I/O-efficient Symbolic Model Checking

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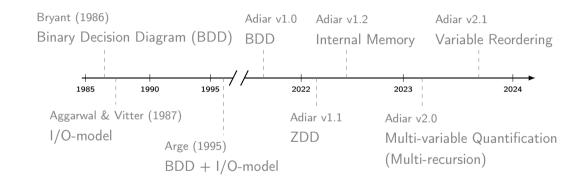
31st of August, 2022

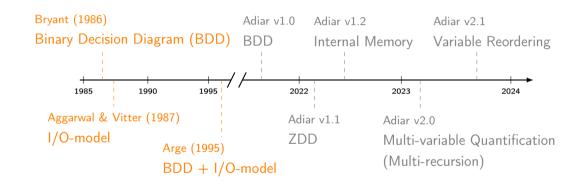


Adiar

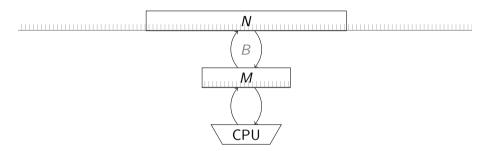
I/O-efficient Decision Diagrams

github.com/ssoelvsten/adiar





Aggarwal and Vitter '87: I/O-model



The I/O-model by Aggarwal and Vitter '87

Aggarwal and Vitter '87: I/O-model

For any realistic values of N, M, and B we have that

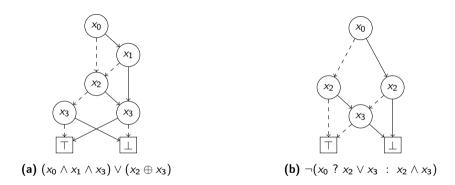
$$N/B < \operatorname{sort}(N) \triangleq N/B \cdot \log_{M/B} N/B \ll N$$
,

Theorem (Aggarwal and Vitter '87) N elements can be sorted in $\Theta(sort(N))$ I/Os.

Theorem (Arge '95)

N elements can be inserted in and extracted from a Priority Queue in $\Theta(sort(N))$ I/Os.

Bryant '86: Binary Decision Diagram

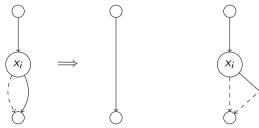


Examples of (Reduced Ordered) Binary Decision Diagrams.

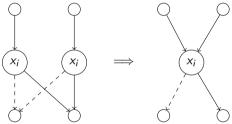
Bryant '86: Binary Decision Diagram

Theorem

For a fixed variable order, if one exhaustively applies the two rules below, then one obtains the Reduced OBDD, which is a unique canonical form of the function.

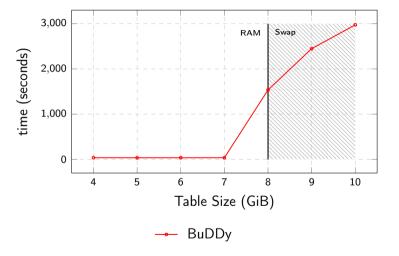




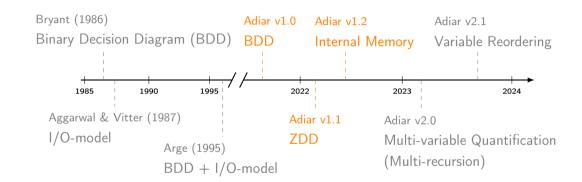


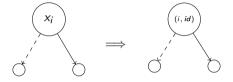
(2) Merge duplicate nodes

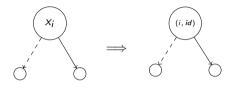
Arge '95 : BDD + I/O-model



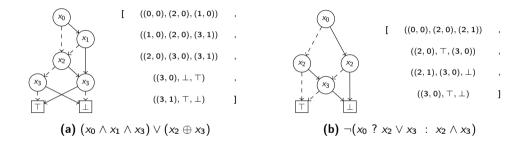
Running time for solving a problem that does not need more than 3 GiB.



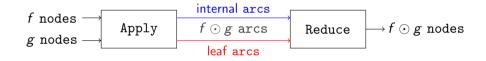


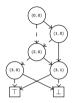


$$(i_1, id_1) < (i_2, id_2) \equiv i_1 < i_2 \lor (i_1 = i_2 \land id_i < id_j)$$



Node-based representation of prior shown BDDs

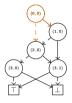




(a) $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$



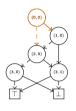
(b)
$$\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$$



(a) $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$



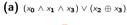
(b) $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$



Priority Queue: *Q_{app:}*1:

- [$(0,0) \xrightarrow{\top} ((1,0),(2,1))$,
 - $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$,

(0,0)

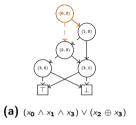




(b)
$$\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$$

J

(c)
$$(a) \wedge (b)$$



(a)
$$(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$$



(b)
$$\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$$

Seek: min((1,0),(2,1))

