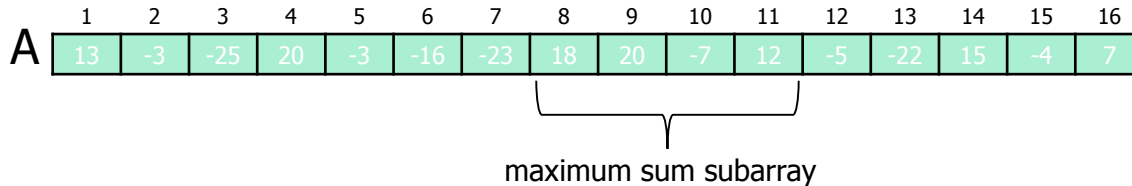




Maximum subarray Problem

Maximum Subarray Problem

- Given an array A with integers,
- Find the contiguous subarray of A whose values have the maximum sum ($A[i] + A[i + 1] + \dots + A[j]$)
- The maximum sum is zero if all the integers are **constraint** negative.
- An example
 - Number of possible ranges $[i, j]$ for n numbers in the array A?



Maximum Subarray Problem

- Possible ranges $[i, j]$

max {
■ $i = 1, j = 1, 2, \dots, n$
■ $i = 2, j = 2, 3, \dots, n$
■ ...
■ ...
■ $i = n-1, j = n-1, n$
■ $i = n, j = n$
} = Optimal Maximum Subarray

- There are $n(n+1)/2$ possible ranges

sum 이 될 수 있는 경우의 수는 $n(n+1)/2$

naive 하게 모든 pair의 sum을 구해 그중 최대인 값을 찾는다.



A Brute-force Solution

- Try every possible pair of range $[i, j]$ and compute $A[i] + A[i+1] + \dots + A[j]$.
- Since we have $\Theta(n^2)$ pairs, it takes $O(n^3)$ time.

A Brute-force Algorithm

■ FIND-MAXIMUM-SUBARRAY1(A)

1. MaxSum = 0

2. **for** i= 1 **to** n

3. **for** j= i **to** n

4. ThisSum = 0

5. **for** k= i **to** j

6. ThisSum = ThisSum + A[k]

7. **if** ThisSum > MaxSum

8. MaxSum = ThisSum

9. **return** MaxSum

← 이 부분을 반복할 필요는 없다.

■ Can you do better?



Actual Running Time

- For $n = 100$, actual time is 0.47 seconds on a particular computer.
- Can use this to estimate time for larger inputs:

$$T(n) = cn^3$$

$$T(10n) = c(10n)^3 = 1000cn^3 = 1000T(n)$$

- Inputs size increases by a factor of 10 means that running time increases by a factor of 1,000.
- For $n = 1,000$, estimate of running time is 470 seconds. (Actual running time was 449 seconds).
- For $n = 10,000$, estimate of running time is 449000 seconds (6 days).



How To Improve

- Remove a loop; not always possible.
- Here it is: innermost loop is unnecessary because it throws away information.
- **ThisSum** for next j is easily obtained from old value of **ThisSum**
 - Need $A[i] + A[i + 1] + \dots + A[j - 1] + A[j]$
 - Just computed $A[i] + A[i + 1] + \dots + A[j - 1]$
 - What we need is (*what we just computed*) + $A[j]$



A Better Brute-force Algorithm

- FIND-MAXIMUM-SUBARRAY2(A)

1. MaxSum = 0
2. **for** i= 1 **to** n
3. ThisSum = 0
4. **for** j= i **to** n
5. ThisSum = ThisSum + A[j]
6. **if** ThisSum > MaxSum
7. MaxSum = ThisSum
8. **return** MaxSum

- Can you do better?



Analysis

- Same logic as before: now the running time is quadratic, or $O(n^2)$.
- As we will see, this algorithm is still usable for inputs in the tens of thousands.
- Recall that the cubic algorithm was not practical for this amount of input.



Actual running time

- For $n = 100$, actual time is 0.011 seconds on the same particular computer.
- Can use this to estimate time for larger inputs:

