I/O-efficient Manipulation of Binary Decision Diagrams

Steffan Christ Sølvsten

S. C. Sølvsten, J. van de Pol, A. B. Jakobsen, and M. W. B. Thomasen. *Adiar: Binary Decision Diagrams in External Memory.* 2022



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What are Binary Decision Diagrams?

Why do they break?

How can we fix it?

CountPaths

Apply

Equality Checking

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What are Binary Decision Diagrams?

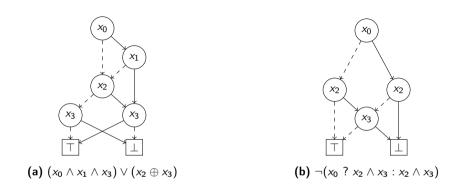
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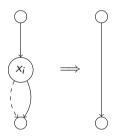
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Examples of (Reduced Ordered) Binary Decision Diagrams.

Theorem (Bryant '86)For a fixed variable order, if one exhaustively applies the two rules below, then one obtains the Reduced OBDD, which is a unique canonical form of the function.



 X_i

(1) Remove redundant nodes

(2) Merge duplicate nodes

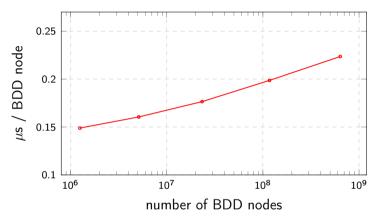
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\begin{array}{ll} \operatorname{bdd\_apply}\left(f,\ g\ ,\ \otimes\right): \\ & \text{if}\ f,g\in\{\bot,\top\} \\ & \text{then}\ f\otimes g \\ & \text{else let}\ i = \operatorname{top}\left(f.(\mathit{var}),\ g.\mathit{var}\right) \\ & t = \operatorname{bdd\_apply}\left(f[x_i := \top],\ g[x_i := \top],\ \otimes\right) \\ & e = \operatorname{bdd\_apply}\left(f[x_i := \bot],\ g[x_i := \bot],\ \otimes\right) \\ & \text{in make\_node}\left(i,t,e\right) \end{array}
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\begin{array}{l} \operatorname{bdd\_apply}\left(f,\ g\ ,\ \otimes\right):\\ & \text{if}\ f,g\in\{\bot,\top\}\\ & \text{then}\ f\otimes g\\ & \text{else let}\ i=\operatorname{top}(f.(\mathit{var}),\ g.\mathit{var})\\ & t=\operatorname{bdd\_apply}\left(f[x_i:=\top],\ g[x_i:=\top],\ \otimes\right)\\ & e=\operatorname{bdd\_apply}\left(f[x_i:=\bot],\ g[x_i:=\bot],\ \otimes\right)\\ & \text{in make\_node}\left(i,t,e\right) \end{array}
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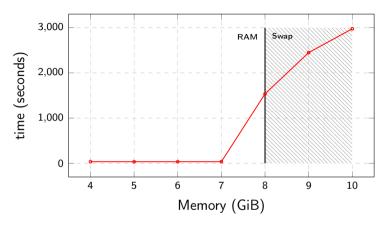
Theorem

bdd_apply runs in $O(N_f \cdot N_g)$ time.

- Memoisation (*Computation Cache*) ensures each recursion is computed only once.
- Reduction Rules can be maintained within make_node(i,t,e) in O(1) time.
 - 1 Redundancy is resolved with an if-statement.
 - 2 Duplication is avoided with a hash table (*Unique Node Table*).



Running time of *BuDDy* for the *N*-Queens problem.



Running time of BuDDy for Tic-Tac-Toe with N=21.

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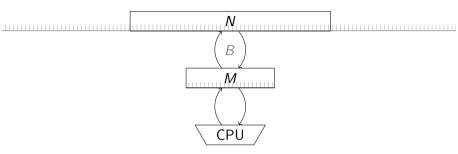
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The I/O model by Aggarwal and Vitter '87

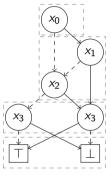
For any realistic values of N, M, and B we have that

$$N/B < \operatorname{sort}(N) \triangleq N/B \cdot \log_{M/B} N/B \ll N$$
,

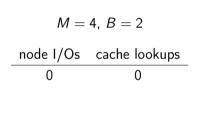
Theorem (Aggarwal and Vitter '87) N elements can be sorted in $\Theta(sort(N))$ I/Os.

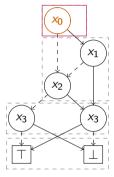
Theorem (Arge '95)

N elements can be inserted in and extracted from a Priority Queue in $\Theta(sort(N))$ I/Os.

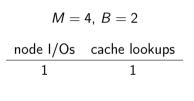


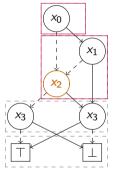
(a)
$$(x_0 \land x_1 \land x_3) \lor (x_2 \oplus x_3)$$



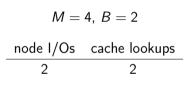


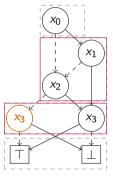
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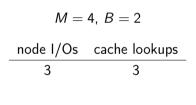


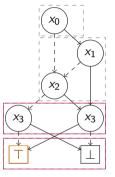
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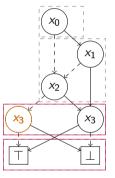
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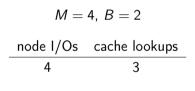


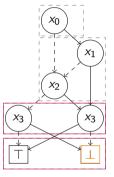
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, $B = 2$
node I/Os cache lookups
4 3

