# I/O-efficient Symbolic Model Checking

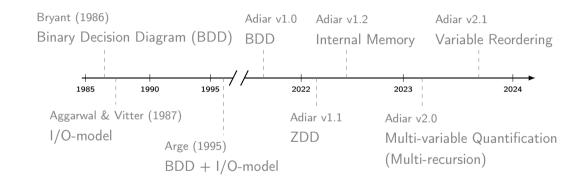
Steffan Christ Sølvsten, Jaco van de Pol

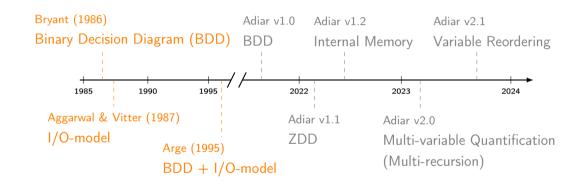
31st of August, 2022



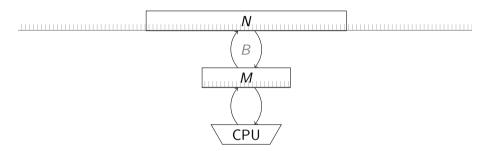
# **A**diar

github.com/ssoelvsten/adiar





# Aggarwal and Vitter '87: I/O-model



The I/O-model by Aggarwal and Vitter '87

# Aggarwal and Vitter '87: I/O-model

For any realistic values of N, M, and B we have that

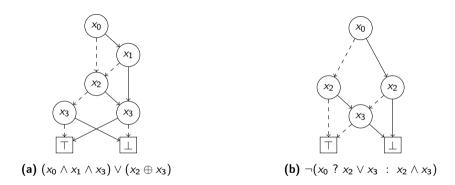
$$N/B < \operatorname{sort}(N) \triangleq N/B \cdot \log_{M/B} N/B \ll N$$
,

Theorem (Aggarwal and Vitter '87) N elements can be sorted in  $\Theta(sort(N))$  I/Os.

# Theorem (Arge '95)

N elements can be inserted in and extracted from a Priority Queue in  $\Theta(sort(N))$  I/Os.

# Bryant '86: Binary Decision Diagram

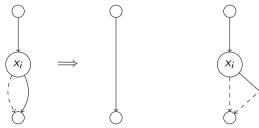


Examples of (Reduced Ordered) Binary Decision Diagrams.

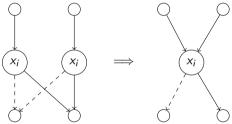
# Bryant '86: Binary Decision Diagram

#### **Theorem**

For a fixed variable order, if one exhaustively applies the two rules below, then one obtains the Reduced OBDD, which is a unique canonical form of the function.

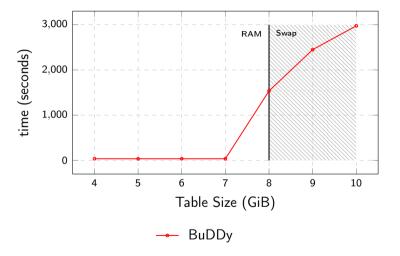




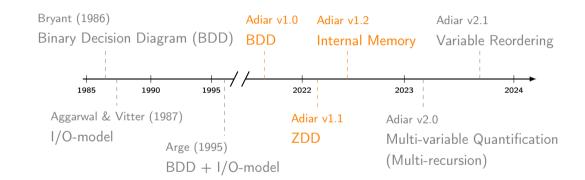


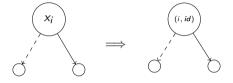
(2) Merge duplicate nodes

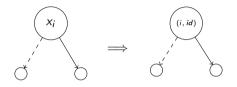
# Arge '95 : BDD + I/O-model



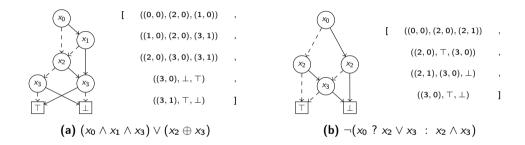
Running time for solving a problem that does not need more than 3  $\,\mathrm{GiB}.$ 



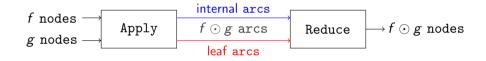


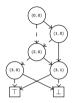


$$(i_1, id_1) < (i_2, id_2) \equiv i_1 < i_2 \lor (i_1 = i_2 \land id_i < id_j)$$



Node-based representation of prior shown BDDs

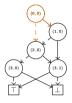




(a)  $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$ 



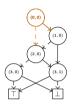
(b) 
$$\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$$



(a)  $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$ 



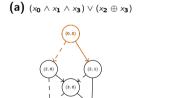
(b)  $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$ 



Priority Queue: Q<sub>app:1</sub>:

- [  $(0,0) \xrightarrow{\top} ((1,0),(2,1))$  ,
  - $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$  ,



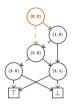


J

**(c)** (a) ∧ (b)

(b)  $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$ 

9



(a)  $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$ 



(b)  $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$ 

Seek: min((1, 0), (2, 1))

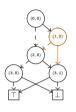
Priority Queue: Qapp:1:

 $[ (0,0) \xrightarrow{\top} ((1,0),(2,1)) ,$  $(0,0) \xrightarrow{\bot} ((2,0),(2,0)) ,$ 

