## **Lab: Abstract Interpretation**

(Week 8)

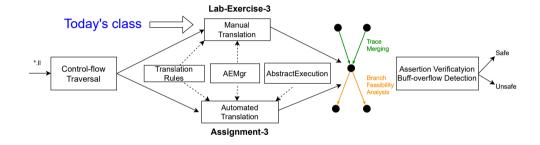
Yulei Sui

School of Computer Science and Engineering University of New South Wales, Australia

## **Lab-2 Marks and Lab-3 Code Template**

- Lab-2 marks are out and let us go through Quiz-2 and Exercise-2!
- Remember to git pull or docker pull to get the code template for Lab-Exercise-3

# Today's class



### Quiz-3 + Lab-Exercise-3

- Quiz-3 (5 points) (due date: 23:59, Tuesday, Week 10)
  - Abstract domain and soundness
  - Handling loops with widening and narrowing
- Lab-Exercise-3 (5 points) (due date: 23:59, Tuesday, Week 10)
  - Goal: Coding exercise to manually update abstract trace based on abstract execution rules and verify the assertions embedded in the code.
  - Specification: https://github.com/SVF-tools/ Software-Security-Analysis/wiki/Lab-Exercise-3

# Lab-3 Exercise: Manual Translation to Compute Abstract States

- Let us look at how to write abstract execution code to analyze examples of a loop-free and a loop C-like code by manually collecting abstract states at each program statement to form the abstract trace
- You will need to finish all the coding tests in AEMgr.cpp under Lab-Exercise-3

```
1 struct A{int f0;};
2 void main() {
     struct A * p;
     int * q :
    int x
    p = malloc;
     q = \&(p \rightarrow f0);
     *a = 10:
     x = *a:
10
    svf_assert(x == 10);
11 }
```

```
1 NodeID p = getNodeID("p");
2 NodeID q = getNodeID("q");
3 NodeID x = getNodeID("x");
4 ...
```

```
-----Var and Value-----
```

AEState:printAbstractState()

Source code

Translation for Abstract execution

```
1 struct A{int f0;};
2 void main() {
     struct A * p:
     int*q:
     int x:
    p = malloc;
     q = \&(p \rightarrow f0);
     *a = 10:
     x = *a:
10
    svf_assert(x == 10);
11 }
```

```
NodeID p = getNodeID("p");
NodeID q = getNodeID("q");
NodeID x = getNodeID("x");
NodeID malloc = getNodeID("malloc");
as[p] = AddressValue(getMemObjAddress("malloc"));
...
```

```
Var1 (p) Value: 0x7f000004
------
0x7f000004 (or 2130706436 in decimal)
```

represents the virtual memory address of this object Each SVF object starts with 0x7f + its ID.

Source code Translation for Abstract execution

```
1 struct A{int f0;};
2 void main() {
     struct A * p:
     int*q:
     int x:
     p = malloc;
     q = \&(p \rightarrow f0);
     *a = 10:
     x = *a:
10
    svf_assert(x == 10);
11 }
```

```
NodeID p = getNodeID("p");
NodeID q = getNodeID("q");
NodeID x = getNodeID("x");
NodeID malloc = getNodeID("malloc");
Sas[p] = AddressValue(getMemObjAddress("malloc"));
as[q] = AddressValue(getGepObjAddress("p", 0));
...
```

```
Var2 (q) Value: 0x7f000001
Var1 (p) Value: 0x7f000004
```

getGepObjAddress returns the field address of the aggregate object p The virual address also in the form of 0x7f... + VarID