$$T(n) = cn^2$$

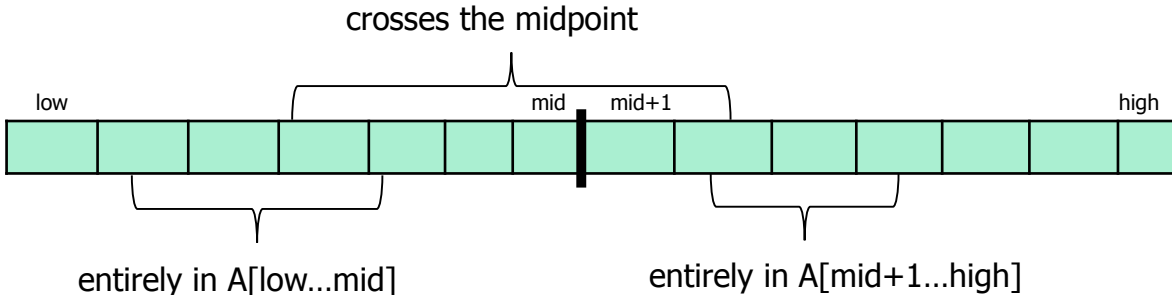
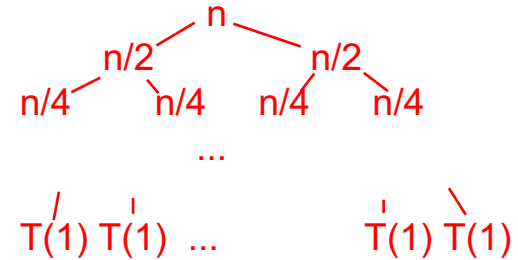
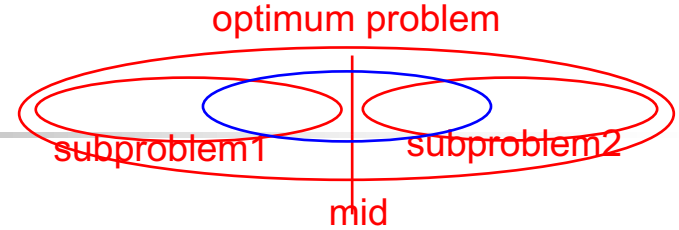
$$T(10n) = c(10n)^2 = 100cn^2 = 100T(n)$$

- Inputs size increases by a factor of 10 means that running time increases by a factor of 100.
- For $N = 1,000$, estimate of running time is 1.11 seconds. (Actual was 1.12 seconds).
- For $N = 10,000$, estimate of running time is 111 seconds.

Divide-and-Conquer Algorithm

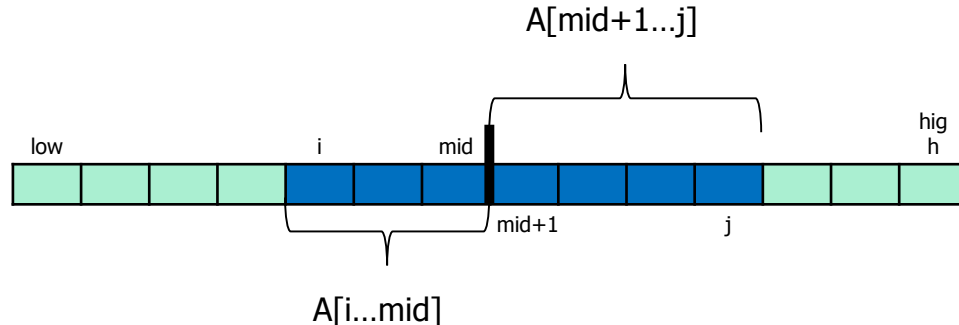
- The maximum subsequence either
 - lies entirely in the first half
 - lies entirely in the second half
 - starts somewhere in the first half, goes to the last element in the first half, continues at the first element in the second half, ends somewhere in the second half.
- Compute all three possibilities, and use the maximum.
- First two possibilities easily computed recursively.

*optimal sol = $\max\{\text{subproblem1}, \text{subproblem2}, \text{crossing problem}\}$
 $T(n/2)$ $T(n/2)$ n



Computing the Third Case

- Easily done with two loops.
- For maximum sum that starts in the first half and extends to the last element in the first half, use a right-to-left scan starting at the last element in the first half.
- For the other maximum sum, do a left-to-right scan, starting at the first element in the first half.





Analysis

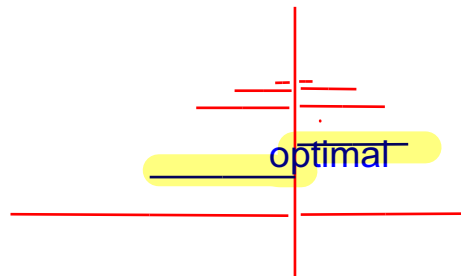
- Let $T(n)$ = the time for an algorithm to solve a problem of size N .
- Then $T(1) = 1$ (1 will be the quantum time unit; remember that constants don't matter).
- $T(n) = 2T\left(\frac{n}{2}\right) + n$
 - Two recursive calls, each of size $n/2$. The time to solve each recursive call is $T(n/2)$ by the above definition
 - Case three takes $O(n)$ time

