Code Verification Using Symbolic Execution

(Week 7)

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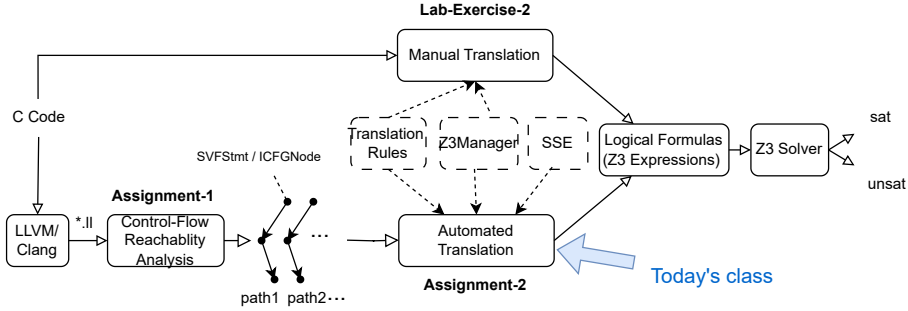
Code Verification Using Static Symbolic Execution



- We will detail the algorithms of translating branches and calls/returns.
- We will showcase branches and interprocedural examples.

Translate Branches and Calls/Returns

Algorithm 1: `translatePath(path)`

```
1 foreach edge ∈ path do
2     if intra ← dyn_cast⟨Intra⟩(edge) then
3         if handleIntra(intra) == false then
4             return false
5     else if call ← dyn_cast⟨CallEdge⟩(edge) then
6         handleCall(call)
7     else if ret ← dyn_cast⟨RetEdge⟩(edge) then
8         handleRet(ret)
9     return true
```

Algorithm 2: `handleIntra(intraEdge)`

```
1 if intraEdge.getCondition() then
2     if !handleBranch(intraEdge) then
3         return false;
4     else
5         return handleNonBranch(intraEdge);
6 else
7     return handleNonBranch(intraEdge);
```

Algorithm 3: `handleCall(callEdge)`

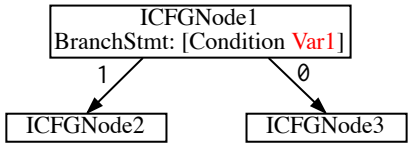
```
1 // Your code starts here ..;
2
3 // For each CallPE lhs = rhs:
4 // (1) Retrieve two z3 expressions ⟨[c], rhs⟩ and
  ⟨[c′], lhs⟩ under contexts c (caller's context) and
  c' (callee's context); Note that pushCallingCtx
  will make c become c'.
5 // (2) Add ⟨[c], rhs⟩ ≡ ⟨[c′], lhs⟩ into the solver.
```

Algorithm 4: `handleRet(retEdge)`

```
1 // Your code starts here ..;
2
3 // For each RetPE lhs = rhs:
4 // (1) Retrieve two z3 expressions ⟨[c′], rhs⟩ and
  ⟨[c], lhs⟩ under contexts c (caller's context) and
  c' (callee's context); Note that popCallingCtx
  will make c' become c.
5 // (2) Add ⟨[c′], rhs⟩ ≡ ⟨[c], lhs⟩ into the solver.
```

`getCondition()` and `getSuccessorCondValue()`

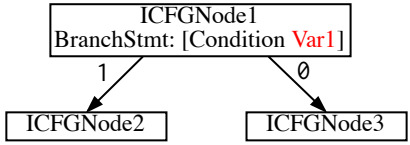
Given a `if/else` branch on the ICFG as the following:



- `edge` → `getCondition()` returns the branch condition (of type `SVFValue*` which is a **boolean** (for if/else) or a **numeric** (for switch).
  - `edge` → `getCondition()` returns `nullptr` if this IntraCFGEdge is not a branch.
- Given the condition value, you could obtain the ID of the corresponding SVFVar (`Var1`) via `edge` → `getCondition()` → `getId()`.

`getCondition()` and `getSuccessorCondValue()`

Given a `if/else` branch on the ICFG as the following:



- `edge` → `getSuccessorCondValue()` returns the actual condition `succCondValue`, whose value is either 1 (for if branch to execute) or 0 (for the else branch to execute).
  - For example, the `succCondValue` is 1 on the edge from ICFGNode1 to ICFGNode2, and 0 on the edge from ICFGNode1 to ICFGNode3.
- When evaluating the feasibility of a branch edge (e.g., ICFGNode1 to ICFGNode2) given an ICFG path, check `sat` of `Var1` == `succCondValue` against the solver's existing constraints.

Example 4: Branches

```c
void main(int x){
    int y;
    if(x > 10) {
        y = x + 1;
    }
    else {
        y = 10;
    }
    svf_assert(y >= x + 1);
}
```

Source code

```
define void @main(i32 %x) #0 {
entry:
  %cmp = icmp ugt i32 %x, 10
  br i1 %cmp, label %if.then, label %if.else

if.then:
  %add = add i32 %x, 1
  br label %if.end

if.else:
  br label %if.end

if.end: = %if.else, %if.then
  %y.0 = phi i32 [%add, %if.then], [10, %if.else]
  %add1 = add i32 %x, 1
  %cmp2 = icmp uge i32 %y.0, %add1
  call void @svf_assert(i1 zeroext %cmp2)
  ret void
}
```

LLVM IR

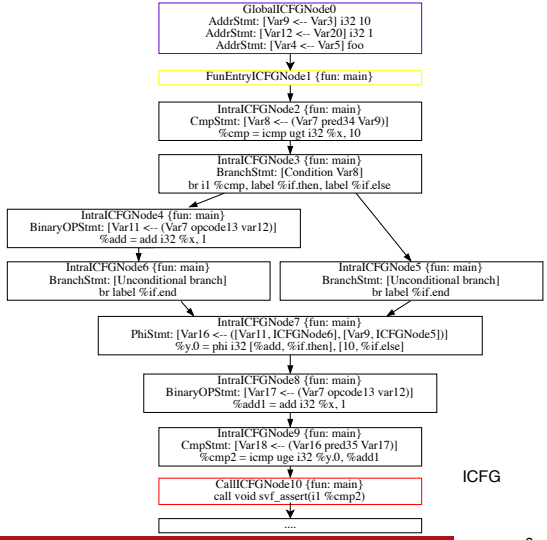
Example 4: Branches

```
1  define void @main(i32 %x) #0 {
2  entry:
3    %cmp = icmp ugt i32 %x, 10
4    br i1 %cmp, label %if.then, label %if.else
5
6  if.then:
7    %add = add i32 %x, 1
8    br label %if.end
9
10 if.else:
11   br label %if.end
12
13 if.end: = %if.else, %if.then
14   %y.0 = phi i32 [%add, %if.then], [10, %if.else]
15   %add1 = add i32 %x, 1
16   %cmp2 = icmp uge i32 %y.0, %add1
17   call void @svf_assert(i1 zeroext %cmp2)
18   ret void
19 }
```

LLVM IR

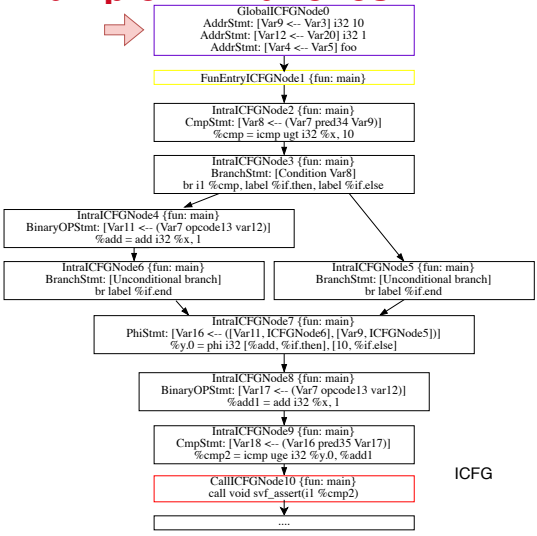
Two ICFG paths:

- `if.then branch`:
  0 → 1 → 2 → 3 → 4 → 6 → 7 → 8 → 9 → *svf_assert*
- `if.else branch`:
  0 → 1 → 2 → 3 → 5 → 7 → 8 → 9 → *svf_assert*



ICFG

Example 4: Branches



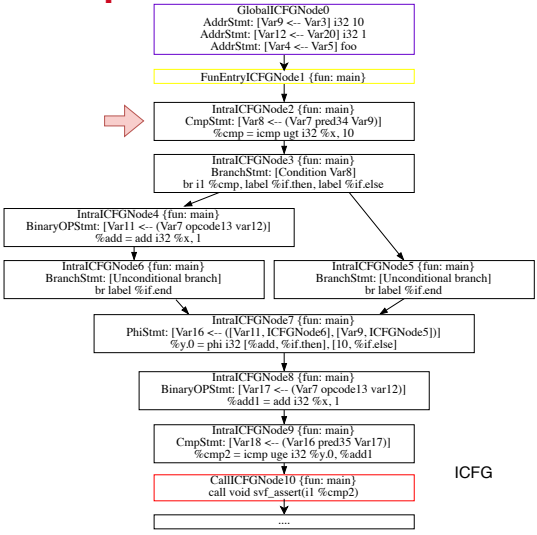
ICFG

Verifying ICFG path: 0 → 1 → 2 → 3 → 4 → 6 → 7 → 8 → 9 → *svf_assert* (if.then branch)

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var9} \equiv 10 \wedge \texttt{Var12} \equiv 1 \wedge \texttt{Var4} \equiv \texttt{0x7f000005}$

```
-------------SVFVar and Value-------------
ObjVar5 (0x7f000005)     Value: NULL
ValVar4                  Value: 0x7f000005
ValVar9                  Value: 10
ValVar12                 Value: 1
                  ...
------------------------------------------
```

Example 4: Branches



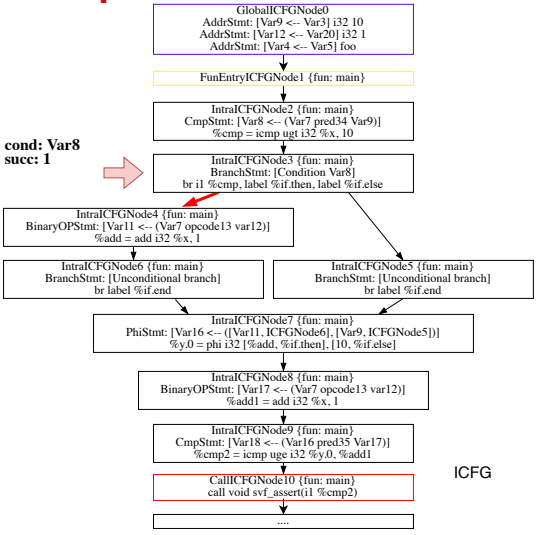
ICFG

Verifying ICFG path: $0 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 6 \rightarrow 7 \rightarrow 8 \rightarrow 9 \rightarrow$ *svf_assert* (if.then branch)

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var9} \equiv 10 \wedge \texttt{Var12} \equiv 1 \wedge \texttt{Var4} \equiv \texttt{0x7f000005}$
ICFGNode 2	$\wedge \quad \texttt{Var8} \equiv \texttt{ite(Var7} > \texttt{Var9, 1, 0)}$

```
-------------SVFVar and Value-------------
ObjVar5 (0x7f000005)    Value: NULL
ValVar4                 Value: 0x7f000005
ValVar9                 Value: 10
ValVar12                Value: 1
ValVar7                 Value: 11
ValVar8                 Value: 1
                  ...
------------------------------------------
```

Example 4: Branches



ICFG

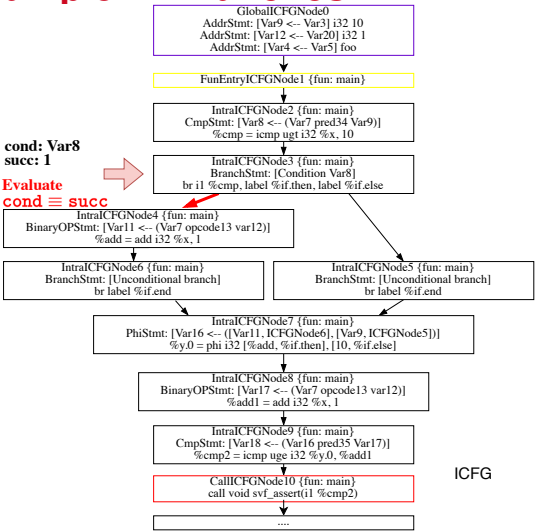
Algorithm 5: 3 `handleIntra(intraEdge)`

```
2  if intraEdge.getCondition() && !handleBranch(intraEdge) then
4      return false;
6  else
8      handleNonBranch(edge);
```

Algorithm 6: `handleBranch(intraEdge)`

```
1  cond = ... // Retrieve branch condition boolean var;
2  succ = ... // Edge's succ condition value (0 or 1);
3  // Evaluate path feasibility;
4  // The current path/edge is feasible if 'cond' and
     'succ' have the same value, when evaluated in the
     solver; otherwise this branch edge is infeasible;
5  if the brach edge is infeasible then
6      return false;
7  else
8      // add (cond=succ) to solver if it holds;
9      return true;
```

Example 4: Branches



ICFG

Algorithm 7: 3 `handleIntra(intraEdge)`

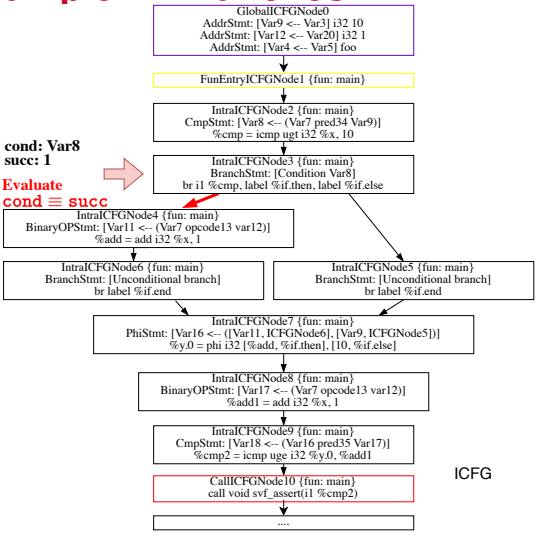
```
2  if intraEdge.getCondition() && !handleBranch(intraEdge) then
4      return false;
6  else
8      handleNonBranch(edge);
```

Algorithm 8: `handleBranch(intraEdge)`

```
1  cond = ... // Retrieve branch condition boolean var;
2  succ = ... // Edge's succ condition value (0 or 1);
3  // Evaluate path feasibility;
4  // The current path/edge is feasible if 'cond' and
     'succ' have the same value, when evaluated in the
     solver; otherwise this branch edge is infeasible;
5  if the brach edge is infeasible then
6      return false;
7  else
8      // add (cond≡succ) to solver if it holds;
9      return true;
```

Note: `getSolver().push()` creates a new stack frame for maintaining the newly added Z3 constraints.

