I/O-efficient Manipulation of Binary Decision Diagrams

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S. C. Sølvsten, J. van de Pol, A. B. Jakobsen, and M. W. B. Thomasen. *Adiar: Binary Decision Diagrams in External Memory.* 2022



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Why do they break?

How can we fix it?

CountPaths

Apply

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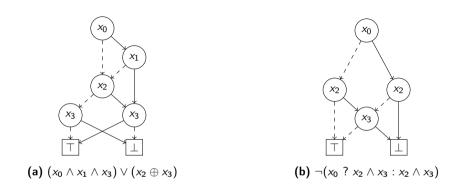
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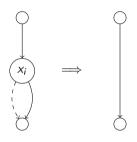
Apply

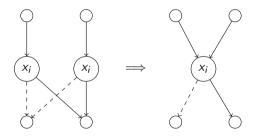
Equality Checking



Examples of (Reduced Ordered) Binary Decision Diagrams.

Theorem (Bryant '86)For a fixed variable order, if one exhaustively applies the two rules below, then one obtains the Reduced OBDD, which is a unique canonical form of the function.





(1) Remove redundant nodes

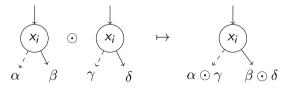
(2) Merge duplicate nodes

 $bdd_apply(f,g,\odot)$

Base Case $(f, g \in \mathbb{B})$:



Inductive Case:

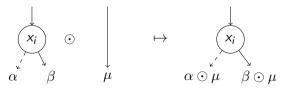


 $bdd_apply(f,g,\odot)$

Base Case $(f, g \in \mathbb{B})$:



Inductive Case:



$$bdd_apply(f,g,\odot)$$

Let N_f , N_g be the size of the BDDs for f and g.

Let T be the $O(N_f \cdot N_g)$ size of the BDD for $f \odot g$.

Theorem

 $bdd_apply(f,g,\odot)$ runs in $O(N_f + N_g + T)$ time

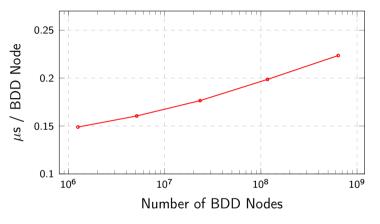
- Memoisation (*Computation Cache*) ensures each (t_f, t_g) is only computed once.
- Reduction Rules can be maintained with a make_node(i, t, e) in O(1) time.
 - 1 Redundancy is resolved with an if-statement.
 - 2 Duplication is avoided with a hash table (*Unique Node Table*).

Corollary

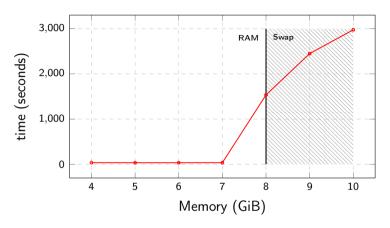
 $bdd_apply(f,g,\odot)$ runs in O(1) time per BDD node.

Adiar

github.com/ssoelvsten/adiar



Running time of *BuDDy* for the *N*-Queens problem.



Running time of BuDDy for 3D Tic-Tac-Toe with N=21.

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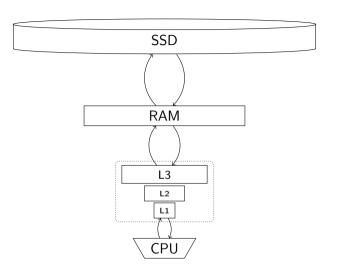
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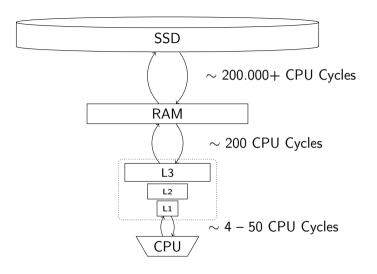
How can we fix it?

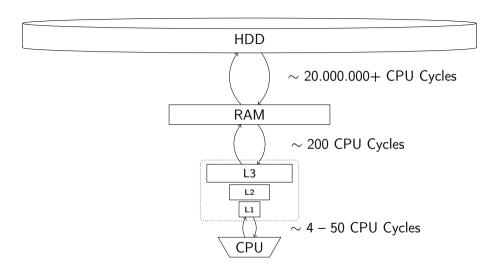
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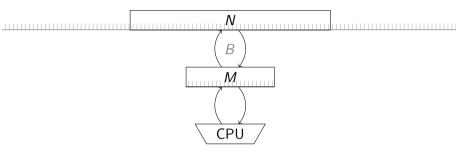
Apply

Equality Checking









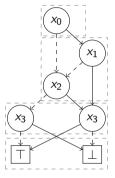
The I/O model by Aggarwal and Vitter '87

For any realistic values of N, M, and B we have that

$$N/B < \operatorname{sort}(N) \triangleq N/B \cdot \log_{M/B} N/B \ll N$$
,

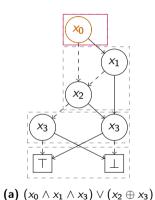
Theorem (Aggarwal and Vitter '87) N elements can be sorted in $\Theta(sort(N))$ I/Os.

Theorem (Arge '95) A Priority Queue can do N insertions and extractions in $\Theta(sort(N))$ I/Os.

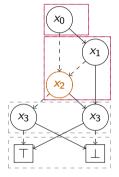


(a)
$$(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$$

$$M = 4$$
, $B = 2$
node I/Os cache lookups
0 0

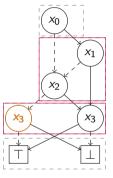


M = 4, B = 2



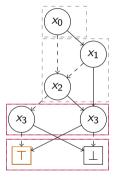
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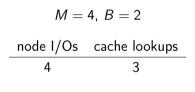


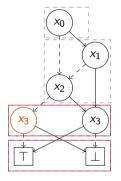
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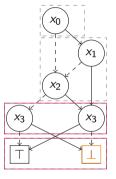
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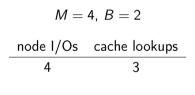


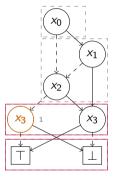
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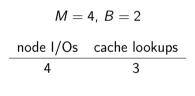


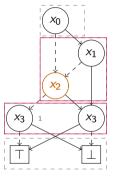
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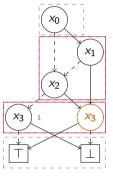
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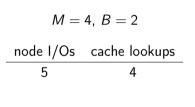


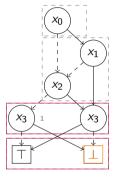
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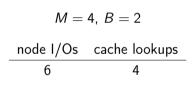