A Divide-and-Conquer Algorithm

■ FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

$$\max_{mid}^{low} \text{sum} = \text{leftsum}$$



모든 leftsum을 계산해보고 최대인것을 택함
모든 rightsum ""

low				mid	mid+1				high
2	-3	5	-1	-2	-4	10	7	-2	-3

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. $\text{leftSum} = -\infty$
2. $\text{sum} = 0$
3. **for** $i = \text{mid}$ **downto** low
4. $\text{sum} = \text{sum} + A[i]$
5. **if** $\text{sum} > \text{leftSum}$
6. $\text{leftSum} = \text{sum}$
7. $\text{rightSum} = -\infty$
8. $\text{sum} = 0$
9. **for** $i = \text{mid}+1$ **to** high
10. $\text{sum} = \text{sum} + A[i]$
11. **if** $\text{sum} > \text{rightSum}$
12. $\text{rightSum} = \text{sum}$
13. **return** $\text{leftSum} + \text{rightSum}$

$\text{leftSum} = -\infty$
 $\text{sum} = 0$

low				mid		mid+1			high	
2	-3	5	-1	-2	-4	10	7	-2	-3	

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = $-\infty$
sum = 0

low			mid		mid+1	high			
2	-3	5	-1	-2	-4	10	7	-2	-3

A Divide-and-Conquer Algorithm

■ FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = $-\infty$
sum = -2

low				mid		mid+1			high	
2	-3	5	-1	-2	-4	10	7	-2	-3	

A Divide-and-Conquer Algorithm

■ FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = -2
sum = -2

low				mid		mid+1			high
2	-3	5	-1	-2	-4	10	7	-2	-3

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = -2
sum = -3

low				mid		mid+1			high
2	-3	5	-1	-2	-4	10	7	-2	-3

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = -2
sum = -3

low				mid		mid+1			high
2	-3	5	-1	-2	-4	10	7	-2	-3

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = -2
sum = 2

low				mid		mid+1			high
2	-3	5	-1	-2	-4	10	7	-2	-3

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = 2
sum = 2

low				mid		mid+1			high
2	-3	5	-1	-2	-4	10	7	-2	-3

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = 2
sum = -1

low				mid		mid+1			high	
2	-3	5	-1	-2	-4	10	7	-2	-3	

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = 2
sum = -1

low				mid		mid+1			high
2	-3	5	-1	-2	-4	10	7	-2	-3

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = 2
sum = 1

low				mid		mid+1			high	
2	-3	5	-1	-2	-4	10	7	-2	-3	

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = 2
sum = 1

low				mid		mid+1			high	
2	-3	5	-1	-2	-4	10	7	-2	-3	

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = 2
rightSum = $-\infty$
sum = 0

low				mid		mid+1			high	
2	-3	5	-1	-2	-4	10	7	-2	-3	

A Divide-and-Conquer Algorithm

■ FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = 2
rightSum = $-\infty$
sum = 0

low				mid		mid+1			high
2	-3	5	-1	-2	-4	10	7	-2	-3

A Divide-and-Conquer Algorithm

■ FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = 2
rightSum = $-\infty$
sum = -4

low				mid		mid+1			high
2	-3	5	-1	-2	-4	10	7	-2	-3

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = 2
rightSum = -4
sum = -4

low				mid		mid+1			high
2	-3	5	-1	-2	-4	10	7	-2	-3

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = 2
rightSum = -4
sum = 6

low				mid		mid+1			high
2	-3	5	-1	-2	-4	10	7	-2	-3

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = 2
rightSum = 6
sum = 6

low				mid		mid+1			high
2	-3	5	-1	-2	-4	10	7	-2	-3

A Divide-and-Conquer Algorithm

■ FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = 2
rightSum = 6
sum = 13

low				mid		mid+1			high	
2	-3	5	-1	-2	-4	10	7	-2	-3	

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = 2
rightSum = 13
sum = 13

low				mid		mid+1			high
2	-3	5	-1	-2	-4	10	7	-2	-3

A Divide-and-Conquer Algorithm

■ FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = 2
rightSum = 13
sum = 11

low				mid		mid+1			high
2	-3	5	-1	-2	-4	10	7	-2	-3

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = 2
rightSum = 13
sum = 11

low				mid		mid+1			high
2	-3	5	-1	-2	-4	10	7	-2	-3

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = 2
rightSum = 13
sum = 8

low				mid		mid+1			high
2	-3	5	-1	-2	-4	10	7	-2	-3

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

leftSum = 2
rightSum = 13
sum = 8

low				mid		mid+1			high
2	-3	5	-1	-2	-4	10	7	-2	-3

A Divide-and-Conquer Algorithm

- FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)

