M2 – Instruction Set Architecture

Module Outline

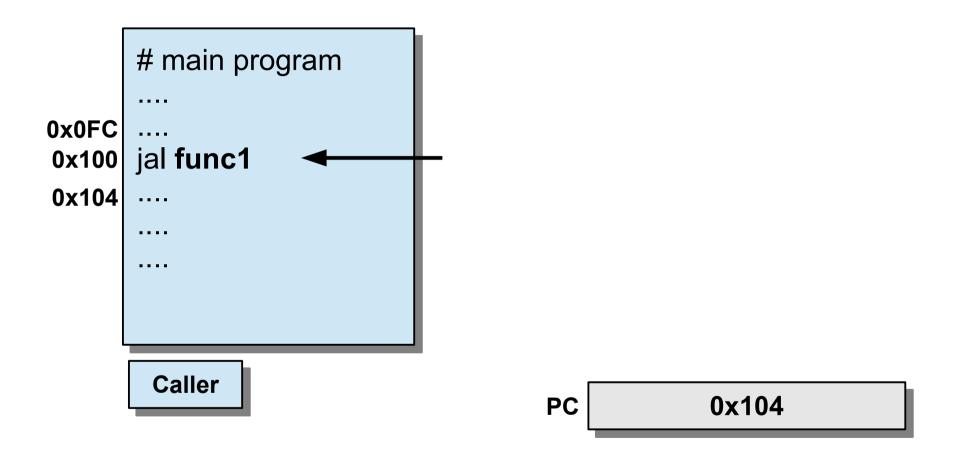
- Addressing modes. Instruction classes.
- MIPS-I ISA.
- Translating and starting a program.
- High level languages, Assembly languages and object code.
- Subroutine and subroutine call. Use of stack for handling subroutine call and return.

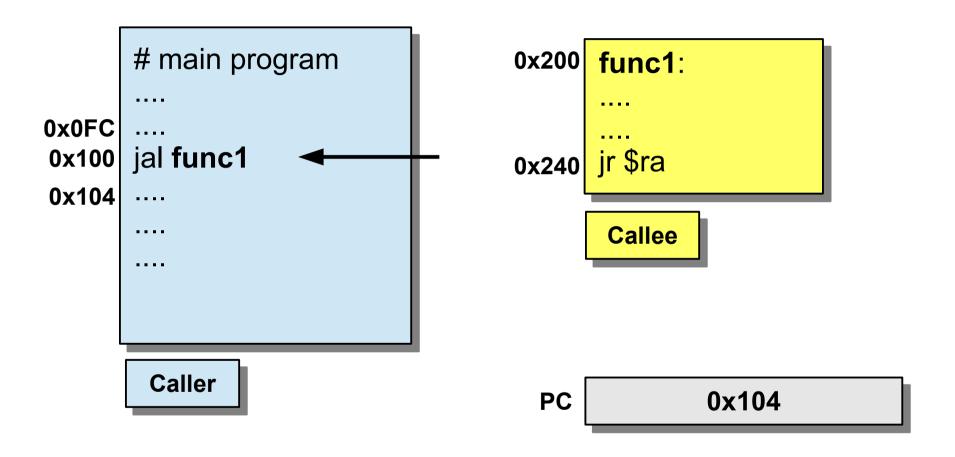
Subroutine Calls

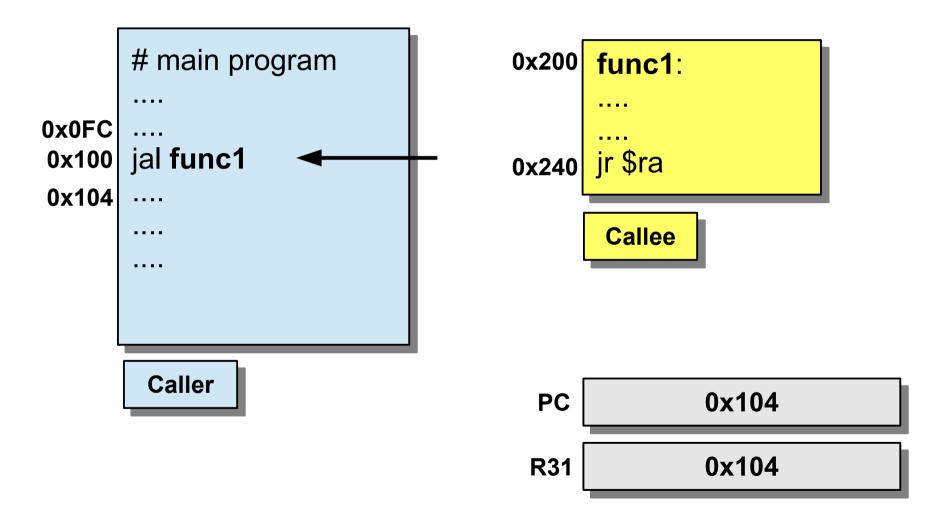
- Subroutine Call jal subname
 - Saves return address in R31 (\$ra) and jumps to subroutine entry label subname

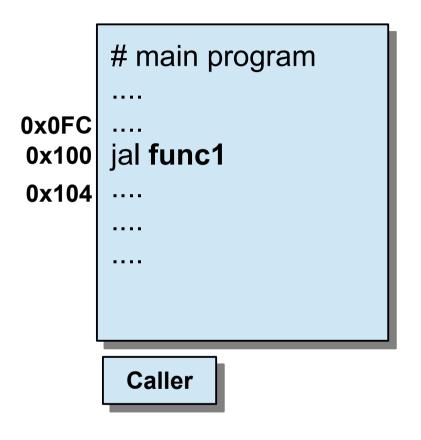
- Subroutine Call jal subname
 - Saves return address in R31 (\$ra) and jumps to subroutine entry label subname
- Subroutine Return jr \$31
 - Loads PC with return address in \$31

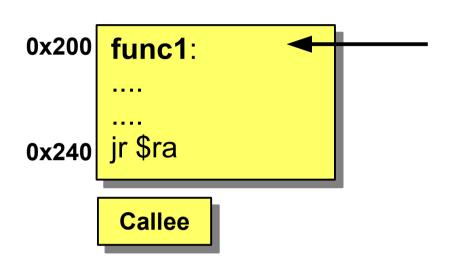
```
# main program
....
0x0FC
0x100 jal func1
....
....
....
```