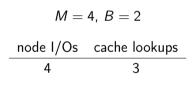


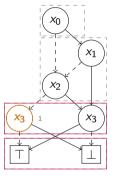
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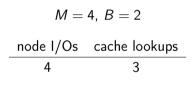


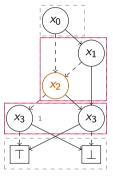
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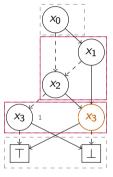
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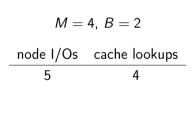


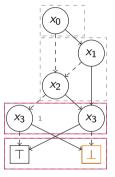
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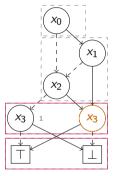
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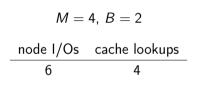


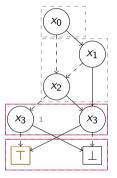
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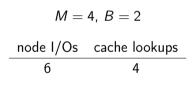


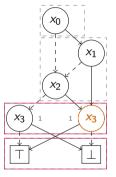
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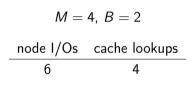


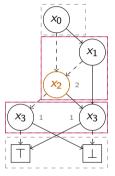
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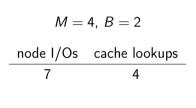


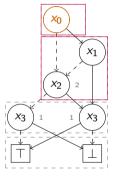
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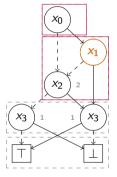
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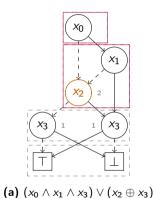
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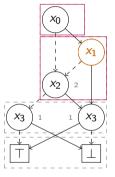


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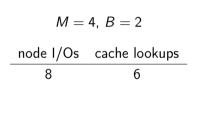
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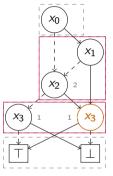


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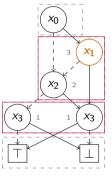


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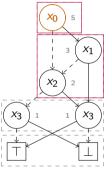




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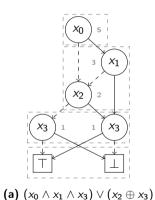


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Algorithm	Time Complexity
bdd_pathcount	$O(N_f)$
bdd_not	$O(N_f)$
bdd_restrict	$O(N_f)$
bdd_apply	$O(N_f \cdot N_g)$
bdd_equal	O(1)

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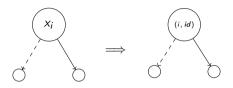
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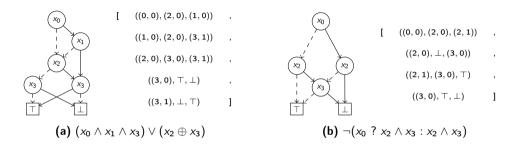
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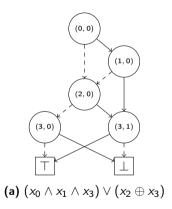
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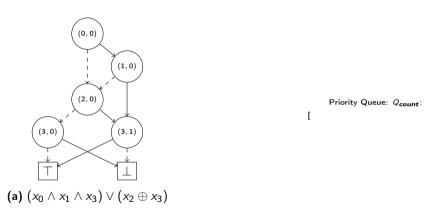


$$(i_1, id_1) < (i_2, id_2) \equiv i_1 < i_2 \lor (i_1 = i_2 \land id_i < id_j)$$

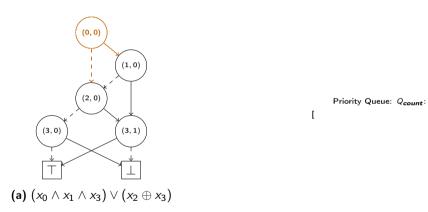


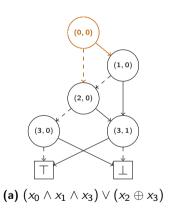
Node-based representation of prior shown BDDs



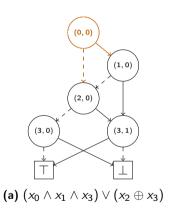


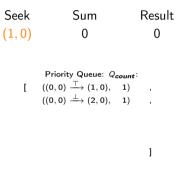
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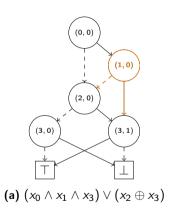


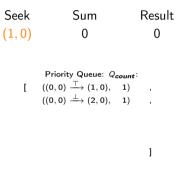


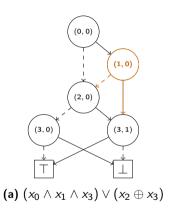
Priority Queue:
$$Q_{count}$$
: [$((0,0) \xrightarrow{\top} (1,0), \quad 1)$, $((0,0) \xrightarrow{\bot} (2,0), \quad 1)$,