(0,0)

J

**(c)** (a) ∧ (b)



(a)  $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$ 



**(b)**  $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$ 

Seek: min((1,0),(2,1))

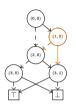
Priority Queue: Q<sub>app:1</sub>:

[ 
$$(0,0) \xrightarrow{\top} ((1,0),(2,1))$$
 ,

$$(0,0) \xrightarrow{\perp} ((2,0),(2,0))$$

(0,0)

J



(a) 
$$(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$$



(b)  $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$ 

# $\begin{array}{c} \text{Seek:} \\ \min((1,0),(2,1)) \end{array}$

Priority Queue: Qapp:1:

 $(0,0)\stackrel{\top}{\longrightarrow} ((1,0),(2,1))$  ,

 $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$ 

 $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$  ,

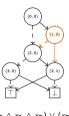
 $(1,0) \xrightarrow{\top} ((3,1),(2,1))$ 

(0,0)

(1,0)

J

**(c)** (a) ∧ (b)



(a) 
$$(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$$



**(b)** 
$$\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$$

Seek: min((1,0),(2,1))

#### Priority Queue: $Q_{app:1}$ :

- $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$
- $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$
- $(1,0) \xrightarrow{\top} ((3,1),(2,1))$  ,

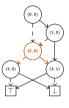
Output:

 $(0,0) \xrightarrow{\top} (1,0)$ 



]

**(c)** (a) ∧ (b)



(a) 
$$(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$$



**(b)** 
$$\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$$

Seek: min((2,0),(2,0))

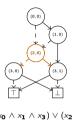
#### Priority Queue: Qapp:1:

- $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$
- $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$
- $(1,0) \xrightarrow{\top} ((3,1),(2,1))$  ,

Output:



]



(a)  $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$ 



(b)  $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$ 

Seek:

 $\min((2,0),(2,0))$ 

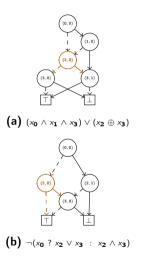
Priority Queue:  $Q_{app:1}$ :

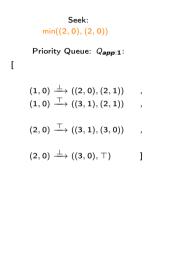
- $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$
- $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$
- $(1,0) \xrightarrow{\top} ((3,1),(2,1))$
- $(2,0) \xrightarrow{\top} ((3,1),(3,0))$
- $(2,0) \xrightarrow{\perp} ((3,0),\top)$

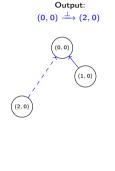
Output:

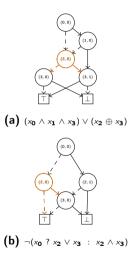


(2,0)

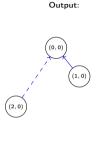


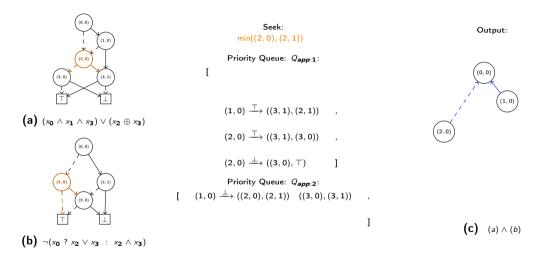


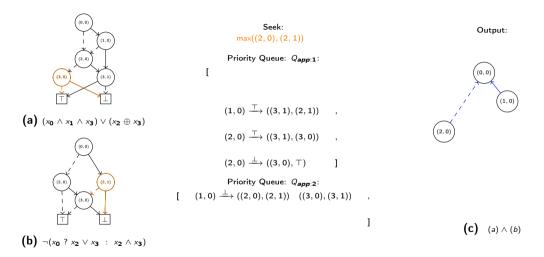


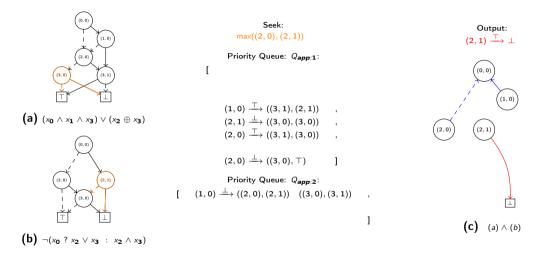


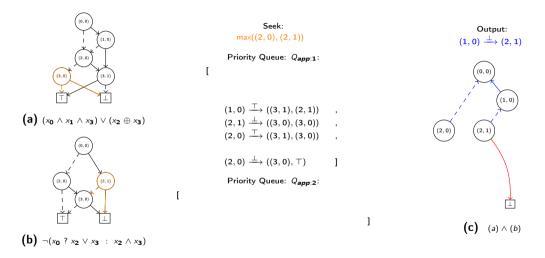
Seek: min((2,0),(2,1))Priority Queue: Qapp:1:  $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$  $(1,0) \xrightarrow{\top} ((3,1),(2,1))$  $(2,0) \xrightarrow{\top} ((3,1),(3,0))$  $(2,0) \xrightarrow{\perp} ((3,0),\top)$ 

