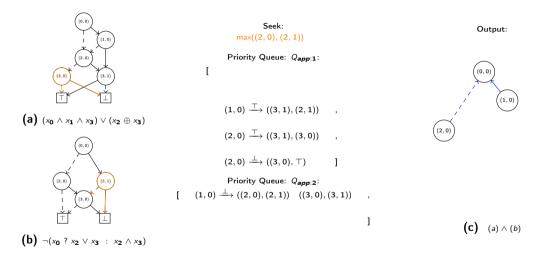


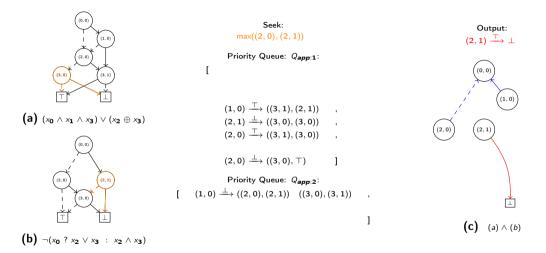
Output:

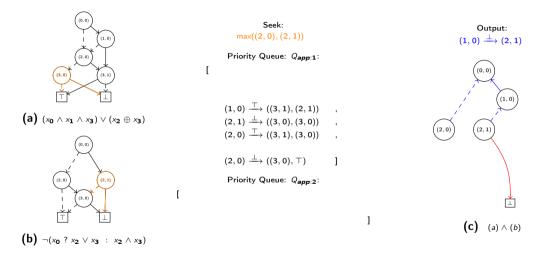


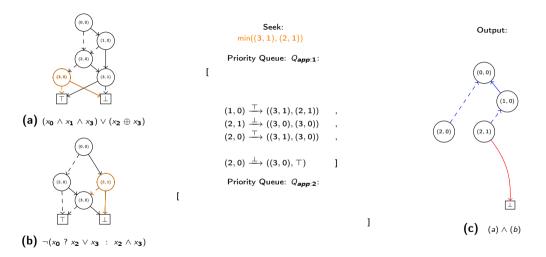
(c) (a) ∧ (b)