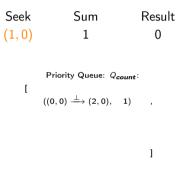


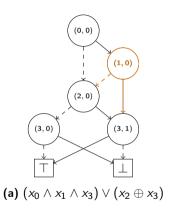


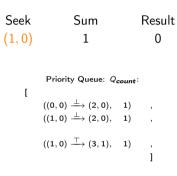


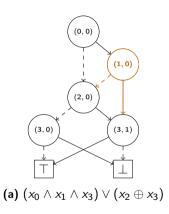


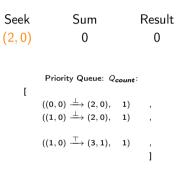


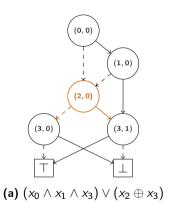


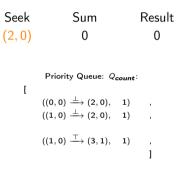


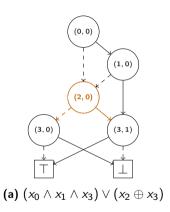


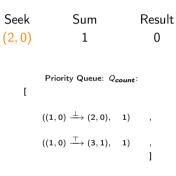


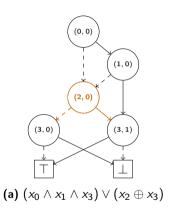


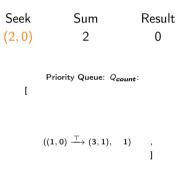


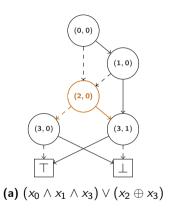


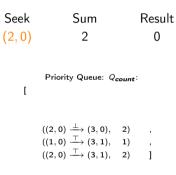


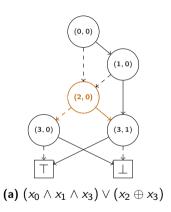


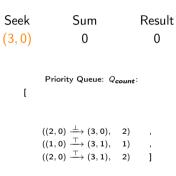


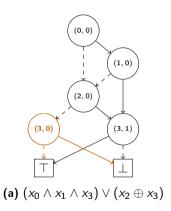


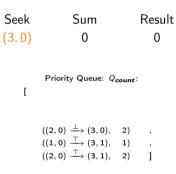


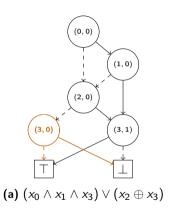


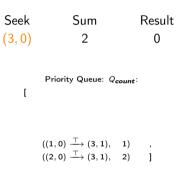


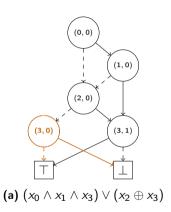


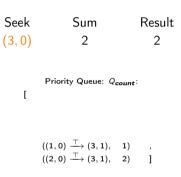


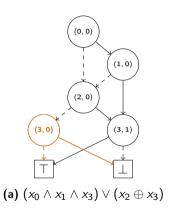


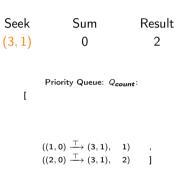


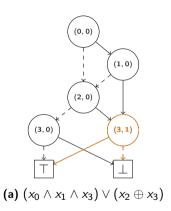


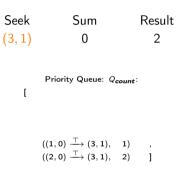


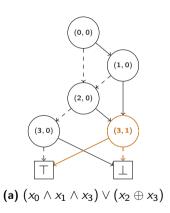


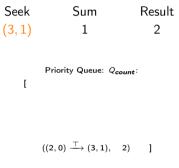


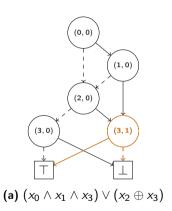


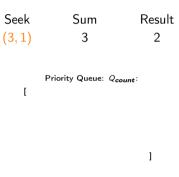


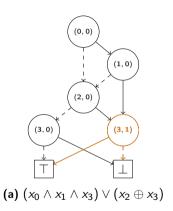


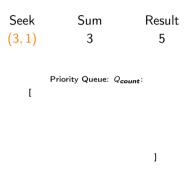


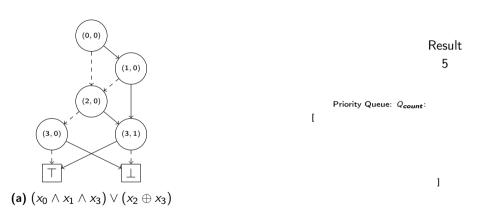












Contents

What are Binary Decision Diagrams?

Why do they break?

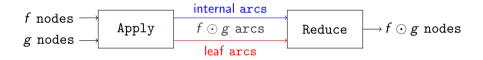
How can we fix it?

CountPaths

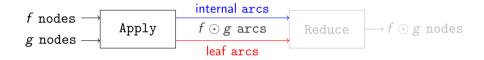
Apply

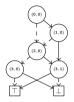
Equality Checking

Apply



Apply



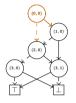


(a) $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$



(b) $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$

(c) $(a) \wedge (b)$

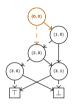


(a) $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$



(b) $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$

(c) $(a) \wedge (b)$



Priority Queue: Qapp:1:

- [$(0,0) \xrightarrow{\top} ((1,0),(2,1))$,
 - $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$

(0,0)

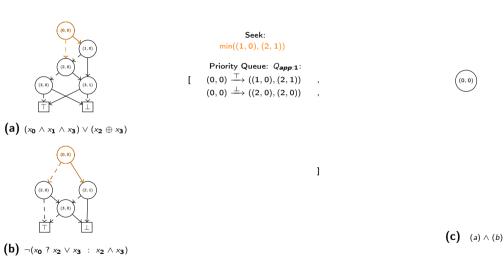
(a) $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$