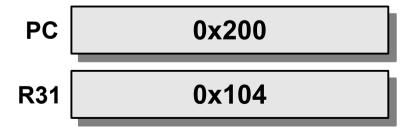


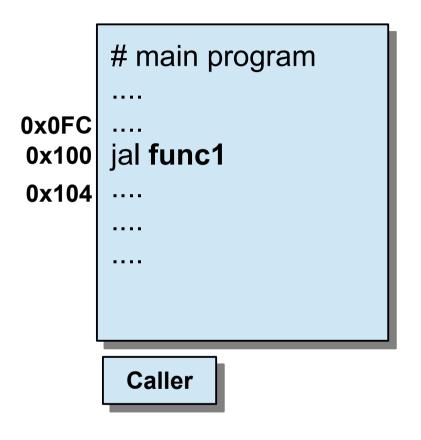


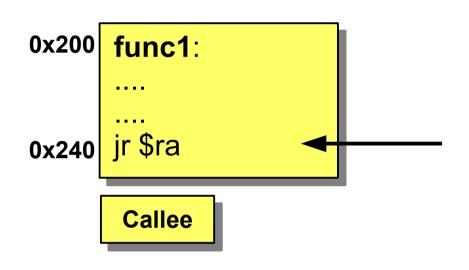


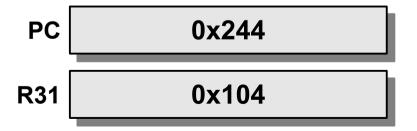


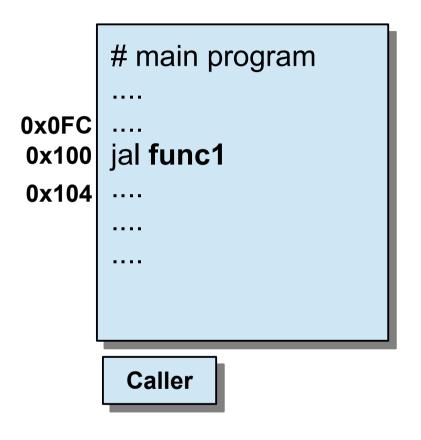


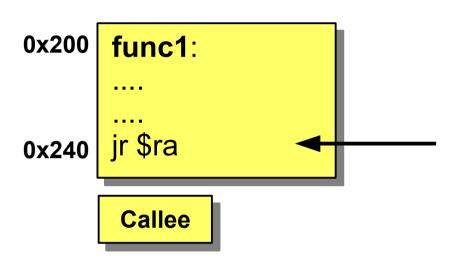


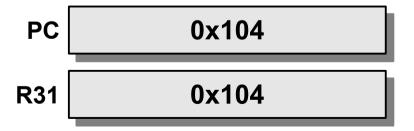


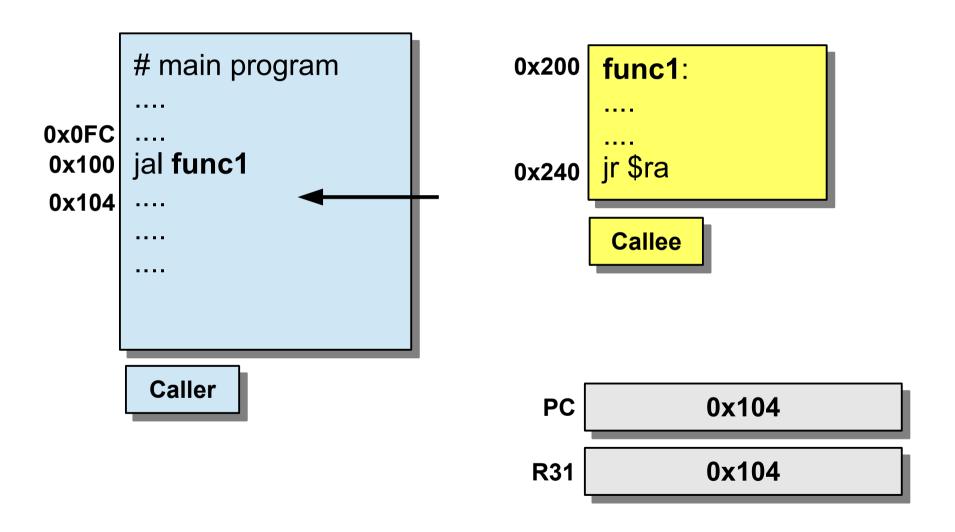












Registers Usage Convention

Name	Register number	Usage	Preserved on call?
\$zero	0	The constant value 0	n.a.
\$v0-\$v1	2–3	Values for results and expression evaluation	no
\$a0-\$a3	4–7	Arguments	no
\$t0-\$t7	8–15	Temporaries	no
\$s0 - \$s7	16–23	Saved	yes
\$t8-\$t9	24–25	More temporaries	no
\$gp	28	Global pointer	yes
\$sp	29	Stack pointer	yes
\$fp	30	Frame pointer	yes
\$ra	31	Return address	yes

```
# main program
add R4, R0, R16
add R5, R0, R17
jal func1
```

```
0x200 func1:
....
0x240 jr $ra
```

```
# main program
add $a0, $zero, $s0
add $a1, $zero, $s1
jal accArray
print $v0
```

```
accArray:
add $v0, $zero, $zero
loop:
beq $a0, $zero, done
lw $t0, 0($a1)
add $v0, $v0, $t0
addiu $a1, $a1, 4
addi $a0, $a0, -1
j loop
done:
jr $ra
```

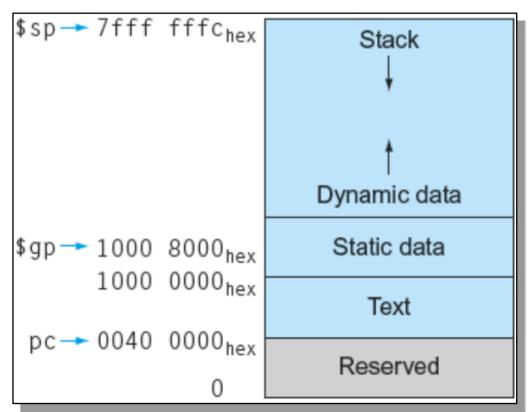
Caller saves parameters in \$a0 - \$a3

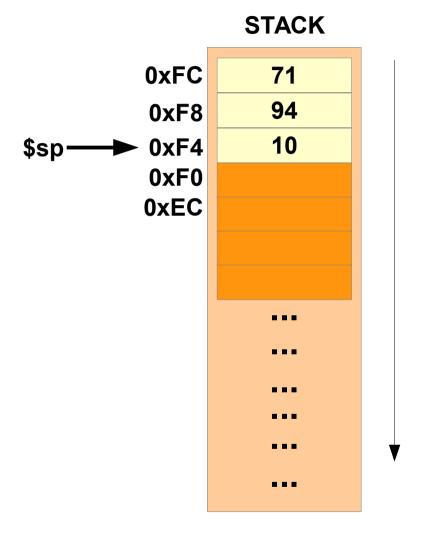
- Caller saves parameters in \$a0 \$a3
- Callee stores results in \$v0, \$v1.

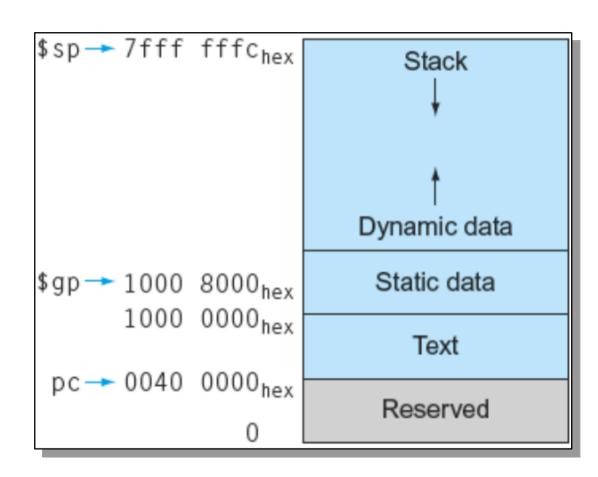
- Caller saves parameters in \$a0 \$a3
- Callee stores results in \$v0, \$v1.
- How does the caller pass more than 4 parameters to the callee?

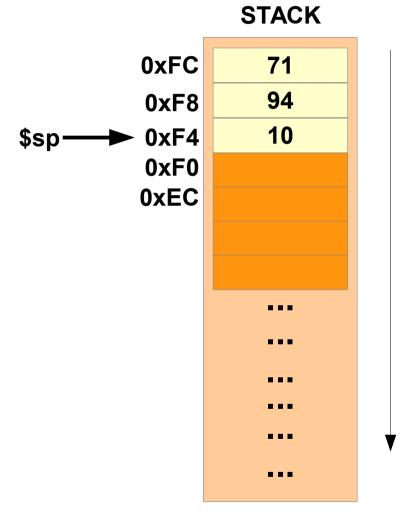
- Caller saves parameters in \$a0 \$a3
- Callee stores results in \$v0, \$v1.
- How does the caller pass more than 4 parameters to the callee?
- Program stack

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- How does the caller pass more than 4 parameters to the callee?
- Program stack

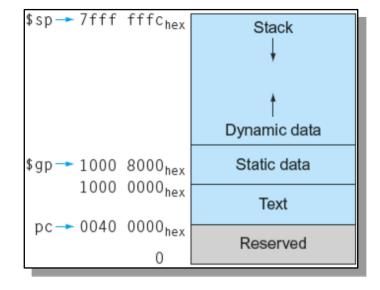




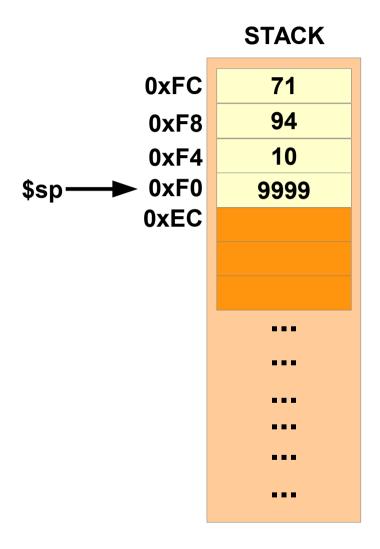




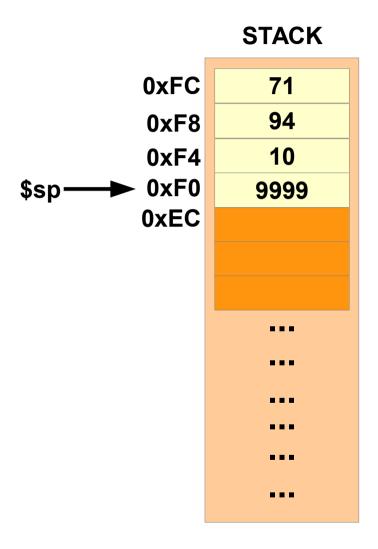




Push the value in \$t0 on the stack

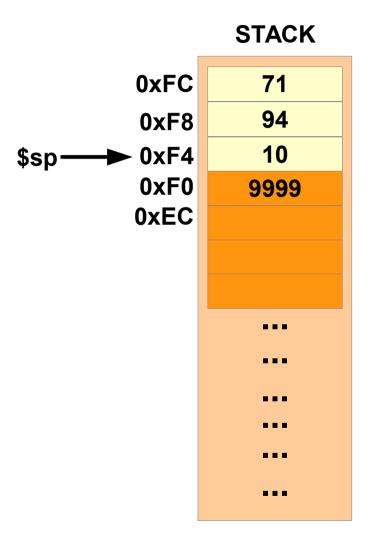


Push the value in \$t0 on the stack

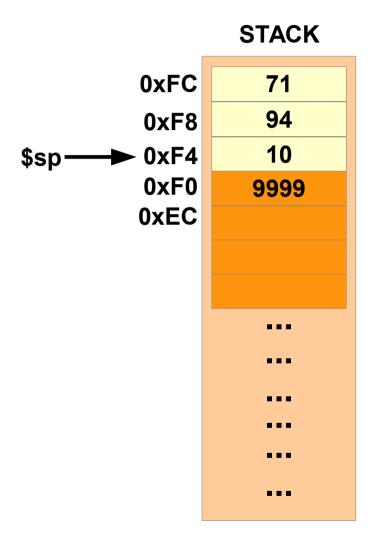


Push the value in \$t0 on the stack

```
Push
addi $sp, $sp, -4
sw $t0, 0($sp)
```