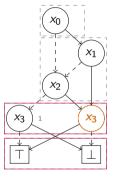
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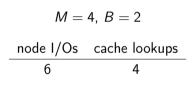


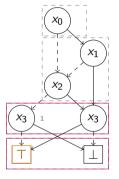
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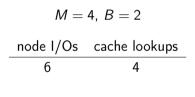


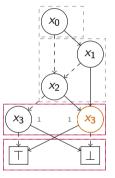
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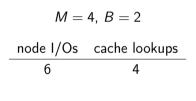


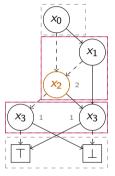
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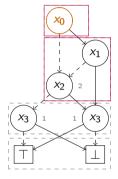
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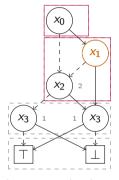
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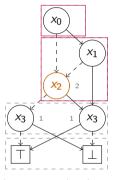
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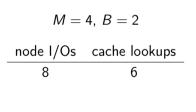


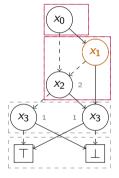
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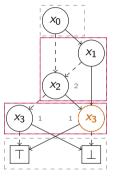


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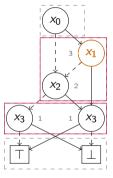
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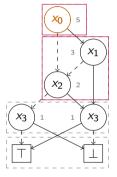


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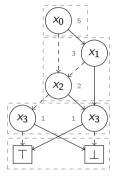


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Algorithm	Time Complexity
bdd_pathcount	$O(N_f)$
bdd_not	$O(N_f)$
bdd_restrict	$O(N_f)$
bdd_apply	$O(N_f \cdot N_g)$
bdd_equal	O(1)

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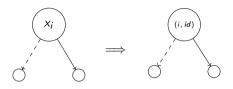
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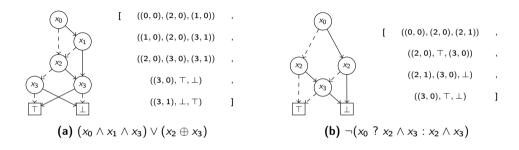
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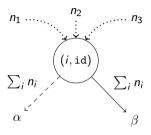


$$(i_1, id_1) < (i_2, id_2) \equiv i_1 < i_2 \lor (i_1 = i_2 \land id_i < id_j)$$



Node-based representation of prior shown $\ensuremath{\mathsf{BDDs}}$

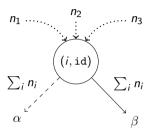
CountPaths



Idea

Count the number of in-going paths to each node.

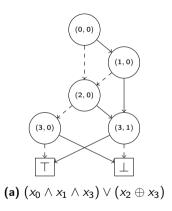
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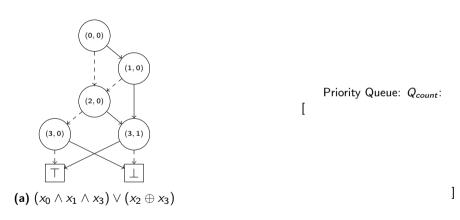


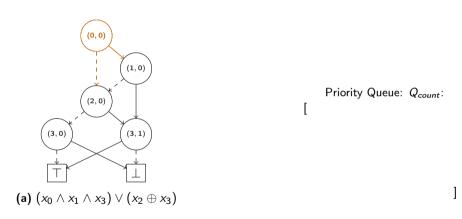
Time-Forward Processing

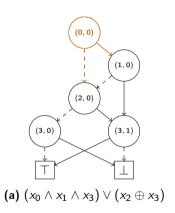
Defer work with Q_{count} : PriorityQueue $\langle (s \to t, \mathbb{N}) \rangle$ sorted on t in ascending order.

$$((i, \mathrm{id}) \xrightarrow{\perp} \alpha, \quad \sum_{i} n_{i}), \qquad ((i, \mathrm{id}) \xrightarrow{\top} \beta, \quad \sum_{i} n_{i})$$

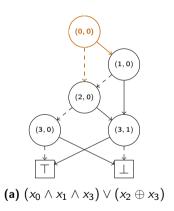




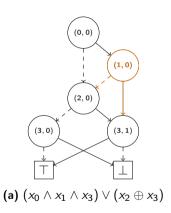


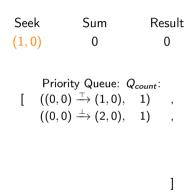


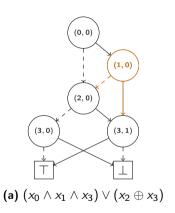
Priority Queue:
$$Q_{count}$$
: [$((0,0) \xrightarrow{\top} (1,0), \quad 1)$, $((0,0) \xrightarrow{\bot} (2,0), \quad 1)$,

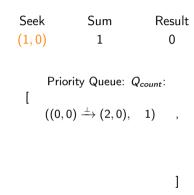


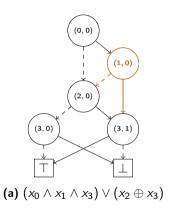
Seek (1, 0		Resul 0	
[Priority Queue: $((0,0) \xrightarrow{\top} (1,0),$ $((0,0) \xrightarrow{\bot} (2,0),$	1)	,
			1

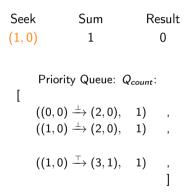


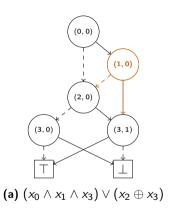


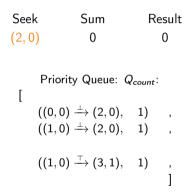


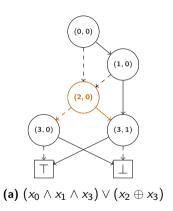


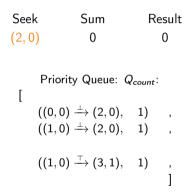


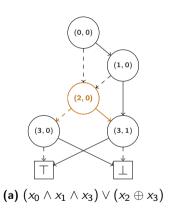


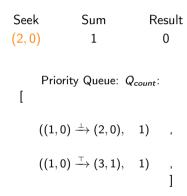


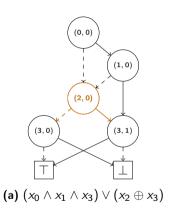


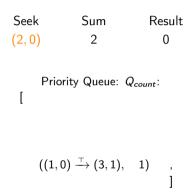


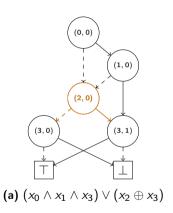




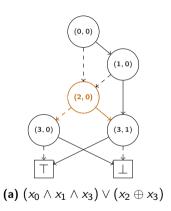




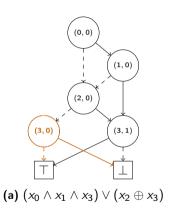




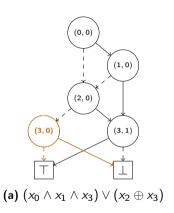
```
Seek
                                   Sum
                                                                      Result
(2,0)
                                        2
                                                                             0
              Priority Queue: Qcount:
             \begin{array}{cccc} ((2,0) \xrightarrow{\bot} (3,0), & 2) & , \\ ((1,0) \xrightarrow{\top} (3,1), & 1) & , \\ ((2,0) \xrightarrow{\top} (3,1), & 2) & ] \end{array}
```

