## Efficient External Memory Algorithms for Binary Decision Diagram Manipulation

**Steffan Christ Sølvsten**, Jaco van de Pol, Anna Blume Jakobsen, and Mathias Weller Berg Thomasen November 26, 2021



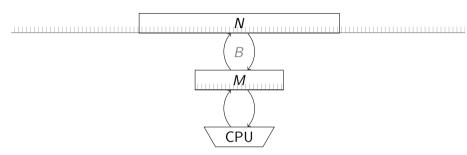


Figure 1: The I/O model by Aggarwal and Vitter '87

For any realistic values of N, M, and B we have that

$$N/B < \operatorname{sort}(N) \triangleq N/B \cdot \log_{M/B} N/B \ll N$$
,

Theorem (Aggarwal and Vitter '87) N elements can be sorted in  $\Theta(sort(N))$  I/Os.

## Theorem (Arge '95)

N elements can be inserted in and extracted from a Priority Queue in  $\Theta(sort(N))$  I/Os.

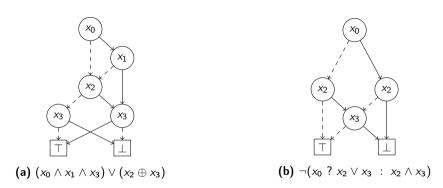
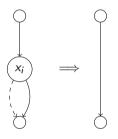


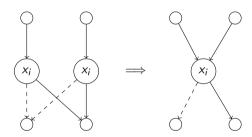
Figure 2: Examples of (Reduced Ordered) Binary Decision Diagrams.

## Theorem (Bryant '86)

For a fixed variable order, if one exhaustively applies the two rules below, then one obtains the Reduced OBDD, which is a unique canonical form of the function.



(a) Rule 1: Remove redundant nodes



(b) Rule 2: Merge duplicate nodes

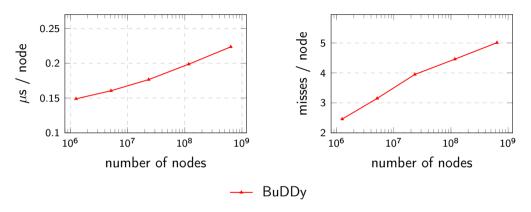
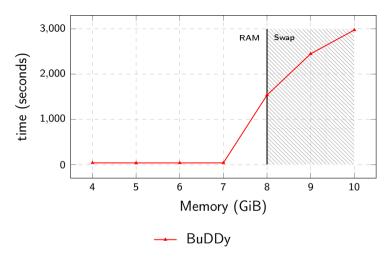


Figure 4: Cache behaviour for the *N*-Queens problem.



**Figure 5:** Running time for *Tic-Tac-Toe* with N = 21.

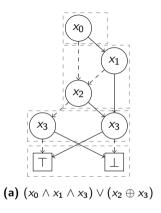


Figure 6: Blocks active in memory

$$M = 4$$
,  $B = 2$ 

node I/Os cache lookups

0 0

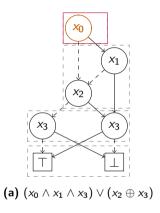


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node I/Os cache lookups

1

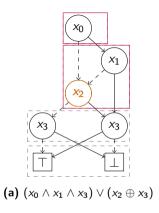


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$$M = 4$$
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node I/Os cache lookups

2

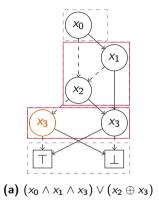


Figure 6: Blocks active in memory

$$M = 4$$
,  $B = 2$ 
node I/Os cache lookups
 $3$ 

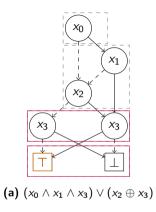


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$$M = 4$$
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node I/Os cache lookups
$$4 3$$

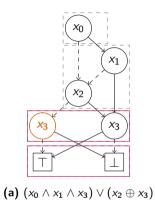


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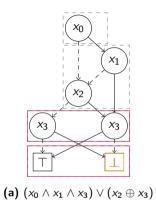


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$$M = 4$$
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node I/Os cache lookups
$$4 3$$

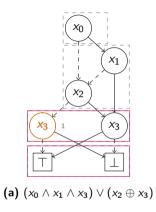


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$$M = 4$$
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node I/Os cache lookups
$$4 3$$

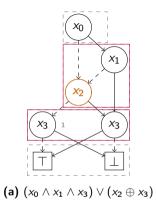


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,  $B = 2$ 

node I/Os cache lookups

 $3$ 

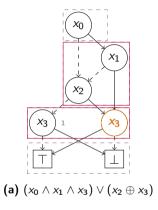


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$$M = 4$$
,  $B = 2$ 
node I/Os cache lookups
$$5 4$$

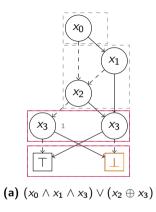


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$$M = 4$$
,  $B = 2$ 

node I/Os cache lookups

6 4

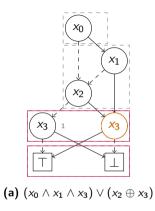


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$$M = 4$$
,  $B = 2$ 

node I/Os cache lookups

6 4

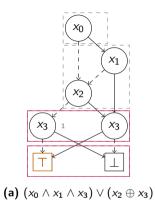


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$$M = 4$$
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node I/Os cache lookups