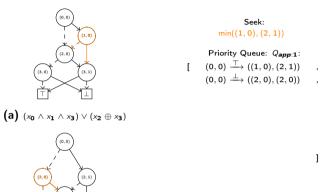
Priority Queue: Qapp:1:

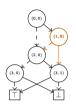
[
$$(0,0) \xrightarrow{\top} ((1,0),(2,1))$$
 ,

$$(0,0) \xrightarrow{\perp} ((2,0),(2,0)) \qquad ,$$

(c) (a) ∧ (b)

(b) $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$





(a)
$$(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$$



(b) $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$

Seek: min((1,0),(2,1))

Priority Queue: Qapp:1:

 $(0,0) \xrightarrow{\top} ((1,0),(2,1))$

 $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$

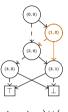
 $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$,

 $(1,0) \xrightarrow{\top} ((3,1),(2,1))$

(0,0)

(1,0)

J



(a)
$$(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$$



(b)
$$\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$$

Seek: min((1, 0), (2, 1))

Priority Queue: $Q_{app:1}$:

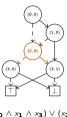
- $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$
- $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$
- $(1,0) \xrightarrow{\top} ((3,1),(2,1))$

Output:

 $(0,0) \xrightarrow{\top} (1,0)$



]



(a)
$$(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$$



(b) $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$

 $\begin{array}{c} \text{Seek:} \\ \min((2,0),(2,0)) \end{array}$

Priority Queue: Qapp:1:

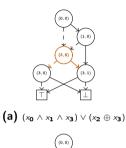
- $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$
- $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$
- $(1,0) \xrightarrow{\top} ((3,1),(2,1))$

Output:



]

(c) (a) ∧ (b)





(b)
$$\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$$

Seek: min((2, 0), (2, 0))

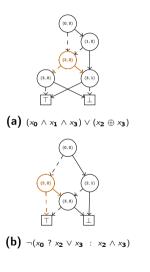
Priority Queue: $Q_{app:1}$:

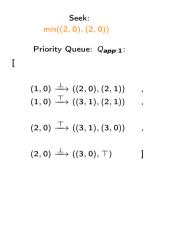
- $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$
- $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$
- $(1,0) \xrightarrow{\top} ((3,1),(2,1))$
- $(2,0) \xrightarrow{\top} ((3,1),(3,0))$
- $(2,0) \xrightarrow{\perp} ((3,0),\top)$]

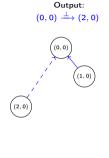
Output:

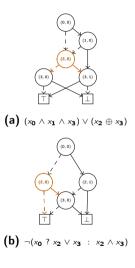


(2,0)





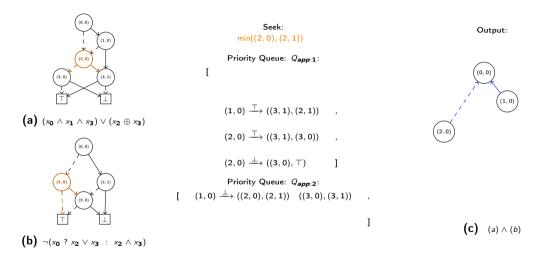


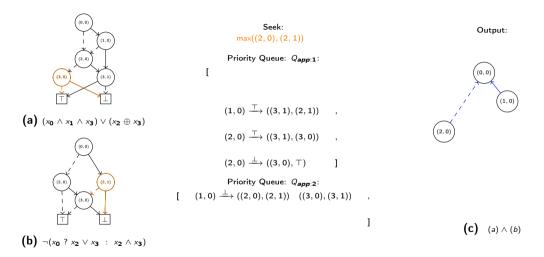


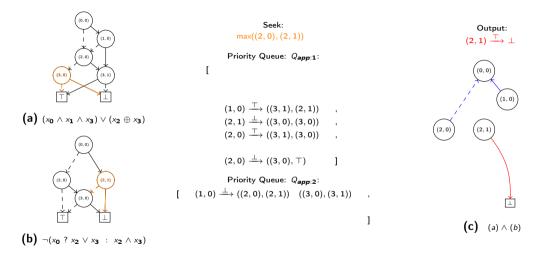
Seek: min((2,0),(2,1))Priority Queue: Qapp:1: $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$ $(1,0) \xrightarrow{\top} ((3,1),(2,1))$ $(2,0) \xrightarrow{\top} ((3,1),(3,0))$ $(2,0) \xrightarrow{\perp} ((3,0),\top)$

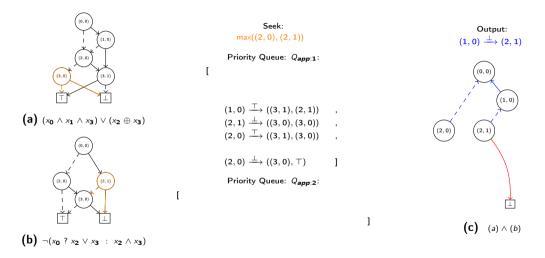
(2,0)