1. leftSum = $-\infty$
2. sum = 0
3. **for** i = mid **downto** low
4. sum = sum + A[i]
5. **if** sum > leftSum
6. leftSum = sum
7. rightSum = $-\infty$
8. sum = 0
9. **for** i = mid+1 **to** high
10. sum = sum + A[i]
11. **if** sum > rightSum
12. rightSum = sum
13. **return** leftSum+rightSum

return 15

leftSum = 2
rightSum = 13
sum = 8

low				mid		mid+1			high	
2	-3	5	-1	-2	-4	10	7	-2	-3	

A Divide-and-Conquer Algorithm

- FIND-MAXIMUM-SUBARRAY3(A, low, high)
 1. **if** high == low
 2. **return** (low, high, A[low]) // base case
 3. **else**
 4. mid = (low + high)/2
 5. leftSum = FIND-MAXIMUM-SUBARRAY3(A, low, mid)
 6. rightSum = FIND-MAXIMUM-SUBARRAY3(A, mid+1, high)
 7. crossSum = FIND-MAX-CROSSING-SUBARRAY(A, low, mid, high)
 8. **return** max(leftSum, rightSum, crossSum)





Incremental Algorithm

- Kadane's Algorithm
- If we know
 - the maximum subarray sum ending at position $A[i]$ (Call it ThisSum)
 - the maximum subarray sum for the range $[1, i]$ (Call it MaxSum)
- What is the maximum subarray sum for the range $[i, i+1]$?
 - $\max(\text{MaxSum}, A[i+1], \text{ThisSum} + A[i+1])$



A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3



A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0

2. ThisSum = 0

3. **for** j = 1 **to** n

4. ThisSum = ThisSum + A[j]

5. **if** ThisSum > MaxSum

6. MaxSum = ThisSum

7. **else if** ThisSum < 0

8. ThisSum = 0

9. **return** MaxSum

MaxSum = 0

ThisSum = 0

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3



A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 0
ThisSum = 0

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3



A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 0
ThisSum = 2
j = 1

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3



A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 2

ThisSum = 2

j = 1

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3





A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 2
ThisSum = -1
j = 2

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3

range of MaxSum





A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 2

ThisSum = -1

j = 2

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3

range of MaxSum





A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 2

ThisSum = 0

j = 2

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3





A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 2

ThisSum = 5

j = 3

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3





A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 5

ThisSum = 5

j = 3

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3





A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 5

ThisSum = 4

j = 4

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3





A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 5

ThisSum = 4

j = 4

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3





A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 5

ThisSum = 4

j = 4

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3





A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 5

ThisSum = 2

j = 5

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3





A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 5

ThisSum = 2

j = 5

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3





A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 5

ThisSum = 2

j = 5

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3





A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 5

ThisSum = -2

j = 6

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3





A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 5
ThisSum = -2
j = 6

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3

range of MaxSum





A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 5

ThisSum = 0

j = 6

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3





A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 5
ThisSum = 10
j = 7

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3

range of MaxSum





A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 10

ThisSum = 10

j = 7

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3



A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 10

ThisSum = 17

j = 8

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3



A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 17

ThisSum = 17

j = 8

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3



A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 17

ThisSum = 15

j = 9

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3



A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 17

ThisSum = 15

j = 9

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3



A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 17

ThisSum = 15

j = 9

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3



A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 17

ThisSum = 12

j = 10

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3



A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 17
ThisSum = 12
j = 10

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3

range of MaxSum



A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 17

ThisSum = 12

j = 10

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3



A Linear-time Algorithm

- FIND-MAXIMUM-SUBARRAY4(A)

1. MaxSum = 0
2. ThisSum = 0
3. **for** j = 1 **to** n
4. ThisSum = ThisSum + A[j]
5. **if** ThisSum > MaxSum
6. MaxSum = ThisSum
7. **else if** ThisSum < 0
8. ThisSum = 0
9. **return** MaxSum

MaxSum = 17

return 17

ThisSum = 12

j = 11

range of MaxSum

1	2	3	4	5	6	7	8	9	10
2	-3	5	-1	-2	-4	10	7	-2	-3



A Linear-time Algorithm

- Linear time (i.e., $O(n)$) algorithm would be best
- Running time is proportional to amount of input
- It makes only one pass through the data
- If the array is on a disk or is being transmitted over the Internet, it can be read sequentially, and there is no need to store any part of it in main memory
- At any point in time, the algorithm can correctly give an answer to the subsequence problem for the data it has already read: We call online algorithm