Example 4: Branches



ICFG

Verifying ICFG path: 0 → 1 → 2 → 3 →4 →
6 → 7 → 8 → 9 → *svf_assert* (if.then branch)

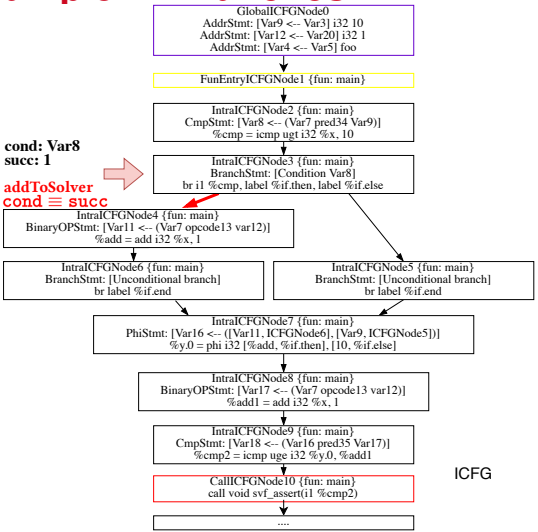
ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var9} \equiv 10 \wedge \texttt{Var12} \equiv 1 \wedge \texttt{Var4} \equiv \texttt{0x7f000005}$
ICFGNode 2	$\wedge \quad \texttt{Var8} \equiv \texttt{ite}(\texttt{Var7} > \texttt{Var9}, 1, 0)$
ICFGEdge 3 → 4	

The constraint $\texttt{Var8} \equiv 1$ is evaluated to be SAT.

The conditional ICFGEdge [ICFGNode3 → ICFGNode4] is feasible.

```
-------------SVFVar and Value-------------
ObjVar5 (0x7f000005)     Value: NULL
ValVar4                  Value: 0x7f000005
ValVar9                  Value: 10
ValVar12                 Value: 1
ValVar7                  Value: 11
ValVar8                  Value: 1
                  ...
------------------------------------------
```

Example 4: Branches



ICFG

Algorithm 9: 3 `handleIntra(intraEdge)`

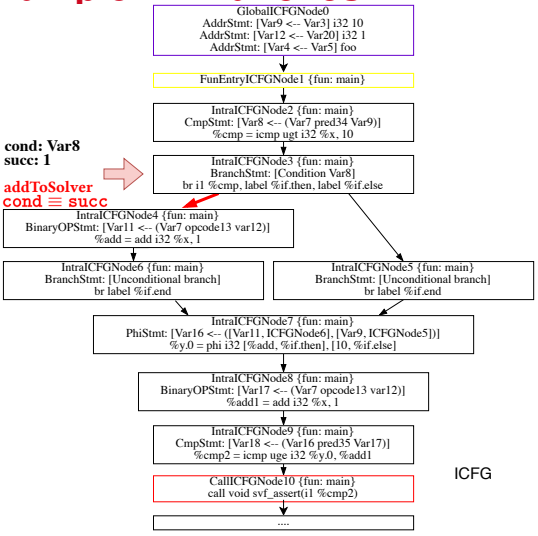
```
1 if intraEdge.getCondition() && !handleBranch(intraEdge) then
2     return false;
3 else
4     handleNonBranch(edge);
```

Algorithm 10: `handleBranch(intraEdge)`

```
1 cond = ... // Retrieve branch condition boolean var;
2 succ = ... // Edge's succ condition value (0 or 1);
3 // Evaluate path feasibility;
4 // The current path/edge is feasible if 'cond' and
    'succ' have the same value, when evaluated in the
    solver; otherwise this branch edge is infeasible;
5 if the brach edge is infeasible then
6     return false;
7 else
8     // add (cond≡succ) to solver if it holds;
9     return true;
```

Note: `res` is `sat`, so the conditional ICFGEdge [`ICFGNode4` ← `ICFGNode3`] is **feasible**!!

Example 4: Branches



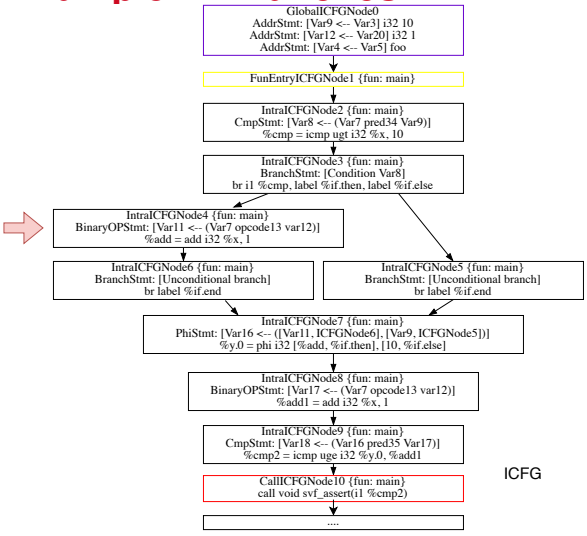
ICFG

Verifying ICFG path: $0 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 6 \rightarrow 7 \rightarrow 8 \rightarrow 9 \rightarrow$ *svf_assert* (if.then branch)

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var9} \equiv 10 \wedge \texttt{Var12} \equiv 1 \wedge \texttt{Var4} \equiv \texttt{0x7f000005}$
ICFGNode 2	$\wedge \quad \texttt{Var8} \equiv \texttt{ite}(\texttt{Var7} > \texttt{Var9}, 1, 0)$
ICFGEdge 3 $\rightarrow$ 4	$\wedge \quad \texttt{Var8} \equiv 1$

```
-------------SVFVar and Value-------------
ObjVar5 (0x7f000005)     Value: NULL
ValVar4                  Value: 0x7f000005
ValVar9                  Value: 10
ValVar12                 Value: 1
ValVar7                  Value: 11
ValVar8                  Value: 1
                  ...
------------------------------------------
```

Example 4: Branches



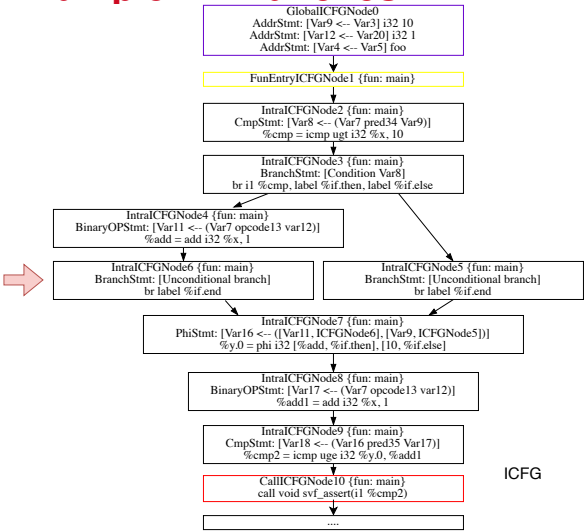
ICFG

Verifying ICFG path: $0 \to 1 \to 2 \to 3 \to 4 \to 6 \to 7 \to 8 \to 9 \to$ *svf_assert* (if.then branch)

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var9} \equiv 10 \wedge \texttt{Var12} \equiv 1 \wedge \texttt{Var4} \equiv \texttt{0x7f000005}$
ICFGNode 2	$\wedge \quad \texttt{Var8} \equiv \texttt{ite}(\texttt{Var7} > \texttt{Var9}, 1, 0)$
ICFGEdge 3 $\to$ 4	$\wedge \quad \texttt{Var8} \equiv 1$
ICFGNode 4	$\wedge \quad \texttt{Var11} \equiv \texttt{Var7} + \texttt{Var12}$

```
------------SVFVar and Value------------
ObjVar5 (0x7f000005)    Value: NULL
ValVar4                 Value: 0x7f000005
ValVar9                 Value: 10
ValVar12                Value: 1
ValVar7                 Value: 11
ValVar8                 Value: 1
ValVar11                Value: 12
                  ...
----------------------------------------
```

Example 4: Branches

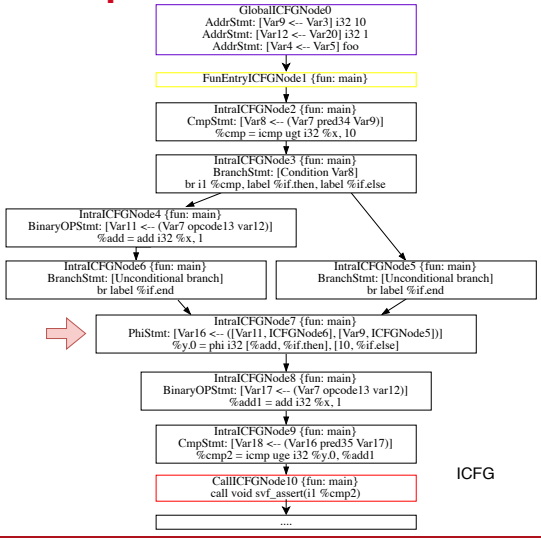


ICFG

Algorithm 11: 3 `handleIntra(intraEdge)`

```
2  if intraEdge.getCondition() &&
     !handleBranch(intraEdge) then
4      return false;
6  else
8      handleNonBranch(edge);
```

Example 4: Branches



ICFG

Verifying ICFG path: $0 \to 1 \to 2 \to 3 \to 4 \to 6 \to 7 \to 8 \to 9 \to$ *svf_assert* (if.then branch)

Algorithm 12: 3 `handlePhi(edge)`

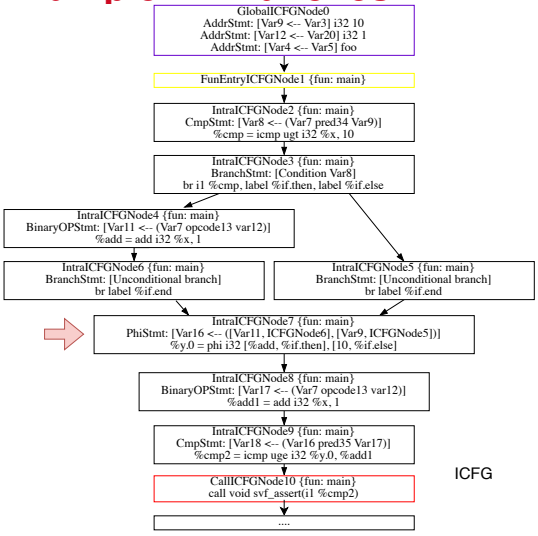
```
1 res ← getZ3Expr(phi.getResID());
2 opINodeFound ← false;
3 for i ← 0 to phi.getOpVarNum() − 1 do
4     if edge.srcNode() postdominates phi.getOpICFGNode(i)
      then
5         ope ← getZ3Expr(phi.getOpVar(i).getId());
6         addToSolver(res == ope);
7         opINodeFound ← true;
```

Given `Var16` ← ([`Var11`, `ICFGNode6`], [`Var9`, `ICFGNode5`]), only `Var16` ≡ `Var11` holds as we traverse the `if.then` branch from `ICFGNode6`, where `Var11`'s definition originates

`edge.srcNode()`: ICFGNode6
`phi.getOpICFGNode(i)`: ICFGNode where *i-th* phi operand var's definition originates.

ICFGNode *m* postdominates *n*: if all paths to the graph's exit starting at *n* must go through *m* (a node postdominates itself).