6 4

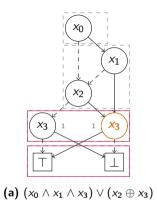


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$$M = 4$$
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6 4

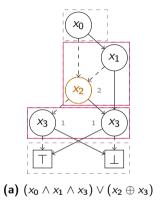


Figure 6: Blocks active in memory

$$M = 4$$
,  $B = 2$ 

node I/Os cache lookups

7

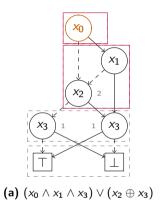


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$$M = 4$$
,  $B = 2$ 
node I/Os cache lookups
8

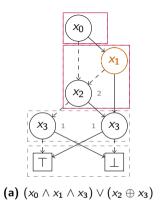


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$$M = 4$$
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node I/Os cache lookups

8 5

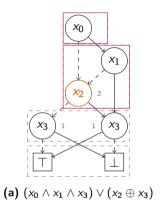


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node I/Os cache lookups

8 6

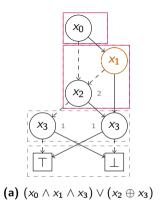


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node I/Os cache lookups

8 6

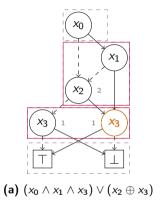


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7

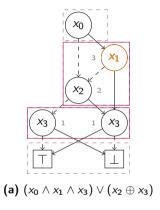


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7

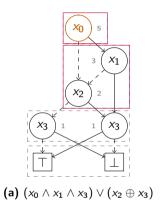


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$$M = 4$$
,  $B = 2$ 
node I/Os cache lookups
$$10 7$$

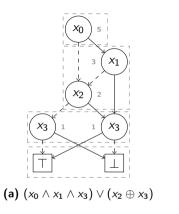


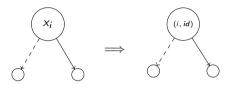
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node I/Os cache lookups

7

Let every node be uniquely identified by a tuple (label, id) :  $\mathbb{N} \times \mathbb{N}$ .



Nodes are ordered based on their *uid* as follows

$$(i_1, id_1) < (i_2, id_2) \equiv i_1 < i_2 \lor (i_1 = i_2 \land id_i < id_j)$$

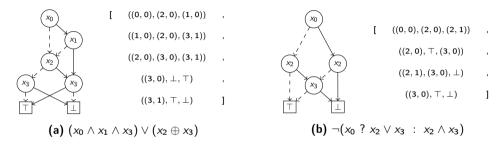


Figure 7: Node-based representation of prior shown BDDs

## CountPaths Example

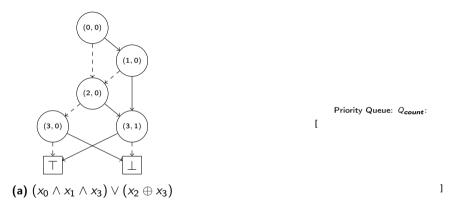


Figure 8: In-order traversal of BDD

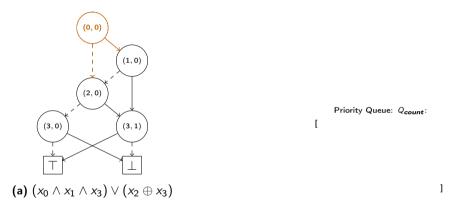


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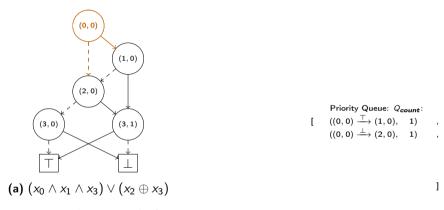


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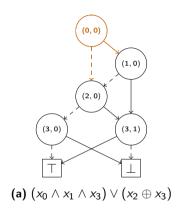


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Seek Sum Result 
$$(1,0)$$
 0 0  $0$ 

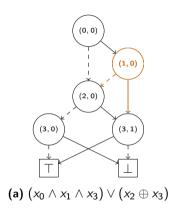
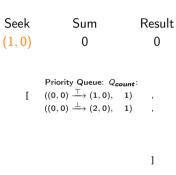


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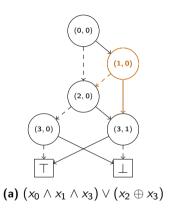
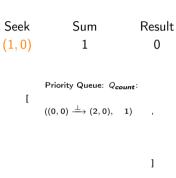


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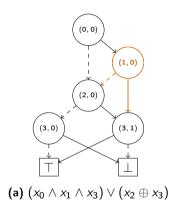
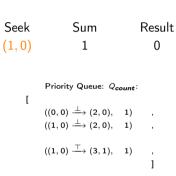


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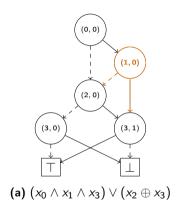
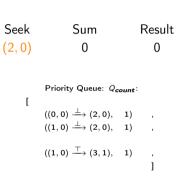


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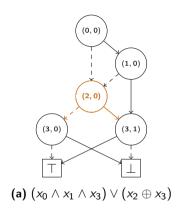
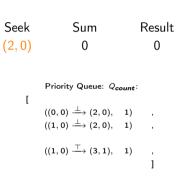