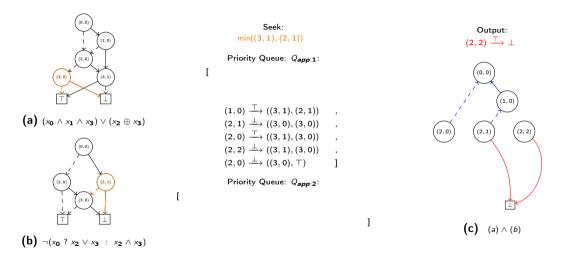


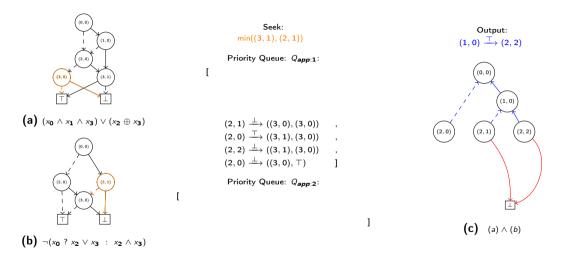


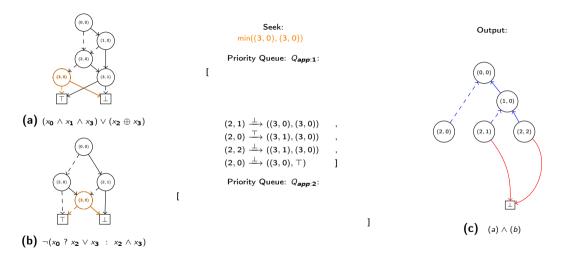


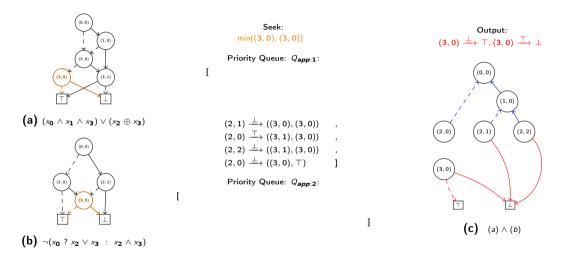


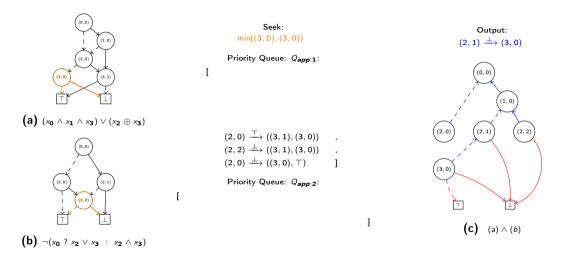


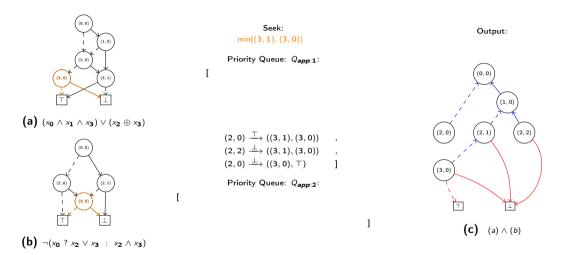


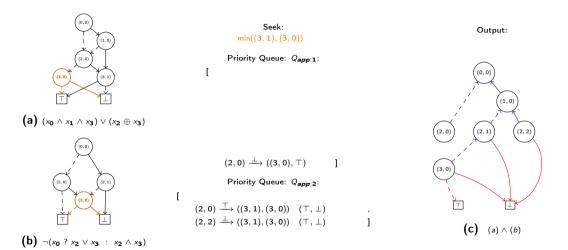


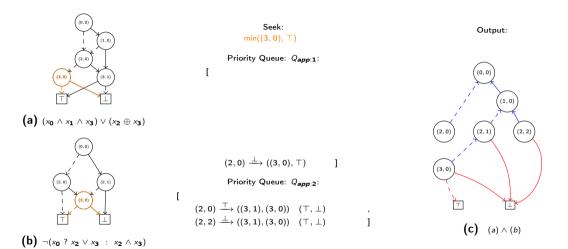


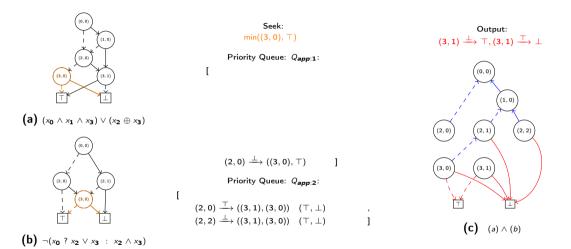


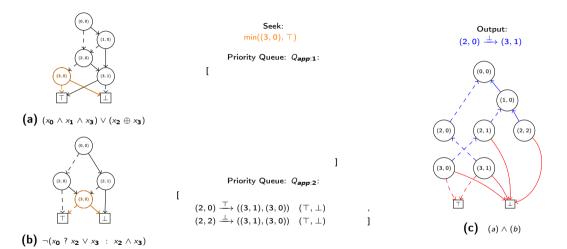


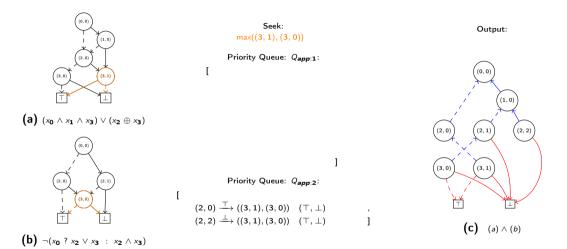


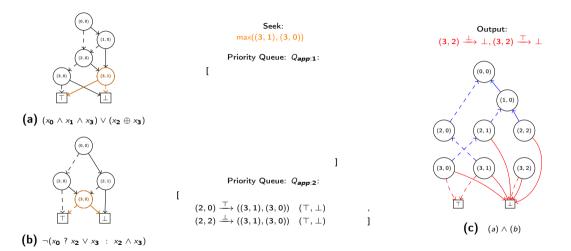


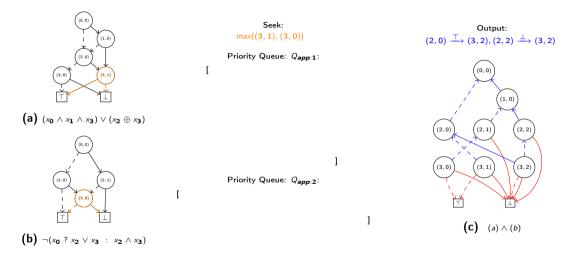


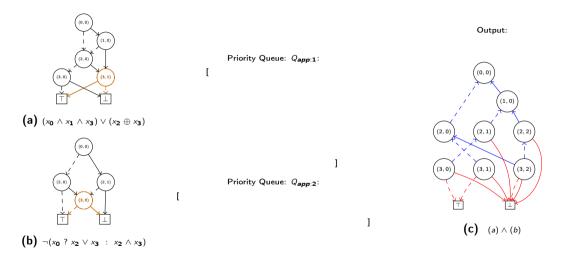


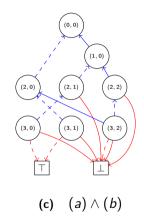


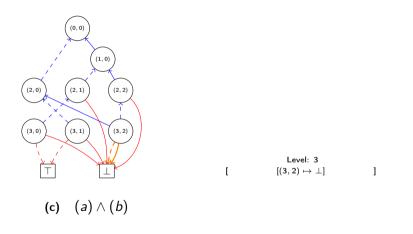


