

Computational statistics, lecture 1

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Course schedule

- Lecture 1: Unidimensional optimization, computer arithmetic
- Math-lecture 1: Basic matrix algebra, analytical optimization, determinants
- Lecture 2: Multidimensional optimization
- Math-lecture 2: Density, cumulative distribution function, integration
- Lecture 3: Random number generation
- Lecture 4: Monte Carlo methods
- Lecture 5: Model selection and hypothesis testing
- Lecture 6: EM algorithm, stochastic optimization

Teaching group: Krzysztof Bartoszek, L5-L6; Frank Miller, examiner, L1-L4, LM1-LM2; Bayu Brahmantio, Héctor Rodriguez Déniz, teaching assistants

Course homepage: https://www.ida.liu.se/~732A90/index.en.shtml; includes schedule, reading material, lecture notes, assignments

Computer labs: For each lecture; exercises to hand-in in **groups of 2**



Evaluation of last course (HT2022)

- 23 students of around 50 submitted the evaluation; average grade 1.65
- Changes to HT2023:
 - Changes in content by deepening central topics (e.g. optimization) and instead removing some topics (e.g. reduced computer arithmetics); lab questions adjusted accordingly
 - Givens and Hoeting textbook slightly more used as basis for course layout (however, Gentle textbook covers topics as well)
 - New lecture slides for large parts of the course
 - Mathematical lectures moved inside the course (after L1 and after L2)



Computational statistics

• When large or huge datasets should be analyzed and/or complex models are used, **statistics depends on effective computational methods**

• We will **learn in this course several algorithms** for optimization, randomization, Monte Carlo integration and **methods to use them**



Today's content

- Optimization
 - Why?
 - Analytic univariate optimisation
 - Bi-section, Newton, and secant methods (univariate)
 - On convergence speed

Computer arithmetics



Optimization in statistics

- Maximum Likelihood
- Minimizing risk in (Bayesian) decision theory
- Minimizing sum of squares (Least Squares Estimate)
- Maximizing information in experimental design
- Machine learning
- Common problem in these examples:
 - x p-dimensional vector, $g: \mathbb{R}^p \to \mathbb{R}$ function
 - We search x^* with $g(x^*) = \max g(x)$
- Minimization problem turns into maximization by considering -g



Least squares estimation (LSE)

- We search a Least Squares estimate $\hat{\beta}$ for β minimising the distance $g(\hat{\beta}) = \|\hat{y} y\|^2$ from $\hat{y} = X\hat{\beta}$ to $y = X \beta + \varepsilon$
- $g(\widehat{\boldsymbol{\beta}}) = \|X\widehat{\boldsymbol{\beta}} y\|^2 = (X\widehat{\boldsymbol{\beta}} y)^T (X\widehat{\boldsymbol{\beta}} y) = \widehat{\boldsymbol{\beta}}^T X^T X \widehat{\boldsymbol{\beta}} 2\widehat{\boldsymbol{\beta}}^T X^T y + y^T y$
- Setting the derivative to $O(\frac{\partial f}{\partial \widehat{\beta}} = 2\mathbf{X}^T\mathbf{X}\widehat{\boldsymbol{\beta}} 2\mathbf{X}^T\mathbf{y} = 0)$, we get $\widehat{\boldsymbol{\beta}} = (\mathbf{X}^T\mathbf{X})^{-1}\mathbf{X}^T\mathbf{y}$
- Optimization problem:
 - $\widehat{\beta}$ p-dimensional vector, $g: \mathbb{R}^p \to \mathbb{R}$ function
 - We search $\widehat{\beta}$ with $g(\widehat{\beta}) = \min g(b)$
- Here, we do not need to iteratively compute this minimum since we have an algebraic solution $\hat{\beta} = (\mathbf{X}^T\mathbf{X})^{-1}\mathbf{X}^T\mathbf{y}$



Variations of least squares estimation

- Algebraic solution exists for the LSE, but not if we vary the problem
- Lasso estimate: $g(\widehat{\beta}) = ||X\widehat{\beta} y||^2 + \lambda ||\widehat{\beta}||_1$
- L₁-estimation: $g(\widehat{\beta}) = ||X\widehat{\beta} y||_1$
- Many further variations of estimates have been considered
- In all cases, we search $\widehat{\beta}$ with $g(\widehat{\beta}) = \min g(b)$
- Recall: Norms for $\mathbf{x} = (x_1, ..., x_p)^T$: $\|\mathbf{x}\| = \|\mathbf{x}\|_2 = \sqrt{x_1^2 + \dots + x_p^2}$ (Euclid), $\|\mathbf{x}\|_1 = |x_1| + \dots + |x_p|$, $\|\mathbf{x}\|_{\infty} = \max\{|x_1|, ..., |x_p|\}$ (max-norm)