Source code

Translation for Abstract execution

```
1 struct A{int f0;};
2 void main() {
     struct A * p:
     int*q:
    int x:
     p = malloc;
     q = \&(p \rightarrow f0);
     *a = 10:
     x = *a:
10
    svf_assert(x == 10);
11 }
```

```
| NodeID p = getNodeID("p");

2 NodeID q = getNodeID("q");

3 NodeID x = getNodeID("x");

4 NodeID malloc = getNodeID("malloc");

5 as[p] = AddressValue(getNemDbjAddress("malloc"));

6 as[q] = AddressValue(getGepDbjAddress("p", 0));

7 as.storeValue(q, IntervalValue(10, 10));

8 as[x] = as.loadValue(q);
```

```
Var3 (x) Value: [10, 10]
Var2 (q) Value: 0x7f000001
Var1 (p) Value: 0x7f000004
Var5 (0x7f000001) Value: [10, 10]
```

store value of 5 to address ox7f000005

load the value from ox7f000005 to x

Source code

Translation for Abstract execution

```
1 struct A{int f0;};
2 void main() {
     struct A * p:
     int*q:
     int x:
     p = malloc;
     q = \&(p \rightarrow f0);
     *q = 10:
     x = *a:
10
    svf_assert(x == 10);
11 }
```

```
1 NodeID p = getNodeID("p", 1);
2 NodeID q = getNodeID("q");
3 NodeID x = getNodeID("x");
5 NodeID x = getNodeID("x");
5 as[p] = AddressValue(getMemObjAddress("malloc"));
6 as[q] = AddressValue(getGepObjAddress("p", 0));
7 as.storeAulue(q, IntervalValue(10, 10));
8 as[x] = as.loadValue(q);
```

 ${\tt svf\_assert}$  checking is done in test.cpp.

```
Var3 (x) Value: [10, 10]
Var2 (q) Value: 0x7f000001
Var1 (p) Value: 0x7f000004
Var5 (0x7f000001) Value: [10, 10]
```

assertion checking

Source code

Translation for Abstract execution

```
1 int main(int argv) {
    // 5 \le argv \le 15
    int x = 10;
2 if(argv > 10)
3    x + +;
4 else
5    x + = 2;
6    svf_assert(x <= 12);
7 }</pre>
```

```
1 NodeID argy = getNodeID("argy");
2 as[argy] = IntervalValue(5, 15);
3 ...
```

```
-----Var and Value-------
Var1 (argv) Value: [5, 15]
```

assume  $5 \le \text{argv} \le 15$ 

Source code

Translation for Abstract execution

```
NodeID argv = getNodeID("argv");
2 as[argv] = IntervalValue(5, 15);
3 NodeID x = getNodeID("x");
4 as[x] = IntervalValue(10, 10);
5 ...
```

#### as:

```
Var1 (argv) Value: [5, 15]
Var2 (x) Value: [10, 10]
```

#### as\_true:

Source code

Translation for Abstract execution

```
1 NodeID argv = getNodeID("argv");
2 as[argv] = IntervalValue(5, 15);
3 NodeID x = getNodeID("x"):
 4 as[x] = IntervalValue(10, 10):
6 AEState as_after_if;
 7 AbstractValue cmp true = as[argv].getInterval() >
                           IntervalValue(10, 10):
9 // feasibility checking
10 cmp true.meet with(IntervalValue(1, 1)):
11 if (!cmp_true.getInterval().isBottom()) {
      AEState as true = as:
      as true[x] = as true[x].getInterval() +
                   IntervalValue(1, 1):
      //Join the states at the control-flow joint point
      as after if.joinWith(as true):
17 }
```

#### as:

```
------Var and Value-------
Var1 (argv) Value: [5, 15]
Var2 (x) Value: [10, 10]
```

#### as\_true:

```
Var1 (argv) Value: [5, 15]
Var2 (x) Value: [11, 11]
```

Source code

Translation for Abstract execution

```
1 int main(int argv) {
    int x = 10:
   if(argv > 10)
     x + +:
    else
     x + = 2:
7
   svf_assert(x \le 12);
```

```
2 AEState as after if:
 3 AbstractValue cmp true = as[argv].getInterval() >
                            IntervalValue(10, 10):
 5 // feasibility checking
 6 cmp true.meet with(IntervalValue(1, 1)):
 7 if (!cmp true.getInterval().isBottom()) {
      AEState as true = as:
      as_true[x] = as_true[x].getInterval() +
                    IntervalValue(1. 1):
      //Join the states at the control-flow joint point
      as after if.joinWith(as true):
13 }
15 AbstractValue cmp_false = as[argv].getInterval() >
                             IntervalValue(10, 10):
  cmp false.meet with(IntervalValue(0, 0)):
  if (!cmp_false.getInterval().isBottom()){
      AEState as false = as:
      as false[x] = as false[x].getInterval() +
                     IntervalValue(2, 2);
      as after if.joinWith(as false):
23 }
24 . . .
```