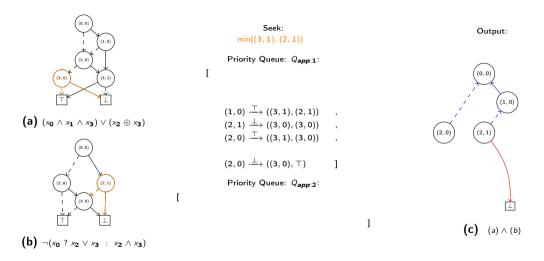


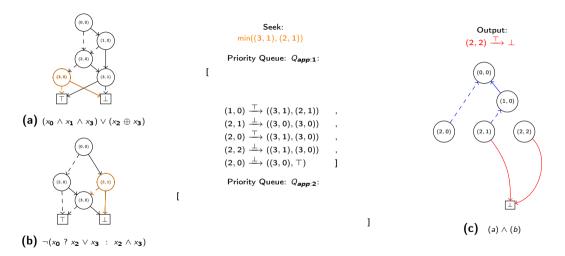


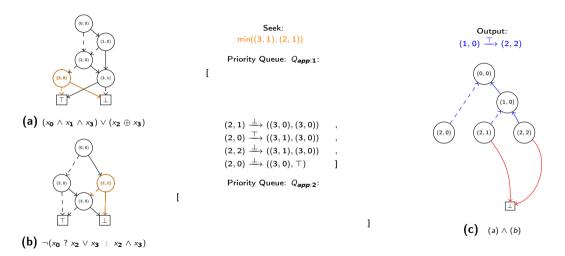


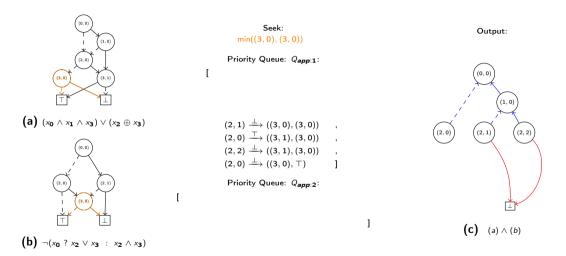


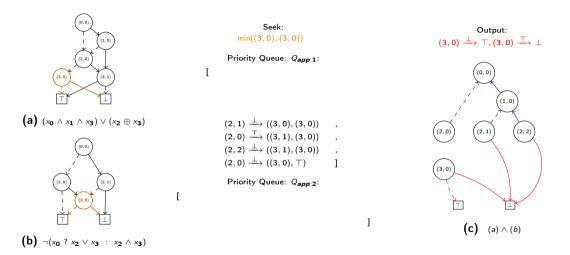


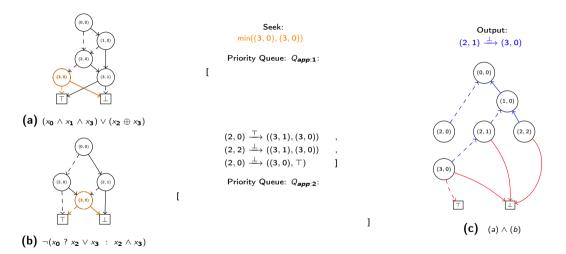


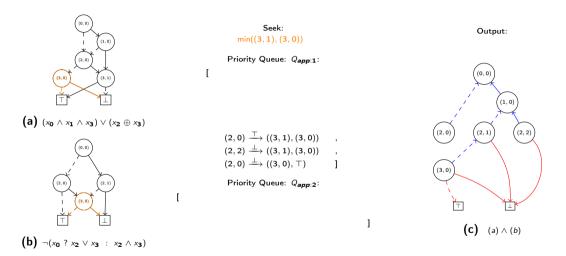


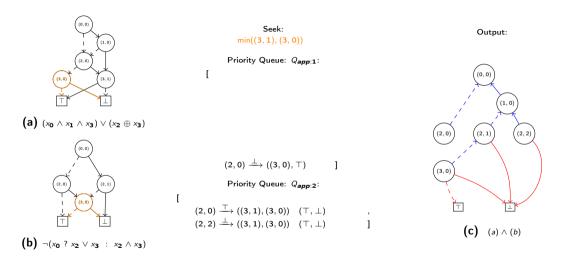


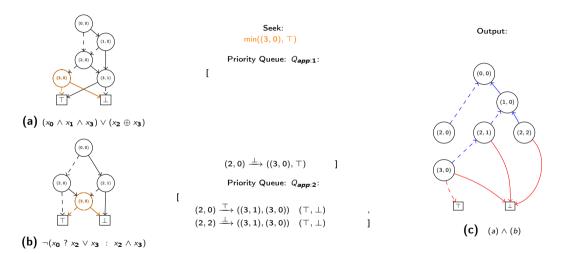


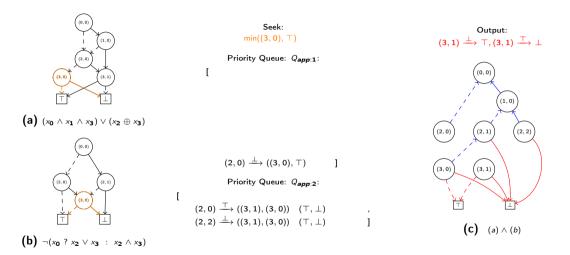


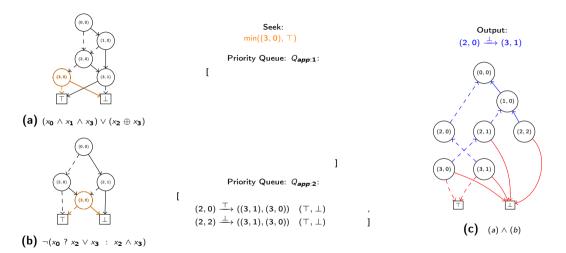


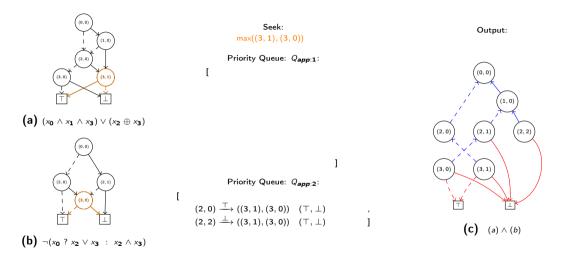


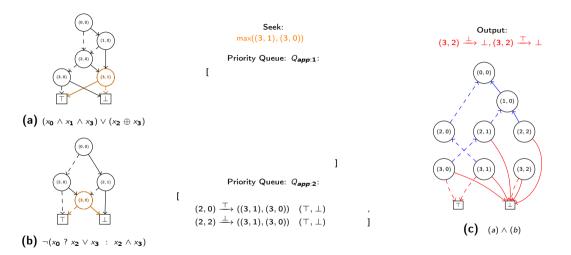


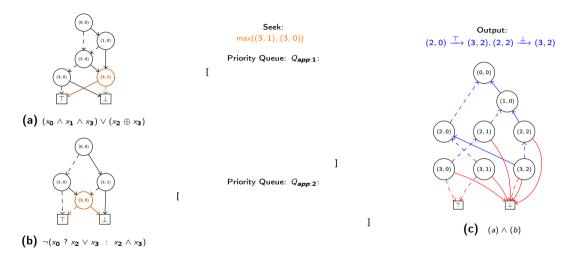


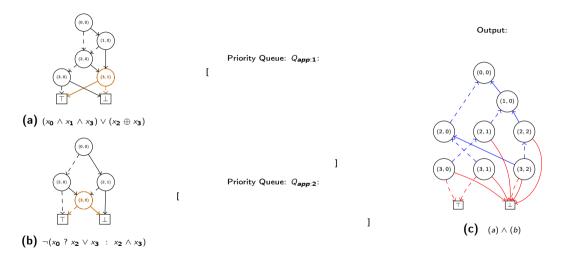


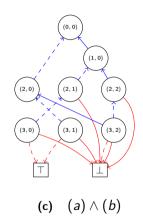


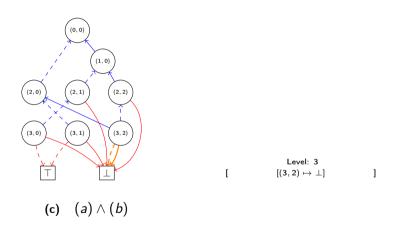


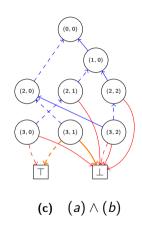


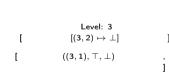


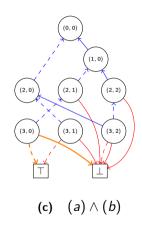


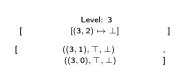


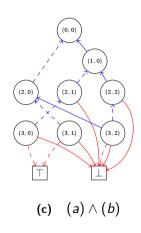


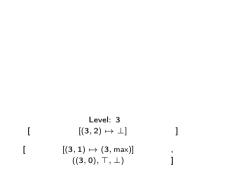




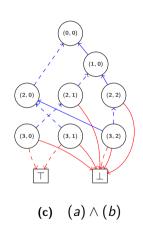


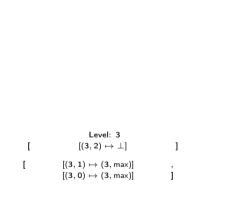


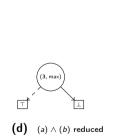