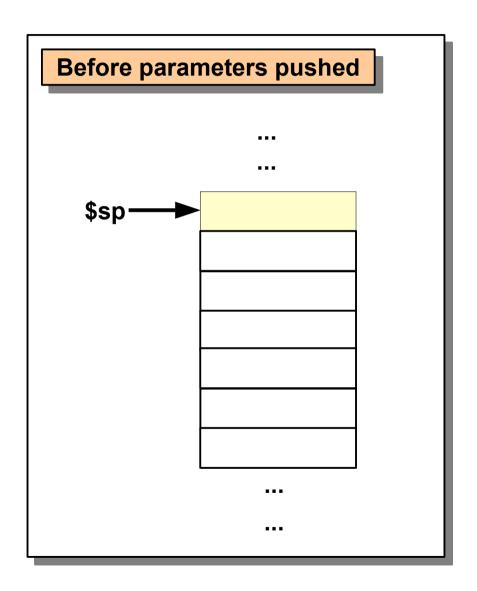
Pop into \$t1



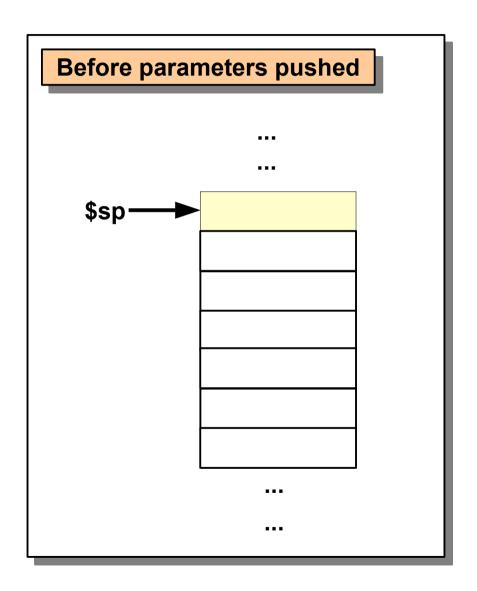
Pop into \$t1

```
Pop
lw $t1, 0($sp)
addi $sp, $sp, +4
```

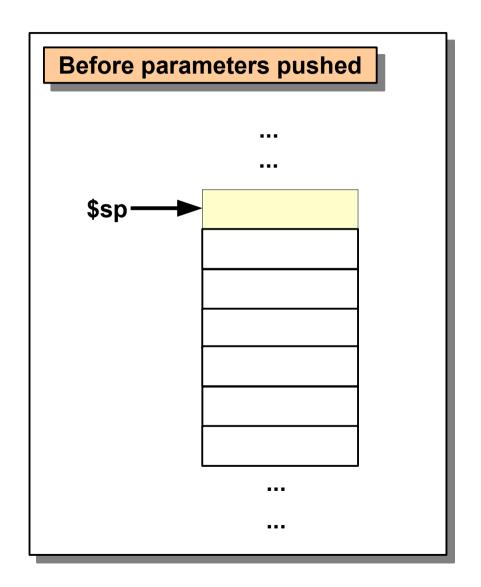
```
# main program
#6 parameters to func1
jal func1
```



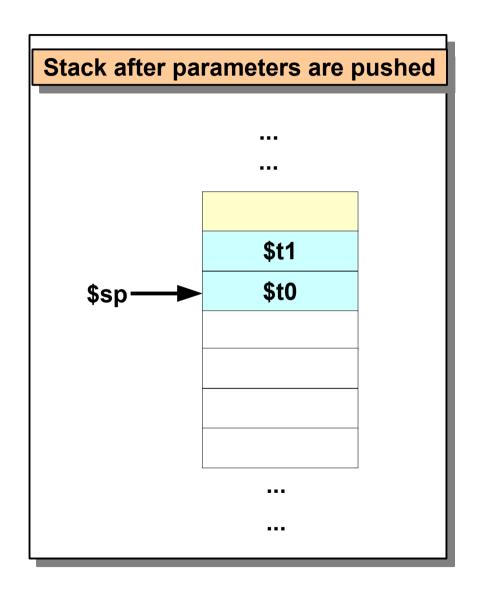
```
# main program
#6 parameters to func1
# 4 args are in $a0 - $a3
# push 2 on stack
jal func1
```



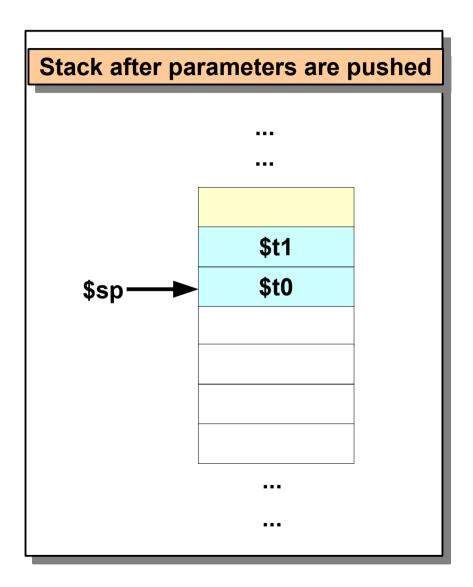
```
# main program
#6 parameters to func1
# 4 args are in $a0 - $a3
# push 2 on stack
addi $sp, $sp, -8
sw $t0, 0($sp)
sw $t1, -4($sp)
jal func1
```

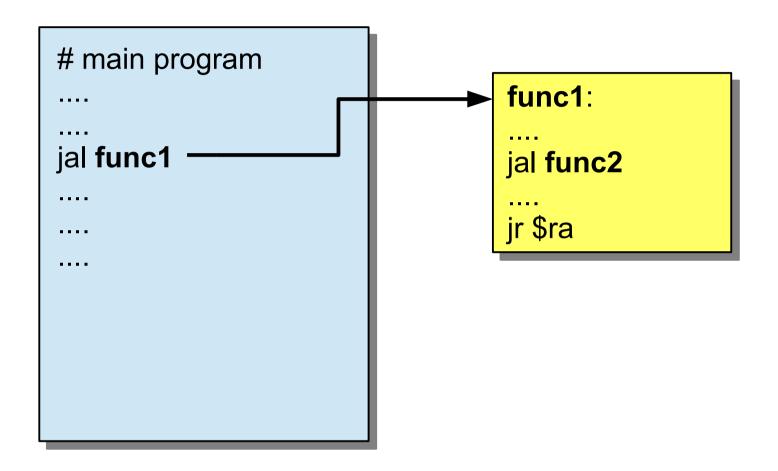


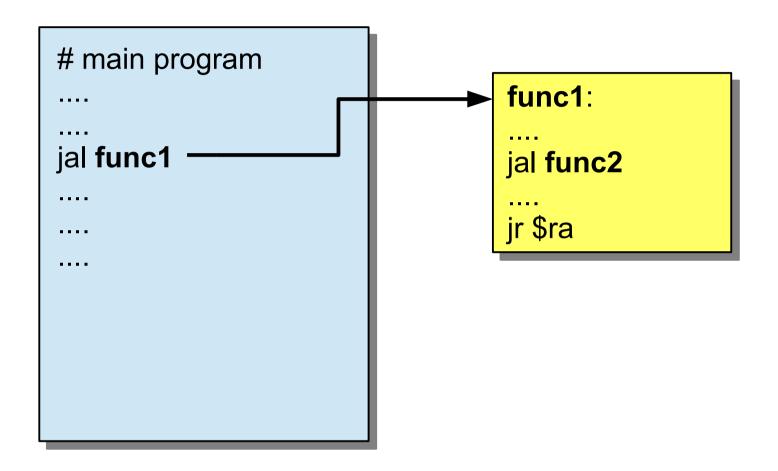
```
# main program
#6 parameters to func1
# 4 args are in $a0 - $a3
# push 2 on stack
addi $sp, $sp, -8
sw $t0, 0($sp)
sw $t1, -4($sp)
jal func1
```



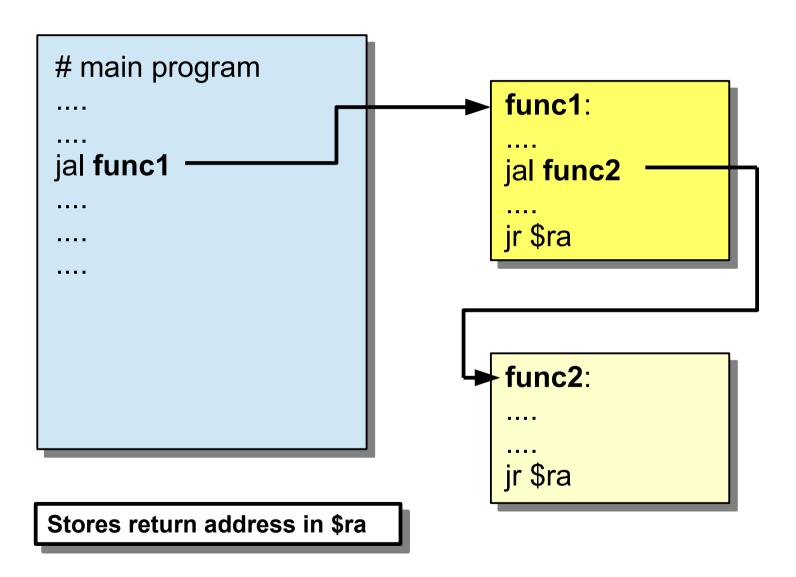
```
# main program
#6 parameters to func1
# 4 args are in $a0 - $a3
# push 2 on stack
addi $sp, $sp, -8
sw $t0, 0($sp)
sw $t1, -4($sp)
jal func1
            func1:
            lw $t4, 0($sp)
            lw $t5, -4($sp)
```