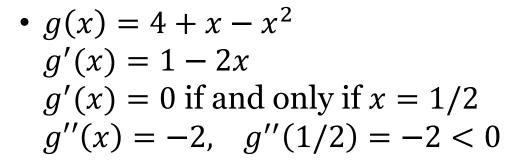
Univariate optimization

- x real number, $g: \mathbb{R} \to \mathbb{R}$ continuously differentiable function
- We search x^* with $g(x^*) = \max g(x)$
- Compute g'(x) and search x^* with $g'(x^*) = 0$
- One has then to check if the result is maximum, minimum, possibly local optimum...

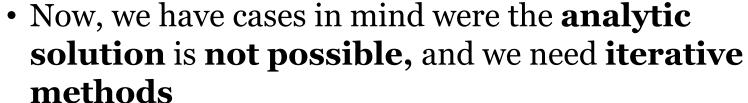


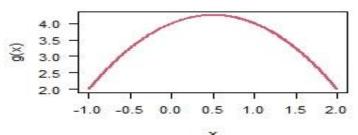
Univariate optimization - analytic solution

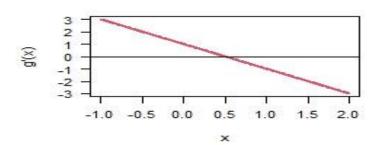
- Compute g'(x) and search x^* with $g'(x^*) = 0$
- Example, where analytic optimization possible:

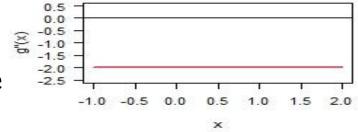








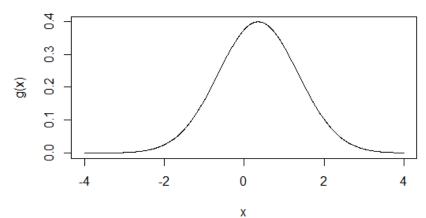


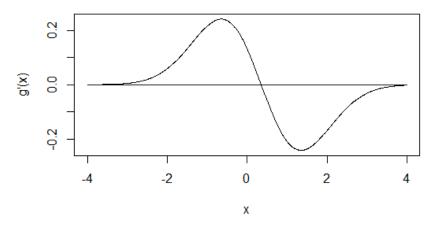




Univariate optimization: bisection

- Search x^* with $g'(x^*) = 0$:
 - 1) Start with interval $[a_0, b_0]$ such that $g'(a_0) \cdot g'(b_0) < 0, t = 0$
 - 2) Set $x^{(t)} = (a_t + b_t)/2$
 - 3) Define next interval $[a_t, b_t]$ by $[a_t, x^{(t)}]$ if $g'(a_t) \cdot g'(x^{(t)}) \leq 0$, $[x^{(t)}, b_t]$ if $g'(x^{(t)}) \cdot g'(b_t) < 0$
 - 4) Set t to t+1 and go to 2)
- See <u>video on course homepage</u>
- **Iteratively** improve approximations for x^* : $x^{(0)} \rightarrow x^{(1)} \rightarrow x^{(2)} \rightarrow ...$







Optimization: convergence criterion

- Compare $x^{(t)}$ and $x^{(t+1)}$ and stop if they are "close enough"
- Absolute convergence criterion:

$$\left| x^{(t+1)} - x^{(t)} \right| < \epsilon$$

• Relative convergence criterion:

$$\frac{\left|x^{(t+1)} - x^{(t)}\right|}{\left|x^{(t)}\right|} < \epsilon$$



Univariate Newton(-Raphson)

- x real number, $g: \mathbb{R} \to \mathbb{R}$ twice differentiable function
- Search x^* with $g(x^*) = \max g(x)$ by searching x^* with $g'(x^*) = 0$
- Taylor expansion around x^* motivates:

$$0 = g'(x^*) \approx g'(x^{(t)}) + (x^* - x^{(t)})g''(x^{(t)})$$
$$-(x^* - x^{(t)})g''(x^{(t)}) \approx g'(x^{(t)})$$
$$x^* \approx x^{(t)} - g'(x^{(t)})/g''(x^{(t)})$$