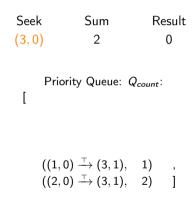


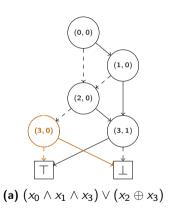
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                                                                      Result
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                                                                             0
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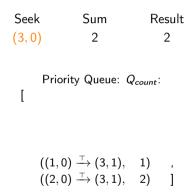


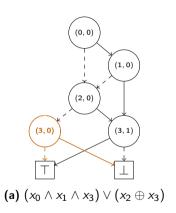
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Seek
                                    Sum
                                                                      Result
(3,0)
                                        0
                                                                             0
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             \begin{array}{cccc} ((2,0) \xrightarrow{\bot} (3,0), & 2) & , \\ ((1,0) \xrightarrow{\top} (3,1), & 1) & , \\ ((2,0) \xrightarrow{\top} (3,1), & 2) & ] \end{array}
```

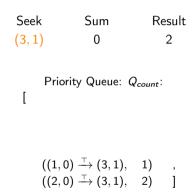


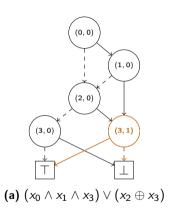


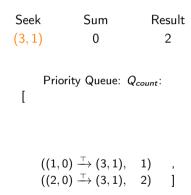


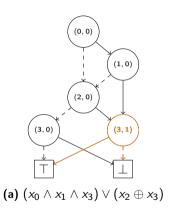


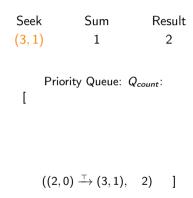


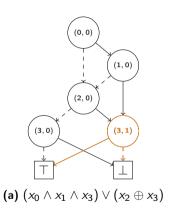


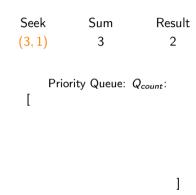


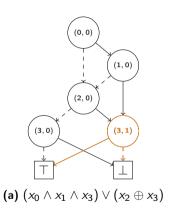


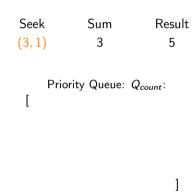




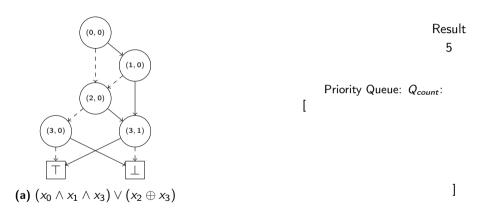








CountPaths : Example



Contents

What are Binary Decision Diagrams?

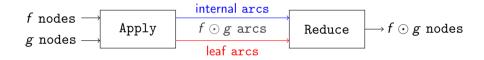
Why do they break?

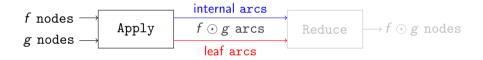
How can we fix it?

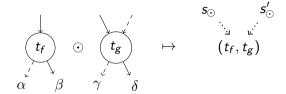
CountPaths

Apply

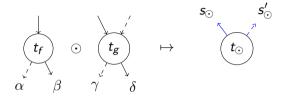
Equality Checking







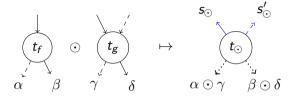
Time-Forward Processing



Observation (semi-tranposition)

 \leftarrow : $s \rightarrow t$ (Internal Arcs) are output at time t and hence sorted by t.

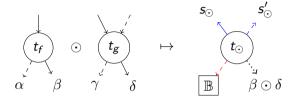
Time-Forward Processing



Observation (semi-tranposition)

 \leftarrow : $s \rightarrow t$ (Internal Arcs) are output at time t and hence sorted by t.

Time-Forward Processing



Observation (semi-tranposition)

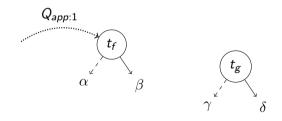
- \leftarrow : $s \rightarrow t$ (Internal Arcs) are output at time t and hence sorted by t.
- \rightarrow : $s \rightarrow \mathbb{B}$ (Terminal Arcs) are output at time s.

Time-Forward Processing

 $Q_{app:1}$: PriorityQueue $\langle (s o (t_f, t_g))
angle$ sorted on $\min(t_f, t_g)$ in ascending order.

Case 1

 $t_f.var() \neq t_g.var()$



