

# Efficient Equality Checking for Non-Shared Binary Decision Diagrams

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$$f \leftrightarrow g \equiv \top$$

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$$\underbrace{O(\text{sort}(N^2))}_{\text{Apply}} + \underbrace{O(\text{sort}(N^2))}_{\text{Reduce}} + \underbrace{O(1)}_{\text{check is } \top} = O(\text{sort}(N^2))$$



**Theorem (Bryant '86)**

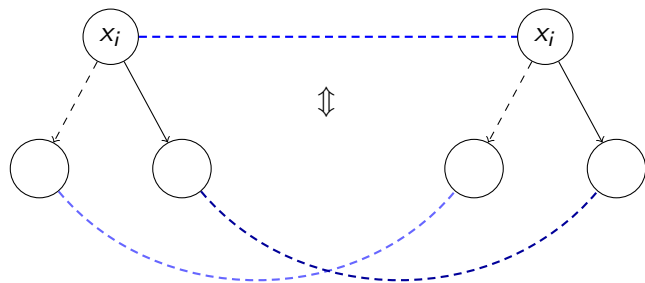
*Let  $\pi$  be a variable order and  $f : \mathbb{B}^n \rightarrow \mathbb{B}$  then there exists a unique (up to isomorphism) Reduced Ordered Binary Decision Diagram representing  $f$  with ordering  $\pi$ .*

### Theorem (Bryant '86)

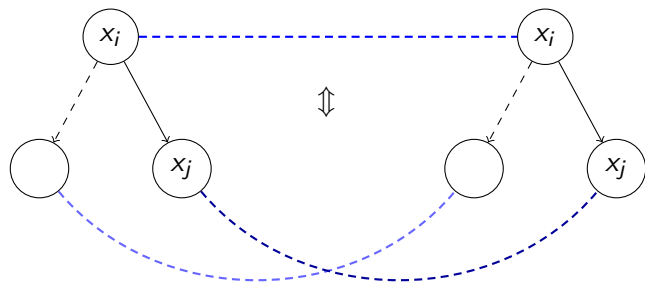
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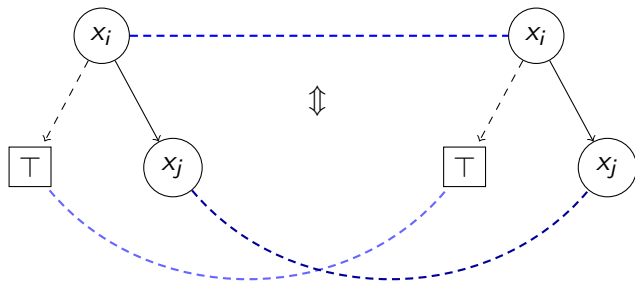
Trivial cases:  $f \not\equiv g$  if there is a mismatch in

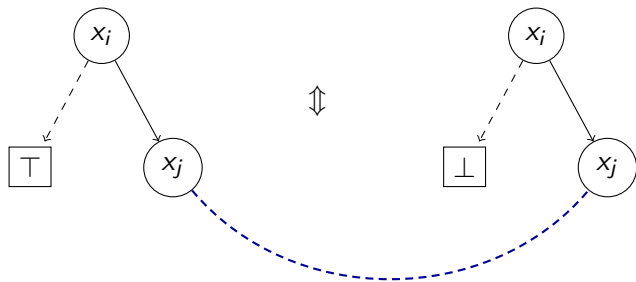
- $N_f \neq N_g$       Number of nodes       $O(1)$  I/Os
- $L_f \neq L_g$       Number of levels       $O(1)$  I/Os
- $N_{f,i} \neq N_{g,i}$       Number of nodes on a level       $O(L/B)$  I/Os
- $L_{f,i} \neq L_{g,i}$       Label of an  $i$ th level       $O(L/B)$  I/Os

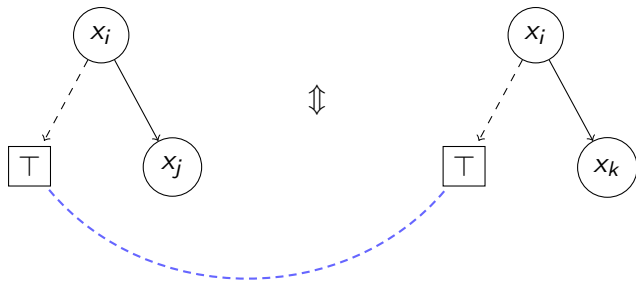


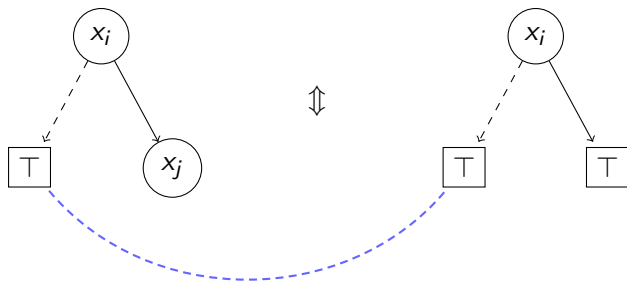












`IsIsomorphic( $f$ ,  $g$ )`

- Check whether root  $v_f$  of  $f$  and root  $v_g$  of  $g$  have a local violation.
- Check  $low(v_f) \sim low(v_g)$  and  $high(v_f) \sim high(v_g)$  “recursively”.