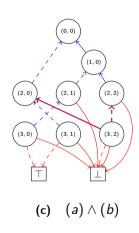


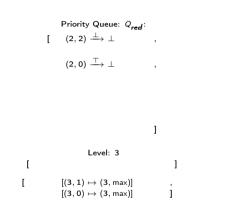


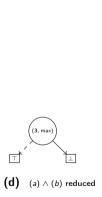


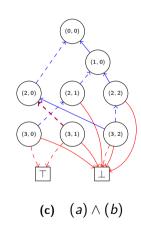


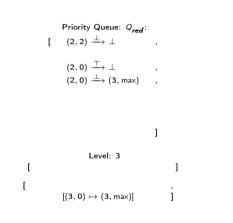


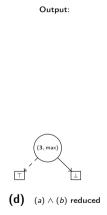


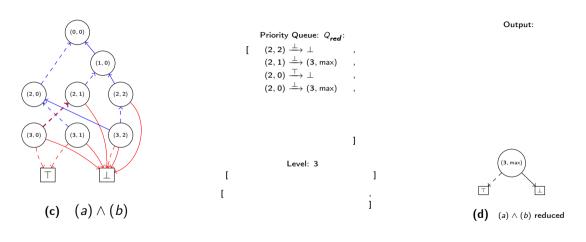


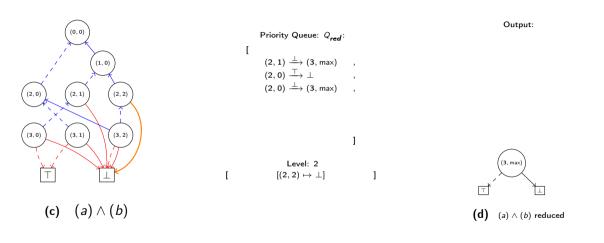


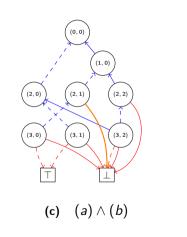


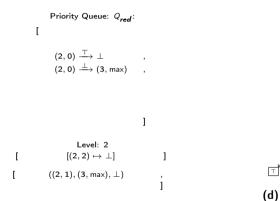




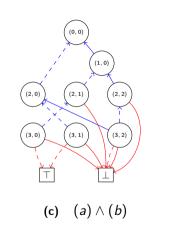


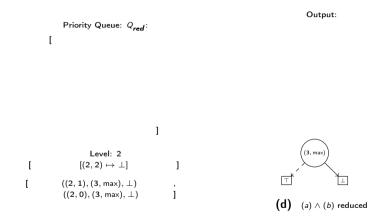


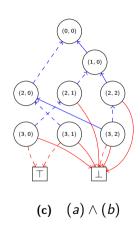


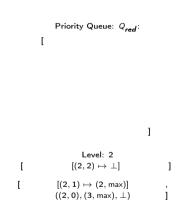


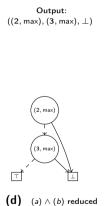


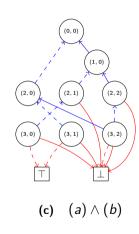


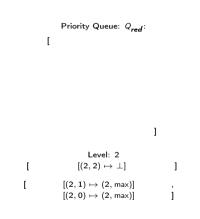


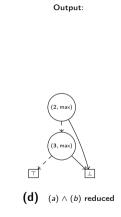


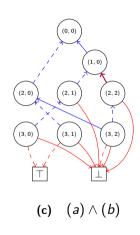


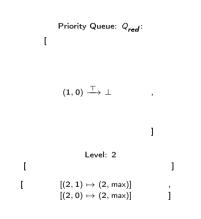


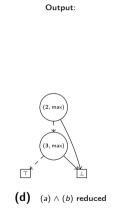


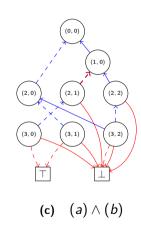


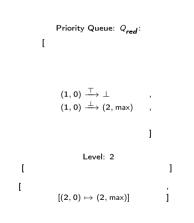


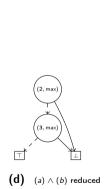


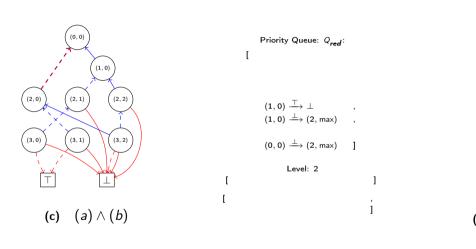