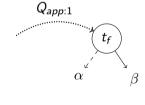
 $Q_{app:1}$: PriorityQueue $\langle (s \rightarrow (t_f, t_g)) \rangle$ sorted on $\min(t_f, t_g)$ in ascending order.

Case 1

 $t_f.var() \neq t_g.var()$

Case 2(a):

 $t_f.var() = t_g.var() \land t_f.id() = t_g.id()$

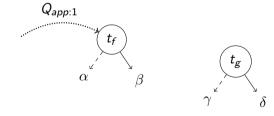




 $Q_{app:1}$: PriorityQueue $\langle (s \rightarrow (t_f, t_g)) \rangle$ sorted on min (t_f, t_g) in ascending order.

Case 1

$$t_f.var() \neq t_g.var()$$



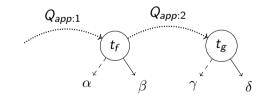
 $Q_{app:2}$: PriorityQueue $\langle (s \to (t_f, t_g), (\alpha, \beta)) \rangle$ sorted on max (t_f, t_g) in ascending order.

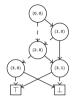
Case 2(a):

$$t_f.var() = t_g.var() \wedge t_f.id() = t_g.id()$$

Case 2(b):

$$t_f.var() = t_g.var() \land t_f.id() \neq t_g.id()$$

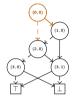




(a) $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$



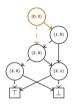
(b)
$$\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$$



(a) $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$



(b) $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$



Priority Queue: Qapp:1:

- [$(0,0) \xrightarrow{\top} ((1,0),(2,1))$,
 - $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$

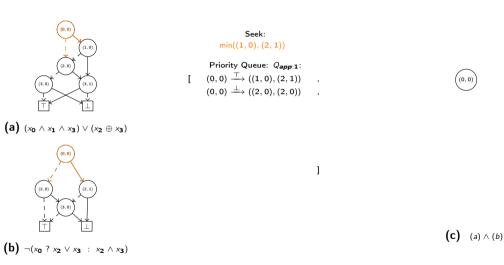
(0,0)

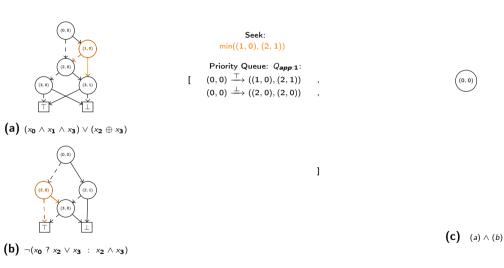
(a) $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$

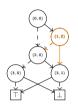


(b) $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$

1







(a)
$$(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$$



(b) $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$

$\begin{array}{c} \text{Seek:} \\ \min((1,0),(2,1)) \end{array}$

Priority Queue: Qapp:1:

 $(0,0) \xrightarrow{\top} ((1,0),(2,1))$

 $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$

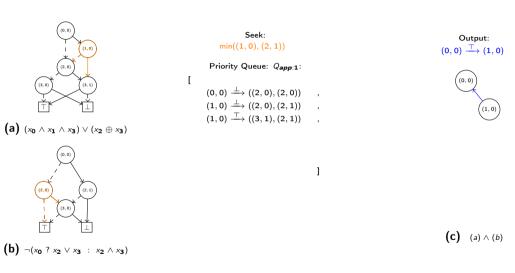
 $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$

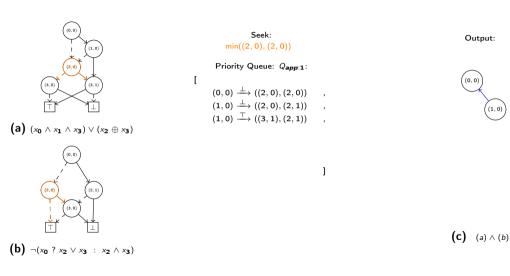
 $(1,0) \xrightarrow{\top} ((3,1),(2,1))$

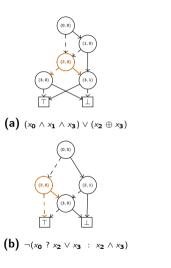
(0,0)

(1,0)

1

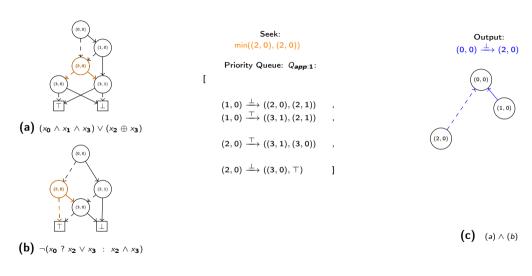


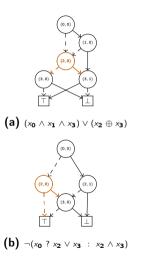




Seek: min((2,0),(2,0))Priority Queue: Qapp:1: $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$ $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$ $(1,0) \xrightarrow{\top} ((3,1),(2,1))$ (2,0) $(2,0) \xrightarrow{\top} ((3,1),(3,0))$ $(2,0) \xrightarrow{\perp} ((3,0),\top)$]

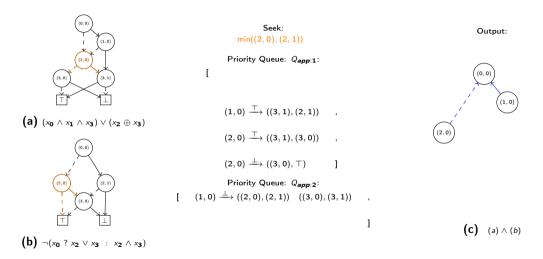


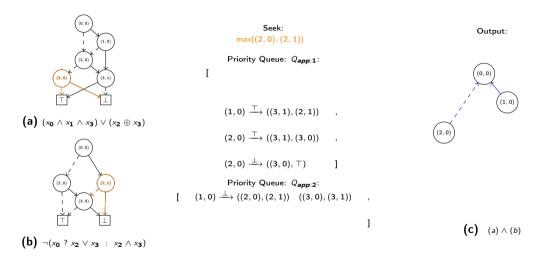


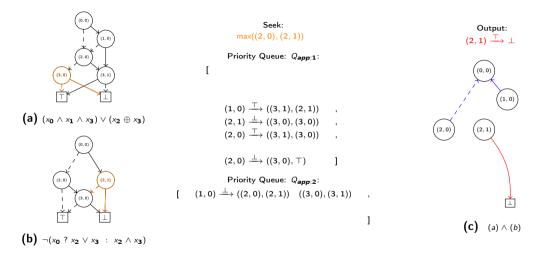