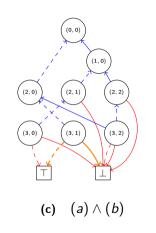


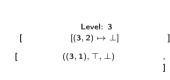


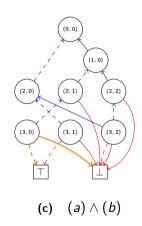


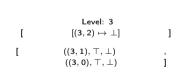


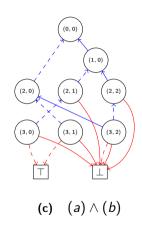


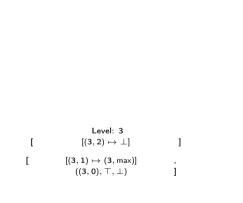


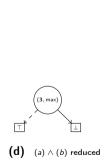




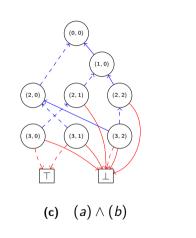


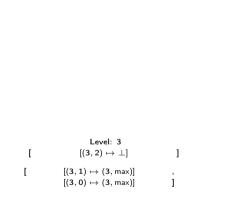


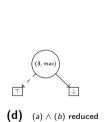




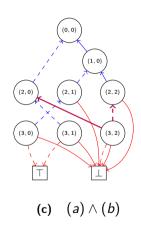
Output: $((3, max), \top, \bot)$

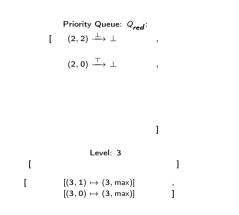




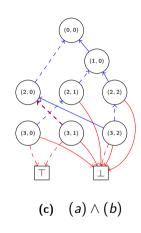


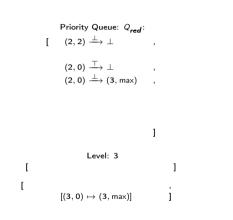
Output:



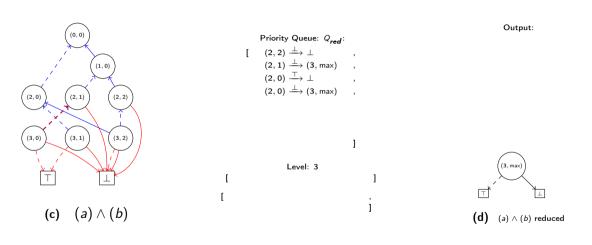


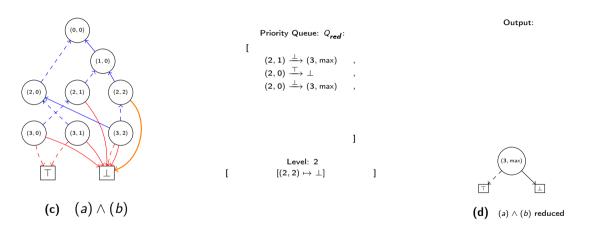


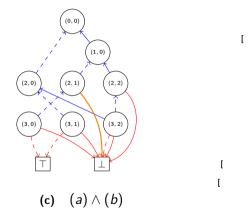


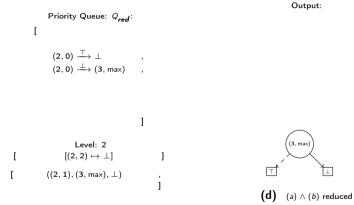


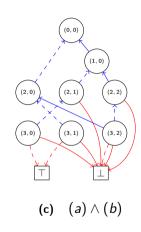


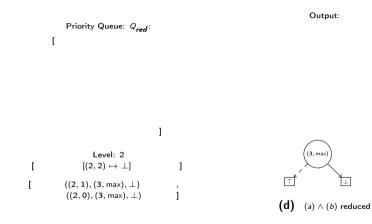


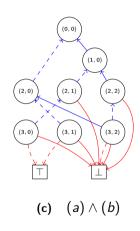


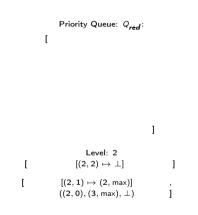


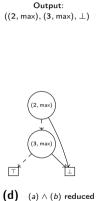


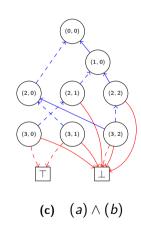


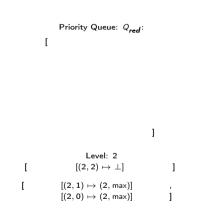


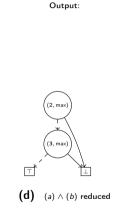


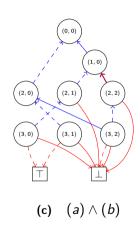


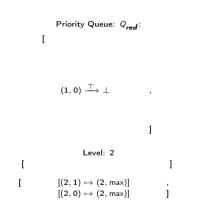


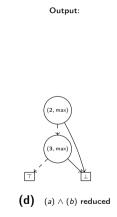


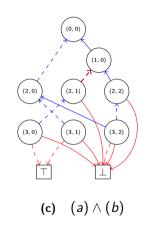


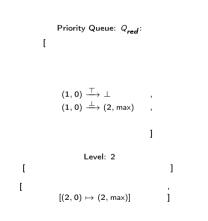


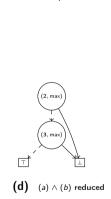




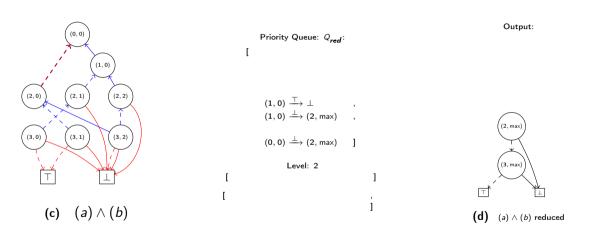


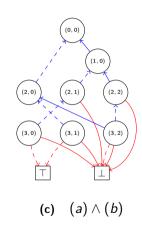


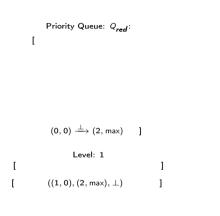


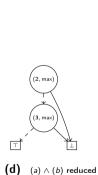


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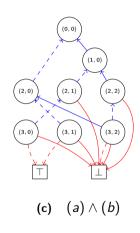


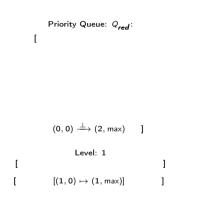


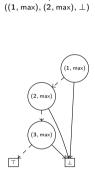




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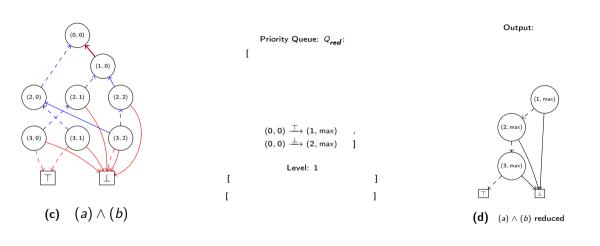


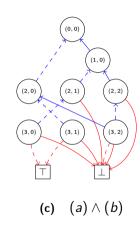


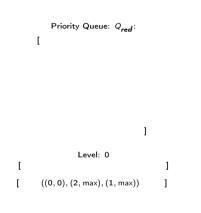


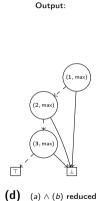
(d) $(a) \wedge (b)$ reduced

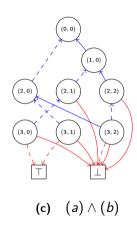
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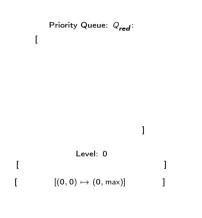


