# I/O-efficient Manipulation of Binary Decision Diagrams

Steffan Christ Sølvsten

S. C. Sølvsten, J. van de Pol, A. B. Jakobsen, and M. W. B. Thomasen. *Adiar: Binary Decision Diagrams in External Memory.* 2022



#### **Contents**

What are Binary Decision Diagrams?

Why do they break?

How can we fix it?

CountPaths

Apply

**Equality Checking** 

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What are Binary Decision Diagrams?

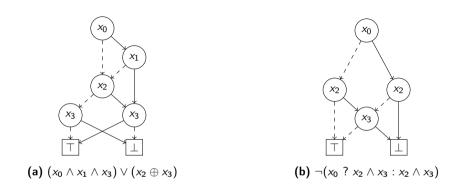
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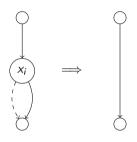
Apply

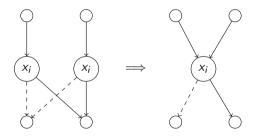
**Equality Checking** 



Examples of (Reduced Ordered) Binary Decision Diagrams.

**Theorem (Bryant '86)**For a fixed variable order, if one exhaustively applies the two rules below, then one obtains the Reduced OBDD, which is a unique canonical form of the function.





(1) Remove redundant nodes

(2) Merge duplicate nodes

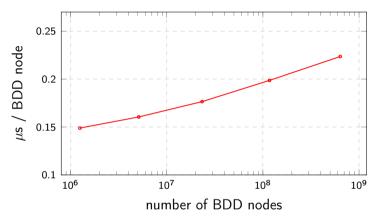
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\begin{array}{ll} \operatorname{bdd\_apply}\left(f,\ g\ ,\ \otimes\right): \\ & \text{if}\ f,g\in\{\bot,\top\} \\ & \text{then}\ f\otimes g \\ & \text{else let}\ i = \operatorname{top}\left(f.(\mathit{var}),\ g.\mathit{var}\right) \\ & t = \operatorname{bdd\_apply}\left(f[x_i := \top],\ g[x_i := \top],\ \otimes\right) \\ & e = \operatorname{bdd\_apply}\left(f[x_i := \bot],\ g[x_i := \bot],\ \otimes\right) \\ & \text{in make\_node}\left(i,t,e\right) \end{array}
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\begin{array}{l} \operatorname{bdd\_apply}\left(f,\ g\ ,\ \otimes\right):\\ & \text{if}\ f,g\in\{\bot,\top\}\\ & \text{then}\ f\otimes g\\ & \text{else let}\ i=\operatorname{top}(f.(\mathit{var}),\ g.\mathit{var})\\ & t=\operatorname{bdd\_apply}\left(f[x_i:=\top],\ g[x_i:=\top],\ \otimes\right)\\ & e=\operatorname{bdd\_apply}\left(f[x_i:=\bot],\ g[x_i:=\bot],\ \otimes\right)\\ & \text{in make\_node}\left(i,t,e\right) \end{array}
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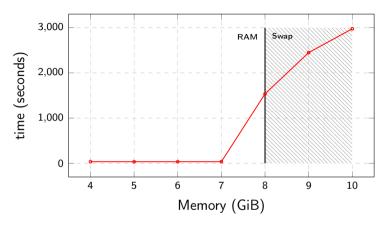
#### Theorem

bdd\_apply runs in  $O(N_f \cdot N_g)$  time.

- Memoisation (*Computation Cache*) ensures each recursion is computed only once.
- Reduction Rules can be maintained within make\_node(i,t,e) in O(1) time.
  - 1 Redundancy is resolved with an if-statement.
  - 2 Duplication is avoided with a hash table (*Unique Node Table*).



Running time of *BuDDy* for the *N*-Queens problem.



Running time of BuDDy for Tic-Tac-Toe with N=21.

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What are Binary Decision Diagrams?

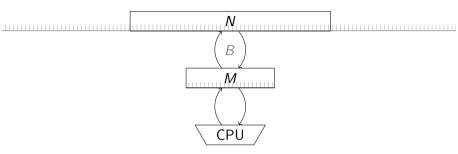
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**Equality Checking** 



The I/O model by Aggarwal and Vitter '87

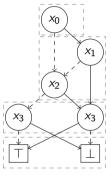
For any realistic values of N, M, and B we have that

$$N/B < \operatorname{sort}(N) \triangleq N/B \cdot \log_{M/B} N/B \ll N$$
,

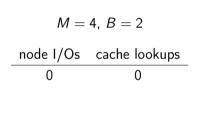
Theorem (Aggarwal and Vitter '87) N elements can be sorted in  $\Theta(sort(N))$  I/Os.

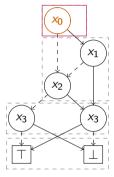
# Theorem (Arge '95)

N elements can be inserted in and extracted from a Priority Queue in  $\Theta(sort(N))$  I/Os.