Example 4: Branches



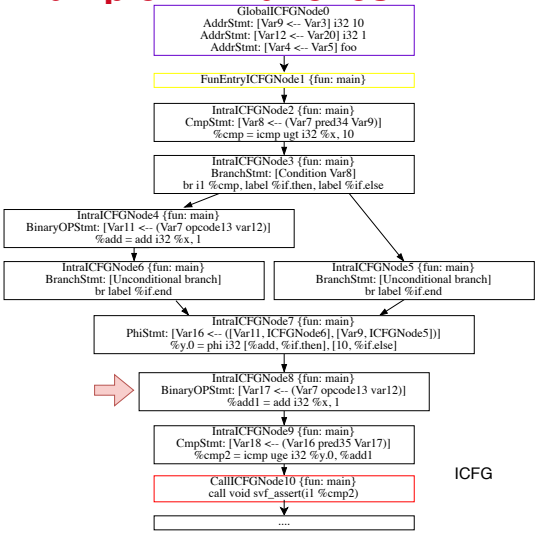
ICFG

Verifying ICFG path: $0 \to 1 \to 2 \to 3 \to 4 \to 6 \to 7 \to 8 \to 9 \to$ *svf_assert* (if.then branch)

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var9} \equiv 10 \wedge \texttt{Var12} \equiv 1 \wedge \texttt{Var4} \equiv \texttt{0x7f000005}$
ICFGNode 2	$\wedge \quad \texttt{Var8} \equiv \texttt{ite}(\texttt{Var7} > \texttt{Var9}, 1, 0)$
ICFGEdge 3 $\to$ 4	$\wedge \quad \texttt{Var8} \equiv 1$
ICFGNode 4	$\wedge \quad \texttt{Var11} \equiv \texttt{Var7} + \texttt{Var12}$
ICFGNode 7	$\wedge \quad \texttt{Var16} \equiv \texttt{Var11}$

```
-------------SVFVar and Value-------------
                    ...
ValVar9                 Value: 10
ValVar12                Value: 1
ValVar7                 Value: 11
ValVar8                 Value: 1
ValVar11                Value: 12
ValVar16                Value: 12
                    ...
------------------------------------------
```

Example 4: Branches



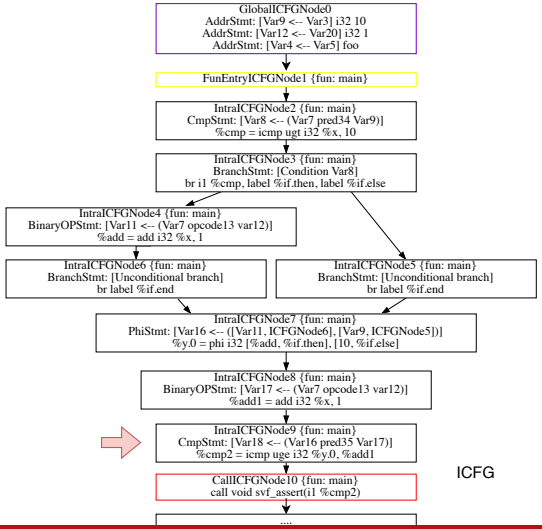
ICFG

Verifying ICFG path: $0 \to 1 \to 2 \to 3 \to 4 \to 6 \to 7 \to 8 \to 9 \to$ *svf_assert* (if.then branch)

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var9} \equiv 10 \wedge \texttt{Var12} \equiv 1 \wedge \texttt{Var4} \equiv \texttt{0x7f000005}$
ICFGNode 2	$\wedge \quad \texttt{Var8} \equiv \texttt{ite}(\texttt{Var7} > \texttt{Var9}, 1, 0)$
ICFGEdge 3 $\to$ 4	$\wedge \quad \texttt{Var8} \equiv 1$
ICFGNode 4	$\wedge \quad \texttt{Var11} \equiv \texttt{Var7} + \texttt{Var12}$
ICFGNode 7	$\wedge \quad \texttt{Var16} \equiv \texttt{Var11}$
ICFGNode 8	$\wedge \quad \texttt{Var17} \equiv \texttt{Var7} + \texttt{Var12}$

```
-------------SVFVar and Value-------------
                    ...
ValVar9                 Value: 10
ValVar12                Value: 1
ValVar7                 Value: 11
ValVar8                 Value: 1
ValVar11                Value: 12
ValVar16                Value: 12
ValVar17                Value: 12
                    ...
------------------------------------------
```

Example 4: Branches



ICFG

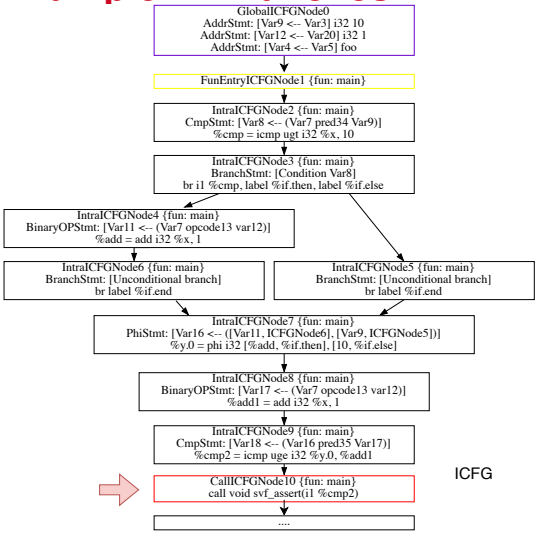
Verifying ICFG path: 0 → 1 → 2 → 3 → 4 →

6 → 7 → 8 → 9 → *svf_assert* (if.then branch)

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var9} \equiv 10 \wedge \texttt{Var12} \equiv 1 \wedge \texttt{Var4} \equiv \texttt{0x7f000005}$
ICFGNode 2	$\wedge \quad \texttt{Var8} \equiv \texttt{ite}(\texttt{Var7} > \texttt{Var9}, 1, 0)$
ICFGEdge 3 → 4	$\wedge \quad \texttt{Var8} \equiv 1$
ICFGNode 4	$\wedge \quad \texttt{Var11} \equiv \texttt{Var7} + \texttt{Var12}$
ICFGNode 7	$\wedge \quad \texttt{Var16} \equiv \texttt{Var11}$
ICFGNode 8	$\wedge \quad \texttt{Var17} \equiv \texttt{Var7} + \texttt{Var12}$
ICFGNode 9	$\wedge \quad \texttt{Var18} \equiv \texttt{ite}(\texttt{Var16} \geq \texttt{Var17}, 1, 0)$

```
------------SVFVar and Value------------
                    ...
ValVar9                 Value: 10
ValVar12                Value: 1
ValVar7                 Value: 11
ValVar8                 Value: 1
ValVar11                Value: 12
ValVar16                Value: 12
ValVar17                Value: 12
ValVar18                Value: 1
                    ...
```

Example 4: Branches



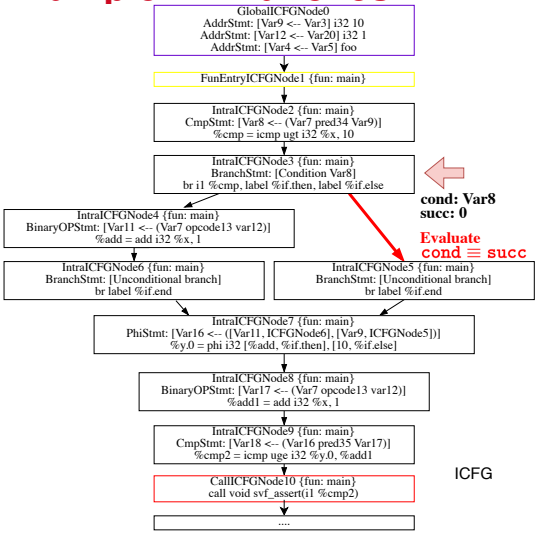
ICFG

Verifying ICFG path: $0 \to 1 \to 2 \to 3 \to 4 \to 6 \to 7 \to 8 \to 9 \to$ *svf_assert* (if.then branch)

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var9} \equiv 10 \wedge \texttt{Var12} \equiv 1 \wedge \texttt{Var4} \equiv \texttt{0x7f000005}$
ICFGNode 2	$\wedge \quad \texttt{Var8} \equiv \texttt{ite}(\texttt{Var7} > \texttt{Var9}, 1, 0)$
ICFGEdge 3 $\to$ 4	$\wedge \quad \texttt{Var8} \equiv 1$
ICFGNode 4	$\wedge \quad \texttt{Var11} \equiv \texttt{Var7} + \texttt{Var12}$
ICFGNode 7	$\wedge \quad \texttt{Var16} \equiv \texttt{Var11}$
ICFGNode 8	$\wedge \quad \texttt{Var17} \equiv \texttt{Var7} + \texttt{Var12}$
ICFGNode 9	$\wedge \quad \texttt{Var18} \equiv \texttt{ite}(\texttt{Var16} \geq \texttt{Var17}, 1, 0)$
ICFGNode 10	$\wedge \quad \texttt{Var18} \equiv 0$ (negation of the assert condition)

Solver yields **UNSAT** (i.e., no counter example), therefore, the assertion is successfully verified!!

Example 4: Branches



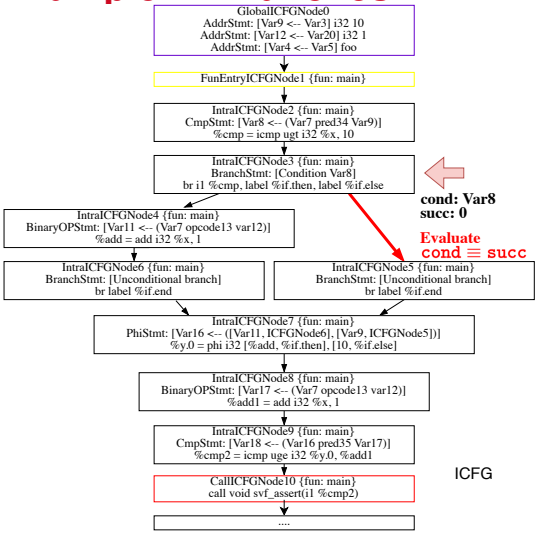
Algorithm 13: 3 `handleIntra(intraEdge)`

```
2  if intraEdge.getCondition() &&
     !handleBranch(intraEdge) then
4      return false;
6  else
8      handleNonBranch(edge);
```

Algorithm 14: `handleBranch(intraEdge)`

```
1  cond = ... // Retrieve branch condition boolean var;
2  succ = ... // Edge's succ condition value (0 or 1);
3  // Evaluate path feasibility;
4  // The current path/edge is feasible if 'cond' and
     'succ' have the same value, when evaluated in the
     solver; otherwise this branch edge is infeasible;
5  if the brach edge is infeasible then
6      return false;
7  else
8      // add (cond≡succ) to solver if it holds;
9      return true;
```

Example 4: Branches



Verifying ICFG path: $0 \to 1 \to 2 \to 3 \to 5 \to 7 \to 8 \to 9 \to$ *svf_assert* (if.else branch)

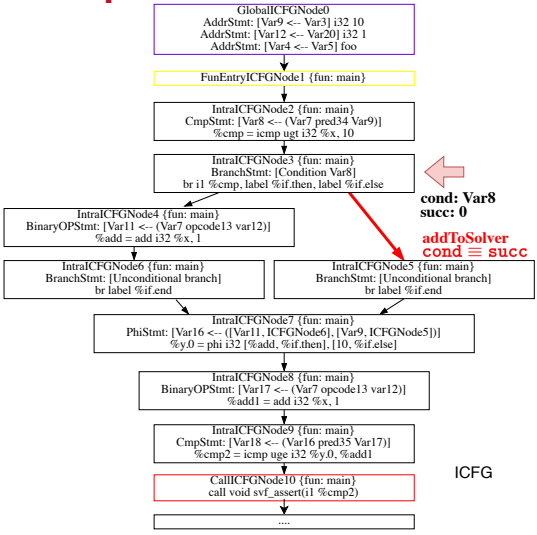
ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var9} \equiv 10 \wedge \texttt{Var12} \equiv 1 \wedge \texttt{Var4} \equiv \texttt{0x7f000005}$
ICFGNode 2	$\wedge \quad \texttt{Var8} \equiv \texttt{ite}(\texttt{Var7} > \texttt{Var9}, 1, 0)$
ICFGEdge 3 → 5	

The constraint $\texttt{Var8} \equiv 0$ is evaluated to be SAT

The conditional ICFGEdge [ICFGNode3 → ICFGNode5] is feasible.

```
------------SVFVar and Value------------
ObjVar5 (0x7f000005)     Value: NULL
ValVar4                  Value: 0x7f000005
ValVar9                  Value: 10
ValVar12                 Value: 1
ValVar7                  Value: 10
ValVar8                  Value: 0
                  ...
----------------------------------------
```