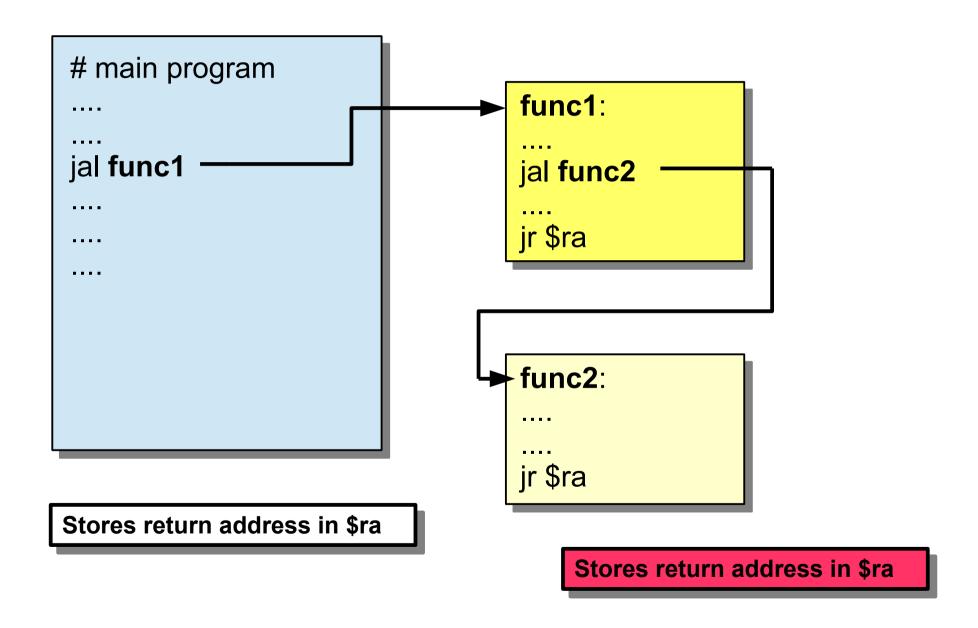




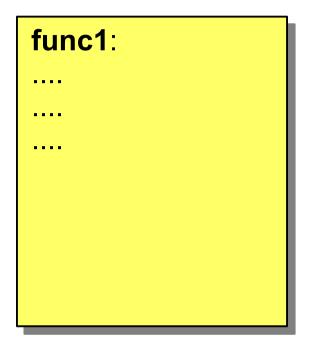


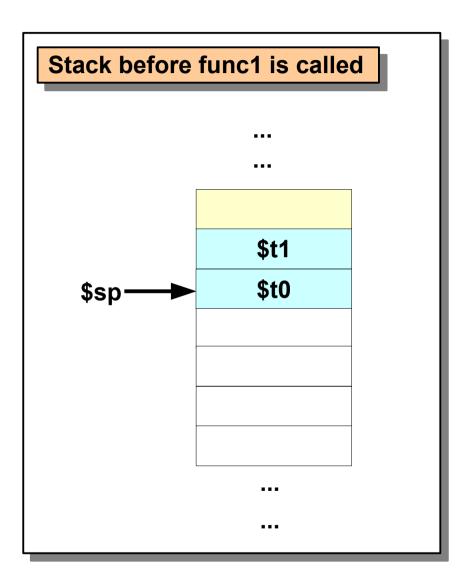
Stores return address in \$ra

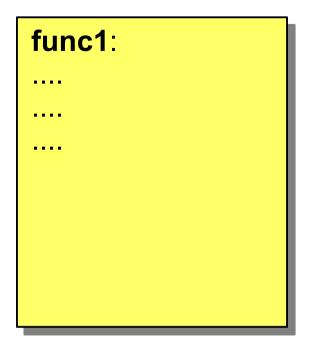


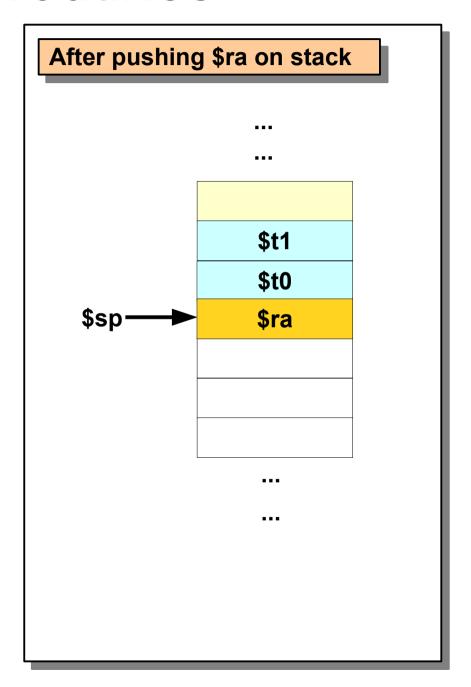


- func1 overwrites return address in \$ra (R31)
- Store the current return address in the program stack

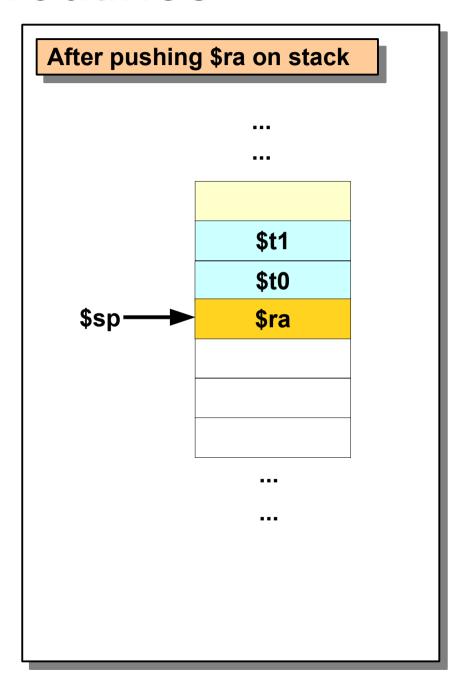






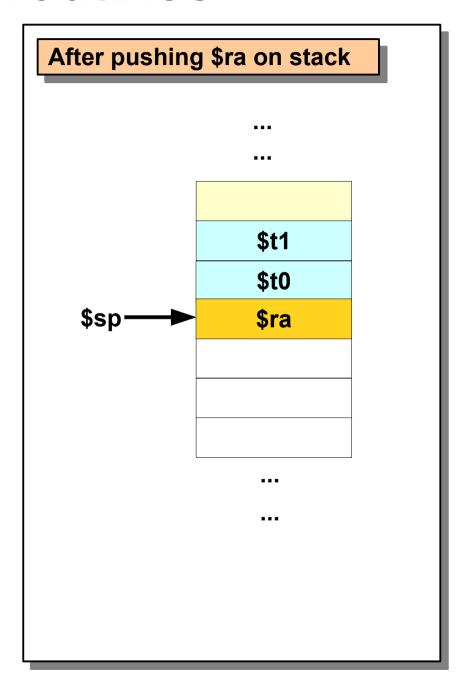


```
func1:
addi $sp, $sp, -4
sw $ra, 0($sp)
....
```



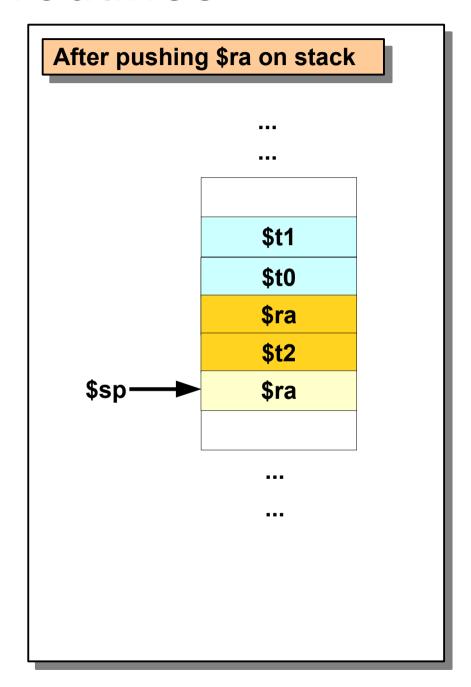
```
func1:
addi $sp, $sp, -4
sw $ra, 0($sp)
....
```

What does the stack look like after func1 passes contents of register \$t2 as a parameter to func2 and calls func2? Show the code changes in func1 and func2.