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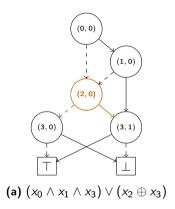
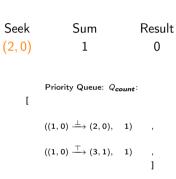


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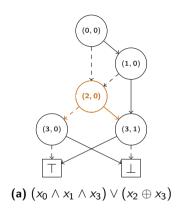
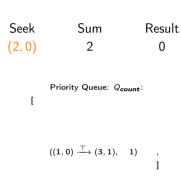


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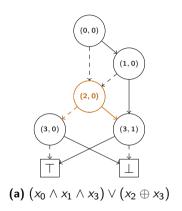
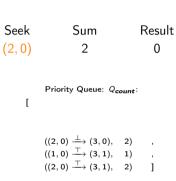


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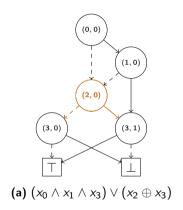
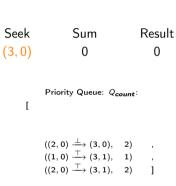


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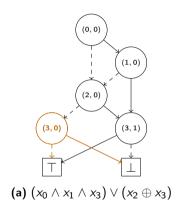
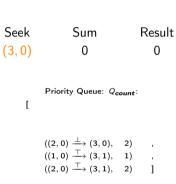


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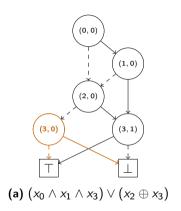
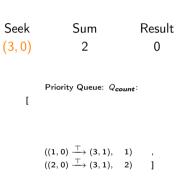


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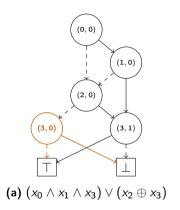
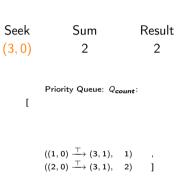


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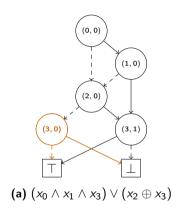
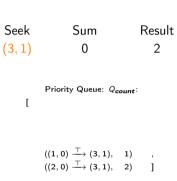


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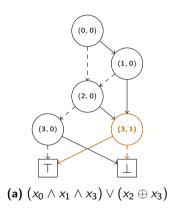
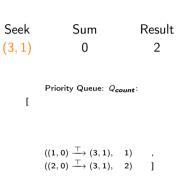


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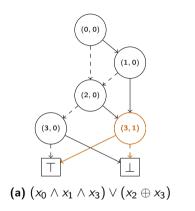
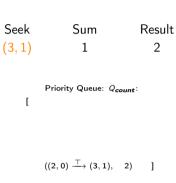


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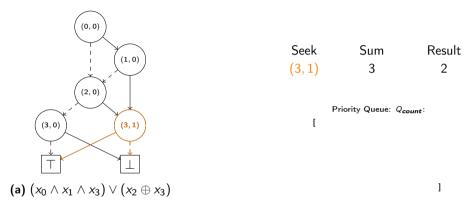


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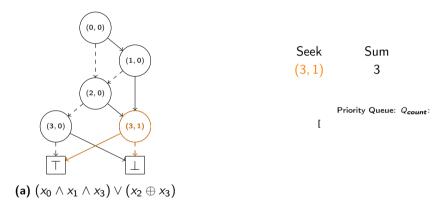
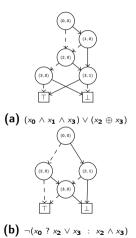
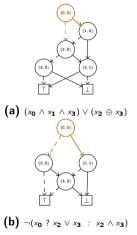


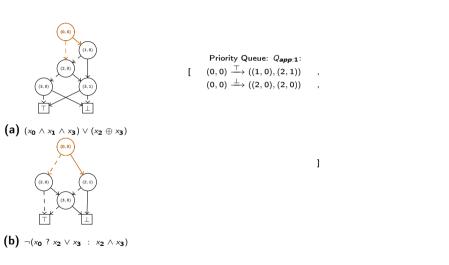
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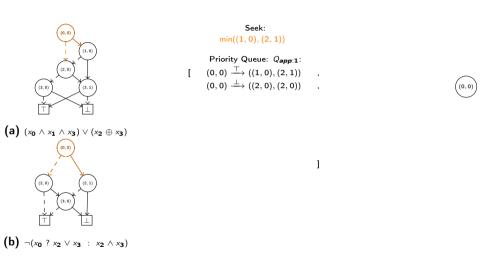
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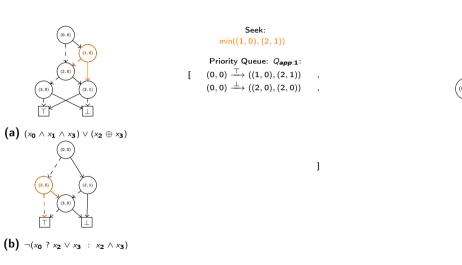
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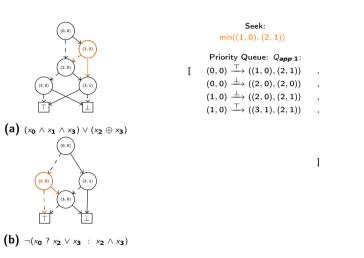