Example 4: Branches



ICFG

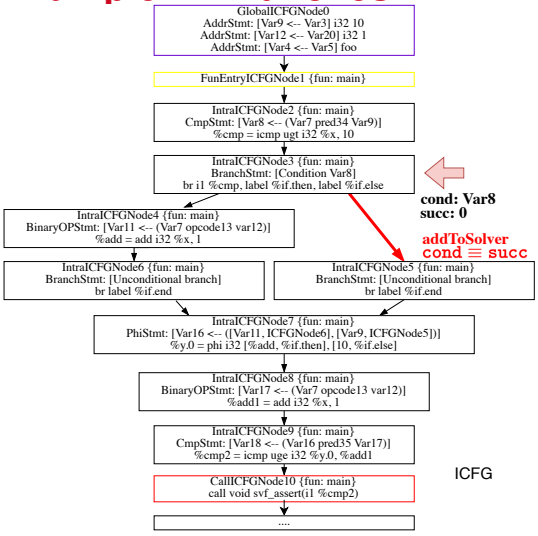
Algorithm 15: 3 `handleIntra(intraEdge)`

```
2  if intraEdge.getCondition() &&
     !handleBranch(intraEdge) then
4      return false;
6  else
8      handleNonBranch(edge);
```

Algorithm 16: `handleBranch(intraEdge)`

```
1  cond = ... // Retrieve branch condition boolean var;
2  succ = ... // Edge's succ condition value (0 or 1);
3  // Evaluate path feasibility;
4  // The current path/edge is feasible if 'cond' and
     'succ' have the same value, when evaluated in the
     solver; otherwise this branch edge is infeasible;
5  if the brach edge is infeasible then
6      return false;
7  else
8      // add (cond≡succ) to solver if it holds;
9      return true;
```

Example 4: Branches

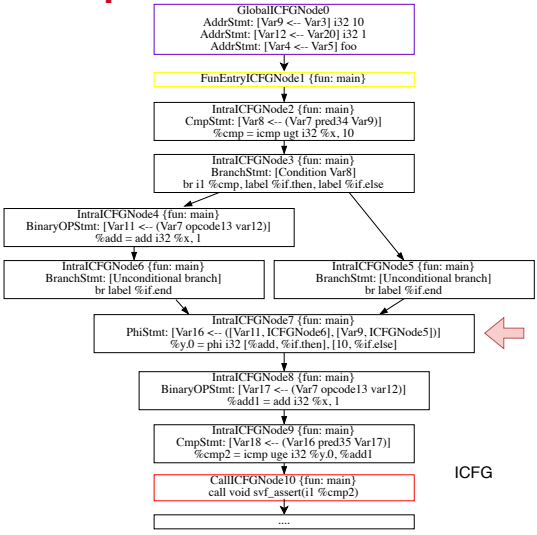


Verifying ICFG path: $0 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow 5 \rightarrow 7 \rightarrow 8 \rightarrow 9 \rightarrow$ *svf_assert* (if.else branch)

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var9} \equiv 10 \wedge \texttt{Var12} \equiv 1 \wedge \texttt{Var4} \equiv \texttt{0x7f000005}$
ICFGNode 2	$\wedge \quad \texttt{Var8} \equiv \texttt{ite}(\texttt{Var7} > \texttt{Var9}, 1, 0)$
ICFGEdge 3 → 5	$\wedge \quad \texttt{Var8} \equiv 0$

```
-------------SVFVar and Value-------------
ObjVar5 (0x7f000005)     Value: NULL
ValVar4                  Value: 0x7f000005
ValVar9                  Value: 10
ValVar12                 Value: 1
ValVar7                  Value: 10
ValVar8                  Value: 0
                  ...
------------------------------------------
```

Example 4: Branches



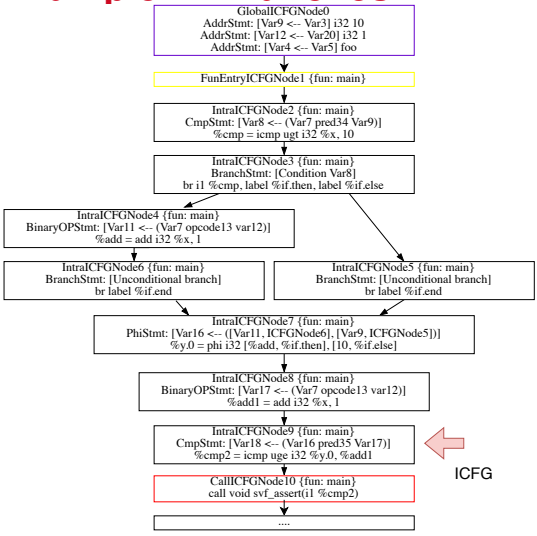
ICFG

Verifying ICFG path: $0 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow 5 \rightarrow 7 \rightarrow 8 \rightarrow 9 \rightarrow$ *svf_assert* (if.else branch)

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var9} \equiv 10 \wedge \texttt{Var12} \equiv 1 \wedge \texttt{Var4} \equiv \texttt{0x7f000005}$
ICFGNode 2	$\wedge \quad \texttt{Var8} \equiv \texttt{ite}(\texttt{Var7} > \texttt{Var9}, 1, 0)$
ICFGEdge 3 $\rightarrow$ 5	$\wedge \quad \texttt{Var8} \equiv 0$
ICFGNode 7	$\wedge \quad \texttt{Var16} \equiv \texttt{Var9}$

```
-------------SVFVar and Value-------------
ObjVar5 (0x7f000005)    Value: NULL
ValVar4                 Value: 0x7f000005
ValVar9                 Value: 10
ValVar12                Value: 1
ValVar7                 Value: 10
ValVar8                 Value: 0
ValVar16                Value: 10
                  ...
------------------------------------------
```

Example 4: Branches

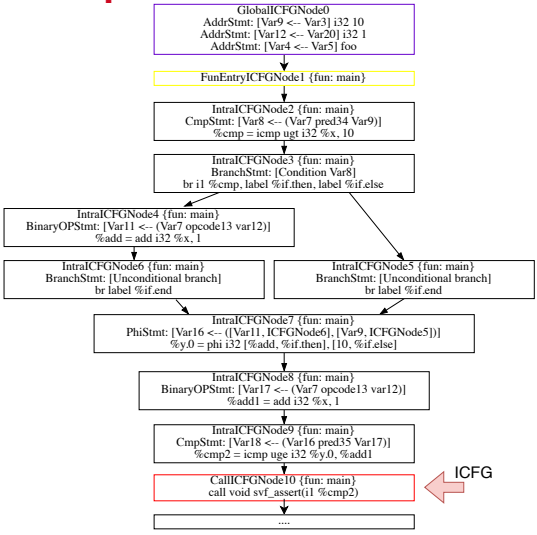


Verifying ICFG path: 0 → 1 → 2 → 3 → 5 → 7 → 8 → 9 → *svf_assert* (if.else branch)

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var9} \equiv 10 \wedge \texttt{Var12} \equiv 1 \wedge \texttt{Var4} \equiv \texttt{0x7f000005}$
ICFGNode 2	$\wedge \quad \texttt{Var8} \equiv \texttt{ite}(\texttt{Var7} > \texttt{Var9}, 1, 0)$
ICFGEdge 3 → 5	$\wedge \quad \texttt{Var8} \equiv 0$
ICFGNode 7	$\wedge \quad \texttt{Var16} \equiv \texttt{Var9}$
ICFGNode 8	$\wedge \quad \texttt{Var17} \equiv \texttt{Var7} + \texttt{Var12}$
ICFGNode 9	$\wedge \quad \texttt{Var18} \equiv \texttt{ite}(\texttt{Var16} \geq \texttt{Var17}, 1, 0)$

```
------------SVFVar and Value------------
ObjVar5 (0x7f000005)    Value: NULL
ValVar4                 Value: 0x7f000005
ValVar9                 Value: 10
ValVar12                Value: 1
ValVar7                 Value: 10
ValVar8                 Value: 0
ValVar16                Value: 10
ValVar17                Value: 11
ValVar18                Value: 0
                  ...
----------------------------------------
```

Example 4: Branches



Verifying ICFG path: $0 \to 1 \to 2 \to 3 \to 5 \to 7 \to 8 \to 9 \to$ *svf_assert* (if.else branch)

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var9} \equiv 10 \wedge \texttt{Var12} \equiv 1 \wedge \texttt{Var4} \equiv \texttt{0x7f000005}$
ICFGNode 2	$\wedge \quad \texttt{Var8} \equiv \texttt{ite}(\texttt{Var7} > \texttt{Var9}, 1, 0)$
ICFGEdge 3 $\to$ 5	$\wedge \quad \texttt{Var8} \equiv 0$
ICFGNode 7	$\wedge \quad \texttt{Var16} \equiv \texttt{Var9}$
ICFGNode 8	$\wedge \quad \texttt{Var17} \equiv \texttt{Var7} + \texttt{Var12}$
ICFGNode 9	$\wedge \quad \texttt{Var18} \equiv \texttt{ite}(\texttt{Var16} \geq \texttt{Var17}, 1, 0)$
ICFGNode 10	$\wedge \quad \texttt{Var18} \equiv 0$ (negation of the assert condition)

Solver yields **SAT**, a counterexample exists:

($\texttt{Var16} \equiv 10 \wedge \texttt{Var17} \equiv 11$).

The assertion is violated and fails!

Example 5: Interprocedural

```
1  int foo(int p) {
2      return p;
3  }
4  int main(int argc) {
5    int x;
6    int y;
7    x = foo(3); //ctx:[ℓ7]
8    y = foo(argc); //ctx:[ℓ8]
9    svf_assert(y == argc);
10 }
```

[ℓ_7]: calling context of `foo` at ℓ_7

[ℓ_8]: calling context of `foo` at ℓ_8

Example 5: Interprocedural

```
1  int foo(int p) {
2      return p;
3  }
4  int main(int argc) {
5    int x;
6    int y;
7    x = foo(3); //ctx:[ℓ7]
8    y = foo(argc); //ctx:[ℓ8]
9    svf_assert(y == argc);
10 }
```

[ℓ_7]: calling context of `foo` at ℓ_7

[ℓ_8]: calling context of `foo` at ℓ_8

Concrete Execution
(Concrete states)

One execution:

`argc` : 0

push calling context (calling foo at ℓ_7)

`p` : 3

calling context pop (returning from foo at ℓ_2)

`x` : 3

push calling context (calling foo at ℓ_8)

`p` : 0

pop calling context (returning from foo ℓ_2)

`y` : 0

Example 5: Interprocedural

```
1  int foo(int p) {
2      return p;
3  }
4  int main(int argc) {
5    int x;
6    int y;
7    x = foo(3); //ctx:[ℓ7]
8    y = foo(argc); //ctx:[ℓ8]
9    svf_assert(y == argc);
10 }
```

$[\ell_7]$: calling context of `foo` at ℓ_7

$[\ell_8]$: calling context of `foo` at ℓ_8

Concrete Execution
(Concrete states)

One execution:

`argc` : 0

push calling context (calling foo at ℓ_7)

`p` : 3

calling context pop (returning from foo at ℓ_2)

`x` : 3

push calling context (calling foo at ℓ_8)

`p` : 0

pop calling context (returning from foo ℓ_2)

`y` : 0

Symbolic Execution
(Symbolic states)

`argc` : `getZ3Expr(argc)`

push abstract calling context (current ctx:$[\ell_7]$)

$\langle[\ell_7], \mathrm{p}\rangle$: 3

`x` : `getZ3Expr(`$\langle[\ell_7], \mathrm{p}\rangle$`)`

pop abstract calling context (current ctx:[ ])

push abstract calling context (current ctx:$[\ell_8]$)

$\langle[\ell_8], \mathrm{p}\rangle$: `getZ3Expr(argc)`

`y` : `getZ3Expr(`$\langle[\ell_8], \mathrm{p}\rangle$`)`

pop abstract calling context (current ctx:[ ])

Example 5: Interprocedural

```
1  int foo(int p) {
2      return p;
3  }
4  int main(int argc) {
5    int x;
6    int y;
7    x = foo(3); //ctx:[ℓ7]
8    y = foo(argc); //ctx:[ℓ8]
9    svf_assert(y == argc);
10 }
```

$[\ell_7]$: calling context of `foo` at ℓ_7

$[\ell_8]$: calling context of `foo` at ℓ_8

Concrete Execution
(Concrete states)

One execution:

`argc` : 0

push calling context (calling foo at ℓ_7)

`p` : 3

calling context pop (returning from foo at ℓ_2)

`x` : 3

push calling context (calling foo at ℓ_8)

`p` : 0

pop calling context (returning from foo ℓ_2)

`y` : 0

Symbolic Execution
(Symbolic states)

`argc` : `getZ3Expr(argc)`

push abstract calling context (current ctx:$[\ell_7]$)

$\langle[\ell_7], \mathrm{p}\rangle$: 3

`x` : `getZ3Expr(`$\langle[\ell_7], \mathrm{p}\rangle$`)`

pop abstract calling context (current ctx:[ ])

push abstract calling context (current ctx:$[\ell_8]$)

$\langle[\ell_8], \mathrm{p}\rangle$: `getZ3Expr(argc)`

`y` : `getZ3Expr(`$\langle[\ell_8], \mathrm{p}\rangle$`)`

pop abstract calling context (current ctx:[ ])

Checking non-existence of counterexamples:

$\psi(N_1) \wedge \psi(N_2) \wedge \ldots \psi(N_i) \wedge \neg\psi(Q)$	Satisfiability	Counterexample
$\langle[\ell_7], \mathrm{p}\rangle \equiv 3 \wedge \mathrm{x} \equiv \langle[\ell_7], \mathrm{p}\rangle \wedge \langle[\ell_8], \mathrm{p}\rangle \equiv \mathrm{argc} \wedge \mathrm{y} \equiv \langle[\ell_7], \mathrm{p}\rangle \wedge \mathrm{y} \neq \mathrm{argc}$	unsat	$\varnothing$

`foo`'s argument `p` needs to be **differentiated and renamed** as $\langle[\ell_7], \mathrm{p}\rangle$ and $\langle[\ell_8], \mathrm{p}\rangle$ due to two calling contexts, $[\ell_7]$ and $[\ell_8]$ to mimic the runtime call stack which holds the local variable `p`.