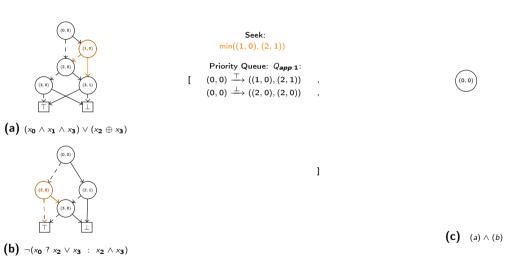


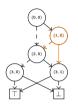
(b) $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$

1

(c) (a) ∧ (b)







(a) $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$



(b) $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$

 $\begin{array}{c} \text{Seek:} \\ \min((1,0),(2,1)) \end{array}$

Priority Queue: Qapp:1:

 $(0,0) \xrightarrow{\top} ((1,0),(2,1))$

 $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$

 $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$

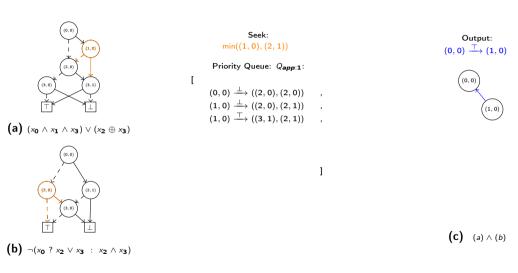
 $(1,0) \xrightarrow{\top} ((3,1),(2,1))$

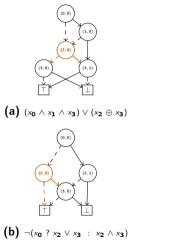
(0,0)

(1,0)

J

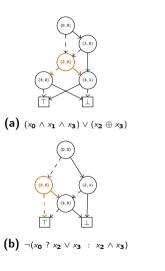
(c) $(a) \wedge (b)$





Seek: min((2,0),(2,0))Priority Queue: Qapp:1: $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$ $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$ $(1,0) \xrightarrow{\top} ((3,1),(2,1))$

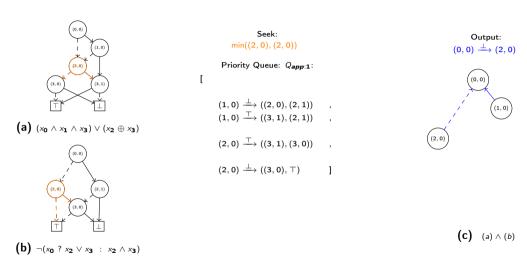
Output: (c) $(a) \wedge (b)$

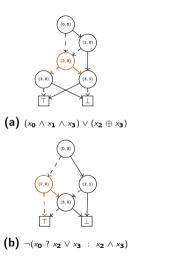


Seek: min((2,0),(2,0))Priority Queue: Qapp:1: $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$ $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$ $(1,0) \xrightarrow{\top} ((3,1),(2,1))$ $(2,0) \xrightarrow{\top} ((3,1),(3,0))$ $(2,0) \xrightarrow{\perp} ((3,0),\top)$]

Output: (2,0)

(c) $(a) \wedge (b)$

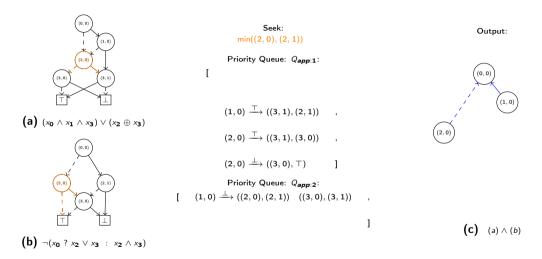


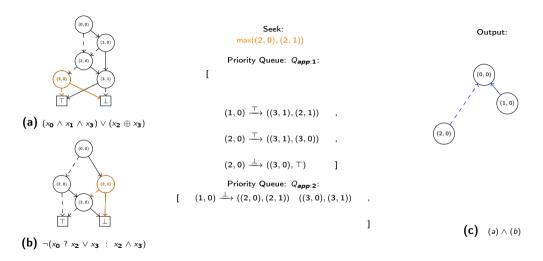


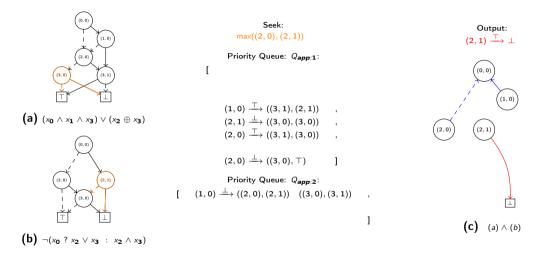
Seek: min((2,0),(2,1))Priority Queue: Qapp:1: $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$ $(1,0) \xrightarrow{\top} ((3,1),(2,1))$ $(2,0) \xrightarrow{\top} ((3,1),(3,0))$ $(2,0) \xrightarrow{\perp} ((3,0),\top)$

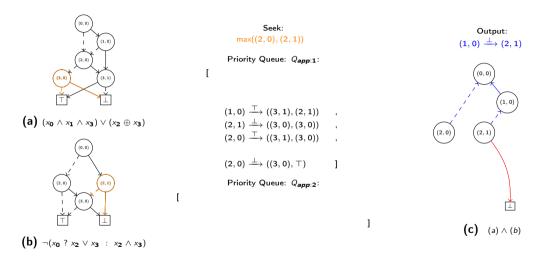
Output: (0,0) (1,0)