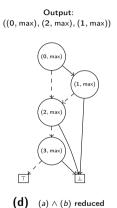


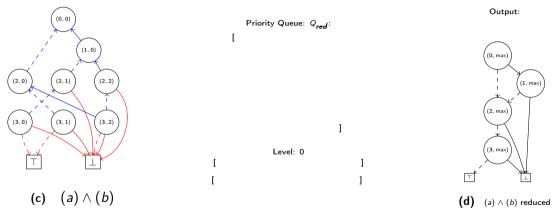


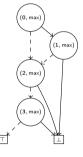


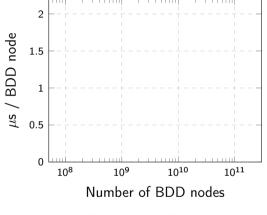








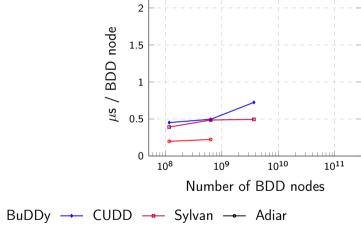




→ BuDDy → CUDD → Sylvan → Adiar

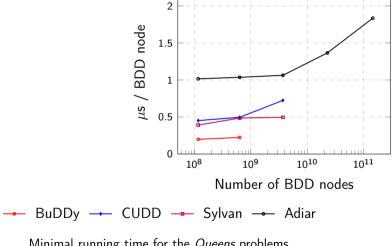
Minimal running time for the *Queens* problems.

Adiar v1.0: BDD



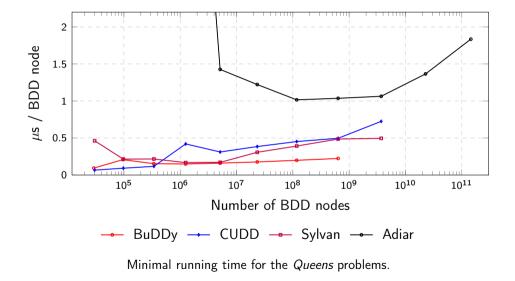
Minimal running time for the *Queens* problems.

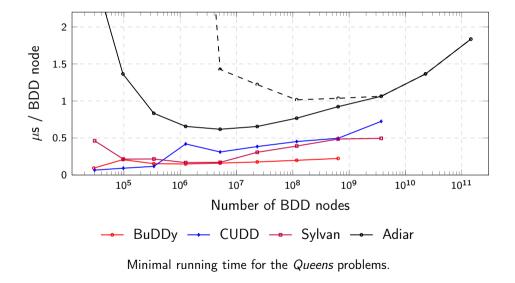
Adiar v1.0: BDD



Minimal running time for the Queens problems.

Adiar v1.0: BDD



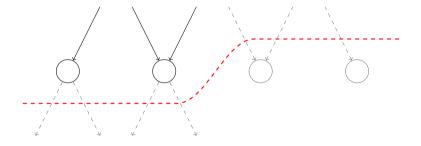


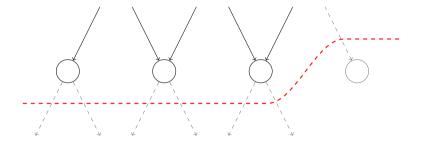
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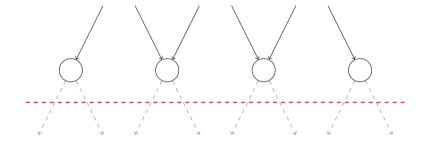




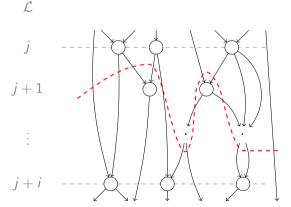








Definition (i-level cut)



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Lemma

The maximum i-level cut problem is in P for $i \in \{1, 2\}$.

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The maximum i-level cut problem is in P for $i \in \{1, 2\}$.

Theorem (Lampis, Kaouri, Mitsou 2011) The maximum i-level cut problem is NP-complete for $i \ge 4$.

Theorem

The maximum (i-level) cut of a BDD with N internal nodes is N+1.