Example 5: Interprocedural

`SSE::getZ3Expr`(SVFVarID) in `Assignment-2`

- Get an Z3 expression based on SVFVarID and the current calling context `callingCtx`
- `callingCtx` is maintained per ICFG path by calling `SSE::pushCallingCtx` and `SSE::popCallingCtx` when handling `CallCFGEdge` and `RetCFGEdge`.

```
1 z3::expr SSE::getZ3Expr(NodeID idx) const {
2     return z3Mgr->getZ3Expr(idx, callingCtx);
3 }
```

Example 5: Interprocedural

```
1  int foo(int p) {
2      return p;
3  }
4  int main(int argc) {
5    int x;
6    int y;
7    x = foo(3);
8    y = foo(argc);
9    svf_assert(y == argc);
10 }
```

Source code

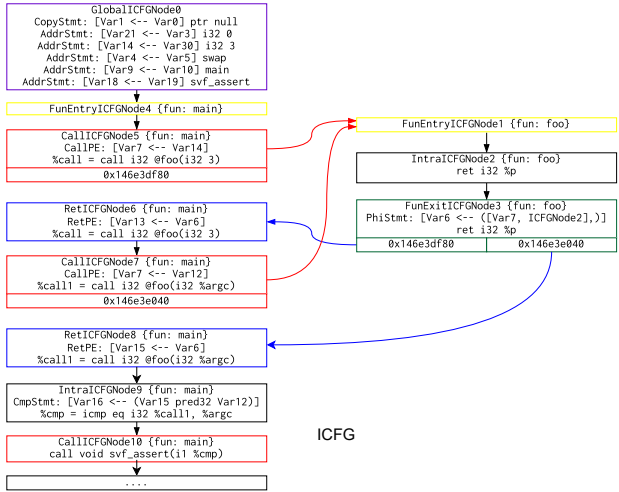
```
1  define i32 @foo(i32 %p) #0 {
2  entry:
3    ret i32 %p
4  }
5
6  define i32 @main(i32 %argc) #0 {
7  entry:
8    %call = call i32 @foo(i32 3)
9    %call1 = call i32 @foo(i32 %argc)
10   %cmp = icmp eq i32 %call1, %argc
11   call void @svf_assert(i1 zeroext %cmp)
12   ret i32 0
13 }
```

LLVM IR

Example 5: Interprocedural

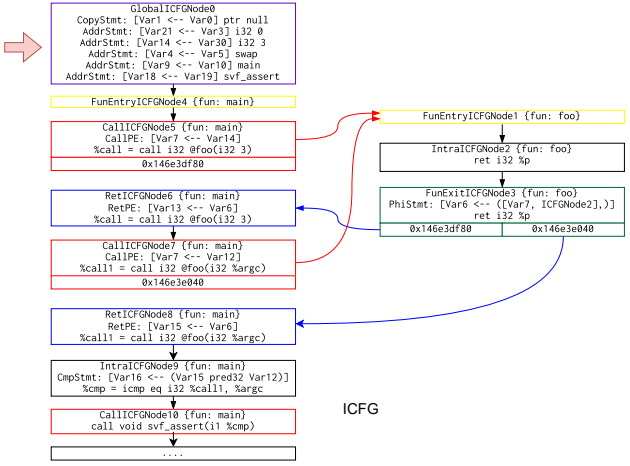
```
1  define i32 @foo(i32 %p) #0 {
2  entry:
3    ret i32 %p
4  }
5
6  define i32 @main(i32 %argc) #0 {
7  entry:
8    %call = call i32 @foo(i32 3)
9    %call1 = call i32 @foo(i32 %argc)
10   %cmp = icmp eq i32 %call1, %argc
11   call void @svf_assert(i1 zeroext %cmp)
12   ret i32 0
13 }
```

LLVM IR



ICFG

Example 5: Interprocedural



ICFG

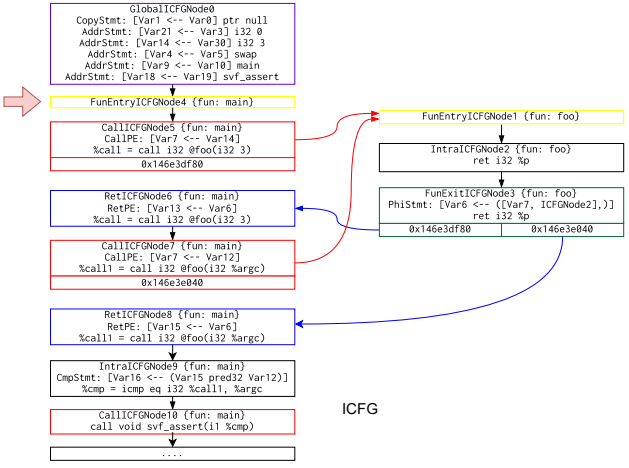
Verifying ICFG path: 0 → 4 → 5 → 1 → 2 → 3 → 6 → 7 → 1 → 2 → 3 → 8 → 9 → *svf_assert*

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	Var21 ≡ 0 ∧ Var14 ≡ 3 ∧ ...

```
-------------SVFVar and Value-------------
ObjVar5 (0x7f000005)     Value: NULL
ObjVar10 (0x7f00000a)    Value: NULL
ObjVar19 (0x7f000013)    Value: NULL
ValVar0                  Value: NULL
ValVar1                  Value: NULL
ValVar21                 Value: 0
ValVar14                 Value: 3
ValVar4                  Value: 0x7f000005
ValVar9                  Value: 0x7f00000e
ValVar18                 Value: 0x7f00001d
                  ...
------------------------------------------
```

The values of Z3 expressions for each SVFVar after analyzing `GlobalICFGNode0` (use `printExprValues()` to print SVFVars and their values)

Example 5: Interprocedural

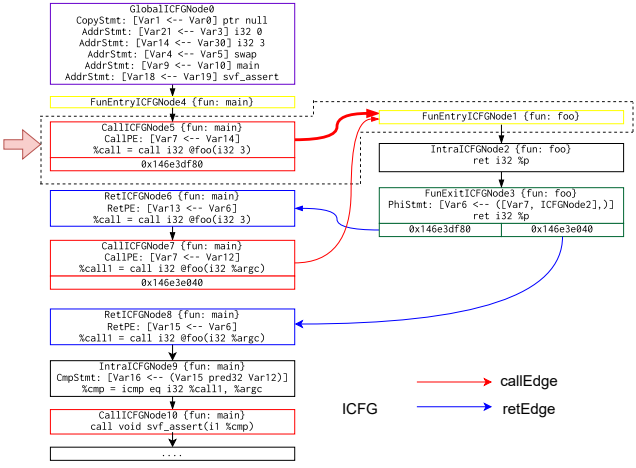


ICFG

Algorithm 17: 2 `translatePath(path)`

```
2   foreach edge ∈ path do
4       if IntraEdge ← dyn_cast⟨IntraCFGEdge⟩(edge)
          then
6           if handleIntra(IntraEdge) == false then
8               return false;
10      else if CallEdge ← dyn_cast⟨CallCFGEdge⟩(edge)
          then
12          handleCall(CallEdge);
14      else if RetEdge ← dyn_cast⟨RetCFGEdge⟩(edge)
          then
16          handleRet(RetEdge);
18      else
20          assert(false &&"what other edges we have?");
21  return true;
```

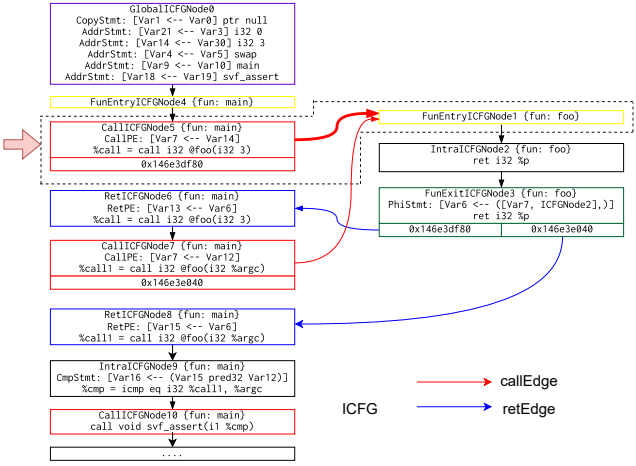
Example 5: Interprocedural



Algorithm 18: `handleCall(callEdge)`

```
1 // (1) For each CallPE lhs = rhs:  retrieve rhs
  expr ⟨[c], rhs⟩ under callingCtx c (caller's
  context)
2 // (2) pushCallingCtx using current calliste
   ICFGNode; Note that pushCallingCtx will make c
   become c'.  ;
3 // (3) For each CallPE lhs = rhs:  retrieve lhs
  expr ⟨[c'], lhs⟩ under callingCtx c' (callee's
  context) ;
4 // (4) For each CallPE lhs = rhs:  add
  ⟨[c], rhs⟩ ≡ ⟨[c'], lhs⟩ into the solver.
5 return true;
```

Example 5: Interprocedural



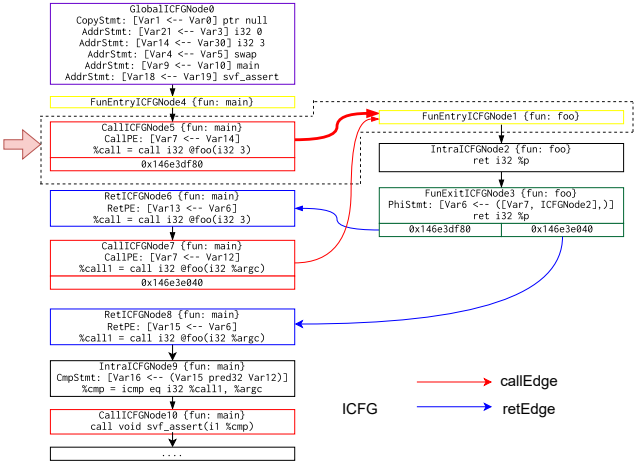
Verifying ICFG path: 0 → 4 → 5 →1 → 2 → 3 →
6 → 7 → 1 → 2 → 3 → 8 → 9 → *svf_assert*

Calling Context

After processing the call edge

`CallICFGNode5` → `FunEntryICFGNode1`

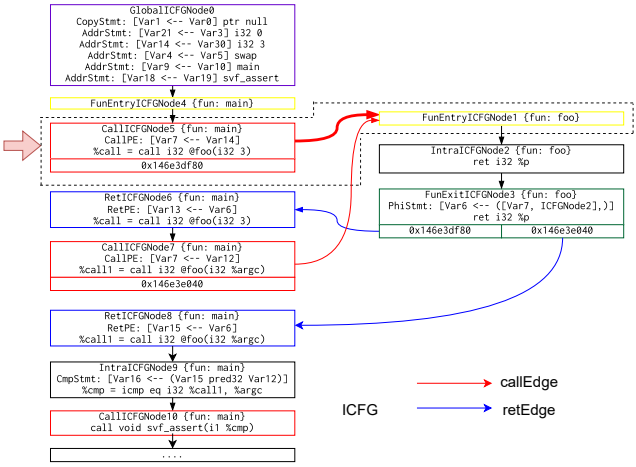
Example 5: Interprocedural



Algorithm 19: `handleCall(callEdge)`

```
1 // (1) For each CallPE lhs = rhs:  retrieve rhs
    expr ⟨[c], rhs⟩ under callingCtx c (caller's
    context)
2 // (2) pushCallingCtx using current calliste
    ICFGNode; Note that pushCallingCtx will make c
    become c'.  ;
3 // (3) For each CallPE lhs = rhs:  retrieve lhs
    expr ⟨[c'], lhs⟩ under callingCtx c' (callee's
    context) ;
4 // (4) For each CallPE lhs = rhs:  add
    ⟨[c], rhs⟩ ≡ ⟨[c'], lhs⟩ into the solver.
5 return true;
```