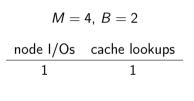


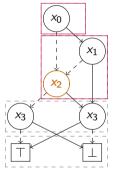
(a) 
$$(x_0 \land x_1 \land x_3) \lor (x_2 \oplus x_3)$$



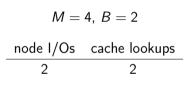


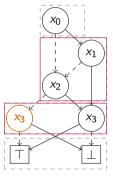
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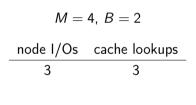


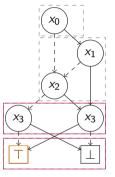
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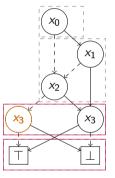
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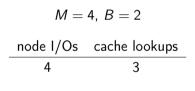


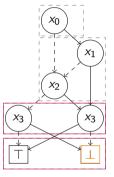
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$$M = 4$$
,  $B = 2$   
node I/Os cache lookups  
4 3

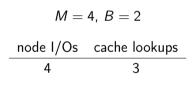


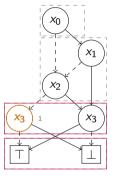
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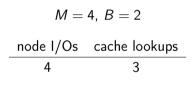


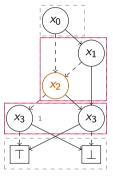
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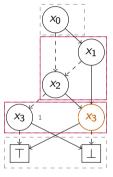
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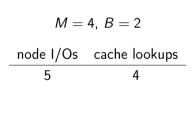


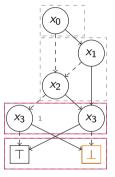
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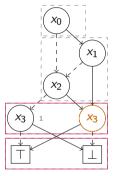
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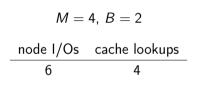


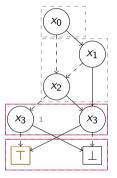
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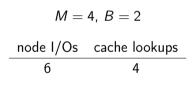


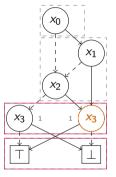
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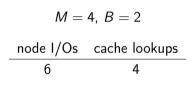


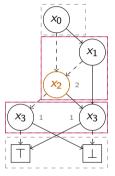
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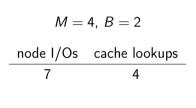


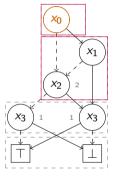
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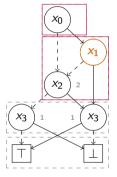
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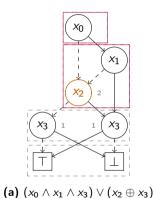
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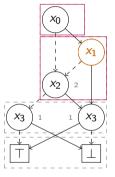


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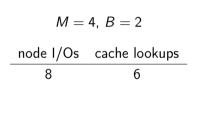
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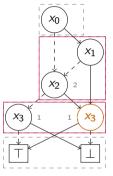


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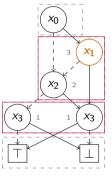


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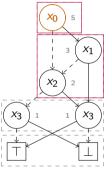




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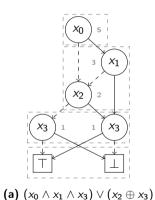


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Algorithm	Time Complexity
bdd_pathcount	$O(N_f)$
bdd_not	$O(N_f)$
bdd_restrict	$O(N_f)$
bdd_apply	$O(N_f \cdot N_g)$
bdd_equal	O(1)

Algorithm	I/O-Complexity
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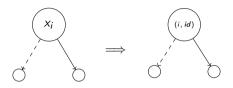
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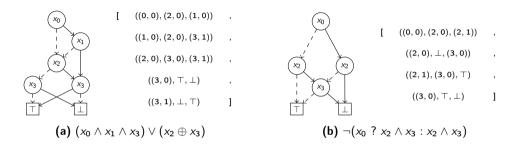
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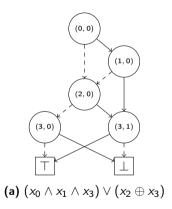
**Equality Checking** 

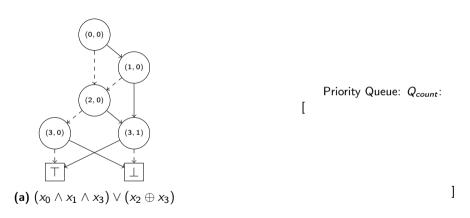


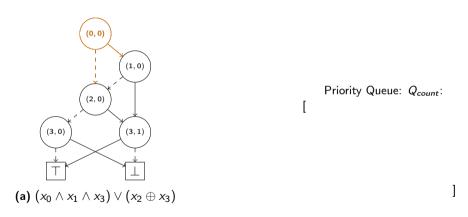
$$(i_1, id_1) < (i_2, id_2) \equiv i_1 < i_2 \lor (i_1 = i_2 \land id_i < id_j)$$

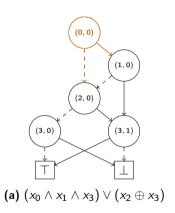


Node-based representation of prior shown BDDs

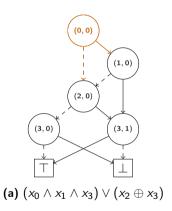


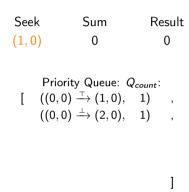


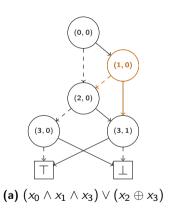


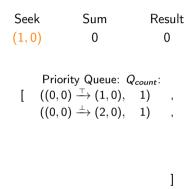


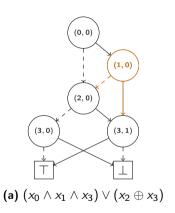
Priority Queue: 
$$Q_{count}$$
: [  $((0,0) \xrightarrow{\top} (1,0), 1)$  ,  $((0,0) \xrightarrow{\bot} (2,0), 1)$  ,