#### as:

```
------Var and Value--------
Var1 (argv) Value: [5, 15]
Var2 (x) Value: [10, 10]
```

#### as\_true:

```
------Var and Value--------
Var1 (argv) Value: [5, 15]
Var2 (x) Value: [11, 11]
```

#### as\_false:

Source code

Translation for Abstract execution

```
2 AEState as_after_if;
 3 AbstractValue cmp true = as[argv].getInterval() >
                            IntervalValue(10, 10):
 5 // feasibility checking
 6 cmp true.meet with(IntervalValue(1, 1)):
 7 if (!cmp true.getInterval().isBottom()) {
      AEState as true = as:
      as_true[x] = as_true[x].getInterval() +
                    IntervalValue(1, 1):
      //Join the states at the control-flow joint point
      as_after_if.joinWith(as_true);
13 }
15 AbstractValue cmp false = as[argv].getInterval() >
                             IntervalValue(10, 10):
  cmp_false.meet_with(IntervalValue(0, 0));
18 if (!cmp_false.getInterval().isBottom()){
      AEState as_false = as:
      as false[x] = as false[x].getInterval() +
                     IntervalValue(2, 2):
      as_after_if.joinWith(as_false);
23 }
24 as = as after if:
```

svf\_assert checking is done in test.cpp.

### Source code

Translation for Abstract execution

### Abstract trace

### as\_after\_if, as:

#### as\_true:

```
------Var and Value-------
Var1 (argv) Value: [5, 15]
Var2 (x) Value: [11, 11]
```

#### as\_false:

### **Before entering loop**

```
int main() {
   int a = 0;
   while(a < 10) {
        a + +;
    }
   svf_assert(a == 10);

return 0;
}</pre>
```

```
AEState entry_as;
2 AEState cur_head_as;
3 AEState body_as;
4 AEState exit_as;
5 u32_t widen_delay = 3;
6 
7 // Compose 'entry_as' (a = 0)
8 NodeID a = getNodeID("a");
9 entry_as[a] = IntervalValue(0, 0);
10 bool increasing = true;
11 for (int cur_iter = 0;; ++cur_iter) {
12 ...
13 }
14
```

### entry\_as

```
-----Var and Value-------Var1 (a) Value: [0, 0]
```

The initialization of a.

### Source code

Translation for Abstract execution

Abstract trace

### Implementation available here:

https://github.com/SVF-tools/Software-Security-Analysis/wiki/Lab-Exercise-3#

 ${\tt 4-widening-and-narrowing-implementation-for-the-below-loop-example-in-lecture-slides}$ 

### Widening delay stage

```
int main() {
  int a = 0;
  while(a < 10) {
    a + +;
  }
  svf_assert(a == 10);
  return 0;
  }
}</pre>
```

#### cur\_head\_as after Line 11:

### body\_as after Line 22:

Source code

Translation for Abstract execution

### Widening delay stage

```
int main() {
  int a = 0;
  while(a < 10) {
    a + +;
  }
  }
  svf_assert(a == 10);
  return 0;
  }
}</pre>
```

#### cur\_head\_as after Line 11:

```
Var1 (a) Value: [0, 1]
```

### body\_as after Line 22:

Source code

Translation for Abstract execution

### Widening delay stage

```
int main() {
  int a = 0;
  while(a < 10) {
    a + +;
  }
  svf_assert(a == 10);
  return 0;
  }
}</pre>
```