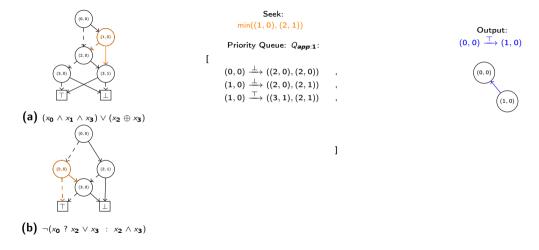


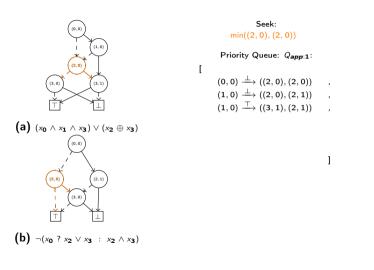






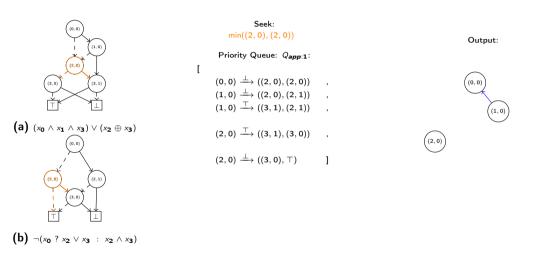
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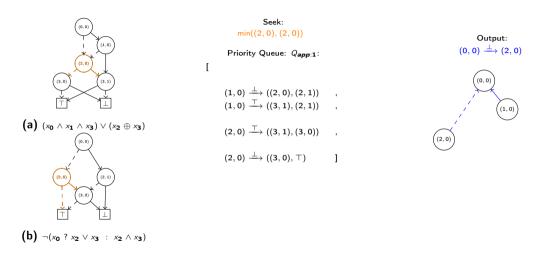


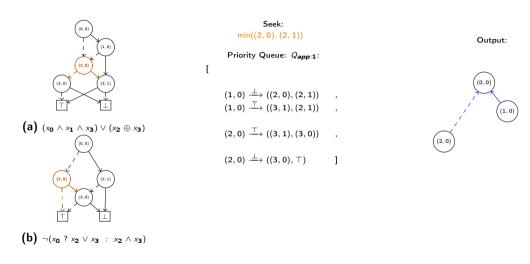


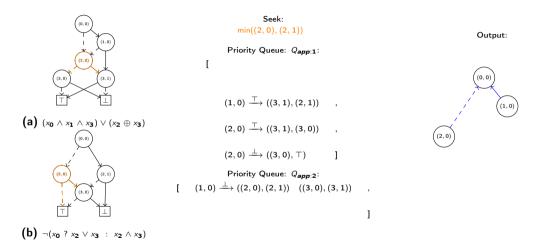
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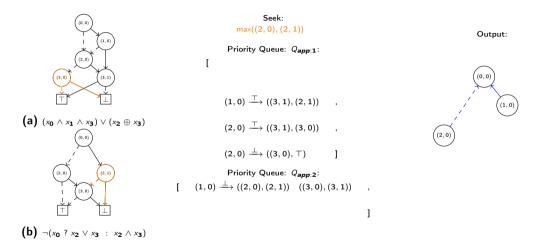


