Experiment 10

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BATCH: K1

(Backtracking)

Aim: Implementation of Sum of subsets.

Theory:

Subset sum problem is the problem of finding a subset such that the sum of elements equals a given number. The backtracking approach generates all permutations in the worst case but in general, performs better than the recursive approach towards subset sum problem.

A subset A of n positive integers and a value **sum** is given, find whether or not there exists any subset of the given set, the sum of whose elements is equal to the given value of sum.

Example:

Problem statement: We are given 'n' distinct positive integers and a target_sum. We have to find the combinations of these numbers which exactly sum up to the target_sum value.

Let n=5 and given positive integers are my_list={1,2,3,4,5}. Let the given target_sum=6. The subsets that produce the total of 6 are $\{1,5\},\{2,4\}$ which is the output of our program. This is because 1+5=2+4=6.

Example 1:

Input: $set[] = \{4, 16, 5, 23, 12\}, sum = 9$

Output = true

Subset $\{4, 5\}$ has the sum equal to 9.

Example 2:

Input: $set[] = \{2, 3, 5, 6, 8, 10\}, sum = 10$

Output = true

There are three possible subsets that have the sum equal to 10.

Subset1: {5, 2, 3}

Subset2: {2, 8}

Subset3: {10}

Algorithm:

- 1. Start with an empty set
- 2. Add the next element from the list to the set
- 3. If the subset is having sum M, then stop with that subset as solution.
- 4. If the subset is not feasible or if we have reached the end of the set, then backtrack through the subset until we find the most suitable value.
- 5. If the subset is feasible (sum of seubset < M) then go to step 2.
- 6. If we have visited all the elements without finding a suitable subset and if no backtracking is possible then stop without solution.

Pseudocode:

```
Algorithm SUB_SET_PROBLEM(i, sum, W, remSum)
// Description : Solve sub of subset problem using backtracking
// Input :
W: Number for which subset is to be computed
i: Item index
sum: Sum of integers selected so far
remSum : Size of remaining problem i.e. (W - sum)
// Output : Solution tuple X
if FEASIBLE_SUB_SET(i) == 1 then
if (sum == W) then
  print X[1...i]
 end
else
 X[i+1] \leftarrow 1
 SUB SET PROBLEM(i + 1, sum + w[i] + 1, W, remSum - w[i] + 1)
 X[i+1] \leftarrow 0 // Exclude the ith item
 SUB\_SET\_PROBLEM(i + 1, sum, W, remSum - w[i] + 1)
end
function FEASIBLE SUB SET(i)
 if (sum + remSum \ge W) AND (sum == W) or (sum + w[i] + 1 \le W) then
  return 0
 end
return 1
```

Time Complexity: O(sum*n), where sum is the 'target sum' and 'n' is the size of array.



Lab Assignment to Complete:

- 1. Given an array of integers and a sum, the task is to have all subsets of given array, total nodes generated with sum equal to the given sum. Input: set[] = {4, 16, 5, 23, 12}, sum = 9.
- 2. Given an array of integers and a sum, the task is to have all subsets of given array, total nodes generated with sum equal to the given sum. Input: set[] = { 2, 3, 5, 6, 8, 10}, sum = 10.

Code:

```
#include <stdio.h>
#define MAX 10
int read sum = 10;
int set[MAX] = \{4, 16, 5, 23, 12\};
int len = 6;
void SOS(int level, int sum, int results[MAX]) {
    int i;
    if (sum==reqd_sum) {
        for (i=0; i<len; i++) {
            if (results[i]==1) {
                printf("%d ", set[i]);
        }
        printf("\n");
    } else if (sum>reqd_sum | level>len) {
        return;
        results[level] = 1;
        SOS(level+1, sum+set[level], results);
        results[level] = 0;
        SOS(level+1, sum, results);
int main()
    int resultset[MAX] = {0};
    printf("Numbers forming the required sum are: \n");
    SOS(0, 0, resultset);
    return 0;
```

Output:

```
Numbers forming the required sum are:
4 5

Numbers forming the required sum are:
2 3 5
2 8
2.
```