#### cur\_head\_as after Line 11:

```
Var1 (a) Value: [0, 2]
```

### body\_as after Line 22:

Source code

Translation for Abstract execution

### **Widening Stage**

```
int main() {
  int a = 0;
  while(a < 10) {
    a + +;
  }
  svf_assert(a == 10);
  return 0;
}</pre>
```

```
2 for (int cur iter = 0:: ++cur iter) {
      if (cur iter >= widen delay) {
          // Handle widening and narrowing after widen delay
          AEState prev head as = cur head as:
          // Update head's state by joining with 'entry_as' and 'body as'
          cur head as = entry as:
          cur head as.joinWith(body as):
          if (increasing) { // Widening phase
              AEState after widen = prev head as.widening(cur head as):
              cur head as = after widen:
              if (cur_head_as == prev_head_as) {
                  increasing = false;
                  continue:
          } else { // Narrow phase after widening
              AEState after narrow = prev head as.narrowing(cur head as):
              cur head as = after narrow:
              if (cur head as == prev head as) //fix-point reached
                  break:
      } else { // Handle widen delay
      // Handle loop body
28 // Handle loop exit
```

### prev\_head\_as after Line 5:

### cur\_head\_as after Line 11:

### body\_as after Line 26 (handle loop body):

```
-----Var and Value-------
Var1 (a) Value: [1, 10]
```

Widening stage where cur\_iter=3.

Source code

Translation for Abstract execution

### **Widening Stage**

```
int main() {
  int a = 0;
  while(a < 10) {
    a + +;
  }
  }
  svf_assert(a == 10);
  return 0;
  }
}</pre>
```

```
2 for (int cur iter = 0:: ++cur iter) {
      if (cur iter >= widen delay) {
          // Handle widening and narrowing after widen delay
          AEState prev head as = cur head as:
          // Update head's state by joining with 'entry_as' and 'body as
          cur head as = entry as:
          cur head as.joinWith(body as):
          if (increasing) { // Widening phase
              AEState after widen = prev head as.widening(cur head as):
              cur_head_as = after_widen;
              if (cur_head_as == prev_head_as) {
                  increasing = false;
                  continue:
          } else { // Narrow phase after widening
              AEState after narrow = prev head as.narrowing(cur head as):
              cur head as = after narrow:
              if (cur head as == prev head as) //fix-point reached
                  break:
      } else { // Handle widen delay
      // Handle loop body
28 // Handle loop exit
```

#### prev\_head\_as after Line 5:

### cur\_head\_as after Line 11:

Widening stage where cur\_iter=4.

Source code

Translation for Abstract execution

### **Narrowing Stage**

```
int main() {
   int a = 0;
   while(a < 10) {
      a + +;
   }
   svf_assert(a == 10);
   return 0;
   }
}</pre>
```

```
2 for (int cur iter = 0:: ++cur iter) {
      if (cur iter >= widen delay) {
          // Handle widening and narrowing after widen delay
          AEState prev head as = cur head as:
          // Update head's state by joining with 'entry_as' and 'body as'
          cur head as = entry as:
          cur head as.joinWith(body as):
          if (increasing) { // Widening phase
              AEState after widen = prev head as.widening(cur head as):
              cur head as = after widen:
              if (cur_head_as == prev_head_as) {
                  increasing = false;
                  continue:
          } else { // Narrow phase after widening
              AEState after narrow = prev head as.narrowing(cur head as):
              cur head as = after narrow:
              if (cur head as == prev head as) //fix-point reached
                  break:
      } else { // Handle widen delay
      // Handle loop body
28 // Handle loop exit
```

#### prev\_head\_as after Line 5:

### cur\_head\_as after Line 11:

```
-----Var and Value------
Var1 (a) Value: [0, 10]
```

### body\_as after Line 26 (handle loop body):

Narrowing stage where cur\_iter=5.