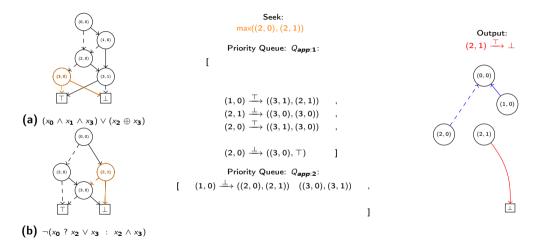


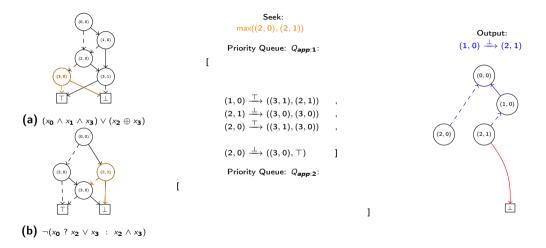


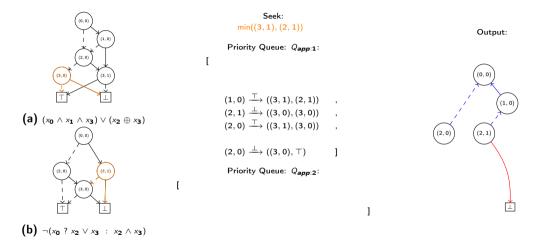


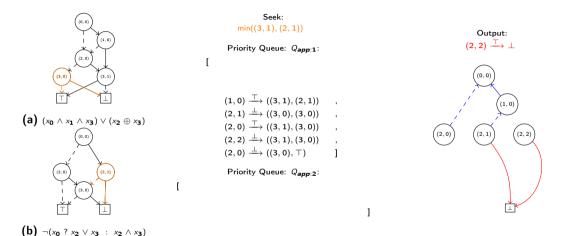


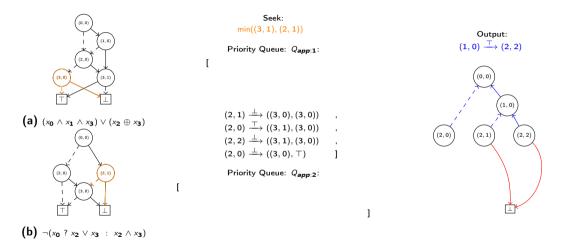


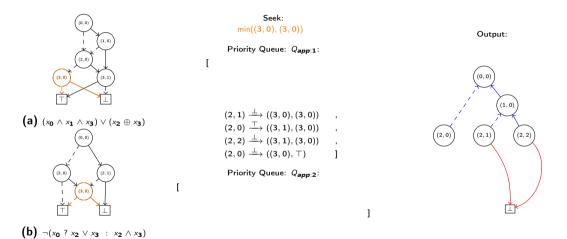


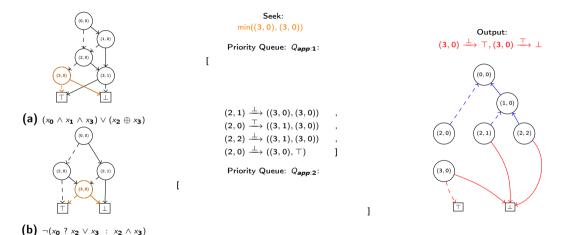


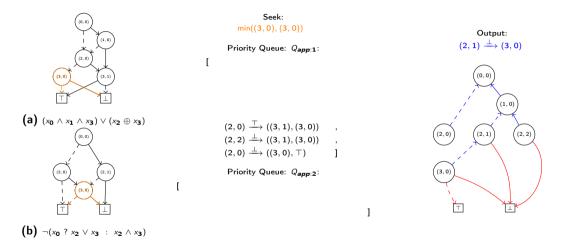


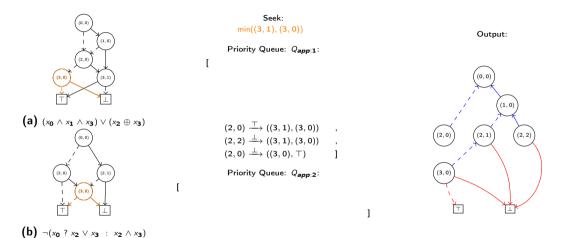


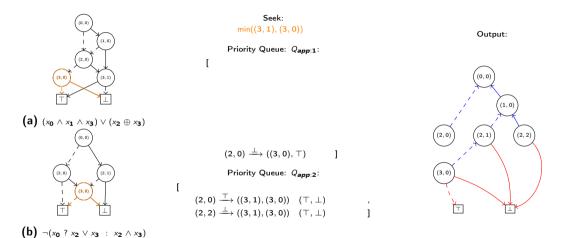


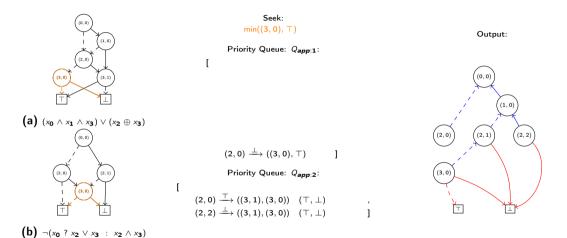


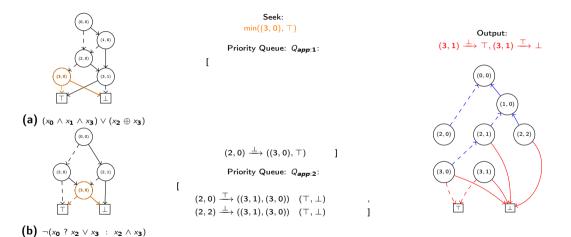


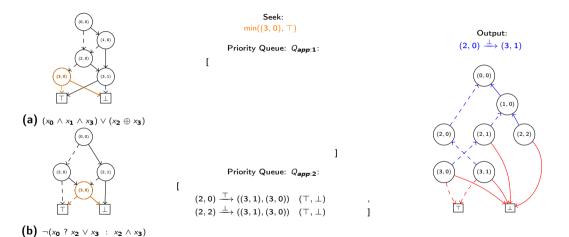


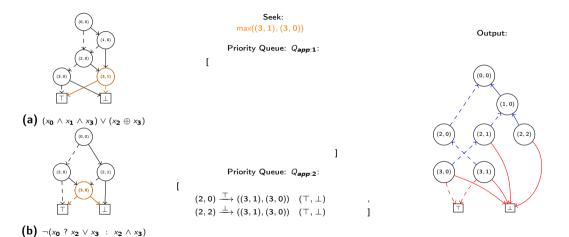


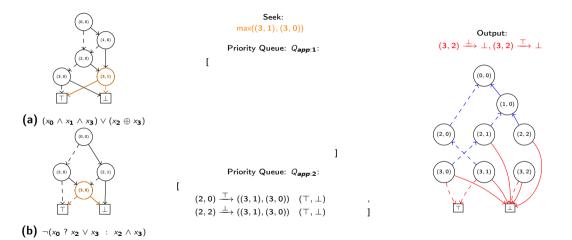