Example 5: Interprocedural



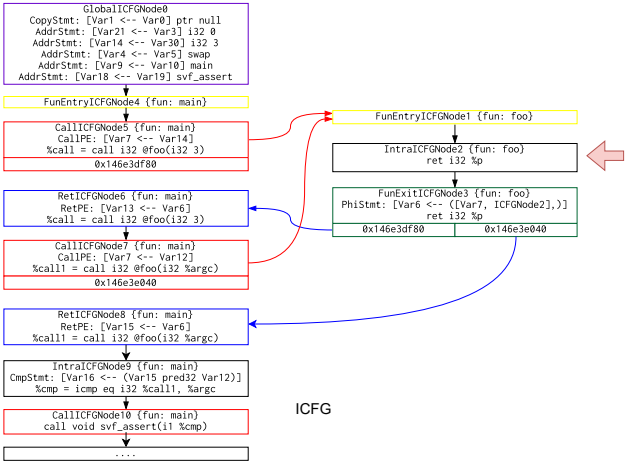
Verifying ICFG path: $0 \to 4 \to 5 \to 1 \to 2 \to 3 \to 6 \to 7 \to 1 \to 2 \to 3 \to 8 \to 9 \to$ *svf_assert*

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var21} \equiv 0 \wedge \texttt{Var14} \equiv 3 \wedge \ldots$
ICFGNode 5	$\wedge \quad \langle [\texttt{ICFGNode5}], \texttt{Var7} \rangle \equiv \langle [\,], \texttt{Var14} \rangle$

```
------------SVFVar and Value-------------
ObjVar5 (0x7f000005)    Value: NULL
ObjVar10 (0x7f00000a)   Value: NULL
ObjVar19 (0x7f000013)   Value: NULL
ValVar0                 Value: NULL
ValVar1                 Value: NULL
ValVar21                Value: 0
ValVar14                Value: 3
ValVar4                 Value: 0x7f000005
ValVar9                 Value: 0x7f00000e
ValVar18                Value: 0x7f00001d
ValVar7                 Value: 3
                ...
----------------------------------------
```

Note: SVFVar and Value table is printed under the calling context [`CallICFGNode5`]

Example 5: Interprocedural

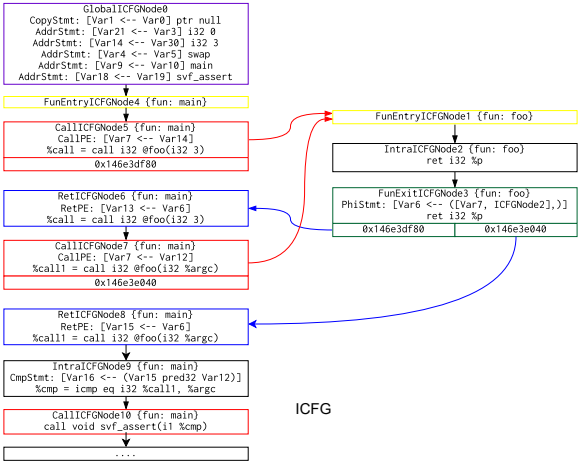


ICFG

Verifying ICFG path: $0 \rightarrow 4 \rightarrow 5 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow 6 \rightarrow 7 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow 8 \rightarrow 9 \rightarrow$ *svf_assert*

`ret i32 %p` instruction.
Nothing needs to be done.

Example 5: Interprocedural

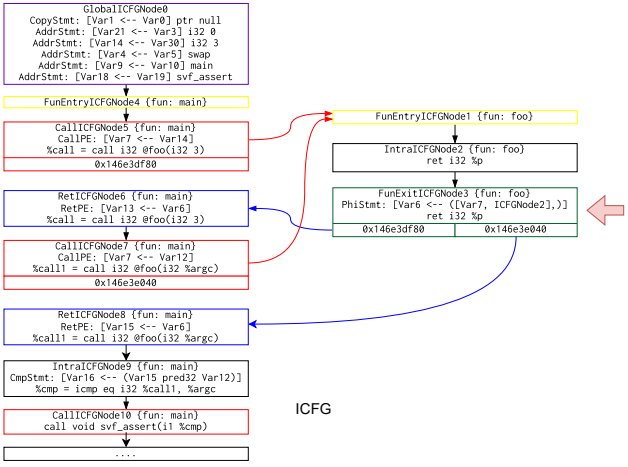


ICFG

Algorithm 20: 3 `handlePhi(edge)`

```
1 res ← getZ3Expr(phi.getResID());
2 opINodeFound ← false;
3 for i ← 0 to phi.getOpVarNum() − 1 do
4     if edge.srcNode() postdominates
        phi.getOpICFGNode(i) then
5         ope ← getZ3Expr(phi.getOpVar(i).getId());
6         addToSolver(res == ope);
7         opINodeFound ← true;
```

Example 5: Interprocedural



ICFG

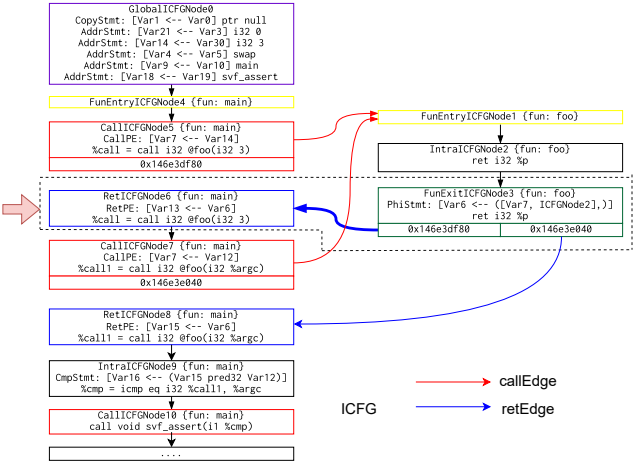
Verifying ICFG path: $0 \to 4 \to 5 \to 1 \to 2 \to 3 \to 6 \to 7 \to 1 \to 2 \to 3 \to 8 \to 9 \to$ *svf_assert*

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var21} \equiv 0 \wedge \texttt{Var14} \equiv 3 \wedge \ldots$
ICFGNode 5	$\wedge \quad \langle [\texttt{ICFGNode5}], \texttt{Var7} \rangle \equiv \langle [\,], \texttt{Var14} \rangle$
ICFGNode 3	$\wedge \quad \langle [\texttt{ICFGNode5}], \texttt{Var6} \rangle \equiv \langle [\texttt{ICFGNode5}], \texttt{Var7} \rangle$

```
-------------SVFVar and Value-------------
ObjVar5 (0x7f000005)     Value: NULL
ObjVar10 (0x7f00000a)    Value: NULL
ObjVar19 (0x7f000013)    Value: NULL
ValVar0                  Value: NULL
ValVar1                  Value: NULL
ValVar21                 Value: 0
ValVar14                 Value: 3
ValVar4                  Value: 0x7f000005
ValVar9                  Value: 0x7f00000e
ValVar18                 Value: 0x7f00001d
ValVar7                  Value: 3
ValVar6                  Value: 3
                  ...
------------------------------------------
```

Note: SVFVar and Value table is printed under the calling context [`CallICFGNode5`]

Example 5: Interprocedural

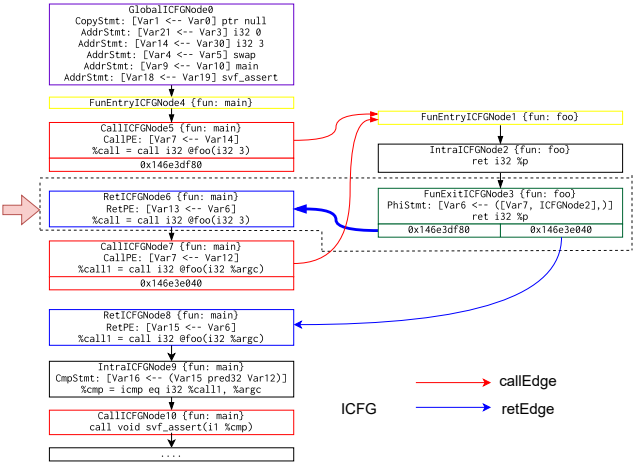


Algorithm 21: `handleRet(retEdge)`

```
1 // (1) For each RetPE lhs = rhs:  retrieve rhs
    expr ⟨[c′], rhs⟩ under callingCtx c' (callee's
    context);
2 // (2) popCallingCtx(); This will make c' become
  c (caller's context)
3 // (3) For each RetPE lhs = rhs:  retrieve lhs
  expr ⟨[c], lhs⟩ under callingCtx c (caller's
  context)
4 // (4) Add ⟨[c′], rhs⟩ ≡ ⟨[c], lhs⟩ into the solver.
5   return true;
```

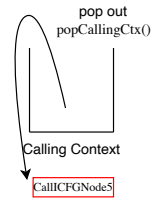
Note: `retPE.getRHSVarID()` returns `ValVar6`
`getZ3Expr(ValVar6)` binds `ValVar6` with the current `callingCtx` `[ICFGNode5]` and returns the Z3 expression for `[ICFGNode5], Var6`

Example 5: Interprocedural



Algorithm 22: `handleRet(retEdge)`

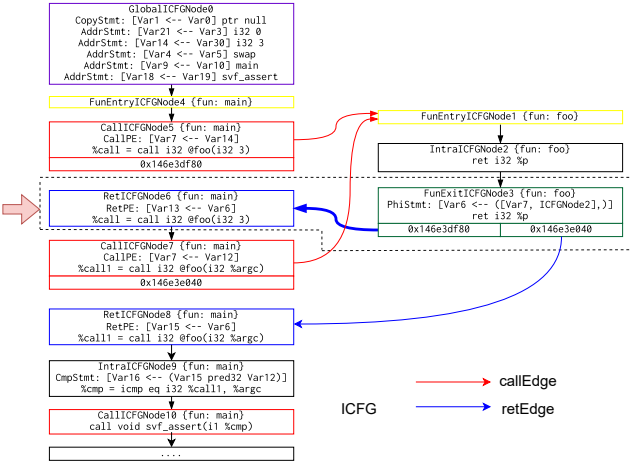
```
1 // (1) For each RetPE lhs = rhs:  retrieve rhs
  expr ⟨[c′], rhs⟩ under callingCtx c' (callee's
  context);
2 // (2) popCallingCtx(); This will make c' become
  c (caller's context)
3 // (3) For each RetPE lhs = rhs:  retrieve lhs
   expr ⟨[c], lhs⟩ under callingCtx c (caller's
   context)
4 // (4) Add ⟨[c′], rhs⟩ ≡ ⟨[c], lhs⟩ into the solver.
5   return true;
```



After processing the return edge

$\texttt{FunExitICFGNode3} \rightarrow \texttt{RetICFGNode6}$

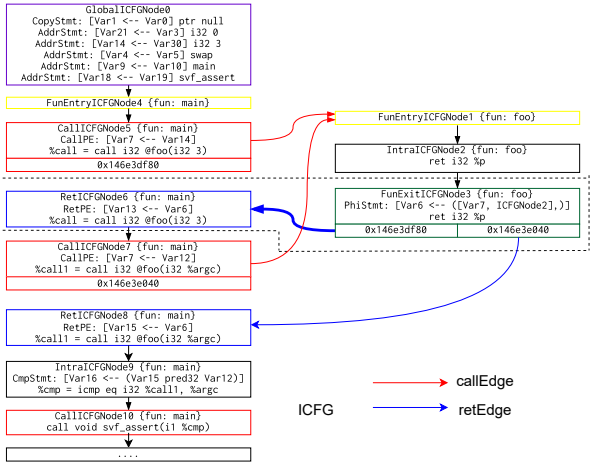
Example 5: Interprocedural



Algorithm 23: `handleRet(retEdge)`

```
1 // (1) For each RetPE lhs = rhs:  retrieve rhs
  expr ⟨[c′], rhs⟩ under callingCtx c' (callee's
  context);
2 // (2) popCallingCtx(); This will make c' become
  c (caller's context)
3 // (3) For each RetPE lhs = rhs:  retrieve lhs
   expr ⟨[c], lhs⟩ under callingCtx c (caller's
   context)
4 // (4) Add ⟨[c′], rhs⟩ ≡ ⟨[c], lhs⟩ into the solver.
5   return true;
```

Example 5: Interprocedural

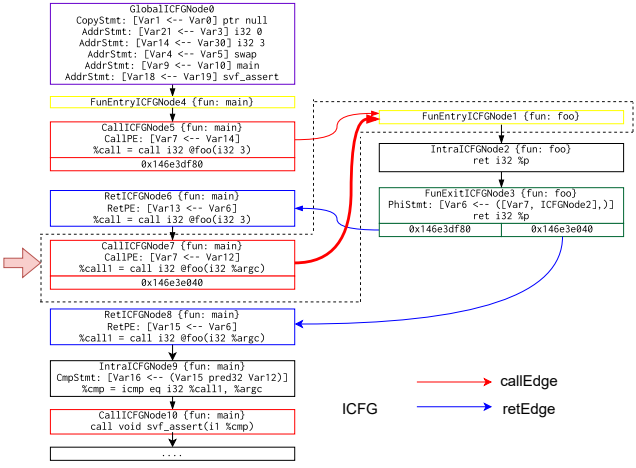


Verifying ICFG path: $0 \to 4 \to 5 \to 1 \to 2 \to 3 \to 6 \to 7 \to 1 \to 2 \to 3 \to 8 \to 9 \to$ *svf_assert*