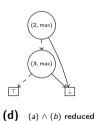


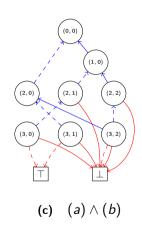


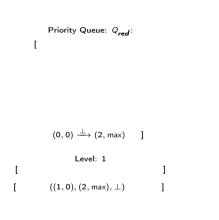


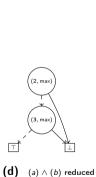


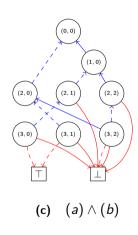


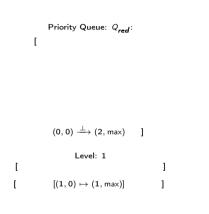


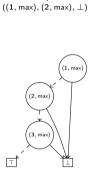




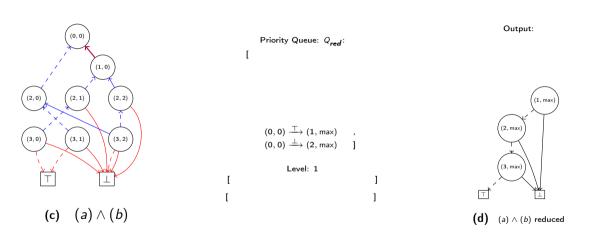


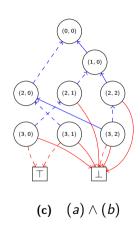


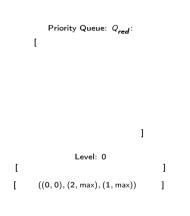


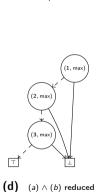


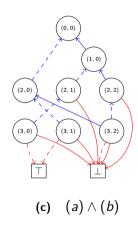
(d)  $(a) \wedge (b)$  reduced

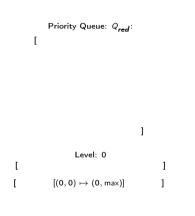


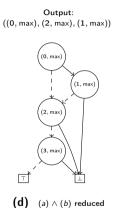


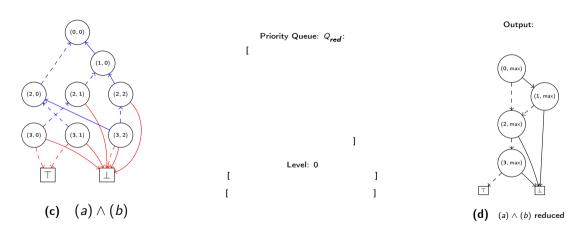


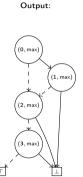


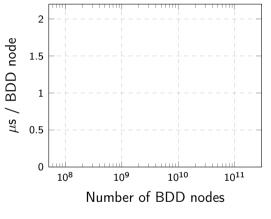






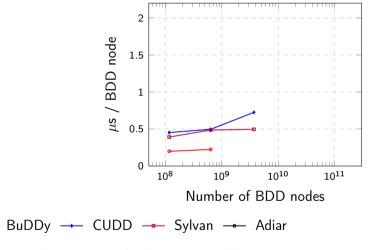




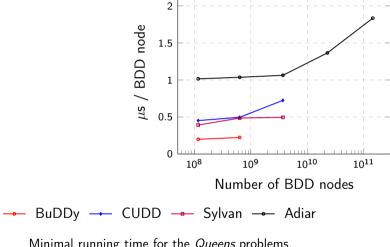


→ BuDDy → CUDD → Sylvan → Adiar

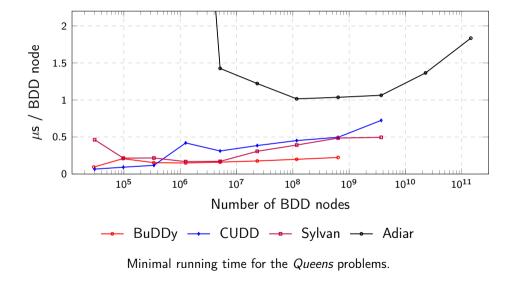
Minimal running time for the *Queens* problems.

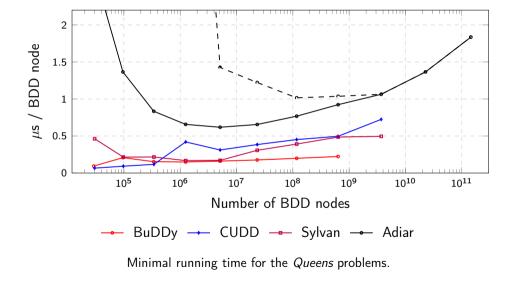


Minimal running time for the *Queens* problems.



Minimal running time for the Queens problems.