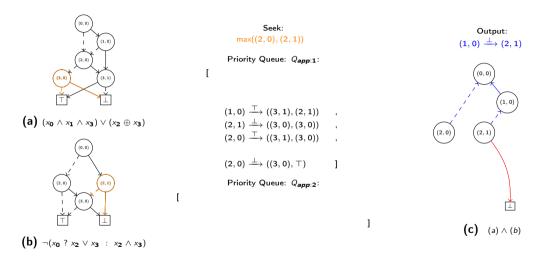
Seek: min((2,0),(2,1))Priority Queue: Qapp:1: $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$ $(1,0) \xrightarrow{\top} ((3,1),(2,1))$ $(2,0) \xrightarrow{\top} ((3,1),(3,0))$ $(2,0) \xrightarrow{\perp} ((3,0),\top)$]

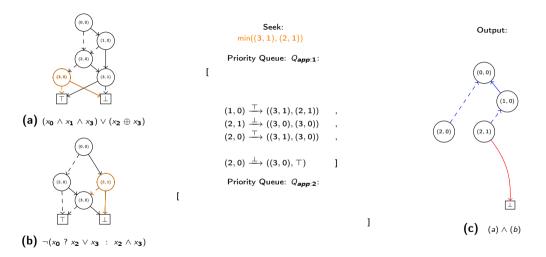
Output: (0,0) (1,0)

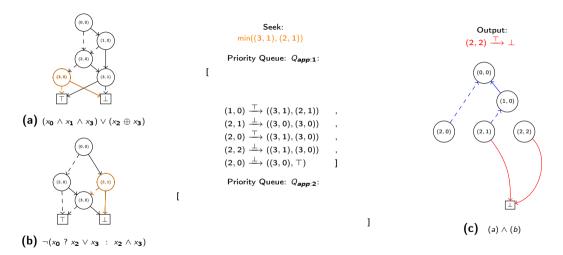


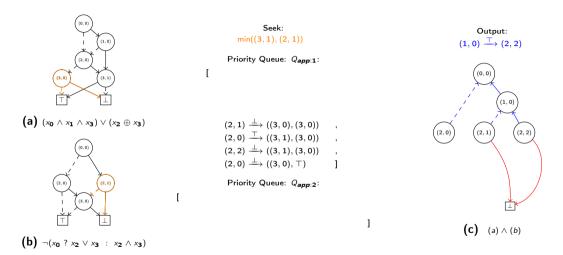


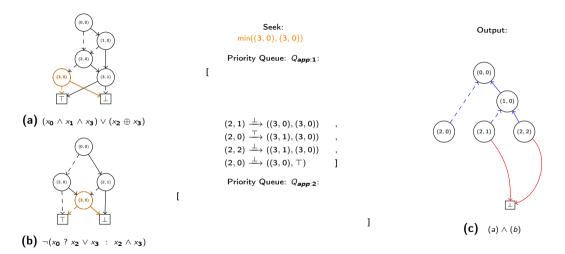


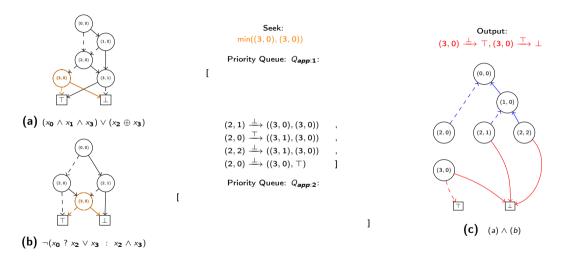


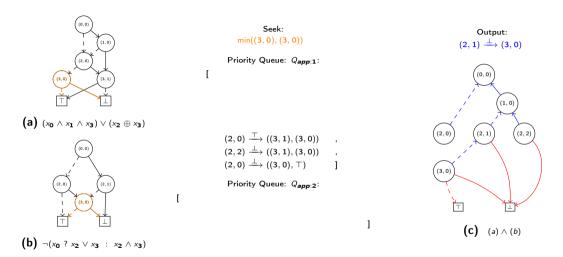


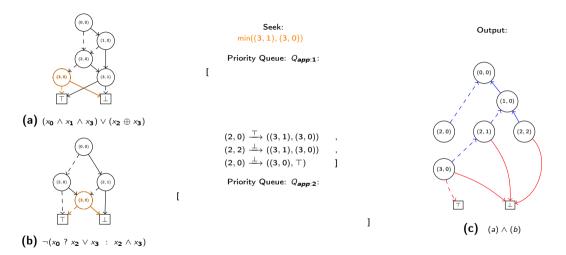


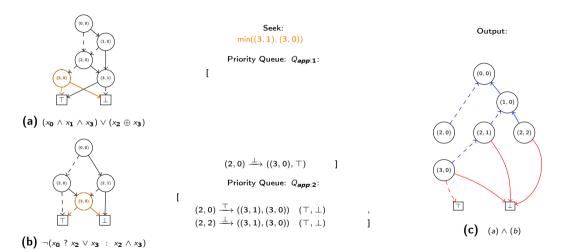


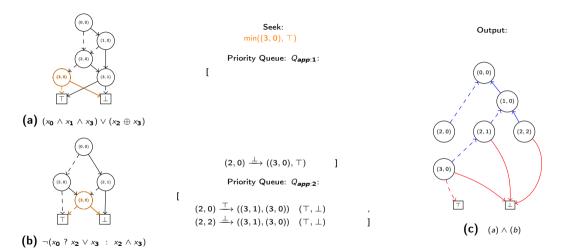


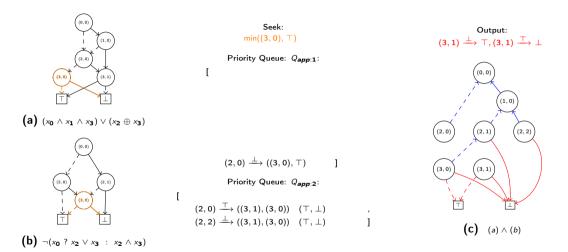


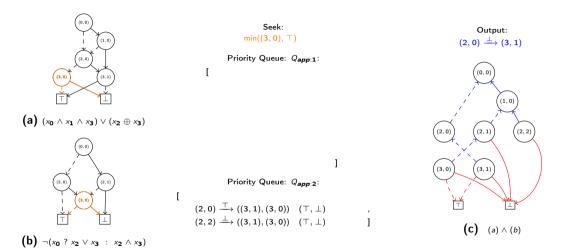


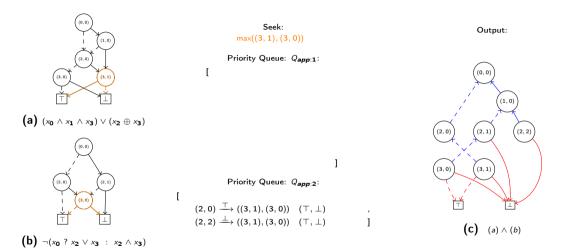


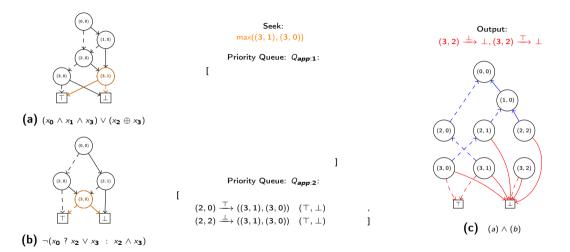


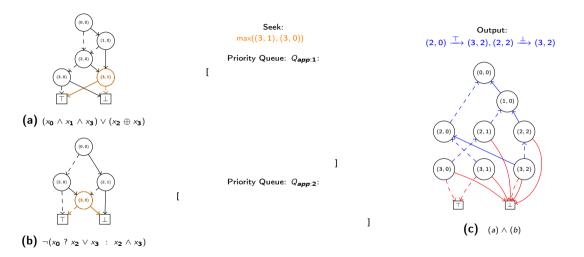


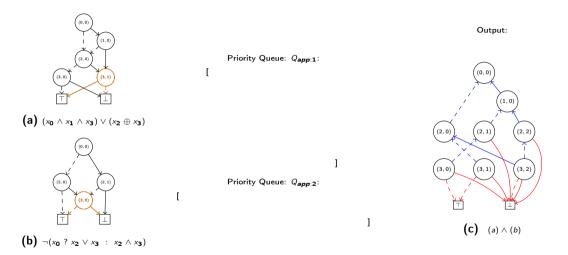




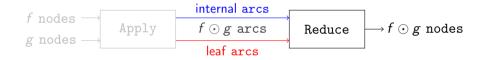


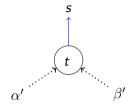






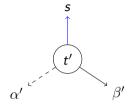
Apply





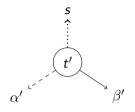
Time-Forward Processing

Send reduction t' with Q_{red} : PriorityQueue $\langle (s o t')
angle$ descending on parent s.



Time-Forward Processing

Send reduction t' with Q_{red} : PriorityQueue $\langle (s o t')
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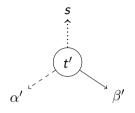


Time-Forward Processing

Send reduction t' with Q_{red} : PriorityQueue $\langle (s o t')
angle$ descending on parent s.

Observation (semi-tranposition)

 \leftarrow : $s \rightarrow t$ (Internal Arcs) provide parents of unreduced node t.

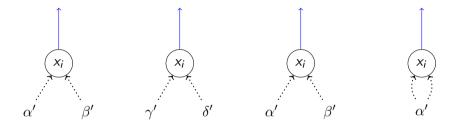


Time-Forward Processing

Send reduction t' with Q_{red} : PriorityQueue $\langle (s o t')
angle$ descending on parent s.

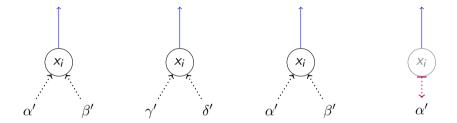
Observation (semi-tranposition)

- \leftarrow : $s \rightarrow t$ (Internal Arcs) provide parents of unreduced node t.
- ightarrow: $s
 ightarrow\mathbb{B}$ (Terminal Arcs) are reduced and already sorted as per Q_{red} .



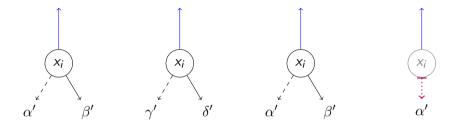
Reduce Level i:

1 Obtain nodes from Q_{red} and terminal arcs. Filter and remember redundant nodes.



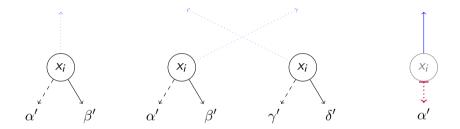
Reduce Level i:

1 Obtain nodes from Q_{red} and terminal arcs. Filter and remember redundant nodes.

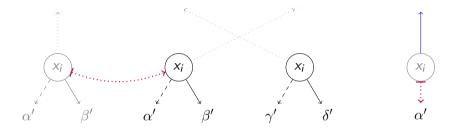


Reduce Level i:

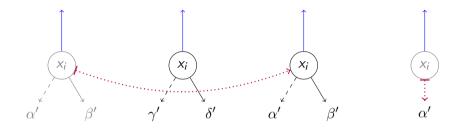
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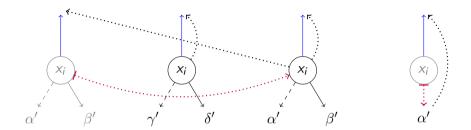