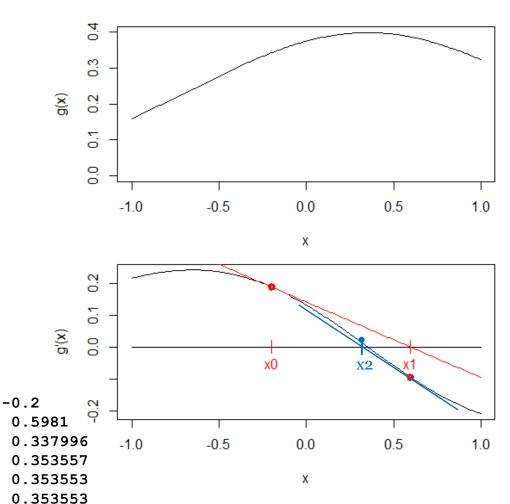
• Therefore, the Newton-iteration works as:

$$x^{(t+1)} = x^{(t)} - g'(x^{(t)})/g''(x^{(t)})$$



Univariate Newton(-Raphson)

- $x^{(t+1)} = x^{(t)} g'(x^{(t)})/g''(x^{(t)})$
- Start with a $x^{(0)}$
- Tangent in $(x^{(0)}, g'(x^{(0)}))$ determines $x^{(1)}$
- Tangent in $(x^{(1)}, g'(x^{(1)}))$ determines $x^{(2)}$
- ...
- until convergence criterion met
- +Newton method is fast
- Requires existence and computation of g"



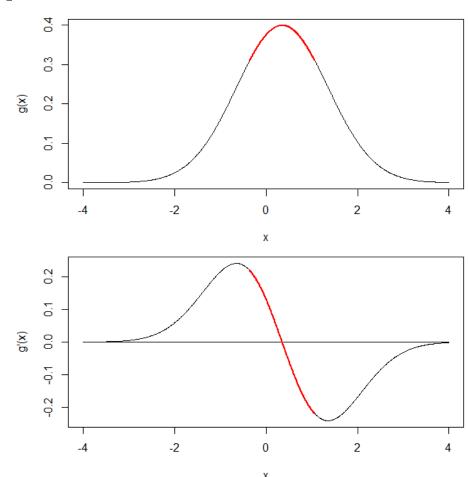
STOP



Univariate Newton(-Raphson)

•
$$x^{(t+1)} = x^{(t)} - g'(x^{(t)})/g''(x^{(t)})$$

• What about the starting value $x^{(0)}$?





Univariate secant method

- *x* real number, $g: \mathbb{R} \to \mathbb{R}$ once differentiable function
- Search x^* with $g(x^*) = \max g(x)$ by searching x^* with $g'(x^*) = 0$
- Recall: The Newton-iteration works as:

$$x^{(t+1)} = x^{(t)} - g'(x^{(t)})/g''(x^{(t)})$$

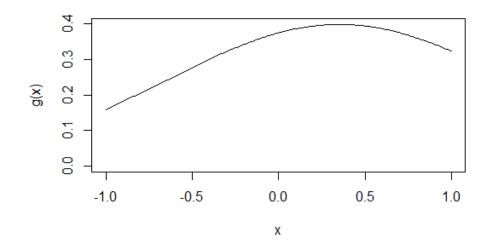
- Need to compute g'' which might be difficult. Instead:
- Approximate $g''(x^{(t)})$ by $[g'(x^{(t)}) g'(x^{(t-1)})]/(x^{(t)} x^{(t-1)})$

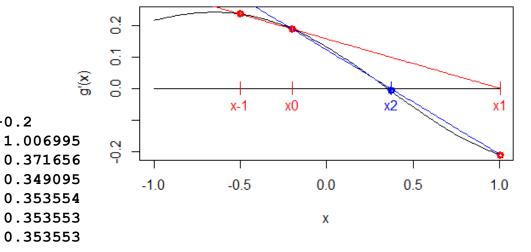


Univariate secant method

•
$$x^{(t+1)} = x^{(t)} - g'(x^{(t)}) \frac{x^{(t)} - x^{(t-1)}}{g'(x^{(t)}) - g'(x^{(t-1)})}$$

- Start with $x^{(0)}$ and $x^{(-1)}$
- Secant through $x^{(0)}$ and $x^{(-1)}$ determines $x^{(1)}$
- Secant through $x^{(1)}$ and $x^{(0)}$ determines $x^{(2)}$
- until stopping crit. fulfilled
- Quite fast
- No 2nd derivative necessary