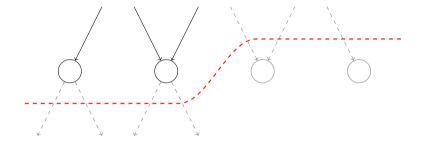


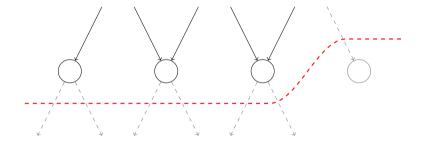


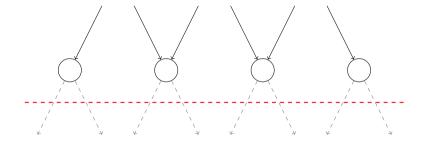




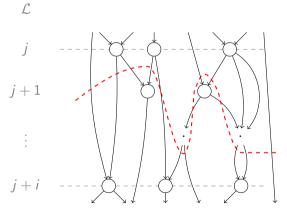








# Definition (i-level cut)



## Definition (i-level cut)

### Lemma

The maximum i-level cut problem is in P for  $i \in \{1, 2\}$ .

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Theorem (Lampis, Kaouri, Mitsou 2011) The maximum i-level cut problem is NP-complete for  $i \ge 4$ .

#### Theorem

The maximum (i-level) cut of a BDD with N internal nodes is N+1.

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For  $i \in \{1,2\}$ , the maximum i-level cut of the (unreduced) output of Apply is upper bounded by the product of the inputs' corresponding i-level cuts.

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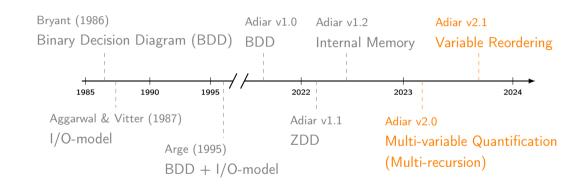
The maximum (i-level) cut of a BDD with N internal nodes is N+1.

#### **Theorem**

For  $i \in \{1,2\}$ , the maximum i-level cut of the (unreduced) output of Apply is upper bounded by the product of the inputs' corresponding i-level cuts.

