# Information provision strategy for congestion mitigation

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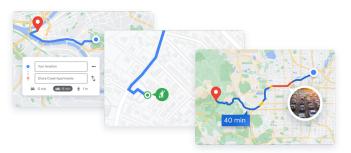
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### 1. Introduction

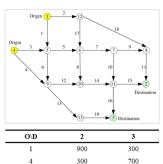
#### Motivations

- Information Paradox (Acemoglu et al., 2018): A larger information path set possibly leads to increased congestion.
- Information penetration: How many drivers should be informed? ⇒ Fairness
- Restricted information set: How many paths should be posted? ⇒ Unawareness
- Differentiated information provision: Users have differing information sets.



### 1. Introduction

#### Nguyen-Dupuis network:



Travel	demand	for	each	OD	pair

O-D	Route	Node sequence	O-D	Route	Link sequence
(1, 2)	1	1-12-8-2		6	1-5-9-13-3
	2	1-5-6-7-8-2	(1, 3)	7	1-5-6-7-11-3
	3	1-5-6-7-11-2		8	1-5-9-10-11-3
	4	1-5-9-10-11-2		9	1-12-6-7-11-3
	5	1-12-6-7-8-2		10	1-5-6-10-11-3
(4, 2)	11	4-9-10-11-2		16	4-9-13-3
	12	4-5-6-7-8-2	(4, 3)	17	4-9-10-11-3
	13	4-5-6-7-11-2		18	4-5-9-13-3
	14	4-5-9-10-11-2		19	4-5-6-7-11-3
	15	4-5-6-10-11-2		20	4-5-9-10-11-3

Potential path set for each OD pair

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Bi-level optimization for deriving optimal information provision strategy:

$$\min \quad F(x,\tau) = \sum_{a \in A} x_a t_a(x_a) + \beta \sum_{r \in R} \sum_{s \in S} \sum_{\pi \in \Pi_u^{r,s}} (1 - \tau_\pi^{r,s})$$

$$\text{s.t.} \quad \begin{cases} \sum_{\pi \in \Pi_u^{r,s}} (1 - \tau_\pi^{r,s}) \ge I & \forall r \in R, \forall s \in S \\ \sum_{\pi \in \Pi_u^{r,s}} (1 - \tau_\pi^{r,s}) \le u & \forall r \in R, \forall s \in S \end{cases}$$

$$(1)$$

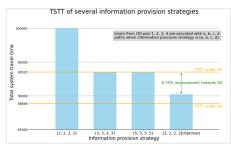
where  $x = x(\tau)$  is implicitly determined by:

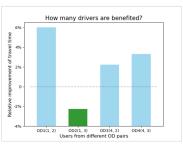
min 
$$f(x,\tau) = \sum_{a \in A} \int_{r \in R}^{\sum \sum \sum \sum \sum \sum r \in I_u^{r,s}} \delta_{\pi}^{r,s} h_{\pi}^{r,s} t_a(x) dx$$
  
s.t. 
$$\begin{cases} \sum_{a \in \Pi_u^{r,s}} h_{\pi}^{r,s} = q^{r,s} & \forall r \in R, \forall s \in S \\ h_{\pi}^{r,s} \ge 0 & \forall r \in R, \forall s \in S, \forall \pi \in \Pi_u^{r,s} \\ \tau_{\pi}^{r,s} h_{\pi}^{r,s} = 0 & \forall r \in R, \forall s \in S, \forall \pi \in \Pi_u^{r,s} \end{cases}$$

$$(2)$$

 $\tau_{\pi}^{r,s}=0$  if user is aware of path  $\pi\in\Pi_{u}^{r,s}$  connecting OD pair (r, s);  $\tau_{\pi}^{r,s}=1$  otherwise.

## 3. Illustrative example





OD pair	(1, 2)	(1, 3)	(4, 2)	(4, 3)
Shortest travel time under UE	42.95	42.39	42.54	41.97
Shortest travel time under OIPS	40.37	43.35	41.59	40.58

The shortest travel time for each OD pair under user equilibrium(UE) and optimal information provision strategy(OIPS) flow pattern

#### 4. Future directions

- Bounded Rationality: Travelers may not choose the shortest path despite receiving path information. Modeling this behavior can effectively mitigate the congestion exacerbated by the concentration caused by traveler's similar choices.
- Costly Information: Fees can be implemented for information of critical paths. This
  encourages travelers to choose suboptimal routes, alleviating congestion on key
  paths and promoting a more balanced traffic distribution, ultimately guiding the
  network toward a system optimal state.
- Differentiated User Type: In addition to varying the quantity of path information, we
  can also differentiate its content. For travelers of type A, information on the 1st,
  2nd, and 4th paths can be provided, while travelers of type B can receive
  information on the 2nd, 3rd, and 5th paths.

## Thanks for your attention!

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### Reference I

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