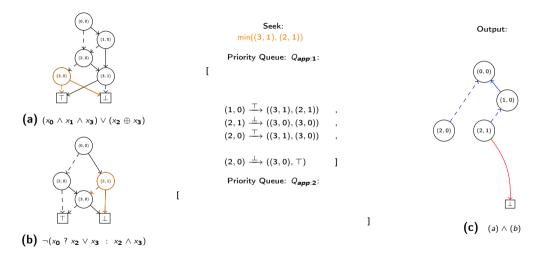
Output:

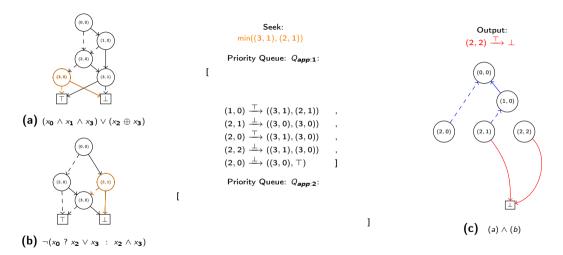


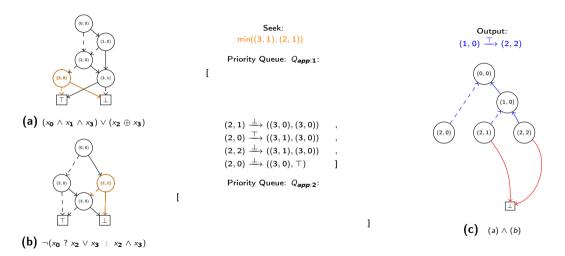


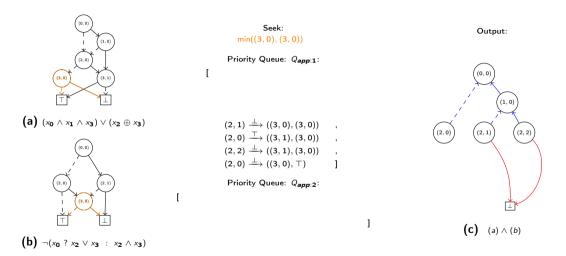


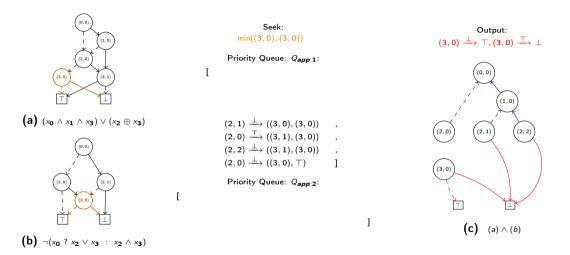


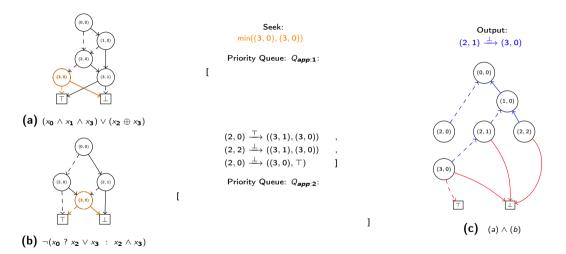


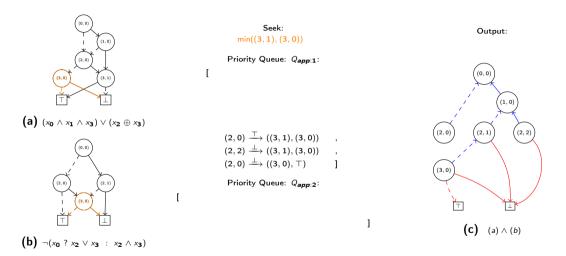


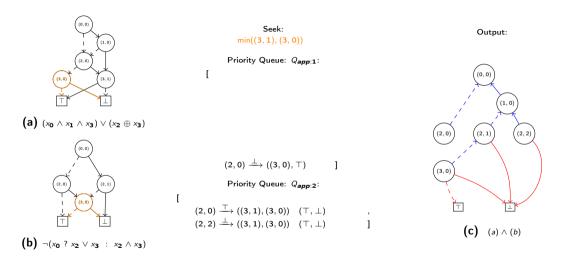


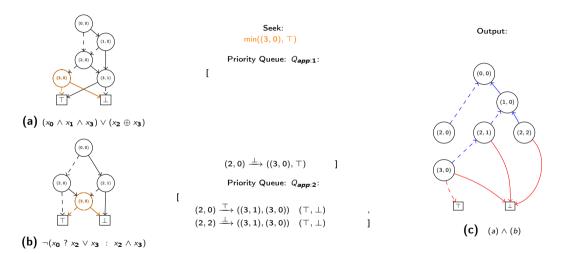


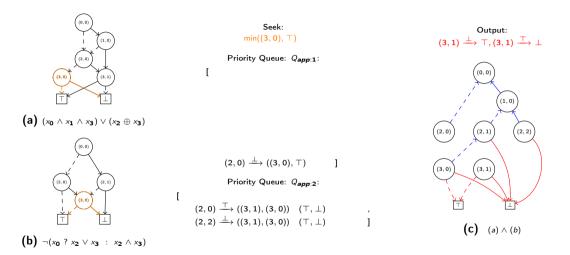


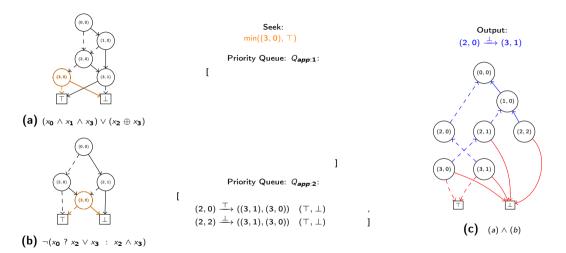


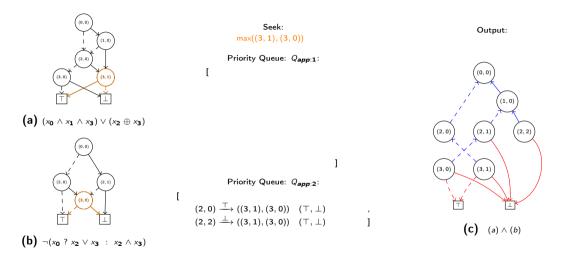


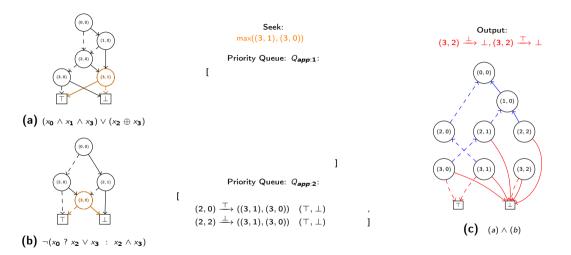


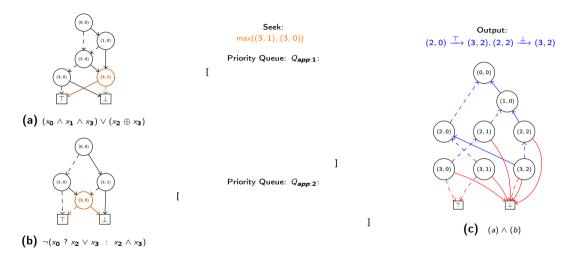


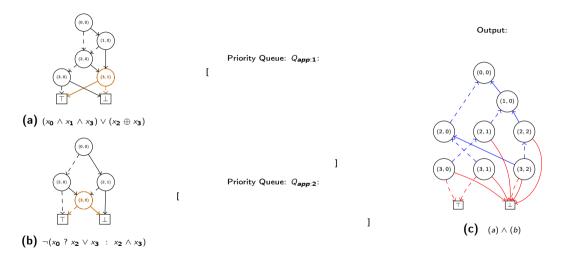


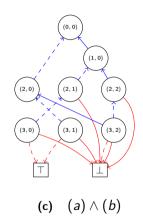


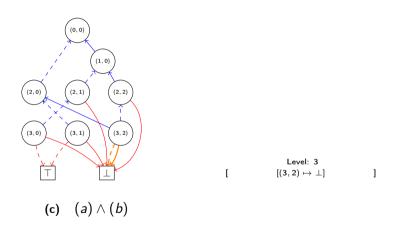


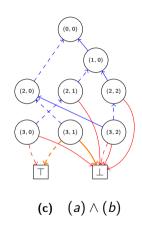


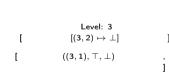


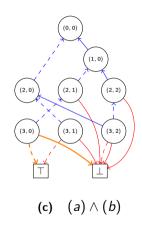


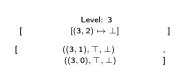


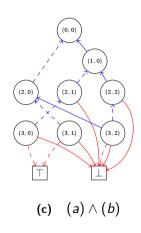


