Assignment -1 What is data structure? What she the various types of data Data structure is method of organizing and storing data in completer's memory to perporm operations on data. operations on data. The Various operations performed are data retrieval, data storage and manipulation of data nata structures are divided into 2 main types is 1) linear data structures 27 Non-linear data structures In linear structures all the elements are stored in a Examples of linear structures
ore:

1) Array A collection of elements that are identified 2) linked lists. All the elements are linked together by 3) Stocks: It polloues last in First out (LIFO). Special 4) anene operations are push, pop

42 anene Follower First In, First out (FIFO). Special sperations are enqueue and dequeue In Non-linear data structures the elements involve structural relationships between each other 1) Trees: Heirarchical method of storing data which has main node and multiple sub nodes. 27 Creraphs: It is a collection of vertices and edges that connect pairs of nodes. 2 What is a structure? How is it different from arrays? Explain different types structure declaration with example. Ans. A structure is a collection of data items, where each item is identified as to its type and name. It can hold data of multiple data types and not a single as serious are isllection of data of same type.

Elements of same type accessed

of structure are accessed using dot operator (,) and pield name.

the types of structure declaration 1) Along with structure definition struct string int i) 3 yar; whe can add palues to it in main part by using vor Ec: Var i = 1) 27 outside structure definition struct strig

inti;

share; Void main () } struct stri var; whe can sold values to it by

wing a valuable with and assign and retrieve Values. Er: Va97. i = 1 ;

Va97. i = 1;

Va97. i = 1; Using typeded beguesord: typeded struct stor 13 int i; shar i; 3 to structure 1; and soid main () structure 1 yar; 1 Wor. 1 = 1 ), Va91 ( = ' 1 ') Define pointers, how to declare and initialize pointers, explain with stores memory address of another directly with memory and help

Syntax por declaring a pointer is where is datatype represents the type of data pointing to. ii) plor represents the name of Ex int \* ptor;

whan \* ptor;

yord \* ptr; A void pointer can be declared but it has to be type casted to make it to point to a variable A pointer is initialized by using an & sperator int \\ \( \partial = 10 \) int \*ptor; pter = & var; Explain dynamic memory allocation in detail for Variables or data structures during runtime of program. It remove can be dynamically allocated using malloc (), calloc (), present in stall h header pile 1) malloc () - It pllocates single large block of memory weith specific size

It returns a pointer of type Void Syntax: pto: ( cost-type \*) mallor ( lyte. Ex: pto: (int +) mallor (100 + siges (int); If the int suge is 4, 400 leytes The pointer holds address of first light in allocated memory.

The space is insufprisent it returns NULL. 2) collor () -> . It dynamically allocates specified no of blocks of memory of specified type.

It initializes each block with depoult value o. Syntax: ptor = (cost-type \*) calloces, uehere n - rige of no of elements. elem-size - size of each element Ex: pter = (ploat +) calloc (>5, size of (plant))
The statement allocates contagious space por 25 elements each with The space insupplicient it returns

NULL pointer

3) reallow () -7. It helps in resigning

of allowated memory.

It increases or decreases

singe

syntax pter = reallor (pter, new sing); ent \* ptor; pto = (int +) mallor (3 \* suggest cents) uptor = reallor (ptor, 4" singer (and)) 4) preech - It is used to pree so deallocate Syramically allocated memory and helps in reducing memory weartage Syntax preecptor); Ex pree (ptor); What is a sparse materia?

Showe with suitable example

sparse materia representation stating as triplets. . A spared matrix is a matrix of mxn order to in which the number of o elements is greater than non-yero elements En: A= [0 0 0 5] 3010 LO 0 0 0 J In this matrix, the No. of o elements are 9 and Non-gero elements are 3 Representation of sparse matrix is done with the help of Duples

This is because if we use 20 array, lot of space well be weasted since no of a entries is .. We use 4 now, cal, value > where row - Total no of rower col - Total no. of columns when - Total no of non-yero values E1: A = [0 0 5 0] 4 0 0 0 0 Consider consider this sparse typeded struct & int rone; ent val; 3 x term; term 45507; When we consider 's which is structures, the representation well 4617 4 [2] A ( 37)

a [4) 1550 a [49] The oth index represent the total no of Iroues, columns, Non-yero values Holloweed by the position of non-zero when from indexing Express the given sparse matrix as buplets and pind its transpose algorithm to transpose a sparse

	67
	Algorithm:
1	
	12 Declare matrices a and l
	where a is original sparse
	matrix & le is transpose sparse
	matrix
	21 peclare sariables in which stores
	no of non-yero volues, i, g and
	urrentl.
	3) Assign print o indexing of row of
	I matrix to a indexing of
	solumn of a matrix.
	4) Similarly for a le 50 J. when
	and be Co T value
	5) Initiato 2 loops i 2 j nohero
	i is from o to a cos il
	jus from to n
	467 compare acj I wal to i and
	if it is town perform the
	transpos and at or end indrement
	surrently by 1
	Talkingto ramost & don't ha
	Howe do you represent a polynomial
	using an array of societies
	and also write a pune to ADD
	2 polynomials
	A polynomial son be represented
	in pollowing manner:
	mack transfers white