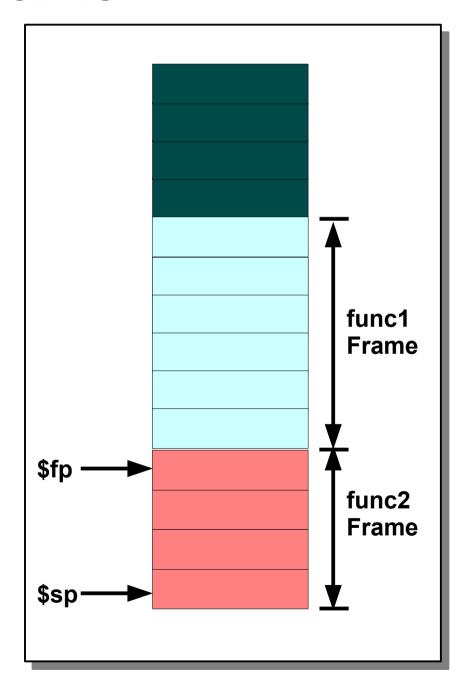


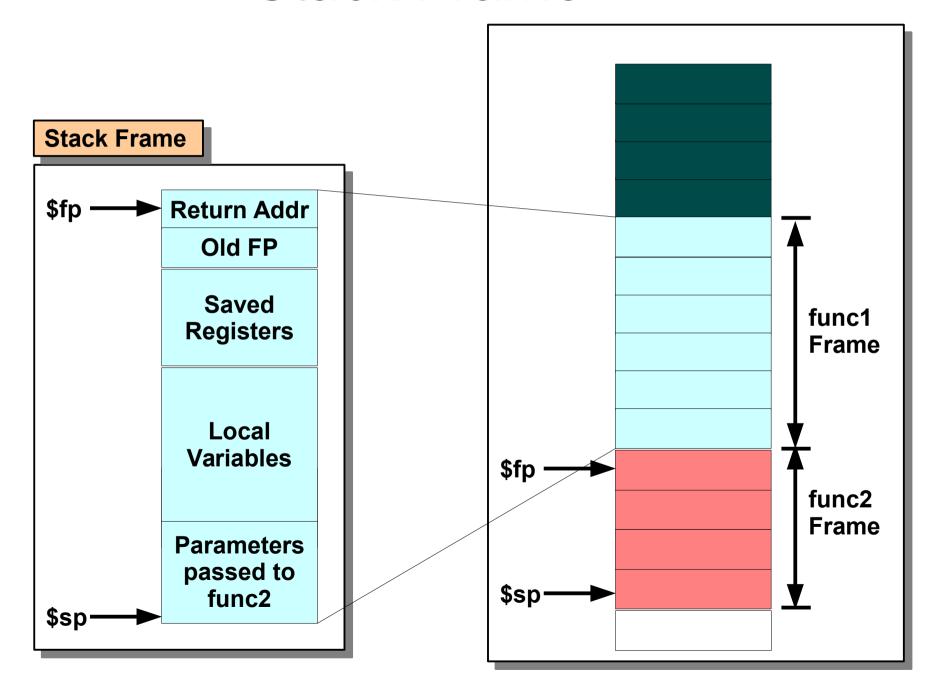
```
func1:
addi $sp, $sp, -4
sw $ra, 0($sp)
....
addi $sp, $sp, -4
sw $t2, 0($sp)
jal func2
....
```

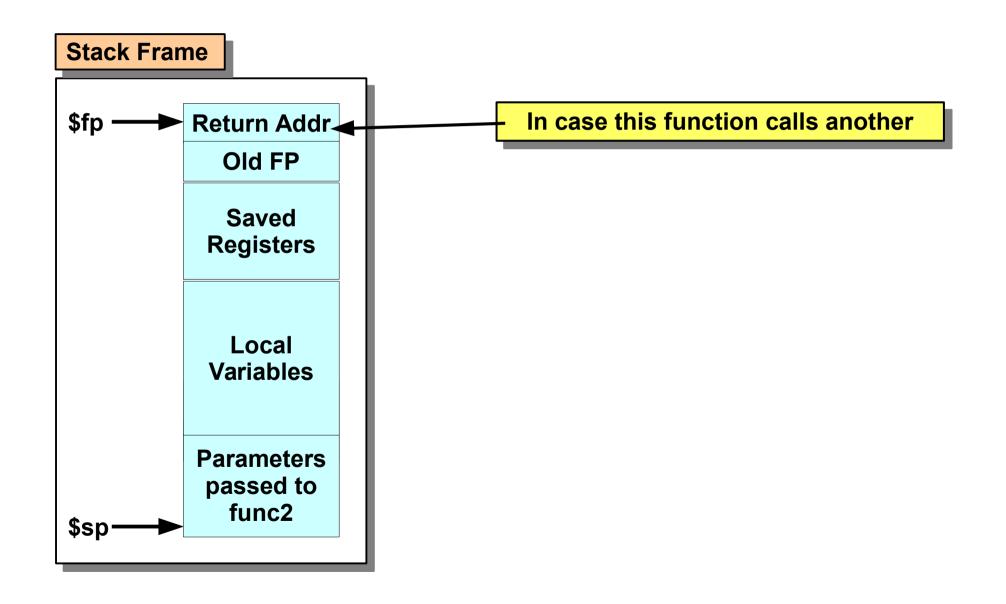
func2: addi \$sp, \$sp, -4 sw \$ra, 0(\$sp)

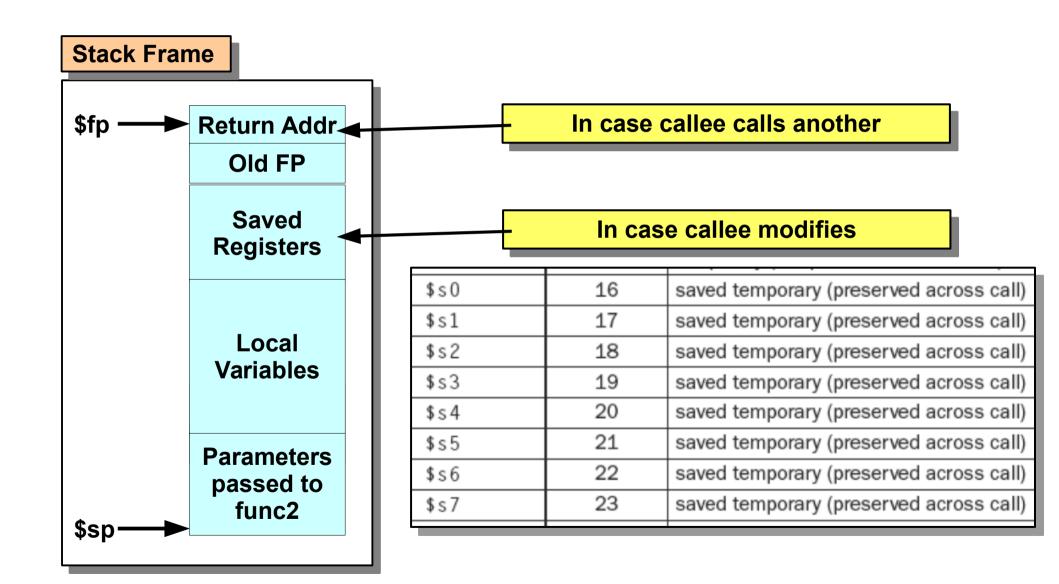


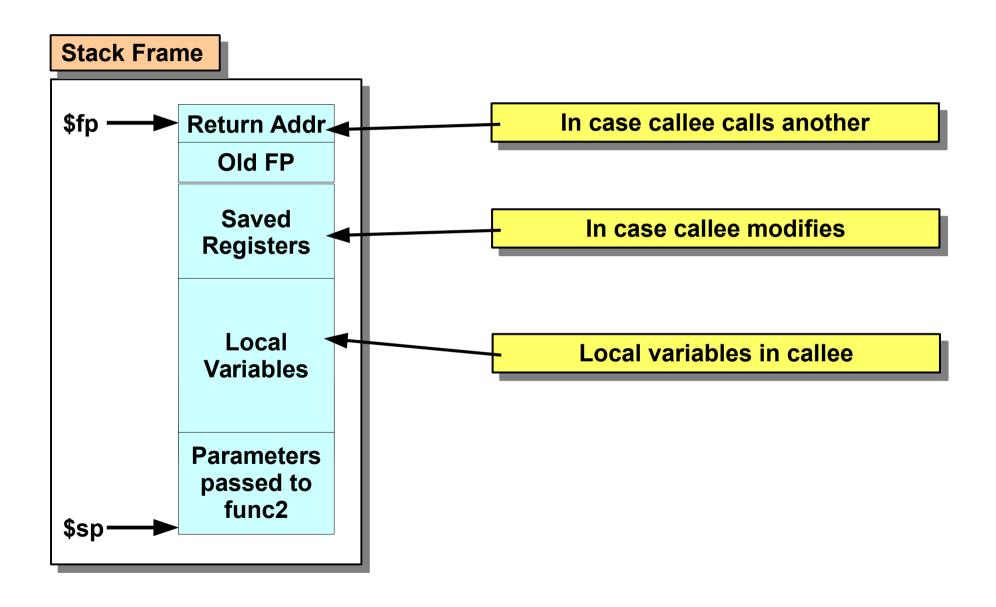
- Stack Frame: Private space for a subroutine allocated on entry and deallocated on exit
- Identified by a Frame Pointer (\$fp (R30))