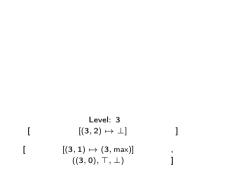




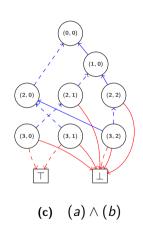


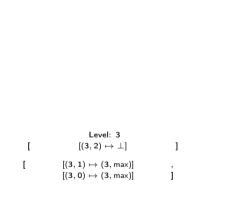


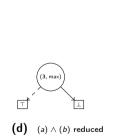


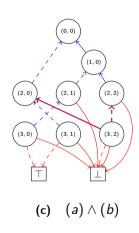


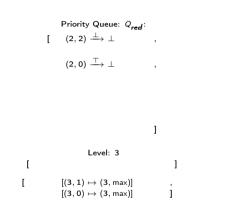


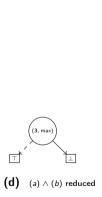


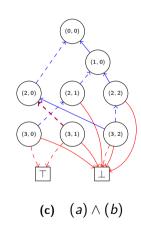


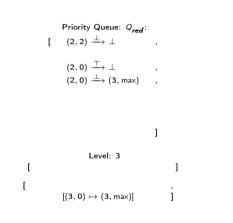


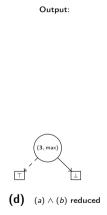


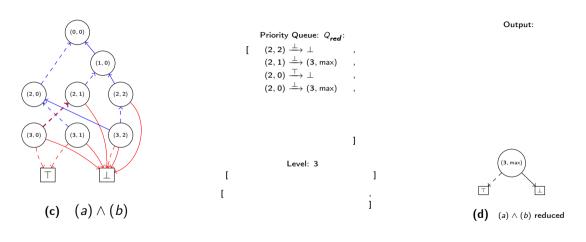


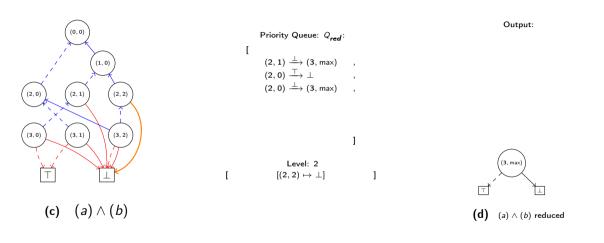


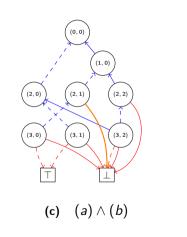


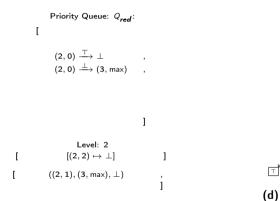




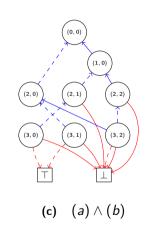


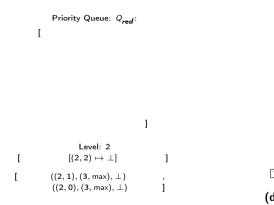




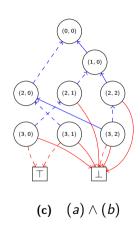


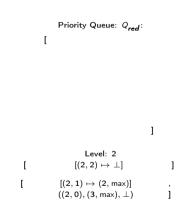


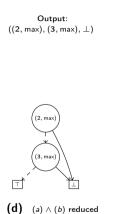


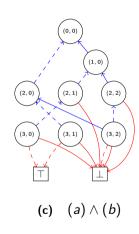


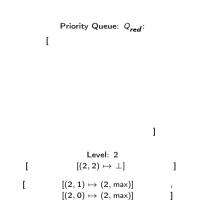


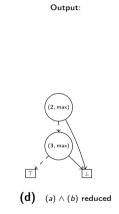