>

struct ?

ploat coep;

int expon;

3 polynomial; typedep The punction to add 2 polynomials polynomial terms (50); soid padd (int starta, int pinisha int starth, int pinishat, int \* starth, int \* pinishat, ploat coeppicient;

\*start = arrail;

while (starta = pinisha & f
start & = pinish &) sweitch (Compare (terms Estarta), expon, terms [start b]. expon case -1:
attach (terms (startl-).
coef, terms (startl-)
expon );
startl ++; soeppicient = terms (startle) attach ( coeppicient, terms (
starta). expon)

starta ++;
startl ++;
lreab;
case -1 x: attach (terms ( start a) rolf, terms [startla]. expon);
starta;
break; por (; starta = pinisha; starta++) attach Cterms [starta] , web, terms (starta) espon? por (; startl == pinish b; startl++) stack ( terms [startle ] rolf, terms [ startl ] expon) \* pinishd = arrail -1) Yord attach (ploat coeff, int expon) if (avail >= 50) prints (" too many terms");
exit(1);

terms [avail] wef = wef;
terms [avail ++] . expon = expon;

8 Write the 'knuth morris pratt algorithm por pattern matching #include & state h > # include & string h > int pmalch bid pail (); int proteh (); int pailwer [100] shar storing [100] char pat [100] int proatch ( share \* string, share \* put) int i=0, j=0; int lens = stolen (storing); int lenp = steller (pat ); while (i Llers 2 ) { Llers 2 } of (string [1] == pat []) i + + ; j + + ; } else if (j==0) i++; else j=pailure (j-1)+1; return ((j == lenp)? (i-lenp):-1

store void pail ( char \* put) and standards at the same and pailure (0) = -1; poer (j=1; j < n; j++) - A KILL ARUE OF BLUDS TO i = pailure (j -1); while ((pat (j)! = pat (i+1)) ) } (1)=0)) i = pailure (i) 4 (pat cj ) = = pat ci+1) pailure (j) = i + 1;
else pailure (j) = -1; COSTELLORD AL LORD in a morety of true years

Write a c program to a compare strings b. concatenate 2 strings Ans a) # include / stdio. h > Hinclude Latting. h > char s1 () = "abc"; char 52[] = "def"; int or = strump (s1, s2); y (2==0) prints (" strongs are equal") else if (result 40) point ("String 1 is less
than string 2"); prints (" String 1 is greater than string 2");

les void strans (char \* s, char \* t, int i) char stericson, \* temp = steri; points (" Position out of bounds ") eaturn; 30 locked my of ( stocler (s)) else if (storler (t)) strong (temp, s, i); stoccat (temp, t); streat (temp, (s+i)); stripy (s, temp);

10 pepine stack the implementation of pop push and display punctions their por empty and pull Ans Stack is a data structure that pollones last in, First out principle.

It is an ordered list in why
insertions and deletions are
made at one end which is
palled top end. Function to push an element onto stack (add element):

Void push (int top,

Void push ()

3

if (top 4 ring-1) printly ("Finter element to be

added: ");

scarf (".1-d", 2 ele);

scrett top ) = ele;

generally ("Stack Overplane in");

3 Function to pop item prom stock

int pop()

if ctop ==-1)

printf('' stock underploye'n")

proturn -1;

3

else

3 return sctop--J; Function to sisplay elements of Void display () if (top ==-1) points (" Stack is empty \n");

else &

prints (" Stack content: ");

por ( i = top; i >= 0; i--)

& points (".1-din", scis);

Write punctions to convert infix expression to postfix and convert the pollowing expression to postfix using stack:

A\*(B+C\*D)+E #include Lotdio h > char inpia [100], postpia [100], be a char 5 [ 50], top = -1; Void push ( shor elem) 3 St ++ top J = elem; char pap () return (sctop -- J) int precedence ( char elem) switch celem? case ( : return ;; pase '+': roturn 2; 1. return 3; default: return o;

store youd connect ( share it) sweitch (ch) case 'C': push (ch); week; uchilo ((ch = pop())!='(') postpic [b++ ) = ch; week) 1-1: dute ' ./. ': case Case while ( precedence (ch) = precedence (s (top)) postpic C++J=pap (); push (ch); break; default: postpix Cb++J=ch; prints ("Enter infin expression:"); gets (inpix); push ('10')

par ( ent i =0; infracio)!= 10'site ¿ connert Cinpir (i ); postpix (b++) = pop()

pointp (''Postpix Exps://s/n'', postpii)

return 0; For the expression: A\*(B+C\*D) +E Input stack Postpix