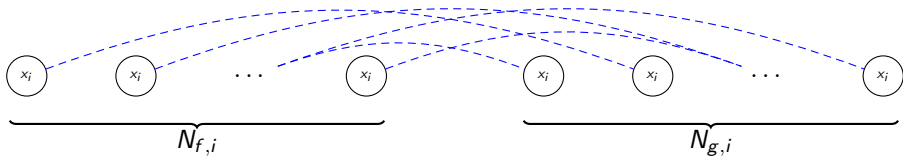
Return `false` on first violation. If there are no violations then return `true`.

IsIsomorphic( $f$ ,  $g$ )

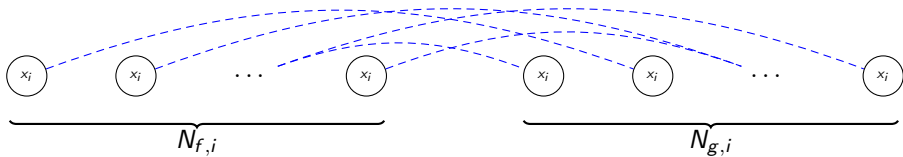
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Return false on first violation. If there are no violations then return true.

$$\underbrace{O(\text{sort}(N^2))}_{\text{Apply'}} + \underbrace{\cancel{O(\text{sort}(N^2))}}_{\text{Reduce}} + \underbrace{\cancel{O(1)}}_{\text{check is T}} = O(\text{sort}(N^2))$$







Return false if more than  $N_{f,i} = N_{g,i}$  pairs of nodes  $(v_f, v_g)$  are checked on level  $i$ .

$$O(\text{sort}(N))$$



## Observation

The output of Reduce has the following properties

- Nodes on level  $i$  are output sorted by their children

$$((i_1, id_1), low_1, high_1) <_{lex(i, low, high)} ((i_2, id_2), low_2, high_2) ,$$

where

$$\forall (i, id) : (i, id) < \perp < \top^1.$$

- Nodes on level  $i$  have their identifiers *consecutively* numbered

$$MAX - N_{f,i} + 1, \dots, MAX - 1, MAX .$$

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<sup>1</sup>Assuming the BDD is not negated. If that is the case then  $(i, id) < \top < \perp$ .

## Theorem

*If  $G_f$  and  $G_g$  are outputs of Reduce.*

$G_f \sim G_g \iff$  For all  $i \in [0; N_f)$  the node  $G_f[i]$  matches  $G_g[i]$  numerically.

## Proof.

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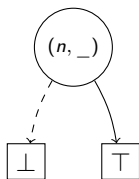
□

## Corollary

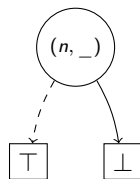
*If  $G_f$  and  $G_g$  are outputs of Reduce then  $f \equiv g$  is computable using  $2 \cdot N/B$  I/Os.<sup>2</sup>*

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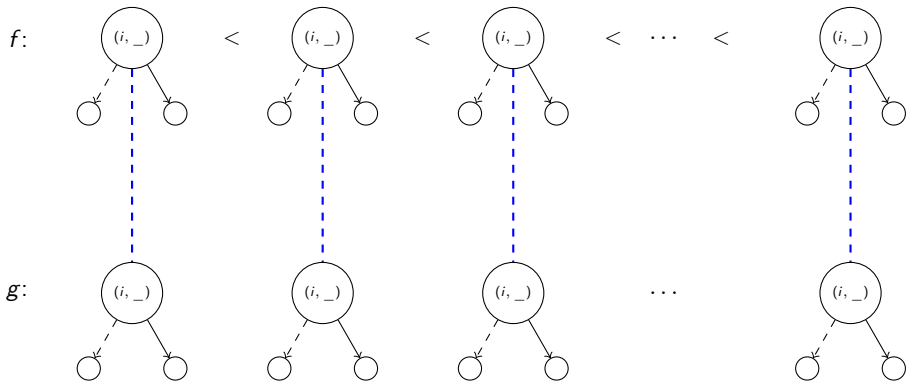
<sup>2</sup>Assuming they are both unnegated (or both negated).



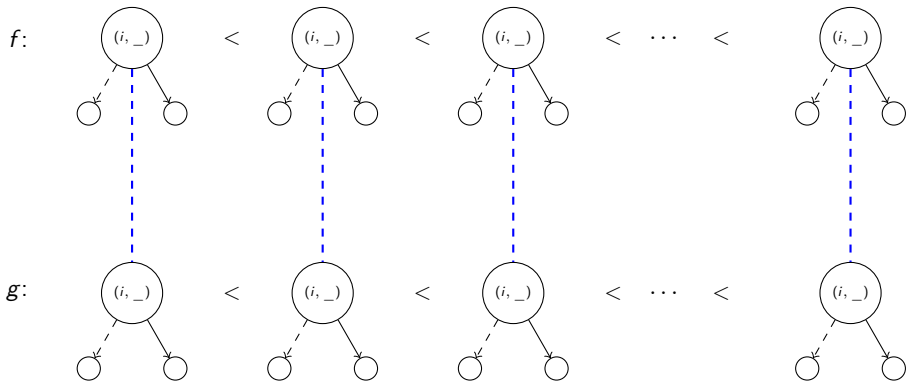
$<$



Base case:  $n$



Induction case:  $i < n$



Induction case:  $i < n$





Algorithm	Time (s)
$f \leftrightarrow g \equiv \top$	0.38
$O(\text{sort}(N))$	0.058
$2 \cdot N/B$	0.006

Checking the (EPFL Benchmark) *voter* circuit's single output gate ( $|N_f| = |N_g| = 5.76$  MiB).