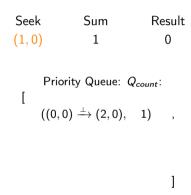


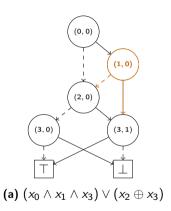


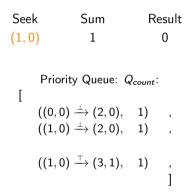


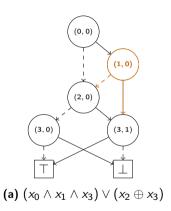


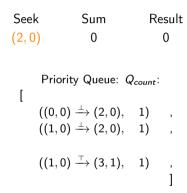


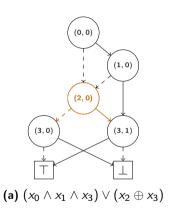


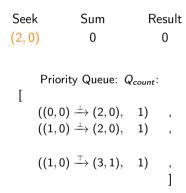


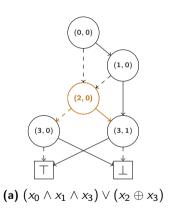


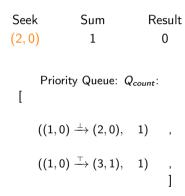


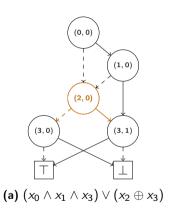


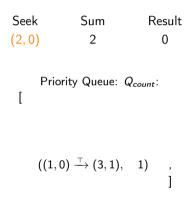


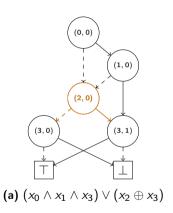




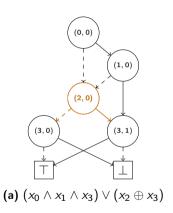




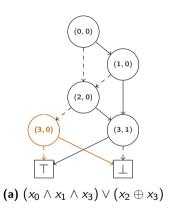




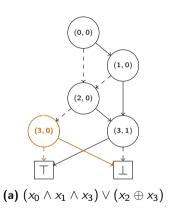
```
Seek
                                   Sum
                                                                      Result
(2,0)
                                        2
                                                                             0
              Priority Queue: Qcount:
             \begin{array}{cccc} ((2,0) \xrightarrow{\bot} (3,0), & 2) & , \\ ((1,0) \xrightarrow{\top} (3,1), & 1) & , \\ ((2,0) \xrightarrow{\top} (3,1), & 2) & ] \end{array}
```

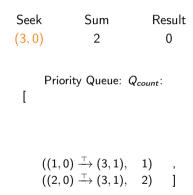


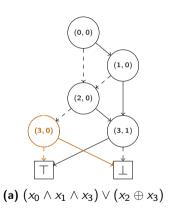
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                                    Sum
                                                                      Result
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                                        0
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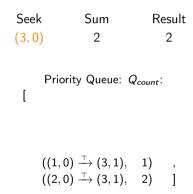


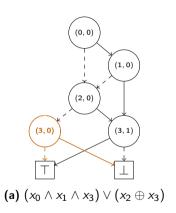
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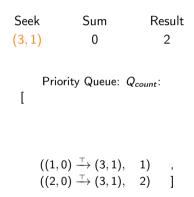


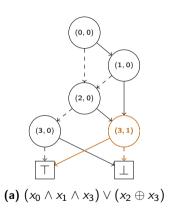


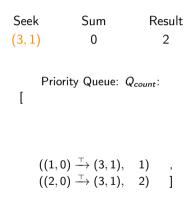


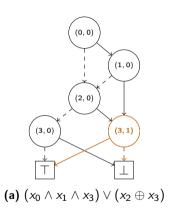


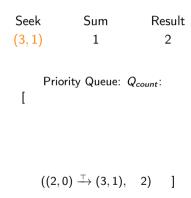


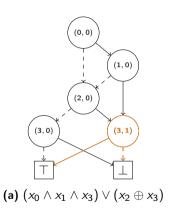


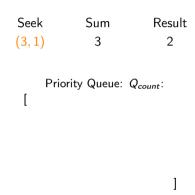


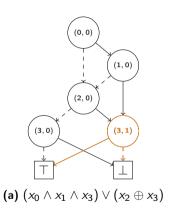


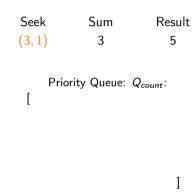


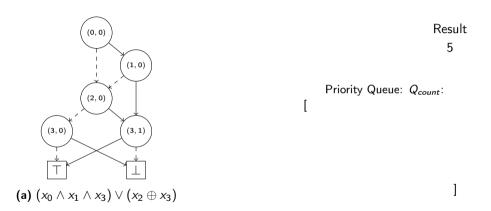












#### **Contents**

What are Binary Decision Diagrams?

Why do they break?

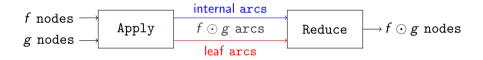
How can we fix it?

CountPaths

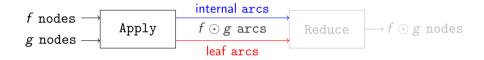
Apply

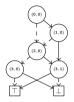
**Equality Checking** 

#### **Apply**



# **Apply**



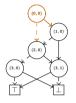


(a)  $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$ 



(b)  $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$ 

(c)  $(a) \wedge (b)$ 

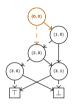


(a)  $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$ 



(b)  $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$ 

(c)  $(a) \wedge (b)$ 



Priority Queue: Qapp:1:

- [  $(0,0) \xrightarrow{\top} ((1,0),(2,1))$  ,
  - $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$

(0,0)

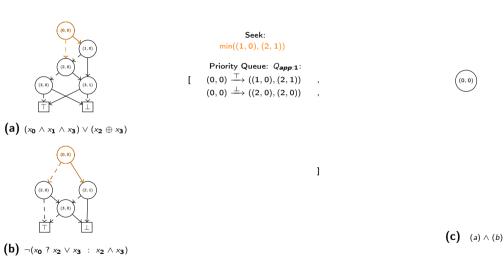
(a)  $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$ 