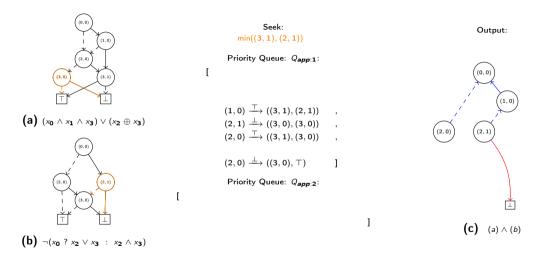
(c) $(a) \wedge (b)$

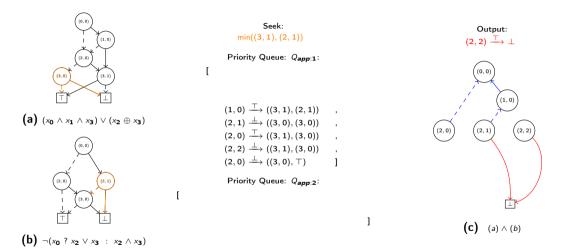


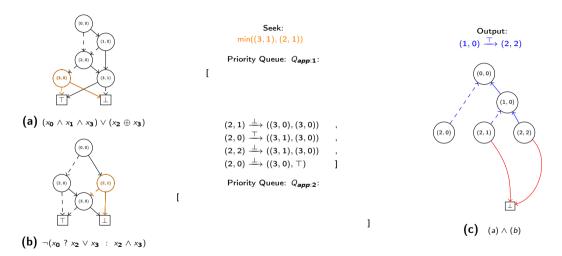


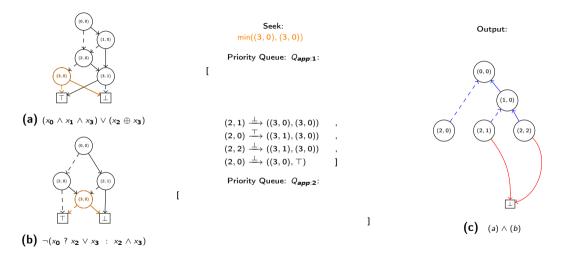


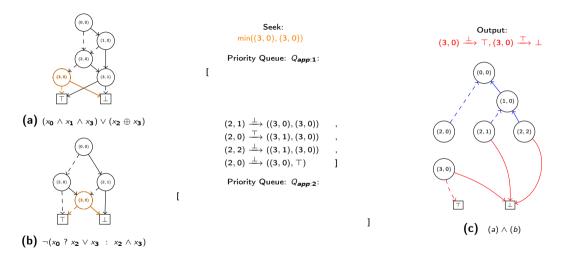


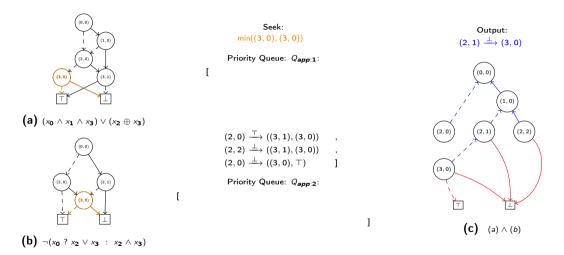


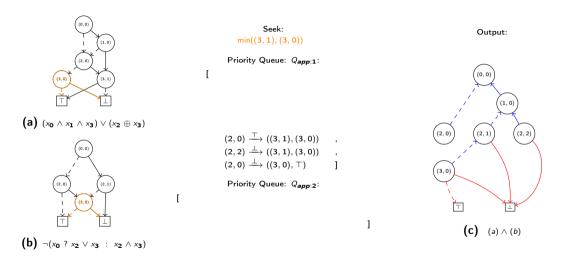


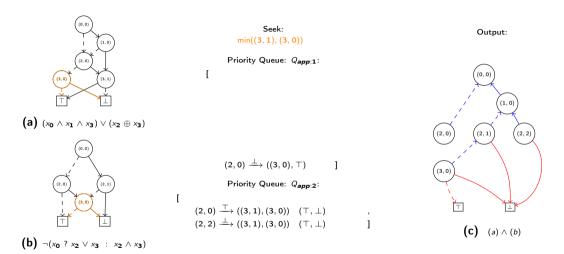


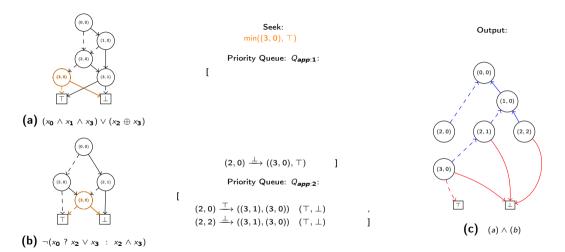


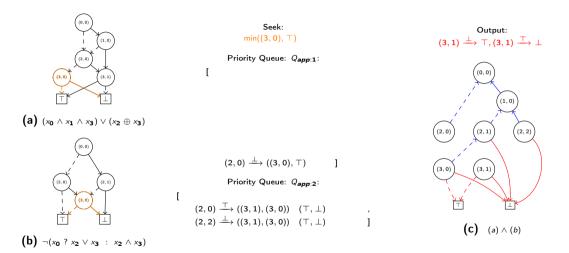


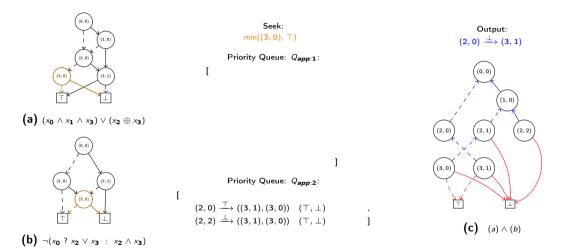


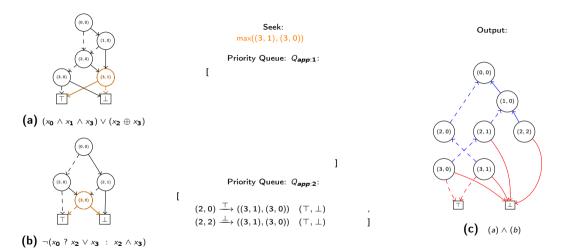


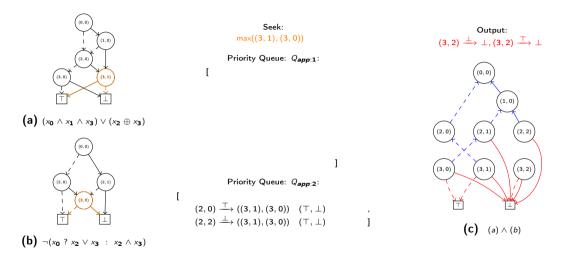


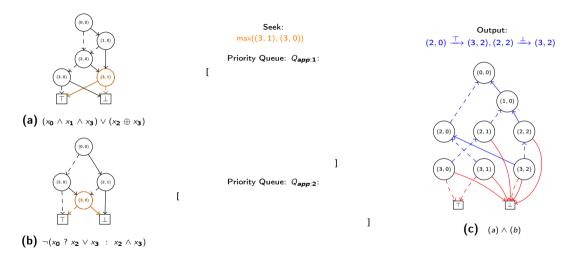


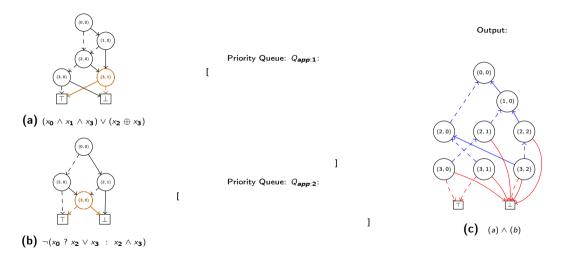




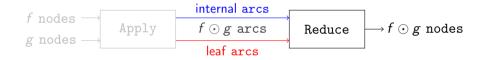


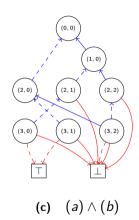