- 1 Obtain nodes from Q_{red} and terminal arcs. Filter and remember redundant nodes.
- Sort remaining nodes by children, output unique nodes, and remember duplications.



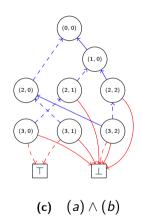
- 1 Obtain nodes from Q_{red} and terminal arcs. Filter and remember redundant nodes.
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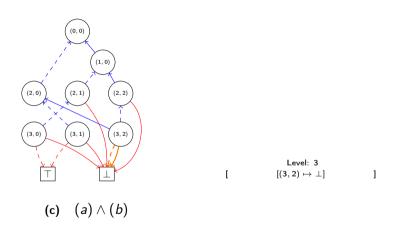


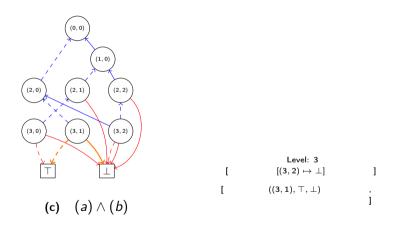
- 1 Obtain nodes from Q_{red} and terminal arcs. Filter and remember redundant nodes.
- Sort remaining nodes by children, output unique nodes, and remember duplications.
- 3 Sort back to match internal arcs and forward to parents with Q_{red} .

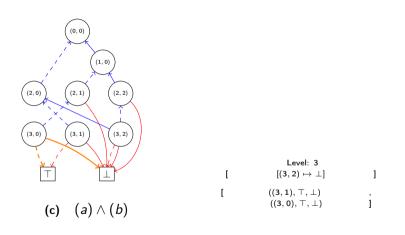


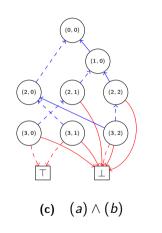
- 1 Obtain nodes from Q_{red} and terminal arcs. Filter and remember redundant nodes.
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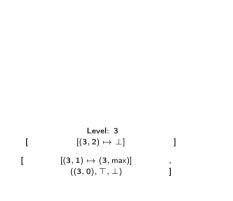




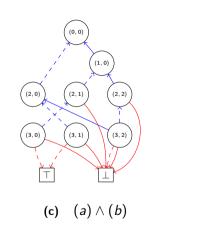


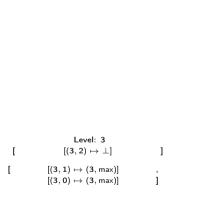


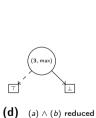


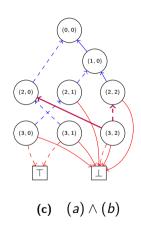


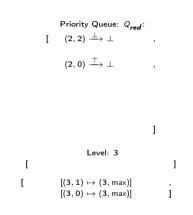






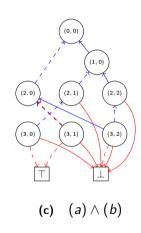


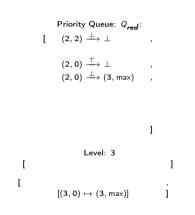


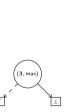




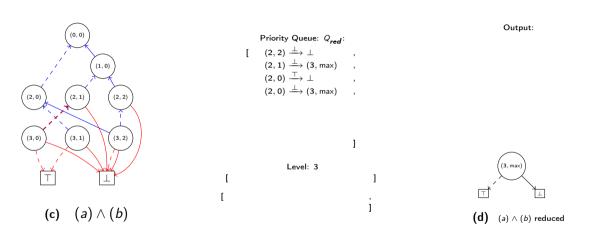


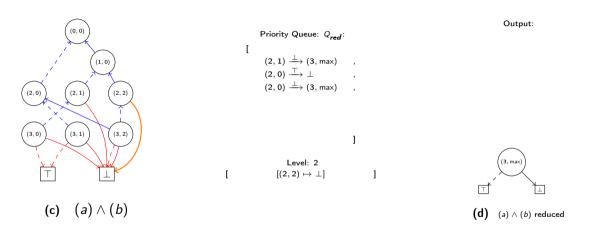


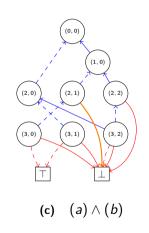


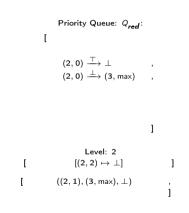


(d) $(a) \wedge (b)$ reduced

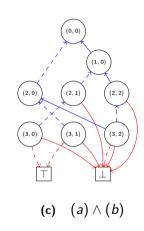


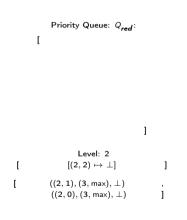




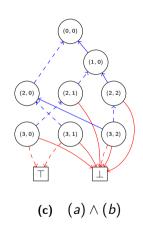


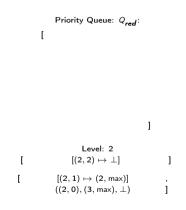


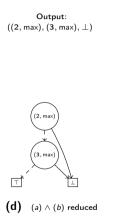


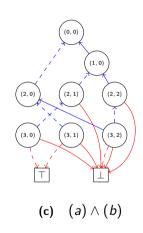


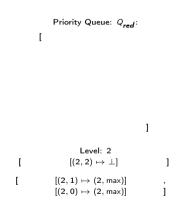


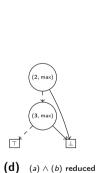


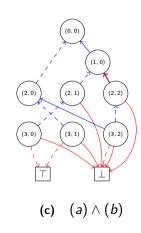


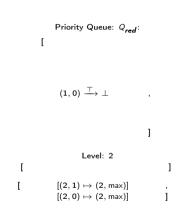


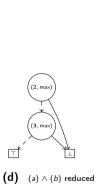


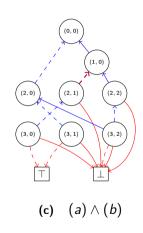


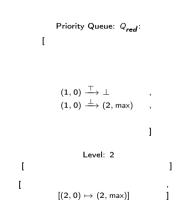


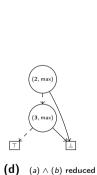


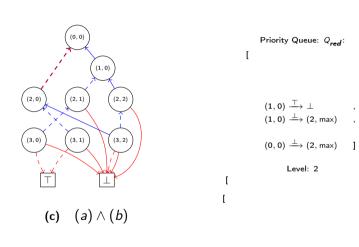


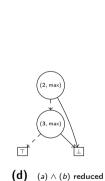


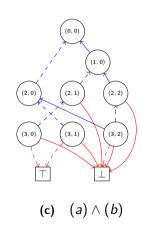


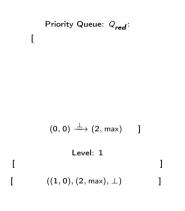


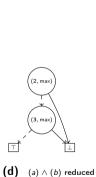


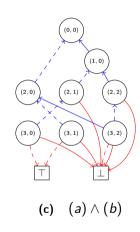


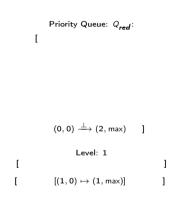


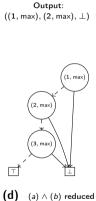


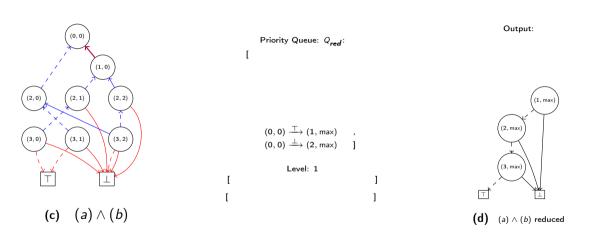


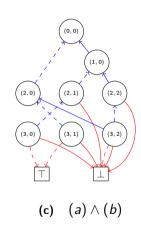


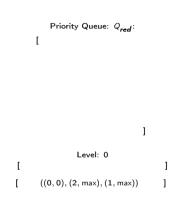


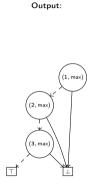




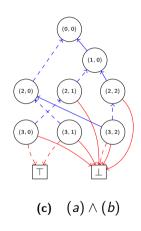