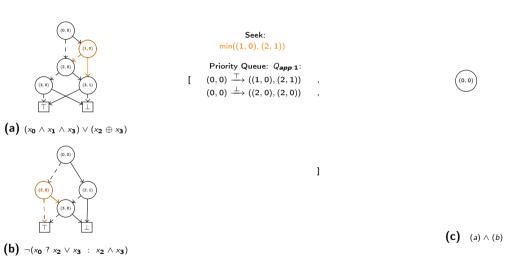


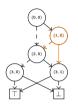
(b)  $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$ 

1

**(c)** (a) ∧ (b)







(a)  $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$ 



**(b)**  $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$ 

 $\begin{array}{c} \text{Seek:} \\ \min((1,0),(2,1)) \end{array}$ 

Priority Queue: Qapp:1:

 $(0,0) \xrightarrow{\top} ((1,0),(2,1))$ 

 $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$ 

 $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$ 

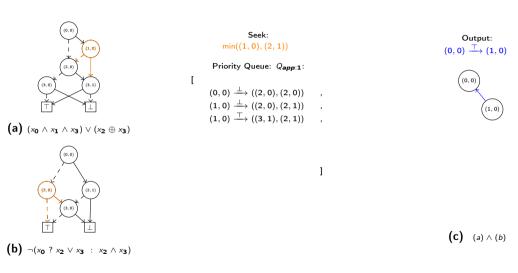
 $(1,0) \xrightarrow{\top} ((3,1),(2,1))$ 

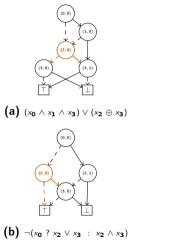
(0,0)

(1,0)

J

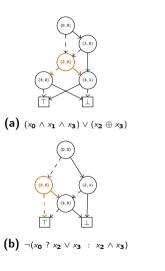
(c)  $(a) \wedge (b)$ 





Seek: min((2,0),(2,0))Priority Queue: Qapp:1:  $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$  $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$  $(1,0) \xrightarrow{\top} ((3,1),(2,1))$ 

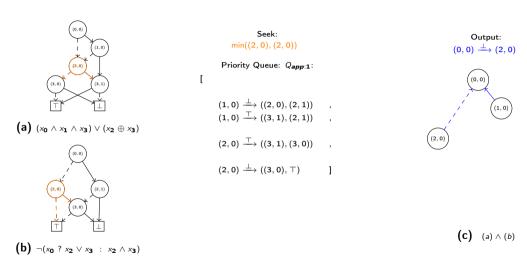
Output: (c)  $(a) \wedge (b)$ 

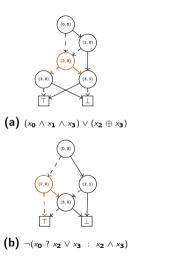


Seek: min((2,0),(2,0))Priority Queue: Qapp:1:  $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$  $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$  $(1,0) \xrightarrow{\top} ((3,1),(2,1))$  $(2,0) \xrightarrow{\top} ((3,1),(3,0))$  $(2,0) \xrightarrow{\perp} ((3,0),\top)$  ]

Output: (2,0)

(c)  $(a) \wedge (b)$ 

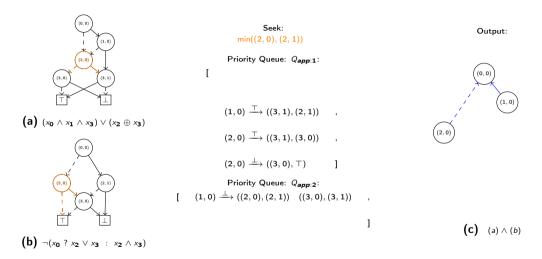


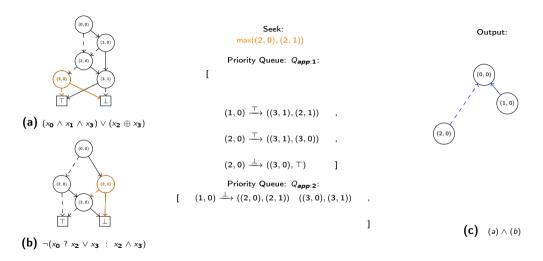


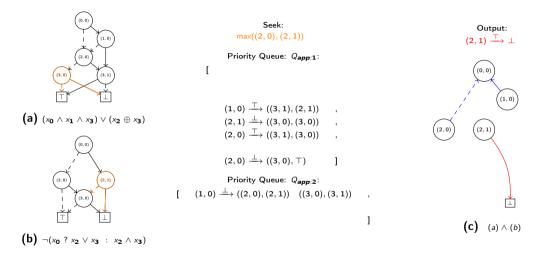
Seek: min((2,0),(2,1))Priority Queue: Qapp:1:  $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$  $(1,0) \xrightarrow{\top} ((3,1),(2,1))$  $(2,0) \xrightarrow{\top} ((3,1),(3,0))$  $(2,0) \xrightarrow{\perp} ((3,0),\top)$ 

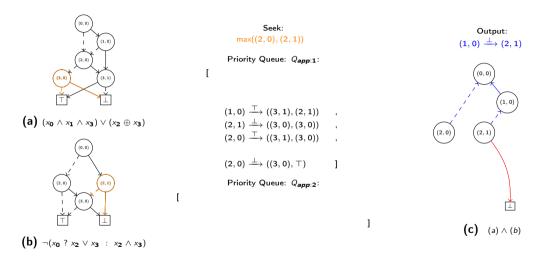
Output: (0,0) (1,0)