#### Lemma

The maximum 2-level cut of a BDD is upper bounded by  $\frac{3}{2}$  its 1-level cut.



$$RelProd(S, T) \equiv (\exists \vec{x}. S(\vec{x}) \land T(\vec{x}, \vec{x'}))[\vec{x'}/\vec{x}]$$

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```
1 exists (f, V)
     if f = 1 \lor f = \top
            then f
    else if V \cap \{i \in \mathbb{N} \mid i \geq \mathsf{top}(f)\} = \emptyset
5
            then f
     else if top(f) \notin V
            then Node { top(f), exists(f.low, V), exists(f.high, V) }
8
     else let low = exists(f.low. V)
                  high = exists(f.high, V)
10
            in or(low, high)
```

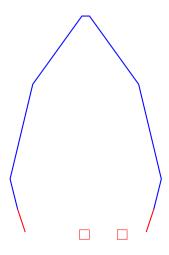
A recursive multi-variable exists operation.

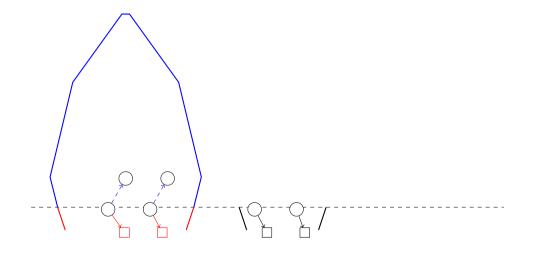
```
1 exists (f, V)
    if f = \bot \lor f = \top
     then f
   else if V \cap \{i \in \mathbb{N} \mid i \geq \mathsf{top}(f)\} = \emptyset
5
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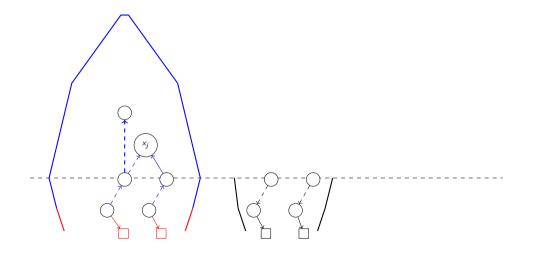
A recursive multi-variable exists operation.

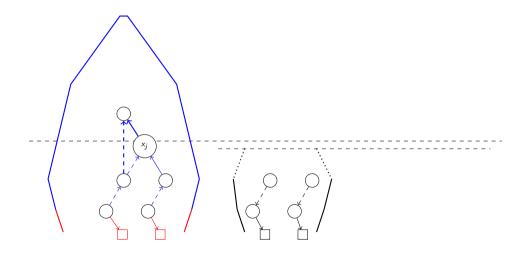
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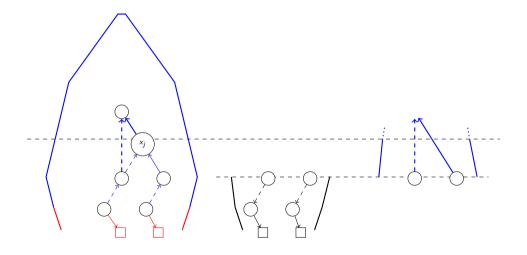
A recursive multi-variable exists operation.