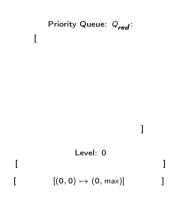


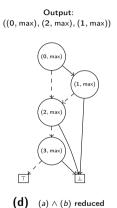


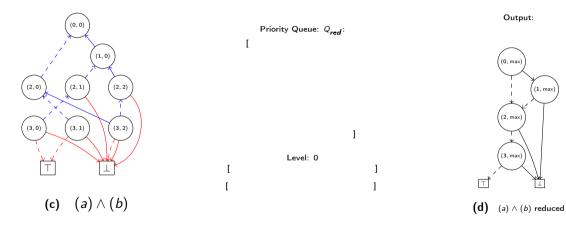


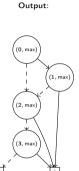
(d) $(a) \wedge (b)$ reduced



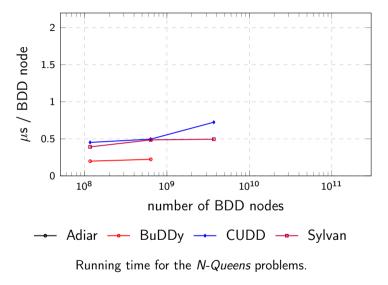


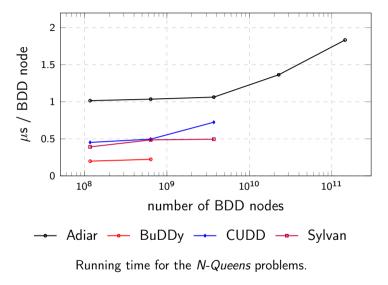






I/O-Complexity
$O(\operatorname{sort}(N_f))$
$2N_f/B$
$O(\operatorname{sort}(N_f))$
$O(\operatorname{sort}(N_f \cdot N_g))$





Contents

What are Binary Decision Diagrams?

Why do they break?

How can we fix it?

CountPaths

Apply

Equality Checking

Algorithm	I/O-Complexity
bdd_pathcount	$O(\operatorname{sort}(N_f))$
bdd_not	$2N_f/B$
bdd_restrict	$O(\operatorname{sort}(N_f))$
bdd_apply	$O(\operatorname{sort}(N_f \cdot N_g))$

Algorithm	I/O-Complexity
bdd_pathcount	$O(\operatorname{sort}(N_f))$
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bdd_equal	?

$$f\leftrightarrow g\equiv \top$$

$$f \leftrightarrow g \equiv \top$$

$$\underbrace{O(\mathsf{sort}(\mathit{N}^2))}_{\mathsf{Apply}} + \underbrace{O(\mathsf{sort}(\mathit{N}^2))}_{\mathsf{Reduce}} + \underbrace{O(1))}_{\mathsf{check is }\top} = O(\mathsf{sort}(\mathit{N}^2))$$

Theorem (Bryant '86)

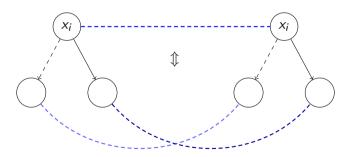
Theorem (Bryant '86)

Let π be a variable order and $f: \mathbb{B}^n \to \mathbb{B}$ then there exists a unique (up to isomorphism) Reduced Ordered Binary Decision Diagram representing f with ordering π .

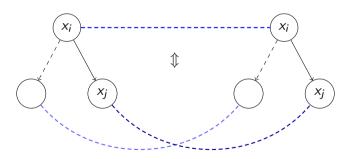
Trivial cases: $f \not\equiv g$ if there is a mismatch in

	$N_f eq N_g$	Number of nodes	O(1) I/Os
•	$L_f eq L_g$	Number of levels	O(1) I/Os
•	$N_{f,i} \neq N_{g,i}$	Number of nodes on a level	O(L/B) I/Os
•	$L_{f,i} \neq L_{g,i}$	Label of an <i>i</i> th level	O(L/B) I/Os

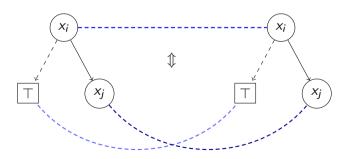
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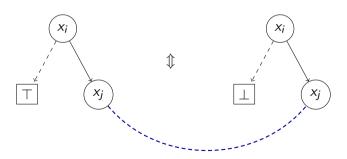
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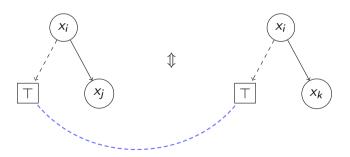
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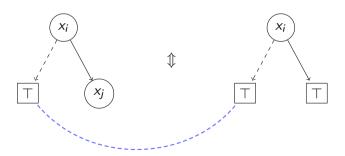
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IsIsomorphic(f, g)

- Check whether root v_f of f and root v_g of g have a local violation.
- Check $low(v_f) \sim low(v_g)$ and $high(v_f) \sim high(v_g)$ "recursively".

Return false on first violation. If there are no violations then return true.

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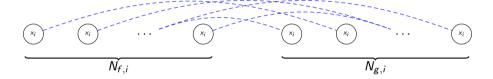
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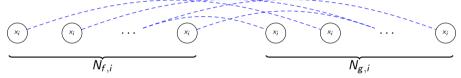