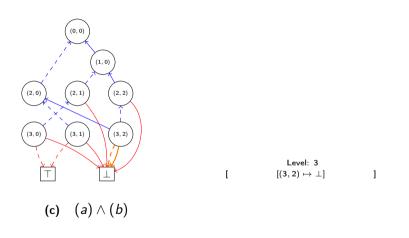


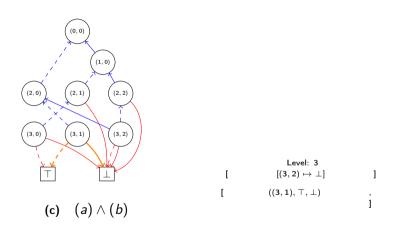


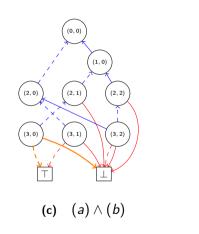
Apply

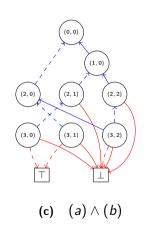


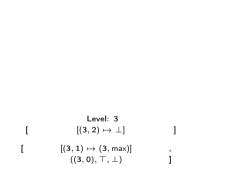




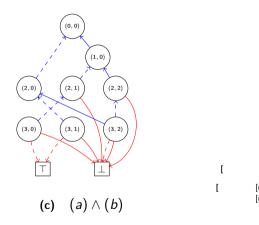


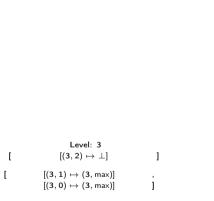


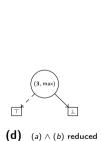




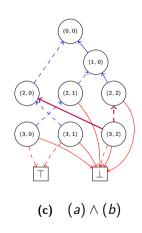


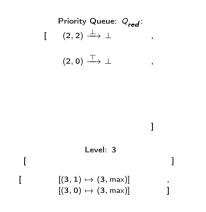






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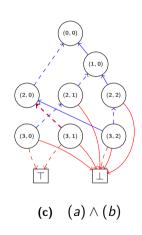


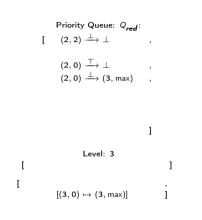




(d) $(a) \wedge (b)$ reduced

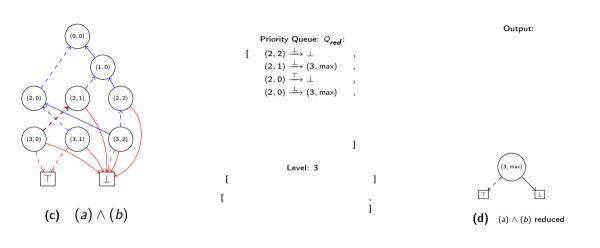
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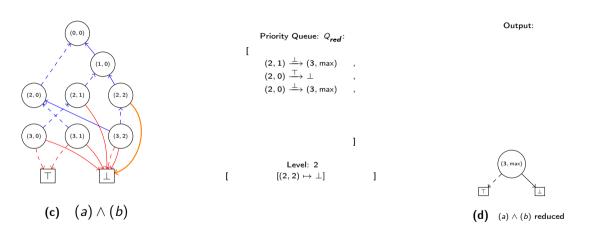