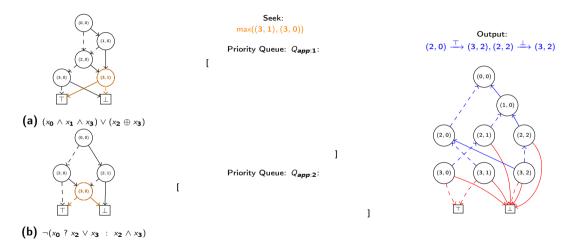


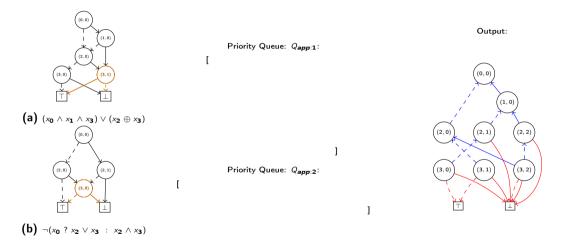


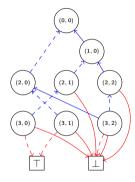


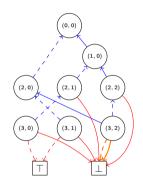




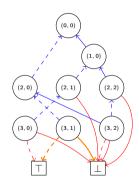




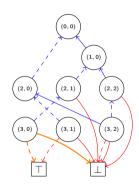




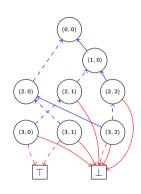
Level: 3 [ $(3,2)\mapsto \bot$ ]

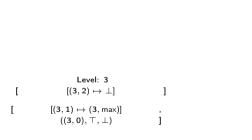


```
Level: 3  [(3,2) \mapsto \bot]   [(3,1), \top, \bot)
```

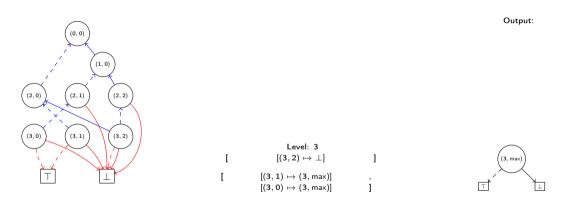


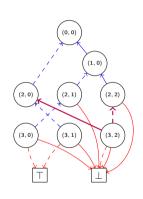
```
Level: 3  [ (3,2) \mapsto \bot ]   [ ((3,1), \top, \bot) \\ ((3,0), \top, \bot) ]
```