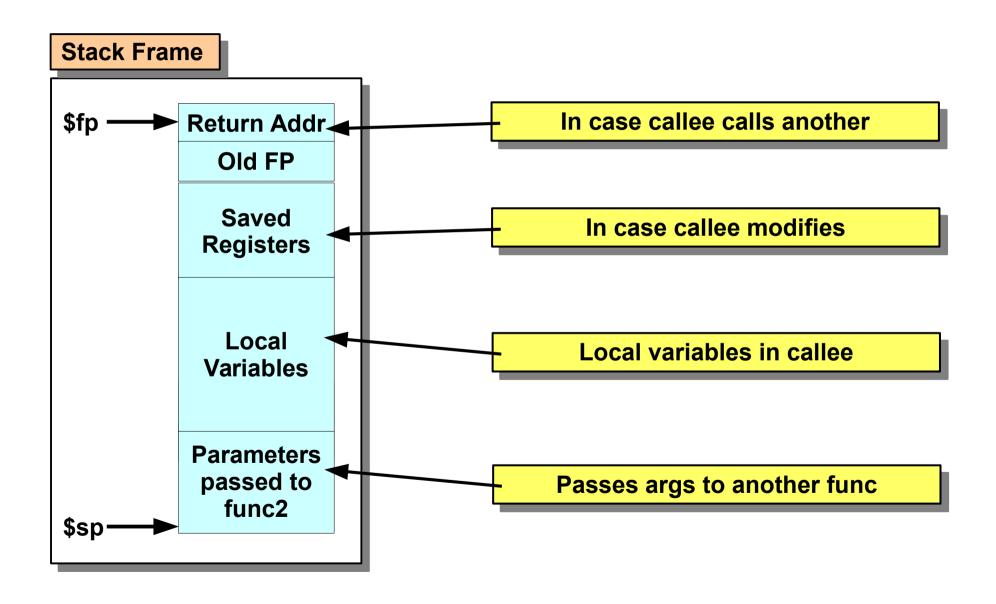


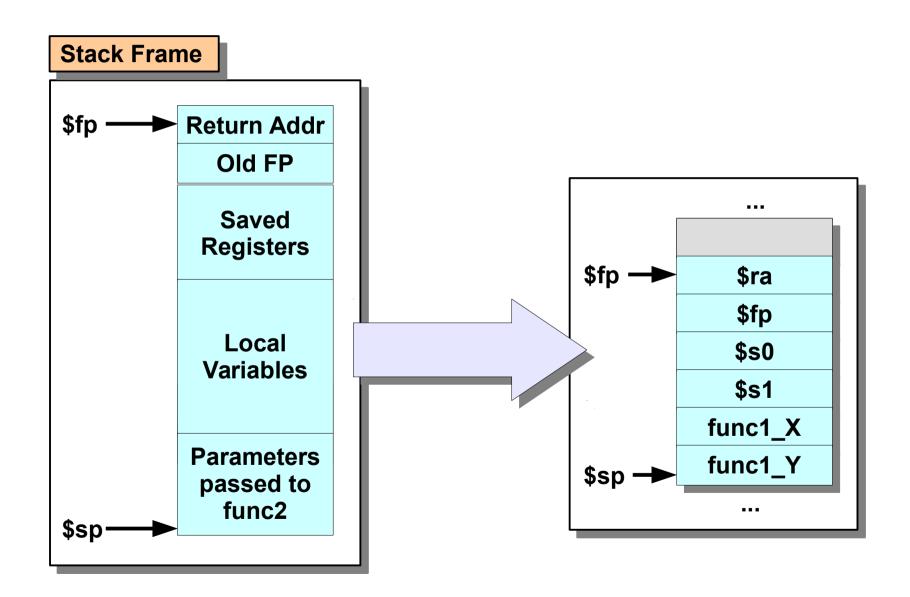




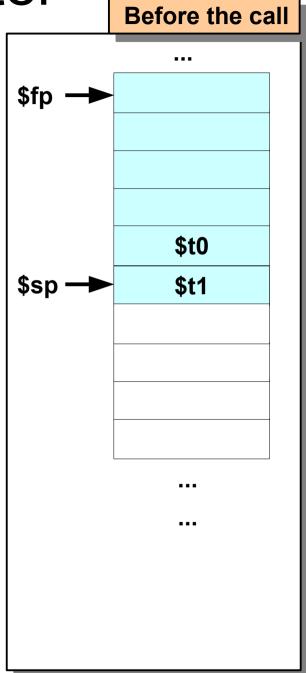




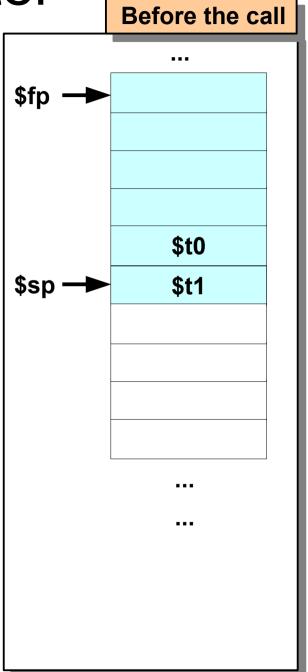




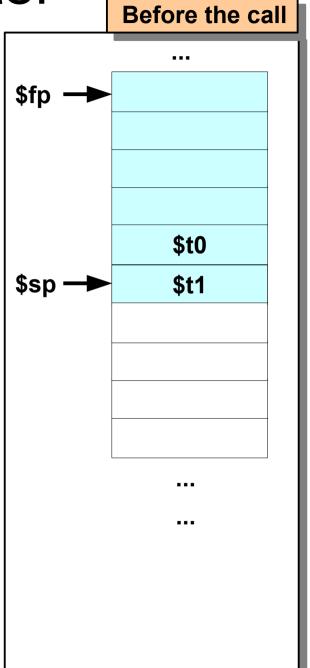
 After entry into a subroutine:



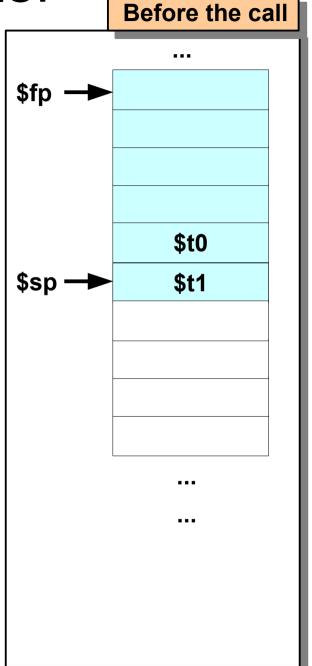
- After entry into a subroutine:
 - Save return address



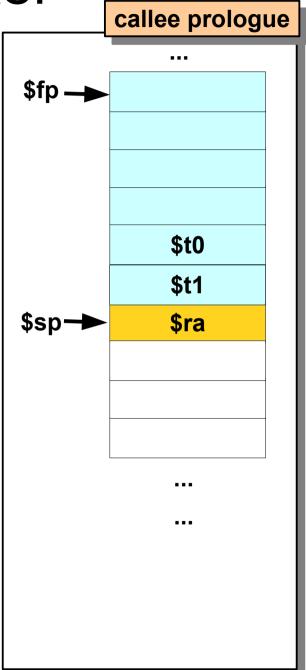
- After entry into a subroutine:
 - Save return address
 - Save frame pointer of the caller function



- After entry into a subroutine:
 - Save return address
 - Save frame pointer of the caller function
 - Point the frame pointer to the first location of stack frame of the current subroutine

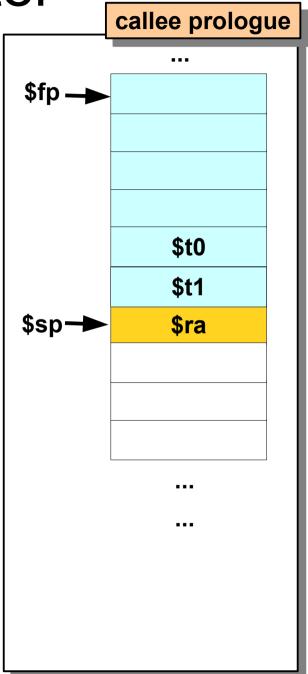


- After entry into a subroutine:
 - Save return address

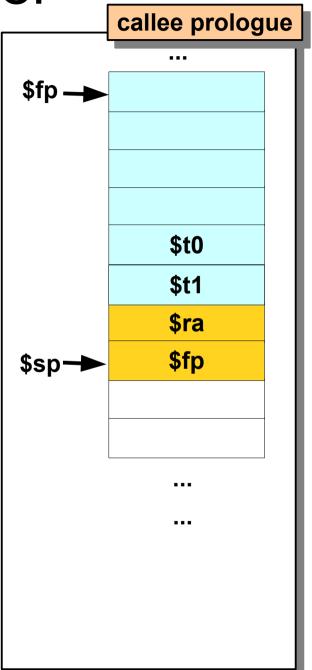


- After entry into a subroutine:
 - Save return address

addi \$sp, \$sp, -4 sw \$ra, 0(\$sp)

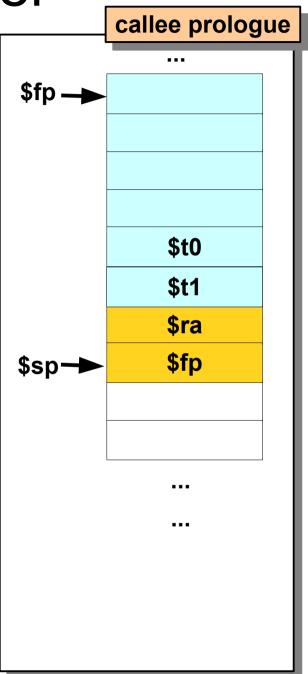


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 - Save return address
 - Save frame pointer of the caller function