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var21} \equiv 0 \wedge \texttt{Var14} \equiv 3 \wedge \ldots$
ICFGNode 5	$\wedge \quad \langle [\texttt{ICFGNode5}], \texttt{Var7} \rangle \equiv \langle [\,], \texttt{Var14} \rangle$
ICFGNode 3	$\wedge \quad \langle [\texttt{ICFGNode5}], \texttt{Var6} \rangle \equiv \langle [\texttt{ICFGNode5}], \texttt{Var7} \rangle$
ICFGNode 6	$\wedge \quad \langle [\,], \texttt{Var13} \rangle \equiv \langle [\texttt{ICFGNode5}], \texttt{Var6} \rangle$

```
------------SVFVar and Value------------
ObjVar5 (0x7f000005)    Value: NULL
ObjVar10 (0x7f00000a)   Value: NULL
ObjVar19 (0x7f000013)   Value: NULL
ValVar0                 Value: NULL
ValVar1                 Value: NULL
ValVar21                Value: 0
ValVar14                Value: 3
ValVar4                 Value: 0x7f000005
ValVar9                 Value: 0x7f00000e
ValVar18                Value: 0x7f00001d
ValVar13                Value: 3
                ...
----------------------------------------
```

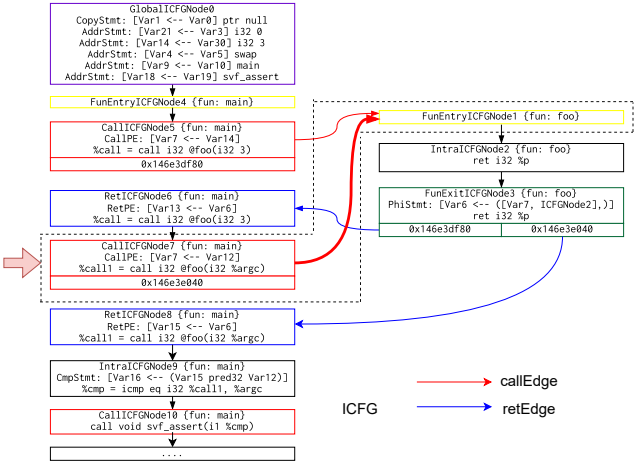
Example 5: Interprocedural



Algorithm 24: `handleCall(callEdge)`

```
1 // (1) For each CallPE lhs = rhs:  retrieve rhs
  expr ⟨[c], rhs⟩ under callingCtx c (caller's
  context)
2 // (2) pushCallingCtx using current calliste
   ICFGNode; Note that pushCallingCtx will make c
   become c'.  ;
3 // (3) For each CallPE lhs = rhs:  retrieve lhs
  expr ⟨[c'], lhs⟩ under callingCtx c' (callee's
  context) ;
4 // (4) For each CallPE lhs = rhs:  add
  ⟨[c], rhs⟩ ≡ ⟨[c'], lhs⟩ into the solver.
5 return true;
```

Example 5: Interprocedural



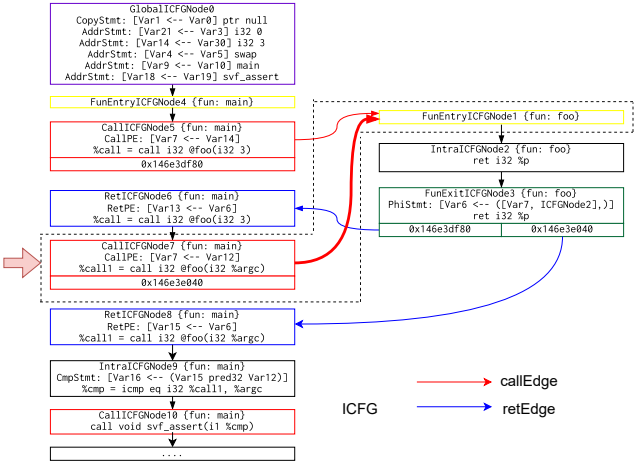
Verifying ICFG path: $0 \rightarrow 4 \rightarrow 5 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow 6 \rightarrow 7 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow 8 \rightarrow 9 \rightarrow$ *svf_assert*

Calling Context

After processing the call edge

`CallCFGNode7` → `FunEntryICFGNode`

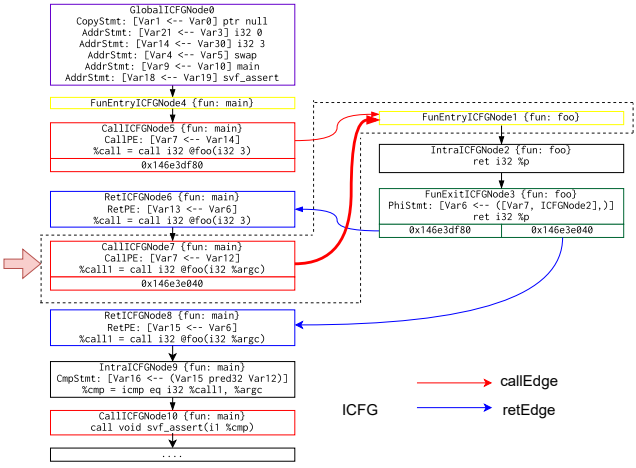
Example 5: Interprocedural



Algorithm 25: `handleCall(callEdge)`

```
1 // (1) For each CallPE lhs = rhs:  retrieve rhs
  expr ⟨[c], rhs⟩ under callingCtx c (caller's
  context)
2 // (2) pushCallingCtx using current calliste
  ICFGNode; Note that pushCallingCtx will make c
  become c'.  ;
3 // (3) For each CallPE lhs = rhs:  retrieve lhs
   expr ⟨[c'], lhs⟩ under callingCtx c' (callee's
   context) ;
4 // (4) For each CallPE lhs = rhs:  add
   ⟨[c], rhs⟩ ≡ ⟨[c'], lhs⟩ into the solver.
5 return true;
```

Example 5: Interprocedural



Verifying ICFG path: 0 → 4 → 5 → 1 → 2 → 3 → 6 → 7 →1 → 2 → 3 → 8 → 9 → *svf_assert*

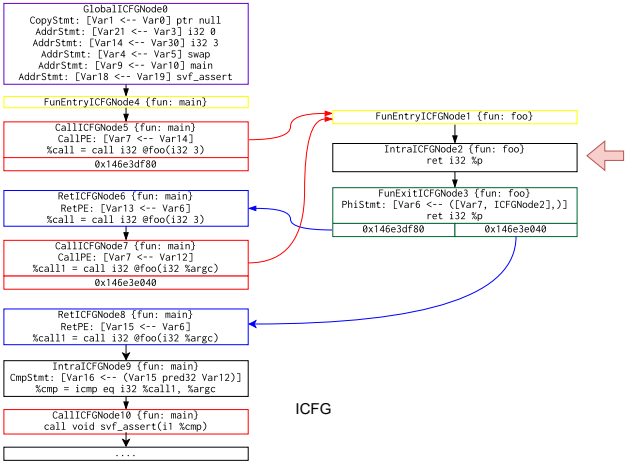
ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var21} \equiv 0 \wedge \texttt{Var14} \equiv 3 \wedge \ldots$
ICFGNode 5	$\wedge \quad \langle[\texttt{ICFGNode5}], \texttt{Var7}\rangle \equiv \langle[\,], \texttt{Var14}\rangle$
ICFGNode 3	$\wedge \quad \langle[\texttt{ICFGNode5}], \texttt{Var6}\rangle \equiv \langle[\texttt{ICFGNode5}], \texttt{Var7}\rangle$
ICFGNode 6	$\wedge \quad \langle[\,], \texttt{Var13}\rangle \equiv \langle[\texttt{ICFGNode5}], \texttt{Var6}\rangle$
ICFGNode 7	$\wedge \quad \langle[\texttt{ICFGNode7}], \texttt{Var7}\rangle \equiv \langle[\,], \texttt{Var12}\rangle$

```
------------SVFVar and Value-------------
                  ...
ValVar21              Value: 0
ValVar14              Value: 3
                  ...
ValVar13              Value 3
ValVar12              Value: 0
ValVar7               Value: 0
                  ...
-----------------------------------------
```

Note: SVFVars and their values in table are under the calling context [`CallICFGNode7`].

`ValVar12` is uninitialized, thus evaluated as 0.

Example 5: Interprocedural



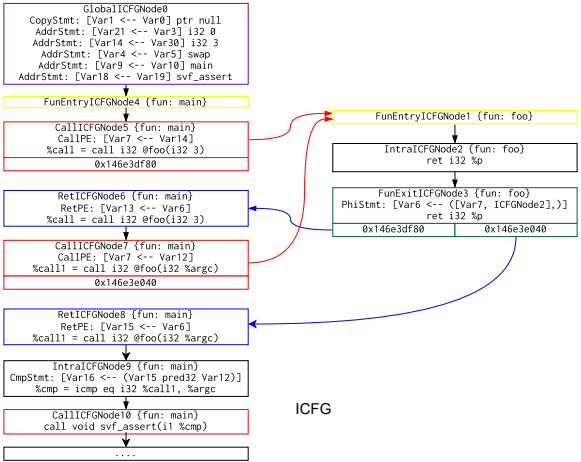
ICFG

Verifying ICFG path: $0 \rightarrow 4 \rightarrow 5 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow 6 \rightarrow 7 \rightarrow 1 \rightarrow 2 \rightarrow 3 \rightarrow 8 \rightarrow 9 \rightarrow$ *svf_assert*

`ret i32 %p` instruction.

Nothing needs to be done.

Example 5: Interprocedural

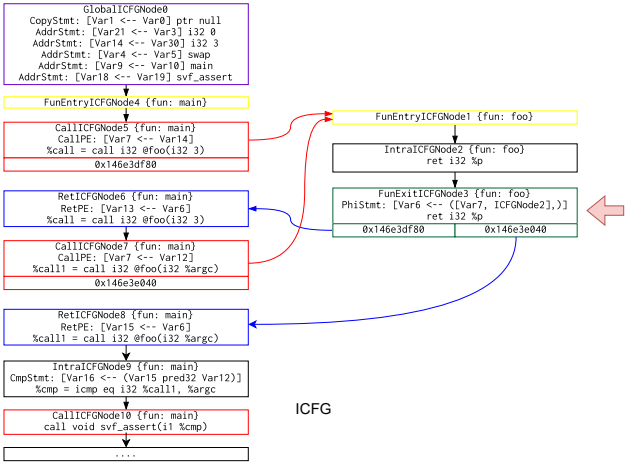


ICFG

Algorithm 26: 3 `handlePhi(edge)`

```
1 res ← getZ3Expr(phi.getResID());
2 opINodeFound ← false;
3 for i ← 0 to phi.getOpVarNum() − 1 do
4     if edge.srcNode() postdominates
        phi.getOpICFGNode(i) then
5         ope ← getZ3Expr(phi.getOpVar(i).getId());
6         addToSolver(res == ope);
7         opINodeFound ← true;
```

Example 5: Interprocedural



ICFG

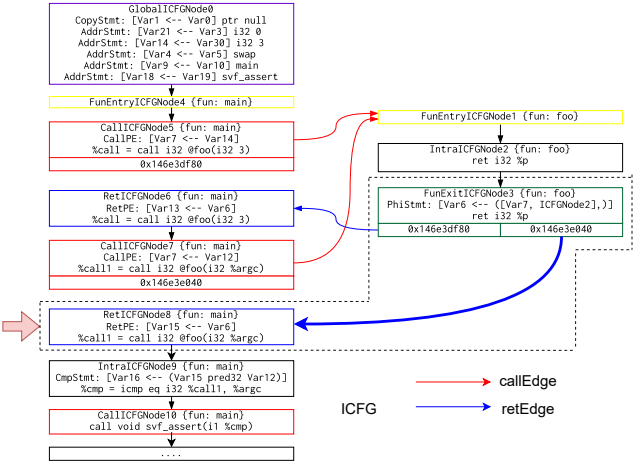
Verifying ICFG path: $0 \to 4 \to 5 \to 1 \to 2 \to 3 \to 6 \to 7 \to 1 \to 2 \to 3 \to 8 \to 9 \to$ *svf_assert*

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var21} \equiv 0 \wedge \texttt{Var14} \equiv 3 \wedge \ldots$
ICFGNode 5	$\wedge \ \langle [\texttt{ICFGNode5}], \texttt{Var7} \rangle \equiv \langle [\,], \texttt{Var14} \rangle$
ICFGNode 3	$\wedge \ \langle [\texttt{ICFGNode5}], \texttt{Var6} \rangle \equiv \langle [\texttt{ICFGNode5}], \texttt{Var7} \rangle$
ICFGNode 6	$\wedge \ \langle [\,], \texttt{Var13} \rangle \equiv \langle [\texttt{ICFGNode5}], \texttt{Var6} \rangle$
ICFGNode 7	$\wedge \ \langle [\texttt{ICFGNode7}], \texttt{Var7} \rangle \equiv \langle [\,], \texttt{Var12} \rangle$
ICFGNode 3	$\wedge \ \langle [\texttt{ICFGNode7}], \texttt{Var6} \rangle \equiv \langle [\texttt{ICFGNode7}], \texttt{Var7} \rangle$

```
-------------SVFVar and Value-------------
                    ...
ValVar21                Value: 0
ValVar14                Value: 3
                    ...
ValVar13                Value: 3
ValVar12                Value: 0
ValVar7                 Value: 0
ValVar6                 Value: 0
                    ...
------------------------------------------
```

Note: SVFVars and their values in table are under the calling context [CallICFGNode7].