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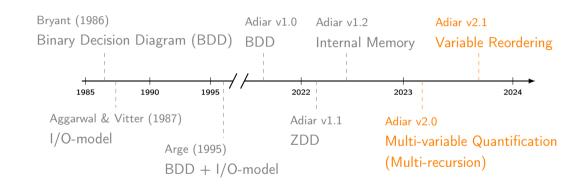
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Theorem

For $i \in \{1,2\}$, the maximum i-level cut of the (unreduced) output of Apply is upper bounded by the product of the inputs' corresponding i-level cuts.

Lemma

The maximum 2-level cut of a BDD is upper bounded by $\frac{3}{2}$ its 1-level cut.



$$RelProd(S, T) \equiv (\exists \vec{x}. S(\vec{x}) \land T(\vec{x}, \vec{x'}))[\vec{x'}/\vec{x}]$$

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1 exists (f, V)
     if f = 1 \lor f = \top
            then f
    else if V \cap \{i \in \mathbb{N} \mid i \geq \mathsf{top}(f)\} = \emptyset
5
            then f
     else if top(f) \notin V
            then Node { top(f), exists(f.low, V), exists(f.high, V) }
8
     else let low = exists(f.low. V)
                  high = exists(f.high, V)
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            in or(low, high)
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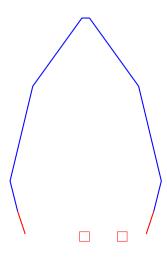
A recursive multi-variable exists operation.

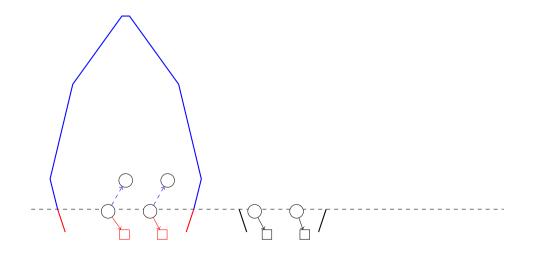
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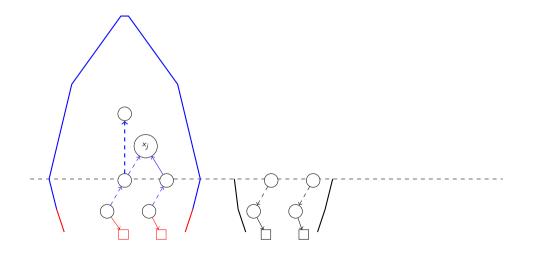
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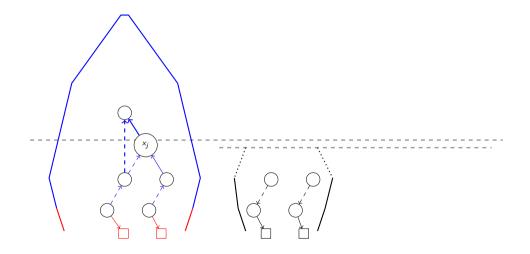
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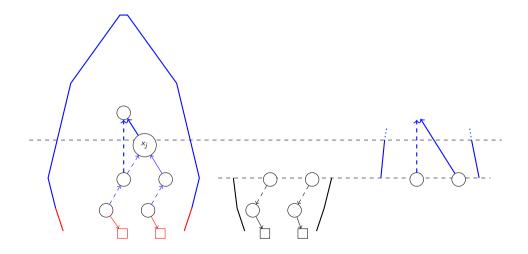
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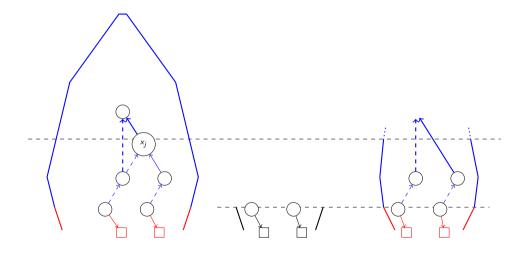