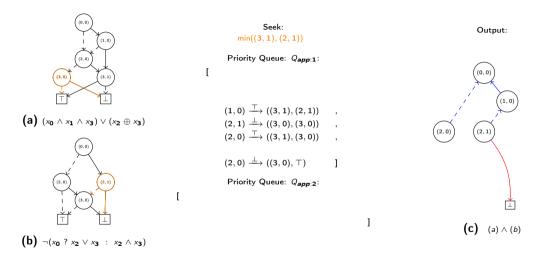
(c)  $(a) \wedge (b)$ 

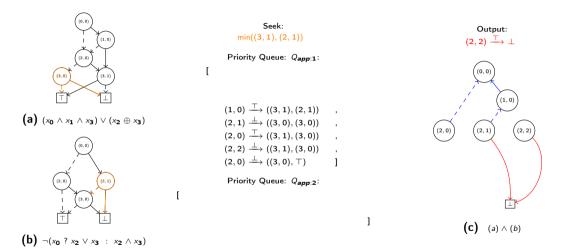


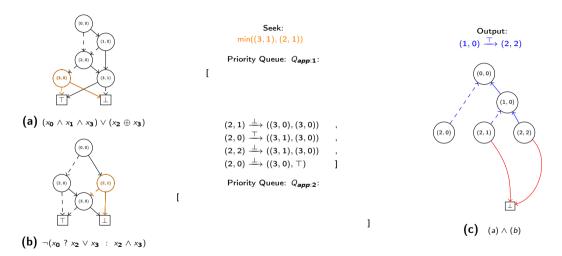


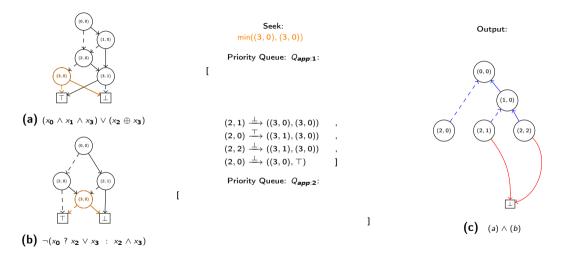


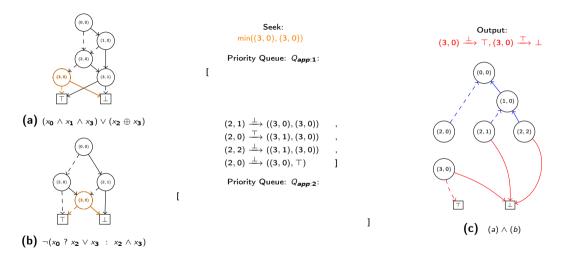


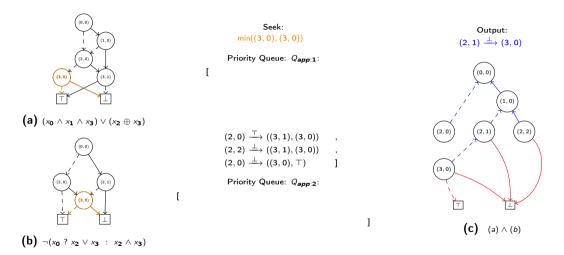


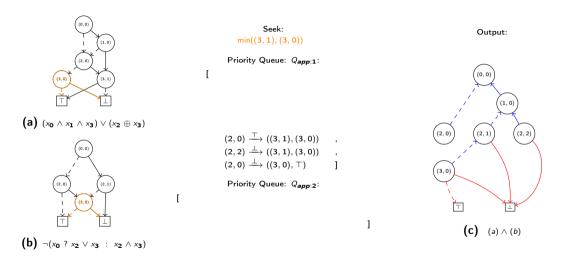


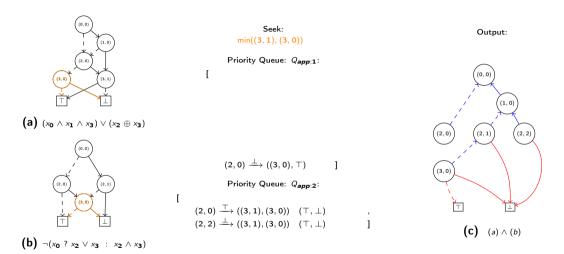


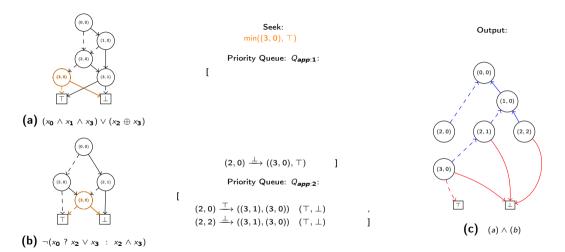


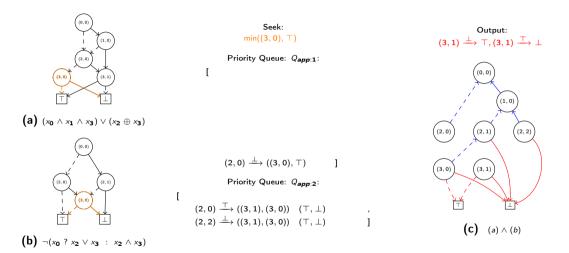


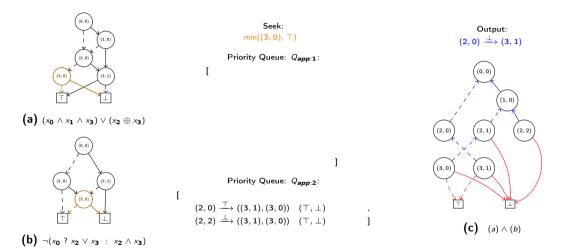


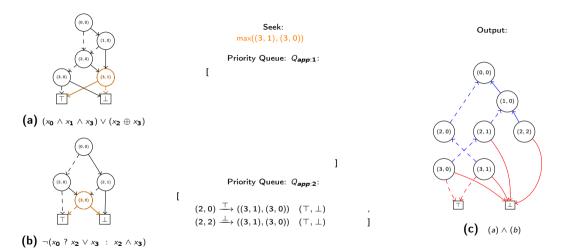


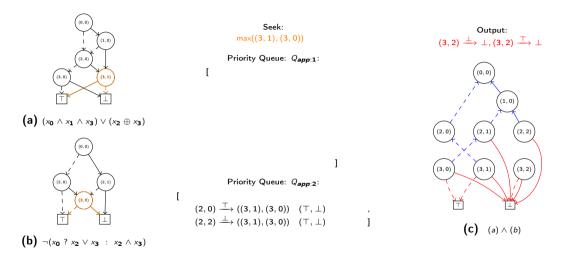


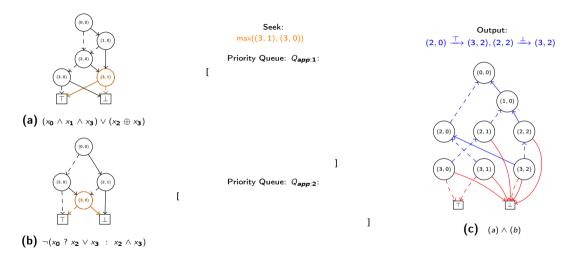


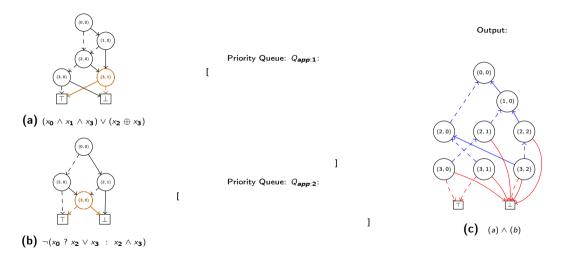




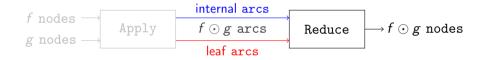