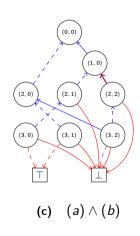


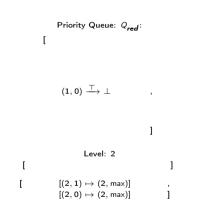


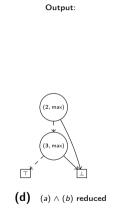


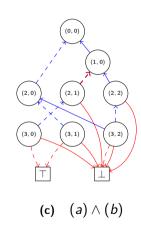


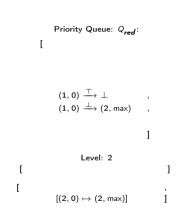


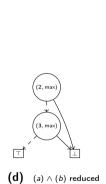


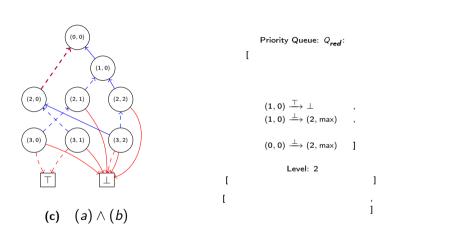


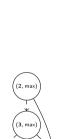




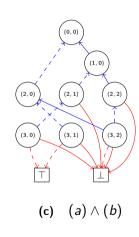


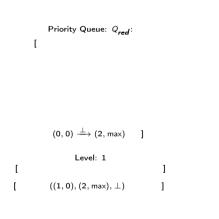


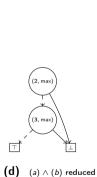


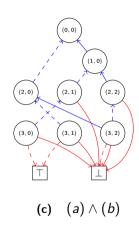


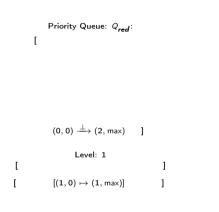
(d) $(a) \wedge (b)$ reduced

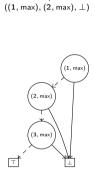




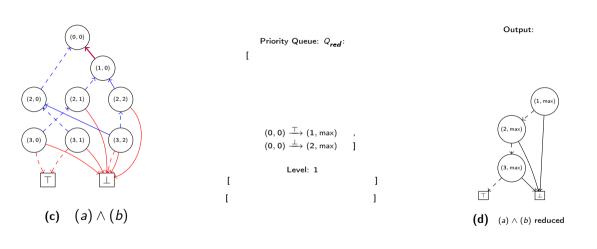


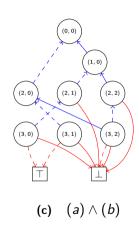


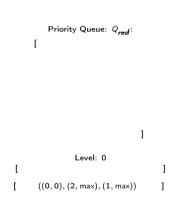


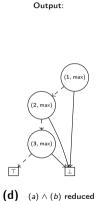


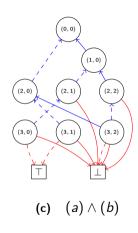
(d) $(a) \wedge (b)$ reduced

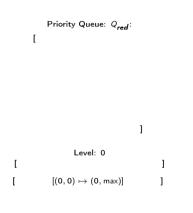