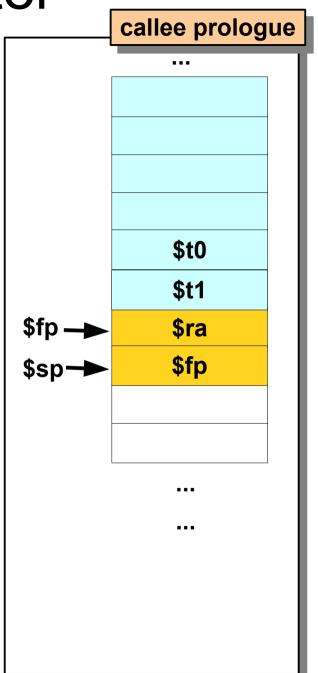


- After entry into a subroutine:
 - Save return address
 - Save frame pointer of the caller function

addi \$sp, \$sp, -4 sw \$fp, 0(\$sp)

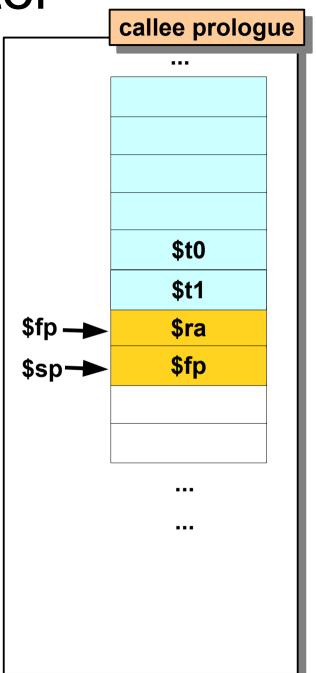


- After entry into a subroutine:
 - Save return address
 - Save frame pointer of the caller function
 - Point the frame pointer to the first location of stack frame of the current subroutine

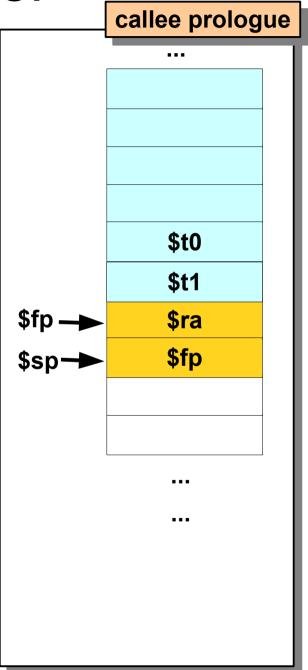


- After entry into a subroutine:
 - Save return address
 - Save frame pointer of the caller function
 - Point the frame pointer to the first location of stack frame of the current subroutine

addi \$fp, \$sp, 4

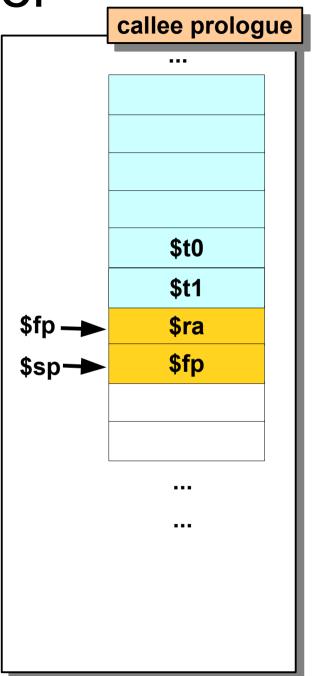


```
func1:
addi $sp, $sp, -8
sw $ra, 4($sp)
sw $fp, 0($sp)
addi $fp, $sp, 4
....
```

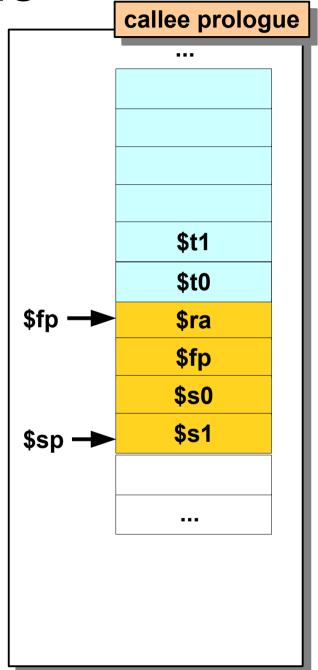


```
func1:
addi $sp, $sp, -8
sw $ra, 4($sp)
sw $fp, 0($sp)
addi $fp, $sp, 4
....
```

Parameters can be accessed in the callee function: 4(\$fp), 8(\$fp)

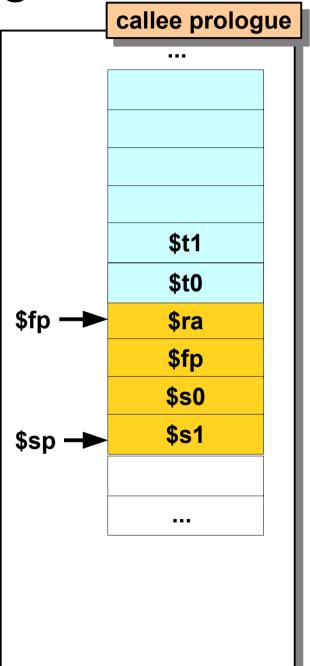


Save change registers

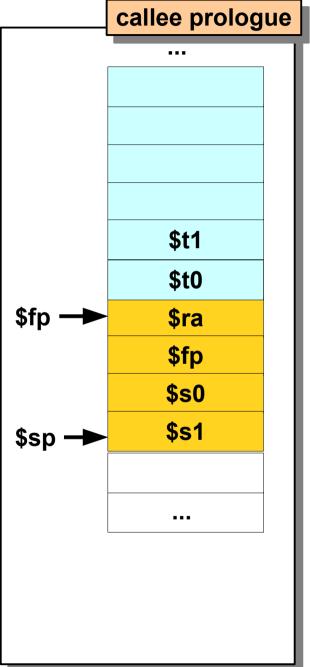


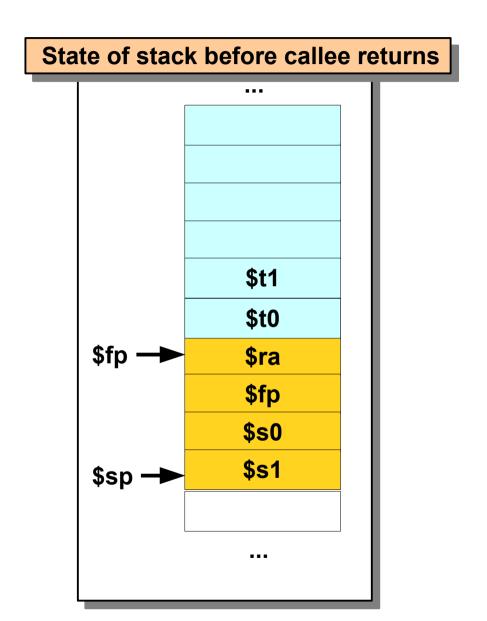
Save change registers

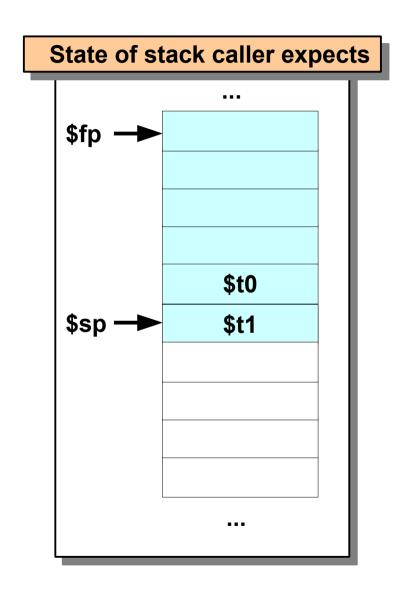
addi \$sp, \$sp, -8 sw \$s0, 4(\$sp) sw \$s1, 0(\$sp)



func1: addi \$sp, \$sp, -8 sw \$ra, 4(\$sp) sw \$fp, 0(\$sp) addi \$fp, \$sp, 4 addi \$sp, \$sp, -8 sw \$s0, 4(\$sp) sw \$s1, 0(\$sp)

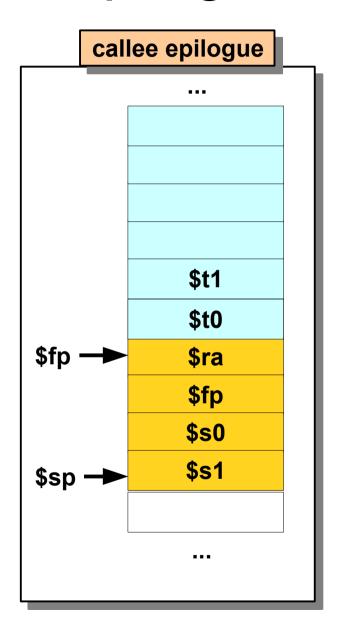






Stack Frame – Callee Epilogue

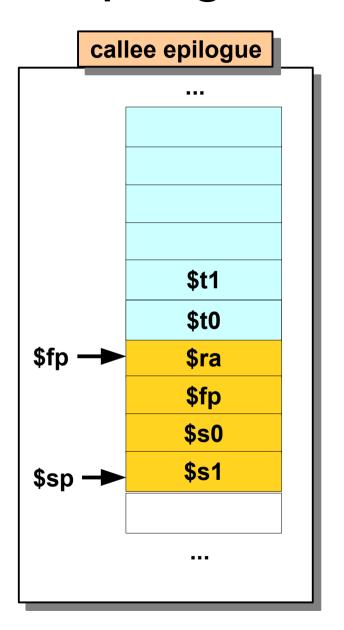
- Before return from callee subroutine:
 - Restore saved regs
 - Restore frame pointer of the caller function
 - Restore return address
 - Return