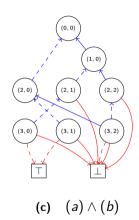


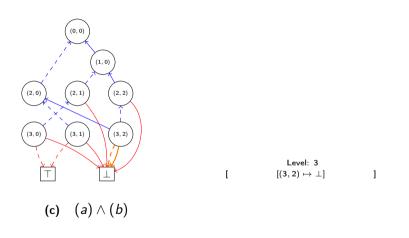


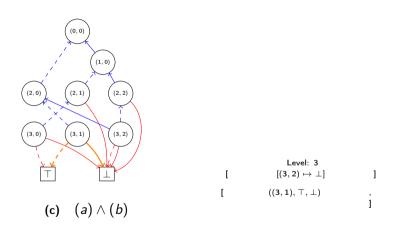


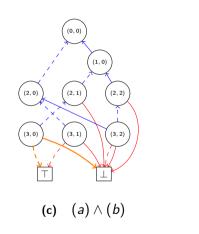
# **Apply**

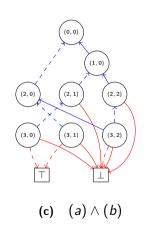


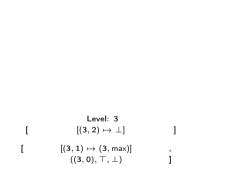




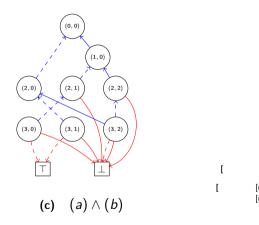


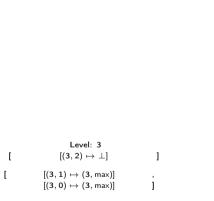


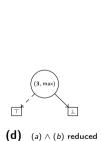




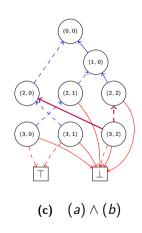


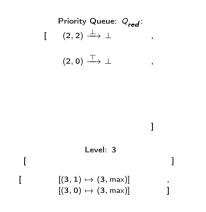






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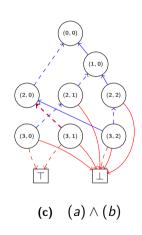


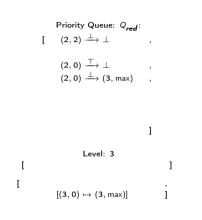




(d)  $(a) \wedge (b)$  reduced

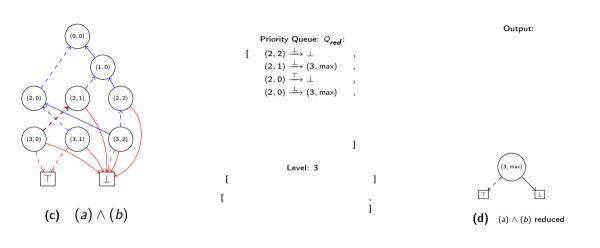
Output:

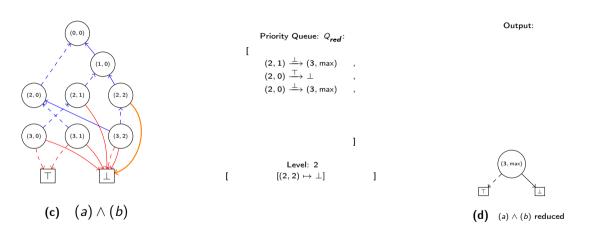


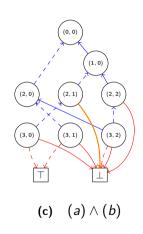


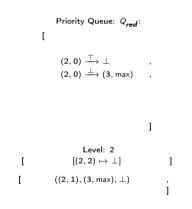




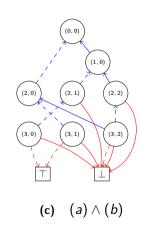


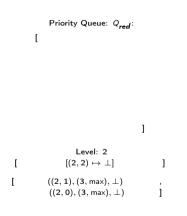




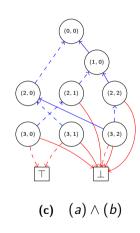


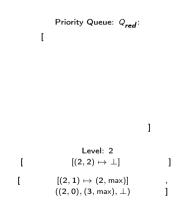


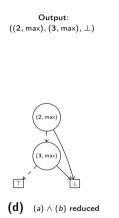


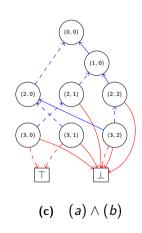


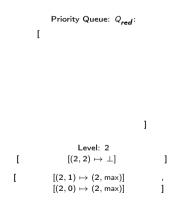


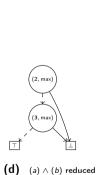


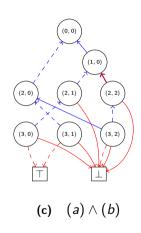


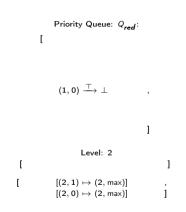


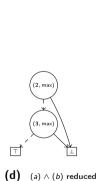


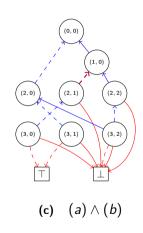


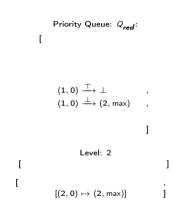


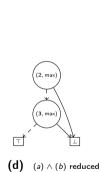


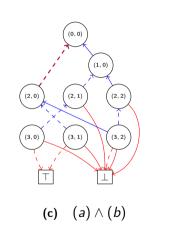


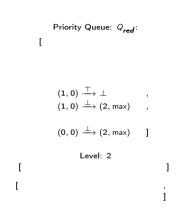


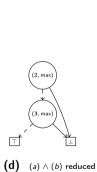


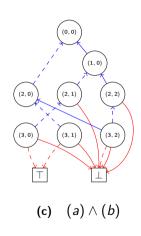


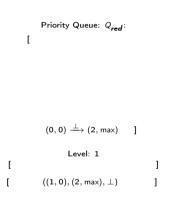


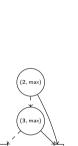




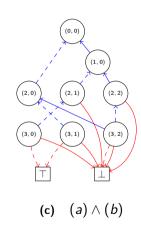


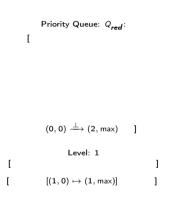


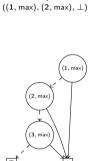


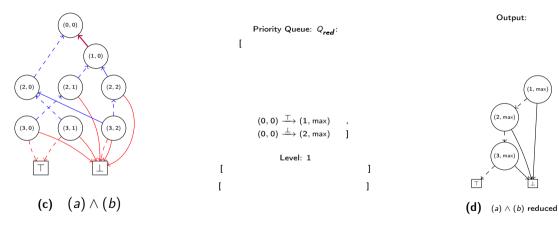


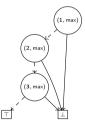
(d)  $(a) \wedge (b)$  reduced

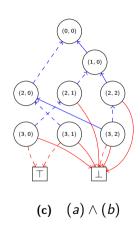


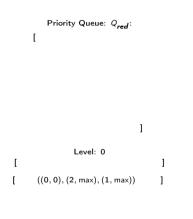


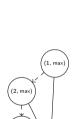


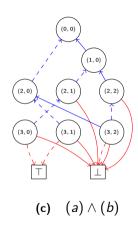


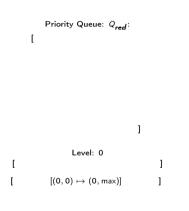


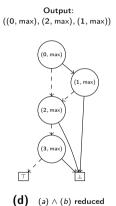


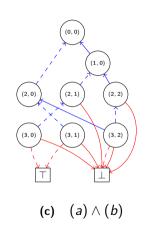


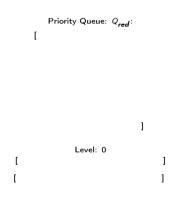


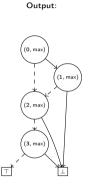






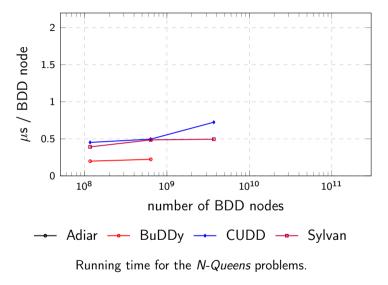


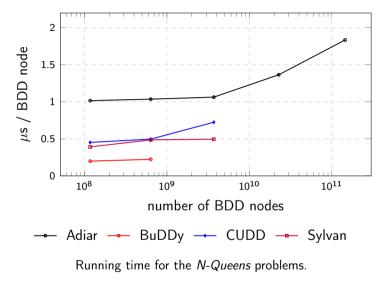




(d)  $(a) \wedge (b)$  reduced

Algorithm	I/O-Complexity	
bdd_pathcount	$O(\operatorname{sort}(N_f))$	
bdd_not	$O(N_f/B)$	
bdd_restrict	$O(\operatorname{sort}(N_f))$	
bdd_apply	$O(sort(\mathit{N_f}\cdot\mathit{N_g}))$	





#### Contents

What are Binary Decision Diagrams?

Why do they break?

How can we fix it?