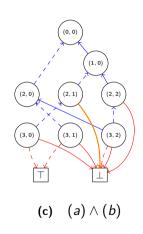


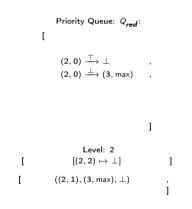




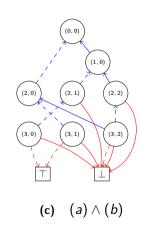


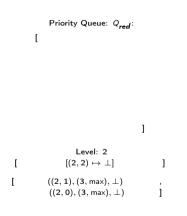




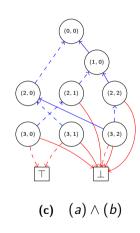


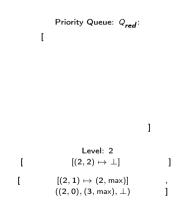


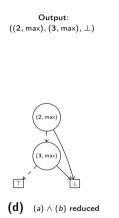


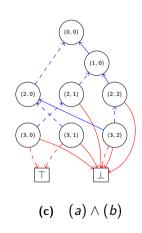


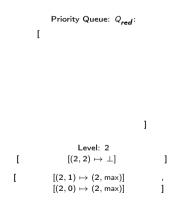


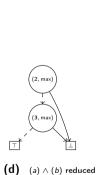


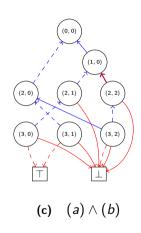


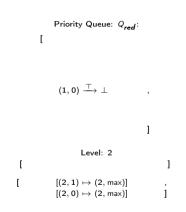


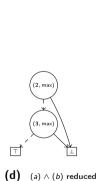


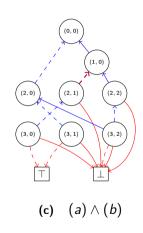


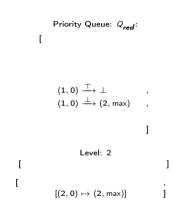


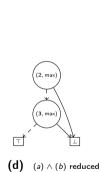


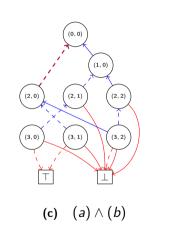


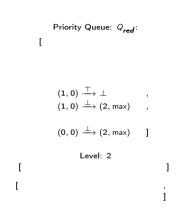


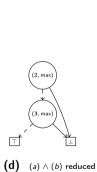


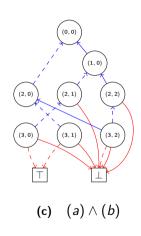


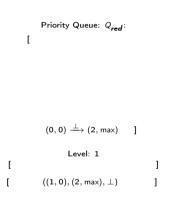


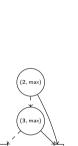




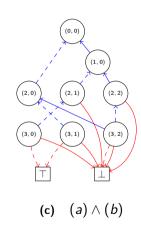


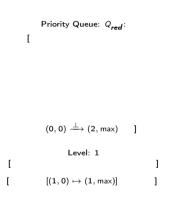


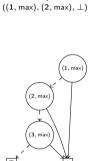


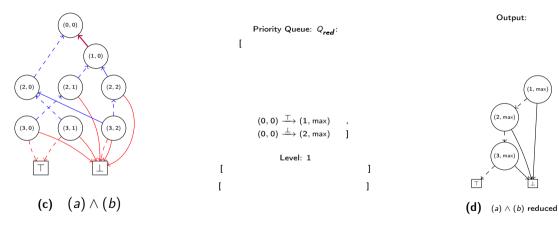


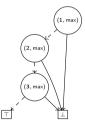
(d) $(a) \wedge (b)$ reduced

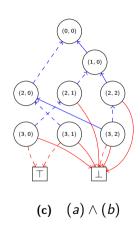


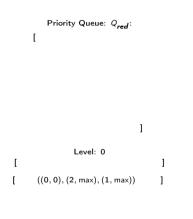


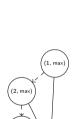


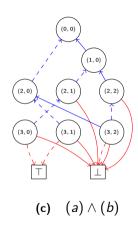


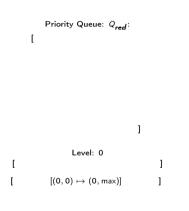


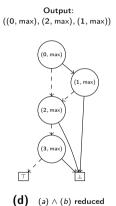


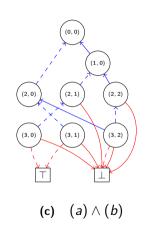


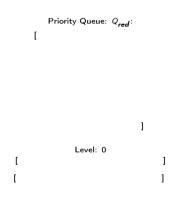


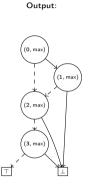












(d) $(a) \wedge (b)$ reduced

Algorithm	I/O-Complexity
bdd_pathcount	$O(\operatorname{sort}(N_f))$
bdd_not	$O(N_f/B)$
bdd_restrict	$O(sort(N_f))$
bdd_apply	$O(sort(\mathit{N_f}\cdot\mathit{N_g}))$

Algorithm	I/O-Complexity
bdd_pathcount	$O(\operatorname{sort}(N_f))$
bdd_not	$O(N_f/B)$
bdd_restrict	$O(\operatorname{sort}(N_f))$
bdd_apply	$O(\operatorname{sort}(N_f \cdot N_g))$
bdd_equal	?

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How can we fix it?