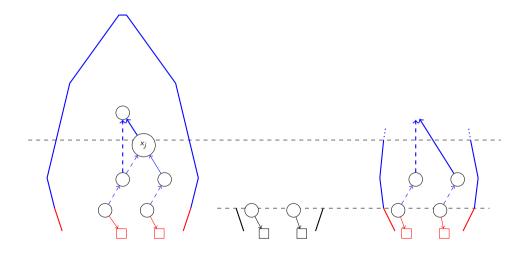


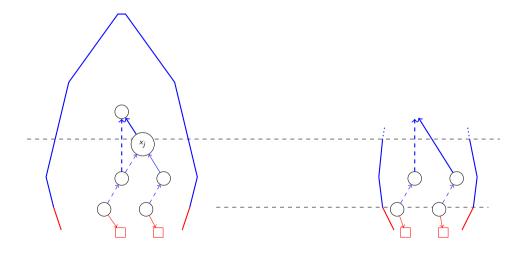


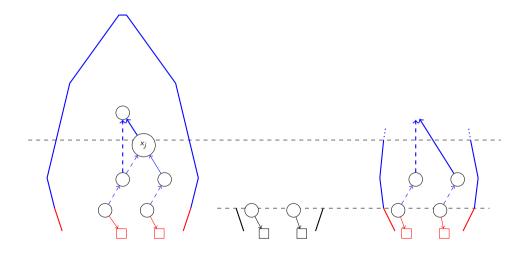


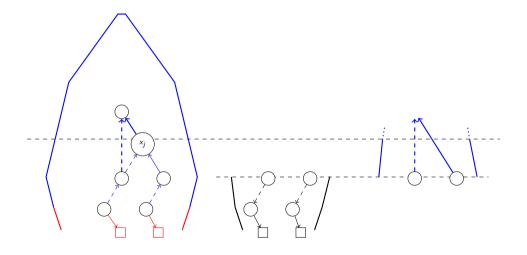


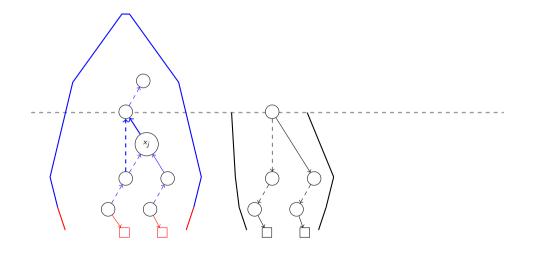


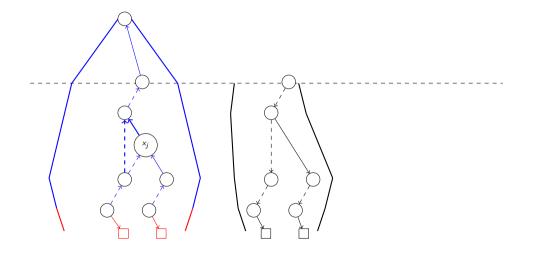


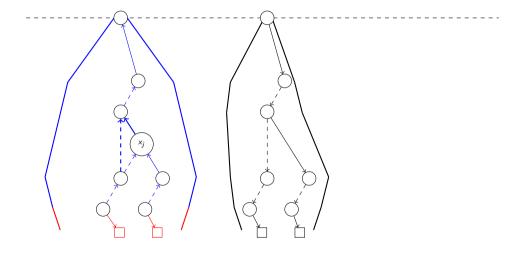




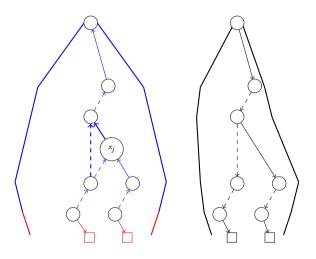






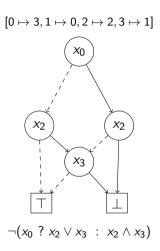


# Adiar v2.0: Multi-variable Quantification

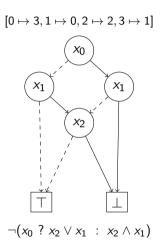


$$RelProd(S, T) \equiv (\exists \vec{x}. S(\vec{x}) \land T(\vec{x}, \vec{x'}))[\vec{x'}/\vec{x}]$$

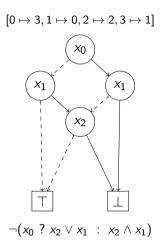
```
1 substitute(f, i_map)
2  let low = substitute(f.low)
3    high = substitute(f.high)
4    i' = i_map[top(f)]
5  in bubble(i', low, high)
```



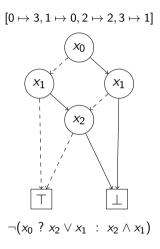
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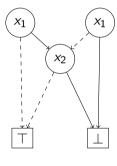
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```
[0\mapsto 3, 1\mapsto 0, 2\mapsto 2, 3\mapsto 1]
```

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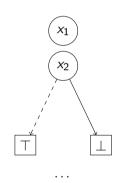
A recursive substitute operation.



. . .

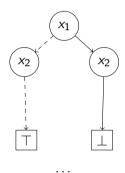
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```



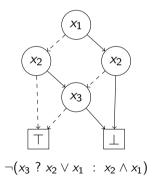
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[0\mapsto 3, 1\mapsto 0, 2\mapsto 2, 3\mapsto 1]
```

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```



```
[0\mapsto 3, 1\mapsto 0, 2\mapsto 2, 3\mapsto 1]
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4  i' = i_map[top(f)]
5  in bubble(i', low, high)
```



```
\frac{\text{Time Space }I/O}{O(NT) O(NT)}
1 substitute (f, i\_map)
2 let low = substitute (f.low)
3 high = substitute (f.high)
4 i' = i\_map[top(f)]
5 in bubble (i', low, high)
```

```
Time
                                                                Space
1 substitute(f, i_map)
    let low = substitute(f.low)
                                                   Complexity of depth-first substitute
          high = substitute(f.high)
          i' = i_map[top(f)]
5
    in bubble(i', low, high)
                                                 Time
                                                               Space
                                             O(NT \log T) \quad O(N+T) \quad O(N \cdot \text{sort}(T))
      A recursive substitute operation.
                                                  Complexity of level-by-level substitute
```

# Problem (Variable Replacement) Given BDD $f_{\pi}$ with variable ordering $\pi$ and remapping of variables $m: \mathbb{N} \to \mathbb{N}$ , construct $f'_{\pi} \equiv f_{\pi}[x/m(x)]$ .

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#### Problem (Static Variable Reordering)

Given BDD  $f_{\pi}$  with variable ordering  $\pi$  and another variable ordering  $\pi'$ , construct  $f_{\pi'} \equiv f_{\pi}$ .

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#### Problem (Static Variable Reordering)

Given BDD  $f_{\pi}$  with variable ordering  $\pi$  and another variable ordering  $\pi'$ , construct  $f_{\pi'} \equiv f_{\pi}$ .

#### Problem (Dynamic Variable Reordering)

Given BDD  $f_{\pi}$  with variable ordering  $\pi$ ,

find  $\pi'$  and construct  $f_{\pi'} \equiv f_{\pi}$  such that  $|f_{\pi'}|$  is minimal.