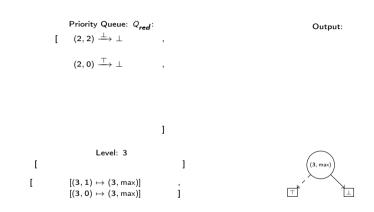


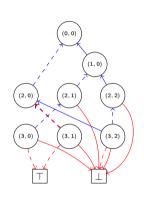


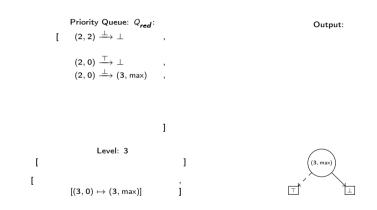


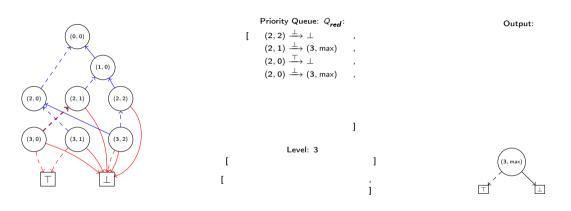


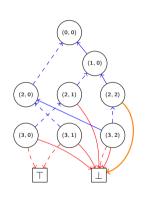


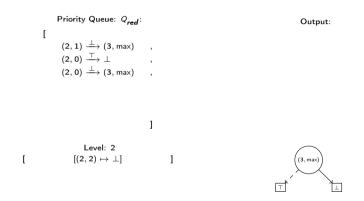


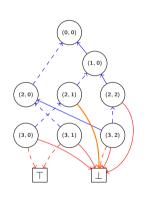


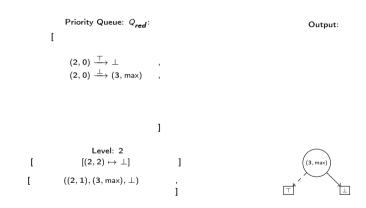


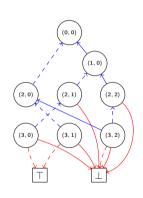


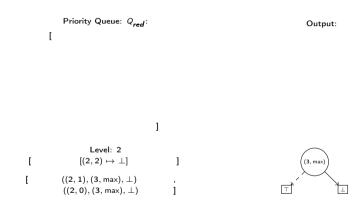


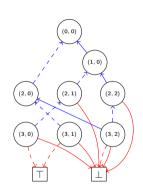


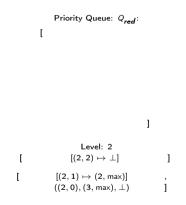


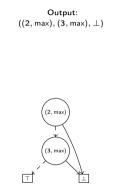


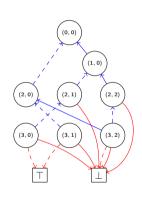


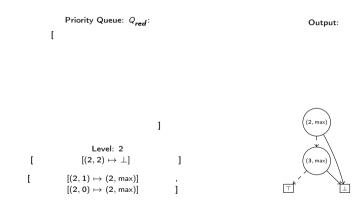


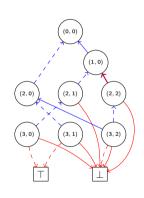


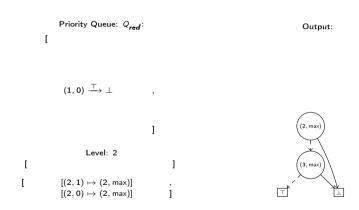


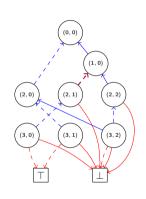


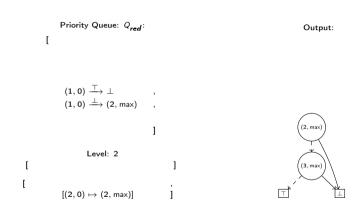


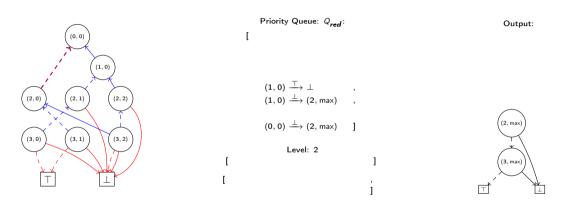


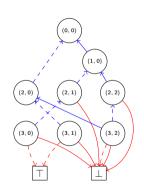


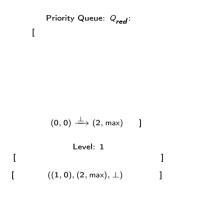


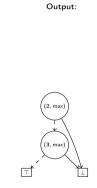


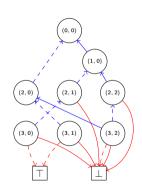


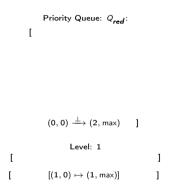


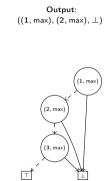


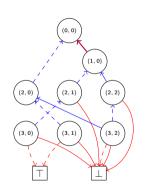


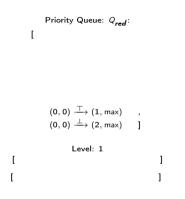


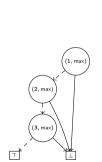




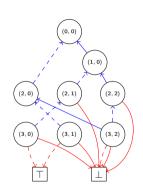


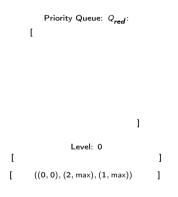


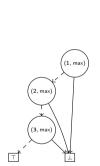




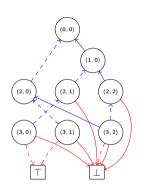
Output:

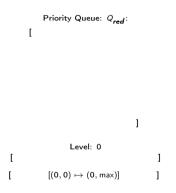


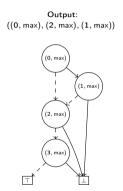




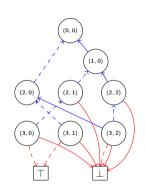
Output:

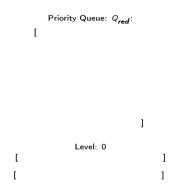


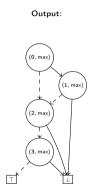




## Reduce Example







Level: 1 Level: 2 Level: 3

$$\left[ (0,0) \xrightarrow{\top} ((1,0),(2,1)) \right]$$

$$[ (0,0) \xrightarrow{\perp} ((2,0),(2,0)) ,$$

Level: 1

$$\left[ (0,0) \xrightarrow{\top} ((1,0),(2,1)) \right]$$

Level: 2

$$\left[ (0,0) \xrightarrow{\perp} ((2,0),(2,0)) \quad , \quad (1,0) \xrightarrow{\perp} ((2,0),(2,1)) \quad , \quad (1,0) \xrightarrow{\top} ((3,1),(2,1)) \quad \right]$$

Level: 1

Level: 2

 $\left[ \quad (0,0) \xrightarrow{\bot} ((2,0),(2,0)) \quad , \quad (1,0) \xrightarrow{\bot} ((2,0),(2,1)) \quad , \quad (1,0) \xrightarrow{\top} ((3,1),(2,1)) \quad \right]$ 

Level: 1

$$\left[ \begin{array}{ccc} (2,0) \xrightarrow{\bot} ((3,0),\top) & , & (2,0) \xrightarrow{\top} ((3,1),(3,0)) & , \end{array} \right.$$

Level: 1

1

Level: 2

,  $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$  ,  $(1,0) \xrightarrow{\top} ((3,1),(2,1))$ 

Level: 3

 $\left[ \begin{array}{ccc} (2,0) \stackrel{\bot}{\longrightarrow} ((3,0),\top) & , & (2,0) \stackrel{\top}{\longrightarrow} ((3,1),(3,0)) \end{array} \right. ,$ 

Level: 1

Level: 2

$$\left[\begin{array}{cccc} & & & & & & \\ & & & & & \\ & & & & \\ & & & & \end{array}, \quad (1,0) \xrightarrow{\top} ((3,1),(2,1)) \quad \right]$$

Level: 1

.