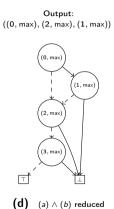




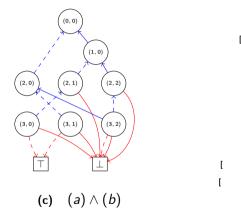


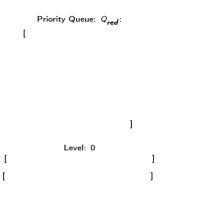


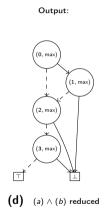


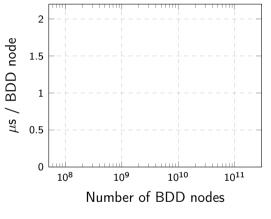


10





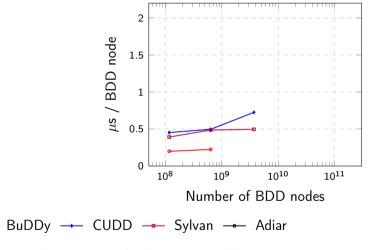




→ BuDDy → CUDD → Sylvan → Adiar

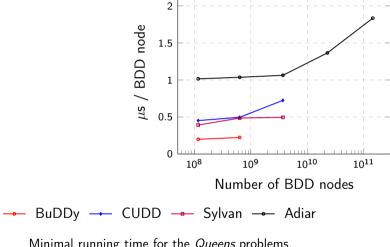
Minimal running time for the *Queens* problems.

Adiar v1.0: BDD



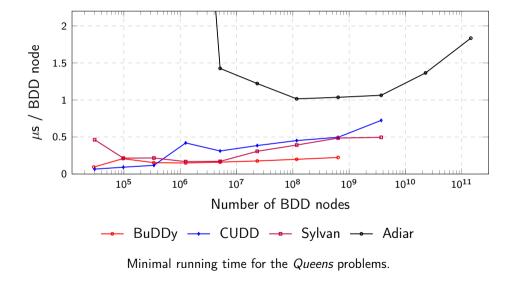
Minimal running time for the *Queens* problems.

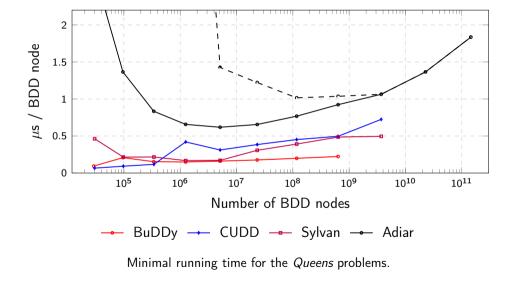
Adiar v1.0: BDD



Minimal running time for the Queens problems.

Adiar v1.0: BDD

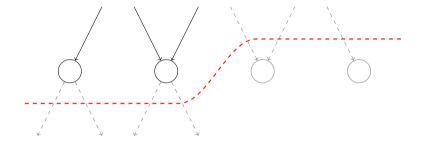


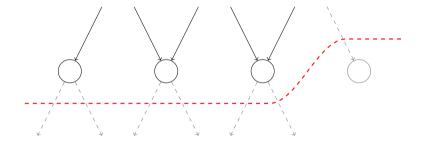


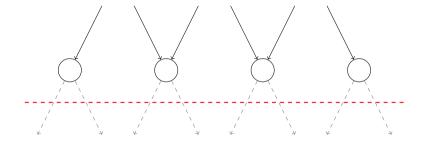




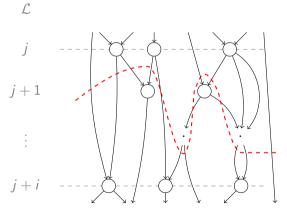








Definition (i-level cut)



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Lemma

The maximum i-level cut problem is in P for $i \in \{1, 2\}$.