-0.2

STOP



Convergence speed of optimization algorithms

Convergence order

- Convergence speed can be quantified by q and c as follows:
 - Let $\varepsilon^{(t)} = x^{(t)} x^*$.

rate

• Find q and c such that $\lim_{t\to\infty} \varepsilon^{(t+1)}/(\varepsilon^{(t)})^q = c$

Intuitively,
$$\varepsilon^{(t+1)} \approx c \cdot (\varepsilon^{(t)})^q$$

- $\varepsilon = 1, 0.5, 0.25, 0.125, 0.063, 0.031, ... \Rightarrow q=1, c=0.5,$
- $\varepsilon = 1, 0.1, 0.01, 0.001, 0.0001, ...$

$$\Rightarrow$$
 q=1, c=0.1

• If q=1, we say that convergence is "linear"

- $\varepsilon = 1, 0.5, 0.125, 0.008, 0.00003, ...$ $\Rightarrow q=2, c=0.5.$
- If q=2, we say that convergence is "quadratic"



Determine empirically convergence rate (and order) of optimization algorithms

• You have a given optimization algorithm and you have determined or know the maximiser x^* . To check convergence speed in an optimization-run, you can calculate

$$D^{(t)} = \frac{\left| x^{(t)} - x^* \right|}{\left| x^{(t-1)} - x^* \right|}$$

(see Givens and Hoeting, 2013, page 101/102, for an example)

• If $D^{(t)} \to 1$, there is not even linear convergence (bad, order q < 1), If $D^{(t)} \to c \in (0,1)$, linear convergence (order q = 1) with rate c, If $D^{(t)} \to 0$, better than linear convergence (order q > 1).



Comparison of univariate optimization methods

Bisection	Secant	Newton
g' required	g' required	$g^{\prime\prime}$ required
finds always an optimum between a_0 and b_0 (but could be local)	converges only when the two starting values "close" to optimum	converges only when starting value "close" to optimum
slow $q=1$	$q = \frac{1 + \sqrt{5}}{2} = 1.62$	fast $q=2$

- There are also algorithms not needing g'
- R-function optimize uses such an algorithm (q=1.324)



Computer arithmetics

- Numbers are represented as binary numbers $(17 = 1 * 2^4 + 1 * 2^1 = "1001")$
- Rational numbers are also represented based on the binary system:

$$\pm 0. d_1 d_2 \dots d_p * 2^e,$$

 $e = \pm e_1 e_2 \dots e_a$

- E.g. p = 52, q = 10, two signs \Rightarrow one number needs 64 bits in the computer
- Limits in representation depending on p and q

```
> 3/5-2/5-1/5
                                                   [1] -5.551115e-17
> 2^1023
                        > 2^{-1074}
                                                   > if (3/5-2/5==1/5)
[1] 8.988466e+307
                       [1] 4.940656e-324
                                                   > print("yes") else
> 2^1024
                        > 2^-1075
                                                       print("no")
[1] Inf
                        [1] 0
                                                   [1] "no"
(overflow)
                        (underflow)
```



Computer arithmetics

- Good to have limitations of computer arithmetics in mind!
- Example: Binomial coefficient (avoid overflow)

$$\binom{n}{k} = \frac{n!}{k!(n-k)!} = \frac{(k+1)(k+2)\cdots(n-1)n}{(n-k)!}$$

$$\binom{200}{2} = \frac{200!}{2!198!} = \frac{3*4*\cdots*199*200}{1*2*\cdots*197*198} = \frac{199*200}{1*2} = 19900$$

```
> n <- 200
> k <- 2
> prod(1:n) / (prod(1:k)*prod(1:(n-k)))
[1] NaN
> prod(((k+1):n) / (1:(n-k)))
[1] 19900
```



Course material, lab, seminar, exam

- Homepage: https://www.ida.liu.se/~732A90/index.en.shtml
 - Lecture notes, lab- and seminar info, exam info
- Submission of 6 labs via LISAM all need to be passed groups of 2
 - First lab: Oct 31 to Nov 7
- Mandatory attendance at 3 seminars and 1 presentation or opposition
- Computer exam: Jan 9, 2024. Own document of 100 pages can be used.
- 10 points to pass (E); 12 or more: D; \geq 14: C; \geq 16: B; \geq 18: A

Literature:

- Gentle JE (2009). Computational Statistics, Springer
- Givens GH, Hoeting JA (2013). Computational Statistics, Wiley

