

L06: GCDs and Congruences

- **Greatest Common Divisor (GCD)**
 - Multiplicative Inverses
 - Solving Linear Congruences
 - The Chinese Remainder Theorem
-
- Reading: Rosen 4.3, 4.4, 4.5

Review of Primary School Knowledge

- **Definition**

A positive integer p greater than 1 is called **prime** if the only positive factors of p are 1 and p . A positive integer that is greater than 1 and is not prime is called *composite*.

- **Theorem** (The Fundamental Theorem of Arithmetic)
Every positive integer greater than 1 can be written uniquely as a prime or as the product of two or more primes where the prime factors are written in order of nondecreasing size.
(Will prove later by induction)

Greatest Common Divisor

- **Definition**

Let a and b be integers, not both zero. The largest integer d such that $d \mid a$ and also $d \mid b$ is called the **greatest common divisor** of a and b , denoted by $\gcd(a, b)$.

- One can find the gcd by prime factorizations

- **Example**

$$120 = 2^3 \cdot 3 \cdot 5 \quad 500 = 2^2 \cdot 5^3$$

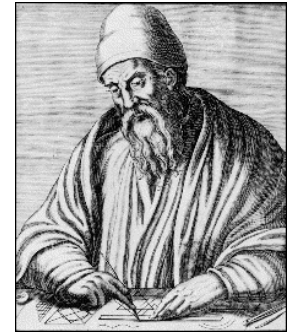
$$\gcd(120, 500) = 2^{\min(3,2)} \cdot 3^{\min(1,0)} \cdot 5^{\min(1,3)} = 2^2 \cdot 3^0 \cdot 5^1 = 20$$

- **Definition**

The integers a and b are **relatively prime** if $\gcd(a, b) = 1$.

- Example: 17 and 22

Euclidean Algorithm



- However, factoring large numbers is hard!
 - No efficient algorithms exist

- **Lemma**

Let $a = bq + r$, where a , b , q , and r are integers. Then $\gcd(a, b) = \gcd(b, r)$.

- **Proof**

- Suppose that d divides both a and b . Then d also divides $a - bq = r$. Hence, any common divisor of a and b must also be any common divisor of b and r .
- Suppose that d divides both b and r . Then d also divides $bq + r = a$. Hence, any common divisor of b and r must also be a common divisor of b and a .
- Therefore, $\gcd(a, b) = \gcd(b, r)$.

Euclidean Algorithm

- Idea: To obtain maximum efficiency, choose the smallest r , i.e., $r = a \bmod b$ (suppose $a > b$), and iterate.

```
gcd( $a, b$ ) :  
   $x \leftarrow a$   
   $y \leftarrow b$   
  while  $y \neq 0$   
     $r \leftarrow x \bmod y$   
     $x \leftarrow y$   
     $y \leftarrow r$   
  return  $x$ 
```

Example:

```
gcd(287, 91)  
= gcd(91, 14)  
= gcd(14, 7)  
= gcd(7, 0)  
= 7
```

- Correctness of algorithm follows from previous lemma
- Termination is obvious
- Running time will be analyzed later

gcds as Linear Combinations

- **Theorem** (Bézout's Theorem)

If a and b are positive integers, then there exist integers s and t such that $\gcd(a, b) = sa + tb$.

- Example

$$\gcd(6, 14) = (-2) \cdot 6 + 1 \cdot 14$$

- Instead of proving this theorem directly, we give an algorithm to find such s and t .

The Extended Euclidean Algorithm

■ Example

Express $\gcd(252, 198)$ as a linear combination of 252 and 198.

■ Solution

■ First find $\gcd(252, 198)$

1) $252 = 1 \cdot 198 + 54$

2) $198 = 3 \cdot 54 + 36$

3) $54 = 1 \cdot 36 + 18$

4) $36 = 2 \cdot 18$

5) $\gcd(252, 198) = 18$

■ Rewriting:

■ $54 = 252 - 1 \cdot 198$

■ $36 = 198 - 3 \cdot 54$

■ $18 = 54 - 1 \cdot 36$

■ Substituting:

$$18 = 54 - 1 \cdot (198 - 3 \cdot 54)$$

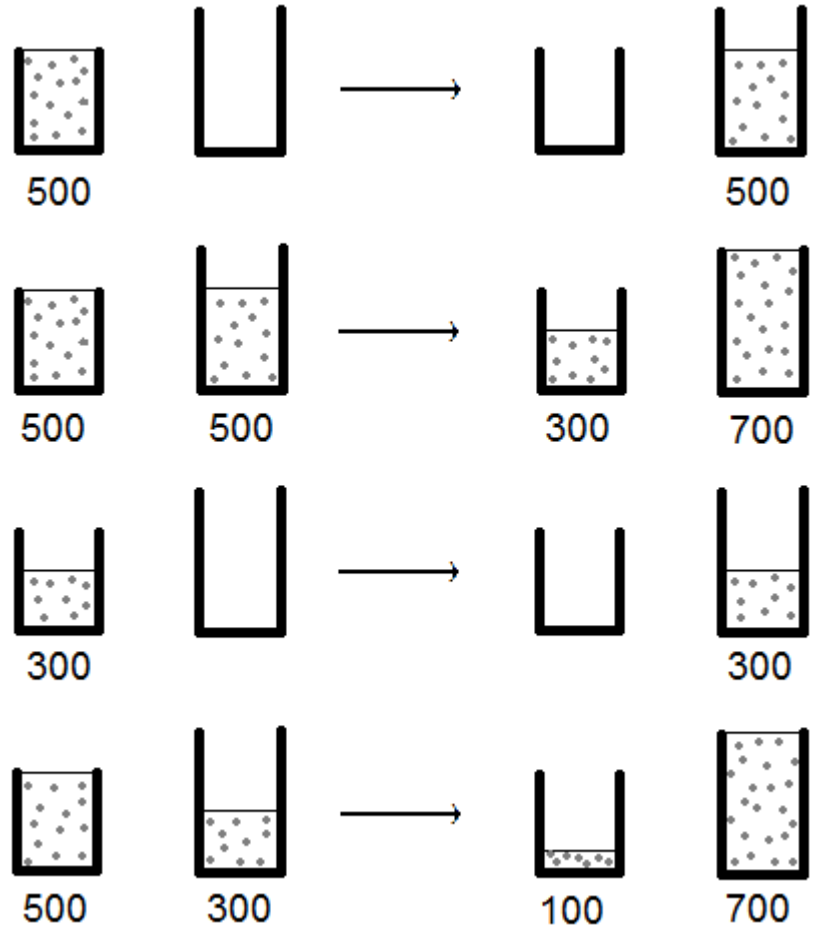
$$= 4 \cdot 54 - 1 \cdot 198$$

$$= 4 \cdot (252 - 1 \cdot 198) - 1 \cdot 198$$

$$= 4 \cdot 252 - 5 \cdot 198$$

Puzzle: Water Measuring

- Given
 - Two bottles: one has volume of 500 ml and the other one 700 ml.
 - Infinite water supply
- Goal: Get exactly 100 ml of water
- This follows exactly from
$$100 = \gcd(500, 700)$$
$$= 3 \times 500 - 2 \times 700$$
- Corollary: Any multiple of the gcd can be obtained.



Outline

- Greatest Common Divisor (GCD)
- **Multiplicative Inverses**
- Solving Linear Congruences
- The Chinese Remainder Theorem

Multiplicative Inverses

■ Definition

The (multiplicative) inverse of a modulo m is some b such that $ab \equiv 1 \pmod{m}$.

- By default “inverse” means “multiplicative inverse”.

■ Examples

\mathbf{Z}_5 :

a	1	2	3	4
a^{-1}	1	3	2	4

\mathbf{Z}_7 :

a	1	2	3	4	5	6
a^{-1}	1	4	5	2	3	6

\mathbf{Z}_6 :

a	1	2	3	4	5
a^{-1}	1	X	X	X	5

\mathbf{Z}_8 :

a	1	2	3	4	5	6	7
a^{-1}	1	X	3	X	5	X	7

Multiplicative Inverses

- **Theorem**

For any $a \in \mathbf{Z}_m, m > 1$, if $\gcd(a, m) = 1$ then a has a unique inverse in \mathbf{Z}_m .

- **Corollary**

For any prime p , every nonzero $a \in \mathbf{Z}_p$ has a multiplicative inverse.

- **Proof of Theorem**

Since $\gcd(a, m) = 1$, by Bézout's Theorem, there are integers s and t such that $sa + tm = 1$.

- Hence, $sa = (-t)m + 1$
- Since $tm \equiv 0 \pmod{m}$, it follows that $sa \equiv 1 \pmod{m}$.
- Since s may not be in \mathbf{Z}_m , we write $(s \bmod m) \cdot_m a = 1$
- Consequently, $s \bmod m$ is the inverse of a in \mathbf{Z}_m .