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Theorem (Lampis, Kaouri, Mitsou 2011) The maximum i-level cut problem is NP-complete for $i \ge 4$.

Theorem

The maximum (i-level) cut of a BDD with N internal nodes is N+1.

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For $i \in \{1,2\}$, the maximum i-level cut of the (unreduced) output of Apply is upper bounded by the product of the inputs' corresponding i-level cuts.

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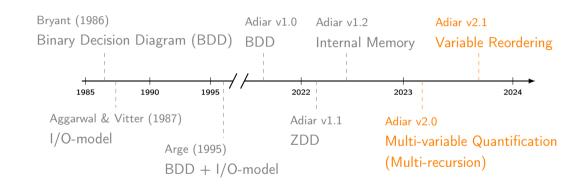
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Theorem

For $i \in \{1,2\}$, the maximum i-level cut of the (unreduced) output of Apply is upper bounded by the product of the inputs' corresponding i-level cuts.

Lemma

The maximum 2-level cut of a BDD is upper bounded by $\frac{3}{2}$ its 1-level cut.



$$RelProd(S, T) \equiv (\exists \vec{x}. S(\vec{x}) \land T(\vec{x}, \vec{x'}))[\vec{x'}/\vec{x}]$$

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```
1 exists (f, V)
     if f = 1 \lor f = \top
            then f
    else if V \cap \{i \in \mathbb{N} \mid i \geq \mathsf{top}(f)\} = \emptyset
5
            then f
     else if top(f) \notin V
            then Node { top(f), exists(f.low, V), exists(f.high, V) }
8
     else let low = exists(f.low. V)
                  high = exists(f.high, V)
10
            in or(low, high)
```

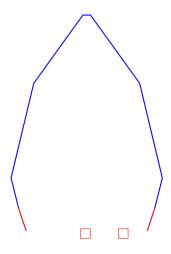
A recursive multi-variable exists operation.

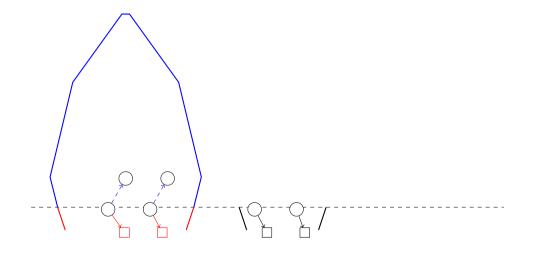
```
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    if f = \bot \lor f = \top
     then f
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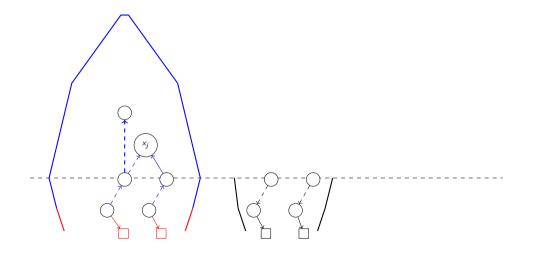
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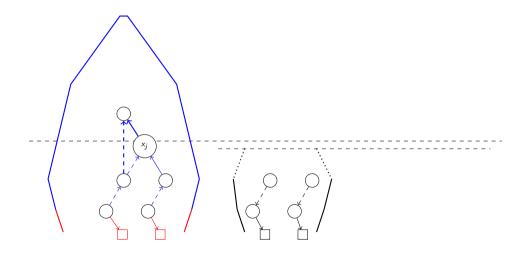
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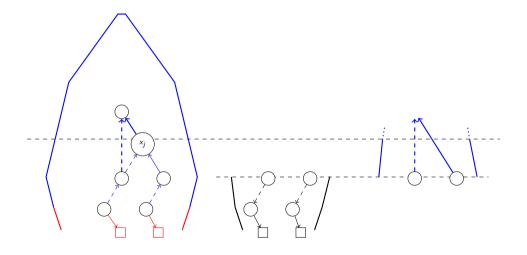
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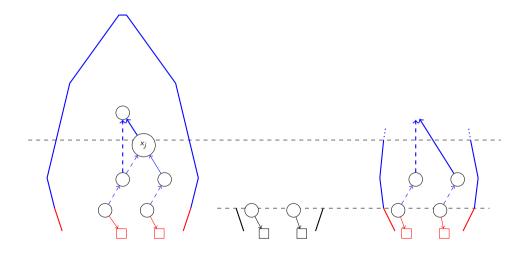