CountPaths

Apply

**Equality Checking** 

Algorithm	I/O-Complexity
bdd_pathcount	$O(\operatorname{sort}(N_f))$
bdd_not	$O(N_f/B)$
bdd_restrict	$O(\operatorname{sort}(N_f))$
bdd_apply	$O(\operatorname{sort}(N_f \cdot N_g))$

Algorithm	I/O-Complexity
bdd_pathcount	$O(\operatorname{sort}(N_f))$
bdd_not	$O(N_f/B)$
bdd_restrict	$O(\operatorname{sort}(N_f))$
bdd_apply	$O(\operatorname{sort}(N_f \cdot N_g))$
bdd_equal	?

$$f\leftrightarrow g\equiv \top$$

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#### Theorem (Bryant '86)

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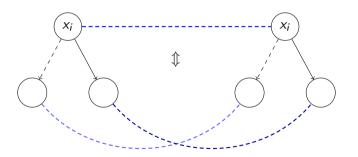
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0(1) 1/0

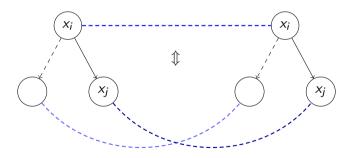
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	$N_f  eq N_g$	Number of nodes	O(1) I/Os
•	$L_f  eq L_g$	Number of levels	<i>O</i> (1) I/Os
•	$N_{f,i} \neq N_{g,i}$	Number of nodes on a level	O(L/B) I/Os
•	$L_{f,i} \neq L_{g,i}$	Label of an <i>i</i> th level	O(L/B) I/Os

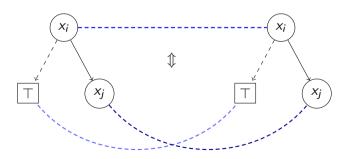
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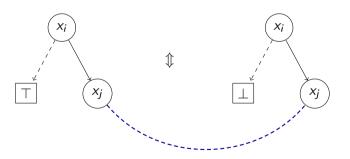
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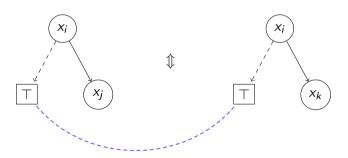
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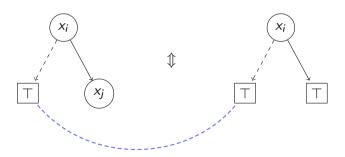
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Let  $\pi$  be a variable order and  $f: \mathbb{B}^n \to \mathbb{B}$  then there exists a unique (up to isomorphism) Reduced Ordered Binary Decision Diagram representing f with ordering  $\pi$ .

IsIsomorphic(f, g)

- Check whether root  $v_f$  of f and root  $v_g$  of g have a local violation.
- Check  $low(v_f) \sim low(v_g)$  and  $high(v_f) \sim high(v_g)$  "recursively".

Return false on first violation. If there are no violations then return true.

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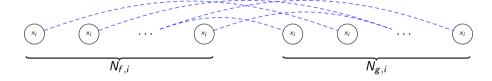
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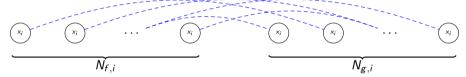
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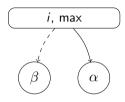
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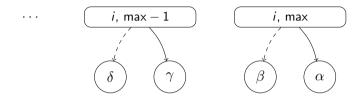
Return false if more than  $N_{f,i} = N_{g,i}$  pairs of nodes are checked on level i.

$$\underbrace{O(\mathsf{sort}(\Sigma_i \ \mathsf{N}_{f,i}))}_{\mathsf{Apply''}} = O(\mathsf{sort}(\mathsf{N}))$$

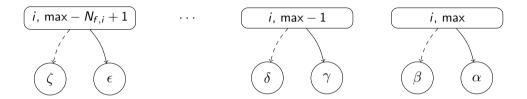
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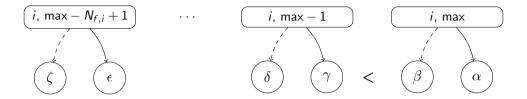
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Each level output by the Reduce algorithm has the following properties:

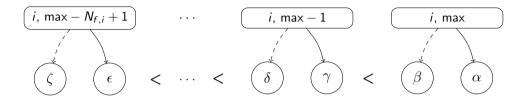
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#### Observation

- Nodes on level *i* have their identifiers *consecutively* numbered.
- Nodes on level *i* are output sorted by their children.

### **Theorem**

If  $G_f$  and  $G_g$  are outputs of Reduce.

 $G_f \sim G_g \iff For \ all \ i \in [0; N_f) \ the \ node \ G_f[i] \ matches \ G_g[i] \ numerically.$ 

#### Proof.

← : Must describe the exact same graph.

 $\Rightarrow$ : Strong induction on BDD levels bottom-up.

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### Corollary

If  $G_f$  and  $G_g$  are outputs of Reduce then  $f \equiv g$  is computable using  $2 \cdot N/B$  I/Os.

Checking the (EPFL Benchmark) *voter* circuit's single output gate ( $|N_f| = |N_g| = 5.76$  MiB).

Algorithm Time (s)
$$f \leftrightarrow g \equiv \top \quad 0.38$$

$$O(\operatorname{sort}(N)) \quad 0.058$$

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Algorithm	Time (s)	
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2N/B	0.006	

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#### Contents

What are Binary Decision Diagrams? Why do they break? How can we fix it? Depth-First Time-Forwarded CountPaths  $O(N_f)$   $O(\operatorname{sort}(N_f))$  $O(N_f \cdot N_g)$   $O(\operatorname{sort}(N_f))$ **Apply Equality Checking** O(1)

2N/B