Advanced R Programming - Lecture 6 Parallel programming

Krzysztof Bartoszek (slides based on Leif Jonsson's and Måns Magnusson's)

Linköping University

krzysztof.bartoszek@liu.se

29 IX, 1 X 2025 (U2)



- Parallelism
- Theoretical limits
- Parallelism in R
- Balance and subsetting





Classical introduction: dining philosophers problem



https://en.wikipedia.org/wiki/Dining_philosophers_problem

Mostly we will not (a.s.) be concerned with such problems.



Naïve distributed leader election

- 1. Each processor has a unique id, the processors form a ring, ids can be ordered and are all comparable.
- 2. Each processor sends its id to its left neighbour.
- 3. If the received id is greater than one's own, then the received one is passed to the left neighbour. Otherwise the receiving processor becomes silent.
- 4. If a processor receives its own id back it declares itself the leader.

Complexity is $O(n^2)$ messages passed. Distributed minimum, maximum algorithm.

See e.g. https://hagit.net.technion.ac.il/da98/notes/for more sophisticated approaches.



Parallelism

What is parallelism?

Parallelism

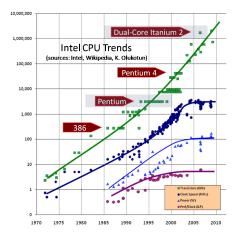
Multiple cores

Each core work with its own part

Cores can exchange information



Why parallelism?



http://www.gotw.ca/publications/concurrency-ddj.htm

Why parallelism?

Single core limits

Handling larger data

Solving problems faster

More robust (if one processor fails, others still can work)

More and more important

Is there any **but** . . . ?



Types of parallelism

Parallelism

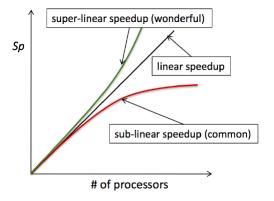
Multicore systems

Distributed systems

Graphical processing units (GPU)



Speedup



https://portal.tacc.utexas.edu/c/document_library/get_file?uuid=

e05d457a-0fbf-424b-87ce-c96fc0077099&groupId=13601



Theoretical limits

Strong scaling: Amdahl's law

Deals with fixed problem size, increasing resources

Weak scaling: Gustafsons law

Deals with increasing size problem along with increasing resources

Amdahl's law

Speedup:
$$S_p = \frac{\text{execution time on 1 processor}}{\text{execution time on } P \text{ processors}}$$

$$S_p \le \frac{1}{f_s + \frac{f_p}{P}}$$

 $f_s =$ serial fraction of computations

 f_p = parallel fraction of computations

P = number of cores

For a *fixed size problem*, single core computation time is fixed= 1!

$$\text{Efficiency}: \quad S_p P^{-1} = \frac{\text{execution time on 1 processors}}{P(\text{execution time on } P \text{ processors})}$$

7 12/ 20

Amdahl's law



https://en.wikipedia.org/wiki/Amdahl's_law



Gustafsons law

Speedup:
$$S_p = \frac{\text{execution time on 1 processor}}{\text{execution time on } P \text{ processors}}$$

$$S_p \leq \frac{\alpha + (1-\alpha)P}{\alpha + (1-\alpha)} = P - \alpha * (P-1)$$

Where:

 $\alpha =$ fraction of time dedicated to serial computations P = number of cores

Problem size scales with P, parallel execution time is fixed = 1!if we only had one core, then the P parallel computations would have to be done on that core with time $(1 - \alpha)P$



Practical problems

Costs of parallelism communication load balancing (NP-hard) scheduling (NP-hard) but 4/3 = k-approximate algorithm

Symmetry:

If
$$\forall_{P_1 \text{ in parallel}} \exists_{P_2 \text{ connected with } P_1} : v(P_1) = v(P_2) \text{ then } v(P_2) + +$$

Byzantine Generals Problem:

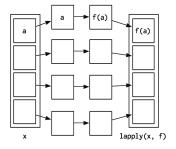
network must reach a decision but some nodes can be corrupt

fine-grained vs embarrassingly parallel (EP)



Parallelism in R (embarassingly parallel)

Based on lapply() iterations inside a loop are indepenendent of each other



(H. Wickham, Advanced R, p. 201)

What about: fctl<-1;for(i in 1:n){fctl<-i*fctl}?



Two approaches:

- 1. mclapply()
- 2. parLapply()

mclapply()

Pros

Simple to use Low overhead (startup)

Cons

Does not work on Windows
Only multi core



parLapply(type="psock")

Pros

Works everywhere Good for testing/developing

Cons

Slow on multiple nodes



Pros

Good for multiple computers Good for production

Cons

Can be used interactively Needs Rmpi package



parSapply: new R on each core

Variables need to be copied explicitely New random seed on each core

```
> library(parallel)
> x<-1
> cl <- makeCluster(getOption("cl.cores", 2))
> parSapply(cl,1:2,function(i){print(x)})
Error in checkForRemoteErrors(val):
    2 nodes produced errors; first error: object 'x' not found
> set.seed(1)
> parSapply(cl,1:2,function(i){print(rexp(1))})
[1] 0.05890067 0.90425878
> parSapply(cl,1:2,function(i){set.seed(1);print(rexp(1))})
[1] 0.7551818 0.7551818
```

All code needs to be sourced on each core



Load balance

lapply(list,...) does in order of list: can result in idle nodes

```
load balancing is NP-hard often we do not know running time for list[[i]] submit to cores in order? split into consecutive equal chunks? Example: all 2^p regression models on p predictors generate all combinations: \emptyset, \{1\}, \ldots, \{p\}, \{1,2\}, \ldots
```

randomize order in list

R's par family has load balancing capabilities: par[L/S]applyLB



Subset methods (non-EP)

glm(): large number of observations

Chunk averaging (estimation)
break data into chunks of rows

to each chunk (in parallel) apply glm()

average the results to obtain single estimate

Observations i.i.d. and model is *decent* Asymptotic equivalance for large samples

Chunks are not identically distributed: first cases then controls randomly permute observations (will not harm an initial random arrangement)



Subset methods (non-EP)

glm(): large number of predictors

Subsetting variables (prediction, *decent* model)

create random subsets of predictors to each subset apply glm() do prediction for each subset combine predictions e.g. average, majority rule

prediction difficult in high dimensions—curse of dimensionality



Example

https://github.com/STIMALiU/AdvRCourse/blob/master/Code/parallel_example.R

Parallel code example

https://github.com/STIMALiU/AdvRCourse/blob/master/Code/parallel_scalarproduct.R

Parallel scalar product example



The End... for today. Questions? See you next time!