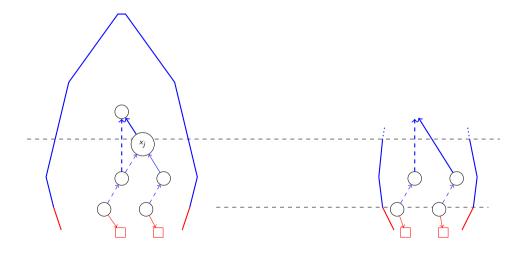


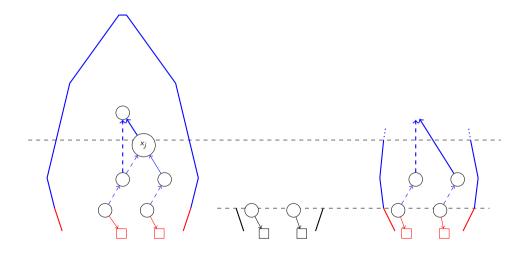


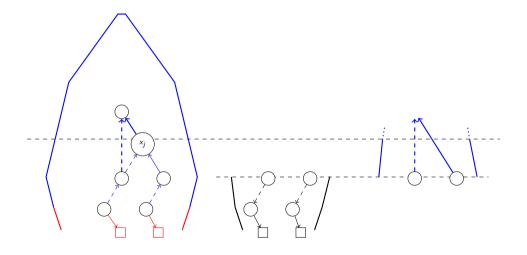


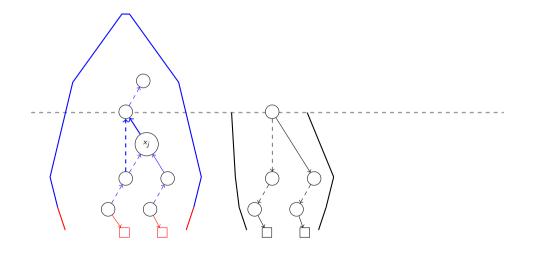


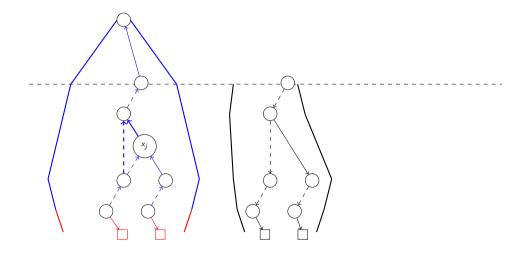


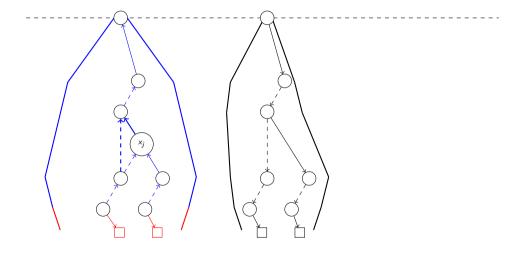


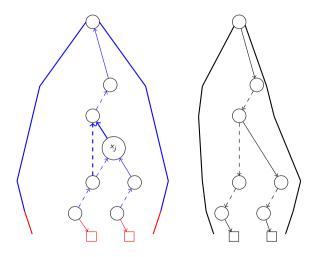






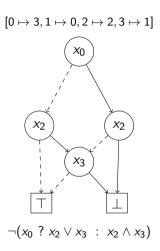




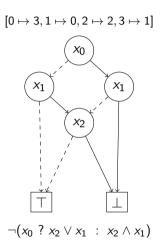


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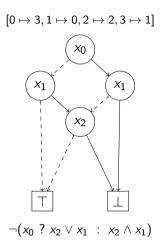
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3    high = substitute(f.high)
4    i' = i_map[top(f)]
5  in bubble(i', low, high)
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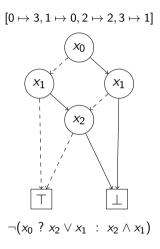
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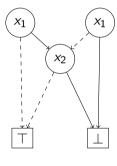
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[0\mapsto 3, 1\mapsto 0, 2\mapsto 2, 3\mapsto 1]
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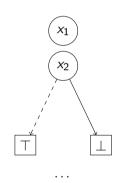
A recursive substitute operation.



. . .

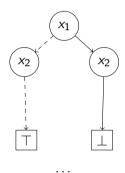
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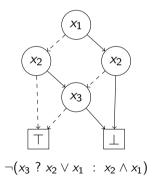
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\frac{\text{Time Space }I/O}{O(NT) O(NT)}
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Problem (Variable Replacement) Given BDD f_{π} with variable ordering π and remapping of variables $m: \mathbb{N} \to \mathbb{N}$, construct $f'_{\pi} \equiv f_{\pi}[x/m(x)]$.

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Given BDD f_{π} with variable ordering π and another variable ordering π' , construct $f_{\pi'} \equiv f_{\pi}$.

Problem (Dynamic Variable Reordering)

Given BDD f_{π} with variable ordering π ,

find π' and construct $f_{\pi'} \equiv f_{\pi}$ such that $|f_{\pi'}|$ is minimal.