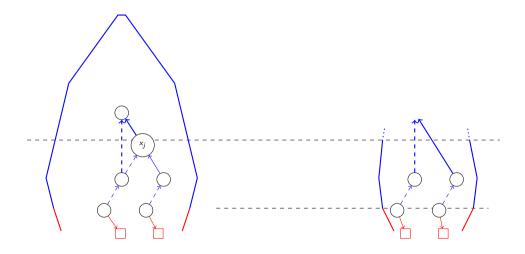


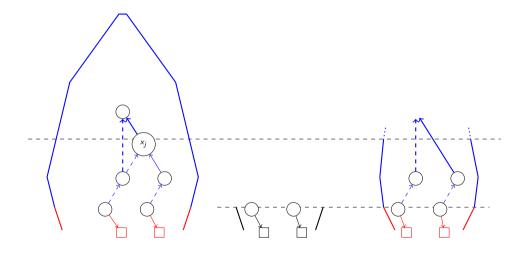


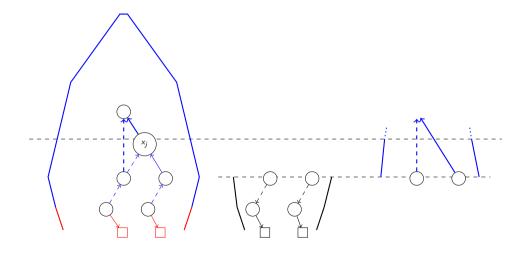


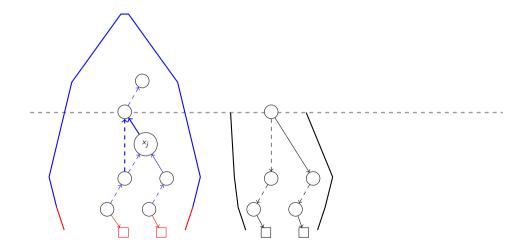


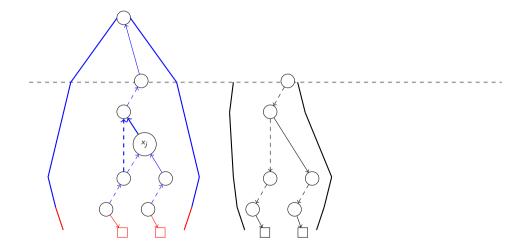


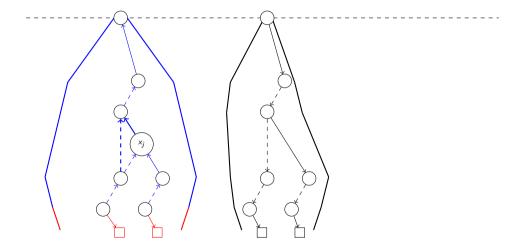


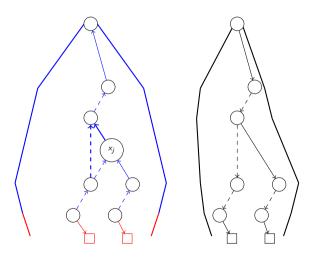






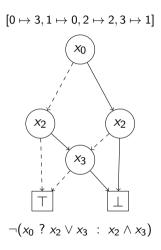




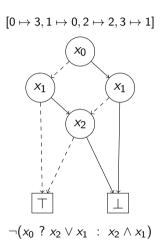


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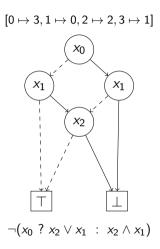
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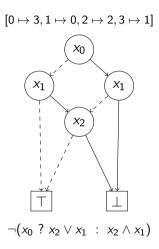
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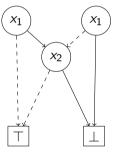
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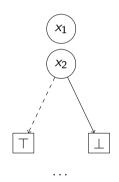
A recursive **substitute** operation.



. . .

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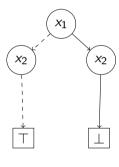
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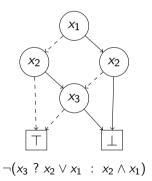
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 \frac{\text{Time Space } I/O}{O(NT) O(NT) O(NT)} 
Complexity of depth-first substitute

 \frac{\text{Time Space } I/O}{O(NT \log T) O(N+T) O(N \cdot sort(T))} 
Complexity of level-by-level substitute
```

Problem (Variable Replacement) Given BDD f_{π} with variable ordering π and remapping of variables $m : \mathbb{N} \to \mathbb{N}$, construct $f'_{\pi} \equiv f_{\pi}[x/m(x)]$.

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Problem (Dynamic Variable Reordering)

Given BDD $\hat{f_{\pi}}$ with variable ordering π ,

find π' and construct $f_{\pi'} \equiv f_{\pi}$ such that $|f_{\pi'}|$ is minimal.