Source code

Translation for Abstract execution

### **Narrowing Stage**

```
int main() {
  int a = 0;
  while(a < 10) {
    a + +;
  }
  svf_assert(a == 10);
  return 0;
  }
}</pre>
```

```
2 for (int cur iter = 0:: ++cur iter) {
      if (cur iter >= widen delay) {
          // Handle widening and narrowing after widen delay
          AEState prev head as = cur head as:
          // Update head's state by joining with 'entry_as' and 'body as
          cur head as = entry as:
          cur head as.joinWith(body as):
          if (increasing) { // Widening phase
              AEState after widen = prev head as.widening(cur head as):
              cur_head_as = after_widen;
              if (cur_head_as == prev_head_as) {
                  increasing = false;
                  continue:
          } else { // Narrow phase after widening
              AEState after narrow = prev head as.narrowing(cur head as):
              cur head as = after narrow:
              if (cur head as == prev head as) //fix-point reached
                  break:
      } else { // Handle widen delay
      // Handle loop body
28 // Handle loop exit
```

#### prev\_head\_as after Line 5:

### cur\_head\_as after Line 11:

Narrowing stage where cur\_iter=6.

Source code

Translation for Abstract execution

### **Handle Loop Exit**

```
1 int main() {
2    int a = 0;
3    while(a < 10) {
4       a + +;
5    }
6    svf_assert(a == 10);
7    return 0;
8 }</pre>
```

### exit\_as after Line 7:

### exit\_as after Line 13:

Exiting loop.

Source code

Translation for Abstract execution

### **Handle Loop Exit**

```
1 int main() {
2    int a = 0;
3    while(a < 10) {
4        a + +;
5    }
6    svf_assert(a == 10);
7    return 0;
8 }</pre>
```

```
1
continuous for (int cur_iter = 0;; ++cur_iter) {
3
3
4
}
5 // Propagate head_as to loop exit
6 exit_as = cur_head_as;
7 // Process loop exit condition (a>=10)
8 exit_as[x].meet_with(IntervalValue(10, plus_infinity()));
9
10 return exit_as;
11
}
```

 $\mathtt{svf}\_\mathtt{assert}$  checking is done in  $\mathtt{test.cpp}$ .

### exit\_as at Line 15:

After analyzing loop.

Source code

Translation for Abstract execution

## **Software Verification Competition (SV-COMP)**

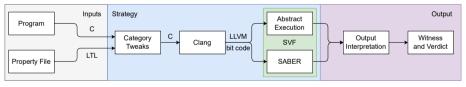
**Optional for Interested Students.** 

Cameron McGowan & Yulei Sui

School of Computer Science and Engineering University of New South Wales, Australia

# Recap - SVF-SVC in SV-COMP 2025

- In the 2025 competition, SVF-SVC participated in SV-COMP for the first time.
  - We built a Python wrapper around SVF which translated C files into an appropriate input format for SVF.
  - We used the specification category information to call SVF with the appropriate flags.
  - We interpreted SVF's output to generate witnesses using a basic format.



• Competition website: https://sv-comp.sosy-lab.org/

- Competition website: https://sv-comp.sosy-lab.org/
- Let's take a deeper look at the competition categories and properties to check. https://sv-comp.sosy-lab.org/2025/rules.php

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- Let's take a deeper look at the competition categories and properties to check. https://sv-comp.sosy-lab.org/2025/rules.php
- Let's look at how programs and witnesses work in SV-COMP.
- Example program: https://github.com/sosy-lab/sv-benchmarks/blob/master/c/loop-acceleration/multivar\_1-1.c
- Example witness: https://github.com/sosy-lab/sv-witnesses/blob/main/multivar\_1-1.c.invariant\_witness.yaml

- Competition website: https://sv-comp.sosy-lab.org/
- Let's take a deeper look at the competition categories and properties to check. https://sv-comp.sosy-lab.org/2025/rules.php
- Let's look at how programs and witnesses work in SV-COMP.
- Example program: https://github.com/sosy-lab/sv-benchmarks/blob/master/c/loop-acceleration/multivar\_1-1.c
- Example witness: https://github.com/sosy-lab/sv-witnesses/blob/main/multivar\_1-1.c.invariant\_witness.yaml
- SVF-SVC 2025 Paper: https://link.springer.com/chapter/10.1007/978-3-031-90660-2\_21
- SVF-SVC 2025 Github: https://github.com/Lasagnenator/svf-svc-comp