Multiplicative Inverses are Unique

- **Proof of uniqueness**

- Suppose b, c are both inverses of a , i.e.,

$$ab \equiv 1 \pmod{m} \quad (1)$$

$$ac \equiv 1 \pmod{m} \quad (2)$$

- Multiply both sides of (1) by c :

$$abc \equiv c \pmod{m}$$

- Multiply both sides of (2) by b :

$$abc \equiv b \pmod{m}$$

- So $b \equiv c \pmod{m}$, i.e., a has a unique inverse in \mathbf{Z}_m .

- The inverse of a is written as a^{-1} .

- Note: It's also true that if $\gcd(a, m) \neq 1$, a^{-1} doesn't exist. (Left as an exercise.)

Finding Inverses

- Given a, m such that $\gcd(a, m) = 1$, how to find the inverse of a in \mathbf{Z}_m ?
- Look at the proof of the previous theorem
 - Use the extended Euclidean algorithm to find s and t such that $sa + tm = 1$
 - $s \bmod m$ is the multiplicative inverse of a in \mathbf{Z}_m .
- **Example**
Find an inverse of 3 modulo 7
- **Solution**
Using the extended Euclidean algorithm: $7 = 2 \cdot 3 + 1$.
we get $-2 \cdot 3 + 1 \cdot 7 = 1$, so $s = -2$.
 $-2 \bmod 7 = 5$ is the inverse of 3 in \mathbf{Z}_7

Finding Inverses

- **Example**

Find the inverse of 101 modulo 4620

Working Backwards:

$$4620 = 45 \cdot 101 + 75$$

$$101 = 1 \cdot 75 + 26$$

$$75 = 2 \cdot 26 + 23$$

$$26 = 1 \cdot 23 + 3$$

$$23 = 7 \cdot 3 + 2$$

$$3 = 1 \cdot 2 + 1$$

$$2 = 2 \cdot 1$$

Since the last nonzero remainder is 1,
 $\gcd(101, 4620) = 1$

$$1 = 3 - 1 \cdot 2$$

$$1 = 3 - 1 \cdot (23 - 7 \cdot 3) = -1 \cdot 23 + 8 \cdot 3$$

$$1 = -1 \cdot 23 + 8 \cdot (26 - 1 \cdot 23) = 8 \cdot 26 - 9 \cdot 23$$

$$1 = 8 \cdot 26 - 9 \cdot (75 - 2 \cdot 26) = 26 \cdot 26 - 9 \cdot 75$$

$$1 = 26 \cdot (101 - 1 \cdot 75) - 9 \cdot 75$$

$$= 26 \cdot 101 - 35 \cdot 75$$

$$1 = 26 \cdot 101 - 35 \cdot (4620 - 45 \cdot 101)$$

$$= -35 \cdot 4620 + 1601 \cdot 101$$

1601 is an inverse of
101 modulo 4620

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Solving Congruences

- Linear congruence

$$ax \equiv b \pmod{m}$$

- Given a, b, m , such that $\gcd(a, m) = 1$. How to find x ?

- Solution:

- Find a^{-1}
- Multiply a^{-1} on both sides

- Example

- Solve $3x \equiv 4 \pmod{7}$
- Find $3^{-1} = 5$
- Multiply 5 on both sides:

$$3^{-1} \cdot 3x \equiv 3^{-1} \cdot 4 \pmod{7}$$

$$x \equiv 5 \quad 4 \equiv 6 \pmod{7}$$

Solving Congruences

- **Corollary**

If $\gcd(a, m) = 1$, the linear congruence

$$ax \equiv b \pmod{m}$$

has a unique solution in \mathbf{Z}_m

- **Proof**

- Existence has already been proved by construction.
- Uniqueness: Suppose it has two solutions x_1, x_2 :

$$ax_1 \equiv b \pmod{m}$$

$$ax_2 \equiv b \pmod{m}$$

Multiply both by a^{-1} :

$$x_1 \equiv ba^{-1} \pmod{m}$$

$$x_2 \equiv ba^{-1} \pmod{m}$$

So, $x_1 \equiv x_2 \pmod{m}$.