CountPaths

Apply

Equality Checking

$$f\leftrightarrow g\equiv \top$$

$$f\leftrightarrow g\equiv \top$$

$$\underbrace{\textit{O}(\mathsf{sort}(\textit{N}^2))}_{\texttt{Apply}} + \underbrace{\textit{O}(\mathsf{sort}(\textit{N}^2))}_{\texttt{Reduce}} + \underbrace{\textit{O}(1))}_{\texttt{check is } \top} = \textit{O}(\mathsf{sort}(\textit{N}^2))$$

Theorem (Bryant '86)

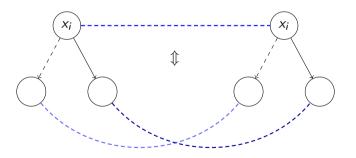
Theorem (Bryant '86)

Let π be a variable order and $f: \mathbb{B}^n \to \mathbb{B}$ then there exists a unique (up to isomorphism) Reduced Ordered Binary Decision Diagram representing f with ordering π .

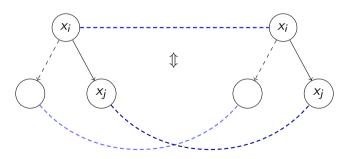
Trivial cases: $f \not\equiv g$ if there is a mismatch in

	$N_f eq N_g$	Number of nodes	<i>O</i> (1) I/Os
•	$L_f eq L_g$	Number of levels	<i>O</i> (1) I/Os
•	$N_{f,i} \neq N_{g,i}$	Number of nodes on a level	O(L/B) I/Os
•	$L_{f,i} \neq L_{g,i}$	Label of an <i>i</i> th level	O(L/B) I/Os

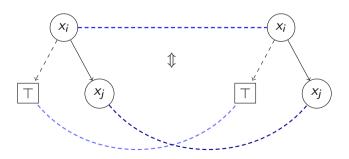
Theorem (Bryant '86)



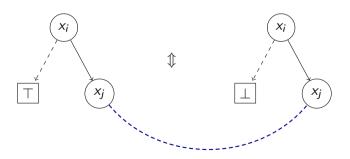
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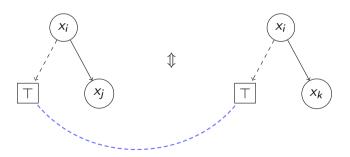
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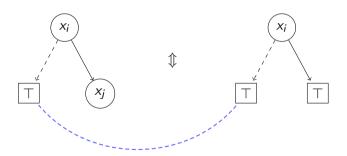
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Let π be a variable order and $f: \mathbb{B}^n \to \mathbb{B}$ then there exists a unique (up to isomorphism) Reduced Ordered Binary Decision Diagram representing f with ordering π .

Is Isomorphic (f, g)

- Check whether root v_f of f and root v_g of g have a local violation.
- Check $low(v_f) \sim low(v_g)$ and $high(v_f) \sim high(v_g)$ "recursively".

Return false on first violation. If there are no violations then return true.

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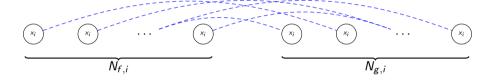
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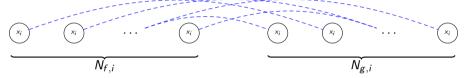
$$\underbrace{O(\mathsf{sort}(\mathit{N}^2))}_{\mathtt{Apply'}} + \underbrace{O(\mathsf{sort}(\mathit{N}^2))}_{\mathtt{Reduce}} + \underbrace{O(1))}_{\mathtt{check is }\top} = O(\mathsf{sort}(\mathit{N}^2))$$

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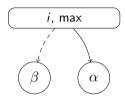
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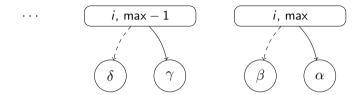
Return false if more than $N_{f,i} = N_{g,i}$ pairs of nodes are checked on level i.

$$\underbrace{O(\mathsf{sort}(\Sigma_i \ N_{f,i}))}_{\mathsf{Apply''}} = O(\mathsf{sort}(N))$$

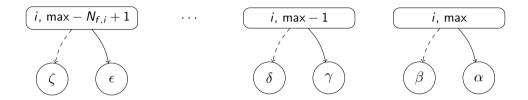
Observation



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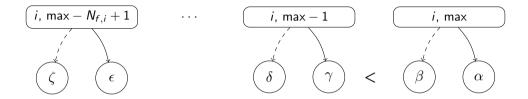
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Each level output by the Reduce algorithm has the following properties:

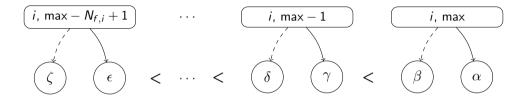
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- Nodes on level *i* have their identifiers *consecutively* numbered.
- Nodes on level *i* are output sorted by their children.

Theorem

If G_f and G_g are outputs of Reduce.

 $G_f \sim G_g \iff For \ all \ i \in [0; N_f) \ the \ node \ G_f[i] \ matches \ G_g[i] \ numerically.$

Proof.

← : Must describe the exact same graph.

 \Rightarrow : Strong induction on BDD levels bottom-up.

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Corollary

If G_f and G_g are outputs of Reduce then $f \equiv g$ is computable using $2 \cdot N/B$ I/Os.

Contents

What are Binary Decision Diagrams?

Why do they break?

How can we fix it?

CountPaths

Apply

Equality Checking $O(N_f)$ $O(N_f)$ $O(N_f)$ $O(Sort(N_f))$ $O(N_f \cdot N_g)$ $O(Sort(N_f))$