Example 5: Interprocedural

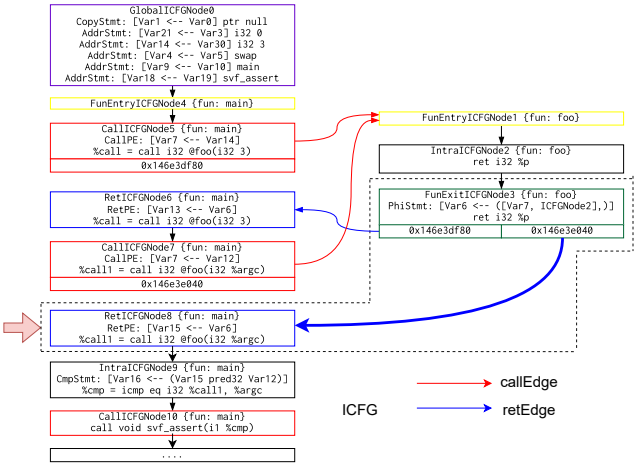


Algorithm 27: `handleRet(retEdge)`

```
1 // (1) For each RetPE lhs = rhs:  retrieve rhs
    expr ⟨[c′], rhs⟩ under callingCtx c' (callee's
    context);
2 // (2) popCallingCtx(); This will make c' become
  c (caller's context)
3 // (3) For each RetPE lhs = rhs:  retrieve lhs
  expr ⟨[c], lhs⟩ under callingCtx c (caller's
  context)
4 // (4) Add ⟨[c′], rhs⟩ ≡ ⟨[c], lhs⟩ into the solver.
5   return true;
```

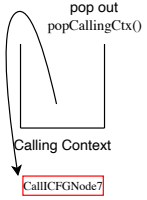
Note: `retPE.getRHSVarID()` returns `ValVar6`
`getZ3Expr(ValVar6)` binds `ValVar6` with the current `callingCtx` [ICFGNode7] and returns the Z3 expression for ⟨[ICFGNode7], Var6⟩

Example 5: Interprocedural



Algorithm 28: `handleRet(retEdge)`

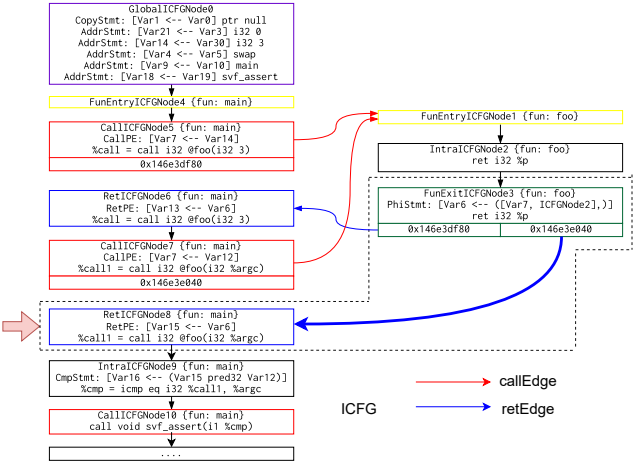
```
1 // (1) For each RetPE lhs = rhs:  retrieve rhs
  expr ⟨[c′], rhs⟩ under callingCtx c' (callee's
  context);
2 // (2) popCallingCtx(); This will make c' become
   c (caller's context)
3 // (3) For each RetPE lhs = rhs:  retrieve lhs
  expr ⟨[c], lhs⟩ under callingCtx c (caller's
  context)
4 // (4) Add ⟨[c′], rhs⟩ ≡ ⟨[c], lhs⟩ into the solver.
5   return true;
```



After processing the return edge

`FunExitICFGNode3` → `RetICFGNode8`

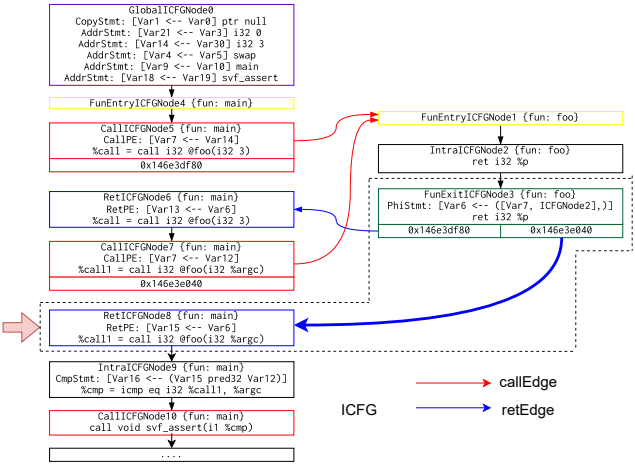
Example 5: Interprocedural



Algorithm 29: `handleRet(retEdge)`

```
1 // (1) For each RetPE lhs = rhs:  retrieve rhs
  expr ⟨[c′], rhs⟩ under callingCtx c' (callee's
  context);
2 // (2) popCallingCtx(); This will make c' become
  c (caller's context)
3 // (3) For each RetPE lhs = rhs:  retrieve lhs
   expr ⟨[c], lhs⟩ under callingCtx c (caller's
   context)
4 // (4) Add ⟨[c′], rhs⟩ ≡ ⟨[c], lhs⟩ into the solver.
5   return true;
```

Example 5: Interprocedural

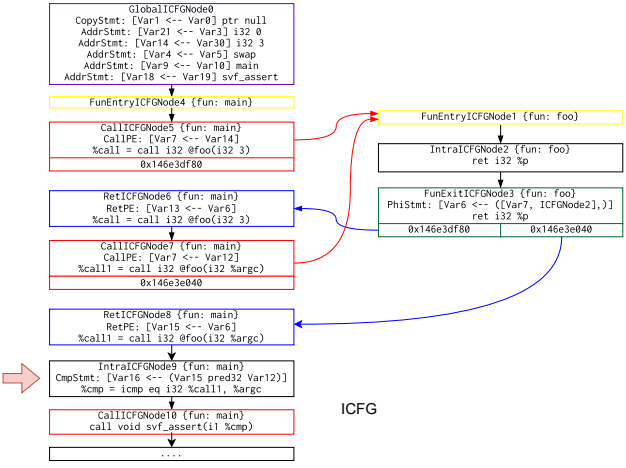


Verifying ICFG path: 0 → 4 → 5 → 1 → 2 → 3 → 6 → 7 → 1 → 2 → 3 →8 → 9 → *svf_assert*

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var21} \equiv 0 \wedge \texttt{Var14} \equiv 3 \wedge \ldots$
ICFGNode 5	$\wedge \ \langle[\texttt{ICFGNode5}], \texttt{Var7}\rangle \equiv \langle[\,], \texttt{Var14}\rangle$
ICFGNode 3	$\wedge \ \langle[\texttt{ICFGNode5}], \texttt{Var6}\rangle \equiv \langle[\texttt{ICFGNode5}], \texttt{Var7}\rangle$
ICFGNode 6	$\wedge \ \langle[\,], \texttt{Var13}\rangle \equiv \langle[\texttt{ICFGNode5}], \texttt{Var6}\rangle$
ICFGNode 7	$\wedge \ \langle[\texttt{ICFGNode7}], \texttt{Var7}\rangle \equiv \langle[\,], \texttt{Var12}\rangle$
ICFGNode 3	$\wedge \ \langle[\texttt{ICFGNode7}], \texttt{Var6}\rangle \equiv \langle[\texttt{ICFGNode7}], \texttt{Var7}\rangle$
ICFGNode 8	$\wedge \ \langle[\,], \texttt{Var15}\rangle \equiv \langle[\texttt{ICFGNode7}], \texttt{Var6}\rangle$

```
------------SVFVar and Value-------------
                    ...
ValVar21                Value: 0
ValVar14                Value: 3
                    ...
ValVar13                Value: 3
ValVar12                Value: 0
ValVar15                Value: 0
                    ...
-----------------------------------------
```

Example 5: Interprocedural



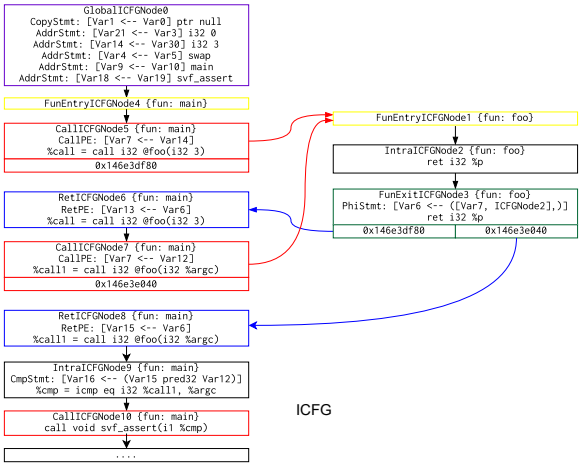
ICFG

Verifying ICFG path: $0 \to 4 \to 5 \to 1 \to 2 \to 3 \to 6 \to 7 \to 1 \to 2 \to 3 \to 8 \to 9 \to$ *svf_assert*

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var21} \equiv 0 \wedge \texttt{Var14} \equiv 3 \wedge \ldots$
ICFGNode 5	$\wedge \quad \langle [\texttt{ICFGNode5}], \texttt{Var7} \rangle \equiv \langle [\,], \texttt{Var14} \rangle$
ICFGNode 3	$\wedge \quad \langle [\texttt{ICFGNode5}], \texttt{Var6} \rangle \equiv \langle [\texttt{ICFGNode5}], \texttt{Var7} \rangle$
ICFGNode 6	$\wedge \quad \langle [\,], \texttt{Var13} \rangle \equiv \langle [\texttt{ICFGNode5}], \texttt{Var6} \rangle$
ICFGNode 7	$\wedge \quad \langle [\texttt{ICFGNode7}], \texttt{Var7} \rangle \equiv \langle [\,], \texttt{Var12} \rangle$
ICFGNode 3	$\wedge \quad \langle [\texttt{ICFGNode7}], \texttt{Var6} \rangle \equiv \langle [\texttt{ICFGNode7}], \texttt{Var7} \rangle$
ICFGNode 8	$\wedge \quad \langle [\,], \texttt{Var15} \rangle \equiv \langle [\texttt{ICFGNode7}], \texttt{Var6} \rangle$
ICFGNode 9	$\wedge \quad \texttt{Var16} \equiv \texttt{ite}(\texttt{Var15} \equiv \texttt{Var12}, 1, 0)$

```
-------------SVFVar and Value-------------
                    ...
ValVar13                Value: 3
ValVar12                Value: 0
ValVar15                Value: 0
ValVar16                Value: 1
                    ...
------------------------------------------
```

Example 5: Interprocedural



ICFG

Verifying ICFG path: $0 \to 4 \to 5 \to 1 \to 2 \to 3 \to 6 \to 7 \to 1 \to 2 \to 3 \to 8 \to 9 \to$ *svf_assert*

ICFG Node/Edge	Constraints in the solver
ICFGNode 0	$\texttt{Var21} \equiv 0 \wedge \texttt{Var14} \equiv 3 \wedge \ldots$
ICFGNode 5	$\wedge \quad \langle [\texttt{ICFGNode5}], \texttt{Var7} \rangle \equiv \langle [\,], \texttt{Var14} \rangle$
ICFGNode 3	$\wedge \quad \langle [\texttt{ICFGNode5}], \texttt{Var6} \rangle \equiv \langle [\texttt{ICFGNode5}], \texttt{Var7} \rangle$
ICFGNode 6	$\wedge \quad \langle [\,], \texttt{Var13} \rangle \equiv \langle [\texttt{ICFGNode5}], \texttt{Var6} \rangle$
ICFGNode 7	$\wedge \quad \langle [\texttt{ICFGNode7}], \texttt{Var7} \rangle \equiv \langle [\,], \texttt{Var12} \rangle$
ICFGNode 3	$\wedge \quad \langle [\texttt{ICFGNode7}], \texttt{Var6} \rangle \equiv \langle [\texttt{ICFGNode7}], \texttt{Var7} \rangle$
ICFGNode 8	$\wedge \quad \langle [\,], \texttt{Var15} \rangle \equiv \langle [\texttt{ICFGNode7}], \texttt{Var6} \rangle$
ICFGNode 9	$\wedge \quad \texttt{Var16} \equiv \texttt{ite}(\texttt{Var15} \equiv \texttt{Var12}, 1, 0)$
ICFGNode 10	$\wedge \quad \texttt{Var16} \equiv 0$ (negation of the assert condition)

Solver yields **UNSAT**, meaning no counter example.

The assertion is verified successfully!!

What's next?

- (1) Understand SSE algorithms and examples in the slides
- (2) Finish implementing the automated translation from code to Z3 formulas using `SSE` and `Z3SSEMgr` in Assignment 2