Solving congruences

- Note: If $\gcd(a, m) \neq 1$, the linear congruence
$$ax \equiv b \pmod{m}$$
may have no solution or multiple solutions in \mathbf{Z}_m

Examples:

- $2x \equiv 1 \pmod{6}$ has no solution in \mathbf{Z}_6
- $2x \equiv 4 \pmod{6}$ has two solutions in \mathbf{Z}_6
$$x = 2, 5$$

Revisiting the String Hash Function

- Consider the case with only 3 characters:

$$h(s) = \left(((s[0] \cdot 31 + s[1]) \cdot 31 + s[2]) \bmod 2^{32} \right) \bmod n$$

- Note $\gcd(31, 2^{32}) = 1$
- Given any $s[2]$ and b , the congruence
$$31x + s[2] \equiv b \pmod{2^{32}}$$
has a unique solution. This means, given $s[2]$, $h(s)$ depends on $x = s[0] \cdot 31 + s[1]$ and every b is possible.
- Similarly, given any $s[1]$, x can possibly take any value depending on $s[0]$.
- Other reasons:
 - Performance: $x * 31 = x \ll 5 - x$
 - Using 31 produces more balanced hashes over English text

Checksums

- **Example**

HKID numbers are of the format X123456(Y), where

- X is one or two letters
- Y is check digit, 0 to 9 or A.

- **How is it computed**

- Replace the first two letters as follows:

A = 10 B = 11 C = 12 D = 13 E = 14 F = 15 G = 16 H = 17 I = 18 J = 19
K = 20 L = 21 M = 22 N = 23 O = 24 P = 25 Q = 26 R = 27 S = 28
T = 29 U = 30 V = 31 W = 32 X = 33 Y = 34 Z = 35 empty = 36

- Denote the resulting two numbers and 6 digits as

x_1, \dots, x_8

- $c = (9x_1 + 8x_2 + 7x_3 + 6x_4 + \dots + 2x_8) \bmod 11$

- Check digit $x_9 = 11 - c$

If $x_9 = 11$, check digit = 0

If $x_9 = 10$, check digit = A

HKID Checksum: Single Error

- Note that for a valid HKID, we have

$$9x_1 + 8x_2 + 7x_3 + \cdots + 2x_8 + x_9 \equiv 0 \pmod{11}$$

- Suppose x_3 is mistyped as $x'_3 \neq x_3$

- Suppose the checksum is still correct, i.e.,

$$9x_1 + 8x_2 + 7x'_3 + \cdots + 2x_8 + x_9 \equiv 0 \pmod{11}$$

- Subtracting one congruence from the other:

$$7(x_3 - x'_3) \equiv 0 \pmod{11}$$

- Since $\gcd(7, 11) = 1$, 7 has an inverse. Multiply both sides by 7^{-1} :

$$x_3 - x'_3 \equiv 0 \pmod{11}$$

- This contradicts with the assumption $x'_3 \neq x_3$ and they are both in $\{0, \dots, 9\}$
- Note: If first or second letter is wrong, it may not be detected!

HKID Checksum: Transposition Error

- Note that for a valid HKID, we have

$$9x_1 + 8x_2 + 7x_3 + \cdots + 2x_8 + x_9 \equiv 0 \pmod{11}$$

- Suppose x_3 and x_5 are swapped, and $x_3 \neq x_5$
- Suppose the checksum is still correct, i.e.,
 $9x_1 + 8x_2 + 7x_5 + 6x_4 + 5x_3 \dots + 2x_8 + x_9 \equiv 0 \pmod{11}$

- Subtracting one congruence from the other:

$$2(x_5 - x_3) \equiv 0 \pmod{11}$$

- Since $\gcd(2, 11) = 1$, 2 has an inverse. Multiply both sides by 2^{-1} :

$$x_5 - x_3 \equiv 0 \pmod{11}$$

- This contradicts with the assumption $x_3 \neq x_5$ and they are both in $\{0, \dots, 9\}$