Stack Frame – Callee Epilogue

- Before return from callee subroutine:
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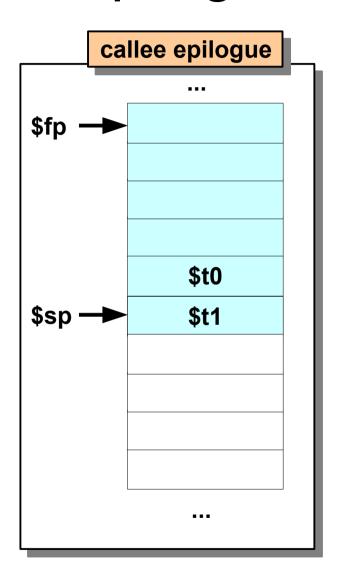
```
Iw $s1, 0($sp)
Iw $s0, 4($sp)
Iw $fp, 8($sp)
Iw $ra, 12($sp)
addi $sp, $sp, 16
jr $ra
```



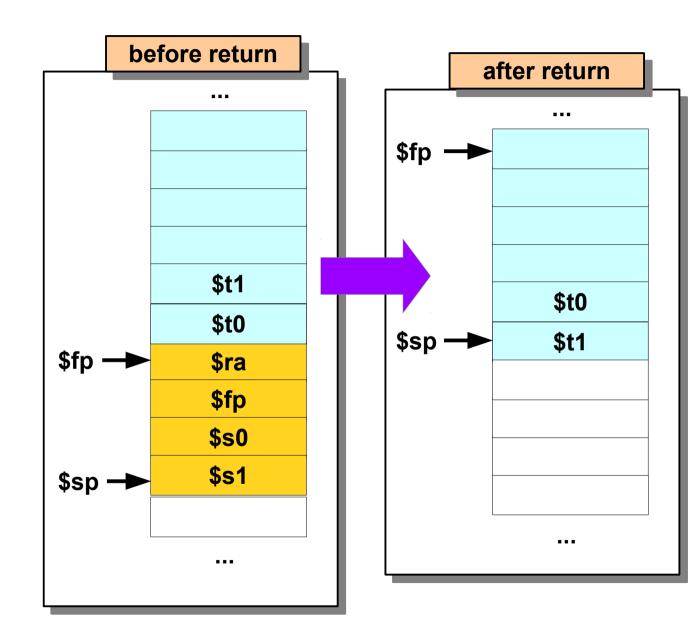
Stack Frame – Callee Epilogue

- Before return from callee subroutine:
 - Restore saved regs
 - Restore frame pointer of the caller function
 - Restore return address
 - Return

```
Iw $s1, 0($sp)
Iw $s0, 4($sp)
Iw $fp, 8($sp)
Iw $ra, 12($sp)
addi $sp, $sp, 16
jr $ra
```



func1: addi \$sp, \$sp, -8 sw \$ra, 4(\$sp) sw \$fp, 0(\$sp) addi \$fp, \$sp, 4 addi \$sp, \$sp, -8 sw \$s0, 4(\$sp) sw \$s1, 0(\$sp) # func1 code lw \$s1, 0(\$sp) lw \$s0, 4(\$sp) Iw \$fp, 8(\$sp) lw \$ra, 12(\$sp) addi \$sp, \$sp, 16 jr \$ra

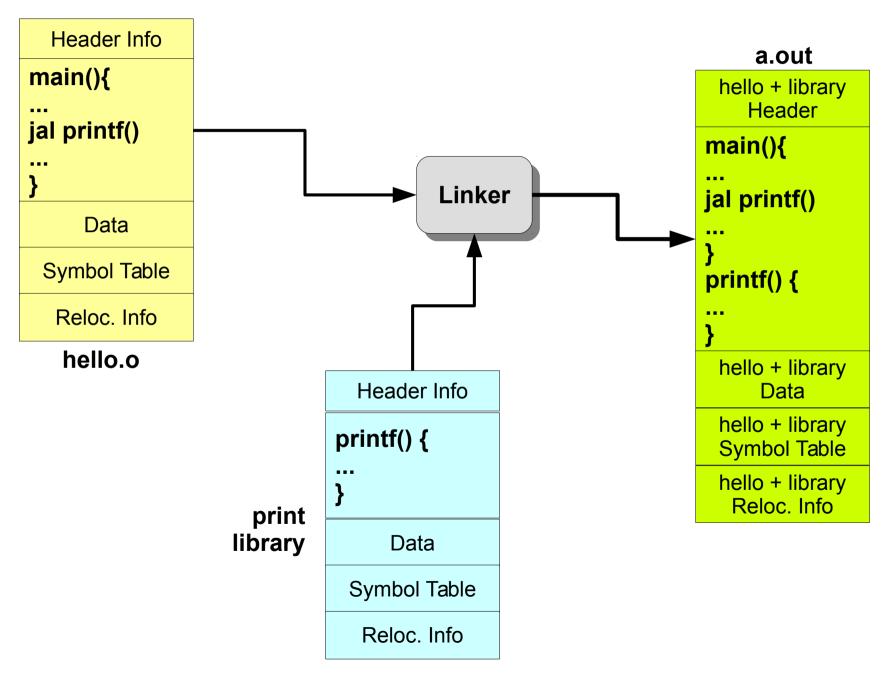


Module Outline

- Addressing modes. Instruction classes.
- MIPS-I ISA.
- Translating and starting a program.
- High level languages, Assembly languages and object code.
- Subroutine and subroutine call. Use of stack for handling subroutine call and return.

Backup

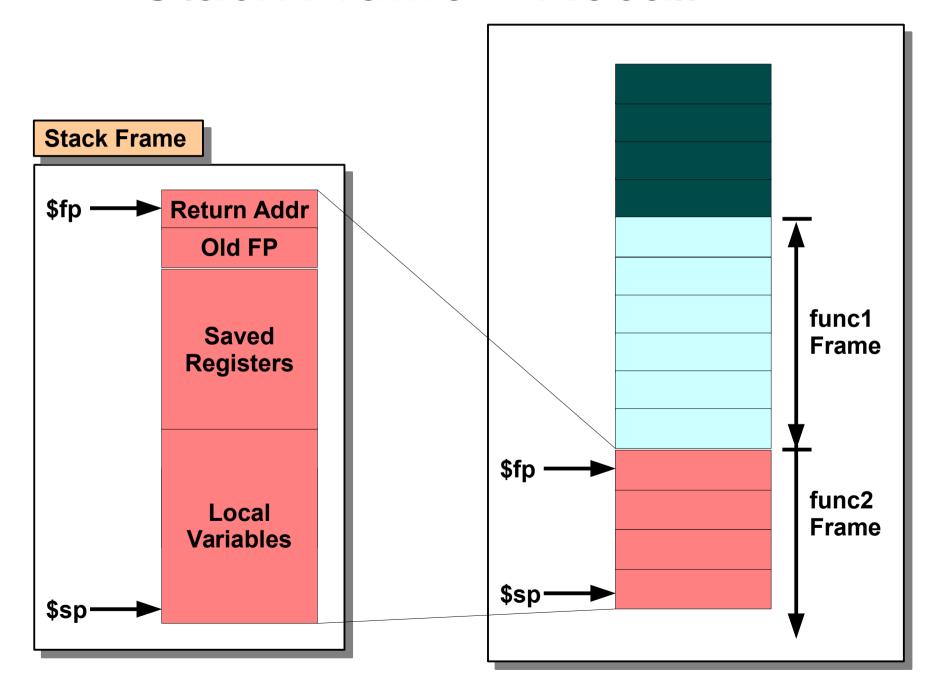
Linking Multiple Modules



The a.out executable

- What does the a.out file contain?
 - Program "code" (machine instructions)
 - Data values (values, size of arrays)
- Other information that is needed for
 - execution
 - debugging
 - Debugging: The stage in program development where mistakes ("bugs") in the program are identified

Stack Frame - Recall



Saved Registers

Registers 16 – 23 are saved across function calls

Name	Register number	Usage	Preserved on call?
\$zero	0	The constant value 0	n.a.
\$v0-\$v1	2–3	Values for results and expression evaluation	no
\$a0-\$a3	4–7	Arguments	no
\$t0-\$t7	8–15	Temporaries	no
\$s0 - \$s7	16–23	Saved	yes
\$t8-\$t9	24–25	More temporaries	no
\$gp	28	Global pointer	yes
\$sp	29	Stack pointer	yes
\$fp	30	Frame pointer	yes
\$ra	31	Return address	yes

Saved Registers

- Registers 16 23 are saved across function calls
- Save registers \$s0 \$s7 if used by the callee
- Example: \$s0, \$s1 are saved

 Local variables are allocated on the stack after the saved registers