Level: 2

,  $(1,0)\stackrel{ op}{\longrightarrow} ((3,1),(2,1))$ 

Level: 3

Level: 2

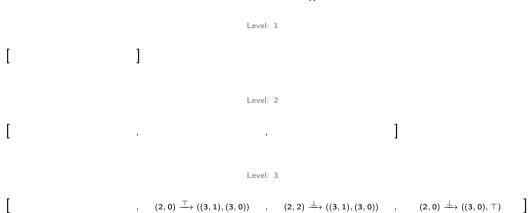
 $(1,0) \xrightarrow{\top} ((3,1),(2,1))$ 

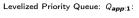
Level: 3

$$(2,0) \xrightarrow{\perp} ((3,0),\top$$

Level: 1 Level: 2 Level: 3

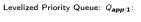
Level: 1 Level: 2 Level: 3 





Level: 1 Level: 2 Level: 3

,  $(2,2)\stackrel{\perp}{\longrightarrow} ((3,1),(3,0))$  ,  $(2,0)\stackrel{\perp}{\longrightarrow} ((3,0),\top)$ 



Level: 1 Level: 2 Level: 3

,  $(2,0) \xrightarrow{\perp} ((3,0),\top)$ 

Level: 1 Level: 2 Level: 3

## Memory layout and efficient sorting

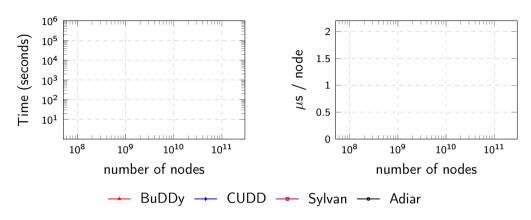
The unique identifier of nodes and leafs can be represented in a single 64-bit integer.



The f bit-flag is used to store the *is\_high* boolean inside of the source of an arc.

# Adiar

github.com/ssoelvsten/adiar



**Figure 11:** Minimum running times for the *N*-Queens problem.

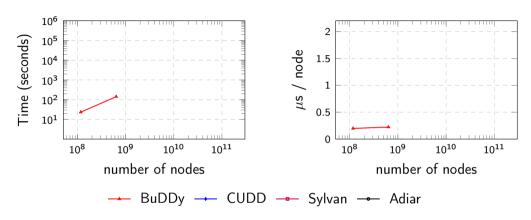
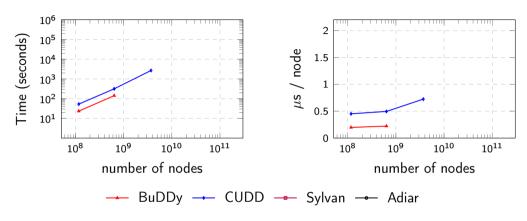
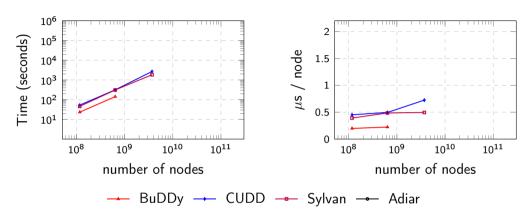


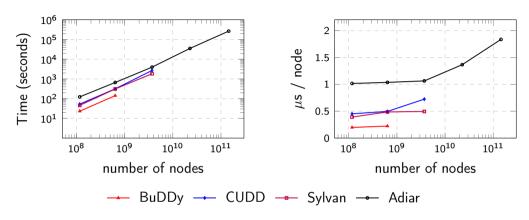
Figure 11: Minimum running times for the *N*-Queens problem.



**Figure 11:** Minimum running times for the *N*-Queens problem.



**Figure 11:** Minimum running times for the *N*-Queens problem.



**Figure 11:** Minimum running times for the *N*-Queens problem.

Algorithm		Depth-first	Time-forwarded		
Reduce		O(N)	$O(\operatorname{sort}(N))$		
BDD Manipulation					
Apply	f ⊙ g	$O(N_f \cdot N_g)$	$O(\operatorname{sort}(N_f \cdot N_g))$		
If-Then-Else	f ? g : h	$O(N_f \cdot N_g \cdot N_h)$	$O(\operatorname{sort}(N_f \cdot N_g \cdot N_h))$		
Restrict	$f _{x_i=v}$	O(N)	$O(\operatorname{sort}(N))$		
Negation	$\neg f$	O(1)	O(1)		
Quantification	$\exists / \forall v : f _{x_i = v}$	$O(N^2)$	$O(\operatorname{sort}(N^2))$		
Counting					
Count Paths	#paths in $f$ to $ op$	O(N)	$O(\operatorname{sort}(N))$		
Count SAT	#x:f(x)	O(N)	$O(\operatorname{sort}(N))$		
Other					
Equality	$f \equiv g$	O(1)	$O(\operatorname{sort}(N))$		
Evaluate	f(x)	O(L)	O(N/B)		
Min/Max SAT	$\min / \max\{x \mid f(x)\}$	O(L)	O(N/B)		

 $\textbf{Table 1:} \ I/O\text{-}complexity of depth-first algorithms compared to our time-forwarded}.$ 