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Sun-Tsu's Problem

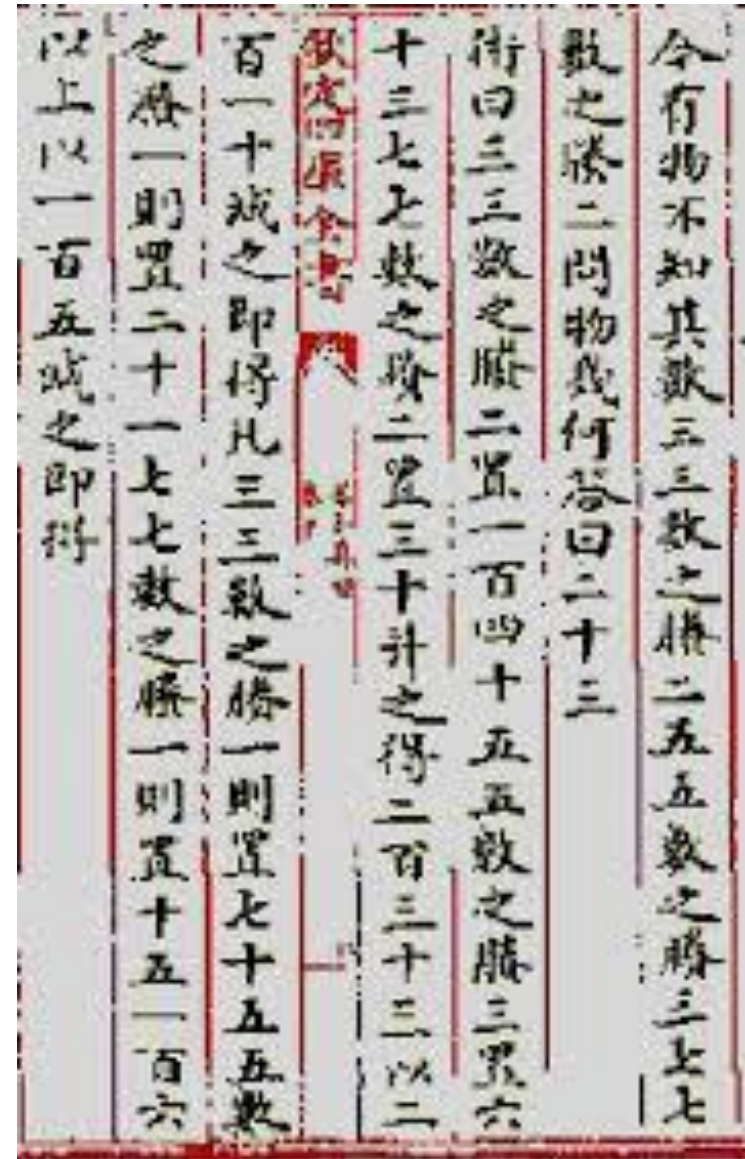
There are certain things whose number is unknown. If we count them by threes, we have two left over; by fives, we have three left over; and by sevens, two are left over. How many things are there?

System of linear congruences:

$$x \equiv 2 \pmod{3},$$

$$x \equiv 3 \pmod{5},$$

$$x \equiv 2 \pmod{7}.$$



The Chinese Remainder Theorem

- **Theorem**

Let m_1, m_2, \dots, m_n be pairwise relatively prime positive integers greater than one and a_1, a_2, \dots, a_n arbitrary integers. Then the system

$$x \equiv a_1 \pmod{m_1}$$

$$x \equiv a_2 \pmod{m_2}$$

...

$$x \equiv a_n \pmod{m_n}$$

has a unique solution modulo $m = m_1 m_2 \cdots m_n$.

- **Proof**

We'll show that a solution exists by describing a way to construct the solution. (Uniqueness proof is left as exercise.)

The Chinese Remainder Theorem

■ Proof

Let $M_k = \frac{m}{m_k}, k = 1, 2, \dots, n$

Since $\gcd(m_k, M_k) = 1$, M_k has an inverse y_k modulo m_k :

$$M_k y_k \equiv 1 \pmod{m_k}$$

We claim that this is a solution:

$$x = a_1 M_1 y_1 + a_2 M_2 y_2 + \dots + a_n M_n y_n$$

Check:

$$x \equiv a_k \pmod{m_k}?$$

$$M_j \equiv 0 \pmod{m_k} \text{ if } j \neq k \Rightarrow a_j M_j y_j \equiv 0 \pmod{m_k} \text{ if } j \neq k;$$

$$M_k y_k \equiv 1 \pmod{m_k} \Rightarrow a_k M_k y_k \equiv a_k \pmod{m_k}$$

The Chinese Remainder Theorem

- Consider the 3 congruences from Sun-Tsu's problem:
 $x \equiv 2 \pmod{3}, x \equiv 3 \pmod{5}, x \equiv 2 \pmod{7}.$
- Let $m = 3 \cdot 5 \cdot 7 = 105$, $M_1 = m/3 = 35$, $M_2 = m/5 = 21$, $M_3 = m/7 = 15$.
- We see that
 - 2 is an inverse of $M_1 \pmod{3}$
 - 1 is an inverse of $M_2 \pmod{5}$
 - 1 is an inverse of $M_3 \pmod{7}$
- Hence,

$$\begin{aligned} x &= a_1 M_1 y_1 + a_2 M_2 y_2 + a_3 M_3 y_3 \\ &= 2 \cdot 35 \cdot 2 + 3 \cdot 21 \cdot 1 + 2 \cdot 15 \cdot 1 = 233 \equiv 23 \pmod{105} \end{aligned}$$