$$\underbrace{O(\mathsf{sort}(\mathit{N}^2))}_{\mathtt{Apply'}} + \underbrace{O(\mathsf{sort}(\mathit{N}^2))}_{\mathtt{Reduce}} + \underbrace{O(1))}_{\mathtt{check is }\top} = O(\mathsf{sort}(\mathit{N}^2))$$

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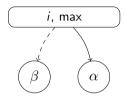
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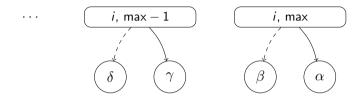
Return false if more than $N_{f,i} = N_{g,i}$ pairs of nodes are checked on level i.

$$\underbrace{O(\mathsf{sort}(\Sigma_i \ \mathsf{N}_{f,i}))}_{\mathsf{Apply''}} = O(\mathsf{sort}(\mathsf{N}))$$

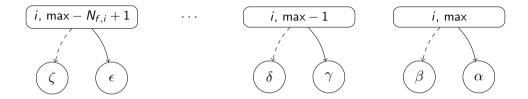
Observation



Observation



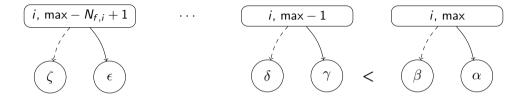
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Observation

Each level output by the Reduce algorithm has the following properties:

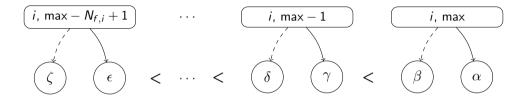
■ Nodes on level *i* have their identifiers *consecutively* numbered.



Observation

Each level output by the Reduce algorithm has the following properties:

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Observation

- Nodes on level *i* have their identifiers *consecutively* numbered.
- Nodes on level *i* are output sorted by their children.

Theorem

If G_f and G_g are outputs of Reduce.

 $G_f \sim G_g \iff For \ all \ i \in [0; N_f) \ the \ node \ G_f[i] \ matches \ G_g[i] \ numerically.$

Proof.

← : Must describe the exact same graph.

 \Rightarrow : Strong induction on BDD levels bottom-up.

Theorem

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Corollary

If G_f and G_g are outputs of Reduce then $f \equiv g$ is computable using $2 \cdot N/B$ I/Os.

$$\begin{array}{c|c} & \text{Algorithm} & \text{Time (s)} \\ \hline f \leftrightarrow g \equiv \top & 0.38 \end{array}$$

Checking the (EPFL Benchmark) voter circuit's single output gate ($|N_f| = |N_g| = 5.76$ MiB).

Algorithm Time (s)
$$f \leftrightarrow g \equiv \top \quad 0.38$$

$$O(\operatorname{sort}(N)) \quad 0.058$$

Checking the (EPFL Benchmark) voter circuit's single output gate ($|N_f| = |N_g| = 5.76$ MiB).

Algorithm	Time (s)
$f\leftrightarrow g\equiv \top$	0.38
O(sort(N))	0.058
2N/B	0.006

Checking the (EPFL Benchmark) voter circuit's single output gate ($|N_f| = |N_g| = 5.76$ MiB).

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■ soelvsten@cs.au.dk

@ssoelvsten

Adiar

github.com/ssoelvsten/adiar

ssoelvsten.github.io/adiar



Algorithm	Depth-First	Time-Forwared
bdd_pathcount	$O(N_f)$	$O(\operatorname{sort}(N_f))$
bdd_not	$O(N_f)$	$2N_f/B$
bdd_restrict	$O(N_f)$	$O(\operatorname{sort}(N_f))$
bdd_apply	$O(N_f N_g)$	$O(\operatorname{sort}(N_f N_g))$
bdd_equal	O(1)	2 N /B